Genetic Algorithm: Traveling Salesman

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**ABSTRACT**

In this paper, we present the impact of the Genetic Algorithm on the Knapsack and the Traveling Salesman problem. First, we use the Genetic algorithm to solve the Knapsack problem. We then use the genetic algorithm to solve the Traveling Salesman problem. There were certain attributes which were changed to improve the performance of the genetic algorithm used for the Traveling Salesman Problem. We compare the results of the optimizations along with the default GA. The genetic algorithm used for the Knapsack problem is compared to the genetic algorithm used for the Traveling Salesman problem.

**CCS Concepts**

• **Symbolic and algebraic manipulation➝ Optimization Algorithm; Combinational Algorithm;**

**Keywords**

Combinatorial optimization; the Knapsack problem;

# INTRODUCTION

The Knapsack problem is a classical problem that searches the highest combinational values from a list of items that consist of cost and values. The knapsack problem is a decision problem such that given a set of items that each consists of a cost and a value, find the highest possible value from each of the item while remaining in the range of the cost that the problem is being constrained. Which means that the cost may only remain less than or equal to the constrained cost limit. The decision form of the Knapsack problem is a NP-complete problem such that a precise solution for a huge input is nearly practically impossible to obtain.

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# THE KNAPSACK PROBLEM BACKGROUND

Suppose a group of hikers is planning on a hiking trip and their plan is to fill their knapsack with items that are considered a necessity for the trip. There are N number of different items that have different **name, weight,** andthe **value** of the item. For example, there are water, sandwich, and more. Each item obviously has their own name, weight, and value. Since the hikers are only able to fit a certain amount of items into the knapsack due to the space or weight limit of how much the knapsack is able to hold, they have to obtain a combination of items that produces a maximum value while staying in within the weight limit of the Knapsack.

## Method Descriptions

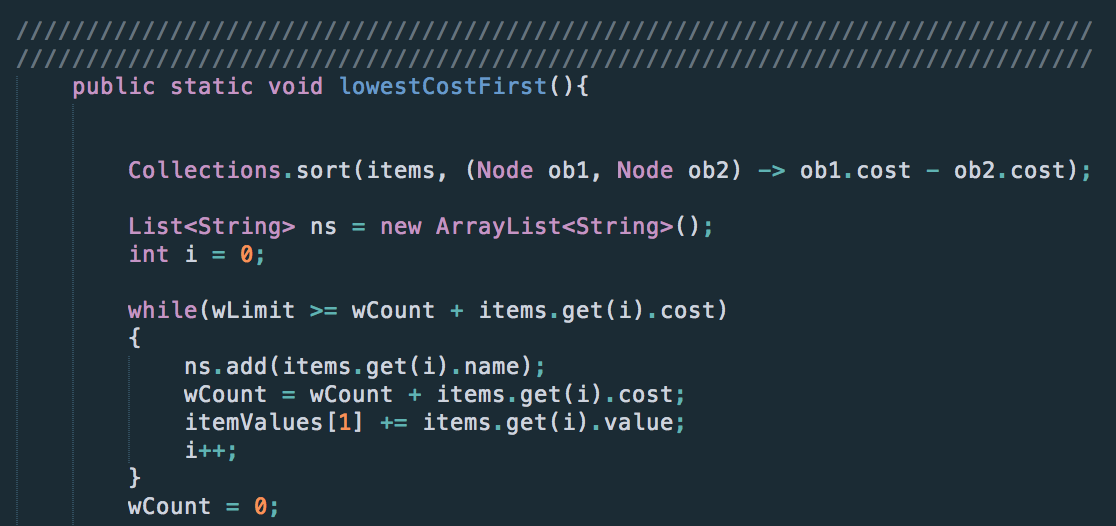
In this section, we will define the Knapsack problem, then we will provide the details of the attempted approaches in which we have tried to get the most optimal results.

### Greedy Solutions

The greedy approach is the primary methods that we have used to obtain the upper bound and the lower bound of the Knapsack problem. There are four type of sorts used to sort the list of items: Sort by highest-value first, sort by lowest-cost-first, sort by highest ratio first, and partial knapsack. However, since it is impossible to determine if the result returned from one of the approaches is the best one, all four approaches mentioned above were used to obtain the most optimal result with the cost and value being at its maximum.

#### Sort by Lowest-cost First.

*Code Snippet:*

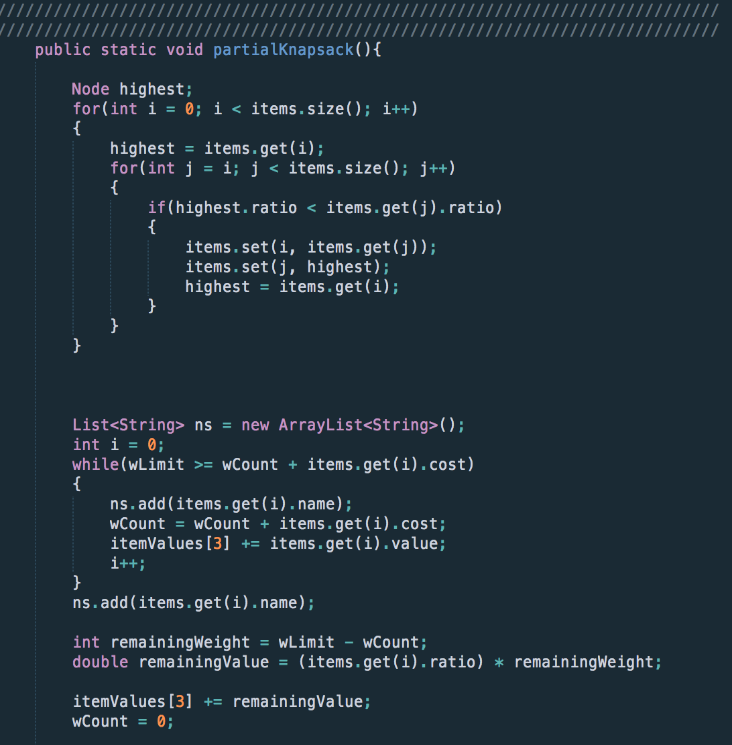


This solution sorts the given list of items from the lowest cost to the highest. By doing so, we are able to keep the cost the lowest for as many iterations as possible, allowing more space for more items to be included in the knapsack. Although this method would work for a list of items with low cost and high values, this method does not produce the optimal answer when the list of items contains both low cost and low values. There were cases where a single item would have a significantly higher value than the result obtained from this greedy approach.

#### Partial Knapsack

This greedy approach first sorts the given list of inputs by the highest cost to value ratio then grabs as many items as possible until the cost is near or equal to the limit that the problem is being constrained to. If the cost is less the constrained limit then we will continue to add part of the cost and value of the next item until we have the cost at its maximum.

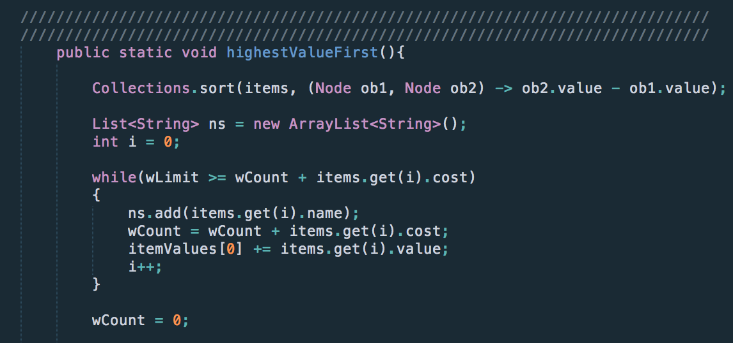
*Code Snippet:*



#### Sort by Highest-Value First

This solution first sorts the given set of items into a collection of highest values first and then sums the values from the top of the list with the highest values while keeping the cost within the range of the limit that the problem is constrained to. Although this algorithm prioritizes the highest values in the list first, there were cases where a more optimal answer was found by looking at the items with lower values with lower cost.

*Code Snippet:*



#### Sort by highest ratio first

This greedy approach sorts the given inputs by the highest ratio first allowing the algorithm to grab the item with the highest cost to value ratio.

*Code Snippet:*



### State Space Search

The state space is an approach that we have used to obtain every possible combination of the given list of items. Thus, allowing us to find the best possible combinations of the items that have its cost at its maximum and the values at its highest.

#### “Dumb” Exhaustive Search

Using a “dumb” exhaustive search, we looked at every possible leaf node of the tree. While looking through each leaf, we calculated the total cost and total value. The last row in the tree consists of all possible combination of items so if we look for the highest total value within the cost limit, clearly the best solution will be found.

#### “Smart” Search

The way we built the smart search will probably be a little different from how most people did it and it involves the way the tree was built. Initially, the tree was built with each node being a certain combination of items. However, with the smart search, the tree was built so that nodes that had already gone over capacity would receive a flag that would indicate that no calculation was needed to be done. This decreased would decrease the time taken to search through leaf nodes for the best solution to the knapsack problem.

# ADDITIONAL OPTIMIZATIONS AND RESULTS

In this section, we will discuss the improved speed on some other method that we have tried to shorten the search time.

## Arrays of Strings Versus Nodes

As mentioned in section 2, the entire state-space was planned to be built with the linked-list nodes approach. However, there were issues that we encountered from that approach: memory and time issues.

By following this approach, a given set of input items with the size of N produces a tree with the size of . An enormous amount of memory would be consumed when N > 18. Consider a regular node in a linked list, a regular node that contained a left and a right pointer takes up about 8 bytes. If we include a string and 2 integers into the node, it all adds up to (number of character in the string \* 2 + 2 + 2 + 8) about 14 bytes per node which total up to about 469762020 bytes = 0.46976202 GB when N = 24.

No calculations were included yet on what had mentioned above. Considering the function we have built has a complexity of O(), time is the other issue that this approach produces.

### Alternative to Trees with Nodes in Linked-list

*Code Snippet:*



The approach that we have taken in an alternative to building a tree which consists of linked-list nodes, we chose to build the tree using an array that consists of only ‘chars’ in each index of the array. Each index of the array represents different nodes in a visualized form of the tree. We are able to move to the left sub-tree by determining if the index is odd (left branch) or even (right branch) as shown in the code included below:

By doing so, it significantly reduced the amount of time taken to compile as compared to building a tree with nodes.

### Calculations After Tree is Built

To shortened the time spent on calculating the cost and values in each of the sub-trees, we’ve decided to perform the calculations after the tree is successfully built. From that, we are able to only look at the leaf indexes of the state-space tree which consists of all combinations since the root index, reducing more than half of the calculations needed to be done as compared to building the tree while calculating at every index.

# RESULTS

Table 1. Results of All Optimizations

|  |  |  |
| --- | --- | --- |
| **Optimization Methods** | **Time taken** | **Memory issue** |
| “Dumb” Exhaustive Search | 25(seconds) | None |
| Greedy | < 1 seconds | None |
| Tree with Nodes (Including Search) | Over a minute | Compile Error; insufficient  ram |
| Tree with Array (Including Search) | 22.43(seconds) | None |

As we can see from Table 1, the time taken to load trees with nodes was significantly longer than the tree with array. The Node class takes up quite a bit of memory, especially when making a tree with Nodes. The Greedy method is seen to have considerably fast search times. The obvious reason for that is because no tree was needed to be built nor requires to search through a huge state space.

# FUTURE OPTIMIZATION

Due to the time constraint of this Knapsack research project, there were optimization ideas that we could not implement in time but thought that it might be worth the time to dig-deeper into. To further optimize the huge memory usage used in order to build a -2) sized tree, while building the tree, we could cache the two most recent rows as row A and row B then build the next row using row B and then assign the new row as row A. Continue to iterate until all combinations are found. By doing so, we are not wasting more space than the program needs to. Therefore, opening up more space for larger lists of items. Another method that would prove to be even more efficient in terms of memory usage would be to delete one row as we are building the other row. For example, as we are building row B using the information from row A, we are also deleting used nodes from row A. This would greatly shorten the cost of memory by 50%. However, the time taken to run the program would significantly increase as well.

# REFERENCES

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