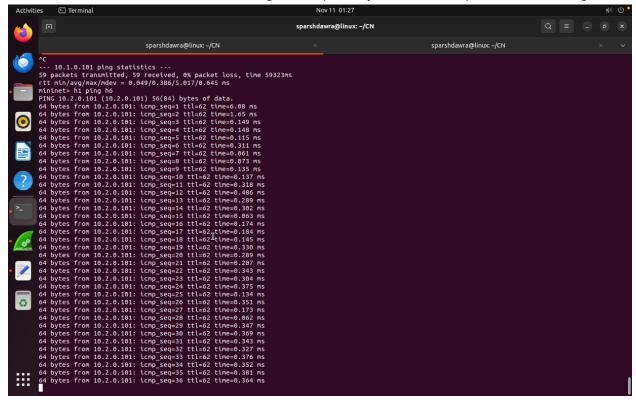
CS433 Assignment-2

Part:1

A) Implementation:

The written code is added in the github repository and here are proof of its working:



PingAll:

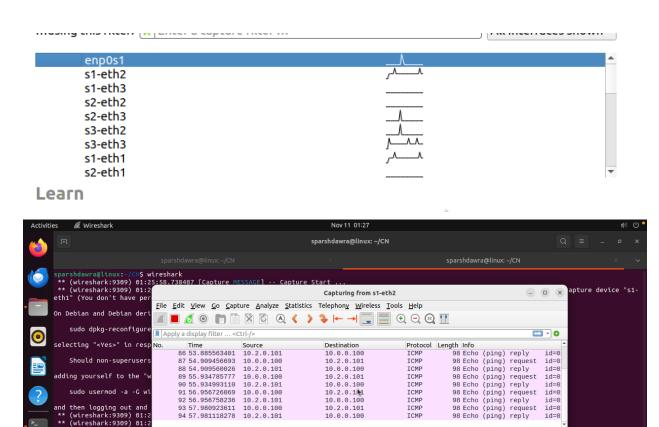
```
*** Starting CLI:
mininet> pingall

*** Ping: testing ping reachability
h1 -> h2 h3 h4 h5 h6 ra rb rc
h2 -> h1 h3 h4 h5 h6 ra rb rc
h3 -> h1 h2 h4 h5 h6 ra rb rc
h4 -> h1 h2 h3 h5 h6 ra rb rc
h5 -> h1 h2 h3 h4 h6 ra rb rc
h6 -> h1 h2 h3 h4 h5 ra rb rc
ra -> h1 h2 h3 h4 h5 ra rb rc
ra -> h1 h2 h3 h4 h5 h6 rb rc
rb -> h1 h2 h3 h4 h5 h6 ra rc
rc -> h1 h2 h3 h4 h5 h6 ra rc
rc -> h1 h2 h3 h4 h5 h6 ra rb

*** Results: 0% dropped (72/72 received)
mininet>
```

B) Observations:

We have pinged H1 to H6 and got the status using wireshark. I have generated the pcapng file for router ra using wireshark and here are some screenshots of it.



10.0.0.100

Frame 1: 70 bytes on wire (560 bits), 70 bytes captured (560 bits) on interface s1-eth2, id 0 Ethernet II, Src: d6:04:ed:67:49:20 (d6:04:ed:67:49:20), Dst: IPv6mcast_02 (33:33:00:00:00:00:02) Internet Protocol Version 6, Src: fe00::d404:edff:fe67:4920, Dst: ff02::2 Internet Control Message Protocol v6

33 -

qΙ

ed 67 49 20 86 dd 60 00 00 00 00 00 00 00 d4 04 00 00 00 00 00 00 00 00 66 14 00 00 00 00 01 01

id=0

id=0:

apture device "s1-

C) I have done the change in the net tabulation for changing the path of h1->ra->rc->h6 to h1->ra->rb->rc->h6. Here are the screenshots for ping as well as iperf.

10.0.0.100

94 57.981118278 10.2.0.101

33 33 00 00 00 02 d6 04 00 00 00 10 3a ff fe 80 ed ff fe 67 49 20 ff 02 00 00 00 00 00 02 85 00 d6 04 ed 67 49 20

and then logging out and

** (wireshark:9309) 01:2

** (wireshark:9309) 01:2

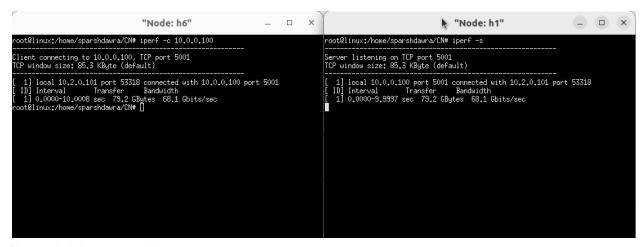
** (wireshark:9309) 01:2
eth1" (You don't have per

On Debian and Debian deri do dpkg-reconfigur selecting "<Yes>" in resu

adding yourself to the

and then logging out and
** (wireshark:9309) 01:2

** (wireshark:9387) 01:2 ** (wireshark:9387) 01:2 ** (wireshark:9387) 01:2



For path h1->ra->rc->h6:

```
*** Starting CLI:
mininet> pingall

*** Ping: testing ping reachability
h1 -> h2 h3 h4 h5 h6 ra rb rc
h2 -> h1 h3 h4 h5 h6 ra rb rc
h3 -> h1 h2 h4 h5 h6 ra rb rc
h4 -> h1 h2 h3 h5 h6 ra rb rc
h5 -> h1 h2 h3 h4 h6 ra rb rc
h6 -> h1 h2 h3 h4 h5 ra rb rc
ra -> h1 h2 h3 h4 h5 ra rb rc
ra -> h1 h2 h3 h4 h5 n6 rb rc
rb -> h1 h2 h3 h4 h5 h6 ra rc
rc -> h1 h2 h3 h4 h5 h6 ra rc
rc -> h1 h2 h3 h4 h5 h6 ra rb

*** Results: 0% dropped (72/72 received)
mininet>
```

```
mininet> h1 ping h6

PING 10.2.0.101 (10.2.0.101) 56(84) bytes of data.

64 bytes from 10.2.0.101: icmp_seq=1 ttl=62 time=6.20 ms

64 bytes from 10.2.0.101: icmp_seq=2 ttl=62 time=1.74 ms

64 bytes from 10.2.0.101: icmp_seq=3 ttl=62 time=0.349 ms

64 bytes from 10.2.0.101: icmp_seq=4 ttl=62 time=0.096 ms

64 bytes from 10.2.0.101: icmp_seq=5 ttl=62 time=0.276 ms

64 bytes from 10.2.0.101: icmp_seq=6 ttl=62 time=0.533 ms

64 bytes from 10.2.0.101: icmp_seq=7 ttl=62 time=0.311 ms

^C

--- 10.2.0.101 ping statistics ---

7 packets transmitted, 7 received, 0% packet loss, time 6099ms

rtt min/avg/max/mdev = 0.096/1.357/6.204/2.041 ms
```

For path h1->ra->rb->rc->h6:

```
*** Starting CLI:
mininet> pingall

*** Ping: testing ping reachability

*** Ping: testing ping: testing ping reachability

*** Ping: testing ping: testing pin
```

```
mininet> h1 ping h6

PING 10.2.0.101 (10.2.0.101) 56(84) bytes of data.

64 bytes from 10.2.0.101: icmp_seq=1 ttl=61 time=5.49 ms

64 bytes from 10.2.0.101: icmp_seq=2 ttl=61 time=1.82 ms

64 bytes from 10.2.0.101: icmp_seq=3 ttl=61 time=0.329 ms

64 bytes from 10.2.0.101: icmp_seq=4 ttl=61 time=0.179 ms

64 bytes from 10.2.0.101: icmp_seq=5 ttl=61 time=0.156 ms

64 bytes from 10.2.0.101: icmp_seq=5 ttl=61 time=0.145 ms

64 bytes from 10.2.0.101: icmp_seq=7 ttl=61 time=0.148 ms

^C

--- 10.2.0.101 ping statistics ---

7 packets transmitted, 7 received, 0% packet ldss, time 6086ms

rtt min/avg/max/mdev = 0.145/1.180/5.487/1.846 ms
```

h1->ra

ra->rb

rb->rc

```
root@linux:/home/sparshdawra/CN# ping 10.2.0.1

PING 10.2.0.1 (10.2.0.1) 56(84) bytes of data.

64 bytes from 10.2.0.1: icmp_seq=1 ttl=64 time=0.334 ms

64 bytes from 10.2.0.1: icmp_seq=2 ttl=64 time=0.035 ms

64 bytes from 10.2.0.1: icmp_seq=3 ttl=64 time=0.075 ms

64 bytes from 10.2.0.1: icmp_seq=4 ttl=64 time=0.047 ms

64 bytes from 10.2.0.1: icmp_seq=5 ttl=64 time=0.135 ms

^C

--- 10.2.0.1 ping statistics ---

5 packets transmitted, 5 received, 0% packet loss, time 4105ms

rtt min/avg/max/mdev = 0.035/0.125/0.334/0.109 ms

root@linux:/home/sparshdawra/CN#
```

rc->h6

```
root@linux:/home/sparshdawra/CN# ping 10.2.0.101
PING 10.2.0.101 (10.2.0.101) 56(84) bytes of data.
64 bytes from 10.2.0.101: icmp_seq=1 ttl=64 time=3.96 ms
64 bytes from 10.2.0.101: icmp_seq=2 ttl=64 time=1.62 ms
64 bytes from 10.2.0.101: icmp_seq=3 ttl=64 time=0.228 ms
64 bytes from 10.2.0.101: icmp_seq=4 ttl=64 time=0.219 ms
64 bytes from 10.2.0.101: icmp_seq=5 ttl=64 time=0.181 ms
^C
--- 10.2.0.101 ping statistics ---
5 packets transmitted, 5 received, 0% packet loss, time 4040ms
rtt min/avg/max/mdev = 0.181/1.241/3.963/1.465 ms
root@linux:/home/sparshdawra/CN#
```

ra->rc

```
root@linux:/home/sparshdawra/CN# ping 10.2.0.1
PING 10.2.0.1 (10.2.0.1) 56(84) bytes of data.
64 bytes from 10.2.0.1: icmp_seq=1 ttl=64 time=0.257 ms
64 bytes from 10.2.0.1: icmp_seq=2 ttl=64 time=0.087 ms
64 bytes from 10.2.0.1: icmp_seq=3 ttl=64 time=0.201 ms
64 bytes from 10.2.0.1: icmp_seq=4 ttl=64 time=0.143 ms
64 bytes from 10.2.0.1: icmp_seq=5 ttl=64 time=0.167 ms
^C
--- 10.2.0.1 ping statistics ---
5 packets transmitted, 5 received, 0% packet loss, time 4074ms
rtt min/avg/max/mdev = 0.087/0.171/0.257/0.056 ms
root@linux:/home/sparshdawra/CN#
```

We can calculate the hop time between h1->ra and ra->rb etc. now,(taking avg value of RTT) Latency1(h1->ra->rb->rc->h6) = RTT(ra)+RTT(rb)+RTT(rc)+RTT(h6) = 2.534 ms Latency2(h1->ra->rc->h6) = RTT(ra)+RTT(rc)+RTT(h6) = 2.406 ms Here, we can see that ra is not directly going to the rc router and it is taking the rb router in its part of connecting h1 to h6.

D) Here are the route tables: For question a:

| Kernel IP routi | .ng table | | | | | | |
|-----------------|-----------|---------------|-------|--------|-----|-----|---------|
| Destination | Gateway | Genmask | Flags | Metric | Ref | Use | Iface |
| 10.0.0.0 | 0.0.0.0 | 255.255.255.0 | U | 0 | 0 | 0 | ra-eth1 |
| 10.0.2.0 | 0.0.0.0 | 255.255.255.0 | U | 0 | 0 | 0 | г1 |
| 10.1.0.0 | 10.0.2.2 | 255.255.255.0 | UG | 0 | 0 | 0 | г1 |
| 10.2.0.0 | 10.2.2.2 | 255.255.255.0 | UG | 0 | 0 | 0 | г5 |
| 10.2.2.0 | 0.0.0.0 | 255.255.255.0 | U | 0 | 0 | 0 | г5 |
| Kernel IP routi | ng table | | | | | | |
| Destination | Gateway | Genmask | Flags | Metric | Ref | Use | Iface |
| 10.0.0.0 | 10.0.2.1 | 255.255.255.0 | UG | 0 | 0 | 0 | г2 |
| 10.0.2.0 | 0.0.0.0 | 255.255.255.0 | U | 0 | 0 | 0 | г2 |
| 10.1.0.0 | 0.0.0.0 | 255.255.255.0 | U | 0 | 0 | 0 | rb-eth1 |
| 10.1.2.0 | 0.0.0.0 | 255.255.255.0 | U | 0 | 0 | 0 | г3 |
| 10.2.0.0 | 10.1.2.2 | 255.255.255.0 | UG | 0 | 0 | 0 | г3 |
| Kernel IP routi | .ng table | | | | | | |
| Destination | Gateway | Genmask | Flags | Metric | Ref | Use | Iface |
| 10.0.0.0 | 10.2.2.1 | 255.255.255.0 | UG | 0 | 0 | 0 | гб |
| 10.1.0.0 | 10.1.2.1 | 255.255.255.0 | UG | 0 | 0 | 0 | г4 |
| 10.1.2.0 | 0.0.0.0 | 255.255.255.0 | U | 0 | 0 | 0 | г4 |
| 10.2.0.0 | 0.0.0.0 | 255.255.255.0 | U | 0 | 0 | 0 | rc-eth1 |
| 10.2.2.0 | 0.0.0.0 | 255.255.255.0 | U | 0 | 0 | 0 | г6 📱 |

For question c:

| Kernel IP rout | ting table | | | | | | |
|----------------|------------|---------------|-------|--------|-----|-----|---------|
| Destination | Gateway | Genmask | Flags | Metric | Ref | Use | Iface |
| 10.0.0.0 | 0.0.0.0 | 255.255.255.0 | U | 0 | 0 | 0 | ra-eth1 |
| 10.0.2.0 | 0.0.0.0 | 255.255.255.0 | U | 0 | 0 | 0 | r1 |
| 10.1.0.0 | 10.0.2.2 | 255.255.255.0 | UG | 0 | 0 | 0 | r1 |
| 10.2.0.0 | 10.0.2.2 | 255.255.255.0 | UG | 0 | 0 | 0 | Γ1 |
| 10.2.2.0 | 0.0.0.0 | 255.255.255.0 | U | 0 | 0 | 0 | r5 |
| Kernel IP rout | ting table | | | | | | |
| Destination | Gateway | Genmask | Flags | Metric | Ref | Use | Iface |
| 10.0.0.0 | 10.0.2.1 | 255.255.255.0 | UG | 0 | 0 | 0 | г2 |
| 10.0.2.0 | 0.0.0.0 | 255.255.255.0 | U | 0 | 0 | 0 | Γ2 |
| 10.1.0.0 | 0.0.0.0 | 255.255.255.0 | U | 0 | 0 | 0 | rb-eth1 |
| 10.1.2.0 | 0.0.0.0 | 255.255.255.0 | U | 0 | 0 | 0 | г3 |
| 10.2.0.0 | 10.1.2.2 | 255.255.255.0 | UG | 0 | 0 | 0 | г3 |
| Kernel IP rout | ting table | | | | | | |
| Destination | Gateway | Genmask | Flags | Metric | Ref | Use | Iface |
| 10.0.0.0 | 10.1.2.1 | 255.255.255.0 | UG | 0 | 0 | 0 | г4 |
| 10.1.0.0 | 10.1.2.1 | 255.255.255.0 | UG | 0 | 0 | 0 | г4 |
| 10.1.2.0 | 0.0.0.0 | 255.255.255.0 | U | 0 | 0 | 0 | г4 |
| 10.2.0.0 | 0.0.0.0 | 255.255.255.0 | U | 0 | 0 | 0 | rc-eth1 |
| 10.2.2.0 | 0.0.0.0 | 255.255.255.0 | U | 0 | 0 | 0 | гб |

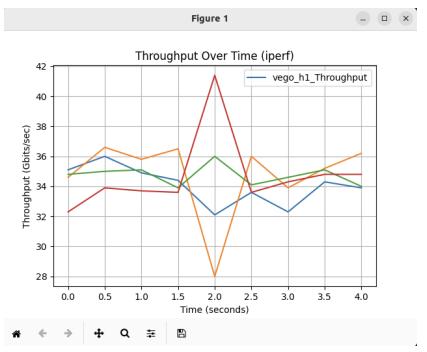
<u> Part:2</u>

A)

```
sparshdawra@linux:~/CN$ sudo python3 Q2.py
h1 h1-eth0:s1-eth1
h2 h2-eth0:s1-eth2
                                   I
h3 h3-eth0:s2-eth1
h4 h4-eth0:s2-eth2
mininet> pingall
*** Ping: testing ping reachability
h1 -> h2 h3 h4
h2 -> h1 h3 h4
h3 -> h1 h2 h4
h4 -> h1 h2 h3
*** Results: 0% dropped (12/12 received)
mininet>
```

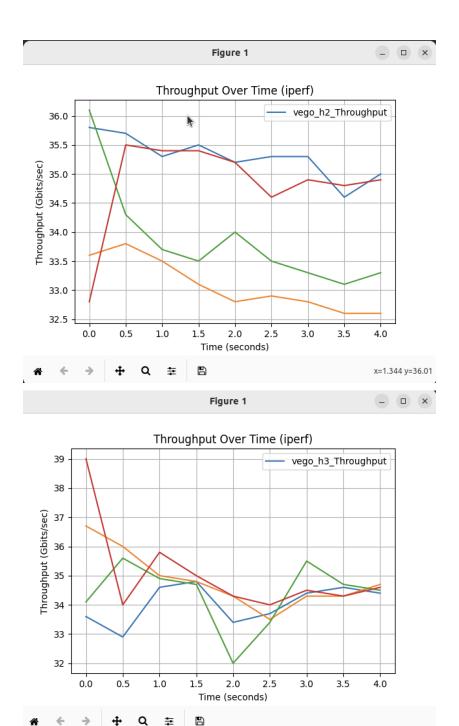
```
CN$ sudo python3 Q2.py --config=b
              h1 h1-eth0:s1-eth1
h2 h2-eth0:s1-eth2
h2 h2-eth0:s1-eth2
h3 h3-eth0:s2-eth1
h4 h4-eth0:s2-eth2
mininet> h1 ping h4
PING 10.0.0.4 (10.0.0.4) 56(84) bytes of data.
64 bytes from 10.0.0.4: icmp_seq=1 ttl=64 time=4.14 ms
64 bytes from 10.0.0.4: icmp_seq=2 ttl=64 time=4.173 ms
64 bytes from 10.0.0.4: icmp_seq=2 ttl=64 time=1.73 ms
64 bytes from 10.0.0.4: icmp_seq=3 ttl=64 time=0.199 ms
64 bytes from 10.0.0.4: icmp_seq=5 ttl=64 time=0.247 ms
64 bytes from 10.0.0.4: icmp_seq=5 ttl=64 time=0.247 ms
64 bytes from 10.0.0.4: icmp_seq=5 ttl=64 time=0.248 ms
64 bytes from 10.0.0.4: icmp_seq=8 ttl=64 time=0.229 ms
64 bytes from 10.0.0.4: icmp_seq=1 ttl=64 time=0.211 ms
64 bytes from 10.0.0.4: icmp_seq=11 ttl=64 time=0.211 ms
64 bytes from 10.0.0.4: icmp_seq=11 ttl=64 time=0.213 ms
64 bytes from 10.0.0.4: icmp_seq=12 ttl=64 time=0.246 ms
64 bytes from 10.0.0.4: icmp_seq=12 ttl=64 time=0.274 ms
64 bytes from 10.0.0.4: icmp_seq=13 ttl=64 time=0.274 ms
64 bytes from 10.0.0.4: icmp_seq=13 ttl=64 time=0.240 ms
64 bytes from 10.0.0.4: icmp_seq=13 ttl=64 time=0.240 ms
64 bytes from 10.0.0.4: icmp_seq=14 ttl=64 time=0.240 ms
64 bytes from 10.0.0.4: icmp_seq=15 ttl=64 time=0.200 ms
64 bytes from 10.0.0.4: icmp_seq=15 ttl=64 time=0.200 ms
64 bytes from 10.0.0.4: icmp_seq=15 ttl=64 time=0.200 ms
64 bytes from 10.0.0.4: icmp_seq=17 ttl=64 time=0.200 ms
64 bytes from 10.0.0.4: icmp_seq=17 ttl=64 time=0.209 ms
64 bytes from 10.0.0.4: icmp_seq=17 ttl=64 time=0.209 ms
64 bytes from 10.0.0.4: icmp_seq=21 ttl=64 time=0.215 ms
64 bytes from 10.0.0.4: icmp_seq=22 ttl=64 time=0.215 ms
64 bytes from 10.0.0.4: icmp_seq=23 ttl=64 time=0.215 ms
64 bytes from 10.0.0.4: icmp_seq=27 ttl=64 time=0.21
                 h3 h3-eth0:s2-eth1
```

- B) In all the graph,
 - Blue->Vegas
 - Orange->Reno
 - Red->Cubic
 - Green->BBR

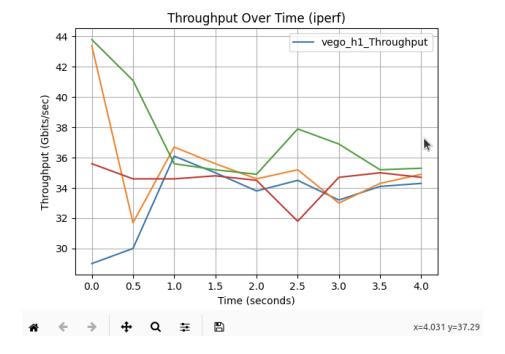


Vegas prioritizes low latency, reacting to round-trip time changes; Reno reacts traditionally to packet loss; Cubic optimizes high-speed network performance with scalable dynamics; BBR dynamically adjusts based on bottleneck bandwidth and round-trip time. Differences arise from varied congestion response algorithms, packet loss interpretation, adaptability to network conditions, and feedback mechanisms. The observed variations in throughput reflect each scheme's unique design for congestion control, demonstrating their distinct approaches and performances in specific network scenarios.

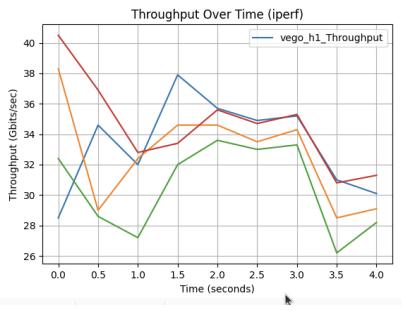
C) H1 client has been shown above. And H2 and H3 as clients are shown below.



According to our Observations, the differences in average throughput among h1, h2, and h3 can be attributed to a combination of path characteristics, initial conditions, randomness in network events, competition for resources, and the dynamic nature of TCP congestion control algorithms. Analyzing the specific conditions of your network scenario and the behavior of each host's TCP connection can provide more insights into the observed throughput variations.



Link-Loss->3%



According to our observation, higher throughput observed in the case of 1% link loss compared to 3% link loss is likely due to the more moderate impact of packet loss on the TCP congestion control mechanisms. The TCP algorithms respond less aggressively to occasional and less severe packet loss, allowing for faster recovery and maintaining higher average throughput.