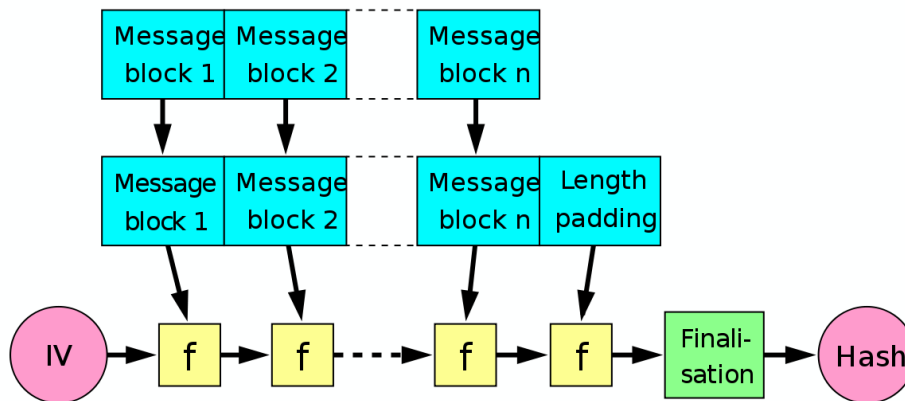
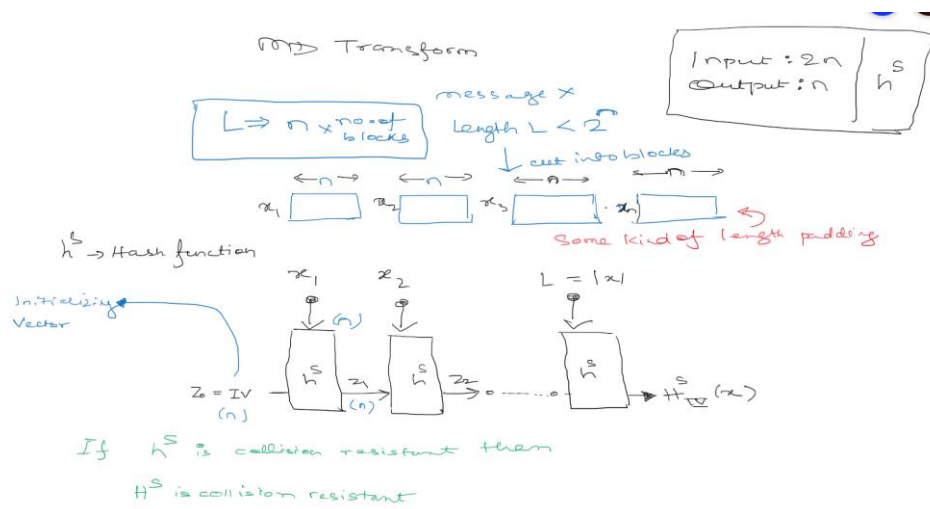


Using MD Transform to obtain a (provably secure) Collision Resistant Hash Function (Theory)

MD Transform or Merkle Damgard Transform is a methodology that is used to constructing a Hash function for longer inputs using a fixed length hash function.



Sketches/Ideas :



Construction:

Once our Fixed length collision resistant hash function is ready to be used as an api. We can further dive in to just use it and turn it into a Provably secure Collision resistant Hash function.

We take two inputs to merkle damgard, an initialization vector and another the message whose hash is to be found.

Let $l(n) = |iv|$

CONSTRUCTION 4.11 The Merkle-Damgård Transform.

Let (Gen_h, h) be a fixed-length hash function with input length $2\ell(n)$ and output length $\ell(n)$. Construct a variable-length hash function (Gen, H) as follows:

- $\text{Gen}(1^n)$: upon input 1^n , run the key-generation algorithm Gen_h of the fixed-length hash function and output the key. That is, output $s \leftarrow \text{Gen}_h$.
- $H^s(x)$: Upon input key s and message $x \in \{0,1\}^*$ of length at most $2^{\ell(n)} - 1$, compute as follows:
 1. Let $L = |x|$ (the length of x) and let $B = \lceil \frac{L}{\ell} \rceil$ (i.e., the number of blocks in x). Pad x with zeroes so that its length is an exact multiple of ℓ .
 2. Define $z_0 := 0^\ell$ and then for every $i = 1, \dots, B$, compute $z_i := h^s(z_{i-1} \| x_i)$, where h^s is the given fixed-length hash function.
 3. Output $z = H^s(z_B \| L)$

At first the message is augmented with zeros at the end to make it a multiple of the size of the initialization vector. Then we encode the length of the message at the end of the message as a binary string. Now all we have to do is setting up the parameters for our fixed length hash function x_1 and x_2 in a loop. We will set x_2 as our message fragment and x_1 is initially our iv and in later stages of propagation becomes the output of our hash function. Both x_1 and x_2 are of size $\ell(n)$, so total data feeded to our hash function $\text{hash}(x_1, x_2)$ is of size $2 \cdot \ell(n)$. Now our gen algorithm as already discussed in previous document generates the group parameters which will then be used along with DLP to generate a hash of 16 bits. After a number of iterations equal to the number of message fragments we will get our final hash.

This version has a limitation because of our fixed length hash function which supports upto 16bits.