



Compilers

LL(1) Parsing Tables

- Construct a parsing table T for CFG G
- For each production $A \rightarrow \alpha$ in G do:
 - For each terminal $t \in \text{First}(\alpha)$ do
 - $T[A, t] = \alpha$
 - If $\epsilon \in \text{First}(\alpha)$, for each $t \in \text{Follow}(A)$ do
 - $T[A, t] = \alpha$
 - If $\epsilon \in \text{First}(\alpha)$ and $\$ \in \text{Follow}(A)$ do
 - $T[A, \$] = \alpha$

LL(1) Parsing Tables

$E \rightarrow T X$

$T \rightarrow (E) \mid \text{int } Y$

$X \rightarrow + E \mid \varepsilon$

$Y \rightarrow * T \mid \varepsilon$

LL(1) Parsing Tables

- If any entry is multiply defined then G is not LL(1)
- Most programming language CFGs are not LL(1)