# Networks: Lectures 5 & 6 Toy Models of Network Formation

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HT 2017

# What you're in for

- Week 1: Introduction and basic concepts
- Week 2: Small worlds
- Week 3: Toy models of network formation
- Week 4: Additional summary statistics and other concepts
- Week 5: Random graphs
- Week 6: Community structure and other mesocopic structures
- Week 7: Dynamical systems on networks
- Week 8: Other topics TBD

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#### Motivation

- In addition to direct measurement of network properties (we've seen some of this and will see more), it is important to study networks using *generative models* (we've seen a bit of this in week 2 as well).
- We'll focus on *preferential attachment* (PA) models as an example.
- PA is a "rich-get-richer" scheme (e.g., highly cited papers are easier to find via Google, which in turn leads them to get even more citations).

#### Attachment kernel

- An attachment kernel gives a rule for how new objects in a network connect to existing objects.
- Suppose that we add one node at a time and that it forms an edge with a single existing node (there must also be a seed network in place or an initial condition).

#### Attachment kernel

New node connects to existing node i with probability

$$q_i = \frac{a_i}{\sum a_i};$$

 $a_i$  is the attachment kernel.

If  $q_i$  depends directly only on the properties of node i, we write:  $a_i = a_i(\alpha_i, \eta_i)$ , where

- $\alpha_i$  = vector of structural properties (e.g., degree, local clustering coefficient, etc).
- $\eta_i$  = vector of other properties (e.g., a fitness value assigned to node i to represent different attributes).

#### Attachment kernel

More generally, could consider beyond node i (e.g., nearest neighbors) or exogenous to the network:  $a_i = a_i(\alpha, \eta, X)$  where

- $\alpha_{ij} = \text{matrix of structural properties being considered}$
- $\eta_{ij} = \text{matrix of other (e.g., "fitness") properties}$
- $X_{ij} = \text{matrix of exogenous properties}$

One could make even more complicated kernels, by including time-dependence. Often such generality is not explored (much simpler situations can exhibit rich behavior).

Note: There has been work on attachment mechanisms with > 1 edge added at a time, rewiring existing nodes, removing edges etc.

### Seeds

To define a toy attachment model, we also need to specify a seed network or an initial condition.

Q: What happens to the "memory" of a seed as the size N of a network increases?

# Linear preferential attachment

Consider an unweighted, undirected network (without self-edges or multi-edges).

Linear preferential attachment by degree has  $a_i = k_i$ .

$$\Rightarrow q_i = \frac{k_i}{\sum_i k_i} = \frac{k_i}{zN}$$

where  $z = \langle k \rangle$  is mean degree.

Can write versions of this for directed networks by considering in- and out-degree separately.

Concerned with citation network, so need to consider a directed network (new paper cites papers that already exist)

#### Network set-up.

- node: paper
- edge: a citation (directed edge from newer paper to older one).
- nodes cannot be removed, so this model gives DAGs.
- let c := mean number of papers cited by new paper (i.e., mean out degree).

In Price's model, we add a node at every time-step. Each new node has on average c outgoing edges.

In this model,

$$q_i = \frac{k_i + a}{\sum_i (k_i + a)},$$

where this  $q_i$  is an affine preferential attachment, and the new bonus parameter a > 0 allows papers without citations to eventually be cited.

Note: a paper can cite another paper multiple times in this model.

Let's write down equations that govern the in-degree (i.e., number of citations) of nodes in terms of c and a. Let  $k_i = k_i^{in}$  denote the in-degree of node i.

Let  $p_k(N) :=$  fraction of nodes with in-deg k when network has N nodes.

Each existing node i attracts an incoming node with probability proportional to its in-degree  $k_i$  plus a, an initial fitness:

$$q_i = \frac{k_i + a}{\sum_i (k_i + a)} = \frac{k_i + a}{Na + N\langle k \rangle} = \frac{k_i + a}{N(a + c)},$$

where  $z = \frac{1}{N} \sum_i k_i = \langle k \rangle$ . Since each new paper cites c papers on average, the expected number of new citations to node i when a new node is added is  $c\left(\frac{a+k_i}{N(a+c)}\right)$ .

 $Np_k(N)$  nodes have in-degree k, so the expected number new citations (i.e., incoming edges) to all nodes with in-degree k is

$$Np_k(N) \times c \times \frac{k+a}{N(c+a)} = p_k(N) \frac{c(k+a)}{c+a}.$$
 (1)

To study dynamics of Eq. 1, we write a master equation for the evolution of the in-degree.

When adding exactly one node to the N-node network, the number of nodes with in-degree k increases by 1 for every node that previously had degree k-1 and receives a new citation. The expected number of such nodes is

$$\frac{c(k-1+a)}{c+a}p_{k-1}(N)$$

Similar terms for nodes that go from in-degree k to in-degree k+1 (i.e. Eq. 1).



The evolution of the number of nodes with degree k > 0 is given by the discrete-time dynamical process:

$$(N+1)p_k(N+1) = \underbrace{Np_k(N)}_{[prev. \#k]} + \underbrace{p_{k-1}(N)\frac{c(k-1+a)}{c+a}}_{\#[(k-1)\to k]} - \underbrace{p_k(N)\frac{c(k+a)}{c+a}}_{\#[(k)\to k+1]}.$$

This is the expected number of nodes with in-degree k when the network has N+1 total nodes.

Note: the maximum degree obtainable at time N is c(N-1).

When k = 0 we have:

$$(N+1)p_0(N+1) = Np_0(N) + 1 - p_0(N)\frac{ca}{c+a}.$$

Note: can't go from  $-1 \to 0$  citations, so the corresponding term isn't present for k=0 case. The new node automatically has in-degree 0.

A master equation is an equation of evolution of the following form: new # things = old # things + ways to add more things - ways to remove things,for each state, e.g., a discrete set as above.

Consider  $N\to\infty$  to calculate an asymptotic form for the degree distribution (abuse of notation  $p_k=p_k(\infty)$  ).

For every N we can obtain  $p_k(N)$ , given  $p_k(0)$  by iteratively solving N equations.

There are also continuous-time versions of these equations, which are often convenient to use.

Case of large N

In the large-graph case (i.e., as  $N \to \infty$ ) we assume that  $p_k(N) \to p_k$  which gives

$$p_0 = 1 - p_0 \frac{ca}{c+a}, \quad k = 0$$

$$p_k = \frac{c}{c+a} \left[ (k-1+a)p_{k-1} - (k+a)p_k \right], \quad k \ge 1.$$

In general, we can't necessarily solve such an equation (might need to do numerical simulations of Eq. 1. Here, however we can rearrange and calculate the solutions:

$$p_0 = \frac{1 + \frac{a}{c}}{a + 1 + \frac{a}{c}},$$

$$p_k = \frac{k + a + 1}{k + a + 1 + \frac{a}{c}} p_{k-1}.$$

Case of large N

Solving iteratively for  $p_k$  we get:

$$p_k = \frac{\prod_{i=1}^k (k+a-i)}{\prod_{i=1}^k (i+1+a+\frac{a}{c})} \frac{1+\frac{a}{c}}{a+1+\frac{a}{c}} = \left(1+\frac{a}{c}\right) \frac{\prod_{i=1}^k (k+a-i)}{\prod_{i=0}^k (i+1+a+\frac{a}{c})}.$$

Recall the gamma function 1:  $\Gamma(x) = \int_0^\infty t^{x-1} e^{-t} dt$  with properties:

$$\Gamma(x+1) = x\Gamma(x) \quad \forall x > 0 \quad (\Rightarrow \Gamma(x) = (x-1)! \text{ for } x \in \mathbb{N}.)$$

$$\frac{\Gamma(x+n)}{\Gamma(x)} = (x+n-1)(x+n-2)\dots x, \text{ (Pochhammer's function)}.$$

Using Pochhammer's function, we can re-write terms of  $p_k$ :

$$\prod_{i=1}^k (k+a-i) = \frac{\Gamma(k+a)}{\Gamma(a)}, \quad \prod_{i=0}^k (i+1+a+\frac{a}{c}) = \frac{\Gamma(1+\frac{a}{c}+a+(k+1))}{\Gamma(1+\frac{a}{c}+a)},$$

$$\Rightarrow p_k = \left(1 + \frac{a}{c}\right) \frac{\Gamma(k+a)\Gamma(a+1+a/c)}{\Gamma(a)\Gamma(k+a+2+a/c)}$$

 $<sup>^1</sup>$ To see more properties of the gamma function see http://dlmf.mist\_gov/ $_{\sim}$ 

Another special function is the beta function

$$B(x,y) = \frac{\Gamma(x)\Gamma(y)}{\Gamma(x+y)}.$$

Using  $\Gamma(x+1) = x\Gamma(x)$  property, we let  $x = 1 + \frac{a}{c}$ , which gives

$$\Rightarrow \Gamma\left(2+\frac{a}{c}\right) = \left(1+\frac{a}{c}\right)\Gamma\left(1+\frac{a}{c}\right).$$

Multiply  $p_k$  by this expression on the numerator and denominator:

$$p_{k} = \underbrace{\left(1 + \frac{a}{c}\right)}^{\Gamma(k+a)\Gamma(a+1+a/c)} \cdot \underbrace{\frac{\Gamma\left(2 + \frac{a}{c}\right)}{\Gamma(a)\Gamma(k+a+2+a/c)}} \cdot \underbrace{\frac{\Gamma\left(k + \frac{a}{c}\right)}{\left(1 + \frac{a}{c}\right)\Gamma\left(1 + \frac{a}{c}\right)}}_{\Gamma(k+a+2+a/c)} \cdot \underbrace{\frac{\Gamma(a+1+a/c)}{\Gamma(a)\Gamma\left(1 + \frac{a}{c}\right)}}_{\Gamma(a)\Gamma\left(1 + \frac{a}{c}\right)},$$

$$= \frac{B(k+a, 2+a/c)}{B(a, 1+a/c)}.$$

Fix a and c, what happens as k gets large? We focus on numerator of  $p_k$ .

Case of large k

To better understand what B(k+a,2+a/c) looks like, we can use Stirling's formula  $\Gamma(x) \sim \sqrt{2\pi}e^{-x}x^{x-\frac{1}{2}}, x \to \infty$  which applied to the expression above gives

$$B(x,y) \sim \frac{e^{-x}x^{x-(1/2)}}{e^{-(x+y)}(x+y)^{x+y-(1/2)}}\Gamma(y), x \to \infty.$$

Now we use

$$(x+y)^{x+y-(1/2)} = x^{x+y-(1/2)} \left[ 1 + \frac{y}{x} \right]^{x+y-(1/2)}$$
$$\sim x^{x+y-(1/2)} e^y, \quad \text{as } x \to \infty$$

$$B(x,y) \sim \frac{e^{-x}x^{x-(1/2)}}{e^{-(x+y)}(x+y)^{x+y-(1/2)}}\Gamma(y) = \Gamma(y)x^{-y}$$
 as  $x \to \infty$ .

- $\therefore$  beta function has a power-law decay for large x.
- ∴ in-degree distribution has power-law tail

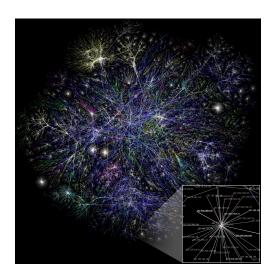
$$p_k \sim (k+a)^{-\beta} \Rightarrow p_k \sim k^{-\beta} \text{ for } k \gg a.$$

The power is  $\beta = 2 + \frac{a}{c} > 2$ .



# Example

The Internet is thought to grow following preferential attachment:



# Barabási-Albert model

A special case of the Price model is known as the Barabási-Albert (BA) model.

$$p_k \sim k^{-(2+\frac{a}{c})}$$

In this model,  $a = c \Rightarrow p_k \sim k^{-3}$ 

It is also has undirected edges rather than directed edges.

We can also study the in-degree distribution as a function of time (not just the infinite-time limit).

Let  $p_k(t, N)$  = mean fraction of nodes created at time t and which have in-degree k when the total number of nodes is N.

We are in event time, so first node is created at t = 1 and most recent nodes are created at t = N.

We'll use a master equation to study the dynamics of  $p_k(t, N)$ .

When we add a new node, the expected number of new edges acquired by existing nodes with in-degree k is:

$$Np_k(t,N) \times c \times \frac{a+k}{N(a+c)} = \frac{c(a+k)}{a+c} p_k(t,N),$$

which is independent of the creation time of those nodes. We obtain the following master equation (for  $k \ge 1$  and k = 0, respectively):

$$\begin{split} (N+1)p_k(t,N+1) &= Np_k(t,N) + p_{k-1}(t,N)\frac{c(k-1+a)}{c+a} - p_k(t,N)\frac{c(k+a)}{c+a},\\ (N+1)p_0(t,N+1) &= Np_k(t,N) + \delta_{t,N+1} - \frac{ca}{c+a}p_0(t,N), \end{split}$$

where 
$$\delta_{t,N+1} = \begin{cases} 1, t = N+1 \text{ i.e., a new in-degree 0 node is created,} \\ 0, \text{else.} \end{cases}$$

There is a problem as  $N \to \infty$  because fraction of nodes created at time t becomes 0 (only one node is created at a given time).

We rescale time using  $\tau := \frac{t}{N} \in [0, 1]$ .

: oldest nodes created at  $\tau=0$  and newest nodes created at  $\tau=1$ 

Since we've rescaled time, we also need to change from  $p_k(t, N)$  to a function  $\pi_k(\tau, N)$  such that  $\pi_k(\tau, N)d\tau$  is the fraction of nodes that have in-degree k that fall in the interval  $[\tau, \tau + d\tau]$ .

The number of nodes in that interval  $d\tau$  is  $Nd\tau \implies \pi_k d\tau = p_k \times Nd\tau$ , giving

$$\pi_k(\tau, N) = Np_k(t, N)$$

Note:  $\pi_k$  does not vanish as  $N \to \infty$ . Also  $\tau$  is not longer constant for a given node.

The master equation becomes:

$$\pi_k\left(\frac{N}{N+1}\tau,N+1\right)=\pi_k(\tau,N)+\frac{\pi_{k-1}(\tau,N)}{N}\frac{c(k-1+a)}{c+a}-\frac{\pi_k(\tau,N)}{N}\frac{c(k+a)}{c+a}.$$

Taking  $N \to \infty$  and abusing notation ( $\pi_k(\tau) = \tau_k(\tau, \infty)$ ) and defining  $\epsilon := \frac{1}{N}$  yields:

$$\frac{\pi_k(\tau) - \pi_k(\tau - \epsilon \tau)}{\epsilon} + \frac{c}{a+c} \left[ (a+k-1)\pi_{k-1}(\tau) - (a+k)\pi_k(\tau) \right] + \mathcal{O}(\epsilon^2) = 0.$$

As  $N \to \infty$ , so LHS  $\to \tau \frac{d\pi_k}{d\tau}$ 

$$\tau \frac{d\pi_k}{d\tau} + \frac{c}{a+c} \left[ (a+k-1)\pi_{k-1}(\tau) - (a+k)\pi_k(\tau) \right] + \mathcal{O}(\epsilon^2) = 0, \ k \ge 1, \quad (2)$$
$$\tau \frac{d\pi_0}{d\tau} - \frac{ca}{a+c}\pi_0(\tau) = 0, \ k = 0.$$

Notice Eqs. 2 is valid for  $\tau = [0, 1)$ . The boundary conditions are the final conditions at  $\tau = 1$ . These are  $\pi_0(1) = 1$  and  $\pi_k(1) = 0$  for  $k \ge 1$ .



Want to solve the k = 0 equation

$$\tau \frac{d\pi_0}{d\tau} - \frac{ca}{a+c}\pi_0(\tau) = 0, \ k = 0,$$

in closed form and ideally we use that to get solutions for  $k \geq 1$  iteratively:

$$\pi_0(\tau) = \tau^{ca/(a+c)}$$

$$\Rightarrow \pi_k(\tau) = \frac{1}{k!} [a(a+1)\cdots(a+k-1)] \tau^{ca/(a+c)} (1 - \tau^{c/(a+c)})^k$$

Again using  $B(x,y) = \frac{\Gamma(x)\Gamma(y)}{\Gamma(x+y)}$  and  $B(x,y) \sim x^{-y}$  as  $x \to \infty$  gives:

$$\pi_k(\tau) \sim k^{a-1} (1 - \tau^{c/(a+c)})^k \text{ as } k \to \infty.$$

$$\pi_k(\tau) \sim k^{a-1} (1 - \tau^{c/(a+c)})^k \text{ as } k \to \infty.$$

We see there is exponential decay (for  $|\tau| < 1$  implies the second term < 1) after some growth from the term  $k^{a-1}$ . Over all  $\tau$  we see power-law behavior as before.

Exponential decay is slower for small  $\tau$ , which means older nodes are more likely to have high in-degree than younger nodes.

Q: Is this true in general (e.g. Yahoo is older than Google in WWW)?

: if we wanted to use this as a model of structures of the WWW, we're clearly missing ingredients (an important one is node fitness)!

# Edge removal

Consider an example in which undirected edges are added preferentially as in the BA model (number of new connections for each new node is exactly  $c = a \in \mathbb{Z}_+$ ).

Then the degree of each node in BA mechanism is exactly c (not the average), so degree of each node is  $k_i = k_i + c$  where  $k_i$  = in-degree from the de Solla Price model (note change of notation).

Also, edges are *removed* uniformly at random during update steps. Probability that one of the removed edge's stubs is attached to node i is then proportional to  $k_i$ .

$$prob(node i looses an edge) = \frac{2k_i}{\sum_i k_i}$$

(factor of 2 because each edge has 2 stubs)



# Edge removal

Suppose at each step, degree c nodes are added to the network and a mean of v edges are removed for each node that is added (assume v < c).

On average, the network has N(c-v) edges (for N nodes).

One can derive the master equation (for  $k \geq 1$ ):

$$(N+1)p_k(N+1) = Np_k(N) + \delta_{kc} + p_{k-1}(N) \frac{c(k-1)}{2(c-v)} + p_{k+1}(N) \frac{2v(k+1)}{2(c-v)} - p_k(N) \frac{(c+2v)k}{2(c-v)},$$

For k = 0, the term proportional to (k - 1) isn't there (because no nodes with degree -1), all other are present.

Take  $N \to \infty$  to study asymptotic properties of degree distribution as usual.



# Nodes with varying quality

It is desirable to develop models in which different nodes have inherently different fitnesses (more than structural features).

e.g., Bianconi-Barabási (BB) model

As in BA model, add nodes one at a time via undirected edges to c existing nodes.

Node i has scalar fitness  $\eta_i \in \mathbb{R}$  assigned when node is created (drawn from a continuous probability distribution  $p(\eta)$ .

The c new stubs each attach to an existing node with probability  $\propto a(k, \eta)$  (e.g.,  $a(k, \eta) = \eta k$ ).

# Network rewiring in bipartite network

So far, we've considered network growth (structural and fitness properties), now rewiring.

e.g., Evans-Plato model

This is a bipartite network with U people and M artifacts.

At time t, each of U people is connected via an undirected, unweighted edge to one of M artifacts (e.g., renting movies).

One can study this kind of setup with a rewiring rule for changing edges.

# Network optimization

Consider network structures that arise through some optimization process. One constructs some sort of *cost function* and see what network structures minimize it.

e.g., Gastner-Newman model for transportation networks: Let  $t_{ij}$  = time to get from city i to city j.

$$= b_1 + b_2 r_{ij}$$

where  $b_1$  is total time spent in airports,  $b_2$  is total time traveled (proportional to distance), and  $r_{ij}$  is geographic distance between i and j.

Location of nodes can come from probability distribution (e.g., placed via some random process in square  $[0,1] \times [0,1]$ ) or from empirical data (e.g., directly from a map).

# Network optimization

Cost function is  $E(m, L) = \lambda m + (1 - \lambda)L$ ,  $\lambda \in [0, 1]$ , where L is the mean geodesic distance between node pairs given by time travel  $t_{ij}$  and m is the number of edges.

# Results (briefly):

- small  $b_1$ , large  $b_2$ : networks "close" to planar (like a road network),
- large  $b_1$ , small  $b_2$ : hub-and-spoke networks (like airline network) with many low-degree nodes and few high-degree nodes.

Q: How does the Gastner-Newman model behave for different values of  $\lambda$ ?

