Data Science in Ten Minutes*

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1 Introduction

Data science is not machine learning. There is no machine learning in this guide, because machine learning is the wrong answer to most problems you're likely to run into. I've never deployed a machine learning system, although I've debugged several misbehaving ones (and will cover how to do that).

A lot of the code examples assume you're running Linux because Linux is the go-to platform for medium/big data processing. It's also ideal for small data due to optimizations in core utilities like sort. If you're running a non-Linux OS and want to try stuff, you have a few options:

- Install Docker and create a container from the ubuntu:18.04 image or similar (I'll assume you have this setup)
- Run a VM like VirtualBox and install Ubuntu Desktop inside it
- Rent a free tier Amazon EC2 instance (you can use the free-tier t2.nano or t2.micro for data science, and they're ideal for learning because they're resource-constrained)

^{*}for extremely large values of ten

 $^{^{1}\}mbox{GNU}$ sort will compress temporary files, unlike the sort that ships with Mac OSX for instance.

• Buy a cheap rack server from eBay and drop Ubuntu Server on it

Realistically, you probably won't want to use OSX or Windows for distributed programming: your binaries won't be portable to cluster machines, and you'll be fighting with things like case-insensitive filesystems, slower core utilities, and general inconsistencies that will increase debugging time.²

1.1 Examples

This guide comes with example code in the original repository. Where necessary, I've included package installation commands required to install dependencies on Ubuntu 18.04.

²All of this is a non-issue for JVM processes, e.g. Hadoop and Spark, but a lot of data science is best done with native tools that use less memory and have lower iteration time.

2 Linux

Oh yes, we are totally going here. Here's why.

Backend programs and processing pipelines and stuff (basically, "big data" things) operate entirely by talking to the kernel, which, in big-data world, is usually Linux; and this is true regardless of the language, libraries, and framework(s) you're using. You can always throw more hardware at a problem³, but if you understand system-level programming you'll often have a better/cheaper option.

Some quick background reading if you need it for the homework/curiosity:

- How to write a UNIX shell
- How to write a JIT compiler
- ELF executable binary spec (more readable version on Wikipedia, and this StackOverflow answer may be useful)
- Intel machine code documentation

2.1 Virtual memory

This is one of the two things that comes up a lot in data science infrastructure. Basically, any memory address you can see is virtualized and may not be resident in the RAM chips in your machine. The kernel talks to the MMU hardware to maintain the mapping between software and hardware pages, and when the virtual page set overflows physical memory the kernel swaps them to disk, usually with an LRU strategy.

Here's where things get interesting. Linux (and any other server OS) gives you a mmap system call to request that the kernel map pages into your program's address space. mmap, however, has some interesting options:

- MAP_SHARED: map the region into multiple programs' address spaces (this
 reuses the same physical page across processes)
- MAP_FILE: map the region using data from a file; then the kernel will load file data when a page fault occurs

MAP_FILE is sort of like saying "swap this region to a specific file, rather than the shared swapfile you'd normally use." The implication is important, though: all memory mappings go through the same page allocation cache, and any page fault has the potential to block your program on disk IO. Clever data structures like Bloom filters, Count-min sketches, and so forth are all designed to give you a way to trade various degrees of accuracy for a much smaller memory footprint.

Sometimes you won't have any good options within the confines of physical RAM, so you'll end up using IO devices to supply data; then the challenge

³Until you can't

becomes optimizing for those IO devices (SSDs are different from HDDs, for instance). I'll get to some specifics later on when we talk about sorting, joins, and compression.

2.2 File descriptors

This is the other thing you need to know about.

Programs don't typically use mmap for general-purpose IO. It's more idiomatic, and sometimes faster, to use read and write on a file descriptor. Internally, these functions ask the kernel to copy memory from an underyling file/socket/pipe/etc into mapped pages in the address space. The advantage is cache locality: you can read a small amount of stuff into a buffer, process the buffer, and then reuse that buffer for the next read. Cache locality does matter; for example:

```
# small block size: great cache locality, too much system calling overhead
$ dd if=/dev/zero count=262144 bs=32768 of=/dev/null
262144+0 records in
262144+0 records out
8589934592 bytes (8.6 GB, 8.0 GiB) copied, 0.774034 s, 11.1 GB/s

# medium block size: great cache locality, insignificant syscall overhead
$ dd if=/dev/zero count=8192 bs=1048576 of=/dev/null
8192+0 records in
8192+0 records out
8589934592 bytes (8.6 GB, 8.0 GiB) copied, 0.574597 s, 14.9 GB/s

# large block size: cache overflow, insignificant syscall overhead
$ dd if=/dev/zero count=2048 bs=$((1048576 * 4)) of=/dev/null
2048+0 records in
2048+0 records out
8589934592 bytes (8.6 GB, 8.0 GiB) copied, 1.14022 s, 7.5 GB/s
```

2.3 Pulling this together: let's write a program

...in machine language. For simplicity, let's write one that prints hello world and then exits successfully (with code 0).

This is also a good opportunity to talk about how we might generate and work with binary data with things like fixed offsets. Two simple functions for this are pack() and unpack(), variants of which ship with both Perl and Ruby.

The first part of any Linux executable is the ELF header, usually followed directly by a program header; here's what those look like as C structs for 64-bit executables (reformatted slightly, and with docs for readability):

```
typedef struct {
```

```
// 0x7f, 'E', 'L', 'F', ...
  unsigned char e_ident[16];
 uint16_t
                                // ET_EXEC = 2 for executable files
                e_type;
                                // EM_X86_64 = 62 for AMD64 architecture
 uint16_t
                e_machine:
                                // EV_CURRENT = 1
 uint32_t
                e_version;
 uint64_t
                e_entry;
                                // virtual address of first instruction
                               // file offset of first program header
 uint64_t
                e_phoff;
                e_shoff;
                               // file offset of first section header
 uint64_t
                                // always zero
 uint32_t
                e_flags;
                               // size of the ELF header struct (this one)
                e_ehsize;
 uint16_t
 uint16_t
               e_phentsize;
                               // size of a program header struct
                                // number of program header structs
 uint16_t
                e_phnum;
                                // size of a section header struct
                e_shentsize;
 uint16_t
                                // number of section header structs
 uint16 t
                e_shnum;
                e_shstrndx;
                                // string table linkage
 uint16_t
} Elf64_Ehdr;
typedef struct {
                                // the purpose of the mapping
  uint32_t p_type;
                                // permissions for the mapped pages (rwx)
 uint32_t p_flags;
                                // file offset of the first byte of data
 uint64_t p_offset;
 uint64_t p_vaddr;
                               // virtual memory offset of the data (NB below)
                               // physical memory offset (usually zero)
 uint64_t p_paddr;
                               // number of bytes from the file
 uint64_t p_filesz;
                               // number of bytes to be mapped into memory
 uint64_t p_memsz;
  uint64_t p_align;
                               // segment alignment
} Elf64_Phdr;
  If we want a very minimal executable, here's how we might write these
headers from Perl:
# elf-header.pl: emit an ELF binary and program header to stdout
use strict;
use warnings;
print pack('C16 SSL',
           0x7f, ord 'E', ord 'L', ord 'F',
           2, 1, 1, 0,
           0, 0, 0, 0,
           0, 0, 0, 0,
           2.
                                        # e_type
                                                    = ET EXEC
                                        # e_machine = EM_X86_64
           62,
           1)
                                        # e_version = EV_CURRENT
    . pack('QQQ',
```

```
0x400078,
                                      \# e_{entry} = 0x400078
       64,
                                      # e_phoff
                                      # e_shoff
       0)
. pack('LSS SSSS',
                                      # e_flags
       0,
       64,
                                      # e_ehsize
                                      # e_phentsize
       56,
       1,
                                      # e_phnum
       0,
                                      # e_shentsize
                                      # e_shnum
       0,
       0)
                                      # e_shstrndx
. pack('LLQQQQQQ',
                                      # p_type = PT_LOAD (map a region)
       1,
       7,
                                      \# p_flags = R|W|X
                                      # p_offset (must be page-aligned)
       0,
       0x400000,
                                      # p_vaddr
                                      # p_paddr
       0,
       0x1000,
                                      # p_filesz: 4KB
                                      # p_memsz: 4KB
       0x1000,
       0x1000);
                                      # p_align: 4KB
```

Now we can generate the ELF header:

```
$ sudo apt install perl  # if perl is missing
$ perl elf-header.pl > elf-header
```

You can verify that the header is correct using file, which reads magic numbers and tells you about the format of things:

```
$ sudo apt install file
$ file elf-header
elf-header: ELF 64-bit LSB executable, x86-64, version 1 (SYSV), statically
linked, corrupted section header size
```

Awesome, now let's get into the machine code.

2.4 hello world in machine code

The first thing to note is that we're talking directly to the kernel here. Our ELF header above is very minimalistic, with no linker instructions or anything else

to complicate things. So we have none of the usual libc functions like printf or exit; it's up to us to define those in terms of Linux system calls. TODO

2.5 Homework (if that's your thing)

1. Write an ELF Linux executable that consumes data from stdin and writes that data to stdout, then exits successfully.