

Samuel Steinberg

CS 465

January 31st, 2020

Homework #3

1. a) This will produce a relation of hotel ID numbers (hotelNo) where the price for a room from the Room table is greater than 50 pounds. Essentially this produce a relation of hotelNo's with a room price greater than 50 pounds.
c) This will produce a natural join of Hotel and Room with a price of a room being greater than 50 pounds. This will produce a relation containing all hotel names with a room price greater than 50 pounds.
f) This will produce a join of Booking and Guest and will essentially produce a relation containing the names of guests and hotel ID's for all tuples of guests with hotels booked in London.

2. b. Relational Algebra: $\Pi_{\text{roomNo}}(\sigma_{\text{price} < 20 \wedge \sigma_{\text{type}} = \text{'single'}})$
Tuple Relational Calculus: $\{R \mid \text{Room}(R) \wedge R.\text{price} < 20 \wedge R.\text{type} = \text{'single'}\}$

- c. Relational Algebra: $\Pi_{\text{guestName}, \text{guestAddress}}(\text{Guest})$
Tuple Relational Calculus: $\{G.\text{guestName}, G.\text{guestAddress} \mid \text{Guest}(G)\}$

- d. Relational Algebra:
 $\Pi_{\text{price}, \text{type}}(\text{Room} \bowtie_{\text{Room.hotelNo}=\text{Hotel.hotelNo}} (\sigma_{\text{hotelName}=\text{'GrosvenorHotel'}}(\text{Hotel})))$

Tuple Relational Calculus:
 $\{R.\text{price}, R.\text{type} \mid \text{Room}(R) \wedge (\exists H)(\text{Hotel } H) \wedge (R.\text{hotelNo} = H.\text{hotelNo}) \wedge (H.\text{hotelName} = \text{'GrosvenorHotel'})\}$

- e. Relational Algebra:
 $\text{Guest} \bowtie_{\text{Guest.guestNo}=\text{Booking.guestNo}} (\sigma_{\text{dateFrom} \leq \text{CURDATE}() \wedge \sigma_{\text{dateTo} > \text{CURDATE}()} (\text{Booking} \bowtie_{\text{Booking.hotelNo}=\text{Hotel.hotelNo}} (\sigma_{\text{hotelName} = \text{'Grosvenor Hotel'}}(\text{Hotel}))))$

Tuple Relational Calculus:
 $\{G \mid \text{Guest}(G) ((\exists B)(\exists H) (\text{Booking}(B) \wedge \text{Hotel}(H) \wedge (B.\text{dateFrom} \leq \text{CURDATE}()) \wedge (B.\text{dateTo} > \text{CURDATE}()) \wedge (G.\text{guestNo} = B.\text{guestNo}) \wedge (B.\text{hotelNo} = H.\text{hotelNo}) \wedge (H.\text{hotelName} = \text{'Grosvenor Hotel'}))))\}$

f. Relational Algebra:

$$\begin{aligned} & \text{Room} \bowtie (\sigma_{\text{hotelName} = \text{'Grosvenor Hotel'}}(\text{Hotel})) \bowtie \Pi_{\text{guestName}, \text{hotelNo}} \\ & \quad (\text{Guest} \bowtie_{\text{Guest.guestNo}=\text{Booking.guestNo}} (\sigma_{\text{dateFrom} \leq \text{CURDATE}()}) \\ & \wedge \sigma_{\text{dateTo} > \text{CURDATE}()}) (\text{Booking} \bowtie_{\text{Booking.hotelNo}=\text{Hotel.hotelNo}} (\sigma_{\text{hotelName}=\text{'Grosvenor Hotel'}}(\text{Hotel}))) \end{aligned}$$

Note to TA: \bowtie was the closest symbol to the textbooks LEFT OUTER JOIN as I could find.

Tuple Relational Calculus:

$$\begin{aligned} & \{R, G, \text{guestName} \mid \text{Room}(R) ((\exists H) (\text{Hotel}(H) \wedge (H.\text{hotelNo} = R.\text{hotelNo}) \\ & \quad \wedge (H.\text{hotelName} = \text{'Grosvenor Hotel'}))) \vee (\text{Guest}(G) ((\exists B)(\exists H)(\text{Booking}(B) \\ & \quad \wedge \text{Hotel}(H) \wedge (H.\text{hotelNo} \\ & \quad = B.\text{hotelNo}) \wedge (H.\text{hotelName} = \text{'Grosvenor Hotel'}) \\ & \quad \wedge (G.\text{guestNo} = B.\text{guestNo}) \wedge (B.\text{dateFrom} \leq \text{CURDATE}() \wedge B.\text{dateTo} \\ & \quad > \text{CURDATE}())) \} \end{aligned}$$

3. a. This would produce a relation containing all hotel names in the city of London.
- b. This would produce a relation containing all hotel names with room prices greater than 50 pounds.
- c. This would produce a relation containing all of the hotel names where guest John Smith has stayed. This would essentially return all hotel names where guest John Smith has had bookings.
- d. This would produce a relation containing the hotel name and guest name where there exist two bookings for the guest at the same hotel.

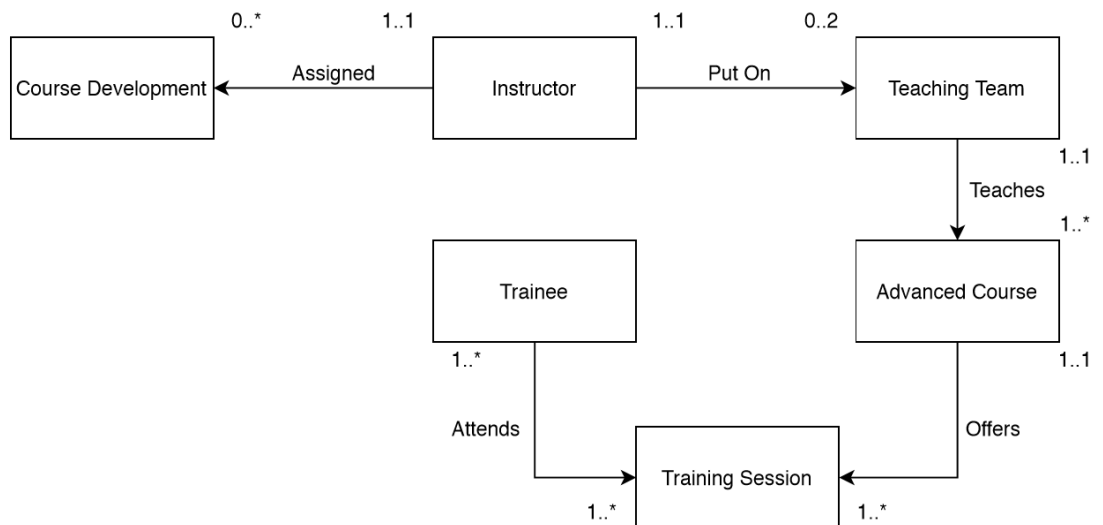
4.

- a. Main Entity Types:
 - i. Instructor
 - ii. Trainee
 - iii. Advanced Technology Course
 - iv. Training Session
 - v. Course Development
 - vi. Teaching Team
- b.

Relationship Types and Multiplicity

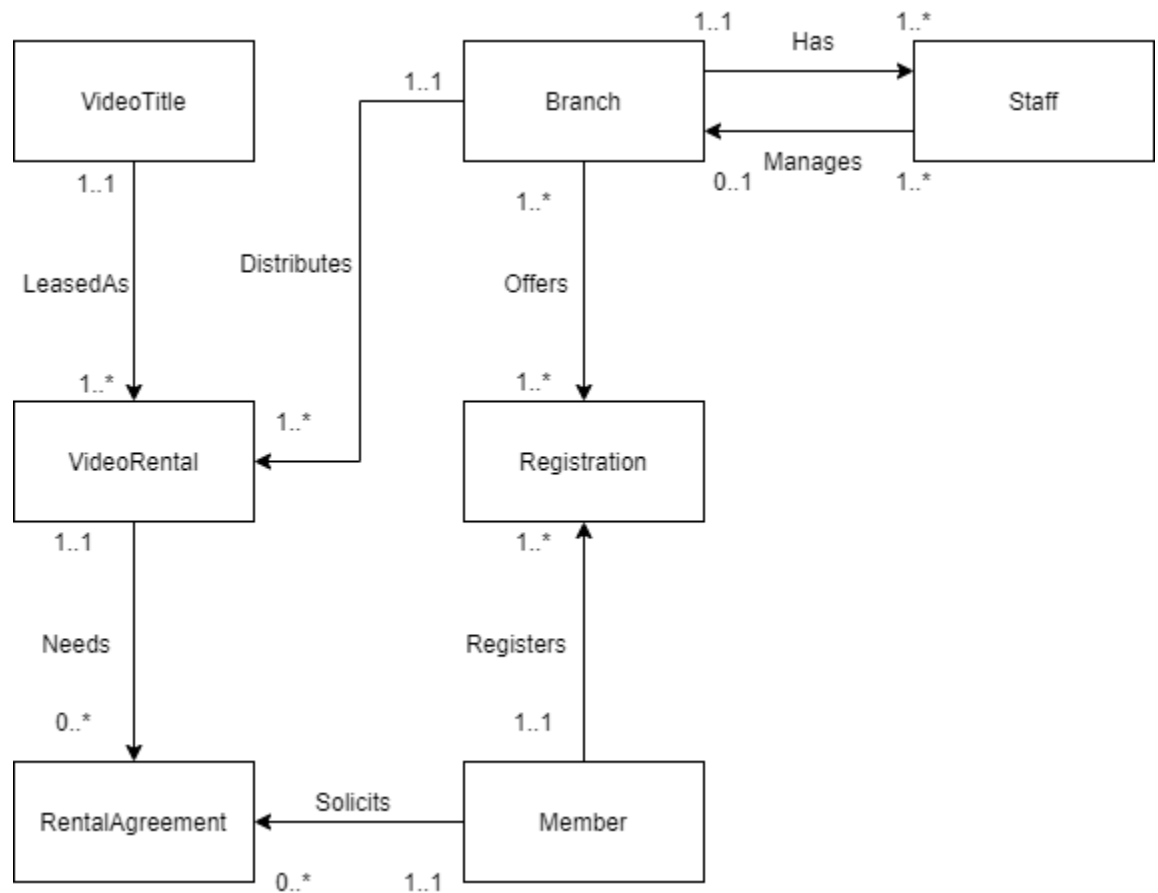
LHS Multiplicity	LHS Entity	Relationship	RHS Entity	RHS Multiplicity
1..*	Trainee	Attends	Training Session	1..*
1..1	Instructor	Put On	Teaching Team	0..2
1..1	Instructor	Assigned	Course Development	0..*
1..1	Advanced Course	Offers	Training Session	1..*
1..1	Teaching Team	Teaches	Advanced Course	1..*

- c. (Made on draw.io)



5.

a. (Made on draw.io)



b. **Red denotes primary keys**

- Branch (**branchNo**, address, phoneNumber, manager)
- Staff (**staffNo**, name, position, salary)
 - Name contains first name and last name
- VideoTitle (**catalogNo**, title, category, dailyRentFee, purchaseCost, actorNames, directorName)
- VideoRental (**videoNum**, category, rentalStatus, catalogNo, title, dailyRentFee, purchaseCost, actorNames, directorName)
- Member (**memberNo**, firstName, lastName, address, registrationDate, branchRegistered)
- Registration (memberNo, branchNo)
- RentalAgreement(**rentalNo**, firstName, lastName, memberNo, videoTitle, dailyRental, dateFrom, dateTo)

c. Primary Keys:

- Branch Entity
 - Primary Key: branchNo
 - Candidate Key: branchNo, address
- Staff Entity

- Primary Key: staffNo
- Candidate Key: staffNo, name
- iii. VideoTitle
 - Primary Key: catalogNo
 - Candidate Key: catalogNo, {title, category}
- iv. VideoRental
 - Primary Key: videoNum
 - Candidate Key: videoNum, {catalogNo, title}
- v. Member
 - Primary Key: memberNo
 - Candidate Key: memberNo, {firstName, lastName}
- vi. Registration
 - Primary Key: None, this is a weak entity type.
 - Candidate Key {memberNo, branchNo}
- vii. RentalAgreement
 - Primary Key: rentalNo
 - Candidate Key: rentalNo, {firstName, lastName, memberNo}