

# Dynamic Programming: Bottum-up Problem Solving

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## Bottom-up Problem Solving

Richard Bellman described this problem solving technique in the 1950's and called it Dynamic Programming to impress his sponsors. Dynamic Programming is an important problem solving paradigm. **Read about Bellman on Wikipedia and summarize the importance of his contributions.**

Write here.

**Describe how dynamic programming works: Solve and memoriz- ing solutions to small problems and use these solutions as building blocks to construct a solution to a larger problem from the bottom- up.**

Dynamic programming is used to optimize solutions of problems that can be solved recursively. This is beneficial because recursion can lead to many different issues because it can be an expensive algorithm to implement. The concept of memorization is relatively simple, you break up a large problem into subproblems where the results are stored to ensure that they are not computed again. This will reduce the time complexity of an application making it a much more optimal solution to many problems.

There are two primary approaches to dynamic programming, top-down and bottom-up. They both utilize breaking a problem into smaller subproblems, however they have different execution techniques. In top-down you will work your way through the subproblems just as they would be laid out in a tree. Meaning that there are different paths a program can take depending on the problem. Whereas, with bottom-up programming the layout of the subproblems is fixed and the program will only work through subproblems that are strictly necessary.

Memorization is used when the result of a subproblem is calculated so that it will not have to be recalculated, this is especially important with dynamic programming because with recursion the same problem occurs more than once. The smaller subproblems all contribute to the overall problem, rather than having an algorithm working to solve every portion of the problem at a single time.

The bottom-up approach to dynamic programming ensures that the order of the subproblems will guarantee that a subproblem will only need to be computed once. This will allow a program to consume less memory.

Both of the primary dynamic programming techniques are optimized versions of recursion that break down a larger problem into smaller subproblems, ensuring that a problem will only be solved when it is necessary. Implementing memorization for any technique will ensure that a solution is saved and will not have to be recalculated bringing down the time complexity of the problem as opposed to a recursive solution.

#### *Example uses of Dynamic Programming*

**Give at least three examples problems that use the dynamic programming paradigm to solve problems. (You should be able to search finding examples and summarize these in your group's words.)**

#### **Example One: (Spencer)**

##### *Coin Change Problem:*

One problem that can utilize dynamic programming to optimize the solution is the coin change problem. I chose this problem because it reminds me of a problem that we had in CSE 1002 and is something that I often think about because I typically handle cash at work.

##### *The problem goes as follows:*

Given an unlimited supply of coins as well as the denominations of coins as input. Find all possible ways the desired change, also given as input, can be returned. There are always different solutions to returning change, however, we typically default to the most convenient, least number of coins.

This problem can be solved using recursion, which would most likely be the go-to solution for most programmers. However, using dynamic programming you can optimize this solution bringing down the time, space and memory complexity of the problem. This being said, using dynamic programming would be the best way to solve this particular problem.

#### **Example Two: (Remi)**

##### *Longest Common Subsequence:*

A second example of a problem which can be solved using dynamic programming is the longest common subsequence problem. This problem is one that I remember learning in data structures & algorithms, so I

chose it for my example.

*The problem goes as follows:*

Given two strings, find the longest subsequence of characters that is shared between the two strings. The dynamic approach to solving this problem is performed in  $O(x * y)$  time where  $x$  and  $y$  are the lengths of the two strings being compared. This is a far better complexity than an approach where every possible subsequence is generated for each string and then compared. This would be an example of a brute-force approach, and it's time complexity is far worse.

**Example Three:** ()

### *Good approximations to hard problems*

There are Hard Problems: Problems that do not (seem) to have efficient solutions. You will explore some of these later in future quests.

There are dynamic programming algorithms that provide good approximate solutions to these hard problems. My memory tells me the 0-1 Knapsack Problem is an example of a hard problem with a good dynamic programming approximation. **Report on this example and at least 2 additional hard problems with good approximate solutions.**

#### **0-1 Knapsack Problem: (Spencer)**

*The problem goes as follows:*

The 0-1 Knapsack problem is a programming problem where as input you are given, an array of items,  $[1, 2, \dots, n-1, n]$ , and their respective weights given as an array, where each index in the array corresponds to its respective item. You are also given a maximum capacity of the the sack that is going to be used to hold the items. You must find the maximum number of items you can fit into the knapsack without exceeding the capacity of the knapsack. However, the quantity of the items that can be put into the bag is either zero or one of each specific item.

*The Solution:*

Using the iterative dynamic programming approach you would define a 2d array, where the rows corresponds to the index of the items and the weights are defined on the columns. For every weight you can either choose to ignore it or use it when constructing the 2d array. Thus, you can calculate the maximum weight that can be held in the knapsack based on the items and their corresponding weights.

By iterating through the 2d array and constructing all possibilities, comparing the maximum for each to the absolute maximum of the knapsack will give you the best solution to the given problem. The use of dynamic programming gives the most optimal space and time complexity as opposed to the brute force approach that could also be used for this problem.

#### **Longest Common subsequence: (Remi)**

*The problem goes as follows:*

The shortest path problem is a programming problem where the goal is to find the quickest route between an origin and goal point. In programming the points travelled would be represented by a graph where, for

example, cities are denoted by nodes and roads are denoted by edges. Each edge is assigned a weight which represents the length of the road between any two nodes. The combination of edges to arrive at the goal with the lowest total value represents the shortest path.

*The Solution:*

One very popular algorithm used to solve this problem is Dijkstra's algorithm, which utilizes a dynamic approach in a graph. Firstly the start node is assigned a distance of 0 and every other node is assigned infinity. First the source node is added to the queue. Each step pops the node with the shortest distance to the source node and checks the distance of all of its neighboring nodes. If the new distance is shorter than the previously assigned one, it is updated. The neighboring nodes are then added to the queue and the process repeats until all nodes have either been assigned a value, or the goal node has been reached.

Dijkstra's algorithm guarantees the optimal path between any two nodes, meaning that it will not waste time producing paths which are not the shortest between the two input nodes. By updating the shortest path to every node at every step, the algorithm will never have to backtrack and reevaluate a distance which has already been assigned. No unnecessary calculations are performed, meaning the time and space complexity is satisfactory.

<b>Name</b>	<b>Section</b>
Remington Greko	Second example of Dynamic Programming & Shortest Path
Tyler Gutowski	
Spencer Hirsch	How Dynamic Programming Works, One Example use of Dynamic Programming & 0-1 Knapsack Problem algorithm example