

Time Complexity:

Big-Oh Notation:

$g(n) = O(f(n))$  if  $g(n) \leq c \cdot f(n)$

Order:

$O(1) < O(\log n) < O(\sqrt{n}) < O(n) < O(n \log n) < O(n^2) < O(n^3) < \dots < O(2^n) < O(n^n)$

1.

```
int n, k=0;
cin>>n;
for(int i=1; i<=n; i++){      //O(n)
    k+=i; //O(1) - > Constant time operation
}
```

2.

```
int n, k=0;
cin>>n;
for(int i=1; i<=n; i++){      //O(n)
    for(int j=0; j<i; j++){ //O(i)
        k++;
    }
} //1 + 2 + 3 + .. + n = O(n^2)
```

3.

```
for(int i=1; i*i<=n; i++){    //O(sqrt(n))
    k+=i; // O(1)
} //Overall O(sqrt(n))
```

4.

```
for(int i=1; i<=n; i++){
    for(int j=0; j<i; j++){
        k++;
    }
}
for(int i=0; i<n; i++){
    q+=i;
}
```

Overall  $\rightarrow O(n^2) \rightarrow O(n^2 + n)$  but  $n$  is ignored w.r.t  $n^2$

5.

```
for(int i=1; i<=n; i++){
    for(int j=0; j<i; j++){
        k++;
    }
}
```

```
for(int i=0; i<m; i++){
    q+=i;
}
```

Overall  $\rightarrow O(n^2 + m) \rightarrow O(\max(n^2, m))$  is technically correct but not conventional to use

6.

```
for(int i=n; i>1; i/=2){
    for(int j=0; j<i; j++){
        res++;
    }
}
```

Overall  $\rightarrow n + (n/2) + (n/4) + \dots + 1 = 2n = O(n)$

$O(n \log n)$  is technically correct but it is possible to have a tighter upperbound of  $O(n)$

7.

```
for(int i=1; i<=sqrt(n); i++){
    res *= i;
}
```

Overall  $\rightarrow O(\text{root}(n) * \log(n))$  because  $\text{sqrt}()$  is called  $\text{root}(n)$  times and time complexity of  $\text{sqrt}()$  is  $O(\log(n))$