Let X_i be a Bernoulli random variable which takes value 1 when the algorithm predicts *incorrectly* in the *ith* attempt. So, for $i \in [N]$, $\delta > 0$, we have:

$$P(X_i = 1) = \frac{1}{2} + \delta$$
$$P(X_i = 0) = \frac{1}{2} - \delta$$

Let Y be the random variable which tracks the total number of incorrect answers given by the algorithm. So, $Y = \sum_{i=1}^{N} X_i$. As majority vote is used to predict the final answer, the answer is incorrect if Y > N/2.

$$\begin{split} P(Y > N/2) &= P(\lambda Y > \lambda N/2) = P(e^{\lambda Y} > e^{\lambda N/2}) & \lambda > 0 \\ &\leq e^{-\lambda N/2} \cdot \mathbb{E}\left[e^{\lambda Y}\right] & \text{Markov's Inequality} \\ &= e^{-\lambda N/2} \cdot \mathbb{E}\left[e^{\lambda \sum_{i=1}^{N} X_i}\right] & X_i\text{'s are independent} \\ &= e^{-\lambda N/2} \cdot \prod_{i=1}^{N} \mathbb{E}\left[e^{\lambda X_i}\right] & X_i\text{'s are identically distributed} \\ &= e^{-\lambda N/2} \cdot \left(\mathbb{E}\left[e^{\lambda X_1}\right]\right)^N & X_i\text{'s are identically distributed} \\ &= e^{-\lambda N/2} \cdot \left(e^{\lambda}\left(\frac{1}{2} - \delta\right) + \left(\frac{1}{2} + \delta\right)\right)^N & \text{MGF of a Bernoulli r.v.} \\ &= \left(e^{\lambda/2}\left(\frac{1}{2} - \delta\right) + e^{-\lambda/2}\left(\frac{1}{2} + \delta\right)\right)^N &= \left(\sqrt{1 - 4\delta^2}\right)^N \end{split}$$

The last step is obtained by finding the tightest bound. For tightest bound, we need the smallest possible value of RHS. As both terms inside the brackets is non-negative, we can apply AM-GM inequality.

$$\left(e^{\lambda/2}\left(\frac{1}{2}-\delta\right)+e^{-\lambda/2}\left(\frac{1}{2}+\delta\right)\right) \ge 2\cdot\sqrt{e^{\lambda/2}\left(\frac{1}{2}-\delta\right)\cdot e^{-\lambda/2}\left(\frac{1}{2}+\delta\right)} = \sqrt{1-4\delta^2}$$

which is obtained for $\lambda = \ln(\frac{1+2\delta}{1-2\delta}) > 0$. So, we get:

$$P(Y > N/2) \le (1 - 4\delta^2)^{N/2} \le e^{-2N\delta^2}$$
 As $1 + x \le e^x, \forall x \in \mathbb{R}$

To obtain a correct answer with probability $1-\epsilon$, i.e., $P(Y < N/2) \ge 1-\epsilon$, we need $P(Y > N/2) \le \epsilon$. To ensure $P(Y > N/2) \le \epsilon$ for any $\epsilon \in (0,1)$, it is enough to ensure $e^{-2N\delta^2} \le \epsilon$:

$$P(Y > N/2) \le e^{-2N\delta^2} \le \epsilon \implies N \ge \frac{1}{2\delta^2} \ln\left(\frac{1}{\epsilon}\right)$$

We have $X = (X_1, ..., X_n)$ such that $X_i \sim \mathcal{N}(0, 1)$ and are independent and we need to estimate $\mathbb{E}[\|X\|_2]$ and $\text{Var}(\|X\|_2)$. Notice that for a random variable $Z \geq 0$, we have:

$$\mathbb{E}[Z^p] = \int_0^\infty pt^{p-1} \cdot P(Z > t)dt \tag{1}$$

Proof. ¹ For a non-negative random variable Z,

$$Z^{p} = \int_{0}^{Z} pt^{t-1}dt = \int_{0}^{\infty} pt^{t-1} \cdot \mathbf{1}(Z > t)dt$$

Taking expectations on both sides,

$$\mathbb{E}[Z^p] = \mathbb{E}\left[\int_0^\infty pt^{t-1} \cdot \mathbf{1}(Z > t)dt\right]$$
$$= \int_0^\infty pt^{t-1} \cdot \mathbb{E}[\mathbf{1}(Z > t)]dt$$
$$= \int_0^\infty pt^{p-1} \cdot P(Z > t)dt$$

The following inequality was proved in the class:

$$P(\left| \|X\|_2 - \sqrt{n} \right| > \epsilon) \le 2e^{-\epsilon^2/8}, \epsilon > 0 \tag{2}$$

(a) Let $Z = |||X||_2 - \sqrt{n}|$ and p = 1,

$$\begin{split} \left| \mathbb{E} \left[\| X \|_2 - \sqrt{n} \right] \right| &\leq \mathbb{E} \left[\left| \| X \|_2 - \sqrt{n} \right| \right] & \text{Jensen's Inequality, i.e., } |\mathbb{E}[X]| \leq \mathbb{E}[|X|] \\ &\stackrel{(1)}{=} \int_0^\infty P(\left| \| X \|_2 - \sqrt{n} \right| > t) dt \\ &\stackrel{(2)}{\leq} \int_0^\infty 2e^{-t^2/8} dt \\ &= \int_0^\infty 4\sqrt{2} \cdot e^{-r^2} dr \\ &= 2\sqrt{2\pi} & \text{Using } \int_0^\infty e^{-t^2} dt = \frac{1}{2} \int_{-\infty}^\infty e^{-t^2} dt = \frac{1}{2} \sqrt{\pi} \end{split}$$

So,
$$\left| \mathbb{E} \left[\|X\|_2 - \sqrt{n} \right] \right| \le 2\sqrt{2\pi} \implies \sqrt{n} - 2\sqrt{2\pi} \le \mathbb{E} \left[\|X\|_2 \right] \le \sqrt{n} + 2\sqrt{2\pi}$$

¹https://math.stackexchange.com/questions/172841

$$\begin{aligned} \text{(b) Let } Z &= |\|X\|_2 - \sqrt{n}| \text{ and } p = 2, \\ \text{Var}(\|X\|_2) &= \text{Var}\left(\|X\|_2 - \sqrt{n}\right) \\ &\leq \mathbb{E}[(\|X\|_2 - \sqrt{n})^2] & \text{As } \text{Var}(X) = \mathbb{E}[X^2] - \mathbb{E}[X]^2 \leq \mathbb{E}[X^2] \\ &= \mathbb{E}\left[\left|\|X\|_2 - \sqrt{n}\right|^2\right] \\ &\stackrel{(1)}{=} \int_0^\infty 2t \cdot P(\left|\|X\|_2 - \sqrt{n}\right| > t) dt \\ &\stackrel{(2)}{=} \int_0^\infty 2t e^{-t^2/8} dt = 2 \cdot \frac{e^{-t^2/8}}{-1/4} \Big|_0^\infty = 8 \\ &\text{So, } \text{Var}(\|X\|_2) \leq 8 \end{aligned}$$

Given $F_d(\boldsymbol{\theta}) : \mathbb{R}^{\binom{d}{2}} \to \mathbb{R}$,

$$F_d(\boldsymbol{\theta}) = \log \left(\sum_{x \in \{\pm 1\}^d} \exp \left(\frac{1}{\sqrt{d}} \sum_{i < j} \theta_{ij} x_i x_j \right) \right)$$

Let $\boldsymbol{\theta}, \boldsymbol{x}$ be $\binom{d}{2}$ -length vectors such that $\boldsymbol{\theta} = (\theta_{ij})_{1 \leq i < j \leq d}$, $\boldsymbol{x} = f(x) = (x_i x_j)_{1 \leq i < j \leq d}$ and $\mathcal{X} = \{\boldsymbol{x} | \boldsymbol{x} = f(x), \forall x \in \{\pm 1\}^d\}$. Note that \mathcal{X} is a multiset, i.e., there could be multiple copies of the same element. So, $F_d(\boldsymbol{\theta})$ can be re-written as follows:

$$F_d(\boldsymbol{\theta}) = \log \left(\sum_{\boldsymbol{x} \in \mathcal{X}} \exp \left(\frac{1}{\sqrt{d}} \boldsymbol{\theta}^{\top} \boldsymbol{x} \right) \right)$$

Lemma 1. The log-sum-exp function ,i.e., $f(x) = \log(e^{x_1} + \cdots + e^{x_n})$ is convex in \mathbb{R}^n .

Proof. We prove f(x) is convex by showing that $\nabla^2 f(x) \succeq 0$. Let $y = [e^{x_1}, \dots, e^{x_n}]^{\top}$

$$\nabla f(x) = \frac{1}{e^{x_1} + \dots + e^{x_n}} [e^{x_1}, \dots, e^{x_n}]^\top = \frac{y}{\mathbf{1}^\top y}$$
$$\left[\nabla^2 f(x) \right]_{ij} = \frac{(\mathbf{1}^\top y) e^{x_i} \cdot \mathbf{1} (i = j) - e^{x_i} e^{x_j}}{(\mathbf{1}^\top y)^2}$$
$$\implies \nabla^2 f(x) = \frac{1}{(\mathbf{1}^\top y)^2} \left(\left(\mathbf{1}^\top y \right) \mathbf{diag}(y) - yy^\top \right)$$
For any $v \in \mathbb{R}^n, v^\top \nabla^2 f(x) v = \frac{1}{(\mathbf{1}^\top y)^2} \left(\left(\sum_{i=1}^n y_i \right) \left(\sum_{i=1}^n y_i v_i^2 \right) - \left(\sum_{i=1}^n v_i y_i \right)^2 \right) \ge 0$

The last line follows by applying Cauchy-Schwarz inequality for the vectors $a = (\sqrt{y_1}, \dots, \sqrt{y_n})$ and $b = (v_1\sqrt{y_1}, \dots, v_n\sqrt{y_n})$. So, $\nabla^2 f(x) \succeq 0 \implies f(x)$ is convex.

- (a) Clearly $F_d(\theta)$ is a log-sum-exp function of the affine transformation of θ . So, using Lemma 1 $F_d(\theta)$ is convex.
- (b) We will first compute the gradient of $F_d(\theta)$ w.r.t θ .

$$\nabla_{\boldsymbol{\theta}} F_d(\boldsymbol{\theta}) = \frac{d}{d\boldsymbol{\theta}} \log \left(\sum_{\boldsymbol{x} \in \mathcal{X}} \exp \left(\frac{1}{\sqrt{d}} \boldsymbol{\theta}^\top \boldsymbol{x} \right) \right)$$

$$= \frac{1}{\left(\sum_{\boldsymbol{x} \in \mathcal{X}} \exp \left(\frac{1}{\sqrt{d}} \boldsymbol{\theta}^\top \boldsymbol{x} \right) \right)} \cdot \sum_{\boldsymbol{x} \in \mathcal{X}} \frac{\boldsymbol{x}}{\sqrt{d}} \cdot \exp \left(\frac{1}{\sqrt{d}} \boldsymbol{\theta}^\top \boldsymbol{x} \right)$$
Rearranging,
$$\nabla_{\boldsymbol{\theta}} F_d(\boldsymbol{\theta}) = \frac{1}{\sqrt{d}} \sum_{\boldsymbol{x} \in \mathcal{X}} \boldsymbol{x} \cdot \frac{\exp \left(\frac{1}{\sqrt{d}} \boldsymbol{\theta}^\top \boldsymbol{x} \right)}{\left(\sum_{\boldsymbol{x} \in \mathcal{X}} \exp \left(\frac{1}{\sqrt{d}} \boldsymbol{\theta}^\top \boldsymbol{x} \right) \right)} = \frac{1}{\sqrt{d}} \mathbb{E}_p[\boldsymbol{x}]$$

where, p is a probability distribution of \boldsymbol{x} such that $p(\boldsymbol{x}) = \frac{\exp\left(\frac{1}{\sqrt{d}}\boldsymbol{\theta}^{\top}\boldsymbol{x}\right)}{\left(\sum_{\boldsymbol{x}\in\mathcal{X}}\exp\left(\frac{1}{\sqrt{d}}\boldsymbol{\theta}^{\top}\boldsymbol{x}\right)\right)}$

Now, we will estimate the norm of the gradient of $F_d(\theta)$ w.r.t θ .

$$\|\nabla_{\boldsymbol{\theta}} F_{d}(\boldsymbol{\theta})\|_{2} = \frac{1}{\sqrt{d}} \|\mathbb{E}_{p}[\boldsymbol{x}]\|_{2}$$

$$\leq \frac{1}{\sqrt{d}} \mathbb{E}_{p}[\|\boldsymbol{x}\|_{2}] \qquad \text{Jensen's inequality for } \ell_{2}\text{-norm, i.e.,convex function}$$

$$= \frac{1}{\sqrt{d}} \cdot \sqrt{\frac{d(d-1)}{2}} \quad \text{As each coordinate of } \boldsymbol{x} \in \{\pm 1\} \text{ and } \boldsymbol{x} \text{ is } \binom{d}{2}\text{-dimensional}$$

$$\leq \sqrt{d/2} \qquad (3)$$

In the rest of the proof, norm refers to ℓ_2 -norm unless stated otherwise. As $F_d(\theta)$ is a convex function, by using the first order condition for convexity, we get:

$$F_{d}(\boldsymbol{\theta}') \geq F_{d}(\boldsymbol{\theta}) + \nabla_{\boldsymbol{\theta}} F_{d}(\boldsymbol{\theta})^{\top} (\boldsymbol{\theta}' - \boldsymbol{\theta})$$

$$\Rightarrow F_{d}(\boldsymbol{\theta}') - F_{d}(\boldsymbol{\theta}) \geq \nabla_{\boldsymbol{\theta}} F_{d}(\boldsymbol{\theta})^{\top} (\boldsymbol{\theta}' - \boldsymbol{\theta}) \qquad (4)$$

$$F_{d}(\boldsymbol{\theta}) \geq F_{d}(\boldsymbol{\theta}') + \nabla_{\boldsymbol{\theta}'} F_{d}(\boldsymbol{\theta}')^{\top} (\boldsymbol{\theta} - \boldsymbol{\theta}')$$

$$\Rightarrow F_{d}(\boldsymbol{\theta}') - F_{d}(\boldsymbol{\theta}) \leq \nabla_{\boldsymbol{\theta}'} F_{d}(\boldsymbol{\theta}')^{\top} (\boldsymbol{\theta}' - \boldsymbol{\theta}) \qquad (5)$$

$$\Rightarrow \nabla_{\boldsymbol{\theta}} F_{d}(\boldsymbol{\theta})^{\top} (\boldsymbol{\theta}' - \boldsymbol{\theta}) \stackrel{(4)}{\leq} F_{d}(\boldsymbol{\theta}') - F_{d}(\boldsymbol{\theta}) \stackrel{(5)}{\leq} \nabla_{\boldsymbol{\theta}'} F_{d}(\boldsymbol{\theta}')^{\top} (\boldsymbol{\theta}' - \boldsymbol{\theta})$$

$$\text{So, } |F_{d}(\boldsymbol{\theta}') - F_{d}(\boldsymbol{\theta})| \leq \max(|\nabla_{\boldsymbol{\theta}} F_{d}(\boldsymbol{\theta})^{\top} (\boldsymbol{\theta}' - \boldsymbol{\theta})|, |\nabla_{\boldsymbol{\theta}'} F_{d}(\boldsymbol{\theta}')^{\top} (\boldsymbol{\theta}' - \boldsymbol{\theta})|)$$

$$\leq \max_{\boldsymbol{\alpha} \in \{\boldsymbol{\theta}, \boldsymbol{\theta}'\}} (\|\nabla_{\boldsymbol{\alpha}} F_{d}(\boldsymbol{\alpha})\| \|\boldsymbol{\theta} - \boldsymbol{\theta}'\|) \qquad \text{As } |a^{\top}b| \leq \|a\| \|b\|$$

$$= \|\boldsymbol{\theta} - \boldsymbol{\theta}'\| \max_{\boldsymbol{\alpha} \in \{\boldsymbol{\theta}, \boldsymbol{\theta}'\}} (\|\nabla_{\boldsymbol{\alpha}} F_{d}(\boldsymbol{\alpha})\|)$$

$$\stackrel{(3)}{\leq} \sqrt{d/2} \|\boldsymbol{\theta} - \boldsymbol{\theta}'\|$$

(c) We state an extension to a theorem discussed in the class (without detailed proof) that which will be pivotal in proving the required inequality.

Theorem 2 (Extension to Theorem 2.26 of MJW²). Let (X_1, \dots, X_n) be a vector of i.i.d. Gaussian random variables such that $X_i \sim \mathcal{N}(0, \sigma^2)$ and $f : \mathbb{R}^n \to \mathbb{R}$ be L-Lipschitz w.r.t. ℓ_2 -norm, then $f(X) - \mathbb{E}[f(X)]$ is sub-Gaussian with parameter at most $L\sigma$.

Proof. The proof can be obtained by following the proof schema as mentioned in MJW for a scaled version of X. Specifically, if $X \sim \mathcal{N}(0, \sigma^2)$, then $Y = \frac{X}{\sigma} \sim \mathcal{N}(0, 1)$.

²Wainwright, Martin J. High-dimensional statistics: A non-asymptotic viewpoint. Vol. 48. Cambridge University Press, 2019.

Given all $\theta_{ij} \sim \mathcal{N}(0, \sigma^2)$ for $1 \leq i < j \leq d$ and are independent, we have,

$$F_{d}(\boldsymbol{\theta}) = \log \left(\sum_{x \in \{\pm 1\}^{d}} \exp \left(\frac{1}{\sqrt{d}} \sum_{i < j} \theta_{ij} x_{i} x_{j} \right) \right)$$

$$\Rightarrow \mathbb{E}[F_{d}(\boldsymbol{\theta})] = \mathbb{E} \left[\log \left(\sum_{x \in \{\pm 1\}^{d}} \exp \left(\frac{1}{\sqrt{d}} \sum_{i < j} \theta_{ij} x_{i} x_{j} \right) \right) \right]$$

$$\leq \log \left(\mathbb{E} \left[\sum_{x \in \{\pm 1\}^{d}} \exp \left(\frac{1}{\sqrt{d}} \sum_{i < j} \theta_{ij} x_{i} x_{j} \right) \right] \right) \quad \text{As } \mathbb{E}[\log X] \leq \log \mathbb{E}[X]$$

$$= \log \left(\sum_{x \in \{\pm 1\}^{d}} \mathbb{E} \left[\exp \left(\frac{1}{\sqrt{d}} \theta_{12} x_{1} x_{2} \right) \right]^{\binom{d}{2}} \right) \quad \text{As } \theta_{ij} \text{'s are i.i.d}$$

$$= \log \left(\sum_{x \in \{\pm 1\}^{d}} \exp \left(\frac{x_{1}^{2} x_{2}^{2} \sigma^{2}}{2d} \cdot \binom{d}{2} \right) \right) \quad \text{MGF for } X \sim \mathcal{N}(0, \sigma^{2})$$

$$= \log \left(\sum_{x \in \{\pm 1\}^{d}} \exp \left(\frac{\sigma^{2} (d-1)}{4} \right) \right)$$

$$= \log \left(2^{d} \cdot \exp \left(\frac{\sigma^{2} (d-1)}{4} \right) \right)$$

$$= d \log 2 + \frac{\sigma^{2} (d-1)}{4} \leq d \log 2 + \frac{\sigma^{2} d}{4}$$

$$\Rightarrow F_{d}(\boldsymbol{\theta}) - \mathbb{E}[F_{d}(\boldsymbol{\theta})] \geq F_{d}(\boldsymbol{\theta}) - d \log 2 - \frac{\sigma^{2} d}{4}$$

$$(6)$$

As (X_1, \ldots, X_n) is a Gaussian vector such that $X_i \sim \mathcal{N}(0, \sigma^2)$ and $F_d(\boldsymbol{\theta})$ is $\sqrt{d/2}$ -Lipschitz, from Theorem 2, we know that $f(X) - \mathbb{E}[f(X)]$ is sub-Gaussian random variable with parameter atmost $\sigma \sqrt{d/2}$. So, using one sided tail bound for sub-Gaussian random variable, we have,

$$P(F_d(\boldsymbol{\theta}) - \mathbb{E}[F_d(\boldsymbol{\theta})] > t) \le \exp\left(\frac{-t^2}{2 \cdot \frac{d}{2}\sigma^2}\right)$$

$$\implies P(F_d(\boldsymbol{\theta}) - \mathbb{E}[F_d(\boldsymbol{\theta})] > dt) \le \exp\left(\frac{-dt^2}{\sigma^2}\right)$$

$$P(F_d(\boldsymbol{\theta}) - d\log_2 - \frac{\sigma^2 d}{4} > dt) \stackrel{(6)}{\le} \exp\left(\frac{-dt^2}{\sigma^2}\right)$$

$$\implies P\left(\frac{F_d(\boldsymbol{\theta})}{d} > \log_2 + \frac{\sigma^2}{4} + t\right) \le \exp\left(\frac{-dt^2}{\sigma^2}\right)$$

(a) Let $X = (X_1, ..., X_{2C})$ be a random vector such that X_i is the number of balls in bin i after all the T balls are thrown randomly. As each of the throws is independent of the other and each of the bins is equally probable to be obtained, X follows a multinomial distribution with parameters $T, p_i = 1/2C, i \in [2C]$.

Let S be the set of 2C-length binary vectors such that each vector has equal number of 0s and 1s. So, $|S| = {2C \choose C}$. Let $v \in S$. Then $v^{\top}X$ is the total number of balls in exactly in C bins. So, the number of balls in the top C bins, i.e., $M_C(T)$ equals the maximum of $v^{\top}X$, for $v \in S$.

$$\begin{split} \mathbb{E}[M_C(T)] &= \mathbb{E}[\max_{v \in \mathcal{S}} v^\top X] \\ &\leq \mathbb{E}\left[\frac{1}{\lambda} \log \left(\sum_{i=1}^{|\mathcal{S}|} \exp\left(\lambda v_i^\top X\right)\right)\right] \qquad \text{summation} \geq \text{maximum} \\ &\leq \frac{1}{\lambda} \log \left(\mathbb{E}\left[\sum_{i=1}^{|\mathcal{S}|} \exp\left(\lambda v_i^\top X\right)\right]\right) \qquad \text{Jensen's inequality} \\ &= \frac{1}{\lambda} \log \left(|\mathcal{S}| \cdot \mathbb{E}\left[\exp\left(\lambda v_i^\top X\right)\right]\right) \qquad v_i \text{s are identical} \\ &= \frac{1}{\lambda} \log \left(|\mathcal{S}| \cdot \left(\frac{Ce^{\lambda} + C}{2C}\right)^T\right) \qquad \text{for multinomial } X, \mathbb{E}\left[e^{t^\top X}\right] = \left(\sum_{i=1}^{2C} p_i e^{t_i}\right)^T \\ &= \frac{1}{\lambda} \log \left(|\mathcal{S}| \cdot e^{\lambda T/2} \cdot \left(\frac{e^{\lambda/2} + e^{-\lambda/2}}{2}\right)^T\right) \\ &\leq \frac{1}{\lambda} \log \left(|\mathcal{S}| \cdot e^{\lambda T/2} \cdot e^{\lambda^2 T/2}\right) \qquad \frac{e^{\lambda/2} + e^{-\lambda/2}}{2} \leq e^{\lambda^2/2}, \lambda \in \mathbb{R} \\ &= \frac{T}{2} + \frac{\log |\mathcal{S}|}{\lambda} + \frac{\lambda T}{2} \end{split}$$

Choosing $\lambda = \sqrt{\frac{2\log|\mathcal{S}|}{T}}$, we get the tightest bound. Substituting the value of λ , we get:

$$\mathbb{E}[M_C(T)] \le \frac{T}{2} + \sqrt{2T\log|\mathcal{S}|}$$

$$= \frac{T}{2} + \sqrt{2T\log\left(\frac{2C}{C}\right)}$$

$$\le \frac{T}{2} + \sqrt{2T\log(2e)^C}$$

$$= \frac{T}{2} + \sqrt{2CT(1 + \log 2)}$$

$$\le \frac{T}{2} + 2\sqrt{CT}$$

$$\binom{n}{k} \le \left(\frac{ne}{k}\right)^k$$

(b) Let us consider the case when C=1. Let the random variable X denote the number of balls in the first bin. Note that $X=\sum_{i=1}^{T}Y_i$ where Y_i is a Bernoulli random variable with p=0.5. So,

$$\mathbb{E}[M_C(T)] = \mathbb{E}[\max(X, T - X)]$$

$$= \frac{T}{2} + \mathbb{E}\left[\max\left(X - \frac{T}{2}, \frac{T}{2} - X\right)\right]$$

$$= \frac{T}{2} + \mathbb{E}\left[\left|X - \frac{T}{2}\right|\right]$$

$$= \frac{T}{2} + \mathbb{E}\left[\left|\sum_{i=1}^{T} Y_i - \frac{T}{2}\right|\right]$$

$$= \frac{T}{2} + \frac{1}{2}\mathbb{E}\left[\left|\sum_{i=1}^{T} (2Y_i - 1)\right|\right]$$

$$= \frac{T}{2} + \frac{1}{2}\mathbb{E}\left[\left|\sum_{i=1}^{T} \epsilon_i\right|\right] \qquad \text{where } \epsilon_i \text{s are i.i.d. Radamacher r.v}$$

$$\stackrel{(a)}{=} \frac{T}{2} + \frac{1}{2} \cdot 2^{1-T} \left[\frac{T}{2}\right] \left(\frac{T}{\lceil \frac{T}{2} \rceil}\right)$$

$$= \frac{T}{2} + \frac{1}{\sqrt{\pi}} \cdot \frac{\Gamma(S + 1/2)}{\Gamma(S)} \qquad \text{where } S = \left[\frac{T}{2}\right], \Gamma(n) = \int_0^{\infty} x^{n-1} e^{-x} dx$$

$$\stackrel{(b)}{=} \frac{T}{2} + \frac{1}{\sqrt{\pi}} \cdot \sqrt{S} \left(1 - \frac{1}{8S} + \frac{1}{128S^2} + \dots\right)$$

$$\stackrel{(c)}{=} \frac{T}{2} + \sqrt{\frac{T}{2\pi}} \left(1 \mp \frac{1}{4T} + \frac{1}{32T^2} \pm \dots\right) \qquad (7)$$

$$\geq \frac{T}{2} + \sqrt{\frac{T}{2\pi}} - \frac{1}{4\sqrt{2\pi T}} \qquad \text{higher order terms are small for } T > 1$$

$$= \frac{T}{2} + \sqrt{\frac{T}{2\pi}} - O\left(\frac{1}{\sqrt{T}}\right) \qquad (8)$$

The expression (a) follows from the expression of expected distance covered in a T-step 1-D random walk³. The step (b) follows from Graham $et\ al.^4$. The expression (c) is obtained by substituting T=2S or T=2S+1 depending upon the parity of T and the top signs are considered for T even and bottom signs for T odd.

Now, let us combine bins $(2k-1,2k), k \in [C]$ and call it a superbin⁵. Let M_i, T_i respectively denote the maximum and the total number of balls in the *i*th superbin. From (8), for $T_i > 0$,

 $^{^3 \}texttt{https://mathworld.wolfram.com/RandomWalk1-Dimensional.html}$

⁴Graham, R. L.; Knuth, D. E.; and Patashnik, O. Answer to problem 9.60 in Concrete Mathematics: A Foundation for Computer Science, 2nd ed. Reading, MA: Addison-Wesley, 1994.

⁵Parts of the subsequent proof closely follows a similar proof in: Bhattacharjee, Rajarshi, Subhankar Banerjee, and Abhishek Sinha. "Fundamental Limits on the Regret of Online Network-Caching." Abstracts of the 2020 SIG-METRICS/Performance Joint International Conference on Measurement and Modeling of Computer Systems. 2020.

we get:

$$\mathbb{E}[M_i|T_i] \ge \frac{T_i}{2} + \sqrt{\frac{T_i}{2\pi}} - \frac{1}{4\sqrt{2\pi T_i}} \tag{9}$$

Note that $M_C(T) \geq \sum_{i=1}^C M_i = CM_1$. The last equality follows as the M_i 's are identical. Furthermore, notice that T_i follows a Binomial distribution with parameters T and 1/C. So, $\mathbb{E}[T_i] = T/C$ and $\operatorname{Var}(T_i) = T\frac{1}{C}\left(1 - \frac{1}{C}\right)$.

$$\mathbb{E}[M_{C}(T)] \geq C\mathbb{E}[M_{1}]$$

$$\stackrel{(a)}{=} C\mathbb{E}[M_{1} \cdot \mathbf{1}(T_{1} > 0)]$$

$$\stackrel{(b)}{=} C\mathbb{E}[\mathbb{E}[M_{1} \cdot \mathbf{1}(T_{1} > 0) | T_{1}]]$$

$$\stackrel{(g)}{\geq} C\mathbb{E}\left[\frac{T_{1} \cdot \mathbf{1}(T_{1} > 0)}{2} + \sqrt{\frac{T_{1} \cdot \mathbf{1}(T_{1} > 0)}{2\pi}} - \frac{\mathbf{1}(T_{1} > 0)}{4\sqrt{2\pi T_{1}}}\right]$$

$$= \underbrace{\frac{C}{2}\mathbb{E}\left[T_{1} \cdot \mathbf{1}(T_{1} > 0)\right]}_{(I)} + \underbrace{\frac{C}{\sqrt{2\pi}}\mathbb{E}\left[\sqrt{T_{1} \cdot \mathbf{1}(T_{1} > 0)}\right]}_{(II)} - \underbrace{\frac{C}{4\sqrt{2\pi}}\mathbb{E}\left[\frac{\mathbf{1}(T_{1} > 0)}{\sqrt{T_{1}}}\right]}_{(III)}$$

The expression (a) follows because if $T_1 = 0 \implies M_1 = 0$. Step (b) describes the law of iterated expectations. As $T_1 \ge 0$,

$$T_{1} = T_{1} \cdot \mathbf{1}(T_{1} = 0) + T_{1} \cdot \mathbf{1}(T_{1} > 0)$$

$$\sqrt{T_{1}} = \sqrt{T_{1}} \cdot \mathbf{1}(T_{1} = 0) + \sqrt{T_{1}} \cdot \mathbf{1}(T_{1} > 0)$$
(10)

From (10), we obtain $\mathbb{E}[T_1] = \mathbb{E}[T_1 \cdot \mathbf{1}(T_1 > 0)]$ and $\mathbb{E}[\sqrt{T_1}] = \mathbb{E}[\sqrt{T_1} \cdot \mathbf{1}(T_1 > 0)]$. So, (I) becomes:

$$\frac{C}{2}\mathbb{E}\left[T_1 \cdot \mathbf{1}(T_1 > 0)\right] = \frac{C}{2}\mathbb{E}\left[T_1\right] = \frac{T}{2}$$
(11)

Consider,

$$\sqrt{x} - \left(1 + \frac{x-1}{2} - \frac{(x-1)^2}{2}\right) = \frac{1}{2}\sqrt{x}(\sqrt{x} - 1)^2(\sqrt{x} + 2) \ge 0$$

$$\implies \sqrt{x} \ge 1 + \frac{x-1}{2} - \frac{(x-1)^2}{2}$$

Let $x = \frac{X}{\mathbb{E}[X]}$, we have almost surely:

$$\begin{split} \sqrt{\frac{X}{\mathbb{E}[X]}} &\geq 1 + \frac{\frac{X}{\mathbb{E}[X]} - 1}{2} - \frac{(\frac{X}{\mathbb{E}[X]} - 1)^2}{2} \\ \Longrightarrow & \mathbb{E}\left[\sqrt{\frac{X}{\mathbb{E}[X]}}\right] \geq \mathbb{E}\left[1 + \frac{\frac{X}{\mathbb{E}[X]} - 1}{2} - \frac{(\frac{X}{\mathbb{E}[X]} - 1)^2}{2}\right] \\ \Longrightarrow & \frac{\mathbb{E}[\sqrt{X}]}{\sqrt{\mathbb{E}[X]}} \geq 1 - \frac{\mathrm{Var}(X)}{2(\mathbb{E}[X])^2} \end{split}$$

$$\implies \mathbb{E}[\sqrt{T_1}] \ge \sqrt{\mathbb{E}[T_1]} \left(1 - \frac{\operatorname{Var}(T_1)}{2(\mathbb{E}[T_1])^2} \right)$$

$$= \sqrt{\frac{T}{C}} \left(1 - \frac{T/C(1 - 1/C)}{2(T/C)^2} \right)$$

$$= \sqrt{\frac{T}{C}} \left(1 - \frac{C - 1}{2T} \right)$$

$$\ge \sqrt{\frac{T}{C}} - \frac{1}{2} \sqrt{\frac{C}{T}}$$

So, (II) becomes:

$$\frac{C}{\sqrt{2\pi}} \mathbb{E}\left[\sqrt{T_1 \cdot \mathbf{1}(T_1 > 0)}\right] = \frac{C}{\sqrt{2\pi}} \mathbb{E}\left[\sqrt{T_1}\right] \ge \sqrt{\frac{CT}{2\pi}} - \frac{C^{3/2}}{2\sqrt{2\pi T}}$$
(12)

$$\mathbb{E}\left[\frac{\mathbf{1}(T_1>0)}{\sqrt{T_1}}\right] \stackrel{(1)}{=} \int_0^\infty P\left(\frac{\mathbf{1}(T_1>0)}{\sqrt{T_1}}>t\right) dt$$

$$= \int_0^1 P\left(\frac{\mathbf{1}(T_1>0)}{\sqrt{T_1}}>t\right) dt + \int_{1^+}^\infty P\left(\frac{\mathbf{1}(T_1>0)}{\sqrt{T_1}}>t\right) dt \qquad (13)$$

Consider the second term of the RHS. Since, $T_1 \sim Binom(T, 1/C) \implies T_1 \in \{0, 1, \dots, T\}$. If $T_1 = 0 \implies \frac{\mathbf{1}(T_1 > 0)}{\sqrt{T_1}} = 0 \implies P(0 > t) = 0$ almost surely. Now, if $T_1 \ge 1$, then $\frac{1}{\sqrt{T_1}} \le 1 \implies P\left(\frac{\mathbf{1}(T_1 > 0)}{\sqrt{T_1}} > t\right) = 0$ almost surely for $t \in (1, \infty)$. So, (13) becomes,

$$\mathbb{E}\left[\frac{\mathbf{1}(T_1 > 0)}{\sqrt{T_1}}\right] = \int_0^1 P\left(\frac{\mathbf{1}(T_1 > 0)}{\sqrt{T_1}} > t\right) dt$$
$$\leq \int_0^1 dt = 1$$

So, (III) becomes,

$$\frac{C}{4\sqrt{2\pi}}\mathbb{E}\left[\frac{\mathbf{1}(T_1>0)}{\sqrt{T_1}}\right] \le \frac{C}{4\sqrt{2\pi}} \implies -\frac{C}{4\sqrt{2\pi}}\mathbb{E}\left[\frac{\mathbf{1}(T_1>0)}{\sqrt{T_1}}\right] \ge -\frac{C}{4\sqrt{2\pi}} \tag{14}$$

Combining (11),(12) and (14), we get:

$$\mathbb{E}[M_C(T)] \ge \frac{T}{2} + \sqrt{\frac{CT}{2\pi}} - \frac{C^{3/2}}{2\sqrt{2\pi T}} - \frac{C}{4\sqrt{2\pi}}$$

$$\implies \mathbb{E}[M_C(T)] \ge \frac{T}{2} + \sqrt{\frac{CT}{2\pi}} - O\left(\frac{1}{\sqrt{T}}\right)$$

Let $s_n = \max\left(\frac{z_n}{n}, 1 - \frac{z_n}{n}\right)$, where z_n is the number of zeros in the sequence be score assigned to a length-n binary sequence. Clearly, $0.5 \le s_n \le 1$, and lower bound equality holds only when the sequence is uninformative, i.e., number of zeros and ones are the same. So, we design our loss function as:

$$\phi_n = 1 - s_n = 1 - \max\left(\frac{z_n}{n}, 1 - \frac{z_n}{n}\right) \tag{15}$$

Consider a sequence $\mathbf{y} = (y_1, y_2, \dots, y_n)$ and $\mathbf{y'} = (y'_1, y'_2, \dots, y'_n)$ obtained by flipping exactly one bit. So,

$$\left|\phi_n(\boldsymbol{y}) - \phi_n(\boldsymbol{y'})\right| = \left|\max\left(\frac{z_n(\boldsymbol{y})}{n}, 1 - \frac{z_n(\boldsymbol{y})}{n}\right) - \max\left(\frac{z_n(\boldsymbol{y'})}{n}, 1 - \frac{z_n(\boldsymbol{y'})}{n}\right)\right| \le \frac{1}{n}$$

So, the candidate ϕ_n is smooth. Consider the case when k = 1. As it is known apriori that there is an imbalance of zeros and ones in the sequence, $s_n > 0.5 \implies \phi_n < 0.5$.

$$\phi_n = 1 - \max\left(\frac{z_n}{n}, 1 - \frac{z_n}{n}\right)$$

$$\implies \mathbb{E}\left[\phi_n\right] = \mathbb{E}\left[1 - \max\left(\frac{z_n}{n}, 1 - \frac{z_n}{n}\right)\right]$$

$$= 1 - \mathbb{E}\left[\max\left(\frac{z_n}{n}, 1 - \frac{z_n}{n}\right)\right]$$

$$= 1 - \mathbb{E}\left[\frac{1}{2} + \max\left(\frac{z_n}{n} - \frac{1}{2}, \frac{1}{2} - \frac{z_n}{n}\right)\right]$$

$$= \frac{1}{2} - \mathbb{E}\left[\left|\frac{z_n}{n} - \frac{1}{2}\right|\right] < \frac{1}{2}$$

As $\mathbb{E}[\phi_n] < 1/2$, by Cover's argument, we can say that there does not exist an algorithm that realises this loss function. So, if add an residual term and design a new loss function ϕ'_n as described in (16), we can prove that exists an algorithm that serves our purpose. Furthermore, it can be easily verified that $\mathbb{E}[\phi'_n] = 1/2$.

$$\phi_n' = \phi_n + \mathbb{E}\left[\left|\frac{z_n}{n} - \frac{1}{2}\right|\right] \tag{16}$$

Now, we extend for the case k > 1. Let x_1, x_2, \ldots, x_k be the k-partitions, $L(x_1), L(x_2), \ldots, L(x_n)$ be the respective lengths of the partitions and $z(x_1), z(x_2), \ldots, z(x_n)$ be the number of zeros in the respective partitions. Let Φ_n be defined as follows:

$$\Phi_n = \max_{x_1, \dots, x_n} \frac{1}{k} \sum_{i=1}^k \phi_{L(x_i)}$$
 (17)

The max refers to maximum over all possible k-partitions of the sequence. Following the steps as earlier, we can show that Φ_n is also smooth.

$$\begin{split} \Phi_n &= \max_{x_1, \dots, x_n} \frac{1}{k} \sum_{i=1}^k \phi_{L(x_i)} \\ \implies \mathbb{E}[\Phi_n] &= \frac{1}{k} \mathbb{E} \left[\max_{x_1, \dots, x_n} \sum_{i=1}^k \phi_{L(x_i)} \right] \\ &= \frac{1}{k} \mathbb{E} \left[\max_{x_1, \dots, x_n} \sum_{i=1}^k \left(1 - \max \left(\frac{z(x_i)}{L(x_i)}, 1 - \frac{z(x_i)}{L(x_i)} \right) \right) \right] \\ &= \frac{1}{k} \mathbb{E} \left[\max_{x_1, \dots, x_n} \sum_{i=1}^k \left(\frac{1}{2} - \max \left(\frac{z(x_i)}{L(x_i)} - \frac{1}{2}, \frac{1}{2} - \frac{z(x_i)}{L(x_i)} \right) \right) \right] \\ &= \frac{1}{k} \mathbb{E} \left[\max_{x_1, \dots, x_n} \sum_{i=1}^k \left(\frac{1}{2} - \left| \frac{z(x_i)}{L(x_i)} - \frac{1}{2} \right| \right) \right] \\ &= \frac{1}{2} + \frac{1}{k} \mathbb{E} \left[\max_{x_1, \dots, x_n} \sum_{i=1}^k - \left| \frac{z(x_i)}{L(x_i)} - \frac{1}{2} \right| \right] \\ &= \frac{1}{2} - \frac{1}{k} \mathbb{E} \left[\min_{x_1, \dots, x_n} \sum_{i=1}^k \left| \frac{z(x_i)}{L(x_i)} - \frac{1}{2} \right| \right] < \frac{1}{2} \end{split}$$

The last inequality follows because one can always choose a k-partition such that there is an imbalance in at least one of the partitions, so the second term > 0. As $\mathbb{E}[\Phi_n] < 1/2$, there is no algorithm that achieves the said loss function. So, we devise a modified loss function Φ'_n in (18) which satisfies $\mathbb{E}[\Phi'_n] = 1/2$.

$$\Phi'_{n} = \Phi_{n} + \frac{1}{k} \mathbb{E} \left[\min_{x_{1}, \dots, x_{n}} \sum_{i=1}^{k} \left| \frac{z(x_{i})}{L(x_{i})} - \frac{1}{2} \right| \right]$$
(18)