

Explanation in Ontology Languages

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ESSLLI 2009, 27-31 July, Bordeaux

With thanks to Matthew Horridge and Johannes Bauer



Outline

Mon Introduction

Tue Computation of justifications

Wed Fine-grained justifications

Thu Lemmata

Fri Model based explanations

Monday Introduction

Schedule for today

- Ontologies and ontology languages
- Background on explanation
- Introduction to justifications and associated services

Ontologies and ontology languages

- Are logical theories
- Capture a domain of interest
- Contain explicit knowledge
- Implicit knowledge can be inferred
- This is possible through the formal semantics underlying ontology languages

Ontology languages

- Often based on Description Logics (DLs)
 - Class/object/role paradigm
 - DLs have clear, formal semantics
 - (Decidable) fragments of first order logic
 - Standard, implemented inference services

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Many services are "automagic" (e.g., classification)

Women are female persons.

Every mother is a female parent.

Every parent is a person.

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Woman \equiv Female \sqcap Person

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Parent

□ Person

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Parent

□ Person

Every parent has at least one child, each of which is a person.

Parent $\sqsubseteq \geqslant 1$ hasChild $\sqcap \forall$ hasChild.Person

Women are female persons.

Woman \equiv Female \sqcap Person

Every mother is a female parent.

Mother \sqsubseteq Female \sqcap Parent

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Parent

□ Person

implies:

Every mother is a woman.

Mother

☐ Woman

Every parent has at least one child, each of which is a person.

Parent $\sqsubseteq \geqslant 1$ hasChild $\sqcap \forall$ hasChild.Person



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DL examples

Parent $\sqsubseteq \geqslant 1$ hasChild $\sqcap \forall$ hasChild.Person



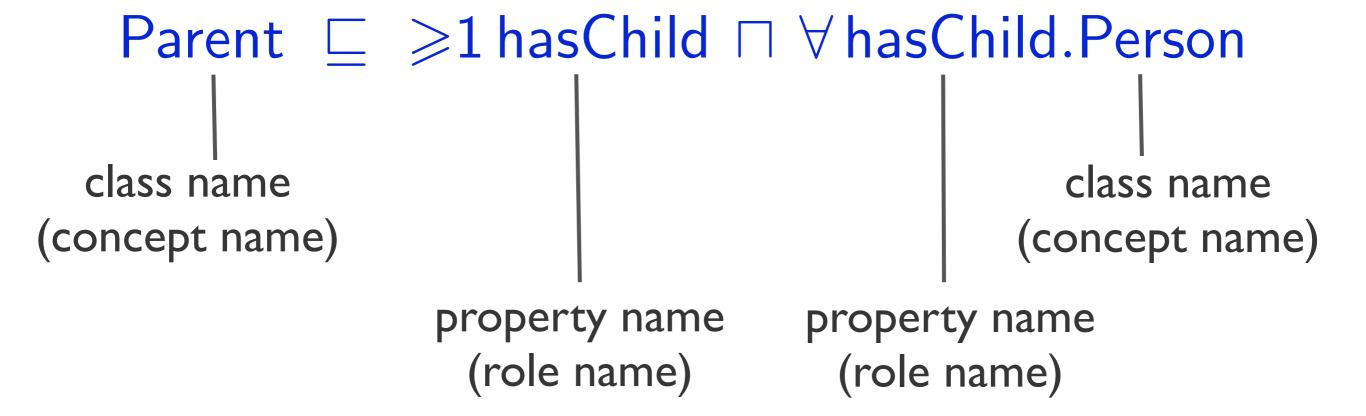
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DL examples

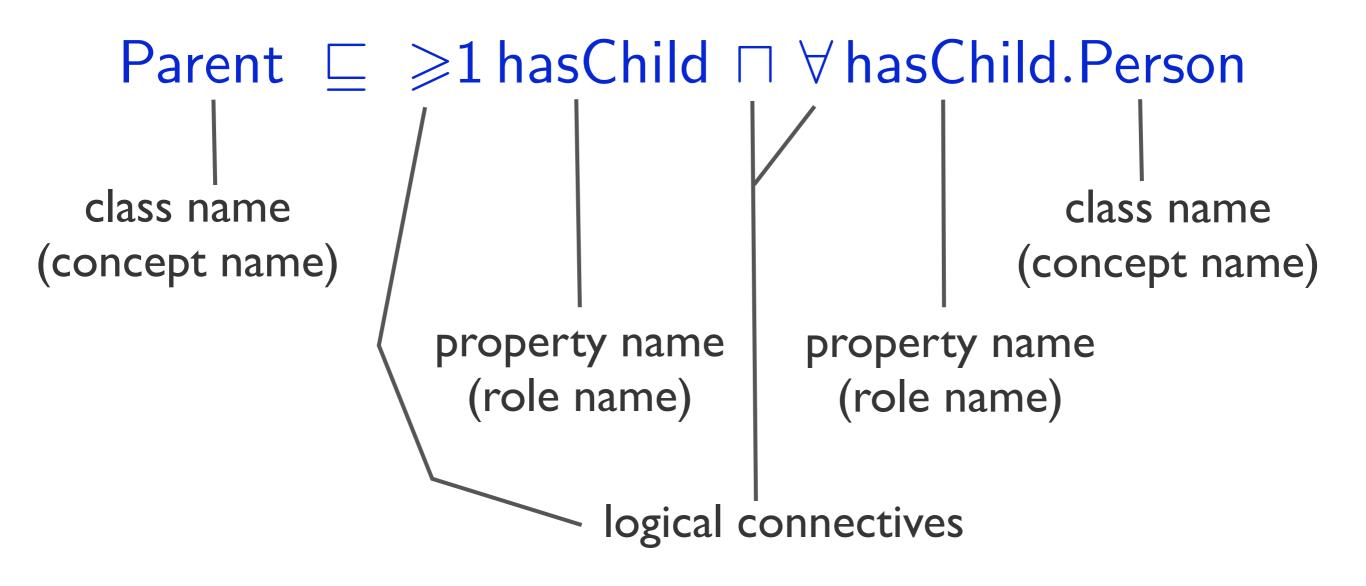
Parent $\sqsubseteq \geqslant 1$ has Child $\sqcap \forall$ has Child. Person



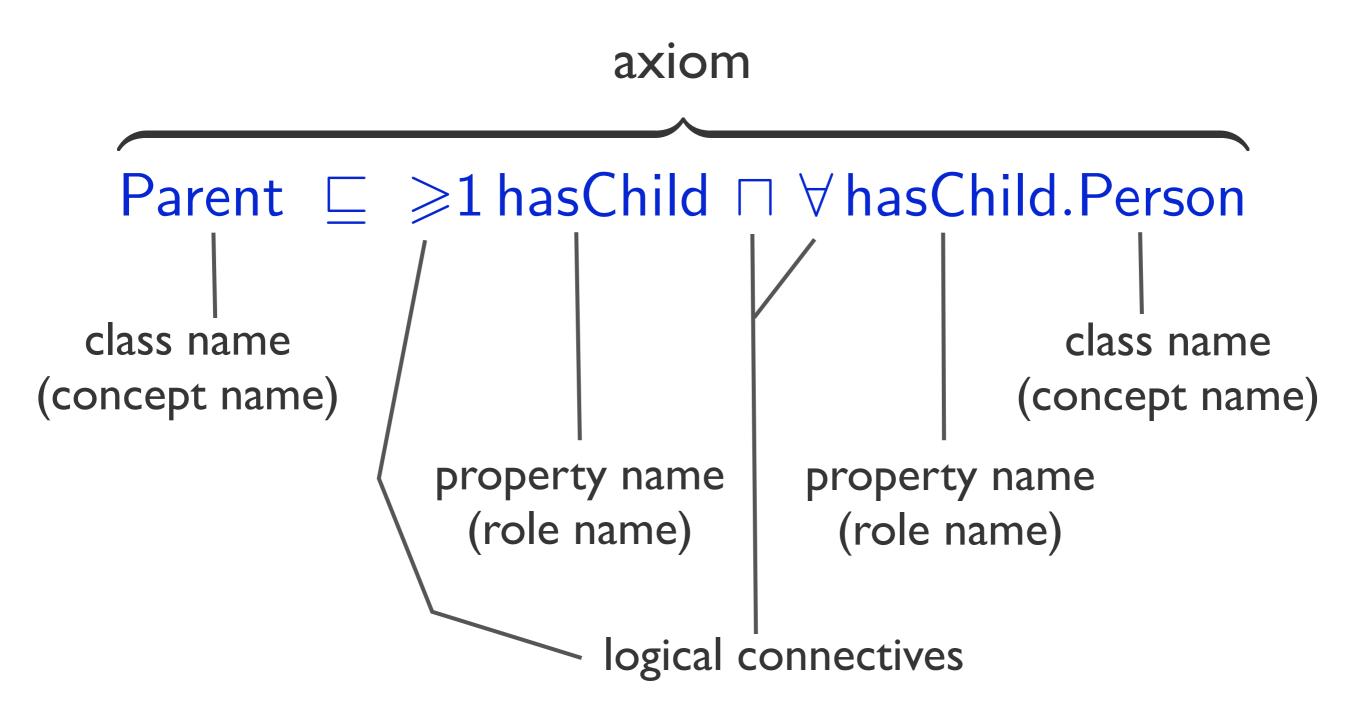
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DL syntax

• class names; property names:

$$A, B, \ldots; P, Q, \ldots$$

complex classes:

T,
$$A$$
, $\neg C$, $C \sqcap D$, $\exists R.C$, $\geqslant n R.C$

A atomic; C, D possibly complex; R atomic property or inverse (P^-) ; $n \in \mathbb{N}$

- convention: A, B atomic classes; C, D, \ldots arbitrary
- syntactic sugar:

$$\bot = \neg \top \qquad \forall R.C = \neg \exists R. \neg C$$

$$C \sqcup D = \neg C \sqcap \neg D \qquad \leqslant n R.C = \neg (\geqslant n+1 R.C)$$

DL syntax

 $C \sqsubseteq D$

subClassOf axiom

 $C \equiv D$

shortcut for $C \sqsubseteq D$ and $D \sqsubseteq C$

 $R_1 \circ \cdots \circ R_n \sqsubseteq S$ subPropertyChain axiom

+ more features, mostly syntactic sugar: disjointness, transitivity, reflexivity, ...

DL syntax vs. OVL Manchester syntax

DL

concept, role

Т, _

 $\neg C$

=nR.C

OWL Manchester syntax

class, property

Thing, Nothing

not C

 $C \sqcap D$, $C \sqcup D$ C and D, C or D

 $\exists R.C, \forall R.C$ R some C, R only C

 $\geqslant nR.C$, $\leqslant nR.C$ R min nC, R max nC

R exactly n C

DL semantics

- Interpretations $\mathcal{I} = (\Delta^{\mathcal{I}}, \cdot^{\mathcal{I}})$ $\Delta^{\mathcal{I}} \neq \emptyset$! with $A^{\mathcal{I}} \subseteq \Delta^{\mathcal{I}}$ for class names A and $P^{\mathcal{I}} \subseteq \Delta^{\mathcal{I}} \times \Delta^{\mathcal{I}}$ for property names P
- Inductive transfer to arbitrary props / classes

$$(P^{-})^{\mathcal{I}} = \{(x, y) \mid (y, x) \in P^{\mathcal{I}}\}$$

$$T^{\mathcal{I}} = \Delta^{\mathcal{I}}$$

$$(\neg C)^{\mathcal{I}} = \Delta^{\mathcal{I}} \setminus C^{\mathcal{I}}$$

$$(C \sqcap D)^{\mathcal{I}} = C^{\mathcal{I}} \cap D^{\mathcal{I}}$$

$$(\exists R.C)^{\mathcal{I}} = \{x \in \Delta^{\mathcal{I}} \mid \text{ for some } y \in C^{\mathcal{I}}, (x, y) \in R^{\mathcal{I}}\}$$

$$(\geqslant n R.C)^{\mathcal{I}} = \{x \in \Delta^{\mathcal{I}} \mid \#\{y \in C^{\mathcal{I}} \mid (x, y) \in R^{\mathcal{I}}\} \geqslant n\}$$

DL semantics

Interpretation \mathcal{I} satisfies ...

- axiom $A \sqsubseteq B$ if $A^{\mathcal{I}} \subseteq B^{\mathcal{I}}$
- axiom $R_1 \circ \cdots \circ R_n \sqsubseteq S$ if:

whenever $xR_1^{\mathcal{I}}z_1R_2^{\mathcal{I}}z_2R_3^{\mathcal{I}}\dots R_{n-1}^{\mathcal{I}}z_{n-1}R_n^{\mathcal{I}}y$, then $xS^{\mathcal{I}}y$

We write $\mathcal{I} \models \alpha$.

We write $\mathcal{I} \models \mathcal{O}$ if $\mathcal{I} \models \alpha$ for all $\alpha \in \mathcal{O}$.

DL semantics

- Interpretation \mathcal{I} is a model of ontology \mathcal{O} if \mathcal{I} satisfies every axiom in \mathcal{O} .
 - We write $\mathcal{I} \models \mathcal{O}$.
- Ontology \mathcal{O} is inconsistent if it has no model.
- Class C is satisfiable if $C^{\mathcal{I}} \neq \emptyset$ for every interpretation \mathcal{I} .
- Axiom α is entailed by \mathcal{O} $(\mathcal{O} \models \alpha)$ if for all \mathcal{I} : $\mathcal{I} \models \mathcal{O} \Rightarrow \mathcal{I} \models \alpha$



(Bio-)Medicine

Chemistry

Astronautics

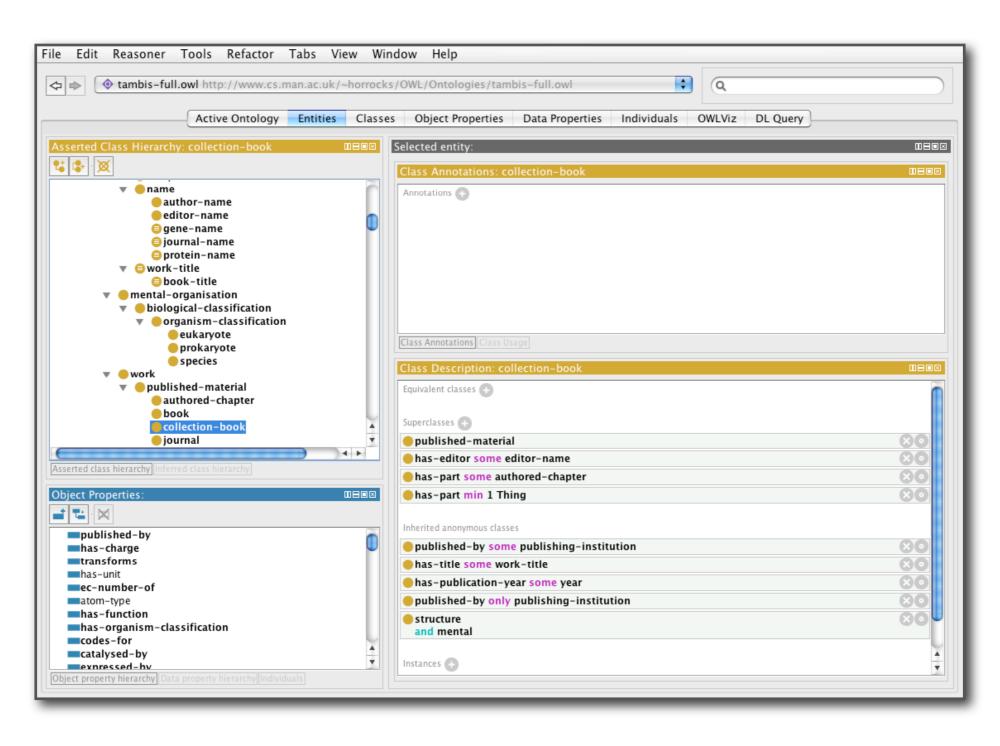
... are used in these fields and others

Linguistics

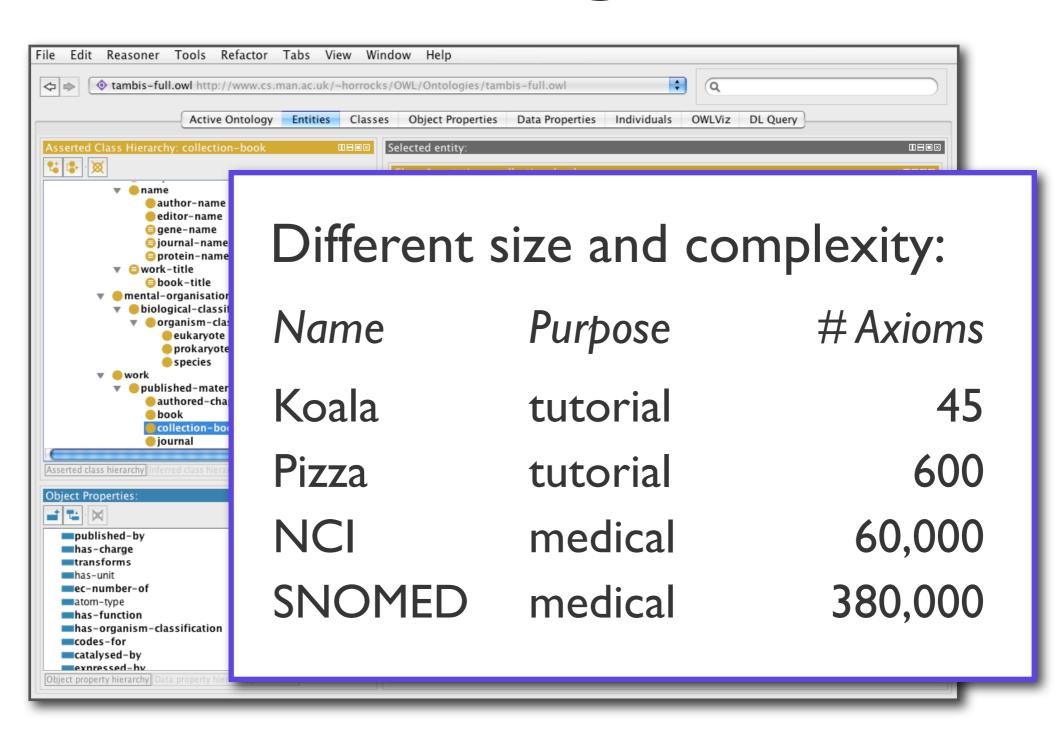
Description of business processes

(Web) Service description

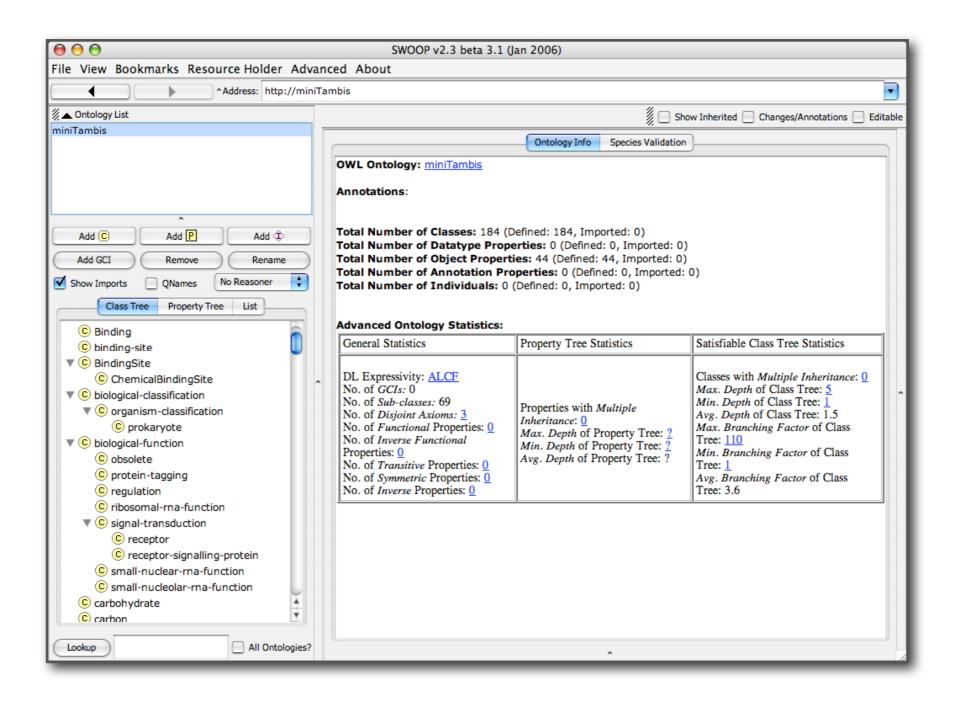




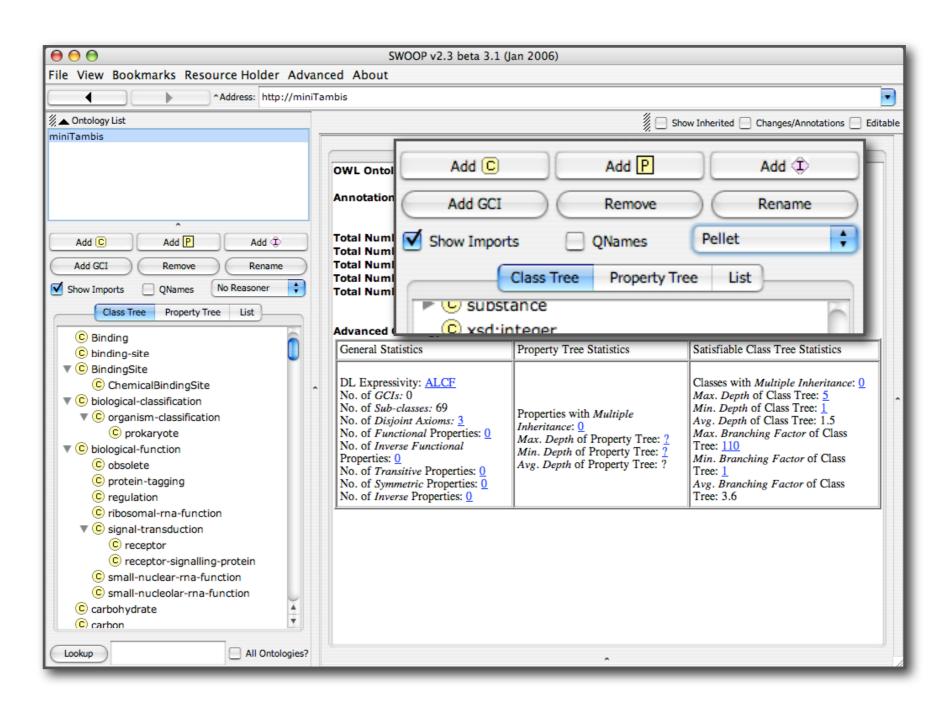




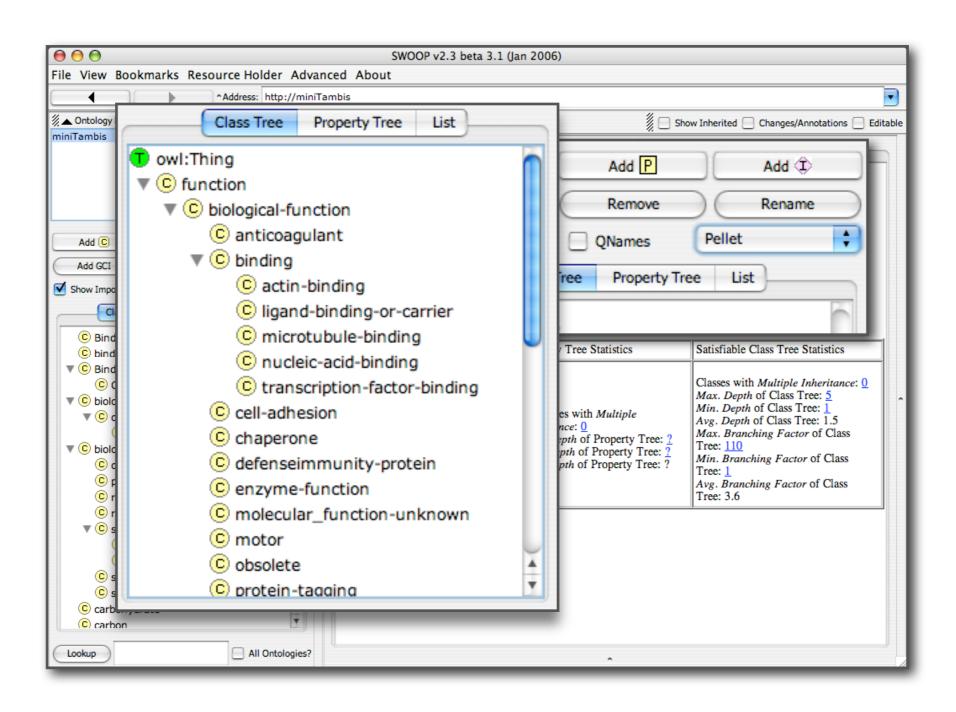




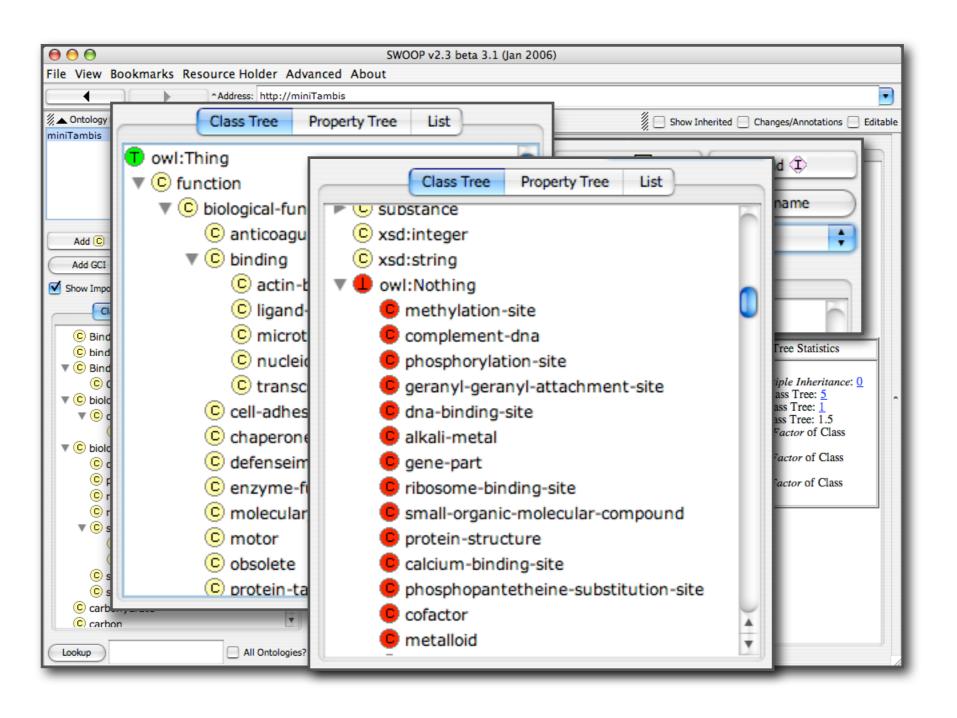












"Large" User Base

- Large and Growing
 - (126,870 registered Protégé Users)
 - (Standardization helps)
 - Bio-medical applications huge
- Useful for research!
 - Diverse users, tasks, issues
 - High impact



Reasoning services

Name

Consistency checking

Subsumption checking

Class hierarchy computation

Unsatisfiability test

Instance checking

Question

$$\mathcal{O} \models \top \sqsubseteq \bot ?$$

$$\mathcal{O} \models C \sqsubseteq D$$
?

$$\mathcal{O} \models A \sqsubseteq B$$
?

$$\mathcal{O} \models C \sqsubseteq \bot$$
?

$$\mathcal{O} \models a : C ?$$

Reasoning services ...

• ... are interreducible, e.g.:

$$\mathcal{O} \models \top \sqsubseteq \bot \iff \mathcal{O} \models A \sqcup \neg A \sqsubseteq \bot$$

$$\mathcal{O} \models C \sqsubseteq D \iff \mathcal{O} \models C \sqcap \neg D \sqsubseteq \bot$$

~> Let's focus on unsatisfiability!

... do not provide explanations

Need for explanations

- Some entailments undesirable (e.g., unsatisfiable classes)
- Debugging necessary
 ~> delete/repair axioms responsible,
 but which axioms?
- Needle(s) in the haystack (of a large ontology)

Background on explanation



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Context!

	Time	
	Development	Deployment
Ontology Developer		
App Programmer		
App user		



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Context!

	Time	
	Development	Deployment
Ontology Developer	Our focus	
App Programmer		
App user		



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OntEng Tasks

Understanding entailments

Debugging and repair

Understanding justifications

Ontology comprehension



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OntEng Tasks

Understanding entailments

Debugging and repair

Understanding justifications

Ontology comprehension



Understanding entailments

- User notices an entailment.
- Decides to obtain an explanation for it in order to find out why it holds.

Understanding justifications

- Justification for an entailment in an ontology has been obtained.
- User wants to understand the justification better.

Ontology comprehension

- User is faced with a new ontology.
- Wants to get a better picture of it.
- Among the possible metrics:
 - → Number of entailments
 - → Average number of justifications for an entailment

...

Debugging and repair

- User is faced with
 - 1. an inconsistent ontology
 - 2. an unsatisfiable class
 - 3. or "some" undesired entailment.
- Determine the cause and debug.



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A lot of research on 1 & 2 can be done without domain knowledge



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1 & 2 have high,visible user value

Manual Debugging

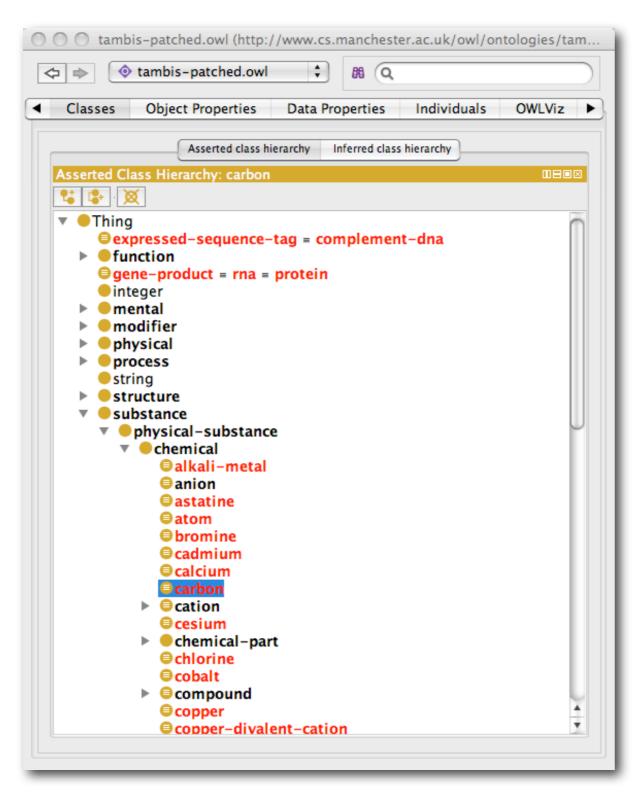
Life without explanations

- Open and classify the TAMBIS ontology tambis-patched.owl in Protégé
- Open the "Asserted class hierarchy" and expand the following terms:

substance physical-substance chemical



Life without explanations



Oops!

144 out of 395 classes are unsatisfiable (red).

Why?

Why are my classes unsatisfiable?

- What have we said in the ontology that causes these classes to be unsat?
- What do we have to change to repair it?
- Where do we start?

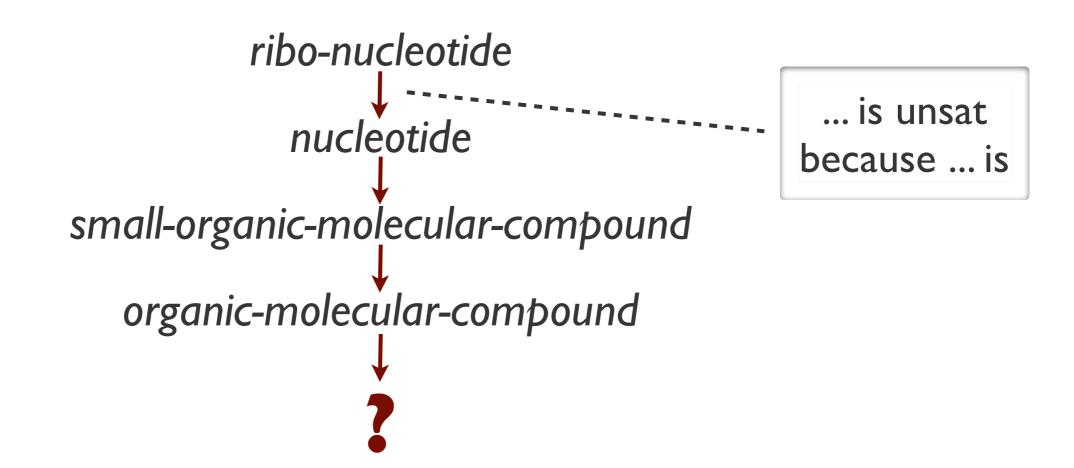
Where do we start? Tracing

- Open and classify the TAMBIS ontology tambis-patched.owl in Protégé
- Open the "Asserted class hierarchy" and expand the following terms:

```
substance
physical-substance
chemical
organic-molecular-compound
small-organic-molecular-compound
nucleotide
ribo-nucleotide
```

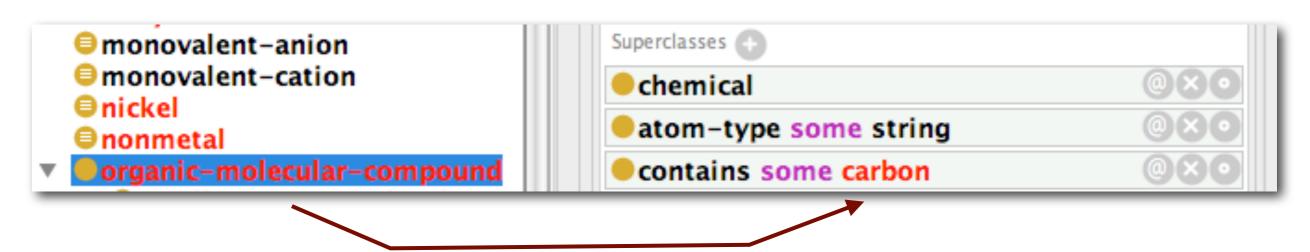
Where do we start? Tracing

Navigate through the subclass hierarchy. If a class is unsat, then so are its subclasses.



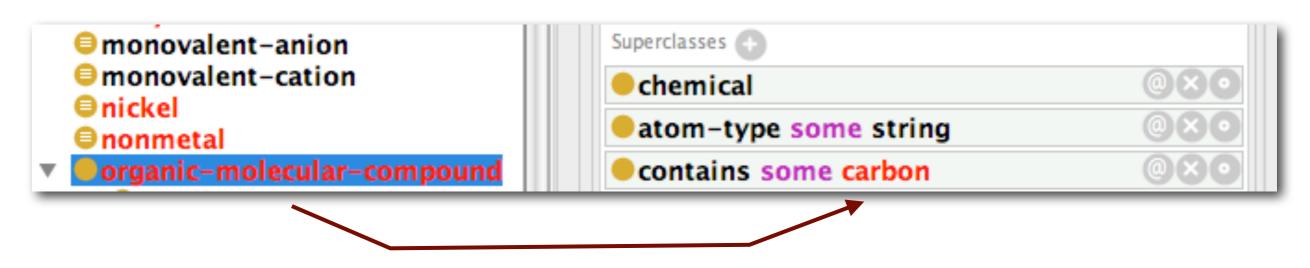
The only way was to manually follow ...

- unsatisfiable named superclasses
- or unsatisfiable fillers of restrictions in anonymous superclasses:



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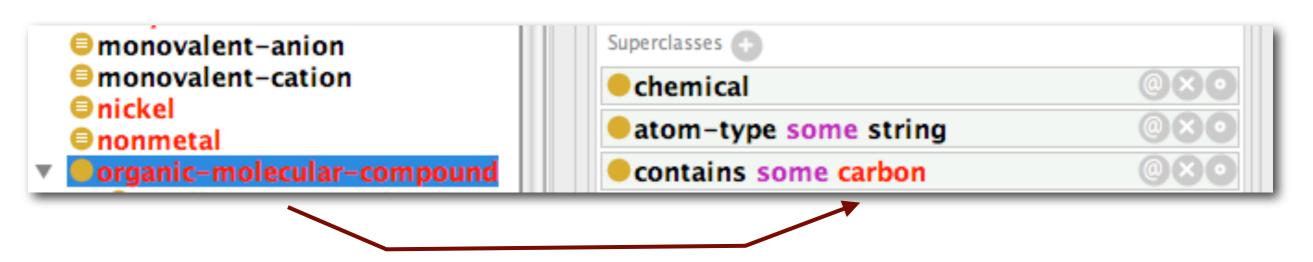
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Tedious for big ontologies! ⁽²⁾

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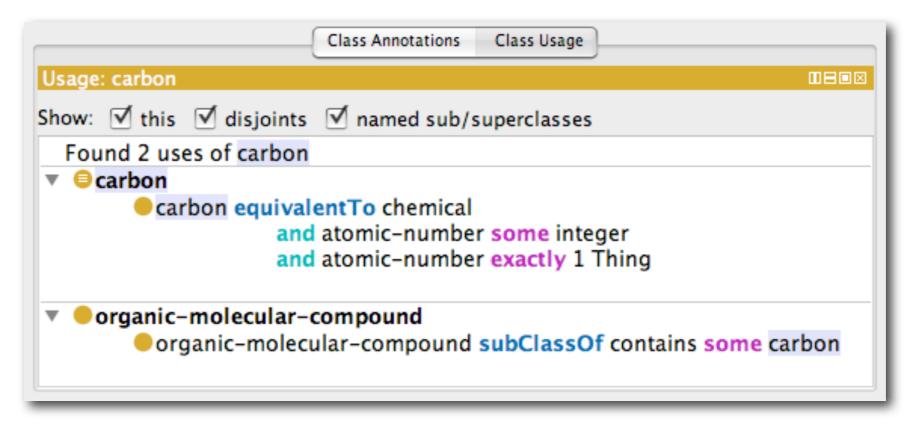


Tedious for big ontologies!

May have comprehension benefits!

Now why is the root carbon unsatisfiable?

- Axioms may be spread throughout the ontology.
- Reasons for unsat may be highly non-obvious.





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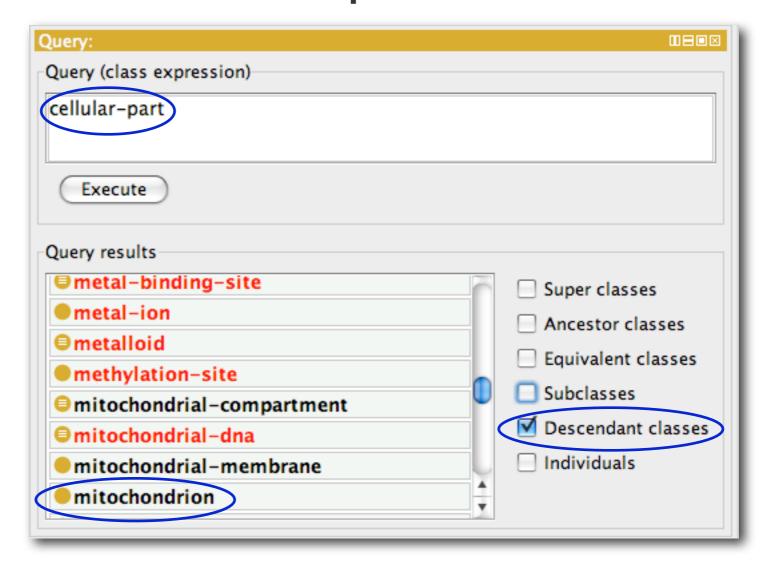
Other entailments

What about inferred subsumptions between two

class names?

mitochondrion

cellular-part



Tracing difficult and not obvious

Tracing Issues

- Identifies problem classes/terms
 - Not problem axioms
- Not all problem axioms "stick" with terms
 - Tracing can't find these
 - Depends on definition view of tool
 - Binary search of ontology?
- Multiple problem classes
 - Multiple sets of problem axioms

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Navigational approaches insufficient

Explanation

- Reasons for entailments are often difficult to understand. ~> Tool support is required.
- Theoretical basis of explanation has been investigated. New services have been defined.
- Tool support has gone from None to Respectable.
- It's now fairly easy to get explanations.
 With focus: Justifications