Programming Assignment 1 - Summary

1. Closest Pairs

a.) Brute Force Closest Pair

For the brute force method, we would compute all the distances between every pair of points and select the minimum distance. If we are at point n, then we have to find the distances between n and n-1 points. This applies to all n points, so the total distance comparisons between all points can be described as n(n-1). In this calculation, we also double counted distances. For example, the distance between point and A and B is the same as the distance between point B and A. Therefore, the actual number of comparisons would be n(n-1)/2. If we simplify the equation we get $O(n^2)$.

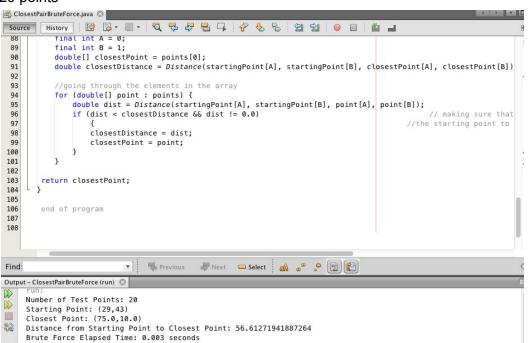
```
Overall we are will be implementing the following:
Closest_Pair(Points P, N)
{
    for i = 0; i < N-1; i++) {
        for (j=1+i; j < N; j++ {
            Get minimum distance(P[i], P[j]);
        }
    }
}
return Closest_Pair</pre>
```

Screenshots:

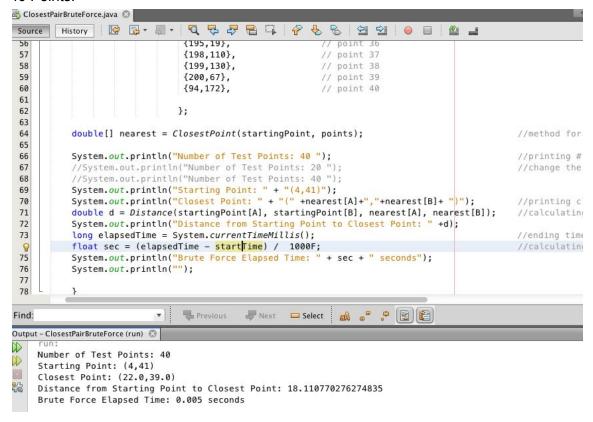
10 Points

```
📆 ClosestPairBruteForce.java 🛭 🛭
Source History
                85
 86
 87
          public static double[] ClosestPoint(double[]startingPoint, double[][]points){
                                                                                             //calculating closes
 88
 89
              final int B = 1;
 90
              double[] closestPoint = points[0];
              double closestDistance = Distance(startingPoint[A], startingPoint[B], closestPoint[A], closestPoint[
 91
 92
 93
              //going through the elements in the array
 94
              for (double[] point : points) {
                  double dist = Distance(startingPoint[A], startingPoint[B], point[A], point[B]);
 95
                                                                                                 // making sure t
 96
                  if (dist < closestDistance && dist != 0.0)</pre>
 97
                                                                                            //the starting point
 98
                     closestDistance = dist;
 99
                     closestPoint = point;
100
101
102
103
          return closestPoint;
104
105
106
       // end of program
107
108
                                           ₩ Next □ Select
                                                            A 6 3 E
Find:
                               Previous
Output - ClosestPairBruteForce (run)
    Number of Test Points: 10
    Starting Point: (4,2)
Closest Point: (34.0,25.0)
    Distance from Starting Point to Closest Point: 33.015148038438355
    Brute Force Elapsed Time: 0.003 seconds
```

20 points



40 Points:



c.) Divide and Conquer Closest Pair

A more efficient way to solve the problem is using the divide & conquer method. In this method, generally a more difficult problem is broken down into many smaller problems. The smaller problems have the same solution as the original large problem. We recursively small the problem and combine all smaller problems at the end.

So for the Divide and Conquer method to find the closest pair we will do the following:

- 1.) split the set of points by a vertical line into two halves
 - a. Sorting points by x and y
 - i. Px, P sorted by x coordinate
 - ii. Py, P sorted by y coordinate
 - b. After x is sorted, we go to the midpoint of x and "draw" a vertical line
 - i. We will have two sets: Left and Right
- 2.) Recursively calculate closest pair in the left and right section
- 3.) Find the closest pair across the vertical dividing line dStrip
 - i. dLeft closest distance in left section dRight closest distance in R
 - ii. d is the minimum of (dLeft, dRight)

- iii. +/- d on both sides of the vertical line will represent the strip = dzone
 - When computing the closest pair across the vertical dividing line will only consider the points within the band
- 4.) Return the shortest distance (dStrip, dLeft, or dRight)

Computing (Px, Py) from P takes O(nlogn) time overall since the sorting time will take O(nlogn) All the steps of the recursive calculation will take O(n) time.

Constructing (Points from Left and Right Regions) from (Px, Py) -- O(n) Constructing Points in strip -- O(n)

Overall: T(n) = 2T(n/2) + n

- The 2T(n/2) for the computation of (Px,Py) from P since we are dividing 2 sets of points of n into n/2.
- The recursive algorithms represents the "n"
- This time complexity represents that of merge sort and so overall T(n) = O(nlogn)

Overall Strategy:

```
Closest_Pair(Px,Py) //assume P split into 2 sorted by x and sorted by y
if(Total Points <= 3) // if you have 3 points or less than calculate using brute force
Calculate Distance using Brute Force Method
Calculate using Divide&Conquer
Construct (points in left and right half)
dLeft = Closest distance in left
dRight = Closest distance in right
Construct all points and find closest distance dStrip
```

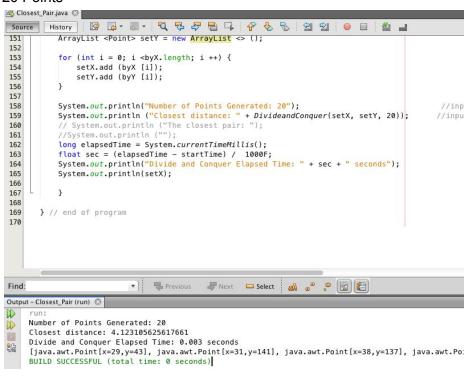
Return dLeft, dRight, or dStrip depending on which is the smallest

Screen Shots:

10 points

```
🖔 Closest_Pair.java 🛭
Source History
                                                     153
                                 for (int i = 0; i <byX.length; i ++) {
154
                                             setX.add (byX [i]);
155
                                             setY.add (byY [i]);
156
157
                                System.out.println("Number of Points Generated: 10");
System.out.println ("Closest distance: " + DivideandConquer(setX, setY, 10));
158
159
                                 // System.out.println ("The closest pair: ");
160
                                 //System.out.println ("");
161
                                 long elapsedTime = System.currentTimeMillis();
162
                                 float sec = (elapsedTime - startTime) / 1000F;
163
                                 System.out.println("Divide and Conquer Elapsed Time: " + sec + " seconds");
164
165
                                 System.out.println(setX);
166
167
168
169
                   } // end of program
170
                                                                                                  Previous
                                                                                                                                                                                                A . . . . .
ind:
                                                                                                                                        Next
                                                                                                                                                                  Select
Output - Closest_Pair (run) 🔘
            Number of Points Generated: 10
            Closest distance: 21.095023109728988
            Divide and Conquer Elapsed Time: 0.006 seconds
            [java.awt.Point[x=34,y=25],\ java.awt.Point[x=42,y=46],\ java.awt.Point[x=53,y=23],\ java.awt.Point[x=54,y=25],\ java.awt.Point[x=54,y=25],\
            BUILD SUCCESSFUL (total time: 0 seconds)
```

20 Points



40 Points

```
🖔 Closest_Pair.java 🛭 🛭
Source History
                 // after calculating midpoint we split the points into a left and right side sorted by
 66
 67
               // each section is put into a list
               ArrayList <Point> xL = new ArrayList <> (); //sorted by x left side
 68
               ArrayList <Point> xR = new ArrayList <> (); //sorted by x right side
ArrayList <Point> yL = new ArrayList <> (); //sorted by y left side
 69
 70
 71
               ArrayList <Point> yR = new ArrayList <> (); //sorted by y right side
 72
 73
               for (int i = 0; i <= mid; i ++) {
 74
75
76
                   xL.add (byX.get (i));
                   yL.add (byY.get (i));
 77
 78
               for (int i = mid + 1; i <n; i ++) {
 79
                   xR.add (byX.get (i));
 80
                   yR.add (byY.get (i));
 81
 82
 83
Find:100
                                Previous
                                            Next - Select
                                                              A 6 5 E
                                                                                                8 matches
Output – Closest_Pair (run) 🛭
    Number of Points Generated: 40
    Closest distance: 4.123105625617661
    Divide and Conquer Elapsed Time: 0.006 seconds
    [java.awt.Point[x=4,y=114], java.awt.Point[x=6,y=153], java.awt.Point[x=13,y=32], java.awt.Point[x=
    BUILD SUCCESSFUL (total time: 0 seconds)
```

100 Points

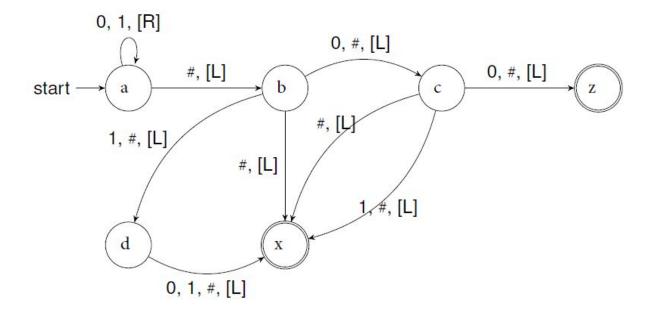
```
Closest_Pair.java 🖾 |
Source History
               ArrayList <Point> setX = new ArrayList <> ();
         ArrayList <Point> setY = new ArrayList <> ();
51
52
         for (int i = 0; i <byX.length; i ++) {</pre>
53
54
             setX.add (byX [i]);
55
             setY.add (byY [i]);
56
57
         System.out.println("Number of Points Generated: 100");
58
59
         System.out.println ("Closest distance: " + DivideandConquer(setX, setY, 100))
60
         // System.out.println ("The closest pair: ");
         //System.out.println ("");
61
         long elapsedTime = System.currentTimeMillis();
62
         float sec = (elapsedTime - startTime) / 1000F;
System.out.println("Divide and Conquer Elapsed Time: " + sec + " seconds");
63
64
65
         System.out.println(setX);
66
67
68
     } // end of program
                            Previous
                                                        AA 60 53 (E) (E)
nd:
                                       Next
                                               Select
utput - Closest_Pair (run) 🔘
   Number of Points Generated: 100
   Closest distance: 1.4142135623730951
   Divide and Conquer Elapsed Time: 0.007 seconds
   BUILD SUCCESSFUL (total time: 0 seconds)
```

Conclusion:

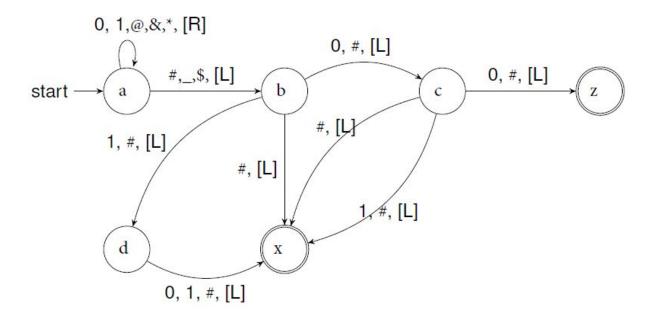
Based on the time complexity of the Brute Force algorithm being linear, I would expect that with the increase in the number of points, the runtime using the brute force strategy would also increase. Whereas, based on the time complexity of the Divide and Conquer algorithm being nlogn, the runtime would eventually level off or increase very slightly with the increase in points. There was a very slight difference between runtimes when running 40 points compared to 100 points. However, there was also a very small change in runtimes when increasing the number of points using the Brute Force algorithm. Overall, I believe the difference between the Brute Force algorithm and Divide and Conquer algorithm would have been more substantial if I was able to test a very large number of points but I was unable to get my divide and conquer program to work with more than 100 points.

2. Turing Machine

Original Turing Machine:



Changes to Turing Machine to implement binary operations:



- Primary changes for addition and multiplication:

State a:

- *
$$\rightarrow$$
 0, [R]

- Changes for subtraction: (not shown on turing machine above)
 - I wasn't able to successfully perform any binary subtraction with the existing turing machine above
 - For subtraction I used additional symbol c under state a in which:

-
$$c \rightarrow c, [R]$$

- for subtraction state a was changed to the following:

-State a:

-
$$1 \rightarrow 0$$
, [R]

$$-$$
 0 \rightarrow 1, [R]

-
$$c \rightarrow c$$
, [R]

- &
$$\rightarrow$$
 0, [R]

-
$$* \rightarrow 0$$
, [R]

ScreenShots

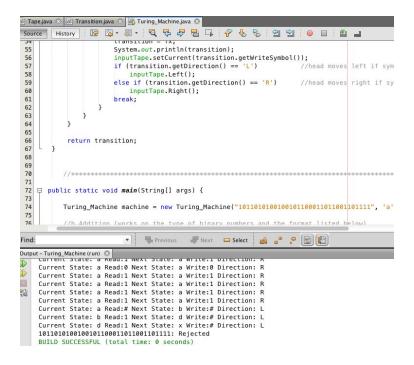
Small input

```
💆 Tape.java 🔞 🙆 Transition.java 🔞 🚳 Turing_Machine.java 🔕
Source History 🕝 👨 🔻 🗸 💆 🔁 📮 🖟 各 🔁 🛂 🥚 🗌
        Sheetal Parikh
Programming Assignment 1 - Question 2
        Turing Machine
 3
4
5
6
7
       Main
       ckage turing_machine;

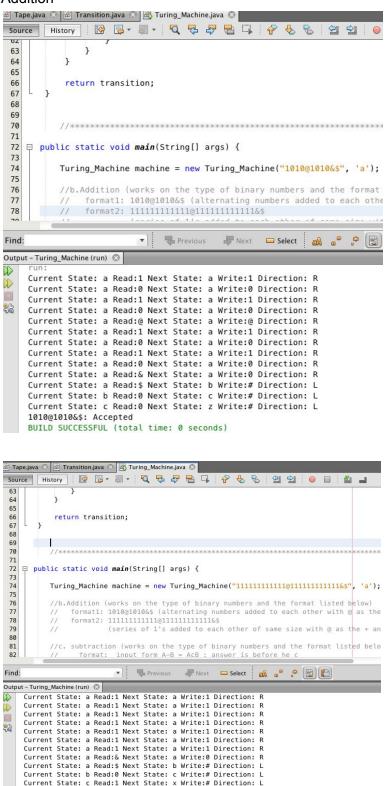
□ port java.util.ArrayList;

10
11
    ₽ *
12
13
14
15
       @author SheetalParikh
       blic class Turing_Machine {
16
94
94
94
21
         private String input;
                                                                                  //tape - use
         private Tape inputTape;
         private ArrayList<Transition> function;
         private char startState;
                                                                                  //state the
         private char currentState;
                                                                                  //the state
22
                              ▼ Previous ₩ Next □ Select ₩ 😅 💝 📳
Find:
Output – Turing_Machine (run) 🛭
     Current State: a Read:1 Next State: a Write:1 Direction: R
    Current State: a Read:0 Next State: a Write:0 Direction: R Current State: a Read:1 Next State: a Write:1 Direction: R
    Current State: a Read:0 Next State: a Write:0 Direction: R
    Current State: a Read:0 Next State: a Write:0 Direction: R Current State: a Read:# Next State: b Write:# Direction: L
     Current State: b Read:0 Next State: c Write:# Direction: L
     Current State: c Read:0 Next State: z Write:# Direction: L
     10100: Accepted
     BUILD SUCCESSFUL (total time: 1 second)
```

Large Input



Addition



```
Subtraction:
Source History 🔯 🖟 . 🖫 . 💆 🔁 🖶 📮 🚱 😓 얼 🗐 .
                     else if (transition.getDirection() == 'R')
 59
                         inputTape.Right();
 61
                     break;
 62
 64
 65
 66
67
          return transition;
 69
 70
 71

□ ublic static void main(String[] args) {
 72
 74
         Turing_Machine machine = new Turing_Machine("01c011", 'a');
 75
 76
         //b.Addition (works on the type of binary numbers and the format listed I
 77
              format1: 1010@1010&$ (alternating numbers added to each other with
 78
              ▼ Previous
                                        ₩ Next □ Select 🔐 😅 🖺
Find:
Output – Turing_Machine (run) 🔕
     Current State: a Read: 0 Next State: a Write: 1 Direction: R
     Current State: a Read:1 Next State: a Write:0 Direction: R
    Current State: a Read:c Next State: a Write:c Direction: R Current State: a Read:0 Next State: a Write:1 Direction: R
     Current State: a Read:1 Next State: a Write:0 Direction: R
     Current State: a Read:1 Next State: a Write:0 Direction: R
     Current State: a Read:# Next State: b Write:# Direction: L
    Current State: b Read:0 Next State: c Write:# Direction: L Current State: c Read:0 Next State: z Write:# Direction: L
     01c011: Accepted
    BUILD SUCCESSFUL (total time: 0 seconds)
🙆 Tape.java 🗵 🙆 Transition.java 🗵 🎒 Turing_Machine.java 🛭 🗀
 Source History 🕼 🖟 🖟 🗸 🗸 🗸 🖧 🔓 🧘 🥱 🖰
  55
                     System.out.println(transition);
                     inputTape.setCurrent(transition.getWriteSymbol());
  57
                     if (transition.getDirection() == 'L')
  58
                         inputTape.Left();
                     else if (transition.getDirection() == 'R')
  59
  60
                         inputTape.Right();
  61
                     break;
  62
  63
  64
  65
        eturn transition;
  66
  67
  68
  69
  70
         ******************
  71
  72
        static void main(String[] args) {
  73
  0
         ring Machine machine = new Turing Machine("011111c000100". 'a
                                    Previous
                                                 ₩ Next
 Find:
                                                          Select
Output - Turing_Machine (run)
      Current State: a Read:1 Next State: a Write:0 Direction: R
      Current State: a Read:1 Next State: a Write:0 Direction: R
 00
      Current State: a Read:c Next State: a Write:c Direction: R
      Current State: a Read:0 Next State: a Write:1 Direction: R
      Current State: a Read:0 Next State: a Write:1 Direction: R
      Current State: a Read:0 Next State: a Write:1 Direction: R
      Current State: a Read:1 Next State: a Write:0 Direction: R
      Current State: a Read:0 Next State: a Write:1 Direction: R
```

Current State: a Read:0 Next State: a Write:1 Direction: R Current State: a Read:# Next State: b Write:# Direction: L Current State: b Read:1 Next State: d Write:# Direction: L Current State: d Read:1 Next State: x Write:# Direction: L

011111c000100: Rejected

Multiplication

