

VALGRIND

Linux x86 memory debugger and performance profiler

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Agenda

Introduction
Detecting memory errors
Heap profiling and leak detection
Data concurrency race conditions
Profiling applications
Summary

What is Valgrind?

- An open source system memory debugger: Initiated by Julian Sward on 2000; released for Linux x86 on 2002 and winner of Open Source Award of 2004.
- <u>Simple and easy to use:</u> does not require recompilation or re-linking: performs dynamic runtime instrumentation running on a synthetic CPU.
- <u>Used to validate many large Projects:</u> KDE, Emacs, Mozilla, OpenOffice and ... SAP HANA.
- Programing language agnostic: used for C/C++, Python, Java, JavaScript development.

\$ valgrind --tool=memcheck \ prog <args>

Uses
LD_PRELOAD
to load first

Load memcheck tool to handle instrumentation

Memcheck tool (vgtool_memcheck.so)

Traps program and runs it on synthetic CPU

prog

Shared objects

Memcheck tool (vgtool_memcheck.so)

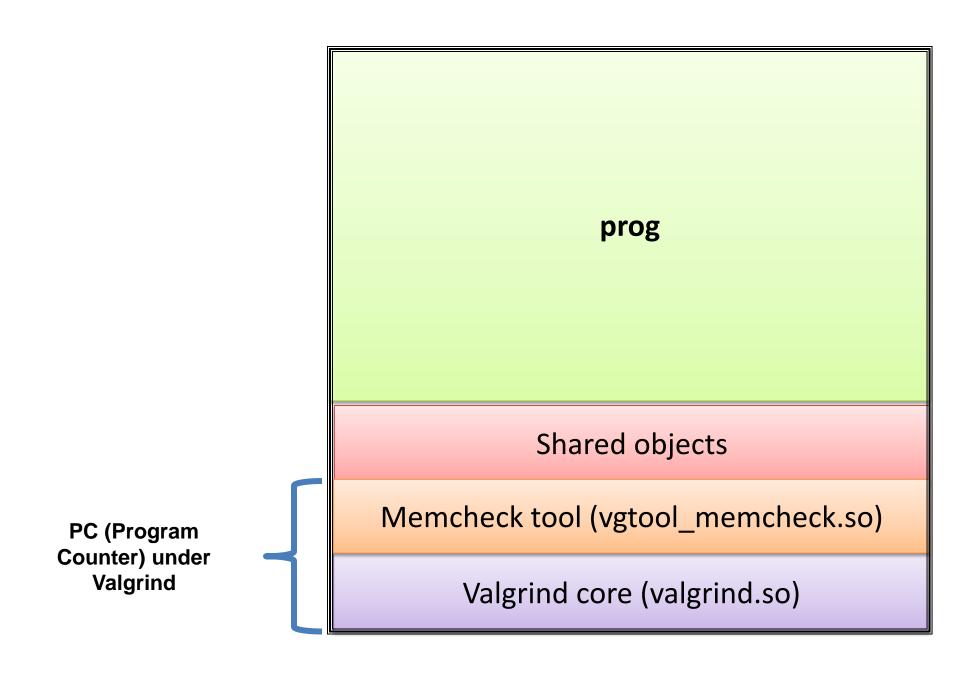
Appears as a normal application

prog

Shared objects

Memcheck tool (vgtool_memcheck.so)

prog **Single Name Space Shared objects** Memcheck tool (vgtool_memcheck.so) Valgrind core (valgrind.so)



prog **PC** (Program Counter) when exists Valgrind **Shared objects** Memcheck tool (vgtool_memcheck.so) Valgrind core (valgrind.so)

Dynamic Binary Instrumentation

"UCode lies at the heart of the x86-to-x86 JITter. The basic premise is that dealing the the x86 instruction set head-on is just too darn complicated, so we do the traditional compiler-writer's trick and translate it into a simpler, easier-to-deal-with form."

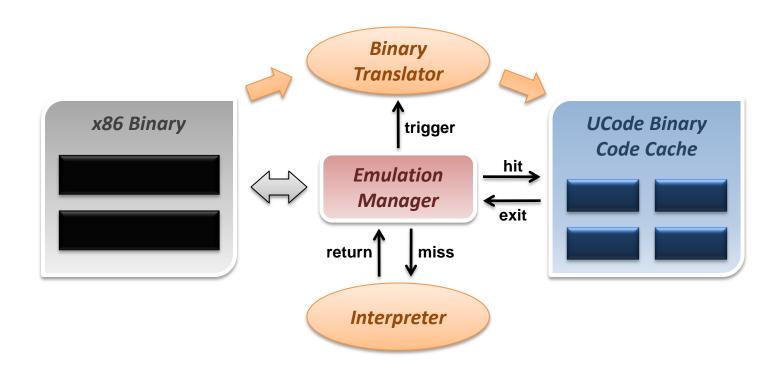
Julian Sward

The Instrumenting JITer

- Parse of an x86 basic block into a sequence of UCode instructions.
- UCode optimization with the aim of caching simulated registers in real registers.
- UCode instrumentation which adds value and address checking code.
- Post-instrumentation cleanup, removing redundant valuecheck computations.
- Register allocation done on UCode.
- Emission of final instrumented x86 code.

Dynamic Binary Translator

- 1. First time execution, no translated code in code cache.
- 2. Miss code cache matching, then directly interpret the guest instruction.
- 3. As a code block discovered, trigger the binary translation module.
- 4. Translate guest code block to host binary, and place it in the code cache.
- 5. Next time execution, run the translated code block from the code cache.



Installing Valgrind

- \$ zypper install valgrind
- \$ yum install valgrind
- \$ apt-get install valgrind

```
$ zypper install -y linux-kernel-headers
```

- \$ wget http://www.valgrind.org/downloads/valgrind-3.9.0.tar.bz2
- \$ tar xvfj valgrind-3.9.0.tar.bz2
- \$ cd valgrind-3.9.0
- \$./configure
- \$ make
- \$ make install

Memcheck

Memcheck

- Use of uninitialized memory
- Reading/writing memory after it has been freed
- Reading/writing off the end of malloc'd blocks
- Reading/writing inappropriate areas on the stack
- Memory leaks -- where pointers to malloc'd blocks are lost forever
- Passing of uninitialized or unaddressible memory to system calls
- Mismatched use of malloc/new/new[] vs free/delete/delete[]
- Some misuses of the POSIX pthreads API

- A-bits: every memory byte is shadowed by a single A (Addressability) bit indicating if the application can legitimately access the byte.
- <u>V-bits:</u> every memory byte is shadowed by eight V (Validity) bits indicating if the values of each bit are defined.
- Heap blocks: records heap blocks in auxiliary hash, enabling to track double-free and other miss-use.

```
$ valgrind --tool=memcheck \
    --num-callers=20 \
    --log-file=vg.log \
    --trace-children=yes \
    prog <args>
```

```
int v, *p = malloc(10 * sizeof(int));
v = p[10]; /* invalid read (heap) */
```

```
int a[10];
```

Not Detected!

```
int a, b;
b = a; /* pass uninitialized */
if (!b) /* use uninitialized */
    printf("non deterministic\n");
```

```
char *p = malloc(100); /* uninitialized */
write(1, p, 100); /* to syscall */
```

```
unsigned char c;
c |= 0x1; /* only 1<sup>st</sup> bit is initialized */
if (c & 0x1) /* bit-1 initialized */
     printf("always be printed\n");
if (c & 0x2) /* bit-2 uninitialized */
     printf("non deterministic\n");
```

Addrcheck

Addrcheck

- Use of uninitialized memory
- Reading/writing memory after it has been freed
- Reading/writing off the end of malloc'd blocks
- Reading/writing inappropriate areas on the stack
- Memory leaks -- where pointers to malloc'd blocks are lost forever
- Passing of uninitialized or unaddressible memory to system calls
- Mismatched use of malloc/new/new[] vs free/delete/delete[]
- Some misuses of the POSIX pthreads API

```
$ valgrind --tool=addrcheck \
    prog <args>
```

Performance slowdown

Program	Time (s)	Nulgrind	Memcheck	Addrcheck	Cachegrind
bzip2	10.7	2.4	13.6	9.1	31.0
crafty	3.5	7.2	44.6	26.5	107.4
gap	0.9	5.4	28.7	14.4	46.6
gcc	1.5	8.5	36.2	23.6	73.2
gzip	1.8	4.4	20.8	14.5	50.3
mcf	0.3	2.1	11.6	5.9	18.5
parser	3.3	3.7	17.4	12.5	34.8
twolf	0.2	5.2	29.2	18.5	53.3
vortex	6.5	7.5	47.9	32.7	88.4
ammp	18.9	1.8	24.8	21.1	47.1
art	26.1	5.9	14.1	11.5	19.4
equake	2.1	5.5	32.7	28.0	49.9
mesa	2.7	4.7	41.9	31.6	64.5
median		5.2	28.7	18.5	49.9

Memory increase

Program	Size (KB)	Nulgrind	Memcheck	Addrcheck	Cachegrind
bzip2	34	5.2	12.1	6.8	9.1
crafty	156	4.5	10.9	5.9	8.2
gap	140	5.6	12.7	7.3	9.7
gcc	564	5.9	13.1	7.6	9.9
gzip	30	5.5	12.6	7.2	9.4
mcf	30	5.7	13.5	7.7	9.9
parser	97	6.0	13.6	7.8	10.1
twolf	114	5.2	12.2	7.0	9.3
vortex	234	5.8	13.2	8.1	10.1
ammp	68	4.7	11.7	7.1	9.5
art	24	5.5	13.0	7.5	9.8
equake	44	5.0	12.2	7.1	9.2
mesa	69	4.8	11.2	6.7	8.9
median		5.5	12.6	7.2	9.5

Getting Accurate Backtraces

gcc main.c && strip a.out

```
==23580== Invalid write of size 1
--23580== at 0x8048622: (within /root/dev/demo/test3)
==23580==
            by 0x8048703: (within /root/dev/demo/test3)
==23580==
            by 0x42015573: __libc_start_main (in /lib/tls/libc-2.3.2.so)
==23580==
            by 0x80484A0: (within /root/dev/demo/test3)
==23580== Address 0x3C03502E is 0 bytes after a block of size 10 alloc'd
==23580==
            at 0x3C01E250: malloc (vg replace malloc.c:105)
==23580==
            by 0x8048615: (within /root/dev/demo/test3)
==23580==
            by 0x8048703: (within /root/dev/demo/test3)
==23580==
            by 0x42015573: __libc_start_main (in /lib/tls/libc-2.3.2.so)
```

gcc main.c

```
==23587== Invalid write of size 1

--23587== at 0x8048622: invalid_write (in /root/dev/demo/test3)

==23587== by 0x8048703: main (in /root/dev/demo/test3)

==23587== Address 0x3C03502E is 0 bytes after a block of size 10 alloc'd

==23587== at 0x3C01E250: malloc (vg_replace_malloc.c:105)

==23587== by 0x8048615: invalid_write (in /root/dev/demo/test3)

==23587== by 0x8048703: main (in /root/dev/demo/test3)
```

gcc -g main.c

```
==23594== Invalid write of size 1

--23594== at 0x8048622: invalid_write (demo_memcheck.c:34)

==23594== by 0x8048703: main (demo_memcheck.c:68)

==23594== Address 0x3C03502E is 0 bytes after a block of size 10 alloc'd

==23594== at 0x3C01E250: malloc (vg_replace_malloc.c:105)

==23594== by 0x8048615: invalid_write (demo_memcheck.c:31)

==23594== by 0x8048703: main (demo_memcheck.c:68)
```

Error Suppression

```
$ valgrind --tool=memcheck \
    --gen-suppressions=yes \
    --logfile=my.supp \
    prog <args>
```

```
{
      <gdk_set_locale>
          Memcheck:Cond
           ...
      fun:gdk_set_locale
}
```

```
{
     libpango>
          Memcheck:Leak
          ...
     obj:/usr/*lib*/libpango*
}
```

```
$ valgrind --tool=memcheck \
    --suppressions=my.supp.* \
    prog <args>
```

Attaching a Debugger

```
#include <malloc.h>
int main()
      int i;
      char *a = malloc(3);
      for (i=0; i<4; i++)
            a[i] = 0;
```

```
$ valgrind --tool=memcheck \
    --db-attach=yes \
    ./prog <args>
```

```
==3513== Memcheck, a memory error detector
==3513== Copyright (C) 2002-2010, and GNU GPL'd, by Julian Seward et al.
==3513== Using Valgrind-3.6.0.SVN-Debian and LibVEX;
==3513== Command: ./a.out
==3513==
==3513== Invalid write of size 1
==3513== at 0x40051C: main (in /home/ortyl/a.out)
==3513== Address 0x51b1043 is 0 bytes after a block of size 3 alloc'd
==3513== at 0x4C2815C: malloc (vg_replace_malloc.c:236)
==3513== by 0x400505: main (in /home/ortyl/a.out)
==3513==
==3513==
==3513== ---- Attach to debugger ? --- [Return/N/n/Y/y/C/c] ---- y
==3513== starting debugger with cma: /usr/pin/gap -nw /proc/3516/td/1014 3516
(gdb)
```

Leak Check

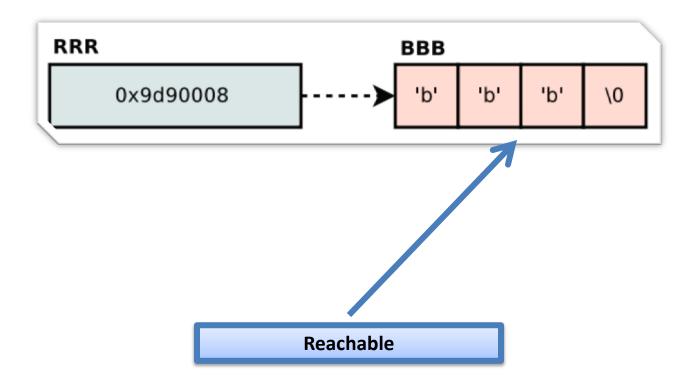
```
$ valgrind --tool=memcheck \
    --leak-check=yes \
    --show-reachable=yes \
    prog <args>
```

HEAP SUMMARY: in use at exit: 4 bytes in 1 blocks total heap usage: 1 allocs, 0 frees, 4 bytes allocated 4 bytes in 1 blocks are still reachable in loss record 1 of 1 at 0x4024C1C: malloc (vg_replace_malloc.c:195) by 0x40B0CDF: strdup (strdup.c:43) by 0x804879B: main (in /home/aleksander/valgrind-memcheck)

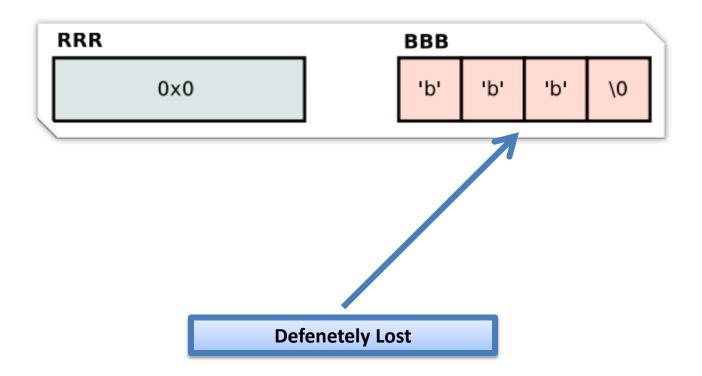
LEAK SUMMARY:

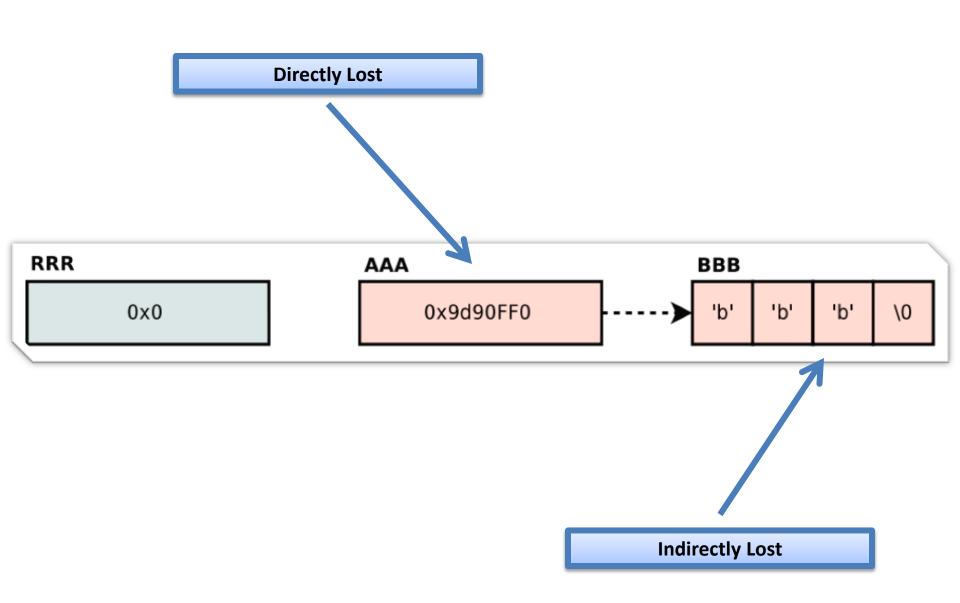
definitely lost: 0 bytes in 0 blocks indirectly lost: 0 bytes in 0 blocks possibly lost: 0 bytes in 0 blocks still reachable: 4 bytes in 1 blocks suppressed: 0 bytes in 0 blocks

```
void *rrr;
int main(void)
     rrr = strdup("bbb");
     return 0;
```

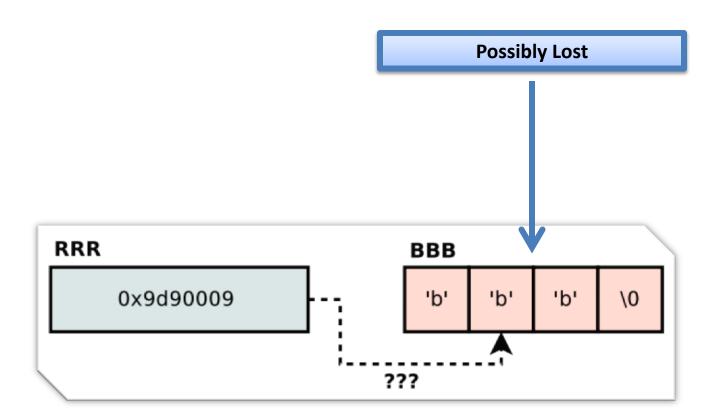


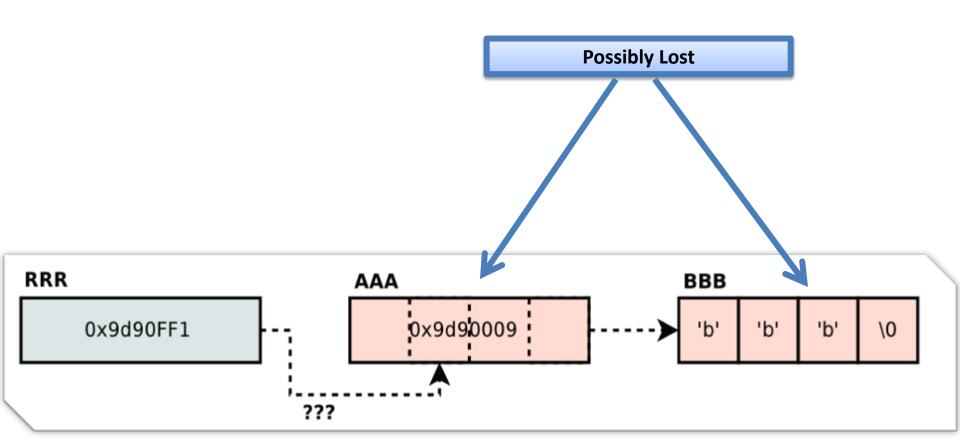
```
void *rrr;
int main(void)
     rrr = strdup("bbb");
     rrr = NULL;
     return 0;
```





```
void *rrr;
int main(void)
      rrr = strdup("bbb");
      rrr = ((char *) rrr) + 1;
      return 0;
```





```
$ valgrind --tool=memcheck \
    --vgdb-error=0 --vgdb=yes \
    prog <args>
```

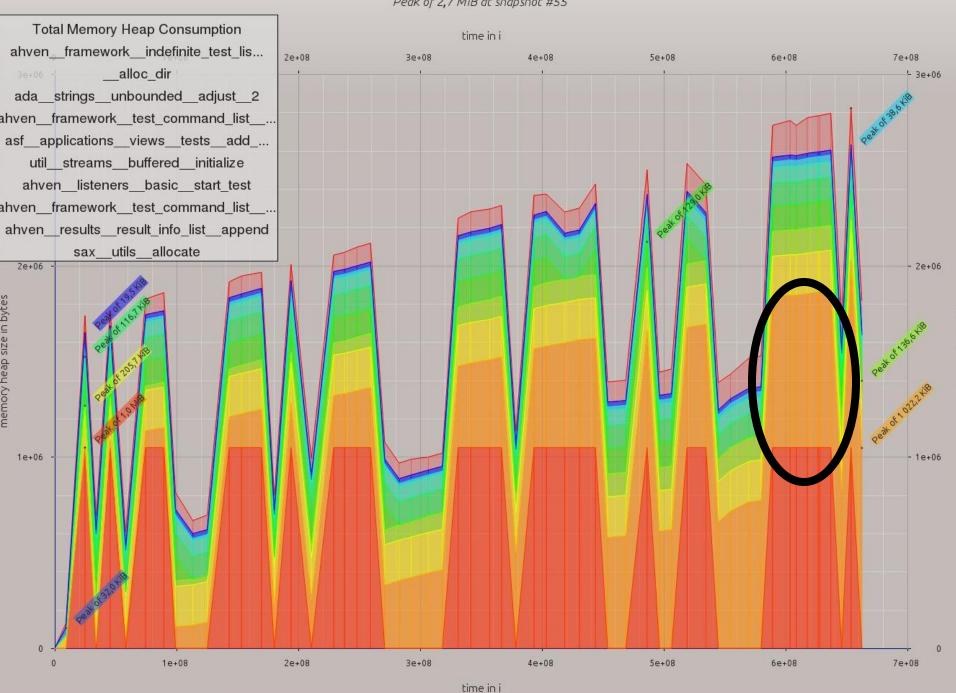
```
(gdb) target remote | vgdb
(gdb) monitor leak check full reachable any
==2418== 100 bytes in 1 blocks are still reachable in loss record 1 of 1
==2418== at 0x4006E9E: malloc (vg_replace_malloc.c:236)
==2418== by 0x804884F: main (prog.c:88)
==2418==
==2418== LEAK SUMMARY:
==2418== definitely lost: 0 bytes in 0 blocks
==2418== indirectly lost: 0 bytes in 0 blocks
==2418== possibly lost: 0 bytes in 0 blocks
==2418== still reachable: 100 bytes in 1 blocks
==2418== suppressed: 0 bytes in 0 blocks
==2418==
(gdb)
```

Massif

Massif is a heap profiler: it measures how much heap memory your program uses. This includes both the useful space, and the extra bytes allocated for book-keeping and alignment purposes. It can also measure the size of your program's stack(s), although it does not do so by default.

Memory consumption of bin/asf_harness

Peak of 2,7 MiB at snapshot #55



Helgrind

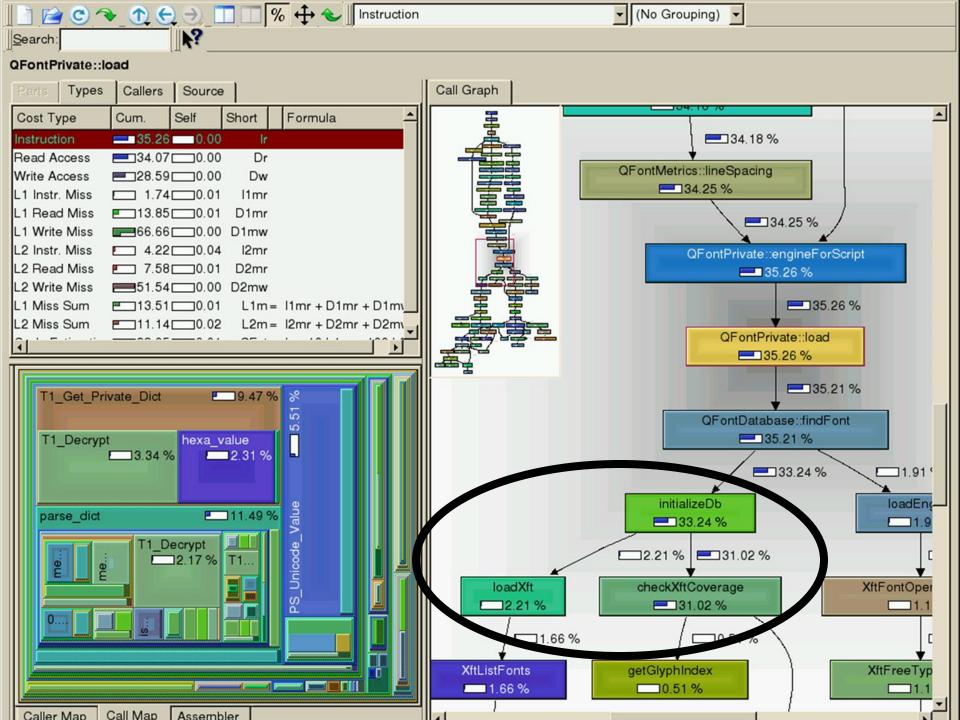
Helgrind is a thread debugger: finds data races in multithreaded programs. It looks for memory locations which are accessed by more than one (POSIX p-) thread, but for which no consistently used (pthread_mutex_) lock can be found.

```
static void *inc_shared (void *v) {
      shared++; /* un-protected access to shared */
      return 0;
pthread ta, b;
pthread create(&a, NULL, inc shared, NULL);
pthread create(&b, NULL, inc shared NULL);
```

Callgrind

Callgrind is a profiling tool: it records the call history among functions in a program's run as a call-graph. Collecting number of instructions executed, their relationship to source lines, the caller/callee relationship between functions, and the numbers of such calls.

- \$ valgrind --tool=callgrind prog <args>
- \$ callgrind_control --zero # zero counters
- ... the application runs ...
- \$ callgrind_control --dump # dump counters
- \$ kcachegrind cachegrind.out.<pid>



Profiling golden rules

- A human has very seldom a clue at all where in the code time is spent, therefore you must use a profiler.
- You should never optimize code unless the application feels slow as optimization always leads to code that is harder to maintain.
- Make sure to profile your application in a realistic context.

Eclipse



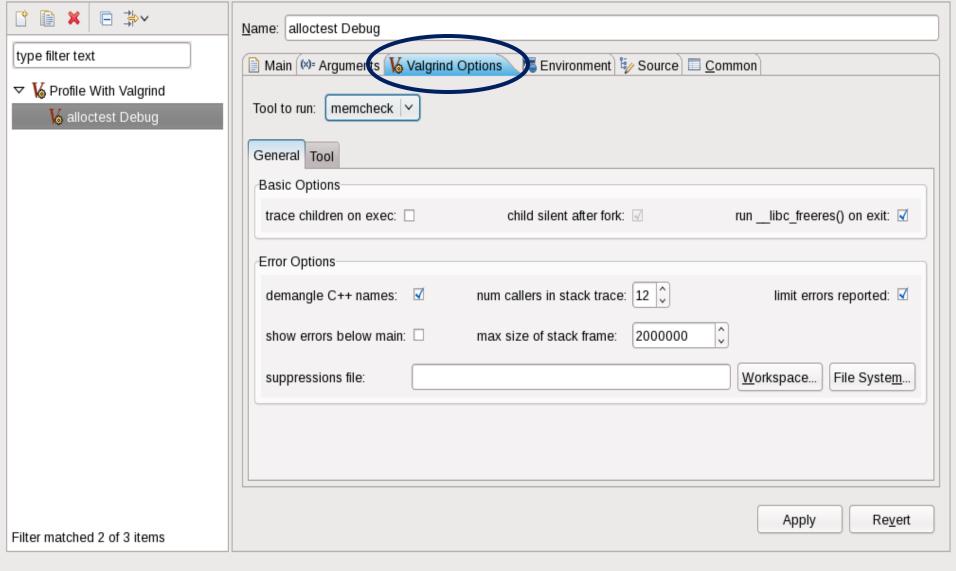


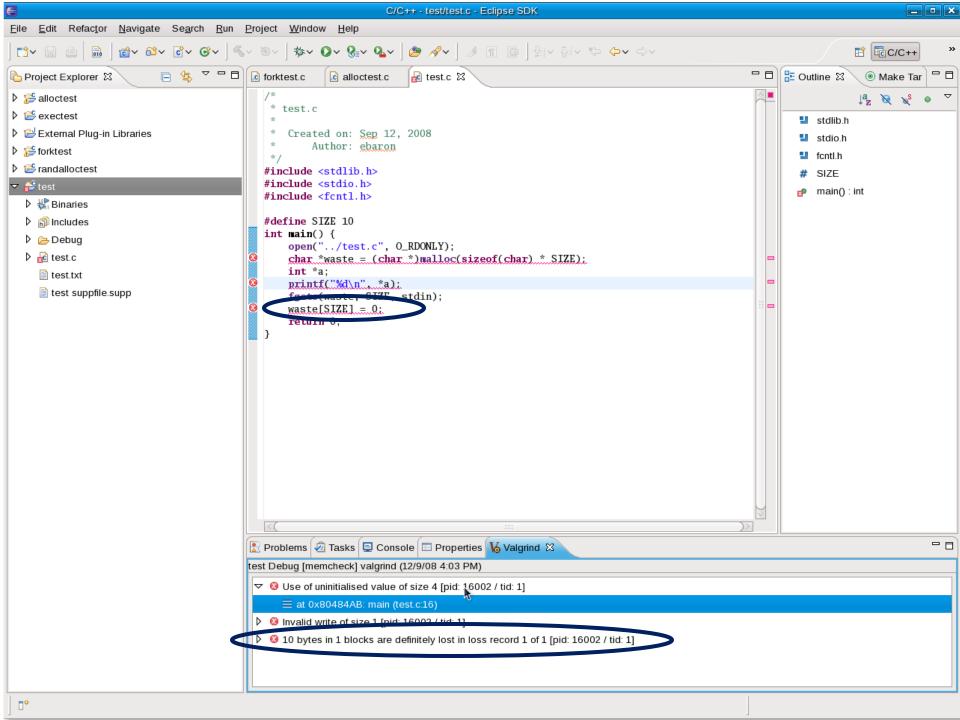
Х

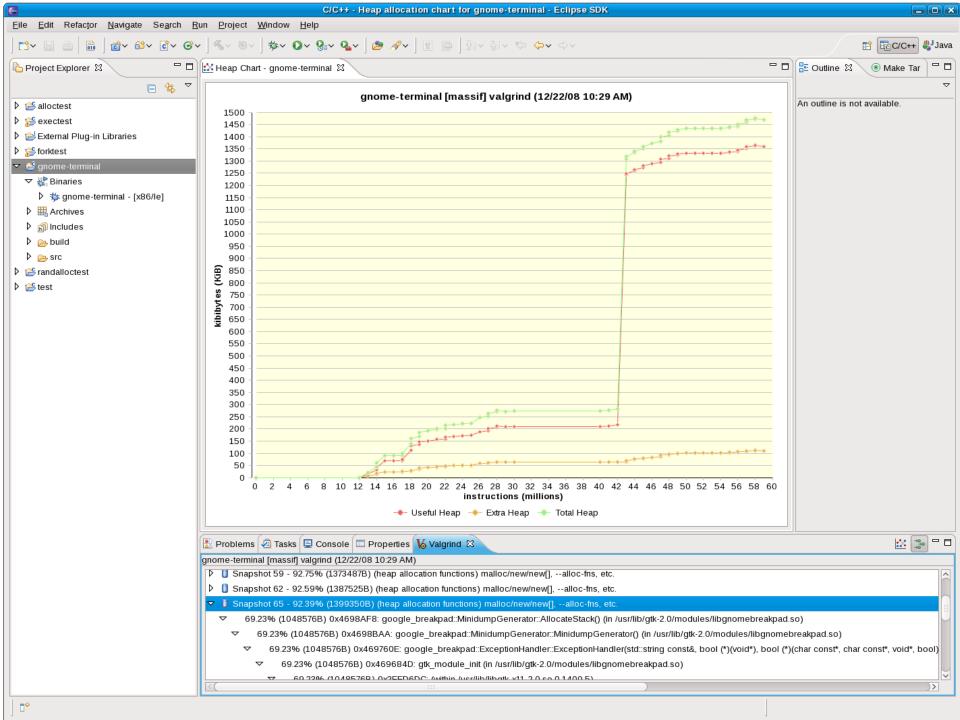
Profile C/C++ Application Using Valgrind

Create, manage, and run configurations









Language Support

Language Survey

• C: 56

• C+: 52

• Fortran: 6

• Java: 3

• asm: 3

Python: 2

• TCL/TK: 1

Objective C: 1

Others: 3

Python

\$ cd python/dist/src

```
$ valgrind --tool=memcheck \
    --suppressions=Misc/valgrind-python.supp \
    ./python -E -tt ./Lib/test/regrtest.py -u \
    bsddb,network
```

Node.js

\$ valgrind --leak-check=yes \
 node foo.js

Java JNI

```
$ valgrind --trace-children=yes \
    --leak-check=full \
    java -Djava.library.path=$(PWD) Foo
```

one slide to go ... ©

- Valgrind is a robust x86-Linux memory debugger.
- Using it regularly enhances product's quality.
- When to use it?
 - All the time.
 - In automatic testing.
 - After big changes.
 - When a bug occurs.
 - When a bug is suspected.
 - Before a release.

THANK YOU!

Valgrind

www.valgrind.org

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References

- Aleksander Morgado. Understanding Valgrind memory leak reports. February 4, 2010
- Nicholas Nethercote and Julian Seward. How to Shadow Every Byte of Memory Used by a Program. 2007.
- Nicholas Nethercote and Julian Seward.
 Valgrind: A Program Supervision Framework.
 2003.