

Last Time: Multirate Operations

Downsampling by Factor 2



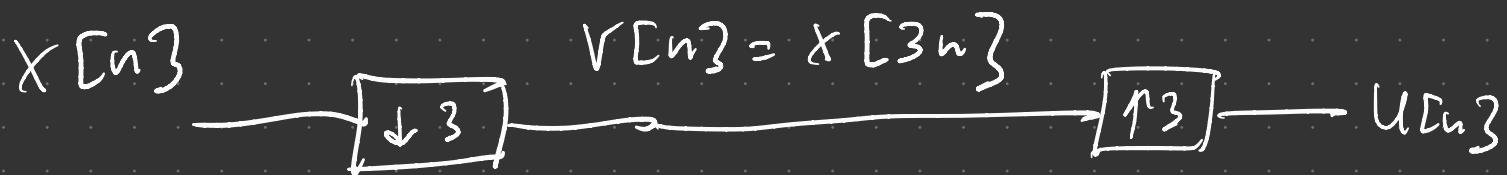
$$\bullet U(e^{j\omega}) = V(e^{j2\omega})$$

$$\bullet U[n] = \frac{1}{2} (x[n] + (-1)^n x[n])$$

$$\Rightarrow U(e^{j\omega}) = \frac{1}{2} (X(e^{j\omega}) + X(e^{j(\omega - \pi)}))$$

$$\Rightarrow V(e^{j\omega}) = \frac{1}{2} \left(X\left(e^{j\frac{\omega}{2}}\right) + X\left(e^{j\left(\frac{\omega - 2\pi}{2}\right)}\right) \right)$$

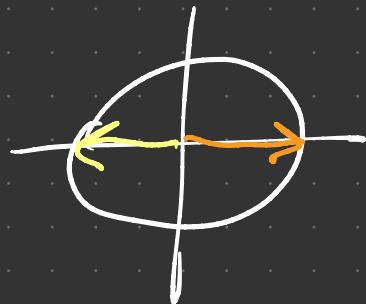
Downsampling by factor 3



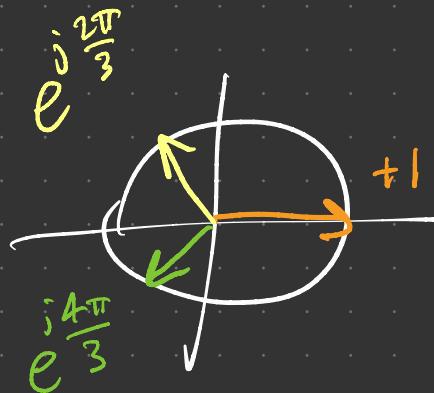
How do we generalize: (by 2)

$$U[n] = \frac{1}{2} \left((+1)^n X[n] + (-1)^n X[n] \right)$$

Obs: ± 1 are the 2nd roots of unity



Claim: 3rd roots of unity



$$\Rightarrow U[n] = \frac{1}{3} \left(X[n] + e^{j\frac{2\pi}{3}n} X[n] + e^{j\frac{4\pi}{3}n} X[n] \right)$$

$$\cdot U(e^{j\omega}) = \frac{1}{3} \sum_{k=0}^2 X\left(e^{j\left(\omega - \frac{2k\pi}{3}\right)}\right)$$

$$\cdot V(e^{j\omega}) = \sqrt{U(e^{j3\omega})}$$

$$V(e^{j\omega}) = \frac{1}{3} \sum_{k=0}^2 X\left(e^{j\left(\frac{\omega - 2k\pi}{3}\right)}\right)$$

Remark: Before down sampling by factor 3,
band limit signal to $\left[-\frac{\pi}{3}, \frac{\pi}{3}\right]$
to avoid aliasing.

Down sample by factor M

$$x[n] \xrightarrow{\downarrow M} v[n] = x[Mn]$$

$$V(e^{j\omega}) = \frac{1}{M} \sum_{k=0}^{M-1} X\left(e^{j\left(\frac{\omega - 2k\pi}{M}\right)}\right)$$

$$V(z) = \frac{1}{M} \sum_{k=0}^{M-1} X\left(z^{\frac{1}{M}} e^{-j\frac{2k\pi}{M}}\right)$$

Exercise: Compare

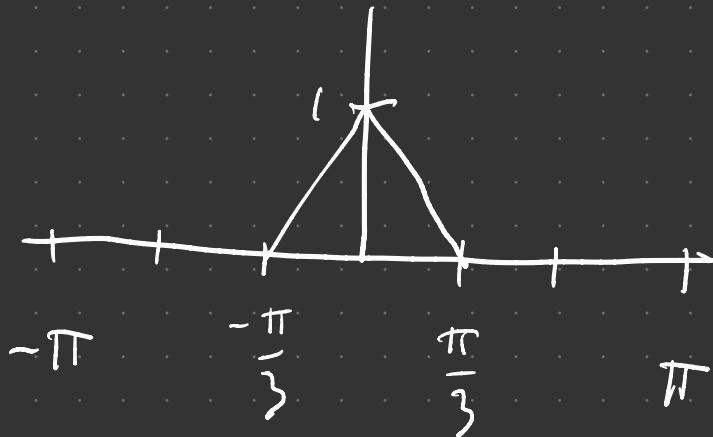


keeps every Mth
coeff. and sets
the rest to 0.

identity

$$\frac{Ex}{x \ln 3} \rightarrow \boxed{\sqrt{3}} \rightarrow V \ln 3$$

$$X(e^{j\omega})$$

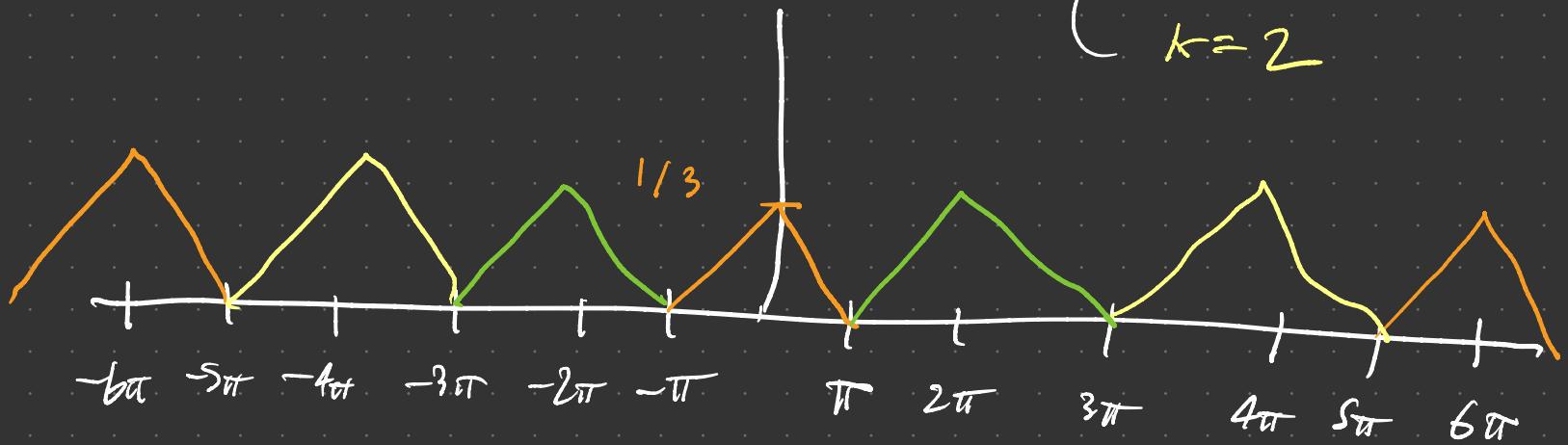


$$V(e^{j\omega}) = \frac{1}{3} \sum_{k=0}^2 X(e^{j(\frac{\omega - 2k\pi}{3})})$$

No aliasing

All together
are valid

$k=0$
Valid?
Not 2π -periodic.
 $k=1$
 $k=2$



Noble Identities

Recall: Multirate DSP is hard because operations don't commute.

- up/down sampling is time-varying

Noble Identities help to quickly analyze multirate systems.

①



②



③

If $\gcd\{M, L\}$:

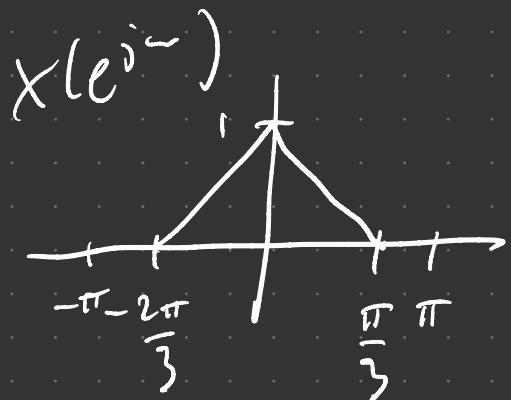
Ex: • $M=3, L=4$ ✓

• $M=2, L=4$ ✗



Proof: Do it at home (Ch. 3, sec. 4).

Exercise :



What is $Y(e^{j\omega})$ in terms of $X(e^{j\omega})$?

Solⁿ: Don't do it with brute force!

$$X(z) \rightarrow \boxed{z^{-3}} \rightarrow z^{-3} X(z) \quad \text{annoying to deal with...}$$

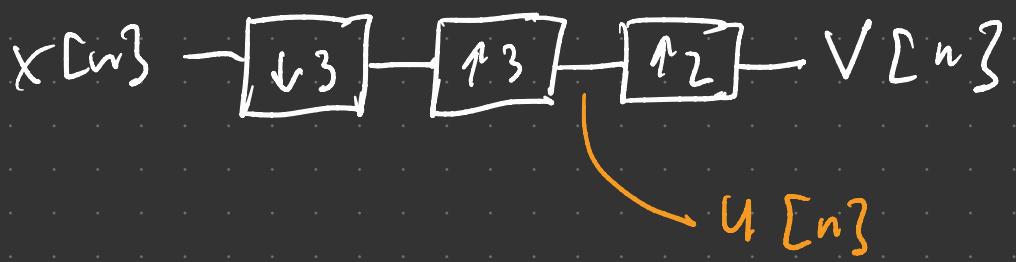
By ③ : $\rightarrow \boxed{z^{-3}} \rightarrow \boxed{\downarrow 3} \rightarrow \boxed{\uparrow 2} \rightarrow \boxed{\uparrow 3} \rightarrow$

By ① : $\rightarrow \boxed{\downarrow 3} \rightarrow \boxed{z^{-1}} \rightarrow \boxed{\uparrow 2} \rightarrow \boxed{\uparrow 3} \rightarrow$

By ② : $\rightarrow \boxed{\downarrow 3} \rightarrow \boxed{\uparrow 2} \rightarrow \boxed{z^{-2}} \rightarrow \boxed{\uparrow 3} \rightarrow$

By ② : $\rightarrow \boxed{\downarrow 3} \rightarrow \boxed{\uparrow 2} \rightarrow \boxed{\uparrow 3} \rightarrow \boxed{z^{-6}} \rightarrow Y[n] = V[n-6]$

$V[n]$



$$U(z) = \frac{1}{3} \sum_{k=0}^2 X(z e^{-j \frac{2k\pi}{3}})$$

$$V(z) = \frac{1}{3} \sum_{k=0}^2 X\left(z^2 e^{-j \frac{2k\pi}{3}}\right)$$

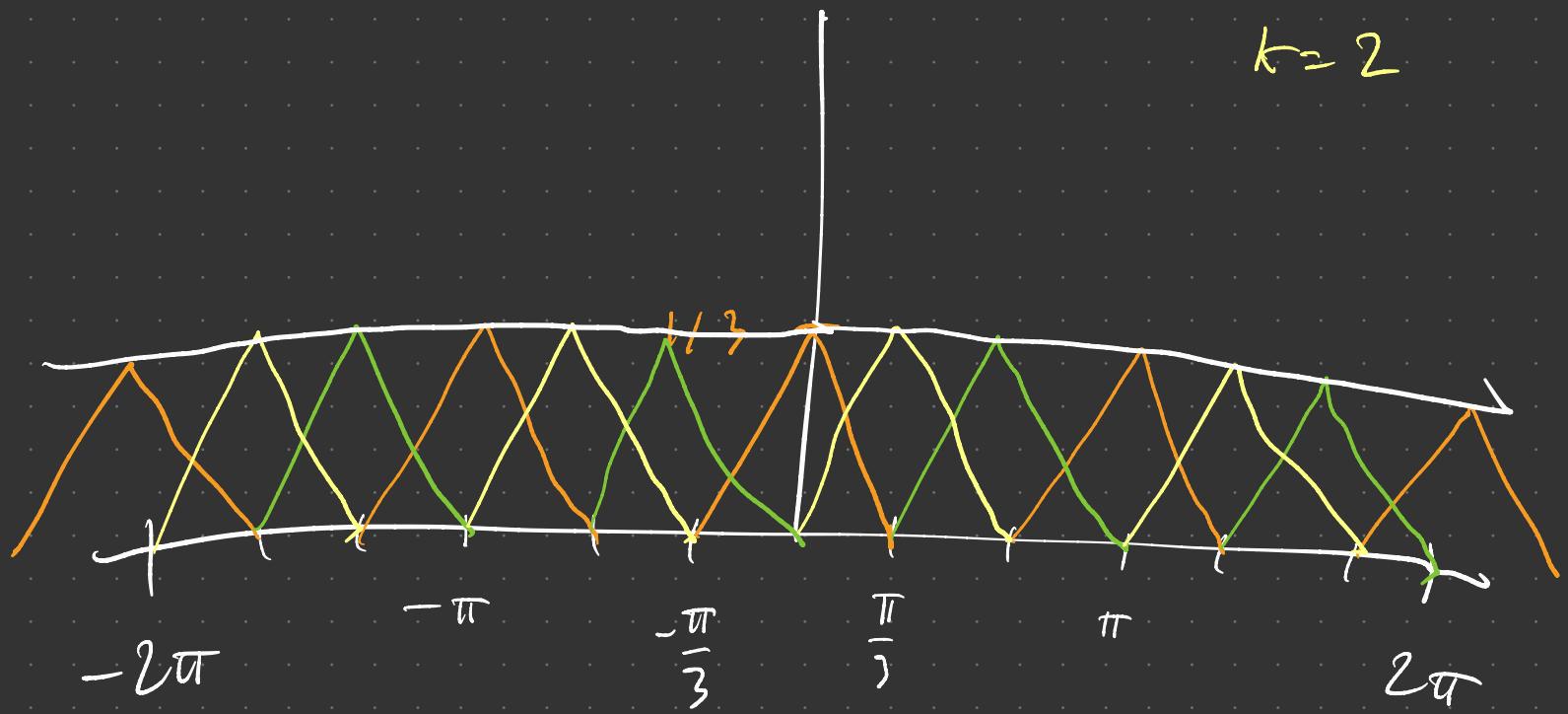
$$V(e^{j\omega}) = \frac{1}{3} \sum_{k=0}^2 X\left(e^{j2(\omega - \frac{k\pi}{3})}\right)$$

$$Y(e^{j\omega}) = e^{-j\omega} \cdot \frac{1}{3} \sum_{k=0}^2 X\left(e^{j2(\omega - \frac{k\pi}{3})}\right)$$

Exer: Try different orders at the Noble identities.

$$|\mathcal{Y}(e^{j\omega})|$$

$k=0$
 $k=1$
 $k=2$



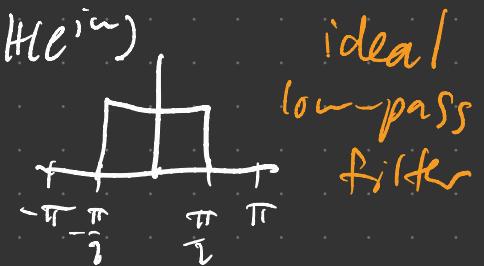
Obs: Severe aliasing?

Q: How could we see this before?

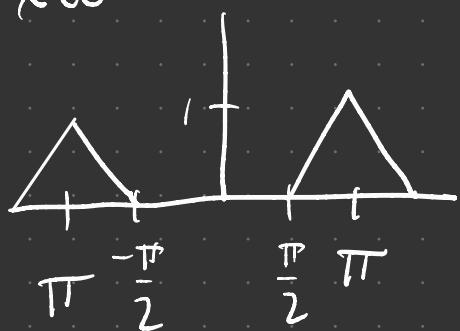
A: Input NOT band limited

to $\left[-\frac{\pi}{3}, \frac{\pi}{3}\right]$ and we immediately downsample by 3.

Exercise:



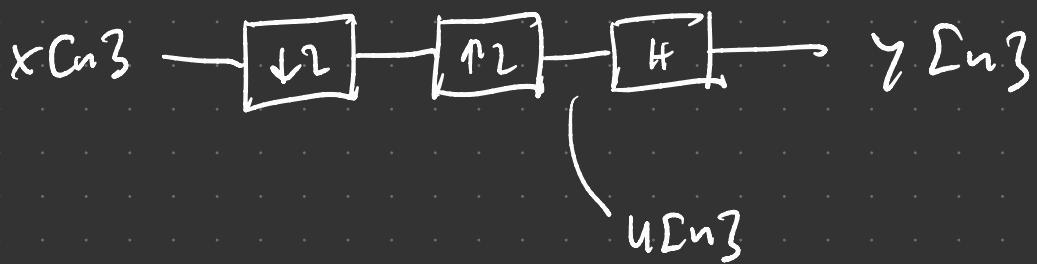
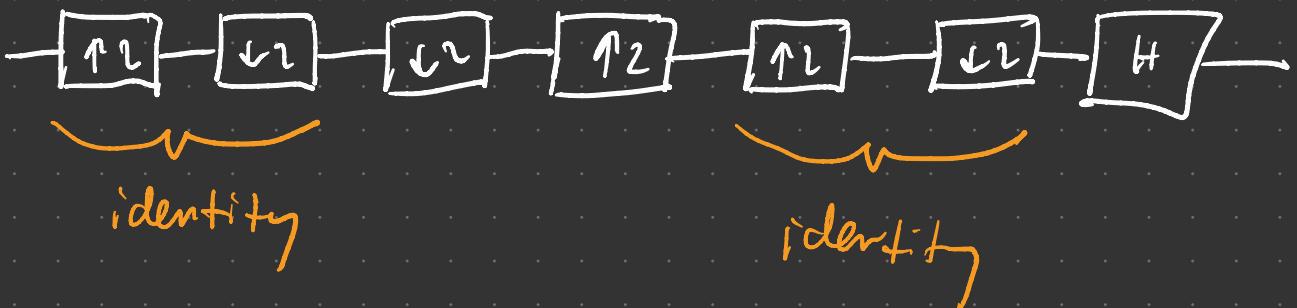
$x(e^{j\omega})$



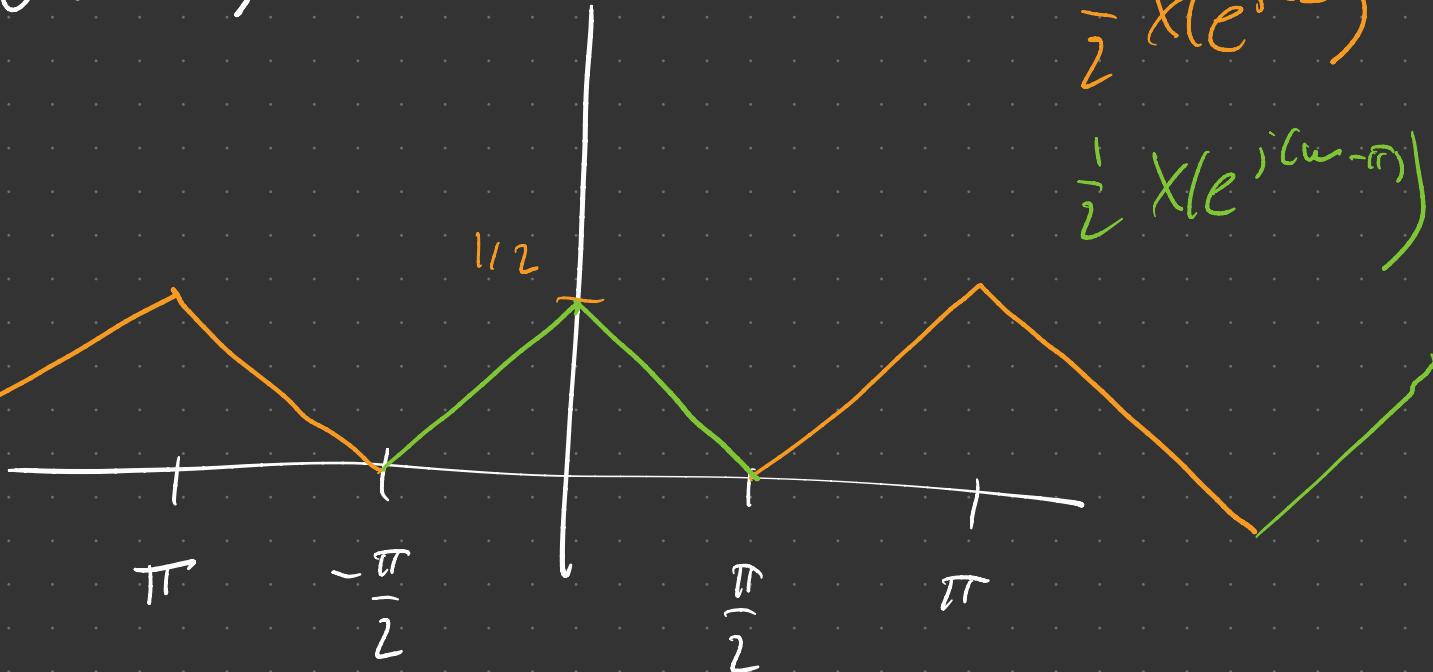
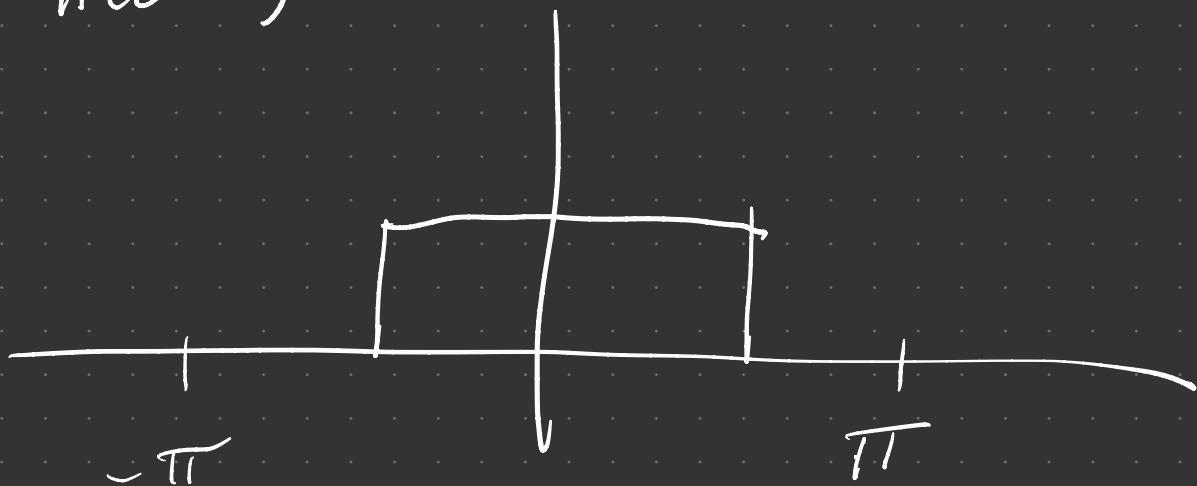
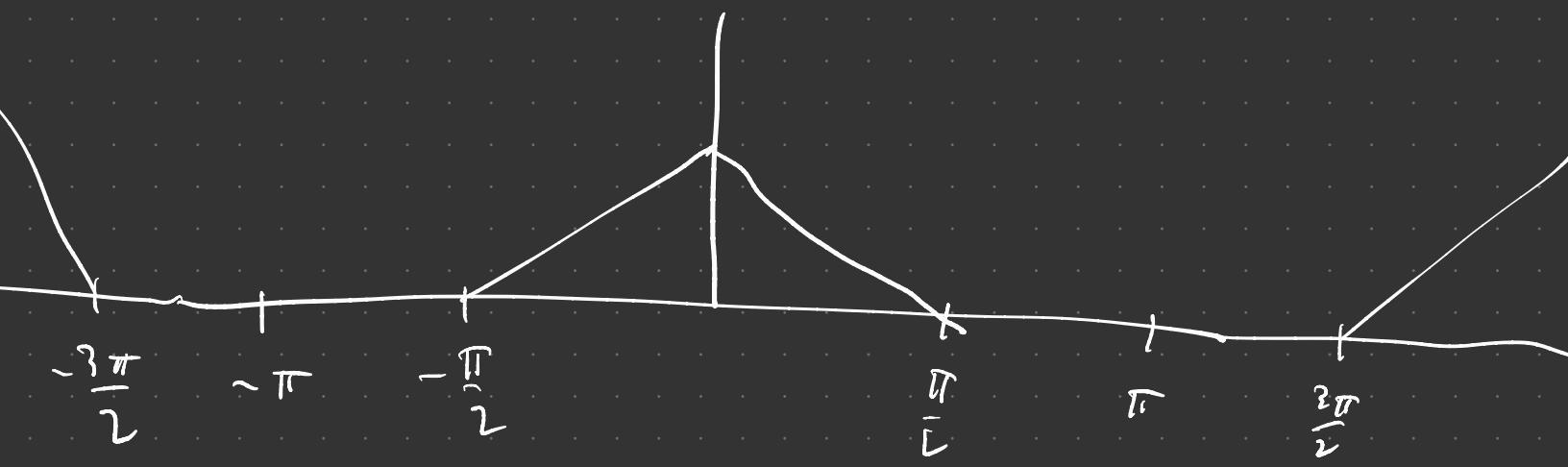
What is $y(e^{j\omega})$

in terms of $x(e^{j\omega})$?

Solⁿ:



$$U(e^{j\omega}) = \frac{1}{2} \left(X(e^{j\omega}) + X(e^{j(\omega-\pi)}) \right)$$

$$U(e^{j\omega})$$

$$H(e^{j\omega})$$

$$Y(e^{j\omega})$$


Polyphase Representation

Recall:



This is only for theory. Extremely wasteful in practice.

Thought Experiment:

Imagine $L = M = 1024$.

- F is processing a lot of zeros
 - Expensive
- Throwing away most of the computations after H .

Q: Can we filter before upsampling?

Q: Can we filter after down sampling?

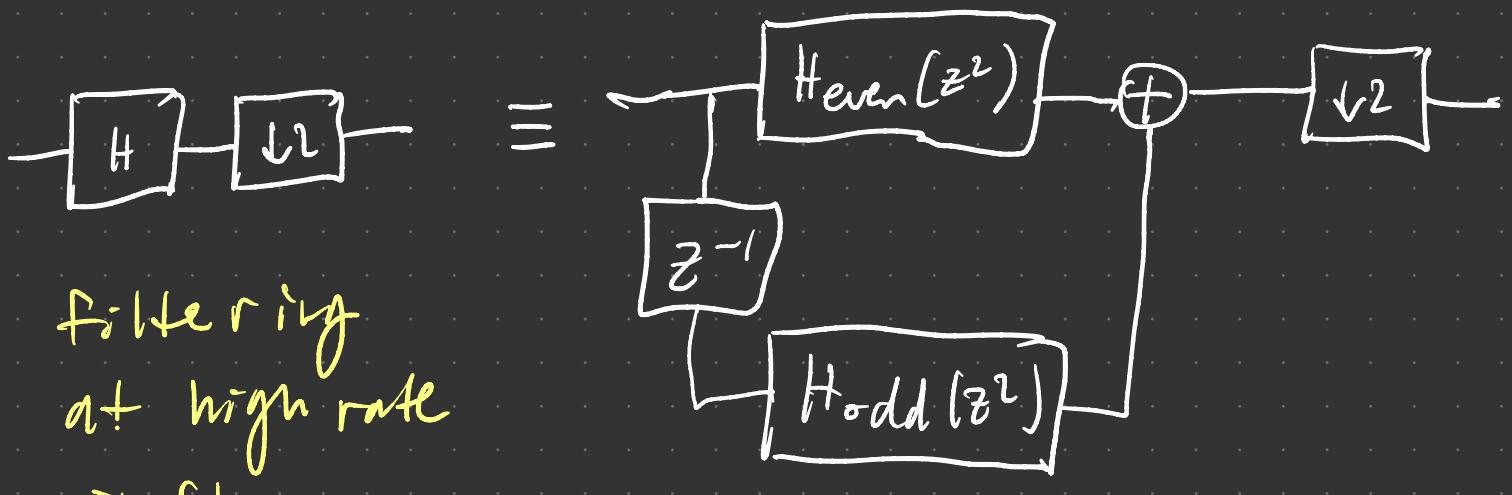
$$\begin{aligned}
 \text{Ex: } H(z) &= 1 + 2z^{-1} + 3z^{-2} + 4z^{-3} \\
 &= (1 + 3z^{-2}) + (2z^{-1} + 4z^{-3}) \\
 &\quad \text{even powers} \qquad \text{odd powers} \\
 &= (1 + 3z^{-2}) + z^{-1}(2 + 4z^{-2}) \\
 &= H_{\text{even}}(z^2) + z^{-1}H_{\text{odd}}(z^2),
 \end{aligned}$$

where

$$H_{\text{even}}(z) = (1 + 3z^{-1})$$

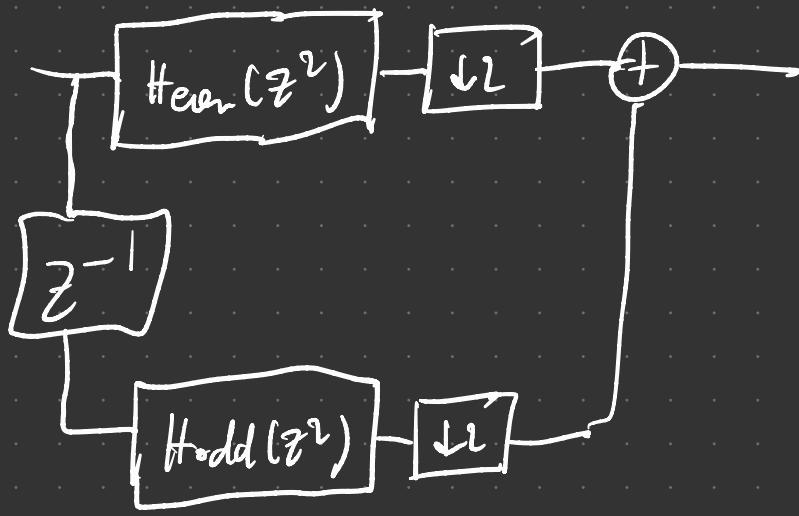
$$H_{\text{odd}}(z) = (2 + 4z^{-1})$$

These are the even/odd polyphases.

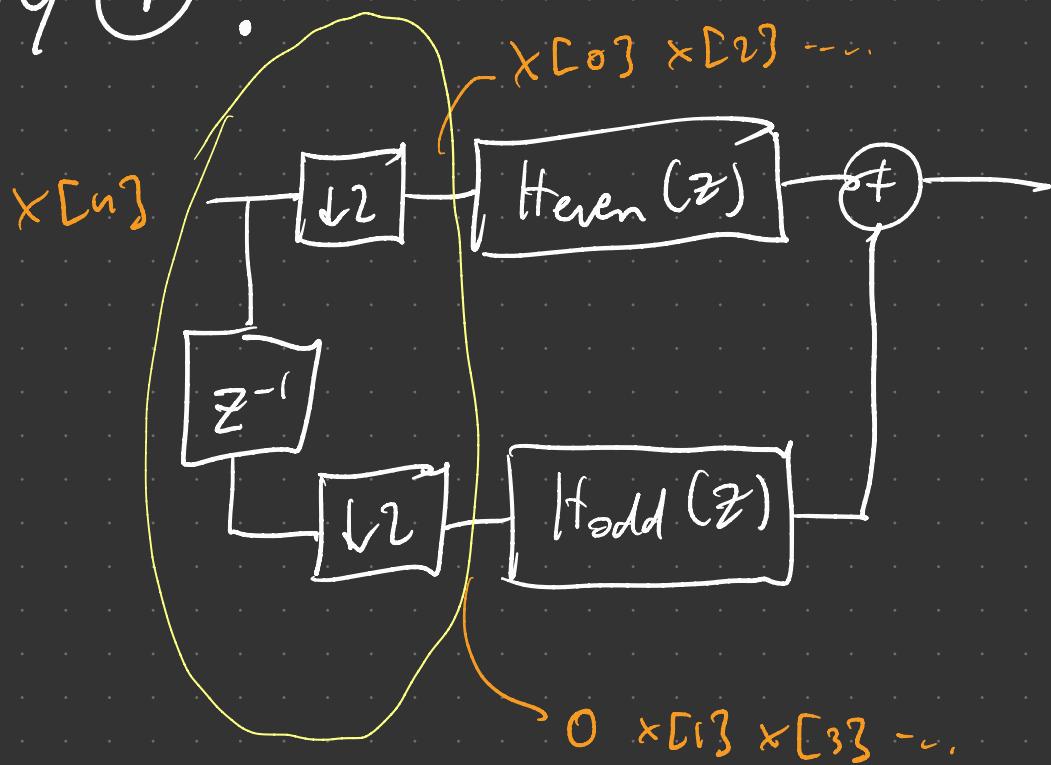


filtering
at high rate
= slow

Since downsampling is linear:



By ① :



filtering at
low rate
= fast

Series-to-parallel
buffer

Q: Why is this useful?

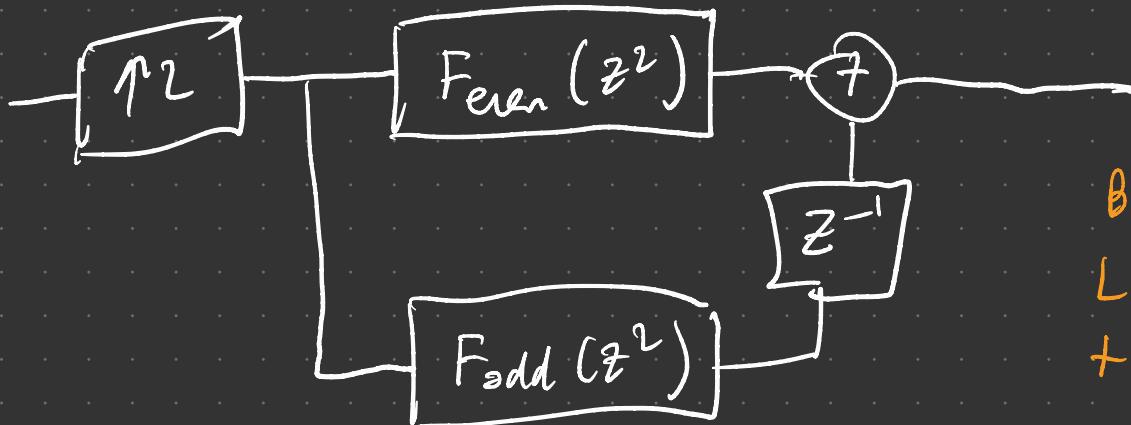
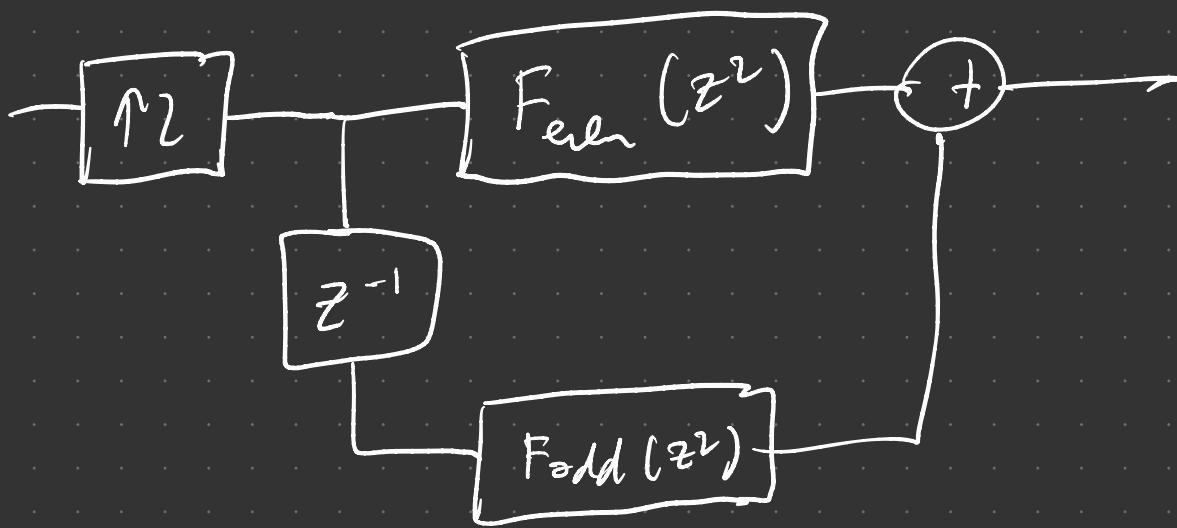
Remark: This is how
every multirate
system is implemented.

Exercise :



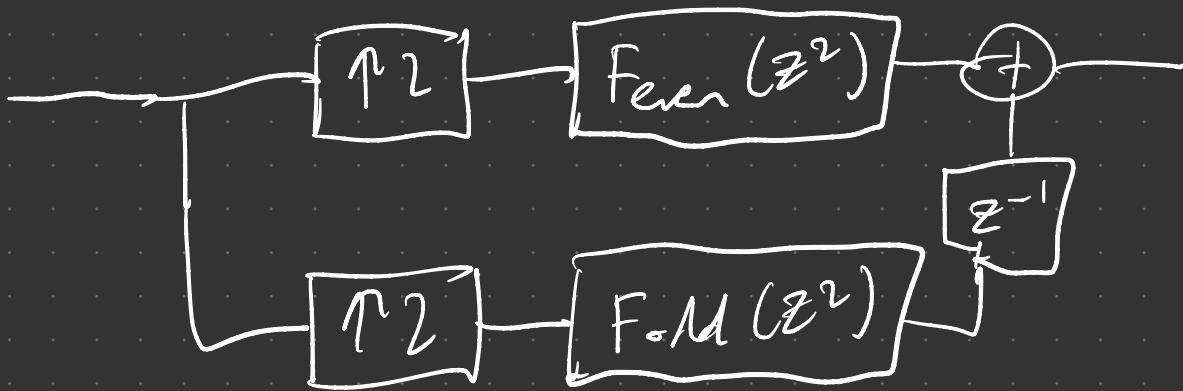
Slow

$$F(z) = F_{\text{even}}(z^2) + z^{-1} F_{\text{odd}}(z^2)$$

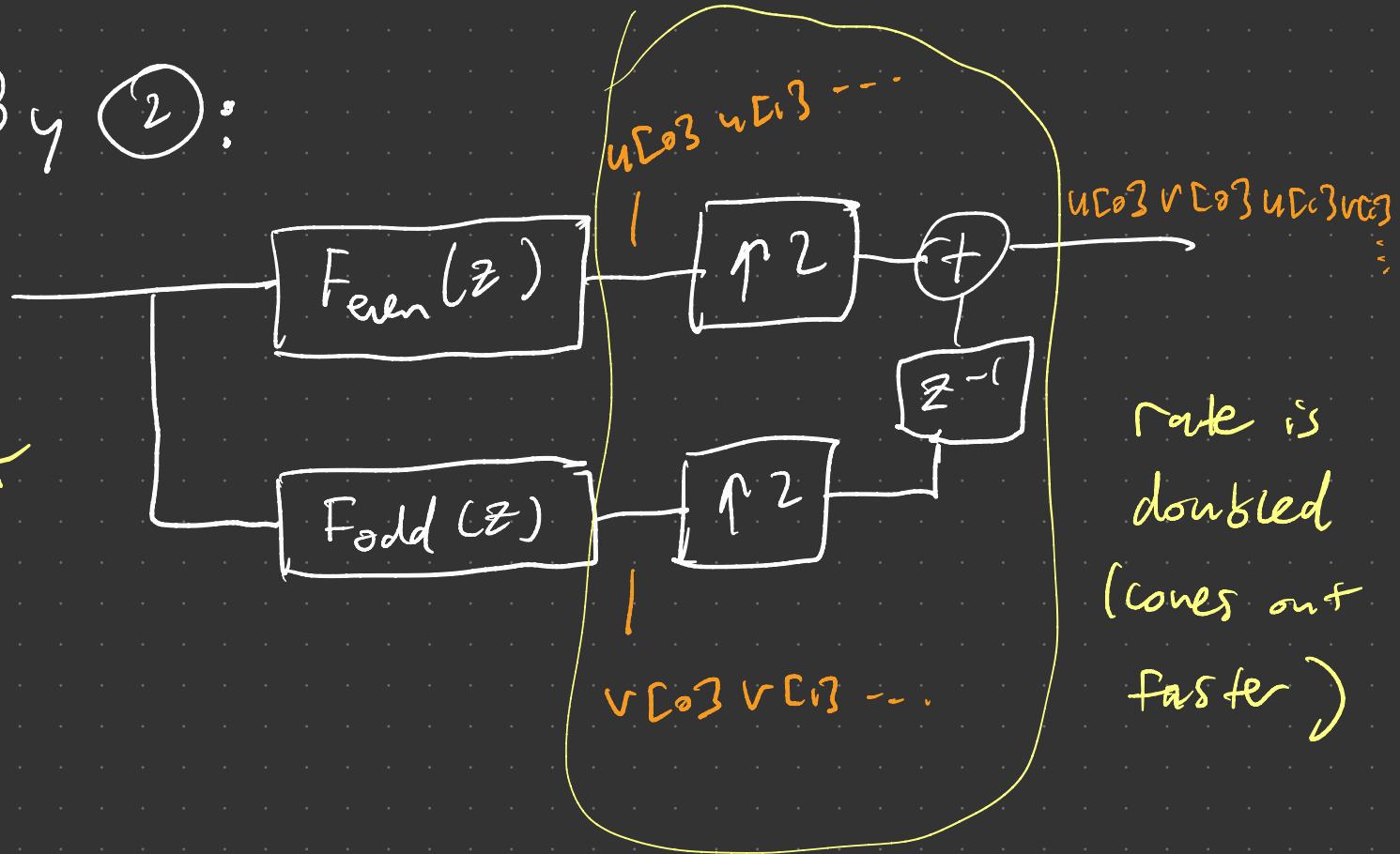


Both are
LTI so
they commute.

Since upsampling is linear:



By (2):



Multiplexing
parallel-to-serial buffer

Remark: Polyphase representations efficiently implement multirate operations.

Obs: For FIR filters it's very clear.

Q: What about IIR filters?

Exercise:

$$H(z) = \frac{1}{1 - \frac{1}{2}z^{-1}}$$



What are the even and odd polyphases?

$$H(z) = \frac{1}{1 - \frac{1}{2}z^{-1}} \cdot \frac{1 + \frac{1}{2}z^{-1}}{1 + \frac{1}{2}z^{-1}}$$

$$= \frac{1 + \frac{1}{2}z^{-1}}{1 - \frac{1}{4}z^{-2}}$$

$$= \left(\frac{1}{1 - \frac{1}{4}z^{-2}} \right) + z^{-1} \left(\frac{\frac{1}{2}}{1 - \frac{1}{4}z^{-2}} \right)$$

even

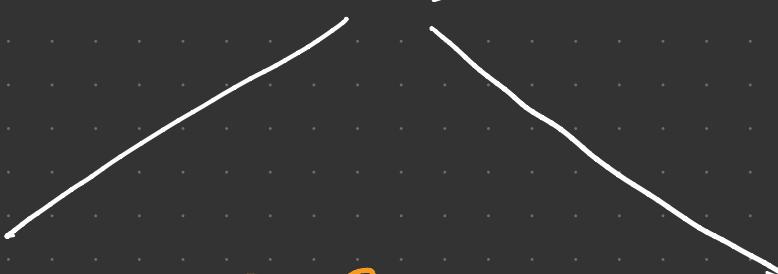
odd

$$H_{\text{even}}(z) = \frac{1}{1 - \frac{1}{4}z^{-1}}, \quad H_{\text{odd}}(z) = \frac{1}{1 + \frac{1}{4}z^{-1}}$$

Obs: The even/odd polyphases of an IIR filter are both IIR.

Time-Domain Characterization

$$h[n]$$



$$h_{\text{even}}[n] = h[2n]$$

$$h_{\text{odd}}[n] = h[2n+1]$$

Ex: $H(z) = \frac{1}{1 - \frac{1}{2}z^{-1}}$

$$h[n] = \left(\frac{1}{2}\right)^n u[n]$$

$$H(z) = \frac{1}{1 - az^{-1}}$$

$$h[n] = a^n u[n]$$

unit step



$$h_{\text{even}}[n] = \left(\frac{1}{2}\right)^{2n} u[2n] = \left(\frac{1}{4}\right)^n u[n]$$

$$H_{\text{even}}(z) = \frac{1}{1 - \frac{1}{4}z^{-1}}$$

$$h_{\text{odd}}[n] = \left(\frac{1}{2}\right)^{2n+1} u[2n+1] = \frac{1}{2} \left(\frac{1}{4}\right)^n u[n]$$

$$H_{\text{odd}}(z) = \frac{\frac{1}{2}}{1 - \frac{1}{4}z^{-1}}$$