Lecture 3: Introduction to Algorithms

BT 3051 - Data Structures and Algorithms for Biology

Karthik Raman

Department of Biotechnology Bhupat and Jyoti Mehta School of Biosciences Indian Institute of Technology Madras

Introduction

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 - Outcomes: Things asserted to be true after the algorithm is executed
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 - Runtime: Time required to execute the algorithm expressed as asymptotic estimate as a function of input size
 - Run space: Space required to execute the algorithm

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History of Algorithms^a

The Vedas
Eratosthenes Greek
al-Biruni Pythagoras Babylonians Chinese
John von Neumann Arabs Hypatia of Alexandria
Al-Khwarizmi Diophantos of Alexandria Egyptians
Heron of Alexandria Aryabhatta Alan Mathison Turing
Ibn Al-Haitham Klaudios Ptolemaeus
Archimedes Alonzo Church Thales of Miletus
Euclid Indians Brahmagupta
Mayans Ada Lovelace

ahttp://cs-exhibitions.uni-klu.ac.at/index.php?id=193

History of Algorithms



"For me, great algorithms are the poetry of computation. Just like verse, they can be terse, allusive, dense, and even mysterious. But once unlocked, they cast a brilliant new light on some aspect of computing."

— Francis Sullivan

Top 10 Algorithms of the 20th Century

- ► Monte Carlo/Metropolis
- Simplex method (LP)
- Krylov Subspace Iteration
- Householder MatrixDecomposition
- ► Fortran Compiler

- QR algorithm
- Quicksort
- ▶ FFT
- Integer Relation Detection
- ► Fast Multipole

Self-assessment Exercise

- Read through the top ten algorithms of the century
- ► Write a paragraph (300–400 words) about one of these algorithms, highlighting the greatness (coolness!) of the algorithm and its applications
- Outcome
 - ► Inspiration!
 - Technical writing practice!

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 - An algorithm must terminate! (Finiteness)
 - An algorithm must be clearly defined (Definiteness)
 - An algorithm must produce the correct result (Correctness
 - Should work on a defined class of inputs, to produce the expected output
 - Algorithms must produce an output.

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Examples

EXAMPLES: FIBONACCI NUMBERS

$$F_1 = 1$$

 $F_2 = 1$
 $F_n = F_{n-1} + F_{n-2}$, for all $n > 2$

- ► Fibonacci numbers appear in Indian and Western math
 - Named after Leonardo Fibonacci (book dating 1202 CE)
- Fibonacci numbers are intimately connected with the golden ratio
- They also appear in biological settings:
 - branching in trees
 - arrangement of leaves on a stem
 - **...**

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How do we compute F_n ?

Closed form solution exists:

$$F_n = \frac{1}{\sqrt{5}} (\phi^n - \hat{\phi}^n)$$

where
$$\phi=\frac{1+\sqrt{5}}{2}$$
 and $\hat{\phi}=\frac{1-\sqrt{5}}{2}$

but ...

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Naïve Method using Recurrence fib_v1.py

```
def fibo(n):
    if n == 1 or n == 2:
        return 1
    else:
        return fibo(n - 1) + fibo(n - 2)
print fibo(10)
```

Output:

```
>>>
55
```

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How much computation does it involve?

Iterative Method

```
def fibo(n):
    fib = [0] * n
    fib[0] = 1
    fib[1] = 1

for i in range(2, n):
        fib[i] = fib[i - 1] + fib[i - 2]

return fib[n - 1]

print fibo(10)
```

Output:

```
>>> 55
```

Iterative Method fib_v2.py

```
def fibo(n):
    fib = [0] * n
    fib[0] = 1
    fib[1] = 1

for i in range(2, n):
        fib[i] = fib[i - 1] + fib[i - 2]

return fib[n - 1]

print fibo(10)
```

Output:

```
>>> 55
```

What's the drawback?

Improved Iterative Method fib_v3.py

```
def fibo(n):
    fibn = 1 # F_n
    fibn1 = 1 # F_{n-1}
    for dummy in range(2,n):
        fib = fibn + fibn1
        fibn1 = fibn
        fibn = fib
    return fib
print(fibo(10))
```

Output:

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def fibo(n):
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Output:

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Trick!

Consider the matrix A:

$$A = \left[\begin{array}{cc} 1 & 1 \\ 1 & 0 \end{array} \right]$$

For $n \geq 2$,

$$A^{n-1} = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}^{n-1} = \begin{bmatrix} F_n & F_{n-1} \\ F_{n-1} & F_{n-2} \end{bmatrix}$$

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So what?

EXAMPLES: EXPONENTIATING A

NUMBER

How to compute x^n ?

Algorithm 1: Naïve exponentiation

```
input : x, n

output: p = x^n

begin

p \longleftarrow 1

for x = 1 to n do

p \longleftarrow p * x
```

How to compute x^n ?

Algorithm 2: Naïve exponentiation

```
input : x, n

output: p = x^n

begin

p \leftarrow 1

for x = 1 to n do

p \leftarrow p * x
```

Can we do better?

How to compute x^n faster?

```
def fastExpo(x, n):
    if n == 1:
        return x
    if n == 2:
        return x * x
    if n % 2 == 0:
        p = fastExpo(x, n // 2)
        return p * p
    else:
        return x * fastExpo(x, n - 1)
```

How to compute x^n faster?

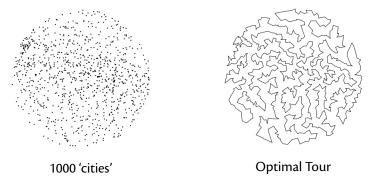
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```

Can we do better?

EXAMPLES: TRAVELLING SALESPERSON PROBLEM

Travelling Salesperson Problem (TSP)

Given *N* cities and the distances between every pair of cities, the goal of a travelling salesperson is to visit all of cities exactly once (and return to the origin) while keeping the total distance travelled as short as possible.



http://www.cs.princeton.edu/courses/archive/fall14/cos126/assignments/tsp.html

Travelling Salesperson Problem

The importance of the TSP does not arise from an overwhelming demand of salespeople to minimize their travel distance, but rather from a wealth of other applications such as vehicle routing, circuit board drilling, VLSI design, robot control, X-ray crystallography, machine scheduling, and computational biology.

Robert Sedgewick (Princeton COS 126)

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How to solve this problem?

Introduction

Examples