Hashing Technique

useful for searching

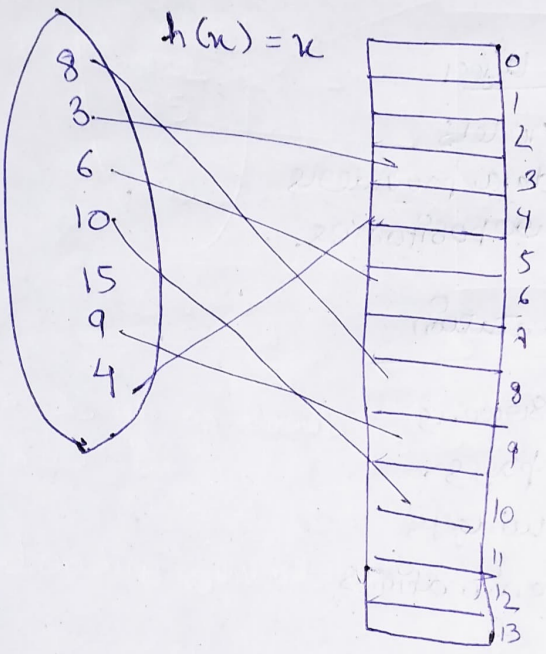
① linear — O(n)

② Binary — O(logn)

O(1)

keys: 8, 3, 6, 10, 15, 9, 4

key space



drawback is that space taken is very large.

mapping

functions

function
- one - one
- one - many
- many - one
- many - many

many to one function

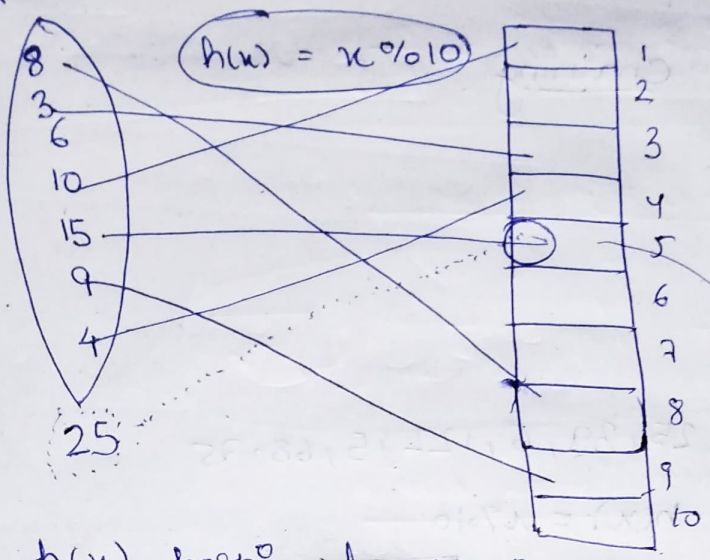
this is

method

Open

Close

many to one function



h(x) hashing function is responsible for space complexity of hash table.

Two keys are mapped in one position this is called collission.

Method to avoid collision

Open Hashing → size can increase.

→ Chaining

Closed Hashing → fixed space.

Open addressing

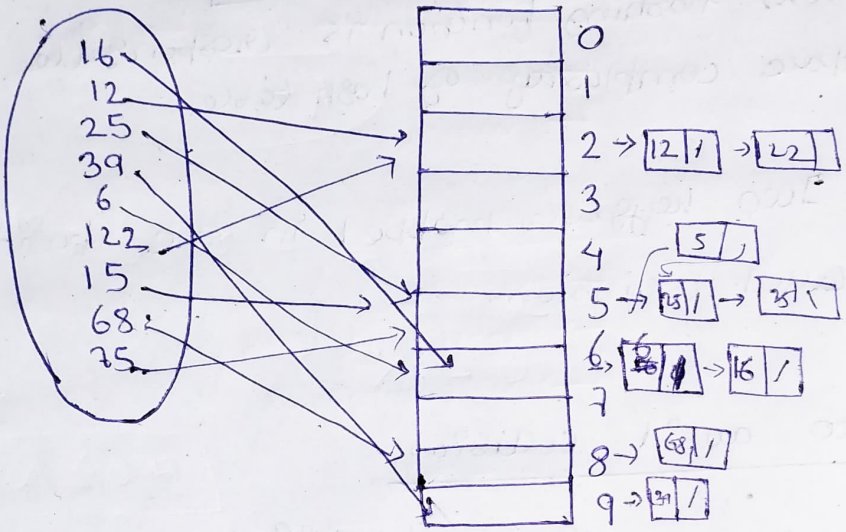
1. linear probing
2. Quadratic probing
3. Double Hashing

Chaining -

1. Insert
2. Search
3. Analysis
4. Delete

keys : 16, 12, 25, 39, 6, 122, 15, 68, 75

$h(x) = x \% 10$



made up of array of chain or array of linked list.

loading factor $\lambda = \frac{n}{size}$

Successful Search

Average successful time

$t = 1 + \lambda/2$

Average unsuccessful time

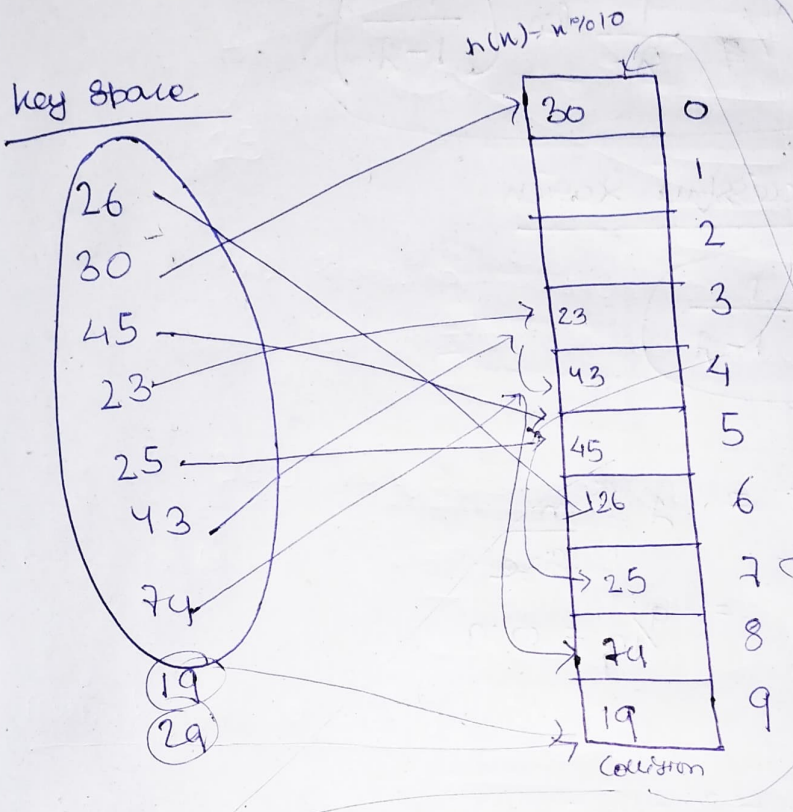
$t = 1 + \lambda$

Hashing Technique

(Linear probing)

$h'(x) = (h(x) + f(i)) \% 10$ where $f(i) = i$

$i = 0, 1, 2, \ldots$



$h'(25) = (h(25) + f(0)) \% 10$

$(5 + 0) \% 10 = 5$ already occupied

$h'(25) = (h(25) + f(1)) \% 10$

$(5 + 1) \% 6 = 6$ already occupied

$h'(25) = (h(25) + f(2)) \% 10$

$(5 + 2) \% 10 = 7$ free

↳ It will go cyclic

Search → take the hash function go to index then go on linearly until key is found or a empty space is found.

loading factor $\lambda = \frac{n}{size}$

Avg successful search

$$t = \frac{1}{\lambda} \ln\left(\frac{1}{1-\lambda}\right)$$

Avg unsuccessful search

$$t = \frac{1}{1-\lambda}$$

loading factor = $\lambda = \frac{n}{size}$

$= 9/10 = 0.9$

☆

$\lambda \leq 0.5$

Drawback

→ cluster formation

→ A block of key at one place.

primary clustering

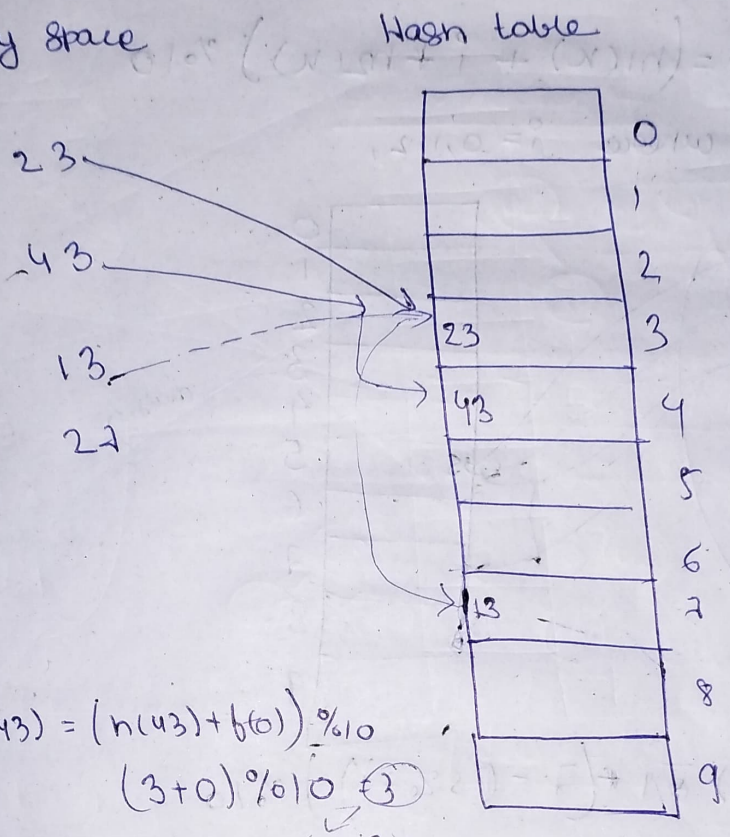
for deleting make a flag

Quadratic Probing

$h'(x) = (h(x) + f(i)) \% 10$ where

$f(i) = i^2$

$i = 0, 1, 2, \ldots$



$h'(43) = (h(43) + f(0)) \% 10$

$(3+0) \% 10 = 3$

collision

$h'(43) = (h(43) + f(1)) \% 10$

$(3+1) \% 10 = 4$

$h'(13) = h(43) + f(2)) \% 10$

$(3+4) \% 10 = 7$

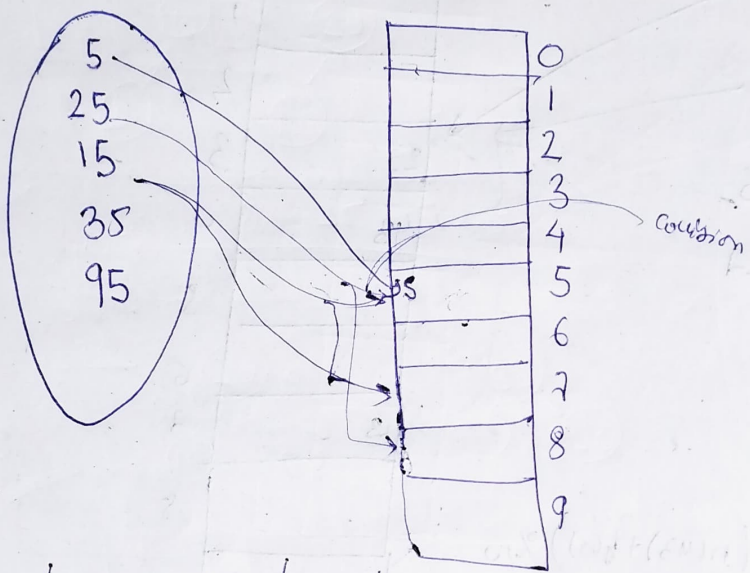
Hashing Technique (Double Hashing)

$h1(x) = x \% 10$

$h2(x) = 7 - (x \% 7)$

$h'(x) = (h1(x) + i * h2(x)) \% 10$

where $i = 0, 1, 2, \ldots$



$(h(25) + 1 * (7 - (25 \% 7))) \% 10$

$5 + 3 = 8$

$h'(15) = (5 + (2 * 6)) \% 10$

Ans 2