

Machine Prog (Assembly Code) ——Basics & Control section

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2023.9.27

Preface

许珈铭

C/C++ 编译

- Pre-processor
- Compiler
- Assembler
- Linker&loader

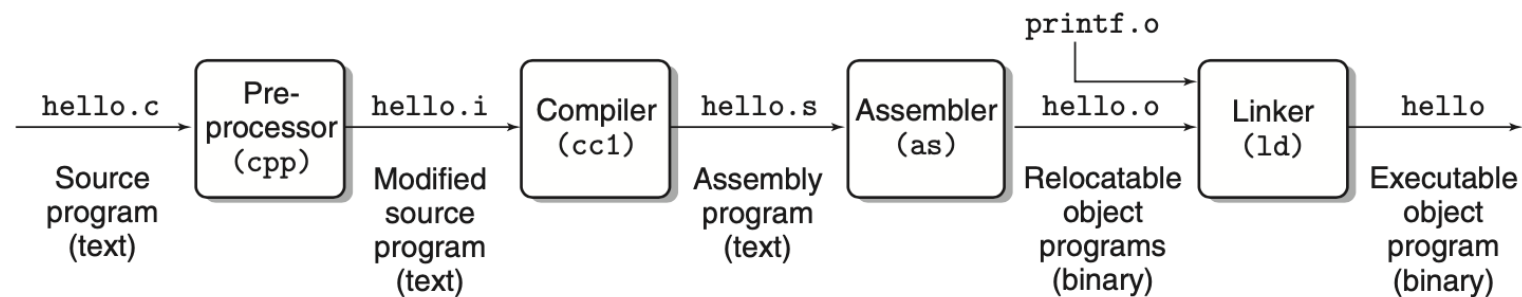


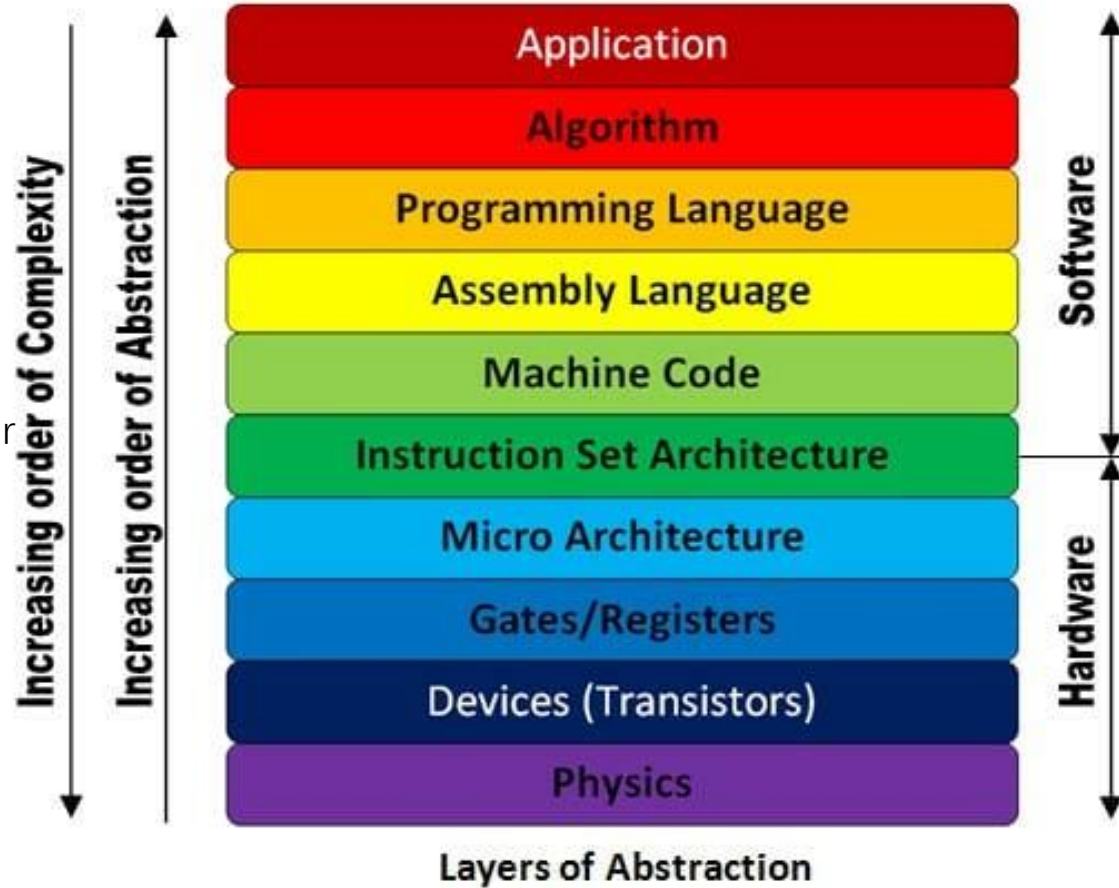
Figure 1.3 The compilation system.

Recall

- C/C++
(.c/.cpp, .h/.hpp)

GCC
↓ Compiler
• Assembly language (.s)
↓ Assembler
• Machine Code (.o, .so, .a)
↓ Linker

CLANG (LLVM)
↓ Compiler
• LLVM IR
↓ LLVM Optimizer
• LLVM Backend
↓ Assembler
↓ Linker



- Java

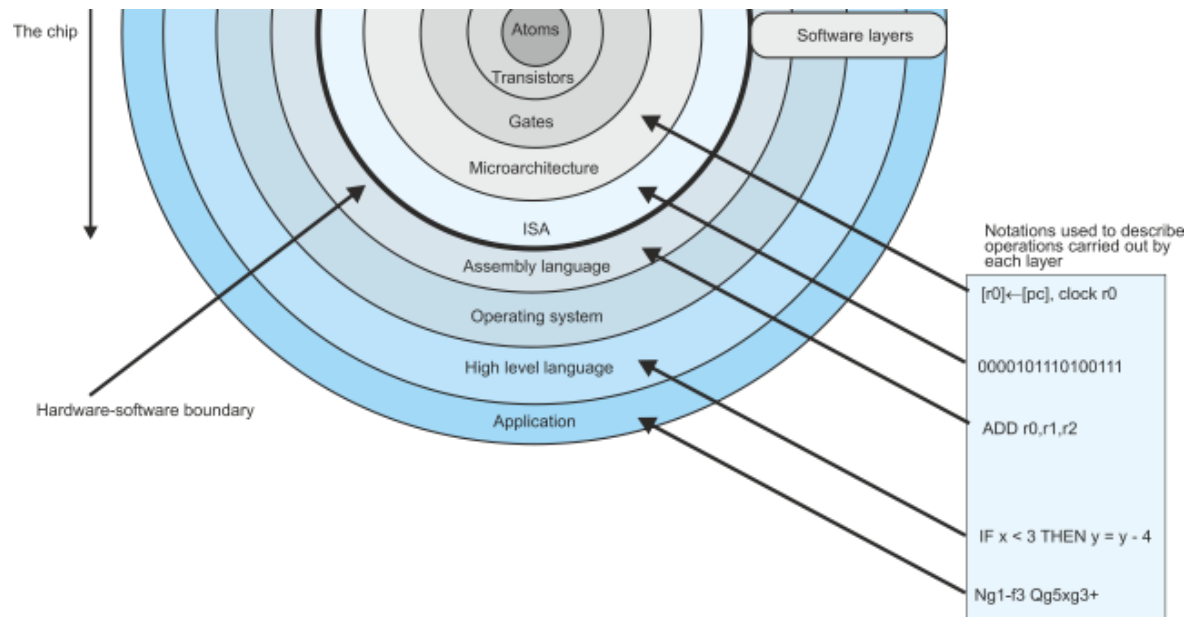
JVM
↓ Compiler
• Java Bytecode
(Java Virtual Machine)
↓ Assembler
↓ Linker
• Machine Code

JIT
↓ Compiler
• Java Bytecode
↓ Just In Time compiler
• Machine code
• (Java Virtual Machine)
↓ Linker
• Machine Code

Two abstractions: ISA

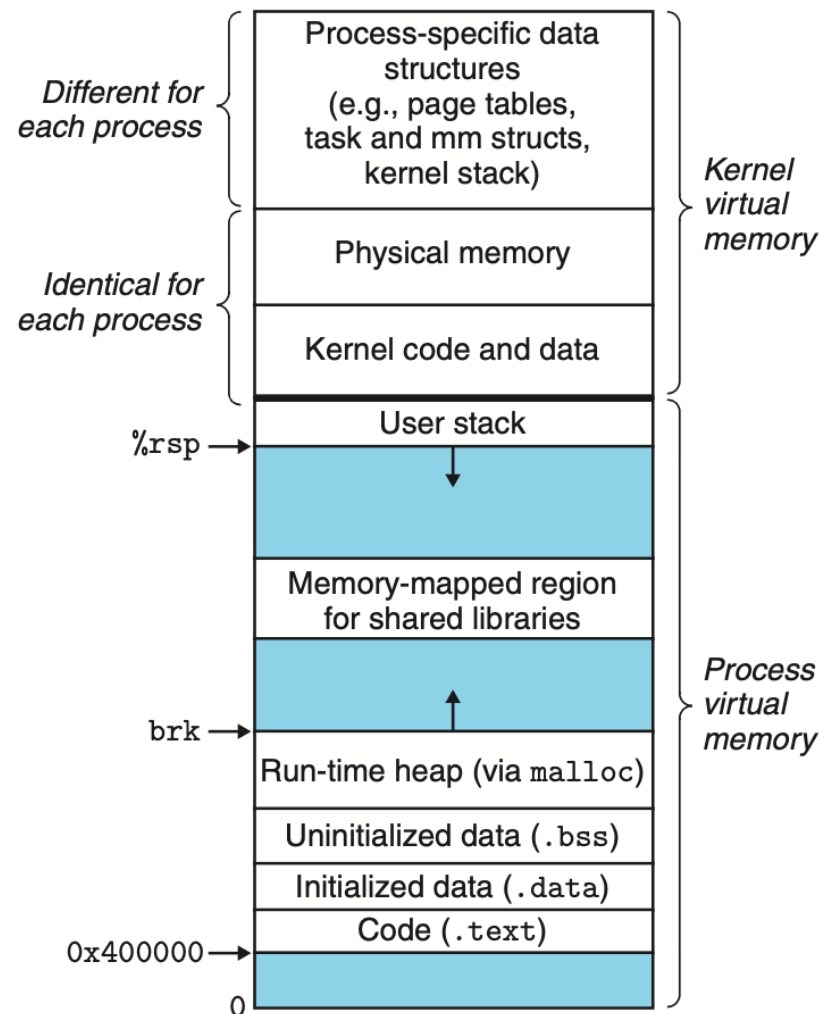
- 定义处理器状态、指令格式和指令对状态的影响
- The hardware/software **interface**
- Examples
 - IA32, **AMD64 (x86-64)**, ARM64, RISC-V, etc.

• ISA \neq 汇编语言



Two abstractions: Virtual Address

- 每个 **进程** 享有同样的虚拟地址空间
- 操作系统的四大抽象之一，随着ICS的学习会逐渐补全拼图



工具链


(1) gcc [-O] -S <file> [-o]

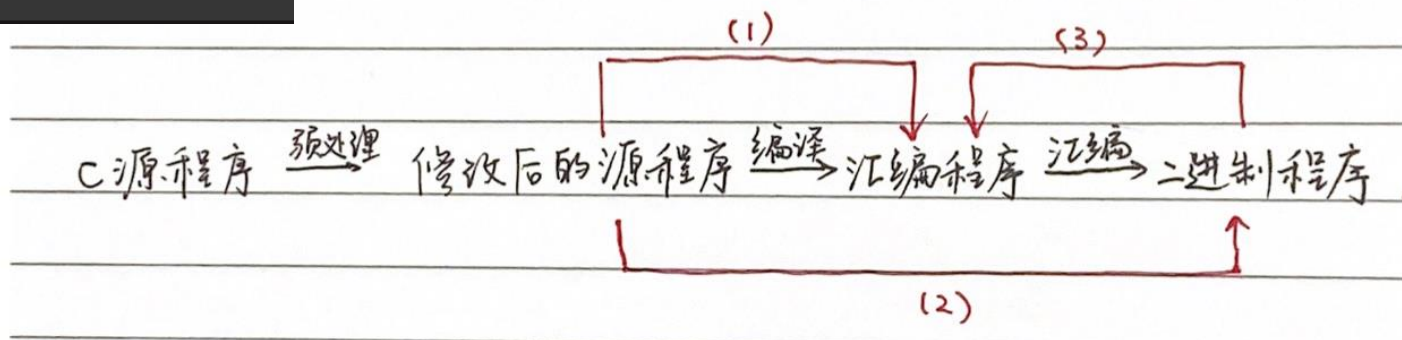
(2) gcc [-O] -c <file> [-o]

(3) objdump -d <file> (> ...)

• gdb and bomblab

(c->s) Compiler Explorer  COMPILER EXPLORER

(o->c) ByteNinja  BINARYNINJA



编译器优化



助教 李翰禹

再具体解释一下优化等级的问题

- O0是几乎所有优化都不做

- O或者-O1是做一部分优化

- Og是-O1优化去除会影响调试器正常工作的优化

- O2是-O1不考虑时间和空间的权衡的情况下做更多优化

- O3是-O2基础上做更多优化，比如一些向量化和循环展开的优化

- Os是-O2优化去除会增加目标代码长度的优化

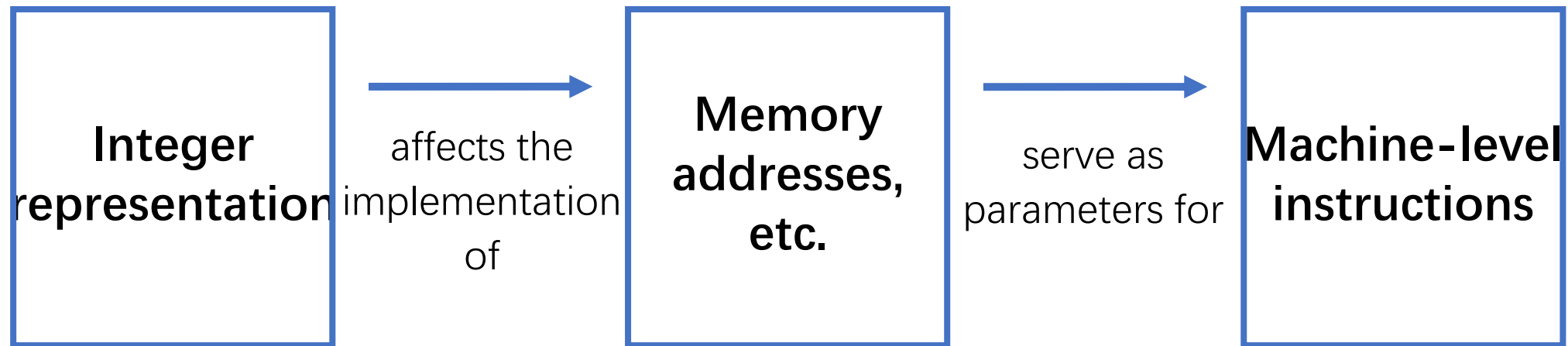
- Ofast是无视严格标准的优化

Basics (CS:APP Ch. 3.1 - 3.5)

贾博暄

Relationships between topics

- (courtesy of TA)



Registers

- Correspondence
- Memory reference in parens: ()
 - **%rax** -> register
 - **(%rax)** -> memory
 - Compare: pointers in C

Note on
naming

b, w, l, q

C declaration	Intel data type	Assembly-code suffix	Size (bytes)
char	Byte	b	1
short	Word	w	2
int	Double word	l	4
long	Quad word	q	8
char *	Quad word	q	8
float	Single precision	s	4
double	Double precision	l	8

Figure 3.1 Sizes of C data types in x86-64. With a 64-bit machine, pointers are 8 bytes long.

63	31	15	7	0	
%rax	%eax	%ax	%al		Return value
%rbx	%ebx	%bx	%bl		Callee saved
%rcx	%ecx	%cx	%cl		4th argument
%rdx	%edx	%dx	%dl		3rd argument
%rsi	%esi	%si	%sil		2nd argument
%rdi	%edi	%di	%dil		1st argument
%rbp	%ebp	%bp	%bpl		Callee saved
%rsp	%esp	%sp	%spl		Stack pointer
%r8	%r8d	%r8w	%r8b		5th argument
%r9	%r9d	%r9w	%r9b		6th argument
%r10	%r10d	%r10w	%r10b		Caller saved
%r11	%r11d	%r11w	%r11b		Caller saved
%r12	%r12d	%r12w	%r12b		Callee saved
%r13	%r13d	%r13w	%r13b		Callee saved
%r14	%r14d	%r14w	%r14b		Callee saved
%r15	%r15d	%r15w	%r15b		Callee saved

Conventions...

- Instructions that generate 4-byte quantities for a register set the upper 4 bytes of the register to zero

1	<code>movabsq \$0x0011223344556677, %rax</code>	<code>%rax = 0011223344556677</code>
2	<code>movb \$-1, %al</code>	<code>%rax = 00112233445566FF</code>
3	<code>movw \$-1, %ax</code>	<code>%rax = 001122334455FFFF</code>
4	<code>movl \$-1, %eax</code>	<code>%rax = 00000000FFFFFFFF</code>
5	<code>movq \$-1, %rax</code>	<code>%rax = FFFFFFFFFFFFFFFF</code>

- Consequences
 - MOV** instructions get a little weird (e.g. there is no **movz1q**)
 - You might see (for example):
 - movzbl** when your intuition says **movzbq**
Textbook: Problem 3.4
 - Writing to **%eax** on one line and quoting **%rax** on the next

Conventions...

- In a memory reference, the **scaling factor** must be either 1, 2, 4, or 8

$Imm(rb, r_i, s)$

9 (%rax, %rdx, 4)

- Why?
 - Useful when referencing array and structure elements

1、某C语言程序中对数组变量a的声明为“int a[10][10];”，有如下一段代码：

```
for (i=0; i<10; i++)  
    for (j=0; j<10; j++)  
        sum+= a[i][j];
```

假设执行到“sum+= a[i][j];”时，sum的值在%rax中，a[i][0]所在的地址在%rdx中，j在%rsi中，则“sum+= a[i][j];”所对应的指令是（ ）。

- A. addl 0 (%rdx, %rsi, 4), %rax
- B. addl 0 (%rsi, %rdx, 4) , %rax
- C. addl 0 (%rdx, %rsi, 2) , %rax
- D. addl 0 (%rsi, %rdx, 2) , %rax

Conventions...

- x86-64 不允许将除 64 位寄存器以外的寄存器作为寻址模式基地址

1. 判断下列 x86-64 ATT 操作数格式是否合法。

(7) () (%ecx)

- Textbook: Problem 3.3

Here is the code with explanations of the errors:

```
movb $0xF, (%ebx)      Cannot use %ebx as address register
```

Conventions...

- Binary instructions' operands are, in order, _Source_ and _Destination_
 - Intuitive (eg. **MOV**)
 - Counter-intuitive (eg. **SUB**, **CMP**)
- the operand pairs CANNOT be of what type? _Any -> Imm, Mem -> Mem_
 - this applies to **MOV** and Arithmetic

Instruction Classes – Move

- **MOV**

- Operand types: $\$Imm, Reg, Mem$
- Operand pairs: Some are forbidden

- Memory reference

- Full form:

Offset (Base, Index, scale factor)

$$Imm + R[rb] + R[ri] \cdot s$$

- Selected omission yields the other formats
- s limited to 1, 2, 4, 8
- No \$ before Imm here

From\To	Imm	Reg	Mem
Imm		✓	✓
Reg		✓	✓
Mem		✓	

Instruction Classes – Move

- **MOV:** `movl` vs `movq` vs `movabsq`

- `movl`

- Moves double word
 - Higher bits automatically set to 0

- `movq`

- Takes 32-bit value, sign-extends to 64-bit (i.e. quad word)
 - **THEN** moves to destination

- `movabsq`

- Takes 64-bit immediate value and moves to destination
 - **Only allows Imm -> Reg**

Instruction Classes – Move

- “Corner cases”

- **MOVZ**

- there is no `movzlw`

- **MOVS**

- **cltq**

- effectively a compact encoding of

`movslq %eax, %rax`

MOVZ



z: zero-extend
s: sign-extend

MOVS



Instruction Classes – Arithmetic

- **leaq**

- (Important!) **leaq** vs **mov**

- Same format

- **leaq** DOES NOT reference memory; purely numerical

B. `leal (%ecx,%ebx,4), %eax`

C. `movl (%ecx,%ebx,4), %eax`

- Unary: **INC, DEC, NEG, NOT**

- Somewhat self-evident

- Binary: **ADD, SUB, IMUL, XOR, OR, AND**

- 2nd operand (*D*) = Destination

- Mnemonic: **Add** *S* to *D*
Subtract *S* from *D*

`ADD S, D` $D \leftarrow D + S$

`SUB S, D` $D \leftarrow D - S$

Instruction Classes – Shifts

- Shifts
 - **SAL(SHL), SAR, SHR**
 - **sh**ifts (logical) vs **s**hifts (**a**rithmetic)

sh: logical shift

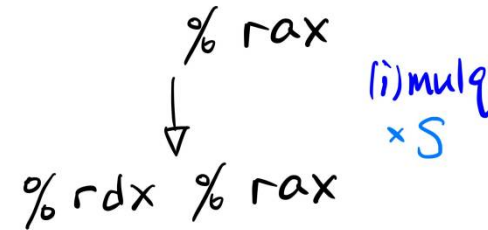
sa: arithmetic shift

- 1st operand (k) = shift amount
 - **must be \$lmm or %c1**
 - actual shift amount (also) depends on type of value to be shifted

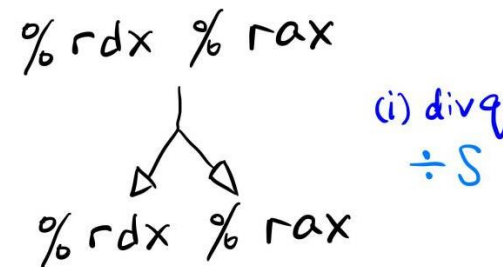
SAL	k, D	$D \leftarrow D \ll k$
SHL	k, D	$D \leftarrow D \ll k$
SAR	k, D	$D \leftarrow D \gg_A k$
SHR	k, D	$D \leftarrow D \gg_L k$

Instruction Classes - 128-bit

- **Single operand**
- **`imulq`, `mulq`**
 - S is multiplicand
 - **`imulq`** S vs **`IMUL`** S, D
- **`idivq`, `divq`**
 - S is divisor
 - Quotient at `%rax`, remainder at `%rdx`
- **`cqto`**
 - Sign extend, presets `%rdx`
 - Compare: **`cltq`**



i: integer



`cltq`:

convert **l** to **q**

`cqto`:

convert **q** to **o**

On to the next section!

Control (CS:APP Ch 3.6)

倪嘉怡

Contents

- 条件码, 跳转指令
- 逆向工程和goto代码
- **if-else, do-while, while, for, switch, ...**
- 条件传送

条件码

- CF Carry Flag：进位；无符号溢出。1进位0不进
- ZF Zero Flag：零。1零0非零
- SF Sign Flag：符号。1负0正
- OF Overflow Flag：溢出；正/负有符号溢出。1溢出0不溢出

- 逻辑运算：进位 CF、溢出 OF 置 0
- 移位运算：溢出 OF 置 0；进位 CF 设置为最后一个被移出的位
- inc, dec：不改变进位标志
- lea：不改变条件码

- cmp = sub -
- test = and &

- <https://www.felixcloutier.com/x86/>
- x86 and amd64 instruction reference

指令	效果	描述
leal S, D	$D \leftarrow \&S$	加载有效地址
INC D	$D \leftarrow D + 1$	加1
DEC D	$D \leftarrow D - 1$	减1
NEG D	$D \leftarrow -D$	取负
NOT D	$D \leftarrow \sim D$	取补
ADD S, D	$D \leftarrow D + S$	加
SUB S, D	$D \leftarrow D - S$	减
IMUL S, D	$D \leftarrow D * S$	乘
XOR S, D	$D \leftarrow D \wedge S$	异或
OR S, D	$D \leftarrow D \vee S$	或
AND S, D	$D \leftarrow D \& S$	与
SAL k, D	$D \leftarrow D \ll k$	左移
SHL k, D	$D \leftarrow D \ll k$	左移（等同于SAL）
SAR k, D	$D \leftarrow D \gg_A k$	算术右移
SHR k, D	$D \leftarrow D \gg_L k$	逻辑右移

	CF	ZF	SF	OF	
LEA	✗	✗	✗	✗	INC, DEC
INC	✗	✓	✓	✓	
DEC	✗	✓	✓	✓	
NEG	✓	✓	✓	✓	算术运算
NOT	✗	✗	✗	✗	
ADD	✓	✓	✓	✓	
SUB	✓	✓	✓	✓	
IMUL	✓	✓	✓	✓	
XOR	0	✓	✓	0	逻辑运算
OR	0	✓	✓	0	
AND	0	✓	✓	0	
SAL	✓	✓	✓	0	移位运算
SHL	✓	✓	✓	0	
SAR	✓	✓	✓	0	
SHR	✓	✓	✓	0	

跳转指令

Jump

if Equal

if Not Equal

if Sign

if Not Sign

if Greater

if Not Greater

if Less

if Not Less

if Above

if Not Above

if Below

if Not Below

指令	同义名	跳转条件	描述
jmp <i>Label</i>		1	直接跳转
jmp <i>*Operand</i>		1	间接跳转
jje <i>Label</i>	jz	ZF	相等/零
jne <i>Label</i>	jnz	~ZF	不相等/非零
js <i>Label</i>		SF	负数
jns <i>Label</i>		~SF	非负数
jg <i>Label</i>	jnle	~(SF ^ OF) & ~ZF	大于 (有符号>)
jge <i>Label</i>	jnl	~(SF ^ OF)	大于或等于 (有符号>=)
jl <i>Label</i>	jnge	SF ^ OF	小于 (有符号<)
jle <i>Label</i>	jng	(SF ^ OF) ZF	小于或等于 (有符号<=)
ja <i>Label</i>	jnbe	~CF & ~ZF	超过 (无符号>)
jae <i>Label</i>	jnb	~CF	超过或相等 (无符号>=)
jb <i>Label</i>	jnae	CF	低于 (无符号<)
jbe <i>Label</i>	jna	CF ZF	低于或相等 (无符号<=)

“*” 相当于C语言指针，汇编的括号

%rax 寄存器
(%rax) 内存
寄存器+括号=内存

逆向工程和goto

- goto代码：描述汇编代码程序控制流的C程序
- 逆向工程：汇编代码->C语言代码

```
x at %ebp+8, y at %ebp+12
1    movl    8(%ebp), %edx    Get x
2    movl    12(%ebp), %eax   Get y
3    cmpl    %eax, %edx       Compare x:y
4    jge     .L2              if >= goto x_ge_y
5    subl    %edx, %eax       Compute result = y-x
6    jmp     .L3              Goto done
7  .L2:                                x_ge_y:
8    subl    %eax, %edx       Compute result = x-y
9    movl    %edx, %eax       Set result as return value
10   .L3:                                done: Begin completion code
```

c). 产生的汇编代码

```
1  int absdiff(int x, int y) {
2      if (x < y)
3          return y - x;
4      else
5          return x - y;
6  }
```

a) 原始的 C 语言代码

goto: 正向视角

```
1  int absdiff(int x, int y) {  
2      if (x < y)  
3          return y - x;  
4      else  
5          return x - y;  
6  }
```

a) 原始的 C 语言代码

```
1  int gotodiff(int x, int y) {  
2      int result;  
3      if (x >= y)  
4          goto x_ge_y;  
5      result = y - x;  
6      goto done;  
7  x_ge_y:  
8      result = x - y;  
9      done:  
10     return result;  
11 }
```

b) 与之等价的 goto 版本

Basic: mov, add, sub...

专用于设置条件码: cmp, test...

跳转: jmp, je, jg...

如何把else转化为汇编?

```
      x at %ebp+8, y at %ebp+12  
1      movl    8(%ebp), %edx    Get x  
2      movl    12(%ebp), %eax   Get y  
3      cmpl    %eax, %edx      Compare x:y  
4      jge     .L2              if >= goto x_ge_y  
5      subl    %edx, %eax       Compute result = y-x  
6      jmp     .L3              Goto done  
7  .L2:  
8      subl    %eax, %edx       Compute result = x-y  
9      movl    %edx, %eax       Set result as return value  
10     .L3:  
      done: Begin completion code
```

c) 产生的汇编代码

goto: 逆向视角

```
x at %ebp+8, y at %ebp+12
1    movl    8(%ebp), %edx
2    movl    12(%ebp), %eax
3    cmpl    %eax, %edx
4    jge     .L2
5    subl    %edx, %eax
6    jmp     .L3
7    .L2:
8    subl    %eax, %edx
9    movl    %edx, %eax
10   .L3:
```

c) 产生的汇编

```
1  int gotodiff(int x, int y) {
2      int result;
3      if (x >= y)
4          goto x_ge_y;
5      result = y - x;
6      goto done;
7      x_ge_y:
8      result = x - y;
9      done:
10     return result;
11 }
```

b) 与之等价的 goto 版本

```
1  int absdiff(int x, int y) {
2      if (x < y)
3          return y - x;
4      else
5          return x - y;
6  }
```

a) 原始的 C 语言代码

循环：do-while

```
do
    body-statement
while (test-expr);
```

```
loop:
    body-statement
    t = test-expr;
    if (t)
        goto loop;
```

```
1  int fact_do(int n)
2  {
3      int result = 1;
4      do {
5          result *= n;
6          n = n-1;
7      } while (n > 1);
8      return result;
9  }
```

a) C 代码

```
Argument: n at %ebp+8
Registers: n in %edx, result in %eax
1      movl    8(%ebp), %edx    Get n
2      movl    $1, %eax        Set result = 1
3      .L2:                    loop:
4          imull %edx, %eax      Compute result *= n
5          subl    $1, %edx      Decrement n
6          cmpl    $1, %edx      Compare n:1
7          jg      .L2           If >, goto loop
Return result
```

c) 对应的汇编代码

循环: while

```
while (test-expr)
    body-statement
```

```
t = test-expr;
if (!t)
    goto done;
loop:
    body-statement
    t = test-expr;
    if (t)
        goto loop;
done:
```

```
1  int fact_while(int n)
2  {
3      int result = 1;
4      while (n > 1) {
5          result *= n;
6          n = n-1;
7      }
8      return result;
9  }
```

a) C 代码

```
1  int fact_while_goto(int n)
2  {
3      int result = 1;
4      if (n <= 1)
5          goto done;
6      loop:
7          result *= n;
8          n = n-1;
9      if (n > 1)
10         goto loop;
11     done:
12         return result;
13 }
```

b) 等价的 goto 版本

Argument: n at %ebp+8
Registers: n in %edx, result in %eax

```
1  movl    8(%ebp), %edx    Get n
2  movl    $1, %eax        Set result = 1
3  cmpl    $1, %edx        Compare n:1
4  jle     .L7             If <=, goto done
5  .L10:
6  imull   %edx, %eax       Compute result *= n
7  subl    $1, %edx        Decrement n
8  cmpl    $1, %edx        Compare n:1
9  jg      .L10            If >, goto loop
10 .L7:
    Return result
done:
```

c) 对应的汇编代码

循环: for

```
for (init-expr; test-expr; update-expr)
    body-statement
```

```
1  int fact_for(int n)
2  {
3      int i;
4      int result = 1;
5      for (i = 2; i <= n; i++)
6          result *= i;
7      return result;
8  }
```

```
init-expr;
while (test-expr) {
    body-statement
    update-expr;
}
```

```
init-expr;
t = test-expr;
if (!t)
    goto done;
loop:
    body-statement
    update-expr;
    t = test-expr;
    if (t)
        goto loop;
done:
```

```
1  int fact_for_goto(int n)
2  {
3      int i = 2;
4      int result = 1;
5      if (!(i <= n))
6          goto done;
7      loop:
8          result *= i;
9          i++;
10         if (i <= n)
11             goto loop;
12     done:
13         return result;
14 }
```

Argument: n at %ebp+8

Registers: n in %ecx, i in %edx, result in %eax

```
1  movl    8(%ebp), %ecx    Get n
2  movl    $2, %edx        Set i to 2      (init)
3  movl    $1, %eax        Set result to 1
4  cmpl    $1, %ecx        Compare n:1      (!test)
5  jle     .L14             If <=, goto done
6  .L17:
7  imull    %edx, %eax      Compute result *= i (body)
8  addl    $1, %edx        Increment i      (update)
9  cmpl    %edx, %ecx       Compare n:i      (test)
10 jge     .L17             If >=, goto loop
11 .L14:
done:
```

switch

• 跳转表内的数组引用

```
1  int switch_eg(int x, int n) {
2      int result = x;
3
4      switch (n) {
5
6          case 100:
7              result += 13;
8              break;
9
10         case 102:
11             result += 10;
12             /* Fall through */
13
14         case 103:
15             result += 11;
16             break;
17
18         case 104:
19         case 106:
20             result += result;
21             break;
22
23         default:
24             result = 0;
25     }
26
27     return result;
28 }
```

a) switch 语句

```
1  int switch_eg_impl(int x, int n) {
2      /* Table of code pointers */
3      static void *jt[7] = {
4          &&loc_A, &&loc_def, &&loc_B,
5          &&loc_C, &&loc_D, &&loc_def,
6          &&loc_D
7      };
8
9      unsigned index = n - 100;
10     int result;
11
12     if (index > 6)
13         goto loc_def;
14
15     /* Multiway branch */
16     goto *jt[index];
17
18     loc_def: /* Default case*/
19     result = 0;
20     goto done;
21
22     loc_C: /* Case 103 */
23     result = x;
24     goto rest;
25
26     loc_A: /* Case 100 */
27     result = x * 13;
28     goto done;
29
30     loc_B: /* Case 102 */
31     result = x + 10;
32     /* Fall through */
33
34     rest: /* Finish case 103 */
35     result += 11;
36     goto done;
37
38     loc_D: /* Cases 104, 106 */
39     result = x * x;
40     /* Fall through */
41
42     done:
43     return result;
44 }
```

b) 翻译到扩展的 C 语言

```
x at %ebp+8, n at %ebp+12
1  movl    8(%ebp), %edx          Get x
2  movl    12(%ebp), %eax        Get n
3                                     Set up jump table access
4  subl    $100, %eax            Compute index = n-100
5  cmpl    $6, %eax              Compare index:6
6  ja      .L2                   If >, goto loc_def
7  jmp     *.L7(,%eax,4)          Goto *jt[index]
8                                     Default case
9  .L2:                          loc_def:
10     movl    $0, %eax            result = 0;
11     jmp     .L8                 Goto done
12                                     Case 103
13     .L5:                          loc_C:
14     movl    %edx, %eax          result = x;
15     jmp     .L9                 Goto rest
16                                     Case 100
17     .L3:                          loc_A:
18     leal    (%edx,%edx,2), %eax  result = x*3;
19     leal    (%edx,%eax,4), %eax  result = x+4*result
20     jmp     .L8                 Goto done
21                                     Case 102
22     .L4:                          loc_B:
23     leal    10(%edx), %eax       result = x+10
24                                     Fall through
25     .L9:                          rest:
26     addl    $11, %eax            result += 11;
27     jmp     .L8                 Goto done
28                                     Cases 104, 106
29     .L6:                          loc_D
30     movl    %edx, %eax          result = x
31     imull   %edx, %eax          result *= x
32                                     Fall through
33     .L8:                          done:
34                                     Return result
```

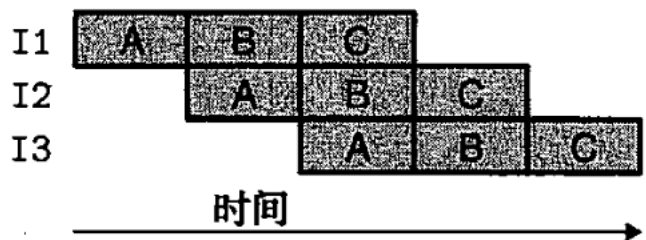
图 3-19 图 3-18 中 switch 语句示例的汇编代码

条件传送

- 条件**控制**转移->条件**数据**传送：可能提高效率
- 计算流水线（p264）：取指—译码—执行—访存—写回

```
1  int absdiff(int x, int y) {  
2      return x < y ? y-x : x-y;  
3  }
```

a) 原始的 C 语言代码



```
1  int cmovdiff(int x, int y) {  
2      int tval = y-x;  
3      int rval = x-y;  
4      int test = x < y;  
5      /* Line below requires  
6          single instruction: */  
7      if (test) rval = tval;  
8      return rval;  
9  }
```

b) 使用条件赋值的实现

Bad Cases for Conditional Move

Expensive Computations

```
val = Test(x) ? Hard1(x) : Hard2(x);
```

Bad Performance

- Both values get computed
- Only makes sense when computations are very simple

Risky Computations

```
val = p ? *p : 0;
```

Unsafe

- Both values get computed
- May have undesirable effects

当test-expr p为False, *p间接引用空指针

Computations with side effects

```
val = x > 0 ? x*=7 : x+=3;
```

Illegal

- Both values get computed
- Must be side-effect free

两个分支同时改变全局变量

Practice

王善上

The End