

I'm going to learn some agda!

```
data Greeting : Set where
  hello : Greeting
```

```
greet : Greeting
greet = hello
```

Defining the natural numbers:

```
data Nat : Set where
  zero : Nat
  suc  : Nat → Nat
```

```
{-# BUILTIN NATURAL Nat #-}
```

```
_+_ : Nat → Nat → Nat
zero + y = y
suc x + y = suc (x + y)
```

```
infixl 6 _+_
```

EXERCISE 1.1: Define the function `halve : Nat → Nat` that computes the result of dividing the given number by 2 (rounded down). Test your definition by evaluating it for several concrete inputs.

```
halve : Nat → Nat
halve 0 = 0
halve 1 = 0
halve (suc (suc n)) = halve n + 1
```

EXERCISE 1.2: Define the function `_*_ : Nat → Nat → Nat` for multiplication of two natural numbers.

```
_*_ : Nat → Nat → Nat
0 * y = 0
suc x * y = y + (x * y)
```

```
infixl 7 _*_
```

EXERCISE 1.3: Define the type `Bool` with constructors `true` and `false`, and define the functions for negation `not : Bool → Bool`, conjunction `_&&_ : Bool → Bool → Bool`, and disjunction `_||_ : Bool → Bool → Bool` by pattern matching.

```
data Bool : Set where
  true  : Bool
  false : Bool
```

```
not : Bool → Bool
not true = false
not false = true
```

```

id : {A : Set} → A → A
id x = x

data List (A : Set) : Set where
  [] : List A
  _::_ : A → List A → List A

infixr 5 _::_

data _×_ (A B : Set) : Set where
  _,_ : A → B → A × B

fst : {A B : Set} → A × B → A
fst (x , _) = x

snd : {A B : Set} → A × B → B
snd (_, y) = y

```

#### EXERCISE 1.4:

```

length : {A : Set} → List A → Nat
length [] = 0
length (x :: xs) = suc (length xs)

_++_ : {A : Set} → List A → List A → List A
[] ++ ys = ys
(x :: xs) ++ ys = x :: (xs ++ ys)

map : {A B : Set} → (A → B) → List A → List B
map f [] = []
map f (x :: xs) = f x :: map f xs

```

#### EXERCISE 1.5:

```

data Maybe (A : Set) : Set where
  nothing : Maybe A
  just : A → Maybe A

lookup : {A : Set} → List A → Nat → Maybe A
lookup [] _ = nothing
lookup (x :: xs) zero = just x
lookup (x :: xs) (suc i) = lookup xs i

data Vec (A : Set) : Nat → Set where
  [] : Vec A 0
  _::_ : {n : Nat} → A → Vec A n → Vec A (suc n)

replicate : {A : Set} → (n : Nat) → A → Vec A n
replicate zero x = []
replicate (suc n) x = x :: replicate n x

```

#### EXERCISE 2.1:

```

downFrom : (n : Nat) → Vec Nat n
downFrom zero = []
downFrom (suc x) = x :: downFrom x

```

```

_++Vec_ : {A : Set} {m n : Nat}
  → Vec A m → Vec A n → Vec A (m + n)
[] ++Vec ys = ys
(x :: xs) ++Vec ys = x :: (xs ++Vec ys)

head : {A : Set} {n : Nat} → Vec A (suc n) → A
head (x :: _) = x

```

#### EXERCISE 2.2:

```

tail : {A : Set} {n : Nat} → Vec A (suc n) → Vec A n
tail (_ :: xs) = xs

```

#### EXERCISE 2.3:

```

dotProduct : {n : Nat} → Vec Nat n → Vec Nat n → Nat
dotProduct [] [] = 0
dotProduct (x :: xs) (y :: ys) = x * y + dotProduct xs ys

data Fin : Nat → Set where
  zero : {n : Nat} → Fin (suc n)
  suc : {n : Nat} → Fin n → Fin (suc n)

zero3 : Fin 3
zero3 = zero

lookupVec : {A : Set} {n : Nat} → Vec A n → Fin n → A
lookupVec (x :: xs) zero = x
lookupVec (x :: xs) (suc i) = lookupVec xs i

```

#### EXERCISE 2.4:

```

putVec : {A : Set} {n : Nat} → Fin n → A → Vec A n → Vec A n
putVec zero x (_ :: xs) = x :: xs
putVec (suc i) x (x1 :: xs) = x1 :: putVec i x xs

data Σ (A : Set) (B : A → Set) : Set where
  _,_ : (x : A) → B x → Σ A B

fstΣ : {A : Set} {B : A → Set} → Σ A B → A
fstΣ (x , _) = x

sndΣ : {A : Set} {B : A → Set} → (z : Σ A B) → B (fstΣ z)
sndΣ (x , y) = y

_×'_ : (A B : Set) → Set
A ×' B = Σ A (λ _ → B)

```

#### EXERCISE 2.5:

```

fromProd : {A B : Set} → A × B → A ×' B
fromProd (x , y) = x , y

toProd : {A B : Set} → A ×' B → A × B
toProd (x , y) = x , y

```

```
List' : (A : Set) → Set
List' A =  $\Sigma$  Nat (Vec A)
```

### EXERCISE 2.6:

```
[]' : {A : Set} → List' A
>[]' = 0 , []
```

```
_::'_ : {A : Set} → A → List' A → List' A
x ::' (n , xs) = suc n , (x :: xs)
```

```
fromList : {A : Set} → List A → List' A
fromList [] = []'
fromList (x :: xs) = x ::' fromList xs
```

```
fromList' : {A : Set} → List' A → List A
fromList' (0 , []) = []
fromList' (suc n , (x :: xs)) = x :: (fromList' (n , xs))
```

### EXERCISE 3.1:

```
data Either (A : Set) (B : Set) : Set where
  left : A → Either A B
  right : B → Either A B
```

```
cases : {A B C : Set} → Either A B → (A → C) → (B → C) → C
cases (left x) f _ = f x
cases (right x) _ f = f x
```

```
data  $\tau$  : Set where
  tt :  $\tau$ 
```

```
data  $\perp$  : Set where
```

```
absurd : {A : Set} →  $\perp$  → A
absurd ()
```

### EXERCISE 3.2

- If  $A$  then ( $B$  implies  $A$ )
 

```
p1 : {A B : Set} → A → (B → A)
p1 x _ = x
```
- If ( $A$  and *true*) then ( $A$  or *false*)
 

```
p2 : {A : Set} → (A ×  $\tau$ ) → (Either A  $\perp$ )
p2 (x , tt) = left x
```
- If  $A$  implies ( $B$  implies  $C$ ), then ( $A$  and  $B$ ) implies  $C$ .
 

```
p3 : {A B C : Set} → (A → (B → C)) → ((A × B) → C)
p3 f (x , y) = f x y
```
- If  $A$  and ( $B$  or  $C$ ), then either ( $A$  and  $B$ ) or ( $A$  and  $C$ ).
 

```
p4 : {A B C : Set} → (A × (Either B C)) → (Either (A × B) (A × C))
p4 (x , left y) = left (x , y)
p4 (x , right z) = right (x , z)
```

- If A implies C and B implies D, then (A and B) implies (C and D).

p5 : {A B C D : Set} → ((A → C) × (B → D)) → ((A × B) → (C × D))

p5 (f , g) (x , y) = f x , g y

proof3 : {P Q R : Set} → (Either P Q → R) → (P → R) × (Q → R)

proof3 f = (λ x → f (left x)) , (λ x → f (right x))

EXERCISE 3.3: Write a function of type {P : Set} → (Either P (P → ⊥) → ⊥) → ⊥.

Assuming (Either P (P → ⊥) → ⊥) then proof3 above says that P → ⊥ and (P → ⊥) → ⊥.

Applying (P → ⊥) → ⊥ to P → ⊥ results in ⊥ which proves the theorem.

constructive-P-or-not-P : {P : Set} → (Either P (P → ⊥) → ⊥) → ⊥

constructive-P-or-not-P f = (λ notx → f (right notx)) (λ x → f (left x))

Some even stuff:

data IsEven : Nat → Set where

zeroIsEven : IsEven zero

sucsucIsEven : {n : Nat} → IsEven n → IsEven (suc (suc n))

6-is-even : IsEven 6

6-is-even = sucsucIsEven (sucsucIsEven (sucsucIsEven zeroIsEven))

7-is-even : IsEven 7 → ⊥

7-is-even (sucsucIsEven (sucsucIsEven (sucsucIsEven ())))

data IsTrue : Bool → Set where

TrueIsTrue : IsTrue true

\_ =Nat\_ : Nat → Nat → Bool

zero =Nat zero = true

suc x =Nat suc y = x =Nat y

\_ =Nat \_ = false

length-is-3 : IsTrue (length (1 :: 2 :: 3 :: []) =Nat 3)

length-is-3 = TrueIsTrue

double : Nat → Nat

double zero = zero

double (suc n) = suc (suc (double n))

double-is-even : (n : Nat) → IsEven (double n)

double-is-even zero = zeroIsEven

double-is-even (suc n) = sucsucIsEven (double-is-even n)

n-equals-n : (n : Nat) → IsTrue (n =Nat n)

n-equals-n zero = TrueIsTrue

n-equals-n (suc n) = n-equals-n n