

Linear programming problems

Linear programming is a mathematical programming technique to optimize performance [ex:- profit or cost] under a set of resource constraints [ex:- machine-hours, man-hours, money, materials, etc.] as specified by an organization.

form of LPP :-

Maximum or minimum $Z = C_1x_1 + C_2x_2 + \dots + C_nx_n$

STC [Subject to the Constraints]

$$a_{11}x_1 + a_{12}x_2 + \dots + a_{1n}x_n \leq b_1$$

$$a_{11}x_1 + a_{12}x_2 + \dots + a_{1n}x_n \leq b_1$$

$$a_{21}x_1 + a_{22}x_2 + \dots + a_{2n}x_n \leq b_2$$

$$a_{n1}x_1 + a_{n2}x_2 + \dots + a_{nn}x_n \leq b_n$$

Where $x_1, x_2, \dots, x_n \geq 0$

Where x_1, x_2, \dots, x_n are variables

c_1, c_2, \dots, c_n are the cost of variables

$a_{11}, a_{12}, \dots, a_{1n}, a_{21}, a_{22}, \dots, a_{2n}, \dots, a_{n1}, a_{n2}, \dots, a_{nn}$

are the number of variables

b_1, b_2, \dots, b_n are the basic feasible

Solutions

Applications of LPP :-

- (1) Product mix problem
- (2) Diet planning problem
- (3) Cargo loading problem
- (4) Capital Budgeting problem
- (5) Manpower planning problem

Models of LPP :- 2 models

(1) Formation of LPP

(2) Solvation of LPP

Methods for Solving of LPP :-

(1) Graphical method

(2) Simplex method

↳ Artificial variable Technique

Big-M-method

Two phase method

(3) Dual simplex method

Assumptions of linear programming

mainly four assumptions are there in LPP

- (1) Linearity (2) Divisibility (3) Non-negativity
- (4) Additivity

→ Linearity :- The amount of resource required for a given activity level is directly proportional to the level of that activity

Ex:- If the number of hours required on a particular machine is 5 hours per unit of that activity, then the total number of hours required on that machine to produce 10 units of that activity is 50 hours.
[* given activity level]

- Divisibility :- This means that fractional values of decision variables are permitted.
- Non-negativity :- This means that the decision variables are permitted to have only the values which are greater than or equal to zero
- Additivity :- This means that the total output for a given combination of activity levels is the algebraic sum of the output of each individual process.

Properties of linear programming Solutions :-

- Feasible Solution
- Optimal Solution
- Alternative optimal solution
- Unbounded Solution
- Infeasible Solution
- Degenerate Solution

2.3 DEVELOPMENT OF LP MODELS

Modelling is an art. One can develop this expertise only by seeing more and more models. In this section, the concept of model building is demonstrated using some example problems.

Example 2.2 A company manufactures two types of products, P_1 and P_2 . Each product uses lathe and milling machine. The processing time per unit of P_1 on the lathe is 5 hours and on the milling machine is 4 hours. The processing time per unit of P_2 on the lathe is 10 hours and on the milling machine, 4 hours. The maximum number of hours available per week on the lathe and the milling machine are 60 hours and 40 hours, respectively. Also the profit per unit of selling P_1 and P_2 are Rs. 6.00 and Rs. 8.00, respectively. Formulate a linear programming model to determine the production volume of each of the products such that the total profit is maximized.

Solution The data of the problem are summarized in Table 2.2.

Table 2.2 Details of Products

(in hour)

Machine	Machine hours/unit		Limit on machine hours
	Product P_1	Product P_2	
Lathe	5	10	60
Milling machine	4	4	40
Profit/unit (Rs.)	6	8	

Let X_1 and X_2 be the production volumes of the products P_1 and P_2 , respectively. The corresponding linear programming model to determine the production volume of each of the products, such that the total profit is maximized, is presented below.

$$\text{Maximize } Z = 6X_1 + 8X_2$$

subject to

$$5X_1 + 10X_2 \leq 60$$

$$4X_1 + 4X_2 \leq 40$$

$$X_1 \text{ and } X_2 \geq 0$$

Example 2.7 (Cargo loading problem) Consider the cargo loading problem, where five items are to be loaded on a vessel. The weight (w_i) and volume (v_i) of each unit of the different items as well as their corresponding returns per unit (r_i) are tabulated in Table 2.7.

Table 2.7 Example 2.7

Item- i	w_i	v_i	r_i
1	5	1	4
2	8	8	7
3	3	6	6
4	2	5	5
5	7	4	4

The maximum cargo weight (W) and volume (V) are given as 112 and 109, respectively. It is required to determine the optimal cargo load in discrete units of each item such that the total return is maximized. Formulate the problem as an integer programming model.

Solution Let, X_i be the number of units of the i th item to be loaded in the cargo, where i varies from 1 to 5. A model to maximize the return is as follows:

$$\text{Maximize } Z = 4X_1 + 7X_2 + 6X_3 + 5X_4 + 4X_5$$

subject to

$$5X_1 + 8X_2 + 3X_3 + 2X_4 + 7X_5 \leq 112$$

$$X_1 + 8X_2 + 6X_3 + 5X_4 + 4X_5 \leq 109$$

X_1, X_2, X_3, X_4 , and $X_5 \geq 0$ and integers

Three products P_1 , P_2 and P_3 . The machine hour

$$X_1, X_2, X_3, X_4, X_5, X_6, X_7, X_8 \text{ and } X_9 \geq 0$$

Example 2.10 A computer company procures cabinets from three different suppliers (*A*, *B* and *C*) located in three different cities. The company has production plants (*P*, *Q* and *R*) in three other cities. The cost of transportation per cabinet for different combinations of supplier and production plant are summarized in Table 2.10. The purchase price per cabinet from different suppliers are also indicated in the same table. The weekly demand at different production plants and the weekly availability of cabinets at different suppliers are also given in the table. Formulate a linear programming model to find the optimal procurement plan for the cabinets such that the total cost is minimized.

Table 2.10 Details of Unit Transportation Cost and Purchase Cost in Rupees

		Production plant			Price/Cabinet
		<i>P</i>	<i>Q</i>	<i>R</i>	
Supplier	<i>A</i>	10	15	8	300
	<i>B</i>	20	21	15	500
	<i>C</i>	12	16	13	100
Demand		100	200	300	75

Solution Table 2.10 summarizes the transportation cost/cabinet and purchase price/cabinet. The purchase price per cabinet from each supplier is added to different cells of the row corresponding to that supplier. The modified totals of the cost of transportation per unit and purchase cost per unit, for different combinations of supplier and production plant are summarized in Table 2.11.

Table 2.11 Details of Totals of Transportation Cost per Unit and Purchase Cost per Unit

		Production plant			Supply
		<i>P</i>	<i>Q</i>	<i>R</i>	
Supplier	<i>A</i>	95	100	93	300
	<i>B</i>	110	111	105	500
	<i>C</i>	87	91	88	100
Demand		100	200	300	

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Let, Y_{ij} be the number of cabinets to be procured from the supplier i for the production plant j for $i = A, B, C$ and $j = P, Q, R$.

A linear programming model to find the optimal procurement plan for the cabinets to minimize the total cost of procurement is presented below.

$$\text{Minimize } Z = 95Y_{AP} + 100Y_{AQ} + 93Y_{AR} + 110Y_{BP} + 111Y_{BQ} + 105Y_{BR} + 87Y_{CP} + 91Y_{CQ} + 88Y_{CR}$$

subject to

$$Y_{AP} + Y_{AQ} + Y_{AR} \leq 300$$

$$Y_{BP} + Y_{BQ} + Y_{BR} \leq 500$$

$$Y_{CP} + Y_{CQ} + Y_{CR} \leq 100$$

$$Y_{AP} + Y_{BP} + Y_{CP} \geq 100$$

$$Y_{AQ} + Y_{BQ} + Y_{CQ} \geq 200$$

$$Y_{AR} + Y_{BR} + Y_{CR} \geq 300$$

$$Y_{AP}, Y_{AQ}, Y_{AR}, Y_{BP}, Y_{BQ}, Y_{BR}, Y_{CP}, Y_{CQ}, \text{ and } Y_{CR} \geq 0$$

① Maximize $Z = 20x_1 + 10x_2$

STC $10x_1 + 5x_2 \leq 50$

STC $6x_1 + 10x_2 \leq 60$

STC $4x_1 + 12x_2 \leq 48$

$x_1, x_2 \geq 0$

Sol:- Step : ① Evaluating the constraints

$\Rightarrow \text{Max } Z = 20x_1 + 10x_2$

STC $10x_1 + 5x_2 = 50 \rightarrow ①$

STC $6x_1 + 10x_2 = 60 \rightarrow ②$

STC $4x_1 + 12x_2 = 48 \rightarrow ③$

Eq ① can be written as.

$2x_1 + x_2 = 10 \rightarrow ④$

Eq ② can be written as

$3x_1 + 5x_2 = 30 \rightarrow ⑤$

Eq ③ can be written as

$x_1 + 3x_2 = 12 \rightarrow ⑥$

Step : ② Construction of points from

④, ⑤ & ⑥

(A) Put $x_1 = 0$ in ④ then $x_2 = 10 \rightarrow (0, 10)$

(B) Put $x_2 = 0$ in ④ then $x_1 = 5 \rightarrow (5, 0)$

From ④ the points $(5, 0), (0, 10)$

$$\textcircled{B} \rightarrow 3x_1 + 5x_2 = 30$$

put $x_1 = 0$ then $x_2 = 6 \Rightarrow (0, 6)$

put $x_2 = 0$ then $x_1 = 10 \Rightarrow (10, 0)$

from \textcircled{B} the points are $(0, 6) \& (10, 0)$

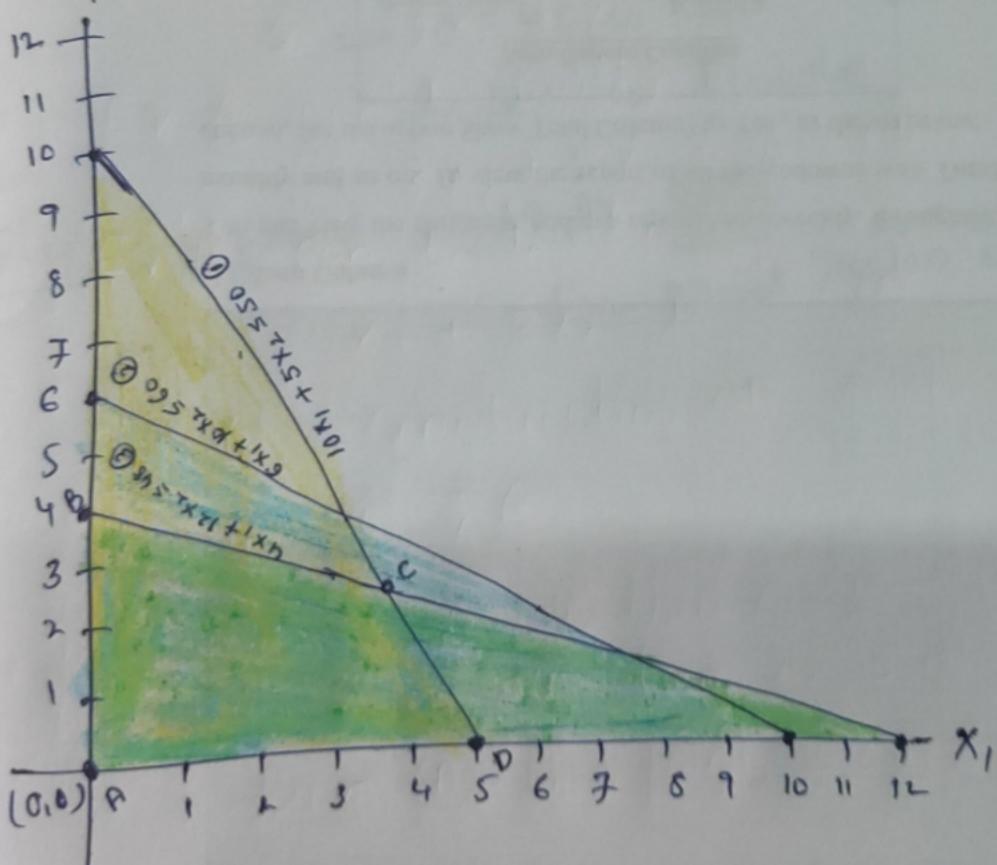
$$\textcircled{C} \rightarrow x_1 + 3x_2 = 12$$

put $x_1 = 0$ then $x_2 = 4 \Rightarrow (0, 4)$

put $x_2 = 0$ then $x_1 = 12 \Rightarrow (12, 0)$

from \textcircled{C} the points are $(12, 0) \& (0, 4)$

Step: \textcircled{D} Draw the Graph



Step : ④

Take constraint ① $10x_1 + 5x_2 \leq 50$

put $x_1 = x_2 = 0$

$\Rightarrow 0 \leq 50$ is true
shaded is Inside



Take constraint ② $6x_1 + 10x_2 \leq 60$

put $x_1 = x_2 = 0$

$\Rightarrow 0 \leq 60$ is true

shaded is Inside



Take constraint ③ $4x_1 + 12x_2 \leq 48$

put $x_1 = x_2 = 0$

$\Rightarrow 0 \leq 48$ is true

shaded is Inside



Step : ⑤ Profit Region is



Another name for profit region is
Permissible Region. Now the Extreme points
of permissible region is A, B, C, D

$A = (0, 0)$; $B = (0, 4)$; $C = ?$ By Solving
① & ③

$D = (5, 0)$

By solving ① & ③ i.e ④ & ⑤

$$2x_1 + x_2 = 10 \rightarrow ④$$

$$x_1 + 3x_2 = 12 \rightarrow ⑤$$

$$④ - 2⑤$$

$$\begin{array}{r} 2x_1 + x_2 = 10 \\ 2x_1 + 6x_2 = 24 \\ \hline \end{array}$$

$$+ 5x_2 = 14$$

$$x_2 = 14/5$$

put x_2 value in ④

$$2x_1 = 10 - x_2$$

$$= 10 - 14/5 = 50/5 - 14/5 = 36/5$$

$$2x_1 = 36/5$$

$$\therefore x_1 = 18/5$$

$$\therefore C = [18/5, 14/5]$$

Now substitute A, B, C, D values

in the profit function

$$\text{Max } Z \geq 20x_1 + 10x_2$$

Step: 6 $Z = 20x_1 + 10x_2$

at A(0,0) $\Rightarrow Z = 20(0) + 10(0)$
 $\Rightarrow Z = 0$

at B(0,4) $\Rightarrow Z = 20(0) + 10(4)$
 $\Rightarrow Z = 40$

at C($\frac{18}{5}, \frac{14}{5}$) $\Rightarrow Z = 20\left(\frac{18}{5}\right) + 10\left(\frac{14}{5}\right)$
 $= 20 \times \frac{18}{5} + 10 \times \frac{14}{5}$
 $= [4 \times 18] + [2 \times 14]$
 $= 72 + 28$
 $= 100$

$\therefore Z = 100$

at ~~D~~(5,0) $\Rightarrow Z = 20(5) + 10(0)$
 $= 100 + 0$
 $\therefore Z = 100$

$\therefore \text{Max } Z = 100$

M.F.S are $x_1 = \frac{18}{5}$ & $x_2 = \frac{14}{5}$
 $(0,4)$
 $x_1 = 5 \quad x_2 = 0$

for calculating of optimal solution two
 Alternative feasible solutions are there.
 so this type of problems will be
 called as Alternative optimal solution
 problem



$$\begin{aligned}
 & \text{(2) Max } Z = 100x_1 + 50x_2 \\
 \text{STC} \quad & 4x_1 + 6x_2 \leq 24 \\
 & x_1 \leq 4 \\
 & x_2 \leq 4/3 \\
 & \left. \begin{array}{l} \\ \\ \end{array} \right\} \text{Graphical method} \\
 & \left. \begin{array}{l} \\ \\ \end{array} \right\} x_1, x_2 \geq 0
 \end{aligned}$$

Sols:- Step: ① Evaluating the constraints

$$\begin{aligned}
 & \text{ie, } \max Z = 100x_1 + 50x_2 \\
 \text{STC} \quad & 4x_1 + 6x_2 \leq 24 \quad \text{--- ①} \\
 & x_1 = 4 \quad \text{--- ②} \\
 & x_2 = 4/3 \quad \text{--- ③}
 \end{aligned}$$

Step: ② Take Eq ①

$$\begin{aligned}
 & 4x_1 + 6x_2 = 24 \\
 \text{Put } x_1 = 0 & \Rightarrow 6x_2 = 24 \\
 & \Rightarrow x_2 = 4 \quad (0, 4) \\
 \text{Put } x_2 = 0 & \Rightarrow 4x_1 = 24 \\
 & \Rightarrow x_1 = 6 \quad (6, 0)
 \end{aligned}$$

\therefore from Eq ① $\Rightarrow (6, 0) (0, 4)$

$$\begin{aligned}
 \text{Take Eq ② } x_1 = 4 & \text{ ie } x_2 = 0 \\
 & \Rightarrow (4, 0)
 \end{aligned}$$

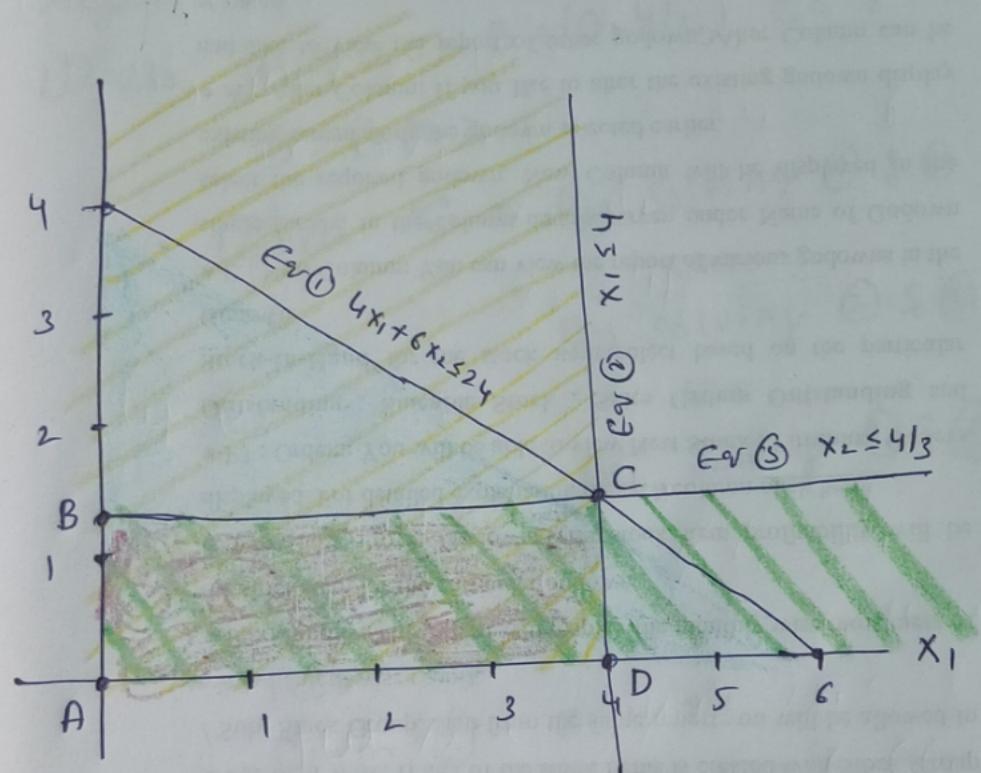
$$\begin{aligned}
 \text{Take Eq ③ } x_2 = 4/3 & \text{ ie } x_1 = 0 \\
 & \Rightarrow (0, 4/3)
 \end{aligned}$$

Step ③: Now draw the Graph by using those points from ①, ②, ③

from ① $\Rightarrow (6, 0), (4, 0)$

from ② $\Rightarrow (4, 0)$

from ③ $\Rightarrow (0, 4/3)$

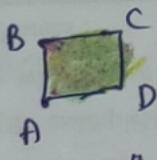


$$E \vee ① \Rightarrow 4x_1 + 6x_2 \leq 24 \Rightarrow 0 \leq 24 \text{ is true}$$

$$E \vee ② \Rightarrow x_1 \leq 4 \Rightarrow 0 \leq 4 \text{ is true } \boxed{\text{ }}$$

$$E \vee ③ \Rightarrow x_2 \leq 4/3 \Rightarrow 0 \leq 4/3 \text{ is true } \boxed{\text{ }}$$

Step: ④ Now the permissible region is



Now the extreme points of permissible region is $A \ B \ C \ D$
Where $A = (0, 0)$ $B = (0, 4/3)$ $C = ?$

$$D = (4, 0)$$

Where $C =$ By solving ① & ②

By solving ② & ③

By solving ① & ③

By solving ② & ③

$$\text{from } ② \quad x_1 = 4$$

$$\text{from } ③ \quad x_2 = 4/3$$

$$\therefore C = (x_1, x_2) = (4, 4/3)$$

For calculating of 'C' i.e, feasible point '3'-Chances are there.

So this type of problem is called as problem of Regeneracy.

$(4, 0) (0, 4)$
 $(100/3, 0) (0, 10/3)$
 $(4, 0) (4, 4)$

Step 5 Now substitute A, B, C, D values
 in profit function $Z = 100x_1 + 50x_2$
 at A(0, 0) $\Rightarrow Z = 0$
 at B(0, 4/3) $\Rightarrow Z = 0 + 50 \times \frac{4}{3} = \frac{200}{3}$
 at C(4, 4/3) $\Rightarrow Z = 100(4) + 50(\frac{4}{3})$
 $= 400 + \frac{200}{3} = \frac{1400}{3}$
 at D(4, 0) $\Rightarrow Z = 100(4) + 50(0)$
 $= 400$

$$\therefore \text{Max } Z = \frac{1400}{3}$$

2 M.F.S are $x_1 = 4$ & $x_2 = \frac{4}{3}$

This problem will be called ~~as~~ as
 Degeneracy

7. Solve the following LP problem graphically:

subject to

$$\text{Maximize } Z = 20X_1 + 80X_2$$

$$4X_1 + 6X_2 \leq 90$$

$$8X_1 + 6X_2 \leq 100$$

$$5X_1 + 4X_2 \leq 80$$

$$X_1 \text{ and } X_2 \geq 0$$

8. Solve the following LP problem graphically:

$$\text{Minimize } Z = 20X_1 + 10X_2$$

subject to

$$X_1 + 2X_2 \leq 40$$

$$3X_1 + X_2 \geq 30$$

$$4X_1 + 3X_2 \geq 60$$

$$X_1 \text{ and } X_2 \geq 0$$

9. Solve the following LP problem graphically:

$$\text{Maximize } Z = 60X_1 + 90X_2$$

subject to

$$X_1 + 2X_2 \leq 40$$

$$2X_1 + 3X_2 \leq 90$$

$$X_1 - X_2 \geq 10$$

$$X_1 \text{ and } X_2 \geq 0$$

10. Solve the following LP problem graphically:

$$\text{Minimize } Z = 45X_1 + 55X_2$$

subject to

$$X_1 + 2X_2 \leq 30$$

$$2X_1 + 3X_2 \leq 80$$

$$X_1 - X_2 \geq 8$$

$$X_1 \text{ and } X_2 \geq 0$$

11. Solve the following LP problem graphically:

$$\text{Maximize } Z = 3X_1 + 2X_2$$

subject to

$$-2X_1 + 3X_2 \leq 9$$

$$X_1 - 5X_2 \geq -20$$

$$X_1 \text{ and } X_2 \geq 0$$

GRAPHICAL METHOD

(1) Solve the following LPP by using Graphical method:-

$$\text{Maximum } Z = 6x_1 + 8x_2$$

$$\text{STC } 5x_1 + 10x_2 \leq 60$$

$$4x_1 + 4x_2 \leq 40$$

$$\& x_1, x_2 \geq 0$$

Sol:- Step:① Equating the constraints

$$\Rightarrow \text{Max } Z = 6x_1 + 8x_2$$

$$\text{STC } 5x_1 + 10x_2 = 60 \quad \text{--- (1)}$$

$$4x_1 + 4x_2 = 40 \quad \text{--- (2)}$$

$$\& x_1, x_2 \geq 0$$

Step:② Take Equation (1)

$$5x_1 + 10x_2 = 60$$

Put $x_1 = 0$ in above equation

$$\Rightarrow 5(0) + 10x_2 = 60$$

$$\Rightarrow 0 + 10x_2 = 60$$

$$\Rightarrow x_2 = 6$$

If $x_1 = 0$ then $x_2 = 6$

The point will be (x_1, x_2) i.e., $(0, 6)$

Put $x_2 = 0$ in above equation

$$5x_1 + 10(0) = 60$$

$$\Rightarrow 5x_1 + 0 = 60$$

$$\Rightarrow x_1 = 12$$

If $x_2 = 0$ then $x_1 = 12$

The point will be (x_1, x_2) i.e., $(12, 0)$

\therefore from equation ① the points will be
 $(12, 0) \& (0, 6)$

Step : ③ Take Equation ②

$$4x_1 + 4x_2 = 40$$

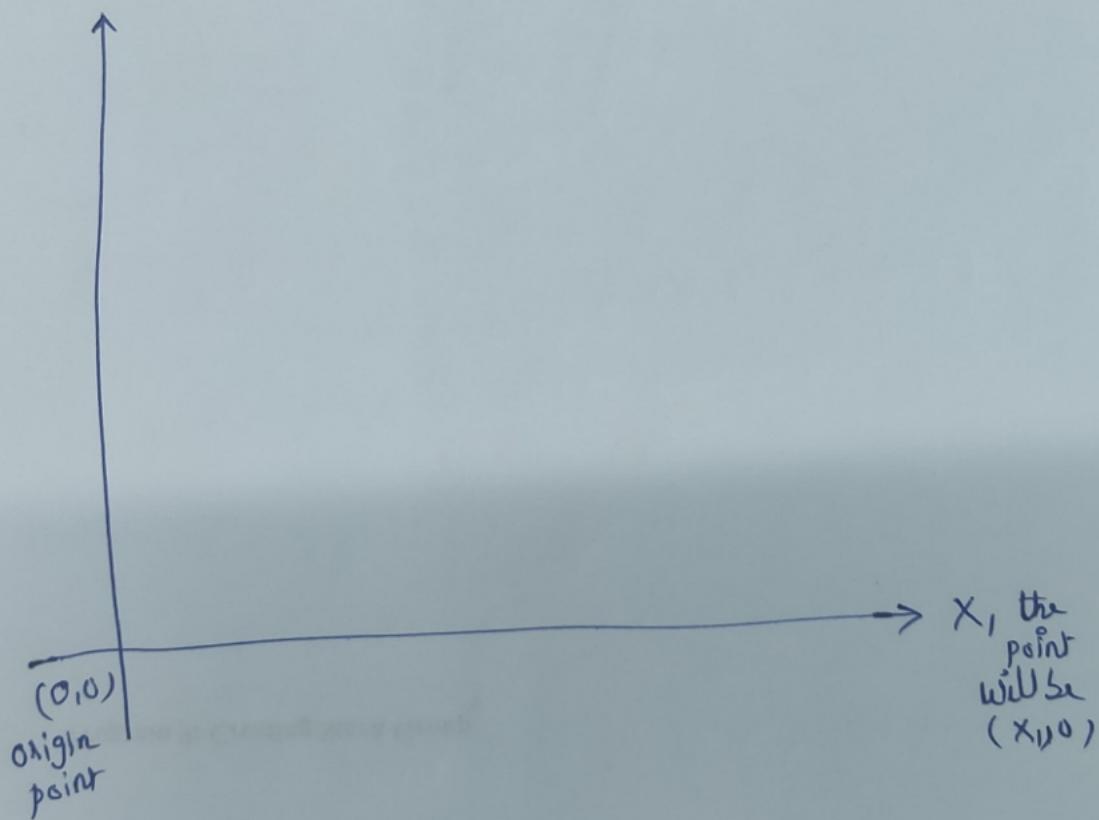
Put $x_1 = 0 \Rightarrow 4x_2 = 40 \quad | : 4$
 $\Rightarrow x_2 = 10 \quad (0, 10)$

$$\text{Put } x_2 = 0 \Rightarrow 4x_1 = 40$$
$$\Rightarrow x_1 = 10 \quad (10, 0)$$

\therefore from Equation ② the points will be
 $(10, 0) \& (0, 10)$

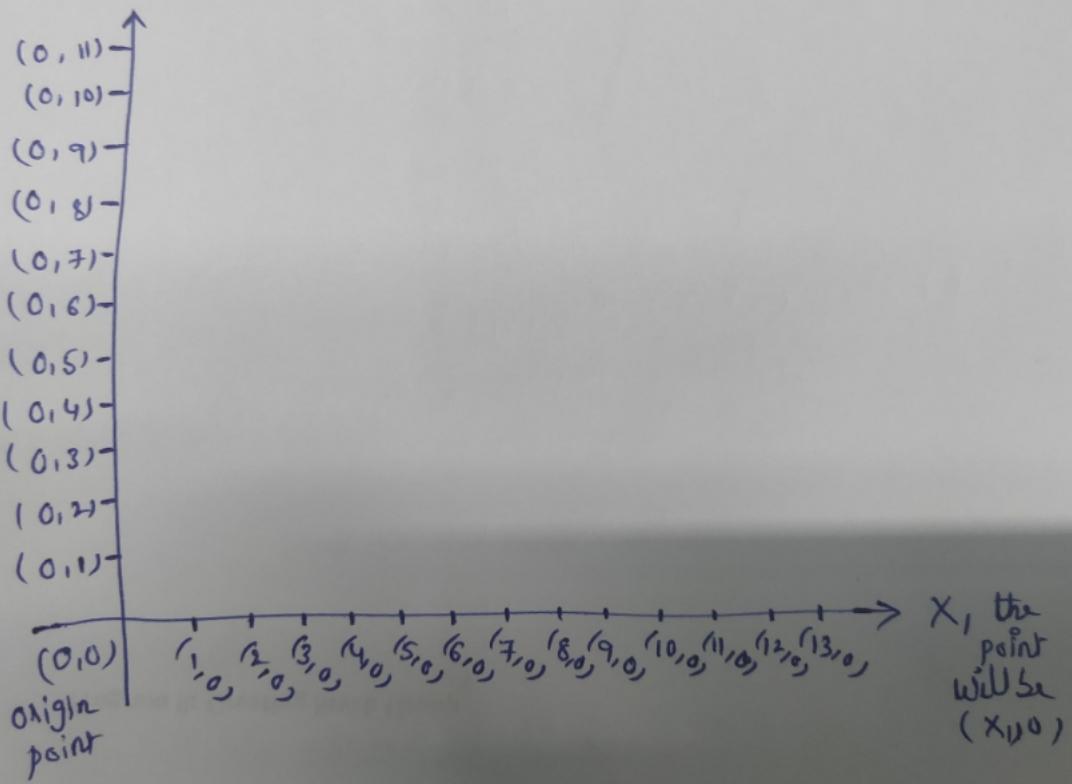
Step : ④ Now draw the Graph by using

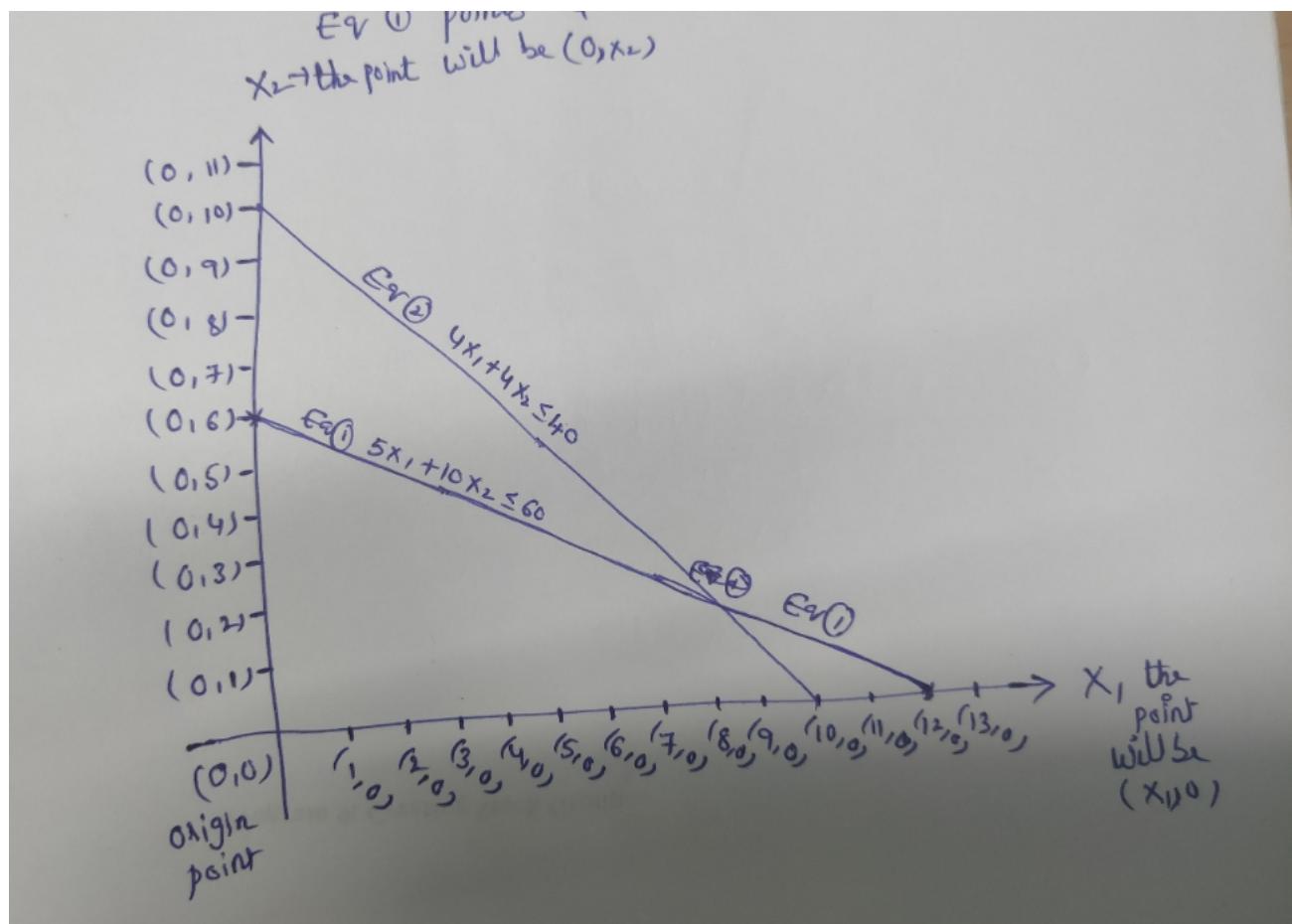
Eq ① points & Eq ② points
 $x_2 \rightarrow$ the point will be $(0, x_2)$



x_1 , the
point
will be
 $(x_1, 0)$

~~Eq ①~~ Eq ① points & Eq ② points
 $x_2 \rightarrow$ the point will be $(0, x_2)$





Step: ⑤ Now check the profit region of each constraint is either inside the region or outside the region.

How can we check?

Take Constraint ①

$$\text{ie., } 5x_1 + 10x_2 \leq 60$$

put $x_1 = x_2 = 0$ in above ~~or~~ Constraint ①

Now Constraint ① becomes

$$5(0) + 10(0) \leq 60$$

$$\Rightarrow 0+0 \leq 60$$

$$\Rightarrow 0 \leq 60$$

Is it true or false.

$\therefore 0 \leq 60$ is True.

→ If the condition is True,

shade the profit region Inside the scale

→ If the condition is false,

shade the profit region outside the scale

for this Constraint ① the condition is True

Now Constraint ① shaded region is Inside

Take Constraint ②

$$\text{ie., } 4x_1 + 4x_2 \leq 10, \text{ put } x_1 = x_2 = 0$$

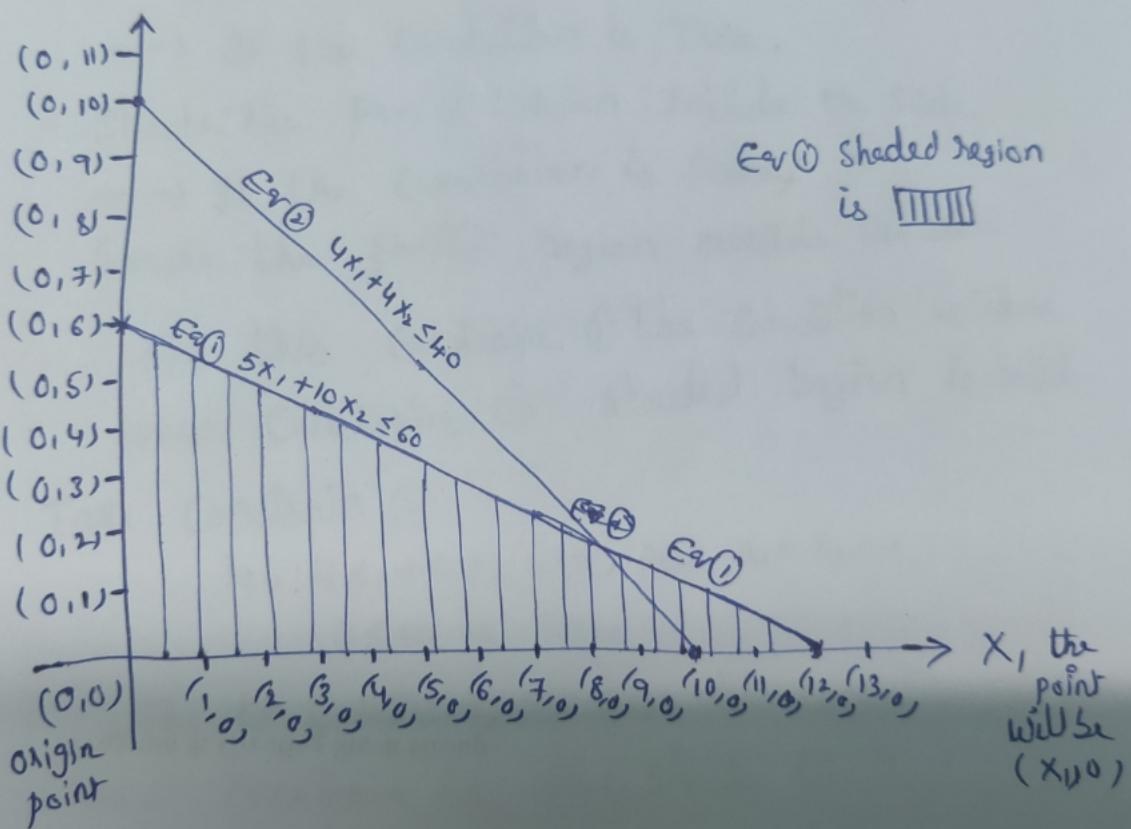
$$0 \leq 10 \text{ is True.}$$

For Constraint ② also the condition is True.

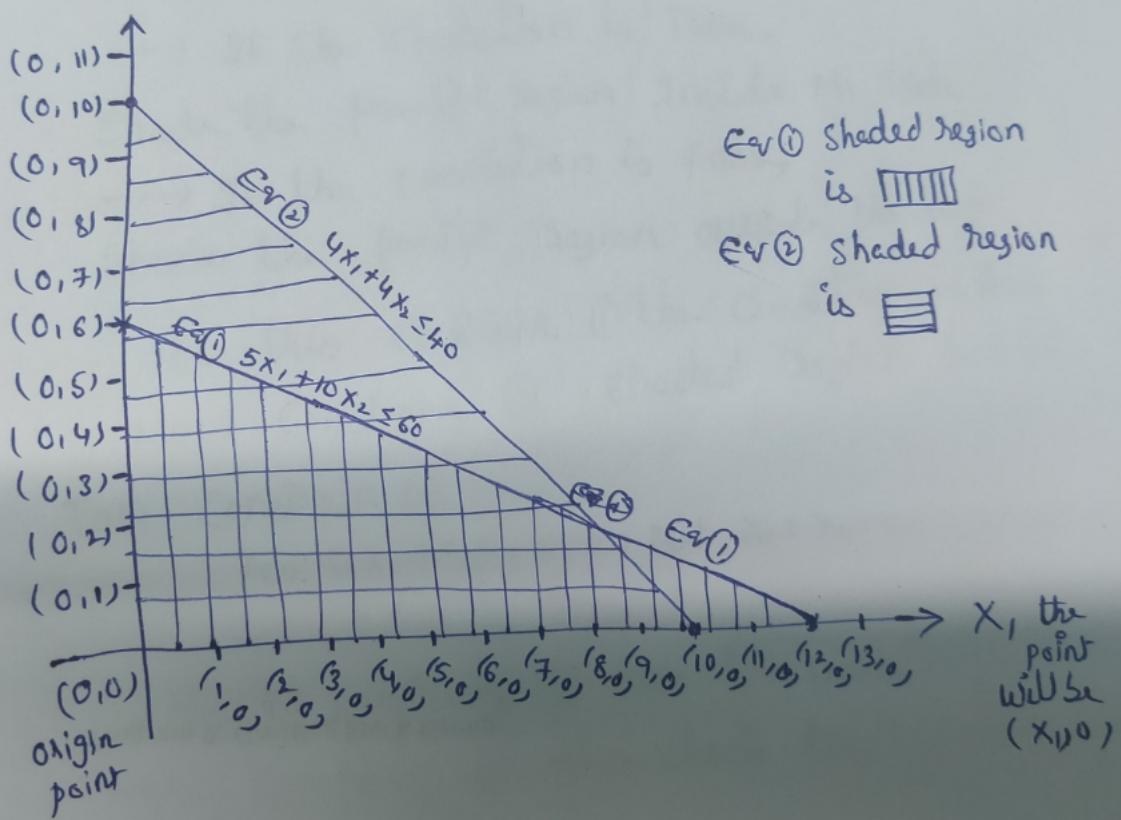
Now Constraint ② also shade the region Inside

only.

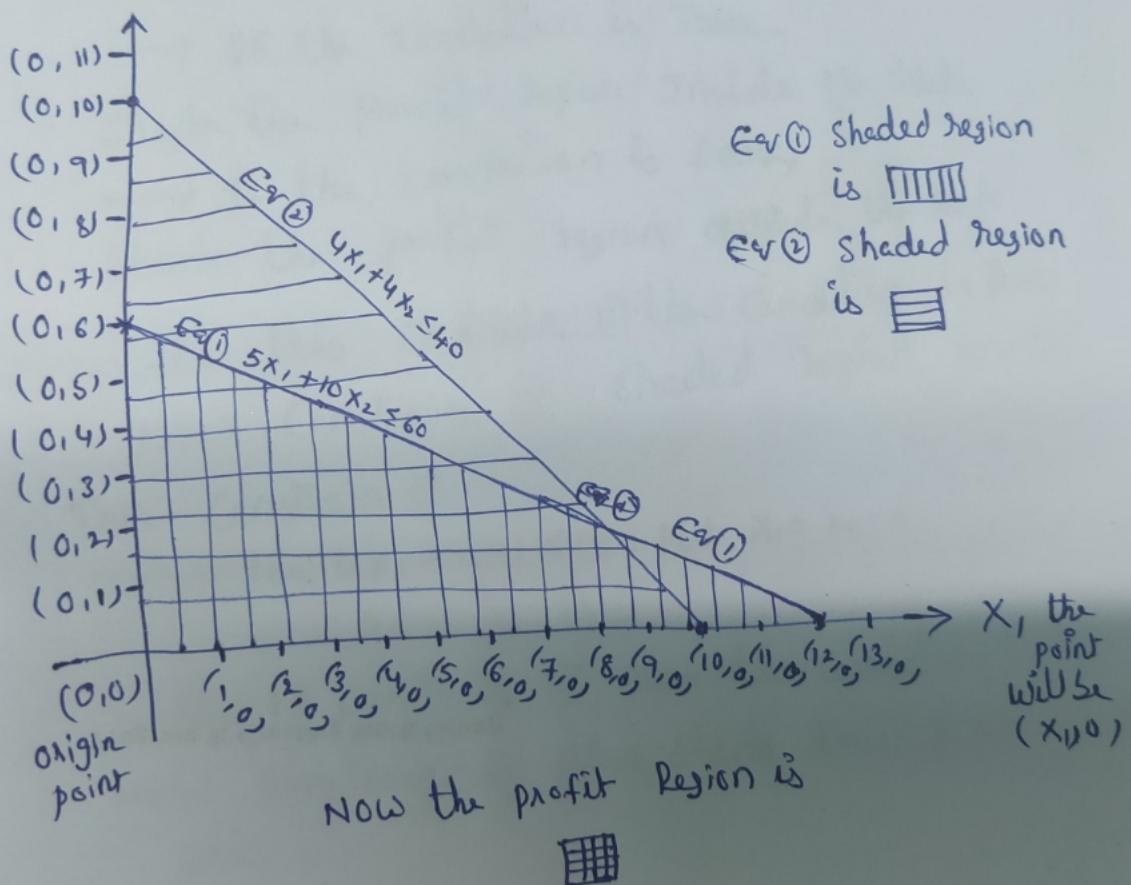
Step 4 Now draw the Graph by using
 Eq ① points & Eq ② points
 $x_2 \rightarrow$ the point will be $(0, x_2)$



Step ④ Now draw the Graph by using
Eq ① points & Eq ② points
 $x_2 \rightarrow$ the point will be $(0, x_2)$

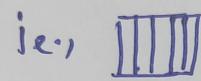


Step ④ Now draw the Graph by using
 Eq ① points & Eq ② points
 $x_2 \rightarrow$ the point will be $(0, x_2)$



Step: ⑥ Now shade the region through Horizontal lines & vertical lines.

→ shade the Eq ① through vertical lines



→ shade the Eq ② through Horizontal lines



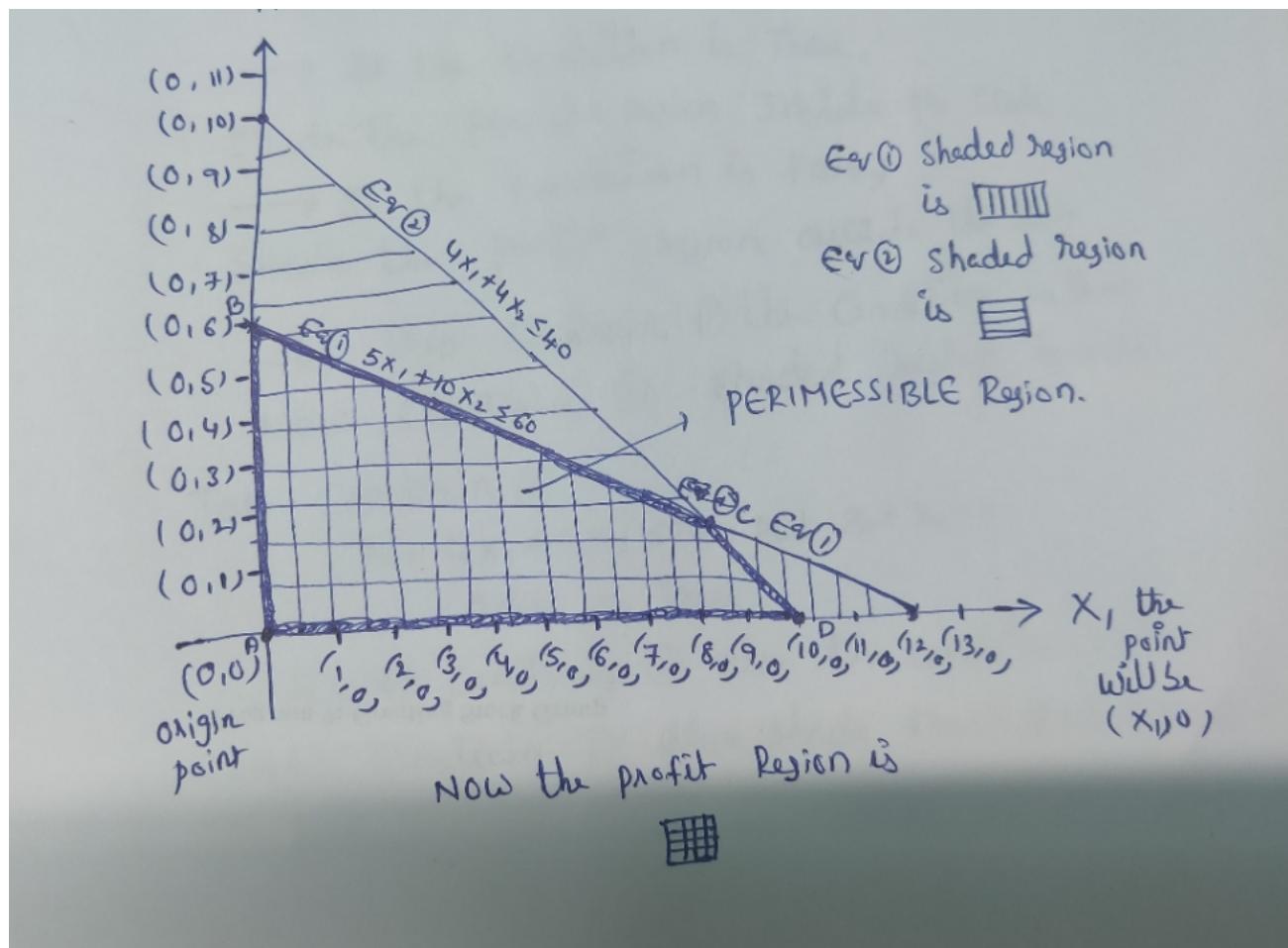
Now the intersecting of these two shaded regions will be called as profit

Area i.e. A hand-drawn diagram of a 3x3 grid of small squares, representing the intersection area.

Another Name for profit Area is

permissible region

.....



permissible region

Step : ⑦ Now calculate the Extreme points of the
permissible region is A, B, C, D

where A = (0,0) on origin

B = (0,6) on x_2 Axis

C = ? ~~Intersection~~ Intersecting the
constraint ① & ②

D = (10,0) on x_1 Axis

Now By solving of Eq ① & Eq ②
we get the value of 'c'

Step: 8 By solving of Eq① & Eq②

$$5x_1 + 10x_2 = 60 \Rightarrow 5[x_1 + 2x_2] = 5(12)$$

$$4x_1 + 4x_2 = 40 \Rightarrow 4[x_1 + x_2] = 4(10)$$

Step: ⑧ By solving of Eq ① & Eq ②

$$5x_1 + 10x_2 = 60 \Rightarrow 5[x_1 + 2x_2] = 5(12)$$

$$4x_1 + 4x_2 = 40 \Rightarrow 4[x_1 + x_2] = 4(10)$$

$$\Rightarrow x_1 + 2x_2 = 12$$

$$\begin{array}{r} x_1 + x_2 = 10 \\ (-) \qquad (-) \qquad (-) \\ \hline \end{array}$$

Step : ⑥ By solving of Eq ① & Eq ②

$$5x_1 + 10x_2 = 60 \Rightarrow 5[x_1 + 2x_2] = 5(12)$$
$$4x_1 + 4x_2 = 40 \Rightarrow 4[x_1 + x_2] = 4(10)$$

$$\begin{aligned} &\Rightarrow x_1 + 2x_2 = 12 \\ &\begin{array}{r} x_1 + x_2 = 10 \\ \hline (-) \quad (-) \quad (-) \end{array} \\ &\qquad\qquad\qquad x_2 = 2 \end{aligned}$$

Put $x_2 = 2$ in $x_1 + x_2 = 10$

$$\begin{aligned} &\Rightarrow x_1 = 10 - x_2 \\ &= 10 - 2 \\ &= 8 \end{aligned}$$

$$\Rightarrow x_1 = 8$$

NOW Take the point $C = (x_1, x_2)$

then $C = (8, 2)$

Step:-⑦ Now take the profit equation & profit points A, B, C, D

Profit equation is $\text{Max } Z = 6x_1 + 8x_2$

& profit points are $A = (0, 0)$

$B = (0, 6)$, $C = (8, 2)$, $D = (10, 0)$

Now substitute profit points in profit equation.

$$\text{at } A(0,0) \Rightarrow Z = 6(0) + 8(0) = 0$$

$$\boxed{\therefore Z = 0 \text{ at } A(0,0)}$$

$$\text{at } B(0,6) \Rightarrow Z = 6(0) + 8(6) \\ = 0 + 48 = 48$$

$$\boxed{\therefore Z = 48 \text{ at } (0,6)}$$

$$\text{at } C(8,2) \Rightarrow Z = 6(8) + 8(2) \\ = 48 + 16 = 64$$

$$\boxed{\therefore Z = 64 \text{ at } (8,2)}$$

$$\text{at } D(10,0) \Rightarrow Z = 6(10) + 8(0) \\ = 60 + 0 = 60$$

$$\boxed{\therefore Z = 60 \text{ at } (10,0)}$$

Step: 10 Now Select the maximum profit
in $Z=0, Z=48, Z=64, Z=60$
 \therefore Maximum profit $Z=64$

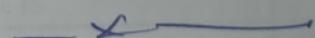
i.e., $\text{Max } Z=64$

∴ Maximum feasible solution are

$$x_1=8 \quad \& \quad x_2=2$$

\therefore The optimal solution is 64

∴ feasible solutions are $x_1=8$
 $x_2=2$



InFeasible Solution

Solve the following LPP Graphically

$$\text{Maximize } Z = 10x_1 + 3x_2$$

$$\text{STC } 2x_1 + 3x_2 \leq 18$$

$$6x_1 + 5x_2 \leq 60$$

$$x_1, x_2 \geq 0$$

Sol:- Step ① Equating the Constraints

$$\text{Max } Z = 10x_1 + 3x_2$$

$$\text{STC } 2x_1 + 3x_2 = 18 \rightarrow ①$$

$$6x_1 + 5x_2 = 60 \rightarrow ②$$

$$x_1, x_2 \geq 0$$

Step ② Take Equation ①

$$2x_1 + 3x_2 = 18$$

$$\text{Put } x_1 = 0 \Rightarrow 3x_2 = 18 \\ \Rightarrow x_2 = 6 \quad (0, 6)$$

$$\text{Put } x_2 = 0 \Rightarrow 2x_1 = 18 \\ \Rightarrow x_1 = 9 \quad (9, 0)$$

∴ from Eq ① the points are (9, 0) & (0, 6)

Step ③ Take Equation ②

$$6x_1 + 5x_2 = 60$$

$$\text{Put } x_1 = 0 \Rightarrow 5x_2 = 60 \Rightarrow (0, 12)$$

$$\Rightarrow x_2 = 12$$

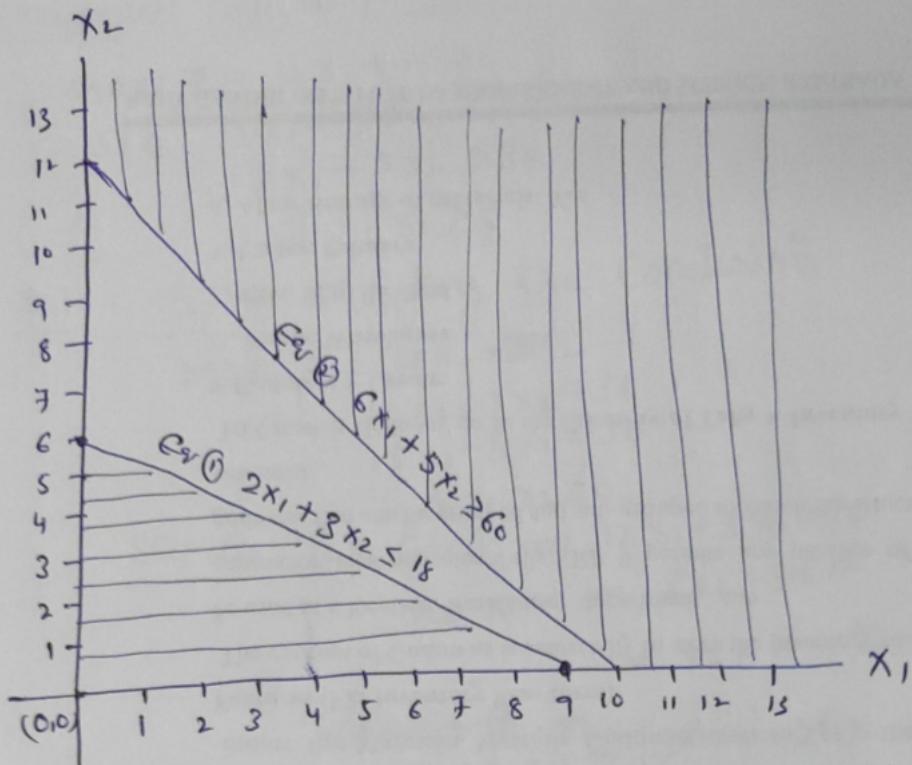
$$\text{Put } x_2 = 0 \Rightarrow 6x_1 = 60 \Rightarrow x_1 = 10 \quad (10, 0)$$

∴ from Eq ② the points are (10, 0) & (0, 12)

Step ④ Now draw a graph by using Eq ①

points & Eq ② points & shade it

by using shaded region for calculating
of profit Region



Ex ① shaded region is \square since Inside the
region and origin scale

Take constraint ① $2x_1 + 3x_2 \leq 18$

put $x_1 = x_2 = 0 \Rightarrow 0 \leq 18$ is True

Ex ② shaded region is \square since outside
the origin scale

Take constraint ② $6x_1 + 5x_2 \geq 60$
put $x_1 = x_2 = 0 \Rightarrow 0 \geq 60$ is ~~False~~.

In the above graph there is no intersecting
profit Area. i.e., No permissible Region

\therefore This type of problems can be called
as Infeasible Solution problems. Since
there is no feasible solution

$\rightarrow \infty$

Unbounded Solution Problem

$$\begin{aligned} \text{Max } Z &= 12x_1 + 25x_2 \\ \text{STC } 12x_1 + 3x_2 &\geq 36 \\ 15x_1 - 5x_2 &\leq 30 \\ x_1, x_2 &\geq 0 \end{aligned}$$

Sol:- Step: ① Equating the Constraints

$$\begin{aligned} \text{Max } Z &= 12x_1 + 25x_2 \\ \text{STC } 12x_1 + 3x_2 &= 36 \\ 15x_1 - 5x_2 &= 30 \\ x_1, x_2 &\geq 0 \end{aligned}$$

Step: ② Take Eq ① $\Rightarrow 12x_1 + 3x_2 = 36$
 Put $x_1 = 0 \Rightarrow 3x_2 = 36 \Rightarrow x_2 = 12 \quad (0, 12)$

$$\begin{aligned} \text{Take Eq ①} &\Rightarrow 12x_1 = 36 \\ \text{Put } x_2 = 0 &\Rightarrow x_1 = 3 \quad (3, 0) \end{aligned}$$

\therefore from Eq ① the points are

$$(3, 0) \text{ & } (0, 12)$$

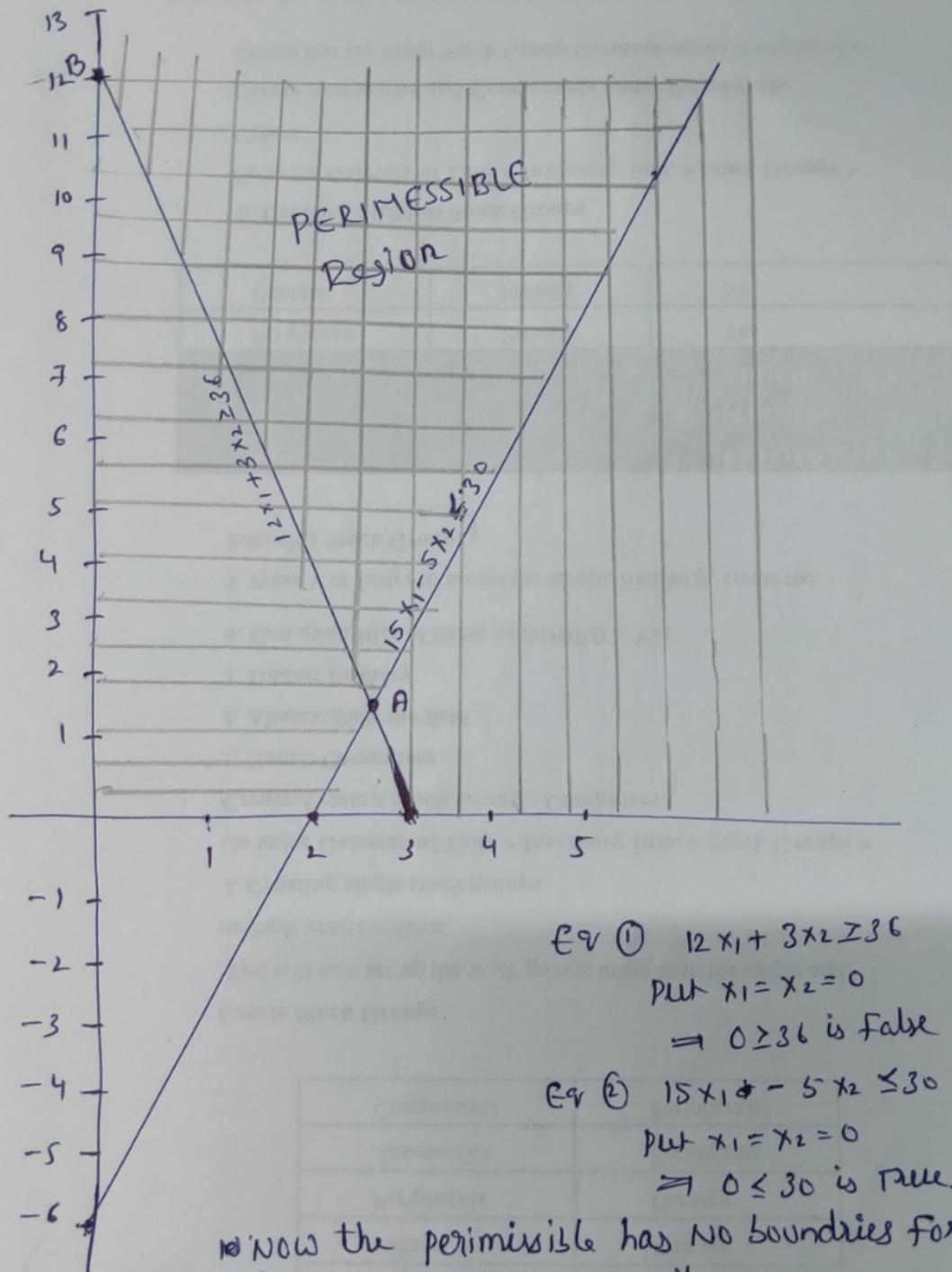
$$\begin{aligned} \text{Step: ③ Take Eq ②} &\Rightarrow 15x_1 - 5x_2 = 30 \\ \text{Put } x_1 = 0 &\Rightarrow -5x_2 = 30 \Rightarrow x_2 = -6 \quad (0, -6) \end{aligned}$$

$$\begin{aligned} \text{Put } x_2 = 0 &\Rightarrow 15x_1 = 30 \Rightarrow x_1 = 2 \quad (2, 0) \end{aligned}$$

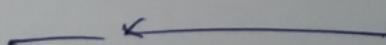
\therefore from Eq ② the points are

$$(2, 0) \text{ & } (0, -6)$$

Step: ④ Now draw the Graph by using
 Eq ① points & Eq ② points and
 shade it



Now the permissible has No boundaries for calculating of maximum profit
 Now this type of problems is called unbounded solution problem



SIMPLEX METHOD

than
 II) Slack variable :- If a constraint is less, or equal to zero (≤ 0) then in order to make it an Equality we have add something positive of LHS. Then that positive variable is called as slack variable.

$$\text{Ex: } \begin{aligned} \text{Max } Z &= 4x_1 + 7x_2 \\ \text{STC } x_1 + x_2 &\leq 6 \\ 2x_1 + x_2 &\leq 5 \\ x_1, x_2 &\geq 0 \end{aligned}$$

Evaluating the Constraints by add slack variables on LHS

$$\Rightarrow \begin{aligned} \text{Max } Z &= 4x_1 + 7x_2 + 0s_1 + 0s_2 \\ \text{STC } x_1 + x_2 + s_1 &= 6 \\ 2x_1 + x_2 + s_2 &= 5 \\ x_1, x_2, s_1, s_2 &\geq 0 \end{aligned}$$

where s_1 & s_2 are slack variables

than
 (2) Surplus variable :- If a constraint is greater, or equal to zero (≥ 0) then in order to make it an Equality we have add something negative on LHS. Then that negative variable is called as surplus variable

$$\text{Ex: Max } Z = 4x_1 + 7x_2 \\ \text{STC } x_1 + x_2 \geq 6 \\ 2x_1 + x_2 \geq 5 \\ x_1 + 3x_2 \leq 7 \\ x_1, x_2 \geq 0$$

Evaluating the Constraints by adding slack & surplus variables on LHS

$$\Rightarrow \begin{aligned} \text{Max } Z &= 4x_1 + 7x_2 + 0s_1 + 0s_2 + 0s_3 \\ \text{STC } x_1 + x_2 - s_1 &= 6 \\ 2x_1 + x_2 - s_2 &= 5 \\ x_1 + 3x_2 + s_3 &= 7 \end{aligned}$$

where $x_1, x_2, s_1, s_2, s_3 \geq 0$

where s_1, s_2 are surplus variables & s_3 is slack variable.



SIMPLEX TABLE:-

C_B	B	X_B	$C_1 \ C_2 \ \dots \ C_n \ 0 \ 0 \ 0 \ -1 \ -1 \ -1$	$A_1 \ A_2 \ \dots \ A_n$	$\theta = \min \left\{ \frac{X_B i}{x_{ik}} \mid x_{ik} > 0 \right\}$
			$x_1 \ x_2 \ \dots \ x_n \ s_1 \ s_2 \ \dots \ s_n$		
$Z_B = \sum C_B X_B$			$Z_J = \sum C_B X_J$ where $J=1, 2, \dots, n$		

$$\text{where } Z_B = \sum C_B X_B$$

Z_B = profit

C_B = cost of basic variable

X_B = basic feasible solution

B = basic variables

optimality test: $Z_J - C_J \geq 0$

when Z_J is construction part

C_J is given problem

i.e., cost of variables

$$Z_J = \sum C_B X_J$$

$$\theta = \min \left\{ \frac{X_B i}{x_{ik}} \mid x_{ik} > 0 \right\}$$

In this simplex two types of variables are there

Incoming variable (or) Incoming vector

outgoing variable (or) outgoing vector

The intersection of Incoming ~~and~~ and

out going vector is called as Key element

key element plays main role for construction of next simplex tables.

problem: ① Solve the following LPP by using simplex method

$$\begin{aligned} \text{Max } Z &= 6x_1 + 8x_2 \\ \text{STC } 5x_1 + 10x_2 &\leq 60 \\ 4x_1 + 4x_2 &\leq 40 \end{aligned}$$

$$x_1, x_2 \geq 0$$

Sol:- Step ① Evaluating the constraints by adding either slack or surplus variables. Here we have to add slack variables only since in the constraints they given " \leq "

$$\begin{aligned} \text{Max } Z &= 6x_1 + 8x_2 + 0s_1 + 0s_2 \\ \text{STC } 5x_1 + 10x_2 + s_1 &= 60 \quad \text{--- (1)} \\ 4x_1 + 4x_2 + s_2 &= 40 \quad \text{--- (2)} \end{aligned}$$

$$\text{where } x_1, x_2, s_1, s_2 \geq 0$$

and s_1, s_2 are slack variables

Step ② write the equations in matrix form

$$\begin{bmatrix} x_1 & x_2 & s_1 & s_2 \\ 5 & 10 & 1 & 0 \\ 4 & 4 & 0 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ s_1 \\ s_2 \end{bmatrix} = \begin{bmatrix} 60 \\ 40 \end{bmatrix}$$

In the above matrix form there is a Basis matrix
is their... [Basic matrix means Identity matrix]
Note: - ① If the Basis matrix is their then start the simplex table
② If the Basis matrix is not their.. Then construct
the Basis matrix Table by using Artificial variables

Step ③ Initial simplex table

CB	B	X _B	x ₁	x ₂	s ₁	s ₂	Z

Initial Simplex Table:-

CB	B	XB	6	8	0	0	0
			x_1	x_2	s_1	s_2	
0	s_1	60	5	10	1	0	$60/10 = 6$
0	s_2	40	4	4	0	1	$40/4 = 10$
$Z_B = \sum (C_B \times B)$			-6	-8	0	0	

$$\Rightarrow [0 \times 60] + [0 \times 40]$$

$$\Rightarrow 0$$

$$\therefore Z_B = 0$$

Here the optimality test condition

$$Z_j - C_j \geq 0 \text{ is not satisfied}$$

$$\Rightarrow \{-6, -8, 0, 0\} \neq 0$$

Here most -ve is -8 ie, in x_2

$\therefore x_2$ is Incoming vector

Min '0' is 6 ie in s_1

$\therefore s_1$ is outgoing vector

'10' is key element.

Second Simplex Table :-

CB	B	XB	6	8	0	0	0
			x_1	x_2	s_1	s_2	
8	s_2	6	1/2	1	1/10	0	$6/1/2 = 12$
0	s_1	16	2	0	-2/5	1	$16/2 = 8$
$Z_B = \sum (C_B \times B)$			-2	0	4/5	0	
$\Rightarrow 8(6) + 0(16) = 48$							

$$NR_1 \rightarrow OR_1/10$$

$$60 \rightarrow 60/10 = 6$$

$$5 \rightarrow 5/10 = 1/2$$

$$10 \rightarrow 10/10 = 1$$

$$1 \rightarrow 1/10 = 1/10$$

$$0 \rightarrow 0/10 = 0$$

$$NR_2 \rightarrow OR_2 - 4NR_1$$

$$40 \rightarrow 40 - 4(6) = 16$$

$$4 \rightarrow 4 - 4(1/2) = 2$$

$$4 \rightarrow 4 - 4(1) = 0$$

$$0 \rightarrow 0 - 4(1/10) = -2/5$$

$$1 \rightarrow 1 - 4(0) = 1$$

Here the optimality test condition
 $Z_J - C_J \geq 0$ is not satisfied

$$\Rightarrow \{-2, 0, 4/5, 0\} \neq 0$$

most negative is -2 is in x_1

$\therefore x_1$ is Incoming vector

min 'θ' is 8 i.e. s_2 is

$\therefore s_2$ is outgoing vector

$\therefore '2'$ is key element

Third simplex table:

CB	B	XB	X ₁	X ₂	S ₁	S ₂	θ
8	X ₂	2	0	1	1/5	-1/4	
6	X ₁	8	1	0	-1/5	1/2	
			0	0	2/5	1	

$$Z_B = \sum C_B X_B$$

$$= 8(2) + 6(8)$$

$$= 16 + 48 = 64$$

$$NR_2 \rightarrow OR_2/2$$

$$16 \rightarrow 16/2 = 8$$

$$2 \rightarrow 2/2 = 1$$

$$0 \rightarrow 0/2 = 0$$

$$-2/5 \rightarrow -2/5/2 = -1/5$$

$$1 \rightarrow 1/2 > 1/2$$

$$NR_1 \rightarrow OR_1 - 1/2 NR_2$$

$$6 \rightarrow 6 - 1/2(8) = 6 - 4 = 2$$

$$1/2 \rightarrow 1/2 - 1/2(1) = \frac{1}{2} - \frac{1}{2} = 0$$

$$1 \rightarrow 1 - 1/2(0) = 1 - 0 = 1$$

$$1/10 \rightarrow 1/10 - 1/2(-1/5) = \frac{1}{10} + \frac{1}{10} = \frac{1}{5}$$

$$0 \rightarrow 0 - 1/2(1/2) = -1/4$$

Here the optimality test condition

$Z_J - C_J \geq 0$ is satisfied

$$\{ 0, 0, 4/5, 1 \} \geq 0$$

$\therefore \text{Max } Z = 64$ & M.F.S are $x_1 = 8$
 $x_2 = 2$

— — —

TRANSPORTATION

Def :- Transporting the goods from one place to another place within the minimum cost is called Transportation.

FORM OF TRANSPORTATION :-

FROM		To	Supply
S ₁		C ₁₁	C ₁₂
S ₂		C ₂₁	C ₂₂
:		:	:
S _n		C _{n1}	C _{n2}
Demand	$\rightarrow D_1 \ D_2 \ \dots \ D_n$		

TYPES OF TRANSPORTATION :- 2 Types

(a) Balanced \Rightarrow Total supply = Total Demand

$$\Rightarrow S_1 + S_2 + \dots + S_n = D_1 + D_2 + \dots + D_n$$

(b) UnBalanced \Rightarrow Total supply \neq Total Demand

$$\Rightarrow S_1 + S_2 + \dots + S_n \neq D_1 + D_2 + \dots + D_n$$

TYPES OF TRANSPORTATION SOLUTIONS :- 2 Types

(a) Initial Basic Feasible Solution
[IBFS]

METHODS FOR CALCULATING OF IBFS

→ NORTH-WEST Corner method

→ LEAST COST Entry method

→ VOGEL'S APPROXIMATION Method

(b) OPTIMAL SOLUTION (OS)

Methods for Calculating of OS

(a) MODI - method

Modified Distribution method

UV - method
(or)

(b) Stepping Stone method

— x —

PROBLEMS :-

(1) Solve the following Transportation and
find the IBFS

	P ₁	P ₂	P ₃	<u>Supply</u>
S ₁	2	7	4	5
S ₂	3	3	1	8
S ₃	5	4	7	7
S ₄	1	6	3	14
<u>Demand</u>	7	9	18	

Sol:- Step:① The given problem is a
Balanced Transportation problem
i.e., Total Supply = Total Demand
 $\Rightarrow 5+8+7+14 = 7+9+18$
 $\Rightarrow 34 = 34$
 $\Rightarrow \text{Total Supply} = \text{Total Demand}$

(a) North-West corner method :-

Supply

2	7	4	5
3	3	1	8
5	4	7	7
1	6	2	14

Demand 7 9 18 / 34

/ 34

34

(a) North-West Corner method :-

	P ₁	P ₂	P ₃	<u>Supply</u>
S ₁	2	7	4	5
S ₂	3	3	1	8
S ₃	5	4	7	7
S ₄	1	6	2	14 / 34
Demand	7	9	18	34 / 34

(a) Start the problem (1,1) Cell

(a) North-West Corner method :-

	P ₁	P ₂	P ₃	<u>Supply</u>
S ₁	2	7	4	5
S ₂	3	3	1	6
S ₃	5	4	7	7
S ₄	1	6	2	14
Demand	7	9	18	34
	7			34

(a) Start the problem (1,1) Cell

(b) Here first row completed

NOW start the problem (2,1) Cell

(a) North-West Corner method :-

	P ₁	P ₂	P ₃	<u>Supply</u>
S ₁	2	7	4	5
S ₂	3	3	1	8
S ₃	5	4	7	7
S ₄	1	6	2	14
Demand	7	9	18	34
			34	

(a) Start the problem (1,1) cell

(b) Here first row completed

NOW start the problem (2,1) cell

(c) Here first row & first column also completed. NOW start (2,2) cell

(a) North-West Corner method :-

	P ₁	P ₂	P ₃	<u>Supply</u>
S ₁	5 2 7 4			5
S ₂	2 3 6 1			8
S ₃	5 4 7			7
S ₄	1 6 2			14
Demand	7 9 18			34
	7 9 18			34

(a) Start the problem (1,1) Cell

(b) Here first Row completed

NOW start the problem (2,1) Cell

(c) Here first Row & first Column also completed. Now start (2,2) Cell

(d) Now start (3,2) Cell

(a) North-West Corner method :-

	P ₁	P ₂	P ₃	<u>Supply</u>
S ₁	2	7	4	5
S ₂	5	3	1	8
S ₃	5	4	7	7
S ₄	3	6	2	4
Demand	7	9	18	34
			34	

(a) Start the problem (1,1) Cell

(b) Here first row completed

Now start the problem (2,1) Cell

(c) Here first row & first column also completed. Now start (2,2) Cell

(d) Now start (3,2) Cell

(e) Now start (3,3) Cell

(a) North-West Corner method :-

	P ₁	P ₂	P ₃	<u>Supply</u>
S ₁	2	7	4	5
S ₂	5	3	1	8
S ₃	5	4	7	7
S ₄	3	6	2	4
Demand	7	9	18	34
			34	

(a) Start the problem (1,1) Cell

(b) Here first row completed

Now start the problem (2,1) Cell

(c) Here first row & first column also completed. Now start (2,2) Cell

(d) Now start (3,2) Cell

(e) Now start (3,3) Cell

(a) North-West Corner method :-

	P ₁	P ₂	P ₃	<u>Supply</u>
S ₁	2	7	4	5
S ₂	3	6	1	8
S ₃	5	4	7	7
S ₄	3	6	2	14
Demand	7	9	18	34
	7	7	14	34

(a) Start the problem (1,1) Cell

(b) Here first Row completed

Now start the problem (2,1) Cell

(c) Here first Row & first Column also completed. Now start (2,2) Cell

(d) Now start (3,1) Cell

(e) Now start (3,2) Cell

Now the Total Cost is

$$\Rightarrow 5 \times 2 + 2 \times 3 + 6 \times 3 + 3 \times 4 + 4 \times 7 + 14 \times 2$$

$$\Rightarrow 10 + 6 + 18 + 12 + 28 + 28$$

$$\Rightarrow 102$$

∴ The Minimum cost for N-W method

is 102

(a) North-West corner method :-

Supply

2	7	4	5
3	3	1	8
5	4	7	7
1	6	2	14

Demand 7 9 18 / 34

/ 34

34

N-W method by step by step process

Step: ①

	P ₁	P ₂	P ₃	S
S ₁	5 ²	7	4	8
S ₂	3	3	1	8
S ₃	5	4	7	7
S ₄	1	6	2	14
D	7 ²	9	18	34

Step: ②

	P ₁	P ₂	P ₃	S
S ₁	3	3	1	8/6
S ₂	2	7	7	7
S ₃	7	4	7	7
S ₄	1	6	2	14
D	7	9	18	29

Step: ③

	P ₂	P ₃	S
S ₁	3	7	8
S ₂	4	7	7
S ₃	6	2	7
S ₄	6	2	14
D	9	18	27

Step: ④

	P ₂	P ₃	S
S ₁	4	7	7/4
S ₂	3	2	14
S ₃	7	2	14
S ₄	3	18	21
D	3	18	

Step: ⑤

	P ₃	
S ₁	7	9
S ₂	2	14
S ₃	18	
S ₄	14	
D	18	14

Step: ⑥

	P ₃	
S ₁	2	14
S ₂	14	14
S ₃	14	
S ₄	14	
D	14	

$$\therefore \text{Min Cost} = 10 + 6 + 18 + 12 + 28 + 28 = 102/-$$

— →

(b) LEAST COST ENTRY METHOD

(a)

Lowest Cost entry method

(b)

Matrix minimum method

(b) LEAST COST ENTRY METHOD

(ON)

Lowest cost entry method

(OR)

Matrix minimum method

	P ₁	P ₂	P ₃	S
S ₁	2	7	4	5
S ₂	3	3	1	8
S ₃	5	4	7	7
S ₄	1	6	2	14
D	7	9	18	34

(a) Start the problem where the minimum cost is there. In this problem the M.C is 'I'. But, 'I' is in two places. Select any one

(Q1)

Matrix minimum method

	P ₁	P ₂	P ₃	S
S ₁	2	7	4	5
S ₂	3	3	1	8
S ₃	5	4	7	7
S ₄	1	6	2	14
D	7	9	18 10	34

- (a) Start the problem where the minimum cost is there. In this problem the M.C is '1'. But, '1' is in two places. Select any one