module HuttonChap16 where

```
open import Haskell.Prelude
open import Haskell.Law.Equality using (sym; begin_; _≡⟨⟩_; step-≡; _■; cong)
++-[] : {a : Set} \rightarrow (xs : List a) \rightarrow xs ++ [] \equiv xs
++-[] [] = begin ([] ++ []) \equiv \langle \rangle []
++-[] (x :: xs) =
    begin
       (x :: xs) ++ []
    ≡⟨⟩ -- Apply ++
      x :: (xs ++ [])
    \equiv \langle cong(x:=)(++-[]xs) \rangle
       x :: xs
++-assoc : \{a : Set\} \rightarrow (xs \ ys \ zs : List \ a)
    \rightarrow (xs ++ ys) ++ zs \equiv xs ++ (ys ++ zs)
++-assoc[] ys zs =
    begin
       ([] ++ ys) ++ zs
    ≡⟨⟩ -- Apply ++
      ys ++ zs
    ≡⟨⟩ -- Unapply ++
      [] ++ (ys ++ zs)
++-assoc (x :: xs) ys zs =
    begin
       ((x :: xs) ++ ys) ++ zs
    ≡⟨⟩ -- Apply ++
       (x :: (xs ++ ys)) ++ zs
    ≡⟨⟩ -- Apply ++
      x :: ((xs ++ ys) ++ zs)
    \equiv \langle cong (x ::_-) (++-assoc xs ys zs) \rangle
      x :: (xs ++ (ys ++ zs))
    ≡⟨⟩ -- Unapply ++
       (x :: xs) ++ (ys ++ zs)
```

Hutton's example of elimination of append from flattening a tree:

```
data Tree (a : Set) : Set where
     Leaf : a → Tree a
     Node : Tree a → Tree a → Tree a
{-# COMPILE AGDA2HS Tree #-}
flatten : {a : Set} → Tree a → List a
flatten (Leaf x) = x :: []
flatten (Node tl tr) = flatten tl ++ flatten tr
{-# COMPILE AGDA2HS flatten #-}
flatten' : {a : Set } → Tree a → List a → List a
flatten' (Leaf x) xs = x :: xs
flatten' (Node t_l t_r) xs = flatten' t_l (flatten' t_r xs)
{-# COMPILE AGDA2HS flatten' #-}
flatten'-flatten : \{a : Set\} \rightarrow (t : Tree \ a) \rightarrow (xs : List \ a)
     → flatten' t xs = flatten t ++ xs
flatten'-flatten (Leaf x) xs = refl
flatten'-flatten (Node t_1 t_r) xs =
  begin
     flatten' (Node t<sub>l</sub> t<sub>r</sub>) xs
  ≡⟨⟩ -- Apply flatten'
     flatten' t<sub>1</sub> (flatten' t<sub>r</sub> xs)
  \equiv \langle cong (flatten' t_1) (flatten'-flatten t_r xs) \rangle
     flatten' t<sub>l</sub> (flatten t<sub>r</sub> ++ xs)
  ≡⟨ flatten'-flatten t₁ (flatten tr ++ xs) ⟩
     flatten t_1 ++ (flatten t_r ++ xs)
  \equiv ( sym (++-assoc (flatten t<sub>1</sub>) (flatten t<sub>r</sub>) xs) \rangle
     (flatten t_1 ++ flatten t_r) ++ xs
  ≡⟨⟩ -- Unapply flatten
     flatten (Node t_1 t_r) ++ xs
flatten'-\equiv-flatten : {a : Set} \rightarrow (t : Tree a)
    → flatten' t [] = flatten t
flatten'-\equiv-flatten (Leaf x) = refl
flatten'-=-flatten (Node t<sub>l</sub> t<sub>r</sub>) =
  begin
     flatten' (Node t<sub>l</sub> t<sub>r</sub>) []
  ≡⟨⟩ -- Apply flatten'
     flatten' t<sub>1</sub> (flatten' t<sub>r</sub> [])
  \equiv \langle \text{ cong (flatten' } t_1) \text{ (flatten'-flatten } t_r \text{ [])} \rangle -- \text{ Inner induction}
     flatten' t_1 (flatten t_r ++ [])
  ≡⟨ flatten'-flatten t₁ (flatten tr ++ []) ⟩ -- Induction again
     flatten t_1 ++ (flatten t_r ++ [])
  \equiv \langle \text{ cong (flatten } t_1 ++- \rangle (++-[] (\text{flatten } t_r)) \rangle -- \text{ Remove trailing } []
     flatten t<sub>l</sub> ++ flatten t<sub>r</sub>
  ≡⟨⟩ -- Unapply flatten
     flatten (Node t<sub>l</sub> t<sub>r</sub>)
```