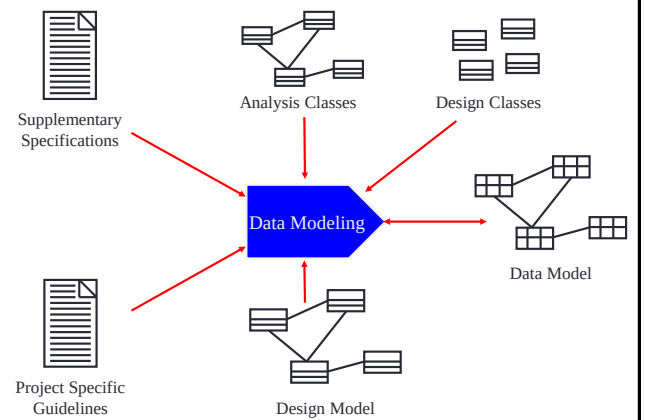
ITSS SOFTWARE DEVELOPMENT

7. DATA MODELING

Some slides extracted from IBM coursewares

1

Data Modeling Overview



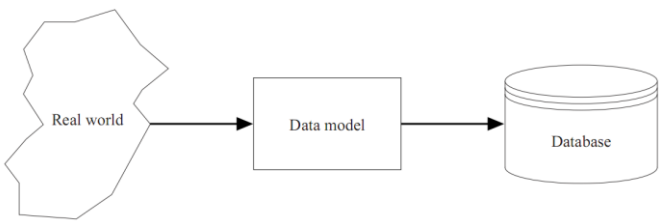
2

Content

3

1. Data Models

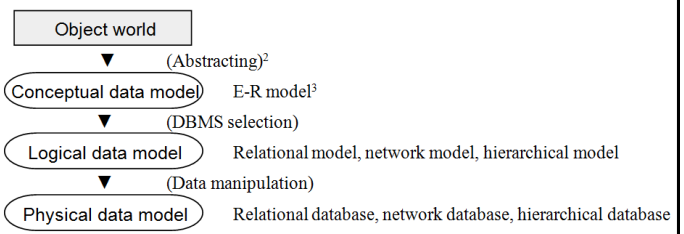
- Data modeling:
  - Abstracting and organizing the structure of real-world information, which is the object to be made into a database, and then expressing it



4

1. Data Models (2)

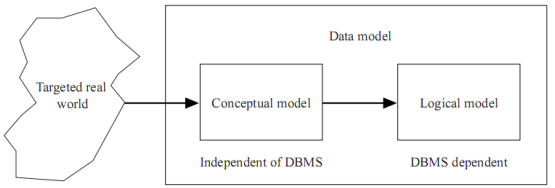
- 3 types of data models



5

1.1. Conceptual data model

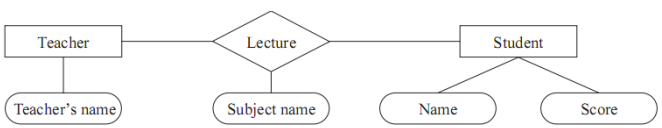
- Naturual expressions without constraints imposed by DBMS
- E-R model
  - Expressed by E-R diagram



6

E-R Diagram

- Three elements
  - Entities
  - Relationships
  - Attributes



7

1.2. Logical Data Model

- 3 types
  - relational model,
  - network model,
  - and hierarchical model

8

9

1.3. Physical Data Model

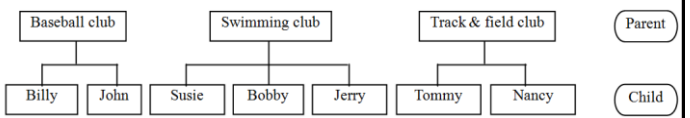
- Logical data models, when they are implemented, become physical data models:
  - relational databases,
  - network databases,
  - or hierarchical databases

9

10

1.3.1. Hierarchical Database (Tree-Structure Database)

- Divides records into parents and children and shows the relationship with a hierarchical structure
- 1-to-many (1:n) correspondences between parent records and child records

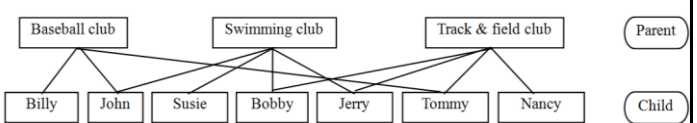


10

11

1.3.2. Network Database

- Parent records and child records do not have 1-to-n (1:n) correspondences; rather, they are in many-to-many (m:n) correspondence
- Sometimes called CODASYL database

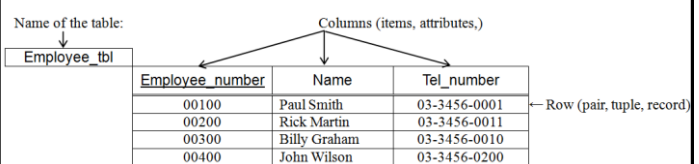


11

12

1.3.3. Rational database

- Data is expressed in a two-dimensional table.
  - Each row of the table corresponds to a record, and each column is an item of the records.
  - The underlined columns indicate the primary key



12

NOSQL DATABASES

Overview, Models, Concepts, Examples

13

What is NoSQL Database?

- NoSQL (cloud) databases
  - Use document-based model (non-relational)
  - Schema-free document storage
    - Still support indexing and querying
    - Still support CRUD operations (create, read, update, delete)
    - Still supports concurrency and transactions
  - Highly optimized for append / retrieve
  - Great performance and scalability
  - NoSQL == “No SQL” or “Not Only SQL”?

14

Relational vs. NoSQL Databases

- Relational databases
  - Data stored as table rows
  - Relationships between related rows
  - Single entity spans multiple tables
  - RDBMS systems are very mature, rock solid
- NoSQL databases
  - Data stored as documents
  - Single entity (document) is a single record
  - Documents do not have a fixed structure

15

Relational vs. NoSQL Models

Relational Model

Name	Svetlin Nakov
Gender	male
Phone	+359333777555
Email	nakov@abv.bg
Site	www.nakov.com

* 1

Street	Al. Malinov 31
Post Code	1729

* 1

Town	Sofia

* 1

Country	Bulgaria

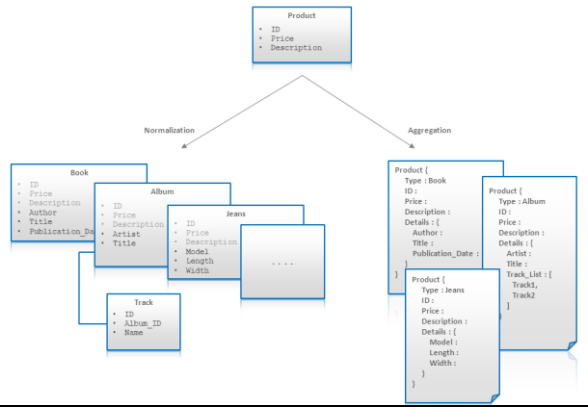
Document Model

Name: Svetlin Nakov
Gender: male
Phone: 359333777555
Address:
- Street: Al. Malinov 31
- Post Code: 1729
- Town: Sofia
- Country: Bulgaria

Email: nakov@abv.bg
Site: www.nakov.com

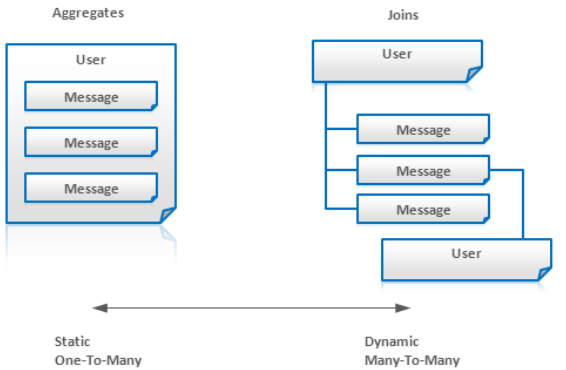
16

Normalization vs. Aggregation



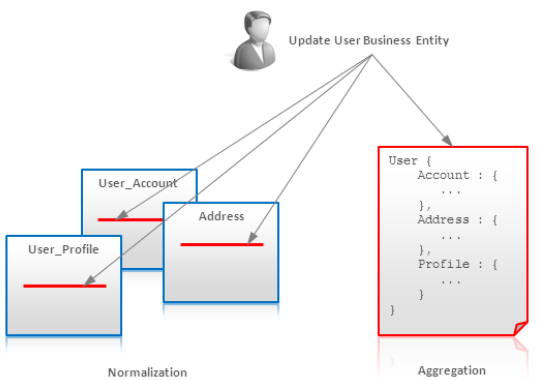
17

Aggregates vs. Joins



18

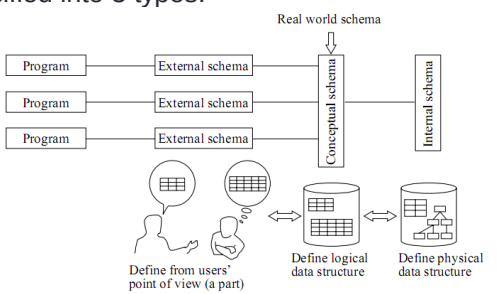
Atomic Aggregates



19

1.3.4. Three-layer schema

- A schema is a description of the framework of a database
- Classified into 3 types:



20

Content

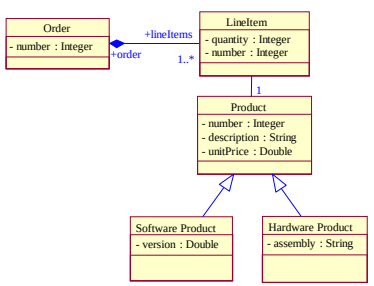
1. Data models
2. Object model and Rational Data Model
3. Mapping class diagram to E-R diagram
4. Normalization

2.1. Relational Databases and OO

- RDBMS and Object Orientation are not entirely compatible
  - RDBMS
    - Focus is on data
    - Better suited for ad-hoc relationships and reporting application
    - Expose data (column values)
  - Object Oriented system
    - Focus is on behavior
    - Better suited to handle state-specific behavior where data is secondary
    - Hide data (encapsulation)

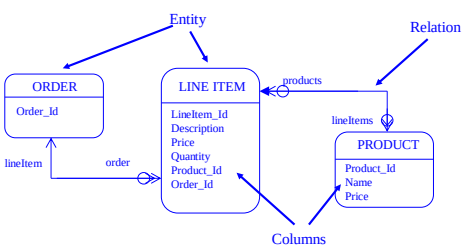
2.2. The Object Model

- The Object Model is composed of
  - Classes (attributes)
  - Relationships
    - Associations
    - Generalization



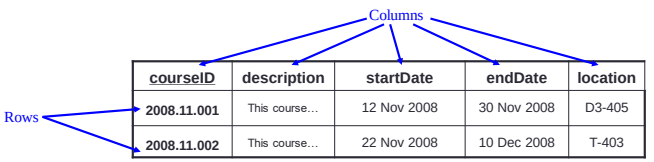
2.3. The Relational Data Model

- Relational data model is composed of
  - Entities - Table
  - Relations - Relationship
  - → **Also called E-R model**



2.3.1. Entities/Tables

- Entities is mapped to table when design physical database
- Including
  - Columns: Attributes
  - Rows: Concrete values of attributes



courseID	description	startDate	endDate	location
2008.11.001	This course…	12 Nov 2008	30 Nov 2008	D3-405
2008.11.002	This course…	22 Nov 2008	10 Dec 2008	T-403

25

2.3.2. Relations/Relationships

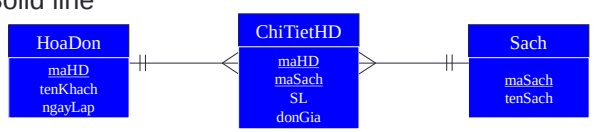
- Relations between entities or relationship between tables
- Multiplicity/Cardinality
  - One-to-one (1:1)
  - One-to-many (1:m)
  - Many-to-one (m:1)
  - Many-to-many (m:n)

(Normally, many-to-many relation is devided to one-to-many and many-to-one relations)

26

Dependency relationships

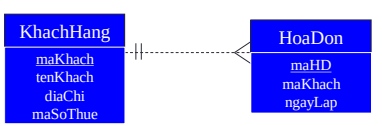
- The child entity can exist only when the parent entity exists
- The child entity has a foreign key referencing to the primary key of the parent entity
- This foreign key is included in the primary key of the child
- Solid line



27

Independency relationships

- The child entity can exist even if the parent entity does not exist
- The child entity has a foreign key referencing to the primary key of the parent entity
- This foreign key is not included in the primary key of the child
- Dash line

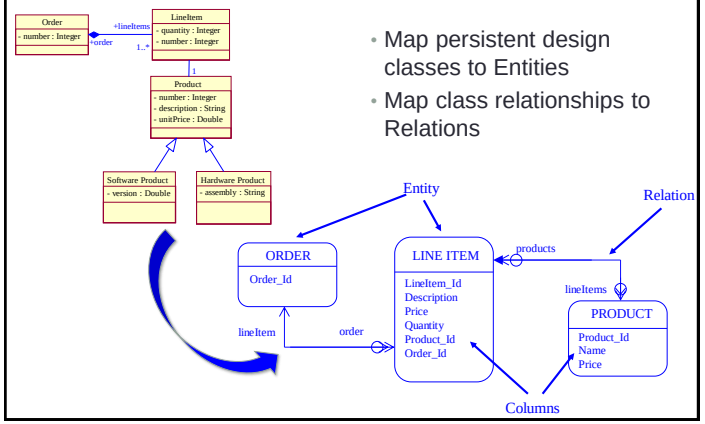


28

Content

29

3. Mapping class diagram to E-R diagram

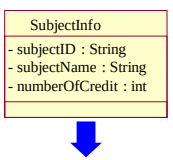


- Map persistent design classes to Entities
- Map class relationships to Relations

30

3.1. Mapping Persistent Design Classes to Entities

- In a relational database
  - Every row is regarded as an object
  - A column in a table is equivalent to a persistent attribute of a class

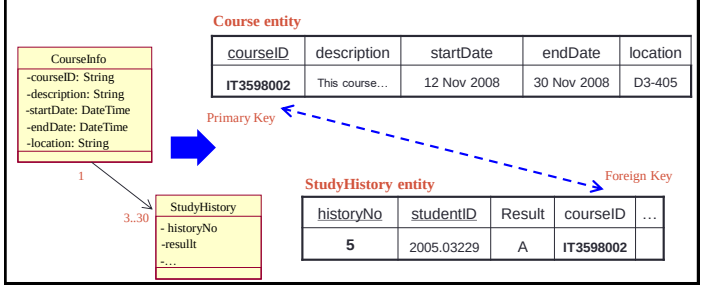


Attributes from object type	subjectID	subjectName	numberOfCredit
Object Instance	IT0001	CS Introduction	4

31

3.2. Mapping Associations Between Persistent Objects

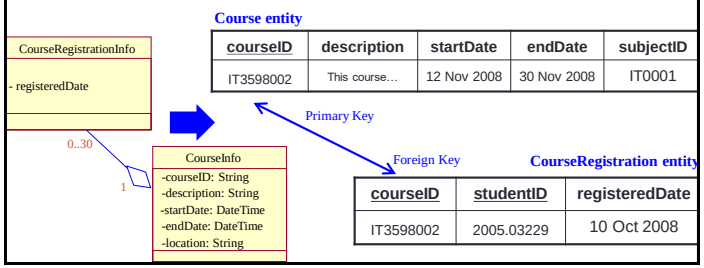
- Associations between two persistent objects are realized as foreign keys to the associated objects.
  - A foreign key (not in primary key) is a column in one table that contains the primary key value of associated object
  - → Independency relationship



32

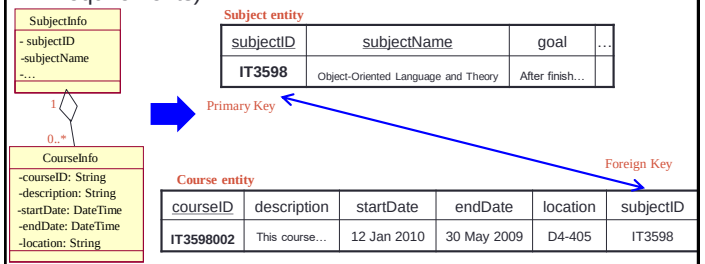
3.3. Mapping Aggregation to the Data Model

- Aggregation is also modeled to dependency relationship using foreign key relationships
  - The use of composition implements a cascading delete constraint

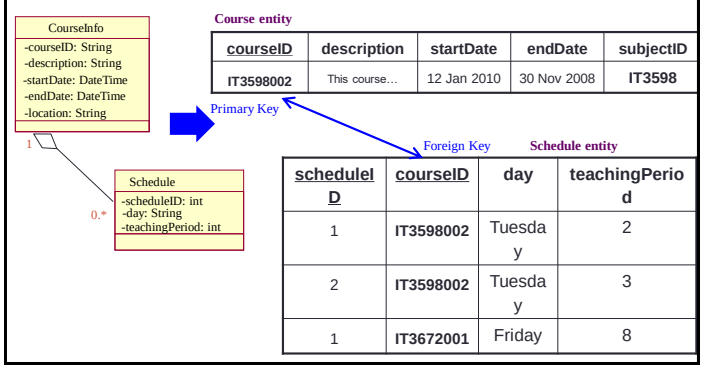


3.3. Mapping Aggregation to the Data Model (2)

- In some case, we can map to independency relationship to simplify the primary key.
- Example: CourseID is the primary key (according the requirements)



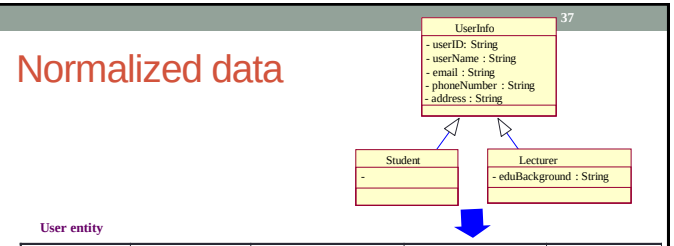
More example in Course Registration



3.4. Modeling Inheritance in the Data Model

- A Data Model does not support modeling inheritance in a direct way
- Two options:
  - Use separate tables (normalized data)
  - Duplicate all inherited associations and attributes (de-normalized data)

Normalized data



User entity

userID	userName	email	phoneNumber	address
2005.03229	Nguyễn Hoàng	student_200503229@hut.edu.vn	0988.394.394	No. 12, …
002.005.00060	Trần Nam Khánh	khanhtn@hut.edu.vn	0912.473.568	No. 157, …

Primary Key

Lecturer entity

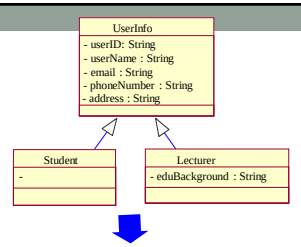
lecturerID	eduBackground
002.005.00060	Master from HUT…

Primary Key

37

37

De-normalized data



Student entity

studentID	studentName	email	phoneNumber	address
2005.03229	Nguyễn Hoàng	student_200503229@hut.edu.vn	0988.394.394	No. 12, …

Replication

Lecturer entity

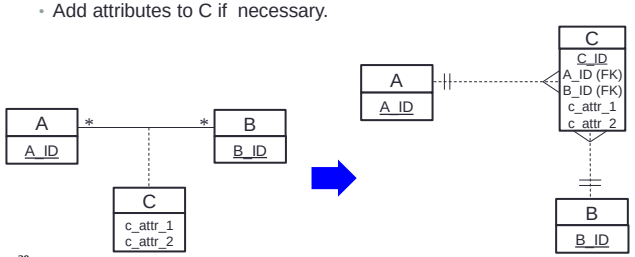
lecturerID	lecturerName	email	phoneNumber	address	eduBackground
002.005.00060	Trần Nam Khánh	khanhtn@hut.edu.vn	0912.473.568	No. 157, …	Master from HUT…

38

38

3.5. Mapping many-to-many cardinality

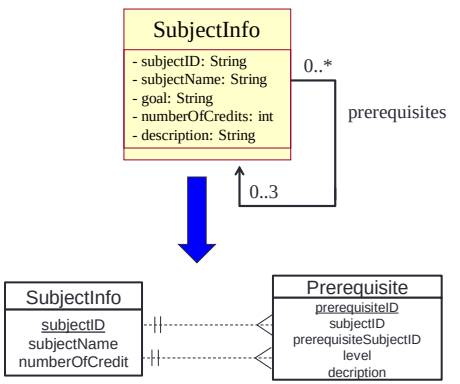
- Use an intermediate entity
- Example: The Cardinality of A and B is many-to-many
  - Add an intermediate entity called “C”
  - Place 2 foreign keys for C, referencing to 2 primary keys of A and B
  - Add attributes to C if necessary.



39

39

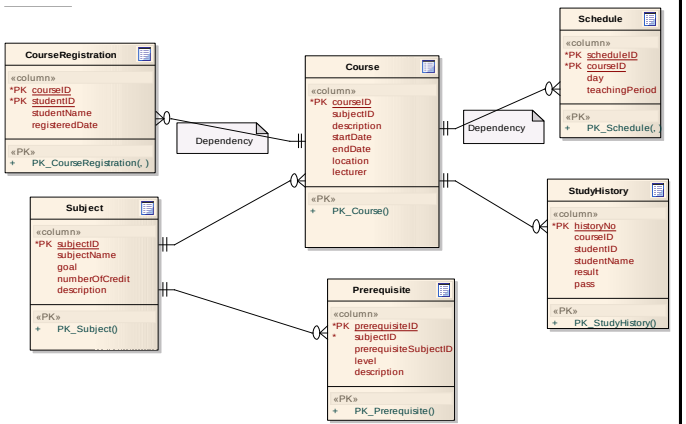
Example in Course Registration CS



40

40

E-R diagram



41

Content

1. Data models
2. Object model and Rational Data Model
3. Mapping class diagram to E-R diagram
4. Normalization

42

4.1. Overview of Normalization

- Normalization: the process of steps that will identify, for elimination, redundancies in a database design.
- Purpose of Normalization: to improve
  - storage efficiency
  - data integrity
  - and scalability

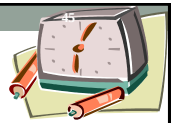
43

4.1. Overview of Normalization (2)

- In relational model, methods exist for quantifying how efficient a database is.
- These classifications are called **normal forms** (or **NF**), and there are algorithms for converting a given database between them.
- Normalization generally involves splitting existing tables into multiple ones, which must be re-joined or linked each time a query is issued

44

4.2. History



- Edgar F. Codd first proposed the process of normalization and what came to be known as the 1st normal form in his paper *A Relational Model of Data for Large Shared Data Banks Codd* stated:

"There is, in fact, a very simple elimination procedure which we shall call normalization. Through decomposition nonsimple domains are replaced by 'domains whose elements are atomic (nondecomposable) values".

45

4.3. Normal Forms

- Edgar F. Codd originally established three normal forms: 1NF, 2NF and 3NF.
- There are now others that are generally accepted, but 3NF is widely considered to be sufficient for most applications.
- Most tables when reaching 3NF are also in BCNF (Boyce-Codd Normal Form).

46

Functionally determines

- In a table, a set of columns X, functionally determines another column Y…

$X \rightarrow Y$

… if and only if each X value is associated with at most one Y value in a table.

- i.e. if you know X then there is only **one** possibility for Y.

47

Normal forms so Far…

- First normal form
  - All data values are atomic, and so everything fits into a mathematical relation.
- Second normal form
  - As 1NF plus no *non-primary-key attribute* is partially dependant on the primary key
- Third normal form
  - As 2NF plus no non-primary-key attribute depends transitively on the primary key

48

Normalization Example

- Consider a table representing orders in an online store
- Each entry in the table represents an item on a particular order. (thinking in terms of records. Yuk.)
- **Columns**
  - Order
  - Product
  - Customer
  - Address
  - Quantity
  - UnitPrice
- Primary key is {Order, Product}

49

Functional Dependencies

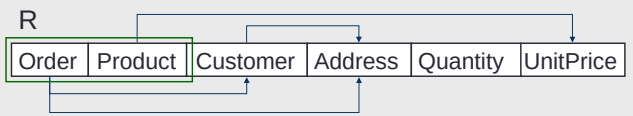
- Each order is for a **single** customer {Order} → {Customer}
- Each customer has a **single** address {Customer} → {Address}
- Each product has a **single** price {Product} → {UnitPrice}
- FD's 1 and 2 are transitive {Order} → {Address}

50

Example – FD Diagram

1NF

R



51

Normalization to 2NF

- Remember 2nd normal form means no partial dependencies on the key. But we have:

  {Order} → {Customer, Address}
  {Product} → {UnitPrice}

  And a primary key of: {Order, Product}

- So to get rid of the first FD we *project* over:

  {Order, Customer, Address}

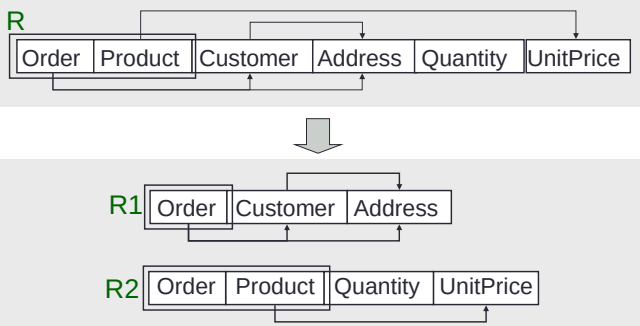
  and

  {Order, Product, Quantity and UnitPrice}

52

Normalization to 2NF

1NF



53

Normalization to 2NF

- R1 is now in 2NF, but there is still a partial FD in R2:

  {Product} → {UnitPrice}

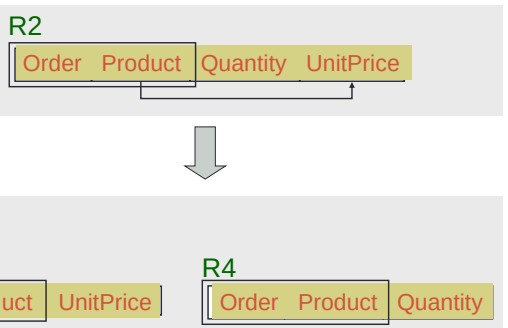
- To remove this we project over:

  {Product, UnitPrice} and {Order, Product, Quantity}

54

Normalization to 2NF

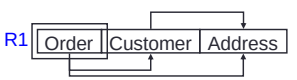
1NF



55

Now let's go 3NF…

- R has now been split into 3 relations - R1, R3, and R4… but R1 has a transitive FD on its key…



{Order} → {Customer} → {Address}

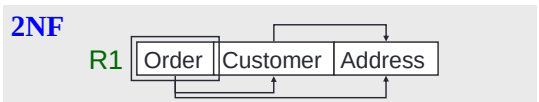
- To remove this problem we project R1 over:

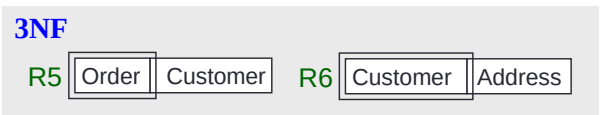
  {Order, Customer} and {Customer, Address}

56

57

So more chopping…





57

58

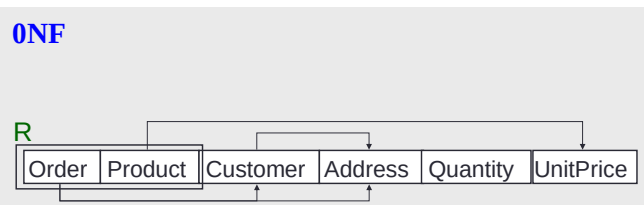
Let's summarize that:

- **1NF:**
  {Order, Product, Customer, Address, Quantity, UnitPrice}
- **2NF:**
  {Order, Customer, Address}
  {Product, UnitPrice}
  {Order, Product, Quantity}
- **3NF:**
  {Product, UnitPrice}
  {Order, Product, Quantity}
  {Order, Customer}
  {Customer, Address}

58

59

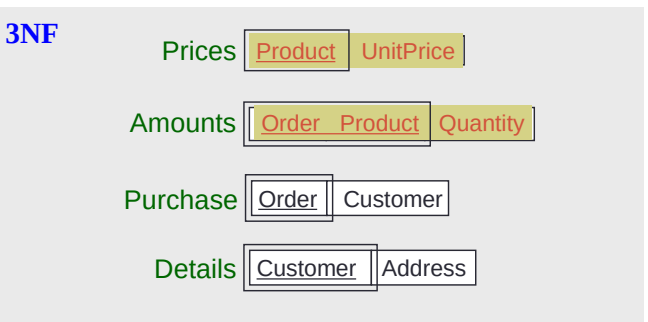
So this…



59

60

has become this…

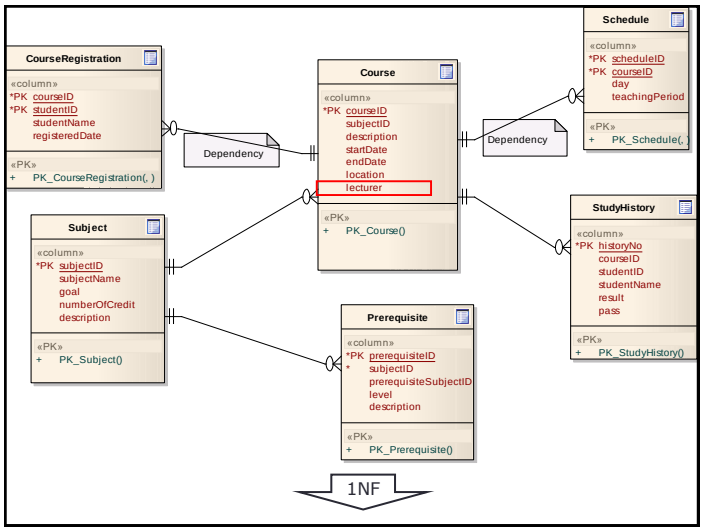


60

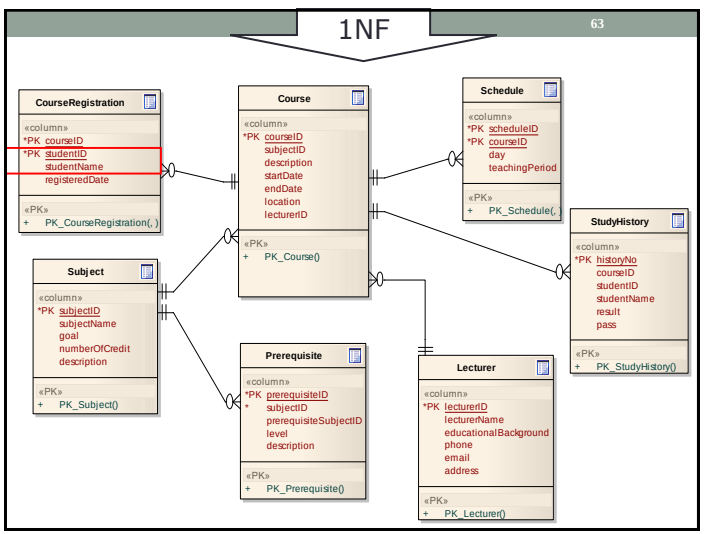
“Register for course” use case

- Make the E-R diagram from the previous step for “Register for course” use case to become:
  - The first normal form
  - The second normal form
  - The third normal form

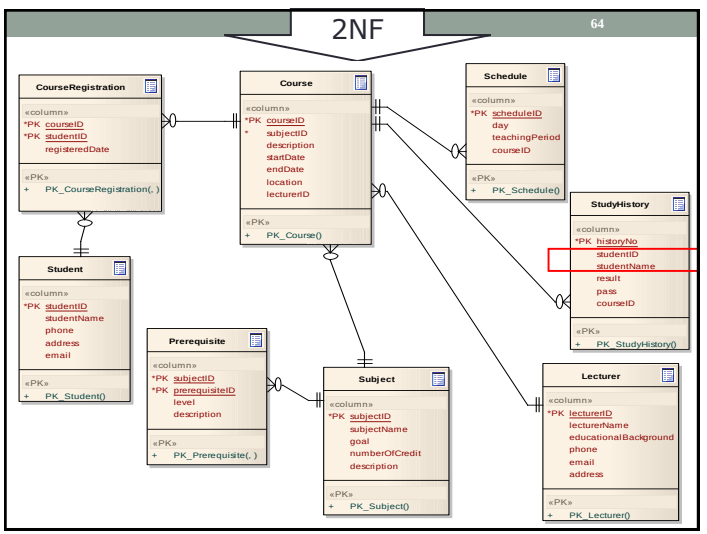
61



62



63



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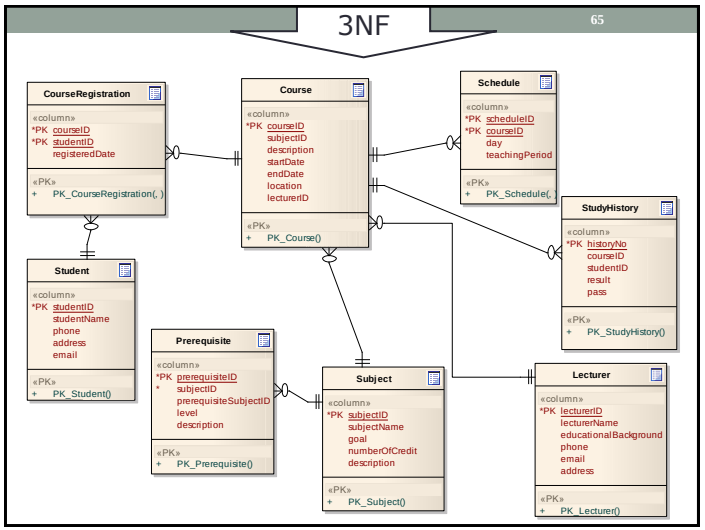
3NF
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CourseRegistration
«column»
*PK courseID
*PK studentID
registeredDate
«PK»
+ PK_CourseRegistration( )
Course
«column»
*PK courseID
subjectID
description
startDate
endDate
location
lecturerID
«PK»
+ PK_Course()
Schedule
«column»
*PK scheduleID
*PK courseID
day
teachingPeriod
«PK»
+ PK_Schedule( )
StudyHistory
«column»
*PK historyNo
courseID
studentID
result
pass
«PK»
+ PK_StudyHistory()
Student
«column»
*PK studentID
studentName
phone
address
email
«PK»
+ PK_Student()
Prerequisite
«column»
*PK prerequisiteID
* subjectID
prerequisiteSubjectID
level
description
«PK»
+ PK_Prerequisite()
Subject
«column»
*PK subjectID
subjectName
goal
numberOfCredit
description
«PK»
+ PK_Subject()
Lecturer
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*PK lecturerID
lecturerName
educationalBackground
phone
email
address
«PK»
+ PK_Lecturer()

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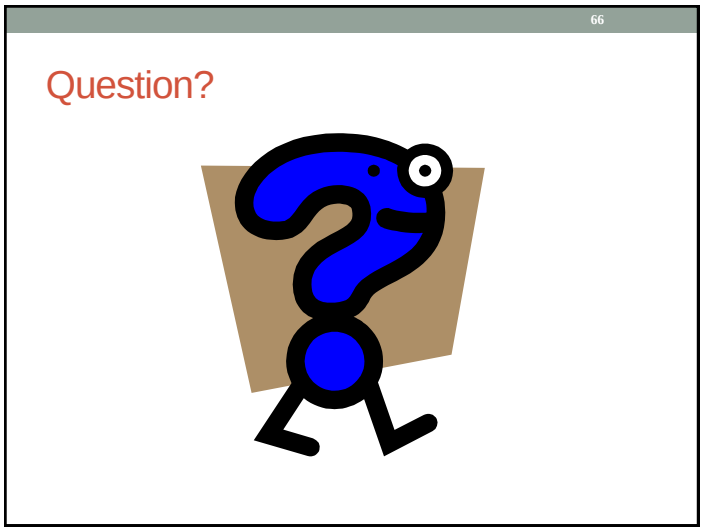
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Question?

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