

# CSE 551 - Foundations of Algorithms

## Assignment 2

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### Solution 1:

For time complexity in the big-O notation i.e. the worst case scenario,

Let us assume that all  $m$  hospitals made proposals to all students and the last students accepted the proposal for the given hospital. This means each hospital made  $n$  offers, so the number of total offers made =  $nm$ .

So,

```
while hospital  $h_i$  has positions -----  $n * m$ 
   $h_i$  offers to student  $s_j$  ----- 1
  //assuming that we are not using any data structure to check students' availability, it will take
   $O(n)$  time to check the availability.
  if  $s_j$  is free then ----- 1
     $s_j$  accepts the offer ----- 1
  else ( $s_j$  is already committed to a hospital  $h_k$ ) ----- 1
    // If a student is committed to a hospital we have to traverse the list to get the preference.
    if ( $s_j$  prefers  $h_k > h_i$ ) then -----
       $s_j$  remains committed to  $h_k$  -----
    else ( $s_j$  becomes committed to  $h_i$ ) -----
      positions at  $h_k + 1$  -----
      positions at  $h_j - 1$  -----
```

}  $n$

$$\begin{aligned}\text{So } T(m,n) &= c_1(n*m) * n((c_2(1+1+1+1) + c_3(n))) \\ &= c_1(n*m) * n(c_{23}(4 + n)) \\ &= c_1 * c_{23}(n*m*n*(4 + n)) \\ &= c_{123}(n^3m + 4n^2m) \\ \Rightarrow T(m,n) &= O(n^3m)\end{aligned}$$

So  $O(n^3m)$  should be the time complexity.

Assumption 1: The data structure is used to maintain the availability of students which reduces the time to check the availability from  $O(n)$  to  $O(1)$

$$\begin{aligned}\text{So } T(m,n) &= c_1(n*m) * (c_2(1+1+1+1) + c_3(n)) \\ &= c_1(n*m) * (c_{23}(4 + n)) \\ &= c_1 * c_{23}(n*m * (4 + n)) \\ &= c_{123}(n^2m + 4nm) \\ &\Rightarrow T(m,n) = O(n^2m)\end{aligned}$$

So  $O(n^2m)$  should be the time complexity.

Assumption 2: If we assume that data structure was used to store the preferences of students in such a way that it takes constant time complexity to check the preference. For this, we will create inverses of hospital preference for students which will take  $O(m)$  time for each student and so  $O(nm)$  for all students.

$$\begin{aligned}\text{So } T(m,n) &= c_0(nm) + c_1(n*m) * (c_2(1+1+1+1) + c_3(1)) \\ &= c_0(nm) + c_1(n*m) (c_{23}(5)) \\ &= c_0(nm) + c_1 * c_{23}(n*m*5) \\ &= c_{023}(6nm) \\ &\Rightarrow T(m,n) = O(nm)\end{aligned}$$

So  $O(nm)$  should be the time complexity.

## Solution 2:

To prove the claim that the student-hospital match in the algorithm is stable, we need to prove that there are no unstable pairs i.e. we will prove by contradiction.

Before moving further we need to note these two observations:

1. Hospitals propose to Students in decreasing order of their preference.
2. Once a student is matched, he/she never becomes unmatched; only “trade-up”.

Let  $S$  be the current Gale-Shapely matching. Consider a hospital  $h_i$  and a student  $s_j$ , such that pair  $(h_i - s_j)$  is an unstable pair in  $S$ .

Then, there exists another Gale-Shapely stable matching  $S^*$ , where both  $h_i$  and  $s_j$  are matched to their more preferable stable partner.

- Case 1:  $h_i$  never proposed to  $s_j$

As we know according to observation 1, hospitals propose to students in decreasing order of their preference. So, If  $(h_i - s_j)$  is unstable pair then it means hospital  $h_i$  prefers its  $S^*$  matching to student  $s_j$  and it shouldn't have proposed to student  $s_j$ . If this was the case then there was no reason for hospital  $h_i$  to leave its partner in  $S^*$ .

Therefore,

- Hospital  $h_i$  prefers student  $s_j$  over its partner in  $S^*$ .
- $(h_i - s_j)$  cannot be unstable pair.

- Case 2:  $h_i$  proposed to  $s_j$

As we know according to observation 2, students only “trade-up”. So, this means hospital  $h_i$  proposed to student  $s_j$ , but was rejected either right away or later. Student  $s_j$  prefers its  $S^*$  partner to  $h_i$ . But if this was the case, why will student  $s_j$  leave  $S^*$  partner to pair up with hospital  $h_i$  (because the student only trades up).

Therefore,

- $(h_i - s_j)$  cannot be unstable pair.

In either case,  $(h_i - s_j)$  is stable pair. Thus the student-hospital match in the algorithm is stable.

### Solution 3:

Let the 6 students be  $s_1, s_2, s_3, s_4, s_5, s_6$  and 2 hospitals be  $h_1$  and  $h_2$ .

Given,  $h_1$  can take 2 students and  $h_2$  can take 3 students.

Let the preference list for hospitals and students as shown below:

<i>Hospital</i>	0th	1st	2nd	3rd	4th	5th
<i><b>h1</b></i>	$s_3$	$s_1$	$s_2$	$s_6$	$s_4$	$s_5$
<i><b>h2</b></i>	$s_5$	$s_2$	$s_3$	$s_4$	$s_1$	$s_6$

<i>Student</i>	0th	1st
<i><b>s1</b></i>	$h_1$	$h_2$
<i><b>s2</b></i>	$h_1$	$h_2$
<i><b>s3</b></i>	$h_2$	$h_1$

<i>s4</i>	h1	h2
<i>s5</i>	h2	h1
<i>s6</i>	h1	h2

***Manually running the algorithm:***

Hospital h1 will propose to student s3. s3 is free and will accept the offer.

<i>Hospital</i>	0th	1st	2nd	3rd	4th	5th
<i>h1</i>	s3	s1	s2	s6	s4	s5
<i>h2</i>	s5	s2	s3	s4	s1	s6

<i>Student</i>	0th	1st
<i>s1</i>	h1	h2
<i>s2</i>	h1	h2
<i>s3</i>	h2	h1
<i>s4</i>	h1	h2
<i>s5</i>	h2	h1
<i>s6</i>	h1	h2

Hospital h1 will propose to student s1. s1 is free and will accept the offer.

<i>Hospital</i>	0th	1st	2nd	3rd	4th	5th
<i>h1</i>	s3	s1	s2	s6	s4	s5
<i>h2</i>	s5	s2	s3	s4	s1	s6

<i>Student</i>	0th	1st
<i>s1</i>	h1	h2
<i>s2</i>	h1	h2
<i>s3</i>	h2	h1
<i>s4</i>	h1	h2
<i>s5</i>	h2	h1
<i>s6</i>	h1	h2

Now, the requirement for hospital h1 is completed and h2 will start proposing. h2 will propose to s5 and s5 being free will accept the offer.

<i>Hospital</i>	0th	1st	2nd	3rd	4th	5th
<i>h1</i>	s3	s1	s2	s6	s4	s5
<i>h2</i>	s5	s2	s3	s4	s1	s6

<i>Student</i>	0th	1st
<i>s1</i>	h1	h2
<i>s2</i>	h1	h2
<i>s3</i>	h2	h1
<i>s4</i>	h1	h2
<i>s5</i>	h2	h1
<i>s6</i>	h1	h2

Now, h2 will propose to s2 now, s2 being free will accept the offer.

<i>Hospital</i>	0th	1st	2nd	3rd	4th	5th
<i>h1</i>	s3	s1	s2	s6	s4	s5
<i>h2</i>	s5	s2	s3	s4	s1	s6

<i>Student</i>	0th	1st
<i>s1</i>	h1	h2
<i>s2</i>	h1	h2
<i>s3</i>	h2	h1
<i>s4</i>	h1	h2
<i>s5</i>	h2	h1
<i>s6</i>	h1	h2

Now, h2 will propose to s3 now, s3 isn't free but prefers h2 over h1 so will accept the offer from h2.

<i>Hospital</i>	0th	1st	2nd	3rd	4th	5th
<i>h1</i>	s3	s1	s2	s6	s4	s5
<i>h2</i>	s5	s2	s3	s4	s1	s6

<i>Student</i>	0th	1st
<i>s1</i>	h1	h2
<i>s2</i>	h1	h2
<i>s3</i>	h2	h1
<i>s4</i>	h1	h2
<i>s5</i>	h2	h1
<i>s6</i>	h1	h2

Now h1 is short by 1 so, it will propose to next preference, s2, and s2 is already committed but will accept the offer from h1 as it prefers h1 over h2.

<i>Hospital</i>	0th	1st	2nd	3rd	4th	5th
<i>h1</i>	s3	s1	s2	s6	s4	s5
<i>h2</i>	s5	s2	s3	s4	s1	s6

<i>Student</i>	0th	1st
<i>s1</i>	h1	h2
<i>s2</i>	h1	h2
<i>s3</i>	h2	h1
<i>s4</i>	h1	h2
<i>s5</i>	h2	h1
<i>s6</i>	h1	h2

Now h2 is short by 1 so, it will propose to next preference, s4, and s4 is free and will accept the offer from h2.

<i>Hospital</i>	0th	1st	2nd	3rd	4th	5th
<i>h1</i>	s3	s1	s2	s6	s4	s5
<i>h2</i>	s5	s2	s3	s4	s1	s6

<i>Student</i>	0th	1st
<i>s1</i>	h1	h2
<i>s2</i>	h1	h2
<i>s3</i>	h2	h1
<i>s4</i>	h1	h2
<i>s5</i>	h2	h1
<i>s6</i>	h1	h2

Output generated by the algorithm is

- Hospital h1 = (s1, s2)
- Hospital h2 = (s5, s3, s4)

#### Solution 4:

**IDE used = Visual Studio Code**

Java Code: Implementing the algorithm

```
import java.util.HashSet;
import java.util.LinkedList;
import java.util.Queue;
import java.util.Scanner;
import java.util.Set;

public class GS_StudentHospital {
    public static void main(String[] args){
        System.out.println("*****");
        System.out.println("Assignment 2: Q4");
        Scanner sc = new Scanner(System.in);
        System.out.println("Enter the number of students");
        int students = sc.nextInt();
```

```

        System.out.println("Enter the number of hospitals");
        int hospitals = sc.nextInt();
        int[][] stuPref = new int[students][hospitals];
        for(int i = 0; i<students; i++){
            System.out.println("Enter the preference of hospitals for student "+
(i+1));
            for(int j=0; j<hospitals; j++)
                stuPref[i][j] = sc.nextInt() - 1;
        }
        int[][] hosPref = new int[hospitals][students];
        for(int i = 0; i<hospitals; i++){
            System.out.println("Enter the preference of students for hospital "+
(i+1));
            for(int j=0; j<students; j++)
                hosPref[i][j] = sc.nextInt() - 1;
        }
        int[] req = new int[hospitals];
        for(int i=0; i<hospitals; i++){
            System.out.println("Enter the requirement for hospital "+ (i+1));
            req[i] = sc.nextInt();
        }
        Set<Integer>[] match = getMatchings(stuPref, hosPref, req);
        System.out.println("Stable matchings are as follows");
        for(int i=0; i<hospitals; i++){
            System.out.print("h" + (i+1)+": ");
            for(int mat : match[i])
                System.out.print("s" + (mat+1)+ " ");
            System.out.println();
        }
        System.out.println("*****");
    }

    public static Set<Integer>[] getMatchings(int[][] stuPref, int[][] hosPref, int[]
req){
        int students = stuPref.length;
        int hospitals = hosPref.length;
        //prefNumHospAt to keep track of proposals made by hospitals
        int[] prefNumHospAt = new int[hospitals];
        //studMatch to keep track of which student is committed to which hospital
        int[] studMatch = new int[students];
        //keeping track of unmatched hospitals and students
        Queue<Integer> freeHosp = new LinkedList<>();
        Set<Integer> freeStud = new HashSet<>();
        for(int i=0; i< hospitals; i++)

```



```

        freeHosp.offer(i);
    }
    for(int i=0; i< students; i++)
        freeStud.add(i);

    Set<Integer>[] match = new HashSet[hospitals];
    inversePref(stuPref);

    while(!freeHosp.isEmpty()){
        int currentHosp = freeHosp.peek();
        int studNum = prefNumHospAt[currentHosp];
        int proposeStu = hosPref[currentHosp][studNum];
        if(studNum==0)
            match[currentHosp] = new HashSet<>();
        prefNumHospAt[currentHosp]++;

        if(freeStud.contains(proposeStu)){
            req[currentHosp]--;
            freeStud.remove(proposeStu);
            match[currentHosp].add(proposeStu);
            studMatch[proposeStu] = currentHosp;
        }
        else{
            int previousHosp = studMatch[proposeStu];
            if(stuPref[proposeStu][previousHosp] >
stuPref[proposeStu][currentHosp]){
                req[previousHosp]++;
                match[previousHosp].remove(proposeStu);
                if(req[previousHosp]==1)
                    freeHosp.offer(previousHosp);
                req[currentHosp]--;
                match[currentHosp].add(proposeStu);
                studMatch[proposeStu] = currentHosp;
            }
        }
        if(req[currentHosp]==0)
            freeHosp.poll();
    }
    return match;
}

public static void inversePref(int[][] stuPref){
    //inverse the student preferences for constant time access.
    int hosp = stuPref[0].length;
    int[] helper = new int[hosp];
    for(int i=0; i<stuPref.length; i++){

```

```

        for(int j = 0; j<hosp; j++)
            helper[j] = stuPref[i][j];
        for(int j = 0; j<hosp; j++)
            stuPref[i][helper[j]] = j;
    }
}
}

```

### Solution 5:

The output from manually running the algorithm matches with the output generated by the program.

**Stable matchings are as follows**

**h1: s1 s2**

**h2: s3 s4 s5**

```

spats@Sannidhyas-MacBook-Air GS_StudentHospital % /usr/bin/env /L
sInExceptionMessages -cp /Users/spats/Library/Application\ Support
tHospital_9f52a38e/bin GS_StudentHospital
*****
Assignment 2: Q4
Enter the number of students
6
Enter the number of hospitals
2
Enter the preference of hospitals for student 1
1 2
Enter the preference of hospitals for student 2
1 2
Enter the preference of hospitals for student 3
2 1
Enter the preference of hospitals for student 4
1 2
Enter the preference of hospitals for student 5
2 1
Enter the preference of hospitals for student 6
1 2
Enter the preference of students for hospital 1
3 1 2 6 4 5
Enter the preference of students for hospital 2
5 2 3 4 1 6
Enter the requirement for hospital 1
2
Enter the requirement for hospital 2
3
Stable matchings are as follows
h1: s1 s2
h2: s3 s4 s5
*****
spats@Sannidhyas-MacBook-Air GS_StudentHospital % █

```