

```
module HuttonChap16 where
```

```
open import Haskell.Prelude
open import Haskell.Law.Equality using (sym; begin_; _≡⟨_⟩_; step-≡; _■; cong)
open import Haskell.Law.Eq.Def using (IsLawfulEq; eqReflexivity)
open import Haskell.Law.Num.Def using (+-assoc; +-comm)
open import Haskell.Law.Num.Int using (iLawfulNumInt)
```

INDUCTION ON NUMBERS

Proving the first fact about replicate:

```
replicate : {a : Set} → Nat → a → List a
replicate zero _ = []
replicate (suc n) x = x :: replicate n x
```

```
len-repl : {A : Set} → (n : Nat) → (x : A) → lengthNat (replicate n x) ≡ n
len-repl zero x = refl
len-repl (suc n) x =
  begin
    lengthNat (replicate (suc n) x)
  ≡⟨ ⟩ -- Apply replicate
    lengthNat (x :: replicate n x)
  ≡⟨ ⟩ -- Apply lengthNat
    suc (lengthNat (replicate n x))
  ≡⟨ cong suc (len-repl n x) ⟩
    suc n
  ■
```

Some facts about append:

```
++-[] : {a : Set} → (xs : List a) → xs ++ [] ≡ xs
++-[] [] = begin ([] ++ []) ≡⟨ ⟩ [] ■
++-[] (x :: xs) =
  begin
    (x :: xs) ++ []
  ≡⟨ ⟩ -- Apply ++
    x :: (xs ++ [])
  ≡⟨ cong (x ::_) (++-[] xs) ⟩
    x :: xs
  ■
```

```

++-assoc : {a : Set} → (xs ys zs : List a)
  → (xs ++ ys) ++ zs ≡ xs ++ (ys ++ zs)
++-assoc [] ys zs =
  begin
    ([] ++ ys) ++ zs
  ≡⟨⟩ -- Apply ++
    ys ++ zs
  ≡⟨⟩ -- Unapply ++
    [] ++ (ys ++ zs)
  ■

++-assoc (x :: xs) ys zs =
  begin
    ((x :: xs) ++ ys) ++ zs
  ≡⟨⟩ -- Apply ++
    (x :: (xs ++ ys)) ++ zs
  ≡⟨⟩ -- Apply ++
    x :: ((xs ++ ys) ++ zs)
  ≡⟨ cong (x ::_) (++-assoc xs ys zs) ⟩
    x :: (xs ++ (ys ++ zs))
  ≡⟨⟩ -- Unapply ++
    (x :: xs) ++ (ys ++ zs)
  ■

```

Hutton's example of elimination of append from flattening a tree:

```

data Tree (a : Set) : Set where
  Leaf : a → Tree a
  Node : Tree a → Tree a → Tree a
{-# COMPILE AGDA2HS Tree #-}

flatten : {a : Set} → Tree a → List a
flatten (Leaf x) = x :: []
flatten (Node tl tr) = flatten tl ++ flatten tr
{-# COMPILE AGDA2HS flatten #-}

flatten' : {a : Set } → Tree a → List a → List a
flatten' (Leaf x) xs = x :: xs
flatten' (Node tl tr) xs = flatten' tl (flatten' tr xs)
{-# COMPILE AGDA2HS flatten' #-}

```

```

flatten'-flatten : {a : Set} → (t : Tree a) → (xs : List a)
  → flatten' t xs ≡ flatten t ++ xs
flatten'-flatten (Leaf x) xs = refl
flatten'-flatten (Node tl tr) xs =
  begin
    flatten' (Node tl tr) xs
  ≡⟨⟩ -- Apply flatten'
    flatten' tl (flatten' tr xs)
  ≡⟨ cong (flatten' tl) (flatten'-flatten tr xs) ⟩
    flatten' tl (flatten tr ++ xs)
  ≡⟨ flatten'-flatten tl (flatten tr ++ xs) ⟩
    flatten tl ++ (flatten tr ++ xs)
  ≡⟨ sym (++-assoc (flatten tl) (flatten tr) xs) ⟩
    (flatten tl ++ flatten tr) ++ xs
  ≡⟨⟩ -- Unapply flatten
    flatten (Node tl tr) ++ xs

```

■

```

flatten'-≡-flatten : {a : Set} → (t : Tree a)
  → flatten' t [] ≡ flatten t
flatten'-≡-flatten (Leaf x) = refl
flatten'-≡-flatten (Node tl tr) =
  begin
    flatten' (Node tl tr) []
  ≡⟨⟩ -- Apply flatten'
    flatten' tl (flatten' tr [])
  ≡⟨ cong (flatten' tl) (flatten'-flatten tr []) ⟩ -- Apply the above equality
    flatten' tl (flatten tr ++ [])
  ≡⟨ flatten'-flatten tl (flatten tr ++ []) ⟩ -- Apply it again
    flatten tl ++ (flatten tr ++ [])
  ≡⟨ cong (flatten tl ++_) (++-[] (flatten tr)) ⟩ -- Remove trailing []
    flatten tl ++ flatten tr
  ≡⟨⟩ -- Unapply flatten
    flatten (Node tl tr)

```

■

COMPILER CORRECTNESS

```

data Expr : Set where
  Val : Int → Expr
  Add : Expr → Expr → Expr
{-# COMPILE AGDA2HS Expr #-}

eval : Expr → Int
eval (Val n) = n
eval (Add x y) = eval x + eval y
{-# COMPILE AGDA2HS eval #-}

Stack = List Int
{-# COMPILE AGDA2HS Stack #-}

data Op : Set where
  PUSH : Int → Op
  ADD : Op
{-# COMPILE AGDA2HS Op #-}

```

```
Code = List Op
{-# COMPILE AGDA2HS Code #-}
```

```
exec : Code → Stack → Stack
exec [] s = s
exec (PUSH n :: c) s = exec c $ n :: s
exec (ADD :: c) (m :: n :: s) = exec c $ n + m :: s
exec (ADD :: c) _ = []
{-# COMPILE AGDA2HS exec #-}
```

```
comp : Expr → Code → Code
comp (Val n) c = PUSH n :: c
comp (Add x y) c = comp x $ comp y $ ADD :: c
{-# COMPILE AGDA2HS comp #-}
```

```
comp-exec-eval : (e : Expr) → (c : Code) → (s : Stack)
→ exec (comp e c) s ≡ exec c (eval e :: s)
```

```
comp-exec-eval (Val n) c s =
  begin
    exec (comp (Val n) c) s
  ≡⟨⟩ -- Apply comp
    exec (PUSH n :: c) s
  ≡⟨⟩ -- Apply exec
    exec c (n :: s)
  ≡⟨⟩ -- Unapply eval
    exec c (eval (Val n) :: s)
```

■

```
comp-exec-eval (Add x y) c s =
  begin
    exec (comp (Add x y) c) s
  ≡⟨⟩ -- Apply comp
    exec (comp x $ comp y $ ADD :: c) s
  ≡⟨ comp-exec-eval x (comp y $ ADD :: c) s ⟩ -- Induction
    exec (comp y $ ADD :: c) (eval x :: s)
  ≡⟨ comp-exec-eval y (ADD :: c) (eval x :: s) ⟩ -- Induction Again
    exec (ADD :: c) (eval y :: eval x :: s)
  ≡⟨⟩ -- Apply exec
    exec c ((eval x) + (eval y) :: s)
  ≡⟨⟩ -- Unapply eval
    exec c (eval (Add x y) :: s)
```

■

```
compile : Expr → Code
compile e = comp e []
{-# COMPILE AGDA2HS compile #-}
```

```
compile-exec-eval : (e : Expr) → exec (compile e) [] ≡ eval e :: []
compile-exec-eval e =
```

```
  begin
    exec (compile e) []
  ≡⟨⟩ -- Apply compile
    exec (comp e []) []
  ≡⟨ comp-exec-eval e [] [] ⟩
    exec [] (eval e :: [])
  ≡⟨⟩ -- Apply exec
    eval e :: []
```

■

EXERCISE 1. Show that $\text{add } n \ (\text{Suc } m) = \text{Suc } (\text{add } n \ m)$ by induction on n

```

+-suc : (n m : Nat) → n + (suc m) ≡ suc (n + m)
+-suc zero m = refl
+-suc (suc n) m =
  begin
    (suc n) + (suc m)
  ≡⟨⟩ -- Apply +
    suc (n + suc m)
  ≡⟨⟩ -- cong suc (+-suc n m)
    suc (suc (n + m))
  ≡⟨⟩ -- Unapply +
    suc (suc n + m)
  ■

```

EXERCISE 2. Using this property, together with $\text{add } n \ \text{zero} = n$, show that addition is commutative, $\text{add } n \ m = \text{add } m \ n$, by induction on n .

```

+-zero : (n : Nat) → n + zero ≡ n
+-zero zero = refl
+-zero (suc n) =
  begin
    suc n + zero
  ≡⟨⟩ -- Apply +
    suc (n + zero)
  ≡⟨⟩ -- cong suc (+-zero n)
    suc n
  ■
+-commut : (n m : Nat) → n + m ≡ m + n
+-commut zero m =
  begin
    zero + m
  ≡⟨⟩ -- Apply +
    m
  ≡⟨⟩ -- sym (+-zero m)
    m + zero
  ■

```

```

+-commut (suc n) m =
  begin
    suc n + m
  ≡⟨⟩ -- Apply +
    suc (n + m)
  ≡⟨⟩ -- cong suc (+-commut n m)
    suc (m + n)
  ≡⟨⟩ -- sym (+-suc m n)
    m + suc n
  ■

```

EXERCISE 3. Complete the proof of the correctness of replicate by showing that it produces a list with identical elements, $\text{all } (= x) \ (\text{replicate } n \ x)$, by induction on $n \geq 0$. Hint: show that the property is always True.

```

all-repl : {a : Type} → {x : a} → {n : Nat} → (x : a)
  → all (λ _ => x) (replicate n x) ≡ True
all-repl zero x = refl
all-repl (suc n) x =
  begin
    all (λ _ => x) (replicate (suc n) x)
  ≡⟨⟩ -- Apply replicate
    all (λ _ => x) (x :: replicate n x)
  ≡⟨⟩ -- Apply all
    (x == x) && (all (λ _ => x) (replicate n x))
  ≡⟨ cong ((x == x) &&_) (all-repl n x) ⟩ -- Induction
    (x == x) && True
  ≡⟨ cong (λ _ => True) (eqReflexivity x) ⟩ -- Reflexivity x == x
    True
  ■

```

EXERCISE 4. This is `++-[]` and `++-assoc` above.

EXERCISE 5. Using the above definition for `++`, together with the definitions for `take` and `drop` show that `take n xs ++ drop n xs = xs`, by simultaneous induction on the integer `n` and the list `xs`. Hint: there are three cases, one for each pattern of arguments in the definitions of `take` and `drop`.

```

take-drop-nat : {a : Set} → (n : Nat) → (xs : List a)
  → takeNat n xs ++ dropNat n xs ≡ xs
take-drop-nat n [] = refl
take-drop-nat zero (x :: xs) =
  begin
    takeNat zero (x :: xs) ++ dropNat zero (x :: xs)
  ≡⟨⟩ -- Apply takeNat and dropNat
    [] ++ x :: xs
  ≡⟨⟩
    x :: xs
  ■

```

```

take-drop-nat (suc n) (x :: xs) =
  begin
    takeNat (suc n) (x :: xs) ++ dropNat (suc n) (x :: xs)
  ≡⟨⟩ -- Apply takeNat and dropNat and ++
    x :: takeNat n xs ++ dropNat n xs
  ≡⟨ cong (x ::_) (take-drop-nat n xs) ⟩
    x :: xs
  ■

```

```

take-drop : {a : Set} → (n : Int) → {x : a} → {xs : List a}
  → take n xs ++ drop n xs ≡ xs
take-drop n xs =
  begin
    take n xs ++ drop n xs
  ≡⟨⟩ -- Apply take and drop
    takeNat (intToNat n) xs ++ dropNat (intToNat n) xs
  ≡⟨ take-drop-nat (intToNat n) xs ⟩
    xs
  ■

```

EXERCISE 6. Given the `Tree` definition above, show that the number of leaves in such a tree is always one greater than the number of nodes, by induction on trees. Hint: start by defining functions that count the number of leaves and nodes in a tree.

```

nLeaves : {a : Set} → Tree a → Int
nLeaves (Leaf x) = 1
nLeaves (Node tl tr) = nLeaves tl + nLeaves tr
{-# COMPILE AGDA2HS nLeaves #-}

nNodes : {a : Set} → Tree a → Int
nNodes (Leaf x) = 0
nNodes (Node tl tr) = 1 + nNodes tl + nNodes tr
{-# COMPILE AGDA2HS nNodes #-}

leaves-nodes : {a : Set} → (t : Tree a)
  → nLeaves t ≡ 1 + nNodes t
leaves-nodes (Leaf x) = refl
leaves-nodes (Node tl tr) =
  begin
    nLeaves (Node tl tr)
  ≡⟨ ⟩
    nLeaves tl + nLeaves tr
  ≡⟨ cong (λ _ → nLeaves tr) (leaves-nodes tl) ⟩
    1 + nNodes tl + nLeaves tr
  ≡⟨ cong ((1 + nNodes tl) +_) (leaves-nodes tr) ⟩
    1 + nNodes tl + (1 + nNodes tr)
  ≡⟨ +-assoc 1 (nNodes tl) (1 + nNodes tr) ⟩
    1 + (nNodes tl + (1 + nNodes tr))
  ≡⟨ cong (1 +_) (sym (+-assoc (nNodes tl) 1 (nNodes tr))) ⟩
    1 + (nNodes tl + 1 + nNodes tr)
  ≡⟨ cong (1 +_) (cong (λ _ → nNodes tr) (+-comm (nNodes tl) 1)) ⟩
    1 + (1 + nNodes tl + nNodes tr)
  ≡⟨ ⟩
    1 + nNodes (Node tl tr)
  ■

```

EXERCISE 7. Verify the functor laws for the Maybe type. Hint: the proofs proceed by case analysis, and do not require the use of induction.

```

open import Haskell.Law.Functor.Def
using (IsLawfulFunctor; identity; composition)

instance
  ILawfulFunctorMaybe : IsLawfulFunctor Maybe
  identity  $\Downarrow$  ILawfulFunctorMaybe  $\Downarrow$  Nothing =
    begin
      fmap id Nothing
     $\equiv$  ( $\langle$ ) -- Apply fmap
      Nothing
     $\equiv$  ( $\langle$ ) -- Unapply id
      id Nothing
    ■

  identity  $\Downarrow$  ILawfulFunctorMaybe  $\Downarrow$  (Just x) =
    begin
      fmap id (Just x)
     $\equiv$  ( $\langle$ ) -- Apply fmap
      Just (id x)
     $\equiv$  ( $\langle$ ) -- Apply id
      Just x
     $\equiv$  ( $\langle$ ) -- Unapply id
      id (Just x)
    ■

  composition  $\Downarrow$  ILawfulFunctorMaybe  $\Downarrow$  Nothing f g =
    begin
      fmap (g  $\circ$  f) Nothing
     $\equiv$  ( $\langle$ ) -- Apply fmap
      Nothing
     $\equiv$  ( $\langle$ ) -- Unapply fmap
      fmap g Nothing
     $\equiv$  ( $\langle$ ) -- Unapply fmap
      fmap g (fmap f Nothing)
     $\equiv$  ( $\langle$ ) -- Unapply  $\circ$ 
      (fmap g  $\circ$  fmap f) Nothing
    ■

  composition  $\Downarrow$  ILawfulFunctorMaybe  $\Downarrow$  (Just x) f g =
    begin
      fmap (g  $\circ$  f) (Just x)
     $\equiv$  ( $\langle$ ) -- Apply fmap
      Just ((g  $\circ$  f) x)
     $\equiv$  ( $\langle$ ) -- Apply  $\circ$ 
      Just (g (f x))
     $\equiv$  ( $\langle$ ) -- Unapply fmap
      fmap g (Just (f x))
     $\equiv$  ( $\langle$ ) -- Unapply fmap
      fmap g (fmap f (Just x))
     $\equiv$  ( $\langle$ ) -- Unapply  $\circ$ 
      (fmap g  $\circ$  fmap f) (Just x)
    ■

```

EXERCISE 8. Given the instance declaration below, verify the functor laws for the `Tree` type, by induction on trees.


```

instance
  iFunctorTree : Functor Tree
  fmap ⟦ iFunctorTree ⟧ f (Leaf x) = Leaf (f x)
  fmap ⟦ iFunctorTree ⟧ f (Node tl tr) = Node (fmap f tl) (fmap f tr)
  _<$_ ⟦ iFunctorTree ⟧ x t = fmap (λ b → x {{b}}) t
  {-# COMPILE AGDA2HS iFunctorTree #-}

```