# CS & IT ENGINEERING

GRAPH THEORY/ LOGIC

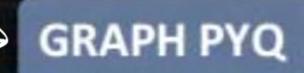
Lecture No: 15



Satish Sir



TOPICS TO BE COVERED



Logic pyq



#### A non-planar graph with minimum number of vertices has (GATE - 92)



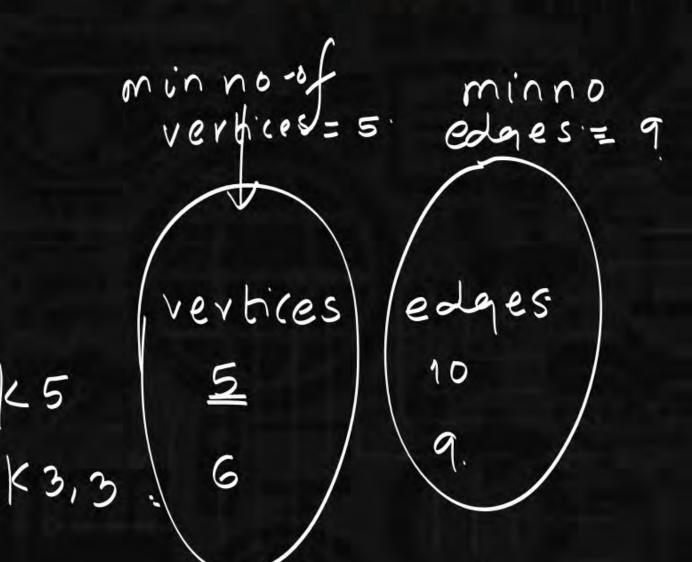
$$\frac{planav}{n=2}$$

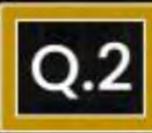
$$n=3$$

$$(102)$$

$$n=4$$

- A. 9 edges, 6 vertices (43/3) n = 4
- B. 6 edges, 4 vertices / nonplanar
- 10 edges, 5 vertices (non planar)
- 9 edges, 5 vertices planow)

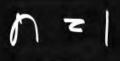




## The number of non-isomorphic simple graphs up to three nodes is (GATE - 94)



- A. 15
- B. 10
- 7 Ans
- D. 9









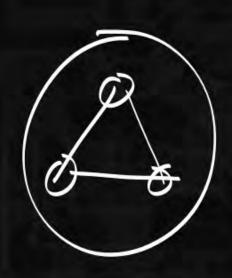
$$n=2$$

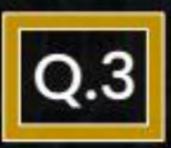












## Maximum number of edges in a n - node undirected graph without self loops is (GATE - 02)

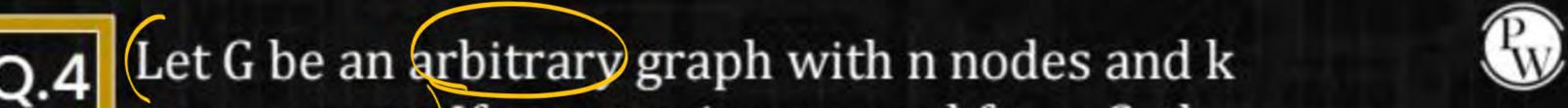


A. 
$$n^2$$

B. 
$$\frac{n(n-1)}{2}$$

$$n-1$$

D. 
$$\frac{(n+1)(n)}{2}$$



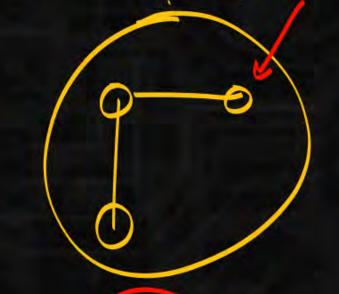


Let G be an arbitrary graph with n nodes and k components. If a vertex is removed from G, the number of components in the resultant graph must

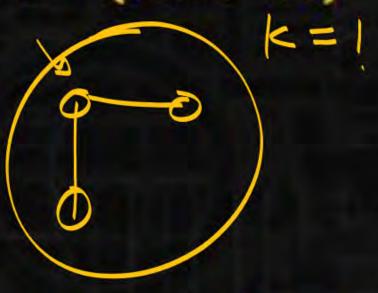
K=1 Remove (GATE-03) necessarily lie between

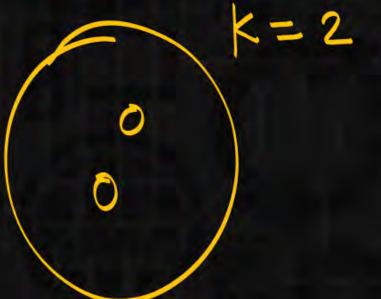


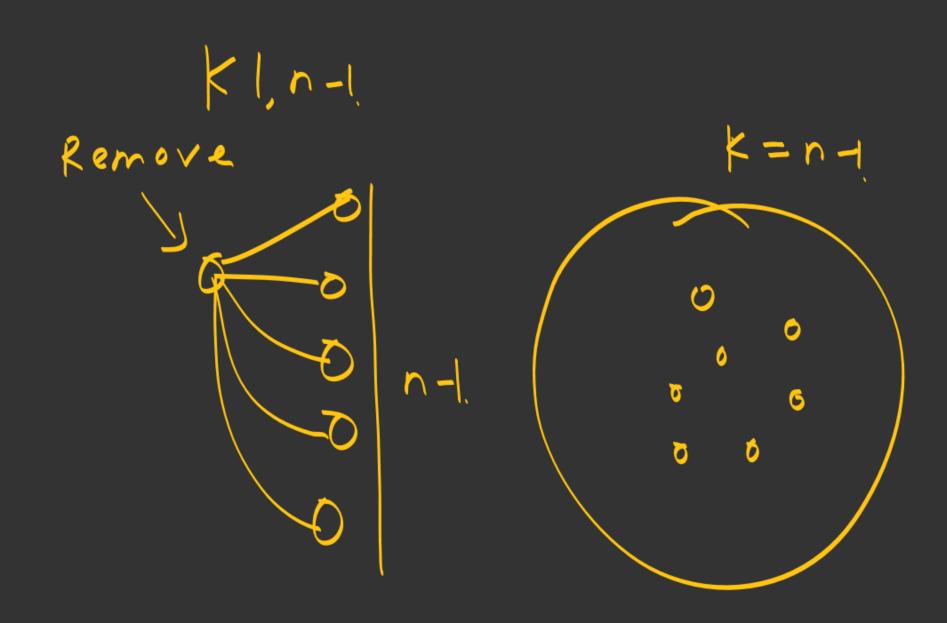
- k and n
- k-1 and k+1
- k-1 and n-1
- k+1 and n-k







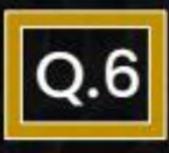




If all the edge weights of an undirected graph are positive, then any subset of edges that connects all the vertices and has minimum total weight is a

(GATE - 06)

- A. Hamiltonian cycle
- B. grid
- c. hypercube
- D. tree

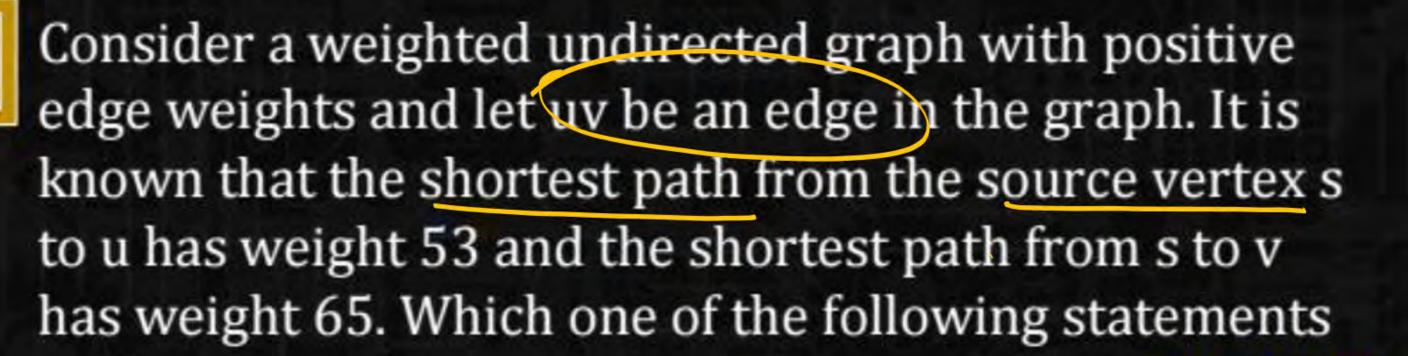


## Let G be the non-planar graph with minimum possible number of edges. Then G has



(GATE - 07)

- A. 9 edges and 5 vertices
- B. 9 edges and 6 vertices
- c. 10 edges and 5 vertices
- D. 10 edges and 6 vertices



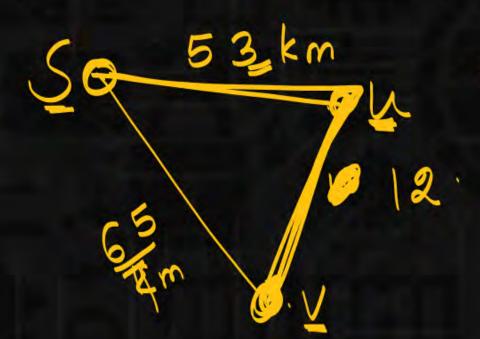
A. weight (u,v)<12

is always true?

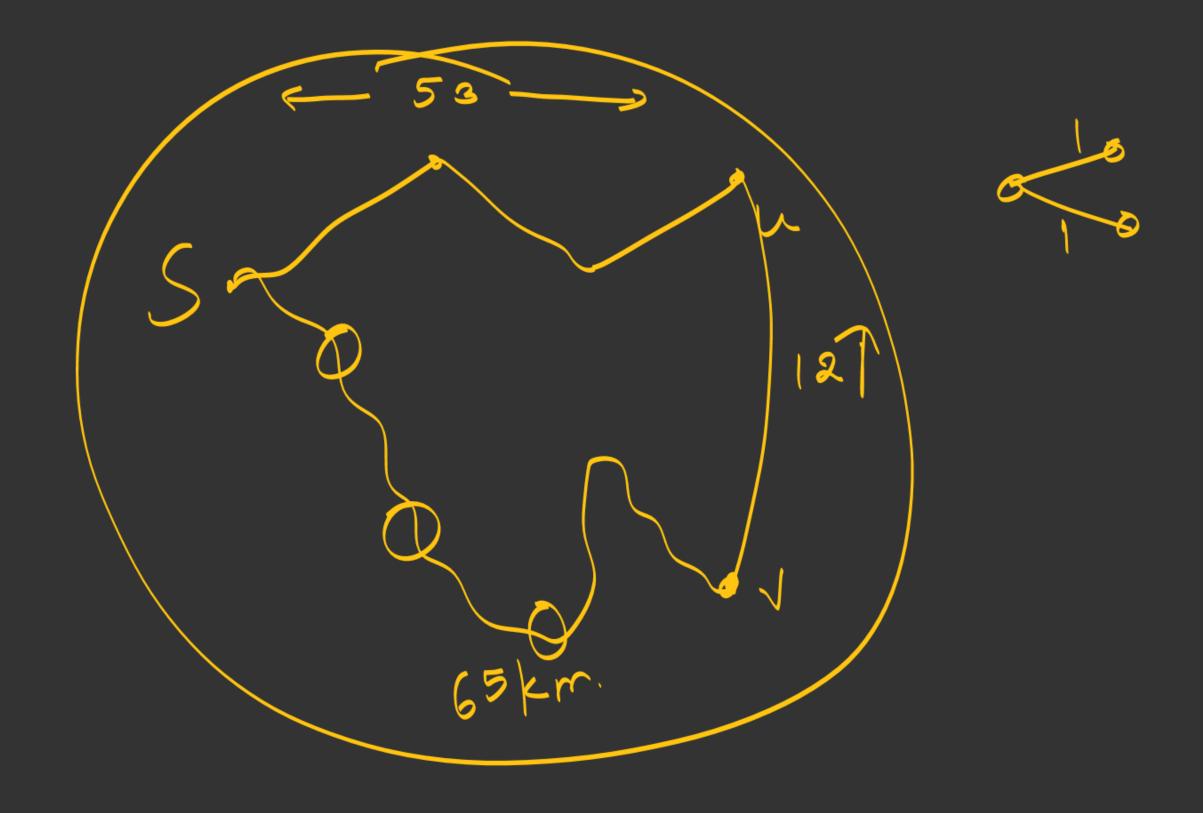
B. weight(u,v)≤12

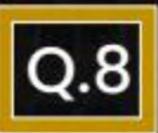
c. weight (u,v)>12

D. weight(u,v)≥12 √



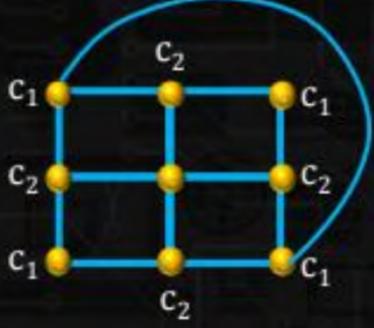






What is the chromatic number of the following graph?





(GATE - 08)

A. 2

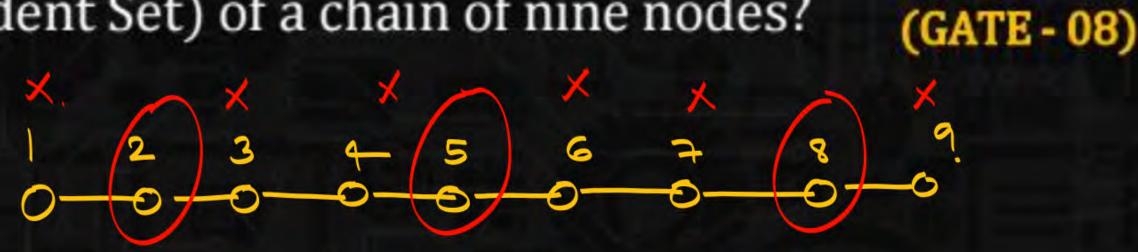
B. 3 V

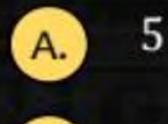
C. 4

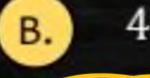
D. 5

What is the size of the smallest MIS (Maximal Independent Set) of a chain of nine nodes?















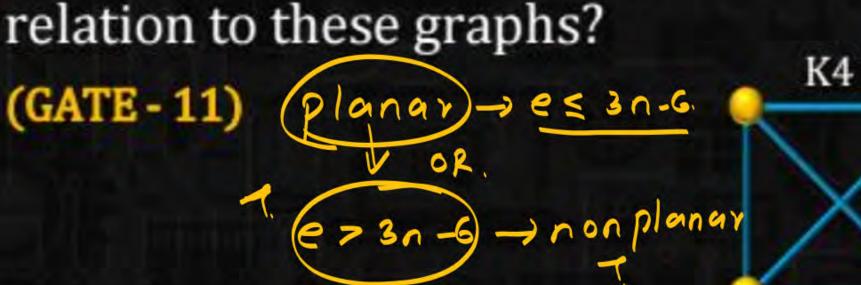
What is the chromatic number of an n-vertex simple connected graph which does not contain any odd TE - 09) length cycle? Assume  $n \ge 2$ .

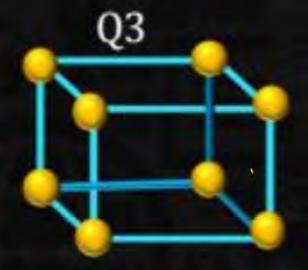


- A. 2
- B. 3
- c. n-1
- D. n

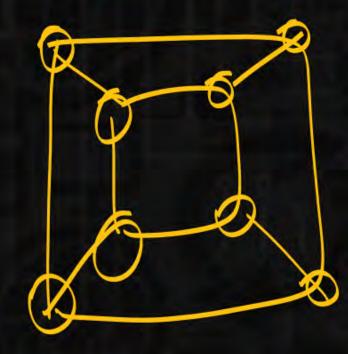
K4 and Q3 are graphs with the following structures. Which one of the following statements is TRUE in

planar -> e < 3n-6.





- A. K4 is planar while Q3 is not
- B. Both K4 and Q3 are planar
- C. Q3 is planar while K3 is not
- D. Neither K4 nor Q3 is planar





Let G be a simple undirected planar graph on 10 vertices with 15 edges. If G is a connected graph, then the number of bounded faces in any embedding of G on the plane is equal to (GATE - 12)



A. 3

B. 4

C. 5

D. 6



Which of the following statements is/are TRUE for undirected graphs?



P: Number of odd degree vertices is even.

Q: Sum of degrees of all vertices is even.

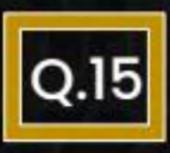
(GATE - 13)

- A. Ponly
- B. Q only
- C. Both P and Q
- D. Neither P nor Q



#### Q.14 The maximum number of edges in a bipartite graph on 12 vertices is GATE-14-Set2





Let G be an undirected complete graph on n vertices, where n>2. Then, the number of different Hamiltonian cycles in G is equal to



(GATE - 19)

$$\frac{(n-1)!}{2}$$

The minimum number of colours required to colour the following graph, such that no two adjacent vertices are assigned the same colour, is



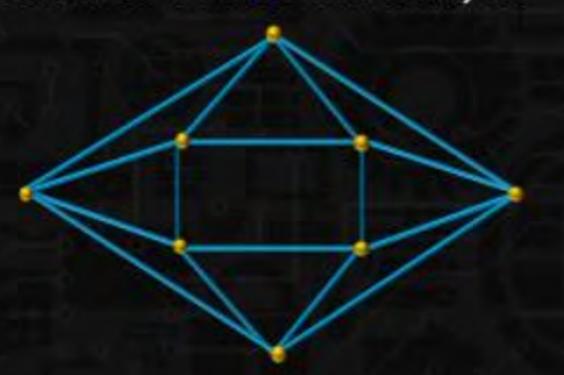
(GATE - 04)

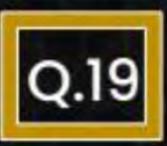
A. 2

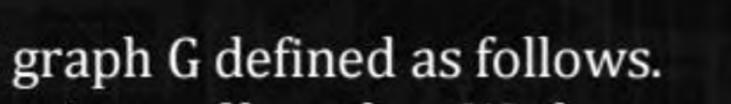
B. 3

c. 4

D.









(GATE - 04)

Consider the undirected graph G defined as follows. The vertices of G are bit strings of length n. We have an edge between vertex u and vertex v if and only if u and v differ in exactly one bit position (in other words, v can be obtained from u by flipping a single bit). The ratio of the chromatic number of G to the

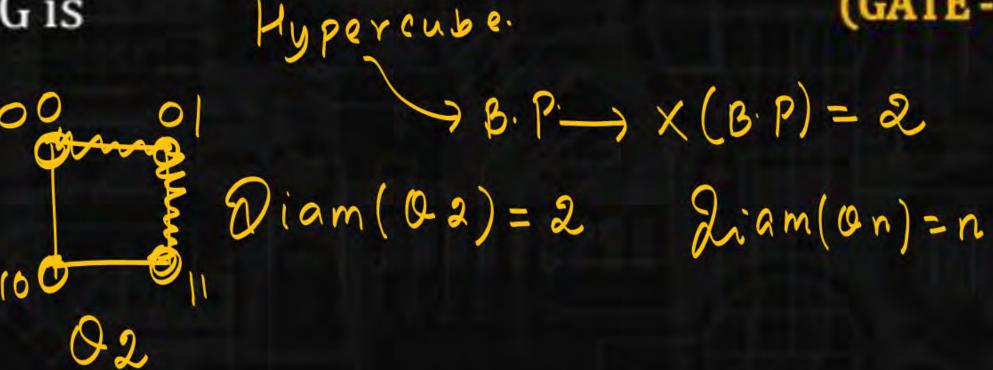
diameter of G is

 $1/2^{n-1}$ 

1/n



3/n



G is a simple undirected graph. Some vertices of G are of odd degree. Add a node v to G and make it adjacent to each odd degree vertex of G. The resultant graph is

sure to be

(GATE - 08)

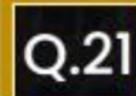
no of odd degr

Regular

complete

Hamiltonian





The degree sequence of a simple graph is the sequence of the degrees of the nodes in the graph in decreasing order. Which of the following sequences cannot be the degree sequence of any graph?

PW

I. 7,6,5,4,4,3,2,1

II. 6,6,6,6,3,3,2,2

III. 7,6,6,4,4,3,2,2

IV. 8,7,7,6,4,2,1,1

A. I and II

B. III and IV

c. IV only

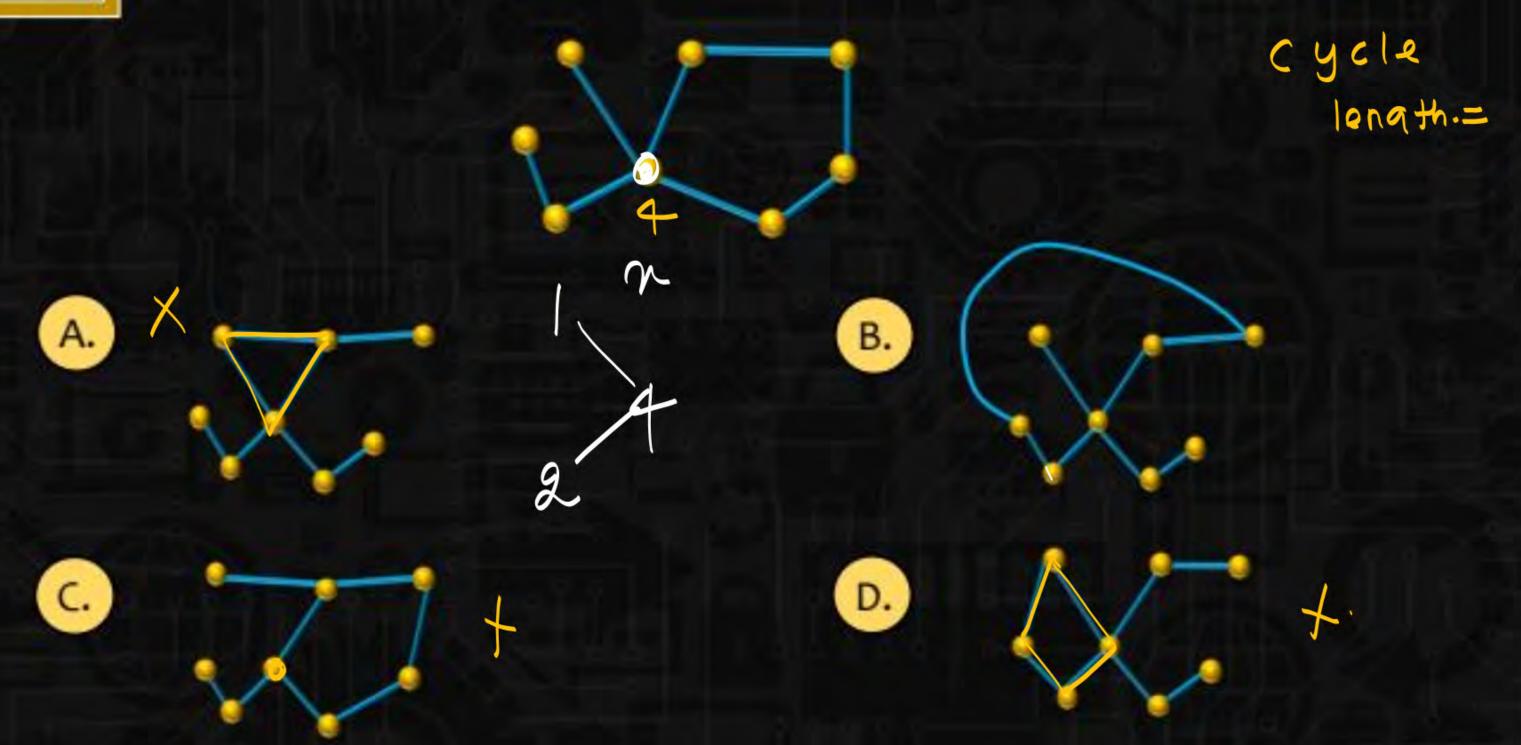
D. II and IV

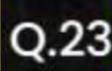
(GATE - 10)

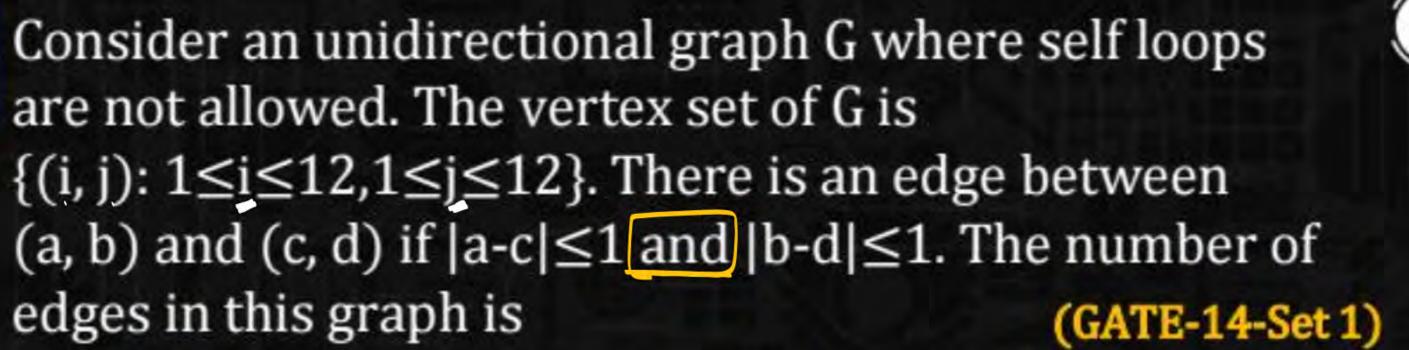
#### Which of the following graph is isomorphic to

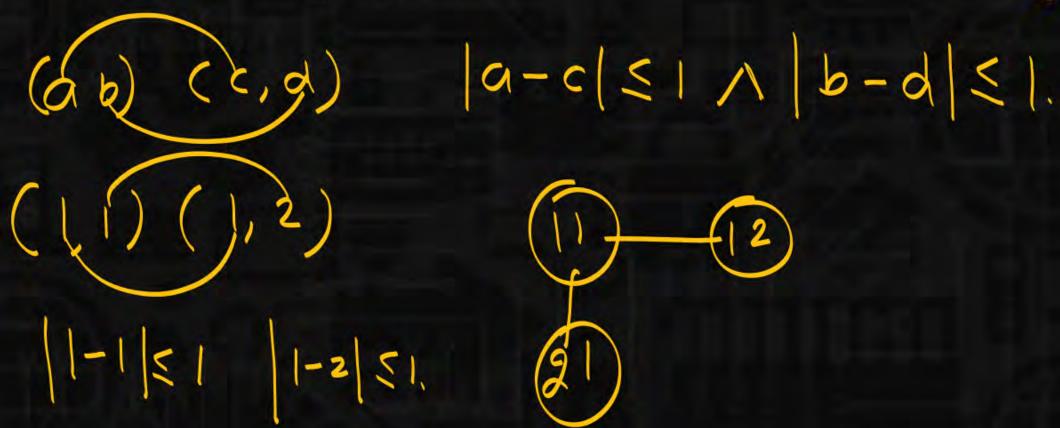
(GATE - 12)

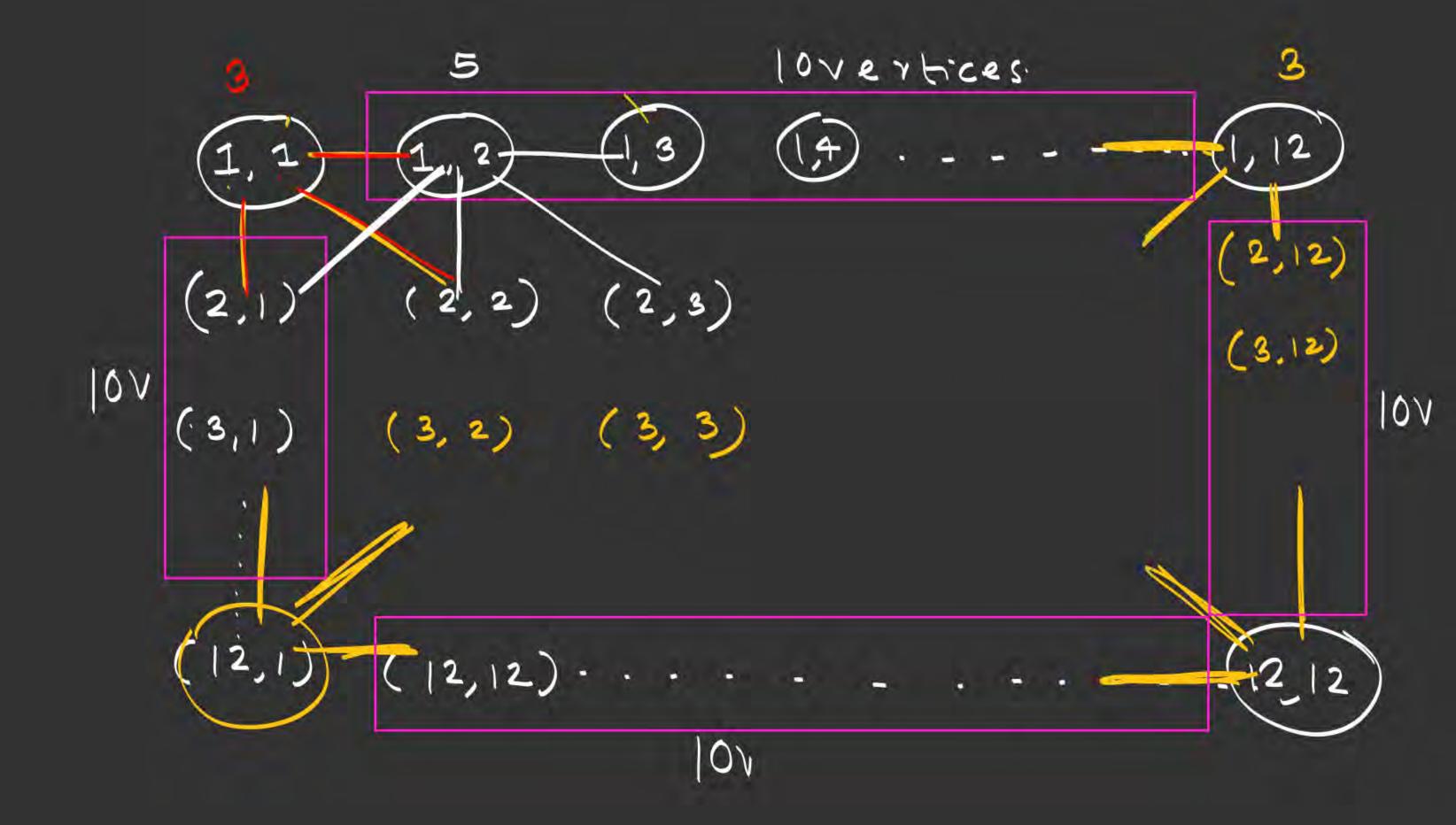


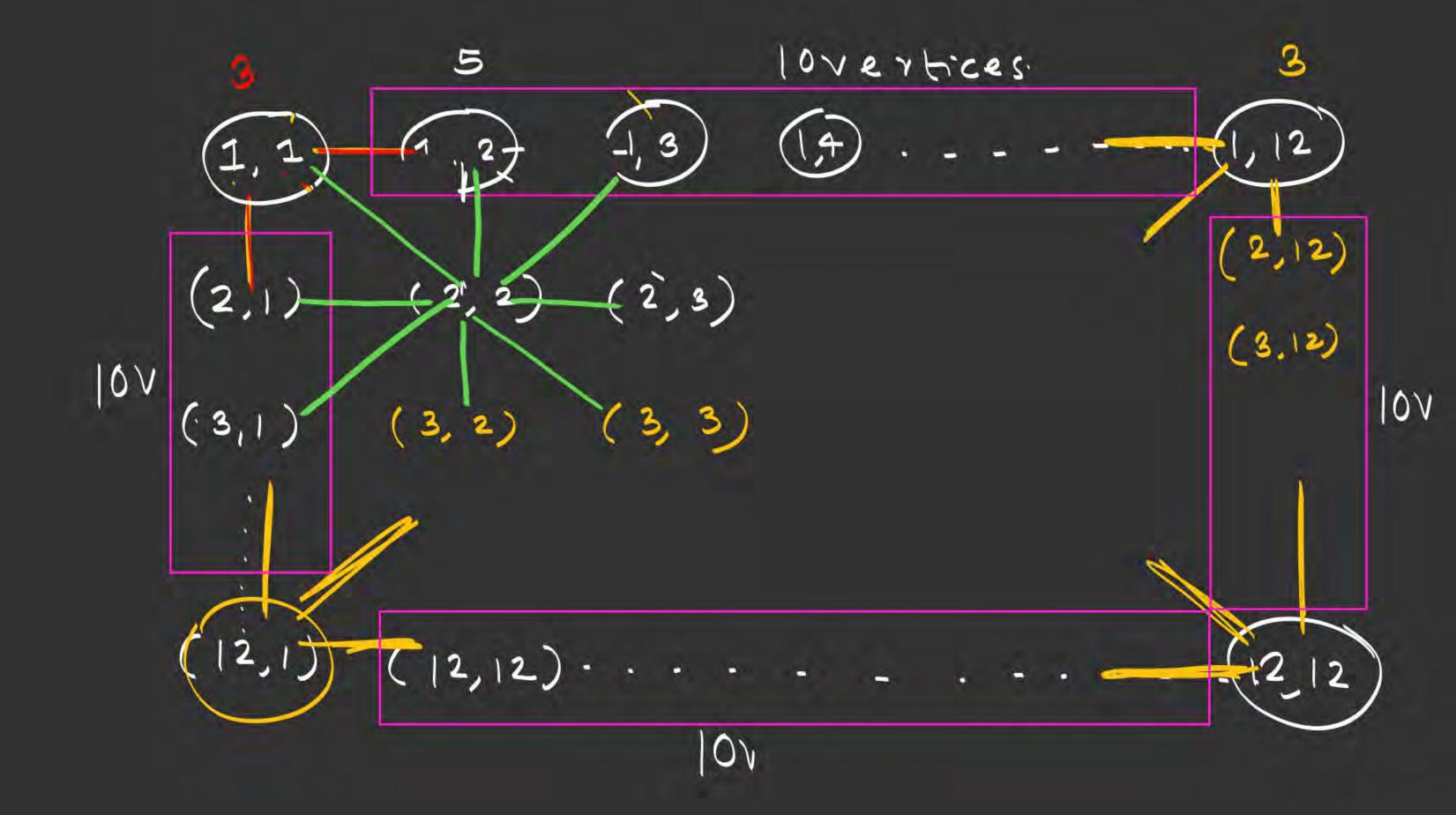








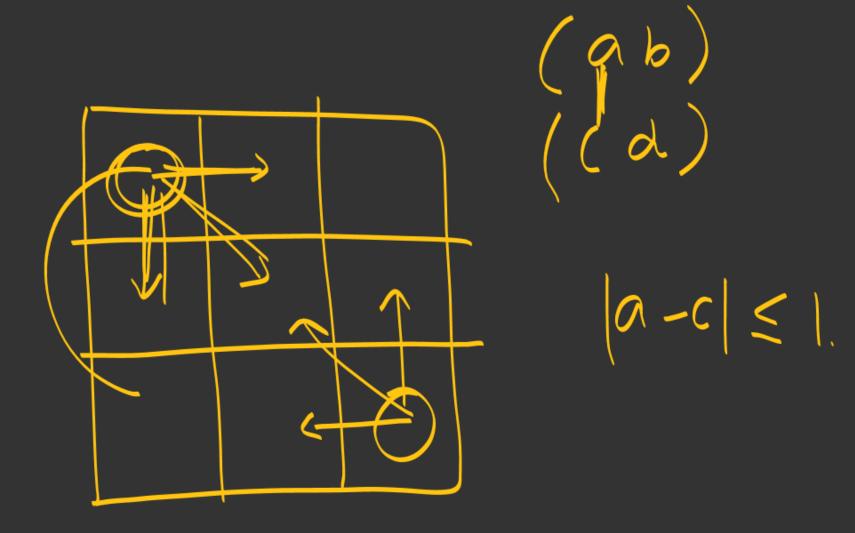


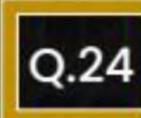


## $4v \times 3 + 40v \times 5 + 100 \times 8 = 2e$

Total vertices = 144.



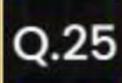


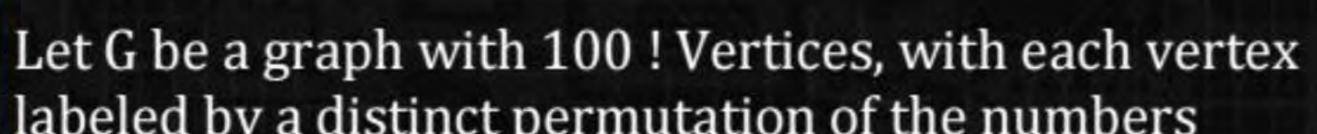




An ordered n-tuple  $(d_1, d_2, ..., d_n)$  with  $d_1 \ge d_2 \ge ... \ge d_n$  is called graphic if there exists a simple undirected graph with n vertices having degrees  $d_1 \ge d_2 \ge ... \ge d_n$  respectively. Which of the following 6-tuples is NOT graphic? (GATE-14-Set 1)

- A. (1,1,1,1,1,1)
- B. (2,2,2,2,2,2)
- (3,3,3,1,0,0)
- D. (3,2,1,1,1,0)

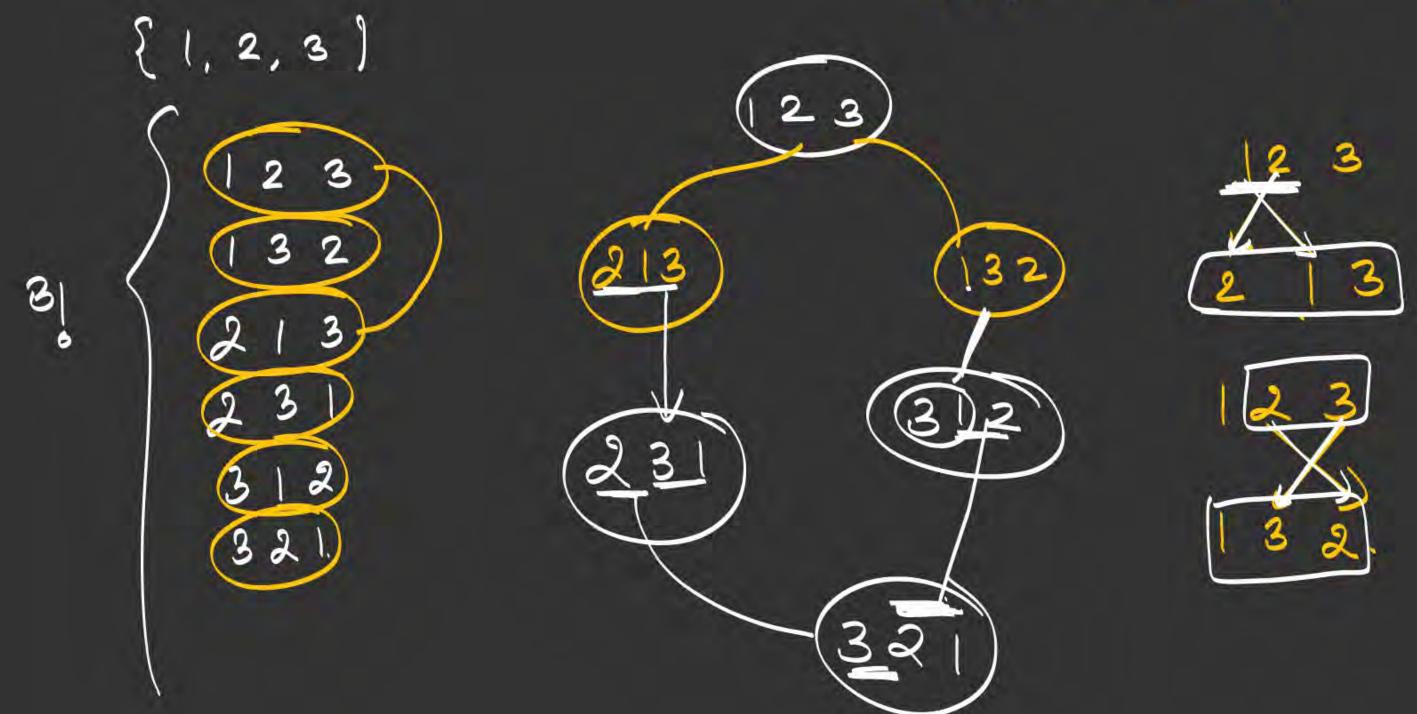


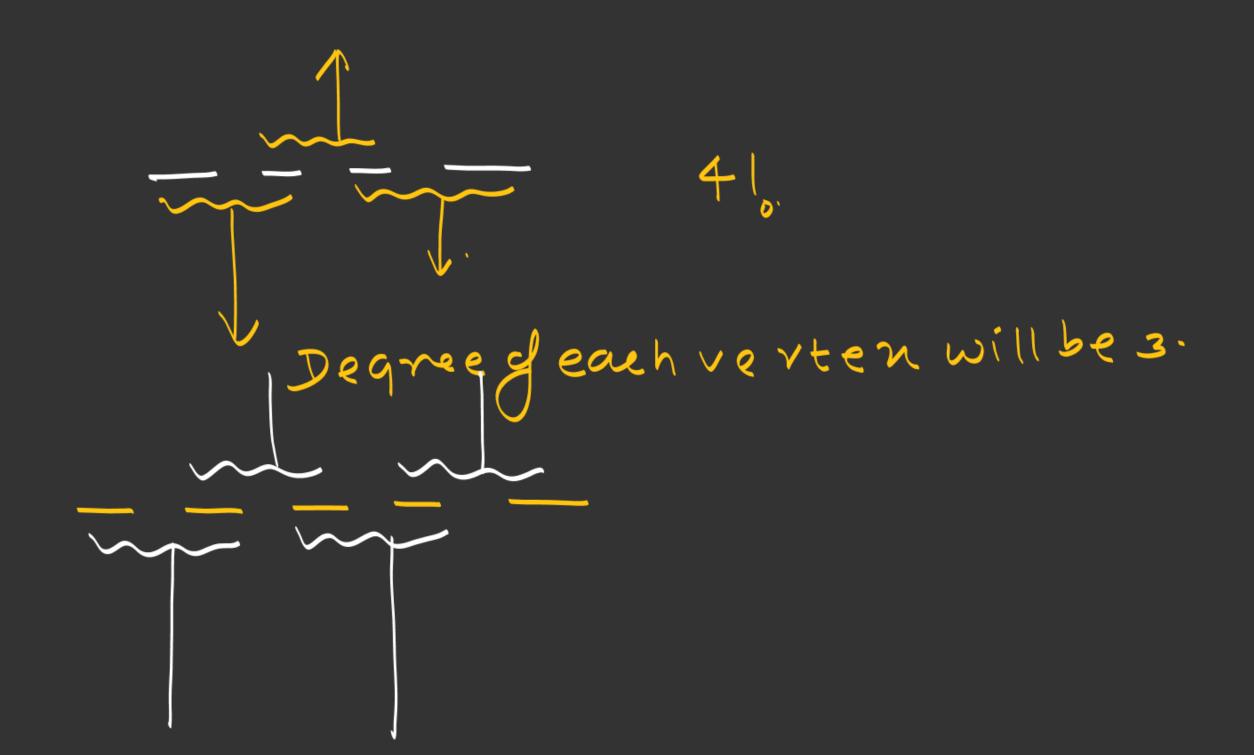


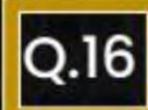


labeled by a distinct permutation of the numbers 1,2,....,100. There is an edge between vertices u and v if and only if the label of u can be obtained by swapping two adjacent numbers in the label of v. Let y denote the degree of a vertex in G, and z denote the number of connected components in G. Then, y+10z=

#### adjacent swap.







Q.16 Let p,q,r,s represent the following propositions.



p:  $x \in \{8, 9, 10, 11, 12\}$ 

q:x is a composite number 4, 6, 8 r. x is a perfect square 8 is perfect.

s:x is a prime number 8 is mi (GATE-16-set 1)

The integer  $x \ge 2$  which satisfies

$$\sim$$
 ((p $\Rightarrow$ q) $\wedge$ ( $\sim$ rV $\sim$ s)) is



Let p,q and r be propositions and the expression  $(p\rightarrow q)\rightarrow r$  be a contradiction. Then, the expression  $(r\rightarrow p)\rightarrow q$  is



(GATE-17-Set 1)

- A. a tautology
- B. a contradiction
- c. always TRUE when p is FALSE
- D. always TRUE when q is TRUE



Let p,q,r denote the statements "It is raining", "It is cold" and "It is pleasant", respectively. Then the statement "It is not raining and it is pleasant, and it is not pleasant only if it is raining and it is cold" is represented by

(GATE-17-Set2)



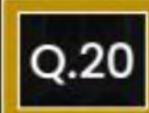
- A.  $(\sim p \land r) \land (\sim r \rightarrow (p \land q))$
- B.  $(\sim p \wedge r) \wedge ((p \wedge q) \rightarrow \sim r)$
- c.  $(\sim p \land r) \lor ((p \land q) \rightarrow \sim 1)$
- D.  $(\sim p \land r) \lor (r \rightarrow (p \land q)) +$

## Which of the following predicate calculus statements is/are valid? (GATE - 92)



$$Au(b(u) \wedge \sigma(u)) \leftarrow Aub(u) \wedge$$

- A.  $((\forall x)P(x)V(\forall x)Q(x))\rightarrow(\forall x)\{P(x)VQ(x)\}$
- B.  $\{(\exists x)P(x)\Lambda(\exists x)Q(x)\}\rightarrow(\exists x)\{P(x)\Lambda Q(x)\}\$
- (C)  $(\forall x)\{P(x)\lor Q(x)\}\rightarrow \{(\forall x)P(x)\lor (\forall x)Q(x)\} \downarrow$
- D.  $(\exists x)\{P(x)VQ(x)\}\rightarrow \sim \{(\forall x)P(x)V(\exists x)Q(x)\}$



Identify the correct translation into logical notation of the following assertion.



Some boys in the class are taller than all the girls Note: taller (x, y) is true if x is taller than y.

(GATE - 04)

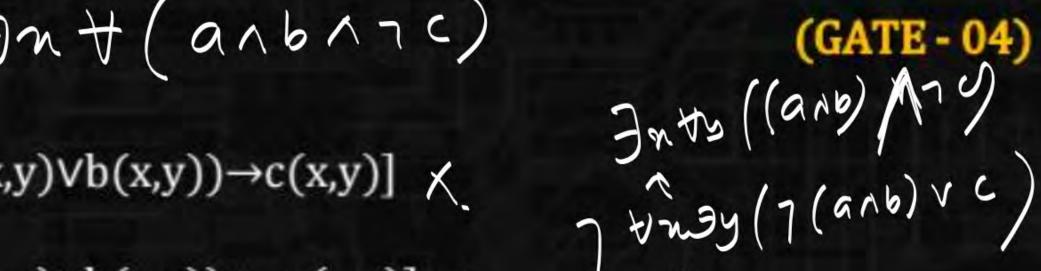
- A.  $(\exists x)(boy(x)\rightarrow(\forall y)(girl(y)\land taller(x, y)))$
- B.  $(\exists x)(boy(x) \land (\forall y)(girl(y) \land taller(x, y)))$
- ( $\exists x$ )(boy(x) $\rightarrow$ ( $\forall y$ )(girl(y) $\rightarrow$  taller(x, y)))
- D.  $(\exists x)(boy(x) \land (\forall y)(girl(y) \rightarrow taller(x, y)))$



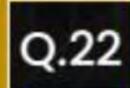
Consider the following statement  $(\exists x)(\forall y)[(a(x,y)\land b(x,y))\land \sim c(x,y)]$ 

Which one of the following is its equivalent?





- $(\forall x)(\exists y)[(a(x,y)\lor b(x,y))\rightarrow c(x,y)]$
- $(\exists x)(\forall y)[(a(x,y)\lor b(x,y))\land \sim c(x,y)]$
- 7 Angy (anb >c)  $\sim [(\forall x)(\exists y)[(a(x,y)\land b(x,y))\rightarrow c(x,y)]$
- $\sim [(\forall x)(\exists y)[(a(x,y)\lor b(x,y))\rightarrow c(x,y)]$



Which one of the following is the most appropriate logical formula to represent the statement:



"Gold and silver ornaments are precious"

The following notations are used:

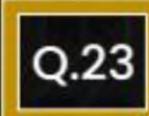
G(x):x is a gold ornament.

S(x):x is a silver ornament.

P(x):x is precious.

TNT

- (a)  $\forall x(P(x)\rightarrow(G(x)\land S(x)))$
- (b)  $\forall x((G(x)\land S(x))\rightarrow P(x))$
- (c)  $\exists x((G(x)\land S(x))\rightarrow P(x))$ (d)  $\forall x((G(x)\lor S(x))\rightarrow P(x))$



## Consider the following well-formed formulae:



I.  $\sim \forall x(P(x))$ 

II.  $\sim \exists x (P(x))$ 

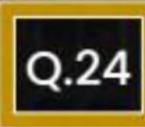
III.  $\sim \exists x (\sim P(x))$ 

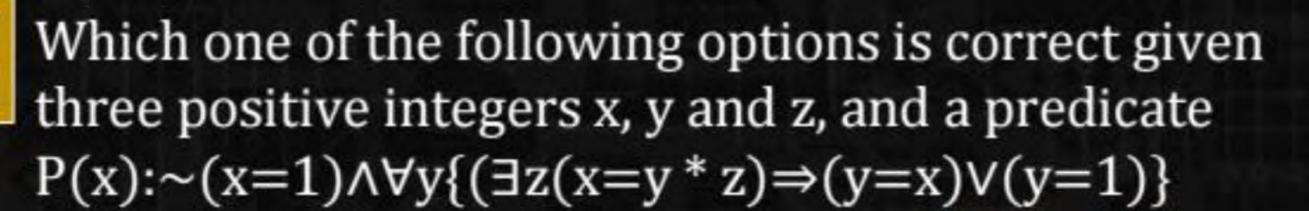
IV.  $\exists x (\sim P(x))$ 

Which of the above are equivalent?

(GATE - 09)

- A. I and II
- B. II and III
- c. I and IV
- D. II and IV

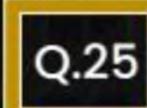




(GATE - 11)



- A. P(x) being true means that x is a prime number
- B. P(x) being true means that x is a number other than 1
- P(x) is always true irrespective of the value of x
- D. P(x) being true means that x has exactly two factors other than 1 and x



## Consider the statement:



"Not all that glitters is gold"
Predicate glitters (x) is true if x glitters and predicate

gold(x) is true if x is gold.

Which one of the following logical formulae

represents the above statement? (GATE-14-Set1)

A.  $\forall x : glitters(x) \Rightarrow \sim gold(x)$ 

B.  $\forall x : gold(x) \Rightarrow glitters(x)$ 

 $\exists x: gold(x) \land \sim glitters(x)$ 

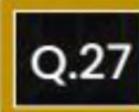
D.  $\exists x : glitters(x) \land \sim gold(x)$ 

## Which one of the following well-formed formulae in predicate calculus is NOT valid?



(GATE-16-Set 2)

- A.  $((\forall x p(x) \Rightarrow \forall x q(x)) \Rightarrow (\exists x \sim p(x) \lor \forall x q(x))$
- B.  $(\exists x p(x) \lor \exists x q(x)) \Rightarrow \exists x (p(x) \lor q(x))$
- (x)  $p x \in \Lambda(x) q x \in \Lambda(x) \Rightarrow (\exists x p(x) \land \exists x q(x))$
- D.  $\forall x(p(x) \lor q(x)) \Rightarrow (\forall x p(x) \lor \forall x q(x))$



Consider the first-order logic sentence  $F: \forall x(\exists y R(x, y))$ . Assuming non-empty logical domain, which of the sentences below are implied by F?



I.  $\exists y(\exists xR(x,y))$ 

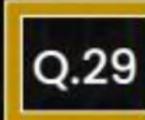
II.  $\exists y(\forall xR(x,y))$ 

III.  $\forall y(\exists xR(x,y))$ 

IV.  $\sim x(\forall y \sim R(x,y))$ 

(GATE-17-Set1)

- A. IV only
- B. I and IV only
- c. II only
- D. II and III only



Consider the first order predicate formula  $\phi$ :



$$(\forall x[(\forall z z|x\Rightarrow((z=x)\lor(z=1)))$$

$$\Rightarrow \exists w(w>x) \land (\forall z \ z|w \Rightarrow ((w=z) \lor (z=1)))])$$

Here 'a|b' denotes that 'a divides b', where a and b are integers.

Consider the following sets:

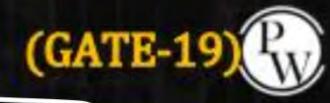
S2. Set of all positive integers

S3. Set of all integers <

Which of the above sets satisfy φ?

- A. (a) S2 and S3
- B. (b) S1, S2 and S3
- c. (c) S1 and S2
- D. (d) S1 and S3

(GATE-19)



 $(\forall x [(\forall z \ z | x \Rightarrow ((z=x) \lor (z=1))) \Rightarrow \exists w(w > x) \land (\forall z \ z | w \Rightarrow ((w=z) \lor (z=1)))])$ 

$$\frac{1}{\sqrt{\kappa}} = \frac{1}{\sqrt{\kappa}} = \frac{1$$

- 1



