

# Laboratory exercises (GPU)

## Eigenvalues and Eigenvectors: Power Method

Introduction HPC, TU Delft, Adapted from  
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### 1 Introduction

#### 1.1 Eigenvalues and Eigenvectors

Let  $A$  be an  $n \times n$  matrix, a non-zero vector  $x$  is an eigenvector of  $A$  if there exists a scalar  $\lambda$  such that

$$Ax = \lambda x \quad (1)$$

The scalar  $\lambda$  is called the eigenvalue of the matrix  $A$ , corresponding to the eigenvector  $x$ .

Let  $\lambda_1, \lambda_2, \lambda_3, \dots, \lambda_n$  be a set of the eigenvalues of an  $n \times n$  matrix  $A$ . If  $|\lambda_1| > |\lambda_i|$  for  $i=2,3,\dots,n$ , then  $\lambda_1$  is called the **dominant eigenvalue of  $A$**  and the eigenvectors corresponding to  $\lambda_1$ , is called the **dominant eigenvector of  $A$** .

Let us consider a system,

$$Ax = \lambda x \quad (2)$$

for which we want to find the eigenvalues and eigenvectors. The standard method for that is to solve for the roots of  $\lambda$  of the characteristic equation

$$||A - \lambda I|| = 0 \quad (3)$$

when  $A$  is large, this method is impractical. Evaluating the determinant of an  $n \times n$  matrix is a huge task, when  $n$  is large, and solving the resulting  $n$ -th degree polynomial equation for  $\lambda$  is another additional task on top of that.

The power method is a simple iteration method that can be used to find  $\lambda_1$  and  $x_1$  for a given matrix, where  $\lambda_1$  is the largest eigenvalue and  $x_1$  is the corresponding eigenvector. Similarly, the inverse power method can be used to find the smallest eigenvalue and its corresponding eigenvector, which is very similar to the power method.

#### 1.2 The Power Method

In mathematics, the power method (also known as power iteration) is an eigenvalue algorithm: given a matrix  $A$ , the algorithm will produce a number  $\lambda$ , which is the greatest (in absolute value) eigenvalue of  $A$ , and a nonzero vector  $x$ , the corresponding eigenvector of  $\lambda$ , such that  $Ax = \lambda x$ . The algorithm is also known as the Von Mises iteration<sup>1</sup>. The power iteration is a very simple algorithm, but it may converge slowly. It does not compute a matrix decomposition, and hence it can be used when  $A$  is a very large sparse matrix.

We first assume that matrix  $A$  has a dominant eigenvalue with the corresponding dominant eigenvectors. As stated before, the power method for approximating eigenvalues is iterative. Hence, we start with an initial approximation  $x_0$  of the dominant eigenvector of  $A$ , which must be non-zero.

<sup>1</sup> Richard von Mises and H. Pollaczek-Geiringer, Praktische Verfahren der Gleichungsauflösung, ZAMM - Zeitschrift für Angewandte Mathematik und Mechanik 9, 152-164 (1929).

Thus, we obtain a sequence of eigenvectors given by the recursive formula as follows,

$$x_{k+1} = \frac{Ax_k}{\|Ax_k\|} \quad (4)$$

So, at every iteration, the vector  $x_k$  is multiplied by the matrix  $A$  and normalized. If we assume  $A$  has an eigenvalue that is strictly greater in magnitude than the other eigenvalues and the starting approximation  $x_0$  of the dominant eigenvector has a nonzero component in the direction of an eigenvector associated with the dominant eigenvalue, then a subsequence  $x_k$  converges to an eigenvector associated with the dominant eigenvalue. Without the two assumptions above, the sequence  $x_k$  does not necessarily converge.

After obtaining the dominant eigenvector using the power method, ideally, one can use Rayleigh quotient<sup>2</sup> to get the associated eigenvalue. This algorithm is

the one used for problems such as *GooglePageRank*<sup>3</sup>. The method can also be used to calculate the spectral radius (or the largest eigenvalue of a matrix) by computing the Rayleigh quotient

$$\frac{x_k^T Ax_k}{x_k^T x_k} = \frac{x_k^T x_{k+1}}{x_k^T x_k} \quad (5)$$

Here, we will give some explicit descriptions about how to implement the aforementioned power method, including the iteration formula and one simple example. As have been mentioned above, a sequence of eigenvectors of matrix  $A$  can be calculated by

$$\begin{aligned} x_1 &= Ax_0 \\ x_2 &= Ax_1 = A(Ax_0) = A^2x_0 \\ x_3 &= Ax_2 = A(A^2x_0) = A^3x_0 \\ &\dots\dots\dots \\ &\dots\dots\dots \\ x_{k+1} &= Ax_k = A(A^{k-1}x_0) = A^kx_0 \end{aligned} \quad (6)$$

When  $k$  is large, we can obtain a good approximation of the dominant eigenvector of  $A$  by properly scaling the sequence.

**Example:** Use the power method to approximate a dominant eigenvalue and

the corresponding eigenvector of  $\begin{pmatrix} 0 & 11 & -5 \\ -2 & 17 & -7 \\ -4 & 26 & -10 \end{pmatrix}$  after 10 iterations.

**Solution:** We begin with an initial nonzero approximation of dominant eigenvector  $x_0$

as  $x_0 = \begin{pmatrix} 1 \\ 1 \\ 0 \end{pmatrix}$  and obtain the following approximation as

$$x_1 = Ax_0 = \begin{pmatrix} 0 & 11 & -5 \\ -2 & 17 & -7 \\ -4 & 26 & -10 \end{pmatrix} \begin{pmatrix} 1 \\ 1 \\ 0 \end{pmatrix} = \begin{pmatrix} 6 \\ 8 \\ 2 \end{pmatrix} = 12.0 \begin{pmatrix} 0.5000 \\ 0.6667 \\ 1.000 \end{pmatrix} \quad (7)$$

<sup>2</sup> Horn, R. A. and C. A. Johnson. 1985. Matrix Analysis. Cambridge University Press. pp. 176–180.

<sup>3</sup> Langville, Amy N.; Meyer, Carl D. (2003). "Survey: Deeper Inside PageRank". Internet Mathematics. 1(3).

As explained in Example 1, we take out the dominant element 12 out of the resultant matrix and the corresponding vector  $\begin{pmatrix} 0.5000 \\ 0.6667 \\ 1.000 \end{pmatrix}$  will be the new vector

for the next approximation. Proceeding in this manner we obtain a series of approximations as follows,

$$\begin{aligned}
 x_2 = Ax_1 &= \begin{pmatrix} 0 & 11 & -5 \\ -2 & 17 & -7 \\ -4 & 26 & -10 \end{pmatrix} \begin{pmatrix} 0.5000 \\ 0.6667 \\ 1.000 \end{pmatrix} = \begin{pmatrix} 2.3337 \\ 3.3339 \\ 5.3342 \end{pmatrix} = 5.3342 \begin{pmatrix} 0.4375 \\ 0.6250 \\ 1.0000 \end{pmatrix} \\
 x_3 &= 4.500 \begin{pmatrix} 0.4167 \\ 0.6111 \\ 1.0000 \end{pmatrix}, \quad x_4 = 4.222 \begin{pmatrix} 0.4079 \\ 0.6053 \\ 1.0000 \end{pmatrix}, \quad x_5 = 4.105 \begin{pmatrix} 0.4038 \\ 0.6026 \\ 1.0000 \end{pmatrix}, \\
 x_6 &= 4.051 \begin{pmatrix} 0.4109 \\ 0.6013 \\ 1.0000 \end{pmatrix}, \quad x_7 = 4.025 \begin{pmatrix} 0.4009 \\ 0.6006 \\ 1.0000 \end{pmatrix},
 \end{aligned} \tag{8}$$

Therefore, the first 7 iteration steps already show a steady convergence to the solution. The dominant eigenvalue is approximately 4.00 and the corresponding eigenvector is  $\begin{pmatrix} 0.5 \\ 0.6 \\ 1.0 \end{pmatrix}$ .

Given an  $n \times n$  matrix  $A$ , a tolerance  $\epsilon > 0$ , and the prescribed maximum allowed number of iterations  $M < \infty$ , the general power method algorithm can be formulated as follows,

$\mathbf{x} := (1, 1, 1, \dots, 1)^T$	<i>initial eigenvector estimate</i>
$\mathbf{x} := \mathbf{x} / \ \mathbf{x}\ $	<i>normalize <math>\mathbf{x}</math></i>
$\lambda := 0$	<i>initialized to any value</i>
$\lambda_0 := \lambda + 2\epsilon$	<i>make sure <math> \lambda - \lambda_0  &gt; \epsilon</math></i>
$k := 0$	
while $ \lambda - \lambda_0  \geq \epsilon$ and $k \leq M$ do	
$\mathbf{y} := A\mathbf{x}$	<i>compute next eigenvector estimate</i>
$\lambda_0 := \lambda$	<i>previous eigenvalue estimate</i>
$\lambda := \mathbf{x}^T \mathbf{y}$	<i>compute new estimate: <math>\lambda \approx \mathbf{x}^T A\mathbf{x}</math></i>
$\mathbf{x} := \mathbf{y} / \ \mathbf{y}\ $	<i>normalize eigenvector estimate</i>
$k := k + 1$	
end while	

Fig. 1: Illustration of the power method algorithm.

If the while-loop terminates with  $k \leq M$ , then we conclude the algorithm has terminated successfully. In this case,  $\lambda$  is the dominant eigenvalue and  $\mathbf{x}$  is the corresponding normalized eigenvector.

The rate of convergence of the power method depends on the difference between the magnitude of the dominant eigenvalue and other eigenvalues. Also, the power method will fail if the matrix does not have any real eigenvalues.

## 2 Parallel Implementation of Power Method in GPUs

### 2.1 Description of Parallelized Power Method

Suppose we have a parallel machine with  $p$  processors. Following the previous parallelized solution of the Poisson equation, we begin by partitioning the problem into many small individual tasks. We can partition the matrix  $A$  into individual rows where  $a_i$  is a  $1 \times n$  row vector that is the  $i^{\text{th}}$  row of  $A$ . Then the product of  $Ax$  becomes

$$y = Ax = \begin{pmatrix} a_1 \\ a_2 \\ a_3 \\ \dots \\ a_n \end{pmatrix} x = \begin{pmatrix} a_1 x \\ a_2 x \\ a_3 x \\ \dots \\ a_n x \end{pmatrix} \quad (9)$$

So the  $i^{\text{th}}$  entry of  $y$  is computed as

$$y_i = a_i x = \sum_{j=1}^n a_{ij} x_j \quad (10)$$

for each entry  $y_i$  in the result vector. Each of these is a scalar since it is the product of a  $1 \times n$  vector and an  $n \times 1$  vector (this is called an *inner product*), and each of these products can be computed in parallel. Since  $n$  will usually be larger than the number of processors  $p$  in a given parallel machine, we can assign multiple rows to each processor during the agglomeration and mapping phases.

Each task must have access to a row of  $A$  and the entire vector of  $x$  in order to compute a single  $y_i$ . Every task must then communicate its  $y_i$  value to all other tasks since the entire vector  $y$  is normalized to become the vector  $x$  for the next iteration.

One obvious way to agglomerate and map is to group tasks together based on the rows of  $A$  they use. The rows may be interleaved or consecutive, but it's probably most natural to consider consecutive rows of  $A$ . In the case of  $p$  processes the matrix  $A$  is blocked into individual blocks  $A_k$ ,  $k=1, \dots, p$ , where each block has approximately  $n/p$  rows. The products  $y = Ax$  can be depicted:

$$y = \begin{pmatrix} A_1 \\ A_2 \\ A_3 \\ \dots \\ A_p \end{pmatrix} x = \begin{pmatrix} A_1 x \\ A_2 x \\ A_3 x \\ \dots \\ A_p x \end{pmatrix} \quad (11)$$

so each process computes the portion of  $y$  given by  $y_k = A_k x$ ,  $k=1, \dots, p$ . Each process must have the entire vector  $y$  before it can compute the estimate of  $\lambda$  and create a normalized eigenvector estimate. Since every process must share its portion of  $y$  with every other process, an *allgather* operation is in order.

The evolution of GPU technology has outperformed the traditional multi-core paradigms, introducing a new opportunity in the field of scientific computing. GPU-based solutions have outperformed the corresponding sequential implementations by leaps and bounds. Particularly, for scientific computing applications, they provide a very valuable resource to parallelize and achieve high performance. This exercise focuses on analyzing its performance on CUDA, and the comparison of the power method between CPU and GPU. The details of the CPU and GPU formats and parts of pseudo-code are described in the following section.

### 2.2 Useful Reference

GPUs (Graphics Processing Units) originally designed for graphics applications become popular for scientific computing and simulations. GPUs consist of

multiprocessor elements that run under the shared-memory threads model. GPUs can run hundreds or thousands of threads in parallel and has their own DRAM. Therefore, GPUs are good at data-parallel processing. That is to say, GPU is suitable for addressing problems that can be expressed as the same operation executed on many data elements in parallel with high arithmetic intensity.

The common programming environments for GPUs are CUDA and OpenCL. In this lab experiment, we use CUDA as a parallel programming environment for GPUs. CUDA (Compute unified device architecture, see the reference on the course website for more details) is a general-purpose parallel computing platform and programming model for managing computations on the GPU. CUDA designed by NVIDIA allows developers to use C as a programming language on their GPUs. A CUDA program given in a file \*.cu consisting of a host program and kernel functions is compiled by the NVIDIA-C-compiler (**nvcc**), which separates both program parts. The execution of a CUDA program starts by executing the host program that calls kernel functions to be processed in parallel on the GPU. After invoking a kernel function, the CPU continues to process the host program which might lead to the call of another kernel function. CUDA extends the C function declaration syntax to distinguish host and kernel functions. In particular, the CUDA-specific keyword **\_\_global\_\_** indicates that the function is a kernel that can be called from a host function to be executed on a GPU. The keyword **\_\_device\_\_** indicates that the function is a kernel to be called from another kernel or device function. A host function is declared using the keyword **\_\_host\_\_**. All functions without any keywords are functions by default.

To begin with, similar to the introduction of parallelization using MPI, we give some descriptions about how to compile the codes and run it in CUDA by a *HelloWorld*. And then, parts of codes about the implementation of parallelized power method separately in CPU and GPUs also will be given. Some additional code should be added to our provided code by the student themselves to do the analysis.

**Code for *HelloWorld*** This code is complete, the students can preliminarily understand how to implement the CUDA codes.

```
// parallel HelloWorld using GPUs
// Simple starting example for CUDA program : this only works on
// arch 2 or higher
// Cong Xiao and Senlei Wang, Modified on Sep 2018

#include <stdio.h>
#include <stdlib.h>
#include <cuda.h>

#define N_THRDS      4 // Nr of threads in a block (blockDim)
#define N_BLKs      4 // Nr of blocks in a kernel (gridDim)

void checkCudaError ( const char *error )
{
    if ( cudaGetLastError () != cuda Success )
    {
        fprintf( stderr , "CUDA : %s\n" , error ) ; exit (EXIT_FAILURE) ;
    }
}

void checkCardVersion ()
{
    cudaDeviceProp prop ;
    cudaGetDeviceProperties(&prop , 0 ) ;
    checkCudaError ( " CUDA Get Device Properties failed " ) ;
}
```

```

    fprintf( stderr, " This GPU has major architecture %d, minor %d
    \n", prop . major, prop . minor ) ; i
    f( prop . major < 2 )
    {
        fprintf( stderr, "Need compute capability 2 or higher.\n"); exit(1);
31    }
}

__global__ void HelloworldOnGPU( void )
{
36    int myid = ( block Idx.x * blockDim.x ) + thread Idx.x ;
    // Each thread simply prints its own string :
    printf( " Hello World, I am thread %d in block with index %d, my
    thread index is %d\n", myid , block Idx.x , thread Idx.x ) ;
}

41 int main ( void )
{
    checkCardVersion ( ) ;
    HelloworldOnGPU <<< N_BLKs, N_THRDS >>> ( ) ;
    cuda Device Synchronize ( ) ; // without using synchronization, output
46                                won't be shown

    return 0 ;
}

```

CUDA is supported by NVIDIA GPUs. You should load the corresponding models before you use the GPUs on the DelftBlue as follows:

**module load 2023r1 nvhpc**

and execute a program **Prog\_file** using srun:

**srun --partition=gpu --gres=gpu --mem=8GB Prog\_file**

PS. if you want to use MPI with GPU (not required for exercise 3), add **--mpi=pmix** to the above command line.

The following BATCH job script **HelloWorld.sh** can be used for submitting a CUDA program using a GPU on the DelftBlue (the part of the compilation is put in comments starting with "#", often not needed).

```

#!/bin/sh -l
#SBATCH --time=00:05:00
#SBATCH --partition=gpu
#SBATCH --gres=gpu
#SBATCH --ntasks=1
5 #SBATCH --cpus-per-task=1
#SBATCH --gpus-per-task=1
#SBATCH --mem-per-cpu=8GB
#SBATCH --account=Education-EEMCS-Courses-IN4049TU
10 # adding the following lines will compile the code
    (actually, a makefile would be more appropriate)
# module load 2023r1 nvhpc
# file='HelloWorld'
# nvcc -o $file $file.cu
# srun ./ $file
srun HelloWorld

```

Submit the job by command: **sbatch HelloWorld.sh**, the result will be in output file **slurm-\*\*\*\*\*.out** (where \*\*\*\*\* is the job number). Alternatively, you could also give a name for the output, e.g., by changing the line in the script into **"srun HelloWorld.sh >HelloWo.out"**

**Description of the sequential code: `power_cpu.cu`** As a starting point, we consider the code `power_cpu.cu` that solves the Power Method problem sequentially in CPU. All the routines in this program are put together in one file. Some modes or functions are explained as follows,

1. `UploadArray` and `InitOne`. In these two routines, the specific matrix and the initial eigenvector are given, respectively.
2. `CPU_AvPoduct`. In this routine a new vector  $x_k$  for the  $k$  iteration step is calculated by multiplying the previously normalized vector  $x_{k-1}$  with the matrix  $A$ .
3. `CPU_NormalizedW`. In this routine, the new estimated vector  $x_k$  for the  $k+1$  iteration step is normalized as Eq.(4).
4. `CPU_ComputeLamda`. In this routine, the estimated largest eigenvalue of this matrix at the  $k$  iteration step is computed as Eq.(5).

And the `RUNCPUPowerMthod` is the main code of the Power Method in CPU. In the next section, a parallel version of Power Method in GPU using CUDA will be described step by step.

### 2.3 Exercise

**Building a parallel program using CUDA** The sequential version of the code that solves a Power Method problem is provided as the `power_cpu.cu`. To make a parallel version from it that uses the CUDA library, modifications have to be made and code has to be added at several places. You will already be familiar with some of these steps since they are similar to the exercises in the introduction to CUDA using `HelloWord`. In this part of the lab, you will build a working parallel code step by step. Meanwhile, you will become familiar with some additional functionality that is offered by the CUDA library. In addition to the initialization of the matrix and vector, e.g, `UploadArray` and `InitOne`, other routines, `CPU_AvPoduct`, `CPU_NormalizedW`, and `CPU_ComputeLamda`, should be recoded to a parallelized version using CUDA. Here, four subroutines, `GPU_AvPoduct`, `GPU_FindNormW`, `GPU_NormalizedW`, and `GPU_ComputeLamda`, are provided to implement the Power Method in GPU. Some tips are given here. The Power method is a matrix-vector multiplication that is computed by parallelization in GPU. However, during the procedure of the Power Method, the normalization step for the vector as Eq.(4) requires to collect of all elements of the resultant vector  $k$  from the threads, which need us to transfer the data from GPU (device) to the CPU (host), and then transfer the normalized vector from the CPU to the GPU. Because this transfer involves a small amount of data and the main parts of matrix-vector multiplication are efficiently implemented in GPU. This feature makes the Power Method very suitable for GPU. In this section, you should understand the provided four subroutines and then finish the missing part of the main code, `power_gpu.cu`. The Workflow for parallelization of the Power Method on GPUs is summarized by a flow chart in Fig.2.

**Performance analysis** After building a parallel code you now will perform some "experiments" with it. The main goal of these experiments is to analyze the performance of GPU. This can be done in a variety of ways, i.e., by adjusting

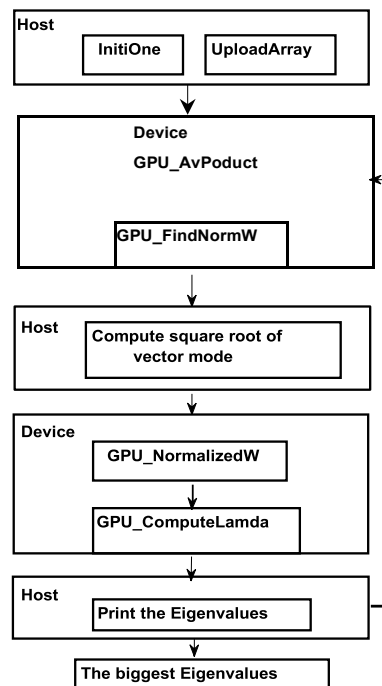


Fig. 2: Workflow for parallelization of Power Method on GPUs

the size of the matrix or the number of threads per block. This implies that it is important to know how much time the program spends in the various phases, and how this time depends on the various parameters like the number of threads and problem size. In addition, some comparisons between the implementation of the Power Method in CPU and GPU also should be conducted to further investigate the advantages of the GPU.

- Step 1: The main task of parallelizing the Power Method on GPUs is to parallelize the matrix-vector multiplication. The provided codes perform matrix-vector multiplications using CUDA with **shared** memory, a type of on-chip memory that is much faster than the **global** memory of the device, however, the limited size of share memory requires an overwhelming programming effort for utilizing the shared memory in practice. You should explore the performance by use of those two different memory access and compare the speed differences. (PS. NVidia V100 is a Volta architecture with tensor cores, it has 5120 CUDA cores and the size of the L1 cache/shared memory is 128KB and configurable up to 96 KB per SM. [NVIDIA Tesla V100 PCIe 16 GB Specs | TechPowerUp GPU Database](#))
- Step 2: Measure the execution time for a matrix  $A$  with different sizes,  $n=50, 500, 2000, 4000$ , and different numbers of threads per block, 32, 64, 100.
- Step 3: Calculate the speedups, (i) excluding the memory copy between the CPU and GPU; (ii) including the time of memory copies; (iii) (Optional, not mandatory) Now try to use Unified Memory (see slides 38, 39&41 of Lecture-week10b-GPU-Cuda) and do the same measurements as in (i) and (ii).
- Step 4: Explain the different performance results from the previous experiments.



