M68HC12 Embedded Application Binary Interface

SYSTEM V APPLICATION BINARY INTERFACE Motorola M68HC12 Processor Supplement

Version 1.0

Copyright © 1998 Motorola, Inc. All rights reserved.

Copyright © 1990-1992 The Santa Cruz Operation, Inc. All rights reserved.

Copyright © 1991,1992 88open Consortium Ltd.. All rights reserved.

This specification includes material copyrighted by 88open Consortium Ltd., which is reproduced with permission.

This specification includes material copyrighted by The Santa Cruz Operation, Inc. which is reproduced with permission.

Important Notice to Users

While every effort has been made to ensure the accuracy of all information in this document, Motorola assumes no liability to any party for any loss or damage caused by errors or omissions or by statements of any kind in this document, its updates, supplements, or special editions, whether such errors are omissions or statements resulting from negligence, accident, or any other cause. Motorola further assumes no liability arising out of the application or use of any information, product, or system described herein; nor any liability for incidental or consequential damages arising from the use of this document. Motorola disclaims all warranties regarding the information contained herein, whether expressed, implied or statutory, *including implied warranties of merchantability or fitness for a particular purpose*. Motorola makes no representation that the interconnection of products in the manner described herein will not infringe on existing or future patent rights, nor do the descriptions contained herein imply the granting or license to make, use, or sell equipment constructed in accordance with this description.

Trademarks

Motorola is a trademark of Motorola, Inc.

The Santa Cruz Operation is a registered trademark of The Santa Cruz Operation, Inc. in the USA and other countries.

88open is a registered trademark of 88open Consortium Ltd.

References

System V Application Binary Interface, Fourth Edition, UNIX System Laboratories, 1995.

System V Application Binary Interface, Motorola 88000 Processor Supplement, 88open Consortium Ltd., 1992.

DWARF Debugging Information Format, Revision 2.0.0, Industry Review Draft, UNIX International, Programming Languages SIG, July 27, 1993.

CPU12 Reference Manual, Motorola, 1996.

	2
THE MOTOROLA M68HC12 PROCESSOR AND THE SYSTEM V ABI	
How to Use the M68HC12 EABI	
EVOLUTION OF THE EABI SPECIFICATION	3
2. SOFTWARE INSTALLATION	4
3. LOW-LEVEL SYSTEM INFORMATION	5
SYSTEM V ABI EXECUTION SUPPORT	5
MACHINE INTERFACE	
Processor Architecture	
Data Representation	
FUNCTION CALLING SEQUENCE	
OPERATING SYSTEM INTERFACE	7
4. OBJECT FILES	8
Introduction	8
Data Representation	8
ELF HEADER	9
Machine Information	
Sections	
Special Sections	
Relocation	11
5. PROGRAM LOADING AND DYNAMIC LINKING	12
6. LIBRARIES	13
7. DWARF DEBUGGING INFORMATION FORMAT	14
DWARF DEFINITION	14
DWARF Register Number Mapping	14
Target-Specific Addressing Information	
Calling Convention Encodings	17
8. S-RECORDS	19
S-Record Address Format	19

1. INTRODUCTION

The Motorola M68HC12 Processor and the System V ABI

The *System V Application Binary Interface*, or *System V ABI*, defines a system interface for compiled application programs. Its purpose is to establish a standard binary interface for application programs on systems that implement the operating system interfaces defined in the *System V Interface Definition, Fourth Edition*.

The M68HC12 Embedded Application Binary Interface (which will be referred to as the EABI), described in this document, supplements the generic System V ABI, and it contains information specific to an implementation on the Motorola M68HC12 processor architecture. Together, these two specifications, the generic System V ABI and the EABI, constitute a complete System V Application Binary Interface for systems that implement the architecture of the M68HC12 processor.

The issues addressed in the *EABI* are considerably narrower in scope than the generic *System V ABI*. The *EABI* was created to meet the unique needs of M68HC12 embedded applications and tools, and does not discuss the many high-level operating system interfaces addressed in the generic *System V ABI*. The focus is primarily on the low-level system information and object files (ELF file format).

The *EABI* specifies DWARF 2.0 as the debugging information format. This format is described in *DWARF Debugging Information Format, Revision 2.0.0*.

Version 1.0 of the *EABI* defines a set of conventions for linked executable files. The goal is to define a standard that will permit a file produced by one vendor's linker to be consumed by another vendor's debugger, programmer, etc.

How to Use the M68HC12 EABI

While the generic *System V ABI* is the prime reference document, the *EABI* contains M68HC12 processor-specific implementation details, some of which supersede information in the generic one.

As with the *System V ABI*, this document refers to other publicly available reference documents (e.g. *CPU12 Reference Manual, DWARF Debugging Information Format Revision 2.0.0*). All the information referenced by the EABI should be considered part of the EABI, and just as binding as the requirements and data it explicitly includes.

Evolution of the EABI Specification

NOTE

Characteristics and behaviors mandated by this version of the *EABI* shall continue to be mandated indefinitely except where this document explicitly states otherwise. All mandates that might be withdrawn or altered in the next edition of the *EABI* are clearly identified by the NOTE identifier:

2. SOFTWARE INSTALLATION

Unlike the *System V ABI*, the *EABI* shall not have required physical media for distribution of *EABI*-conforming application software; a required software format (such as a continuous data stream) on the physical media; a required layout of files on the physical media; or a required format or interpretation of installation data files.

3. LOW-LEVEL SYSTEM INFORMATION

System V ABI Execution Support

An *EABI*-conforming entity, such as an application or a static linker, shall not have requirements pertaining to:

- dynamic linking
- global offset tables
- procedure linkage tables
- shared objects

Machine Interface

Processor Architecture

The *CPU12 Reference Manual* defines the M68HC12 architecture. Programs intended to execute directly on the processor use the CPU12 instruction set, and the instruction encodings and semantics of the architecture.

An application program can assume that all instructions defined by the architecture exist and work as documented.

To be *EABI*-conforming, the processor must implement the instructions of the architecture, perform the specified operations, and produce the expected results. The *EABI* neither places performance on systems nor specifies what instructions must be implemented in hardware. A software emulation of the architecture could conform to the *EABI*.

Data Representation

Byte Ordering

Byte ordering defines how the bytes that make up 16-bit and 32-bit values are ordered in memory. The CPU12 supports significant byte (MSB) ordering, or "Big-Endian" as it is sometimes called, meaning that the most significant byte is located in the lowest addressed byte position in a storage unit.

Fundamental Types

An entity conforming to Version 1.0 of the *EABI* shall not have requirements pertaining to the size or alignment of fundamental types (e.g., char,short,int,long,float,double).

A 4-byte floating-point number shall conform to the single-precision format of the IEEE-754 floating point specification.

An 8-byte floating-point number shall conform to the double-precision format of the IEEE-754 floating point specification.



Future versions of the EABI may define requirements for the size and alignment of fundamental types.

Aggregates and Unions

An entity conforming to Version 1.0 of the *EABI* shall not have requirements pertaining to the alignment or padding of aggregates (structures and arrays) and unions.



Future versions of the EABI may define requirements pertaining to the alignment or padding of aggregates and unions.

Bit-Fields

An entity conforming to Version 1.0 of the *EABI* shall not have requirements pertaining to the width and ranges of the integral objects defined by "bit-fields" in C struct and union definitions. Version 1.0 of the *EABI* does not define whether a "plain" bit-field is signed or unsigned.



Future versions of the EABI may define requirements for the width and range of bitfields.

Function Calling Sequence

An entity conforming to Version 1.0 of the *EABI* shall not have requirements pertaining to:

- calling conventions
- register usage during a function calling sequence
- the stack frame
- parameter passing
- variable argument lists
- return values



Future versions of the EABI may define requirements pertaining to calling conventions, register usage during a function calling sequence, the stack frame, parameter passing, variable argument lists, and return values.

Operating System Interface

Unlike the *System V ABI*, an *EABI*-conforming entity shall not have operating system interface requirements.

4. OBJECT FILES

Introduction

This chapter describes the object file format, called ELF (Executable and Linking Format).

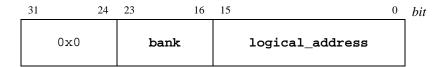


The *System V ABI* describes three main types of object files—*relocatable*, *executable*, and *shared object*. Version 1.0 of the EABI requires that conforming entities support *executable* files. The other types of object files may be addressed in a future version of the EABI.

Data Representation

Certain members of the M68HC12 family incorporate support for addressing a larger memory space than the standard 64-Kb through a paged memory or bank-switching scheme.

The *EABI* defines the *System V ABI* **Elf32_Addr** data type as having the following format in order to support banking:



bank

For logical addresses falling within a bank window, this is the value of the page register. For addresses that do not fall in a bank window, this value is 0x0. For chips in the M68HC12 family that do not support banked addressing, this value is 0x0.

logical_address

This is the address within the 64-Kb address space of the M68HC12. For banked addresses, the value of logical_address determines which page register (PPAGE, DPAGE, or EPAGE) the bank value belongs in. For memory expansion information for a particular chip in the M68HC12 family, please refer to the technical specification for that chip.

Figure 4-1, Example of **Elf32_Addr** format

The *CPU12 Reference Manual* documents that the program window (controlled by the PPAGE register) always occupies the 16-Kb space from 0x8000 to 0xBFFF. Suppose an **Elf32_Addr** address had the value 0x00129000. Its meaning is as follows:

The value of **logical_address** is 0x9000. This specifies the address within the 64-Kb address space of the M68HC12. Since 0x9000 falls within the program window this identifies that the value of **bank** belongs in the PPAGE register.

The value of **bank** is 0x12. This value belongs in the PPAGE register.

ELF Header

Machine Information

For file identification in **e_ident**, the *EABI* requires the values shown in Figure 4-2.

Figure 4-2, M68HC12 Identification, e_ident

Position	Va	lue	Comments
e_ident[EI_CLASS]	1	(ELFCLASS32)	M68HC12 address space < 4 Gbytes
e_ident[EI_DATA]	2	(ELFDATA2MSB)	M68HC12 is Big-Endian
e_ident[EI_VERSION]	1		Original ELF file format

Note that the ELFDATA2MSB value specifies that data objects in a file are encoded as 2's complement values, with the most significant byte occupying the lowest address (Big-Endian).

For the object type member **e_type**, an entity conforming to Version 1.0 of the *EABI* must support the type **ET_EXEC**, which specifies an executable file.

Processor identification resides in the ELF header's **e_machine** member and must have the value decimal 53, defined as the name **EM_68HC12**.



A draft of an ELF specification has been circulated among some tools developers with an unofficial value specified for e_machine. The value specified for the M68HC12 processor was **0x4D12**. While an object file with this value would not conform to the EABI, a consuming tool developer may wish to recognize it.

Sections



While a section header table is not required in a linked executable file, the table and the accompanying section name string table (.shstrtab section) are necessary for a debugger to locate the sections specifying debugging information.

Special Sections

In addition to the special sections defined in the generic *System V ABI*, the *EABI* defines the sections in the list below.

Figure 4-3, Special Sections

Name	Type	Attributes
.debug_abbrev	SHT_PROGBITS	none
<pre>.debug_aranges</pre>	SHT_PROGBITS	none
.debug_frame	SHT_PROGBITS	none
<pre>.debug_info</pre>	SHT_PROGBITS	none
$. { t debug_line}$	SHT_PROGBITS	none
.debug_loc	SHT_PROGBITS	none
$. { t debug_macinfo}$	SHT_PROGBITS	none
$. { t debug_pubnames}$	SHT_PROGBITS	none
.debug_str	SHT_PROGBITS	none

.debug_abbrev	This section holds DWARF 2.0 abbreviation tables.
<pre>.debug_aranges</pre>	This section holds the DWARF 2.0 lookup table to find the debugging
	information for an object by address.
.debug_frame	This section holds DWARF 2.0 virtual unwind information.
<pre>.debug_info</pre>	This section holds DWARF 2.0 debugging information entries.
.debug_line	This section holds DWARF 2.0 line number information.
.debug_loc	This section holds DWARF 2.0 location lists.
<pre>.debug_macinfo</pre>	This section holds DWARF 2.0 macro information.
.debug_pubnames	This section holds the DWARF 2.0 lookup table to find the debugging
	information for an object by name.
.debug_str	This section holds string values for DWARF 2.0 attributes.

Relocation

An entity conforming to Version 1.0 of the *EABI* shall not have requirements pertaining to relocation.



Future versions of the EABI may define requirements pertaining to relocation.

5. PROGRAM LOADING AND DYNAMIC LINKING

Program Loading

An EABI-conforming entity shall not have program loading requirements.

Dynamic Linking

An EABI-conforming entity shall not have dynamic linking requirements.

6. LIBRARIES

An entity conforming to Version 1.0 of the *EABI* shall not have requirements pertaining to library interfaces.



Future versions of the EABI may define requirements pertaining to library interfaces.

7. DWARF Debugging Information Format

DWARF Definition

The *EABI* adopts DWARF Version 2.0 as the debugging format for *EABI*-conforming applications. This format is described in *DWARF Debugging Information Format, Revision* 2.0.0.

DWARF is a specification developed for symbolic, source-level debugging. The debugging information format does not favor the design of any compiler or debugger.

DWARF Register Number Mapping

Figure 7-1 outlines the register number mapping for the M68HC12 processor family. See Section 2.4.2 of *DWARF Debugging Information Format, Revision 2.0.0* for additional information regarding register naming operations.

Figure 7-1, M68HC12 Register Number Mapping

Register Name	Number
Accumulator A	0
Accumulator B	1
Accumulator D	3
Index Register X	7
Index Register Y	8
Stack Pointer	15
Program Counter	16
Condition Code Register	17

Target-Specific Addressing Information

Any debugging information entry representing a pointer or reference type or a subroutine or subroutine type may have a **DW_AT_address_class** attribute, whose value is a constant. The set of permissible values for the M68HC12 processor family is described in Figure 7-2.

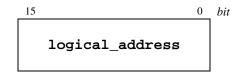
Figure 7-2, Address Class Codes

Code	Value	Meaning
DW_ADDR_none	0	No class specified
DW_ADDR_page_2	1	Page value is in byte 2

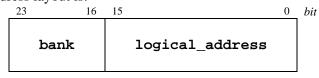
If the **DW_AT_address_class** attribute is not present or has a value of **DW_ADDR_none**, the address has the format described below in Figure 7-3.

Figure 7-3, Address Layout for DW_ADDR_none

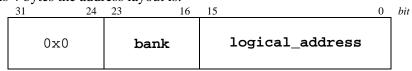
If the pointer is 2 bytes the address layout is:



If the pointer is 3 bytes the address layout is:



If the pointer is 4 bytes the address layout is:



bank

For logical addresses falling within a bank window, this is the value of the page register. For addresses that do not fall in a bank window, this value is 0x0. For chips in the M68HC12 family that do not support banked addressing, this value is 0x0.

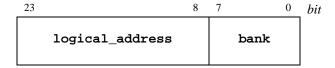
logical_address

This is the address within the 64-Kb address space of the M68HC12. For banked addresses, the value of <code>logical_address</code> determines which page register (PPAGE, DPAGE, or EPAGE) the <code>bank</code> value belongs in. For memory expansion information for a particular chip in the M68HC12 family, please refer to the technical specification for that chip.

If the **DW_AT_address_class** attribute is present and has a value of **DW_ADDR_page_2**, the address has the format described below in Figure 7-4.

Figure 7-4, Address Layout for DW_ADDR_page_2

If the pointer is 3 bytes the address layout is:



If the pointer is 4 bytes the address layout is:

31		24	23		8	7		0	bit
	0x0			logical_address			bank		

bank

For logical addresses falling within a bank window, this is the value of the page register. For addresses that do not fall in a bank window, this value is 0x0. For chips in the M68HC12 family that do not support banked addressing, this value is 0x0.

logical_address

This is the address within the 64-Kb address space of the M68HC12. For banked addresses, the value of <code>logical_address</code> determines which page register (PPAGE, DPAGE, or EPAGE) the <code>bank</code> value belongs in. For memory expansion information for a particular chip in the M68HC12 family, please refer to the technical specification for that chip.

Calling Convention Encodings

Certain members of the M68HC12 family incorporate support for addressing a larger memory space than the standard 64-Kb through a paged memory or bank-switching scheme. Functions can reside in this expanded memory. The *EABI* specifies that the DWARF DW_AT_calling_convention attribute should be used to distinguish the different ways in which a function may be called. This attribute will provide the debugger the necessary information to perform a back trace using the stack.

The encodings for values of the **DW_AT_calling_convention** attribute are given in figure 7-5.

Figure 7-5, Calling Convention Encodings

Code	Value
DW_CC_normal	0x1
DW_CC_nocall	0x3
DW_CC_far	0x40

The **DW_CC_far** value shall be used to identify a banked function (called with the CALL instruction).

The **DW_CC_nocall** value shall be used to identify an interrupt handler function (terminated by RTI instruction).

If the DW_AT_calling_convention attribute is not present or has a value of DW_CC_normal, the function shall be assumed to be non-banked (called with the JSR or BSR instructions).

Target Machine Address Size

The *DWARF Debugging Information Format, Revision 2.0.0* specification makes numerous reference to "the size of an address on the target machine." While the M68HC12 architecture defines a 16-bit logical address space, in order to support banked addressing, the *EABI* defines DWARF "target machine addresses" to be 4 bytes (32 bits) with the following format:

31		24	23		16	15			0	bit
	0x0			bank			logical_	_address		

bank

For logical addresses falling within a bank window, this is the value of the page register. For addresses that do not fall in a bank window, this value is 0x0. For chips in the M68HC12 family that do not support banked addressing, this value is 0x0.

logical_address

This is the address within the 64-Kb address space of the M68HC12. For banked addresses, the value of <code>logical_address</code> determines which page register (PPAGE, DPAGE, or EPAGE) the <code>bank</code> value belongs in. For memory expansion information for a particular chip in the M68HC12 family, please refer to the technical specification for that chip.

Thus, the value of a DWARF attribute with the form class "address" (DW_FORM_addr) is 32 bits in size and has the format described above. This also means that a DWARF 2.0 location expression resolves into a 32-bit address in the above format, and the size of the location expression stack is 32 bits.

8. S-Records

S-record Address Format

The S-record format encodes programs or data files in a printable format for transportation between computer systems.

For code/data in banked memory areas (addresses > 2 bytes), the address field of an S-record for that code/data shall be encoded with the physical address. The physical address is the value resulting from the expansion by the M68HC12 of the logical address and the page register value. Memory expansion for a particular chip in the M68HC12 family is described in the technical reference for that chip.