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Programming Assignment 3 - Summary

(a) Give an efficient algorithm that takes strings s , x , and y and decides if s is an interleaving of x and y . Derive the computational complexity of your algorithm.

In order to ensure a better time complexity and avoid having to continuously compare all the characters of all three strings, we can use dynamic programming to solve the problem. We will use a 2D boolean array to keep track of the parts of the strings we have traversed and move from the end of the strings to the start of the strings. Through this method, we will be able to recursively compare subsequent characters of the strings. As we will be traversing the lengths of string s and string y and comparing them to string x , I believe the complexity of the algorithm will be $O(nm)$. The $n * m$ represents the length of the string x , (which will be represented by n), and length of string y (which will be represented by m).

Pseudocode:

Boolean Function CheckifInterweaving(String x , String y , String s)

```
1      If  $x = \text{null}$  or  $y = \text{null}$  or  $s = \text{null}$ 
2          return false
3      If length of  $x$  and length of  $y$  and length of  $s = 0$ 
4          return true
5      Boolean 2D Array  $A[][] = [\text{length of } x][\text{length of } y]$ 
6       $A[0][0] = \text{true}$ 
7      for  $i$  to length of  $x$  when  $y$  is empty
8          if current character at  $x$  matches  $s$ 
9               $A[i][0] = \text{true}$ 
10     for  $j$  to length of  $y$  when  $x$  is empty
11         if current character at  $y$  matches  $s$ 
12              $A[0][j] = \text{true}$ 
13     for  $i$  to length of  $x$  when  $y$  is not empty
14         for  $j$  to length of  $y$  when  $x$  is not empty
15             if current character at  $x$  matches  $s$ 
16                  $A[i][j] = A[i - 1][j]$ 
17             if current character at  $y$  matches  $s$ 
18                  $A[i][j] = A[i][j] \text{ or } A[i][j - 1]$ 
19     return  $A[\text{length of } x][\text{length of } y]$ 
```

(b) Implement your algorithm above and test its run time to verify your analysis. Remember that CPU time is not a valid measure for testing run time. You must use something such as the number of comparisons.

As stated in part a, we assumed that the time complexity of the algorithm would be $O(nm)$. We get $n * m$ based on n , the length of string x and m , the length of string y. I ran 4 different tests based on different lengths of strings x, y, and s. The length of string s was a minimum of the length of string x plus the length of string y. I implemented a count of comparisons as the method of testing the complexity. I counted one comparison when any of the strings were compared to each other. (If the comparison count was implemented correctly) As seen below, the number of comparisons in the results, as assumed, correlated with the change in the length of x and y.

Results & Screenshots:

```
Output1: x.length = 3
         y.length = 2
         s.Length = 9
```

Predicted # of comparisons = $3 \times 2 = 6$
Actual # of comparisons = 10

[illegible]

```

- InterweavingStrings (run)
run:
String X: 101
String Y: 00
String S: 100010101

S is an interweaving of X and Y
Number of Comparisons: 10

*****

```

```
Output2: x.length = 10
         y.length = 5
         s.Length = 50
```

Predicted # of comparisons = $5 \cdot 10 = 50$
Actual # of comparisons = 55

[illegible]

```
Output3: x.length = 50
         y.length = 50
         s.Length = 100
```

Predicted # of comparisons = $50 \times 50 = 2500$
Actual # of comparisons = 2501

