Pseudo-Boolean Optimization in memory of P.L. Hammer

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Outline

- Pseudo-Boolean Optimization
 - Pseudo-Boolean Functions
 - Representations and Bounds
 - Persistencies and Autarkies
 - Graph Cut Models
 - Implication Networks
- 2 Results
 - Computational Results

Variables and Literals

- Variables: $x_1, x_2, ..., x_n \in \{0, 1\}; V = \{1, 2, ..., n\}$
- Negations: $\bar{x}_i = 1 x_i \in \{0, 1\} \text{ for } i = 1, ..., n$

Pseudo-Boolean Function (PBF):

$$f:\{0,1\}^n\to\mathbb{R}$$

$$f(x_1, ..., x_n) = \sum_{S \subseteq V} a_S \prod_{i \in S} x_i$$

Unconstrained Binary Optimization (PBO

$$\min_{(x_1,...,x_n)\in\{0,1\}^n} f(x_1,...,x_n)$$

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Roof Dual Bound: $C_2(f) \le f$ (Hammer, Hansen and Simeone, 1984)

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Complete Hierarchy of Bounds:

(B, Crama and Hammer, 1990)

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Recent generalizations

Bisubmodular functions: (Kolmogorov, 2010; Kahl and Strandmark, 2011)



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$$S = \{\mathbf{x_1}, \mathbf{x_2}, \mathbf{x_3}\}$$
 and $y = (1, 0, 1)$ is an autarky of the posiform

$$\phi = \mathbf{x_1}\mathbf{x_2} + 5\overline{\mathbf{x_1}}\mathbf{x_3}x_6 + 4\overline{\mathbf{x_2}}\overline{\mathbf{x_3}}x_7 + 4\overline{\mathbf{x_1}}x_4 + 5\mathbf{x_2}x_5 + 6x_4x_5$$

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- hard to test if y is a persistency for a given PBF f.

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For a quadratic function f it is "easy" to find the unique maximal persistent set S(f) provable by quadratic posiforms. (B. and Hammer, 1990)

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- To a submodular QPBF f associate a network G_f as follows

- There is a one-to-one correspondence between values of f and s-t cut values of G_f . (Hammer, 1965)
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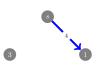
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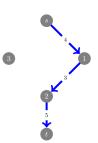




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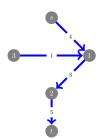
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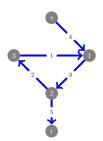
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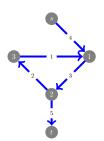
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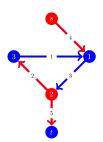
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- A QPBF is submodular IFF all quadratic coefficients are nonpositive. (Doit Yourself, anytime)
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$$\begin{split} f &= 4 - \mathbf{x}_1 + 7\mathbf{x}_2 + \mathbf{x}_3 - 3\mathbf{x}_1\mathbf{x}_2 - \mathbf{x}_1\mathbf{x}_3 - 2\mathbf{x}_2\mathbf{x}_3 \\ &= 4\overline{\mathbf{x}}_1 + 5\mathbf{x}_2 + 3\mathbf{x}_1\overline{\mathbf{x}}_2 + \overline{\mathbf{x}}_1\mathbf{x}_3 + 2\mathbf{x}_2\overline{\mathbf{x}}_3 \\ &= 4s\overline{\mathbf{x}}_1 + 5\mathbf{x}_2\overline{\mathbf{t}} + 3\mathbf{x}_1\overline{\mathbf{x}}_2 + \overline{\mathbf{x}}_1\mathbf{x}_3 + 2\mathbf{x}_2\overline{\mathbf{x}}_3 \end{split}$$

• There is a one-to-one correspondence between values of f and s-t cut values of G_f . (Hammer, 1965)

$$f(0,1,0) = C(\{s,2\},\{1,3,t\}) = 11$$

• Graph cuts (325000) in computer vision: (Greig, Porteous, and Seheult, 1989), (Boykov, Veksler and Zabih, 1998)

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$$f = 10 - 2x_1 - 6x_2 + 2x_1x_2 - 2x_1x_3 + 4x_2x_3$$

$$f = \mathbf{10} - 2\mathbf{x_1} - 6\mathbf{x_2} + 2\mathbf{x_1}\mathbf{x_2} - 2\mathbf{x_1}\mathbf{x_3} + 4\mathbf{x_2}\mathbf{x_3}$$

= $10 - 2x_1 - 6x_2 + 2x_1x_2 - 2x_1(1 - \overline{x_3}) + 4x_2x_3$

$$f = \mathbf{10} - 2\mathbf{x}_1 - 6\mathbf{x}_2 + 2\mathbf{x}_1\mathbf{x}_2 - 2\mathbf{x}_1\mathbf{x}_3 + 4\mathbf{x}_2\mathbf{x}_3$$

= $10 - 2x_1 - 6x_2 + 2x_1x_2 - 2x_1(1 - \overline{x}_3) + 4x_2x_3$
= $10 - 4x_1 - 6x_2 + 2x_1x_2 + 2x_1\overline{x}_3 + 4x_2x_3$

$$f = \mathbf{10} - \mathbf{2x_1} - 6\mathbf{x_2} + \mathbf{2x_1x_2} - \mathbf{2x_1x_3} + 4\mathbf{x_2x_3}$$

$$= 10 - 2x_1 - 6x_2 + 2x_1x_2 - 2x_1(1 - \overline{x_3}) + 4x_2x_3$$

$$= 10 - 4x_1 - 6x_2 + 2x_1x_2 + 2x_1\overline{x_3} + 4x_2x_3$$

$$= 10 - 4(1 - \overline{x_1}) - 6(1 - \overline{x_2}) + 2x_1x_2 + 2x_1\overline{x_3} + 4x_2x_3$$

$$f = \mathbf{10} - 2\mathbf{x}_1 - 6\mathbf{x}_2 + 2\mathbf{x}_1\mathbf{x}_2 - 2\mathbf{x}_1\mathbf{x}_3 + 4\mathbf{x}_2\mathbf{x}_3$$

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$$= 10 - 4(1 - \overline{x}_1) - 6(1 - \overline{x}_2) + 2x_1x_2 + 2x_1\overline{x}_3 + 4x_2x_3$$

$$= 4\overline{x}_1 + 6\overline{x}_2 + 2\mathbf{x}_1\mathbf{x}_2 + 2\mathbf{x}_1\overline{x}_3 + 4\mathbf{x}_2\mathbf{x}_3 = \phi$$

$$f = \mathbf{10} - 2\mathbf{x}_1 - 6\mathbf{x}_2 + 2\mathbf{x}_1\mathbf{x}_2 - 2\mathbf{x}_1\mathbf{x}_3 + 4\mathbf{x}_2\mathbf{x}_3$$

$$= 10 - 2x_1 - 6x_2 + 2x_1x_2 - 2x_1(1 - \overline{x}_3) + 4x_2x_3$$

$$= 10 - 4x_1 - 6x_2 + 2x_1x_2 + 2x_1\overline{x}_3 + 4x_2x_3$$

$$= 10 - 4(1 - \overline{x}_1) - 6(1 - \overline{x}_2) + 2x_1x_2 + 2x_1\overline{x}_3 + 4x_2x_3$$

$$= 4\overline{\mathbf{x}}_1 + 6\overline{\mathbf{x}}_2 + 2\mathbf{x}_1\mathbf{x}_2 + 2\mathbf{x}_1\overline{\mathbf{x}}_3 + 4\mathbf{x}_2\mathbf{x}_3 = \phi$$

$$\mathbf{C}(\phi) = \mathbf{0}$$
 and $\mathbf{S}(\phi) = \emptyset$









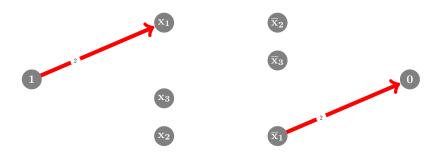




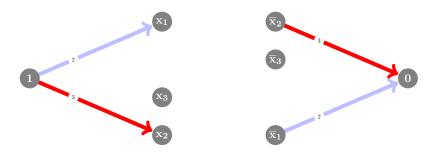
$$\overline{\overline{x}}_1$$

$$f = 4\overline{x}_1 + 6\overline{x}_2 + 2x_1x_2 + 2x_1\overline{x}_3 + 4x_2x_3$$

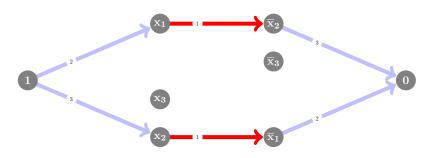




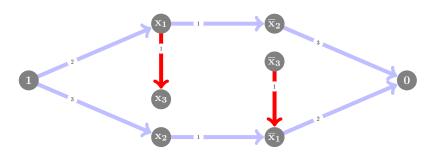
$$f \ = \mathbf{4}\overline{x}_1 + 6\overline{x}_2 + 2x_1x_2 + 2x_1\overline{x}_3 + 4x_2x_3$$



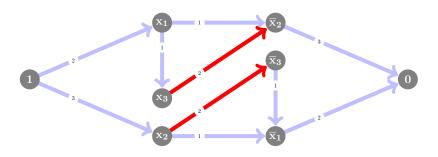
$$f = 4\overline{x}_1 + \textcolor{red}{6}\overline{x}_2 + 2x_1x_2 + 2x_1\overline{x}_3 + 4x_2x_3$$



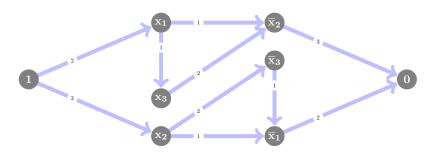
$$f = 4\overline{x}_1 + 6\overline{x}_2 + 2x_1x_2 + 2x_1\overline{x}_3 + 4x_2x_3$$



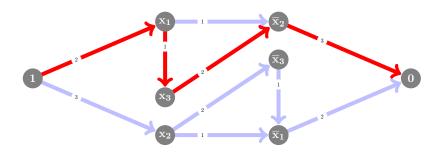
$$f = 4\overline{x}_1 + 6\overline{x}_2 + 2x_1x_2 + \textcolor{red}{2x_1}\overline{x}_3 + 4x_2x_3$$



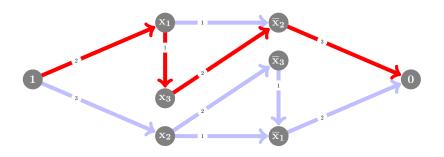
$$f \ = 4\overline{x}_1 + 6\overline{x}_2 + 2x_1x_2 + 2x_1\overline{x}_3 + \textcolor{red}{4x_2x_3}$$



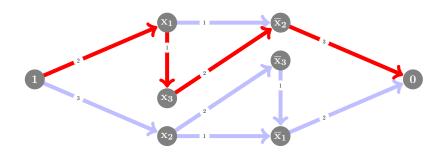
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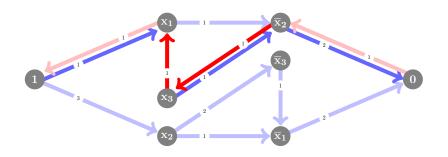
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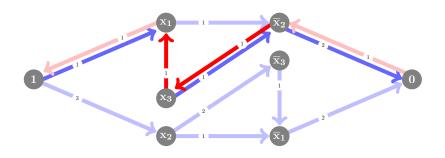
$$\begin{array}{ll} f &= 4\overline{x}_1 + 6\overline{x}_2 + 2x_1x_2 + 2x_1\overline{x}_3 + 4x_2x_3 \\ &= 3\overline{x}_1 + 5\overline{x}_2 + 2x_1x_2 + x_1\overline{x}_3 + 3x_2x_3 \\ &+ (\overline{x}_1 + x_1\overline{x}_3 + x_3x_2 + \overline{x}_2) \end{array}$$



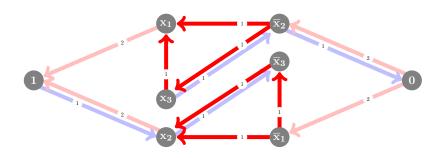
$$\begin{array}{ll} f &= 4\overline{x}_1 + 6\overline{x}_2 + 2x_1x_2 + 2x_1\overline{x}_3 + 4x_2x_3 \\ &= 3\overline{x}_1 + 5\overline{x}_2 + 2x_1x_2 + x_1\overline{x}_3 + 3x_2x_3 \\ &\quad + (\overline{x}_1 + x_1\overline{x}_3 + x_3x_2 + \overline{x}_2) \\ &= 3\overline{x}_1 + 5\overline{x}_2 + 2x_1x_2 + x_1\overline{x}_3 + 3x_2x_3 \\ &\quad + (1 + \overline{x}_1x_3 + \overline{x}_3\overline{x}_2) \end{array}$$



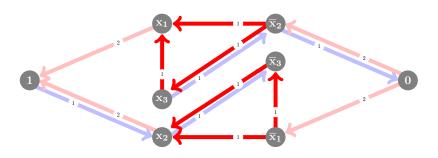
$$\begin{split} f &= 4\overline{x}_1 + 6\overline{x}_2 + 2x_1x_2 + 2x_1\overline{x}_3 + 4x_2x_3 \\ &= 3\overline{x}_1 + 5\overline{x}_2 + 2x_1x_2 + x_1\overline{x}_3 + 3x_2x_3 \\ &+ (\overline{x}_1 + x_1\overline{x}_3 + x_3x_2 + \overline{x}_2) \\ &= 3\overline{x}_1 + 5\overline{x}_2 + 2x_1x_2 + x_1\overline{x}_3 + 3x_2x_3 \\ &+ (1 + \overline{x}_1x_3 + \overline{x}_3\overline{x}_2) \\ &= 1 + 3\overline{x}_1 + 5\overline{x}_2 + 2x_1x_2 + x_1\overline{x}_3 + 3x_2x_3 + \overline{x}_1x_3 + \overline{x}_3\overline{x}_2 \end{split}$$



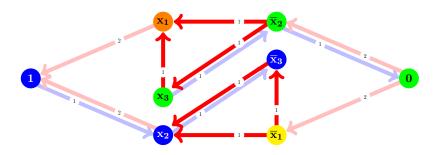
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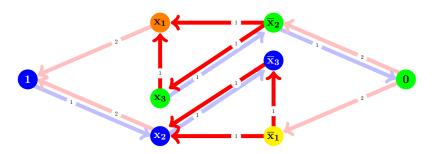
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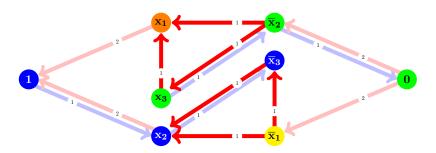
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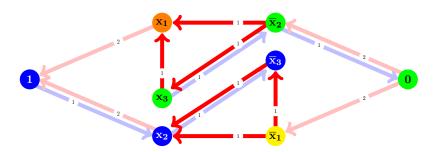
$$\begin{array}{ll} f &= 10 - 2x_1 - 6x_2 + 2x_1x_2 - 2x_1x_3 + 4x_2x_3 \\ &= 4\overline{x}_1 + 6\overline{x}_2 + 2x_1x_2 + 2x_1\overline{x}_3 + 4x_2x_3 \\ &= 4 + 2\overline{x}_2 + 2\overline{x}_1\overline{x}_2 + 2\overline{x}_1x_3 + 2x_2x_3 + 2\overline{x}_2\overline{x}_3 \end{array}$$



- Strong persistency: $x_2 = 1, x_3 = 0$
- Weak persistency: $x_1 = 1$ (or $x_1 = 0$)
- This method computes the unique maximal set of autark variables justifiable by any quadratic posiform of the given quadratic PBF (B. and Hammer, 1990).



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Via Minimization in VLSI Design

| | | Percentage of Variables Fixed by | | | | | |
|---------|------|----------------------------------|---------------------|--------|---------|-------|-------|
| Problem | n | Persist | Persistency Probing | | ALL | Time | |
| | | (strong) | (weak) | (forc) | (equal) | TOOLS | (sec) |
| via.c1y | 829 | 93.6% | 6.4% | 0% | 0% | 100% | 0.03 |
| via.c2y | 981 | 94.7% | 5.3% | 0% | 0% | 100% | 0.06 |
| via.c3y | 1328 | 94.6% | 5.4% | 0% | 0% | 100% | 0.09 |
| via.c4y | 1367 | 96.4% | 3.6% | 0% | 0% | 100% | 0.09 |
| via.c5y | 1203 | 93.1% | 6.9% | 0% | 0% | 100% | 0.08 |
| via.c1n | 828 | 57.4% | 9.6% | 32.4% | 0.6% | 100% | 0.49 |
| via.c2n | 980 | 12.4% | 4.4% | 83.1% | 0.1% | 100% | 7.14 |
| via.c3n | 1327 | 6.8% | 5.7% | 87.3% | 0.2% | 100% | 18.17 |
| via.c4n | 1366 | 11.1% | 1.3% | 87.6% | 0% | 100% | 23.08 |
| via.c5n | 1202 | 3.4% | 1.4% | 95.0% | 0.2% | 100% | 17.13 |

¹S. Homer and M. Peinado. Design and performance of parallel and distributed approximation algorithms for maxcut. Journal of Parallel and Distributed Computing 46 (1997) 48-61.

Vertex Cover in Planar Graphs

| | Averages for 100 graphs in each of the 4 groups | | | | |
|------|---|-------------------|-----------------------|-------------------|--|
| | Variables 1 | Fixed (%) | Time (sec) | | |
| n | A. D. N. ² | \mathbf{QUBO}^3 | A. D. N. ² | \mathbf{QUBO}^3 | |
| 1000 | 68.4 | 100 | 4.06 | 0.05 | |
| 2000 | 67.4 | 100 | 12.24 | 0.16 | |
| 3000 | 65.5 | 100 | 30.90 | 0.27 | |
| 4000 | 62.7 | 100 | 60.45 | 0.53 | |

³Pentium 4, 2.8 GHz, Windows XP, 512 MB

²Alber, Dorn, Niedermeier. Experimental evaluation of a tree decomposition based algorithm for vertex cover on planar graphs. Discrete Applied Mathematics 145 (2005) 219-231; 750 GHz, Linux PC, 720 MB

Jumbo Vertex Cover in Planar Graphs

| | Computing Times (min) ⁴ | | | |
|----------|------------------------------------|-------|-------|--|
| Vertices | Planar Density | | | |
| | 10% | 50% | 90% | |
| 50,000 | 0.7 | 2.3 | 0.9 | |
| 100,000 | 2.9 | 10.2 | 3.9 | |
| 250,000 | 19.5 | 69.8 | 26.3 | |
| 500,000 | 79.3 | 277.3 | 106.9 | |

QUBO fixed all variables for all problems!

⁴Averages over 3 experiments on a Xeon 3.06 GHz, XP, 3.5 GB RAM.

One Dimensional Ising Models

| | | Average Comp | uting Time | (s) |
|----------|-----------------|----------------------------------|-----------------|-------------------|
| σ | Number of Spins | Branch, Cut & Price ⁵ | ${ m BiqMaq^5}$ | \mathbf{QUBO}^6 |
| 2.5 | 100 | 699 | 68 | 1 |
| | 150 | 92 079 | 388 | 3 |
| | 200 | N/A | 993 | 9 |
| | 250 | N/A | 6567 | 14 |
| | 300 | N/A | $34\ 572$ | 21 |
| 3.0 | 100 | 256 | 59 | 1 |
| | 150 | 13 491 | 293 | 2 |
| | 200 | 61 271 | 1 034 | 3 |
| | 250 | 55 795 | 3 594 | 4 |
| | 300 | 55 528 | 8 496 | 5 |

 $^{^5{\}rm F.}$ Rendl, G. Rinaldi, A. Wiegele. (2007). Solving max-cut to optimality by intersecting semidefinite and polyhedral relaxations.

⁶ALL problems were solved by QUBO.

Larger One Dimensional Ising Models

| | | Average of 3 Problems | | |
|----------|------|-----------------------|-------------------|--|
| σ | n | Variables not fixed | QUBO Time $(s)^7$ | |
| 2.5 | 500 | 5 | 13 | |
| | 750 | 22 | 30 | |
| | 1000 | 24 | 53 | |
| | 1250 | 20 | 81 | |
| | 1500 | 32 | 124 | |
| 3.0 | 500 | 0 | 4 | |
| | 750 | 0 | 12 | |
| | 1000 | 0 | 23 | |
| | 1250 | 0 | 37 | |
| | 1500 | 0 | 59 | |

 $^{^{7}}$ Pentium M, 1.6 GHz 760 MB RAM

- How can we extend the above results for general PBF-s?
- How can we find autarkies for a given posiform?
- Quadratization via graph stability (Ebenegger, Hammer, de Werra, 1984): Can we map persistencies?
- Quadratization in higher dimension: given f, find quadratic g such that

$$f(x) = \min_{y \in \{0,1\}^m} g(x,y)$$
 for all $x \in \{0,1\}^n$

Passe it resists? How many new variables do we need? (Greenberg 1972)

Which quadratization provides as with the most associations (8)

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 (Resultant 1975)...
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