

Artificial Intelligence & Machine Learning

UNIT - I

- * AI is defined as the study of systems that act in a way that to any observer would appear to be intelligent.
- * AI involves using methods based on the intelligent behaviour of humans & other animals to solve complex problems.
- * Abilities to achieve AI by a machine :
 - i) Problem solving
 - ii) logical Reasoning
 - iii) Generalized learning .

According to Aristotle a syllogism is defined as "a discourse in which certain things having been stated, something else follows of necessity from their being said".

PROLOG → PROgramming in LOGIC
 LISP → LIST Programming

Rules

contains(x,y) :- made from (x,z), contains(z,y)

B :- A

- 1) i) tasty (cheese) → cheese is tasty
ii) made from (cheese, milk)
Cheese is made from milk.
iii) contains (milk, calcium)
Milk contains calcium.

B :- A

X contains Y
X is made from Z? ⇒ X contains Y.
Z contains Y
⇒ X contains Y

X contains & is made from Z & Z
contains Y then X contains Y

- 2) animals (mammals)
mammals (elephants)
animals (elephants)
→ Mammals are animals
Elephants are mammals
Elephants are animals.
→ Mammals are animals & elephants are
mammals then elephants are animals.

animals (elephants) :- animals (mammals),
mammals (elephants)

animalst

- 3) has (mammals, legs)
 is (dog, mammal)
 has (cat, legs)

mammals - has

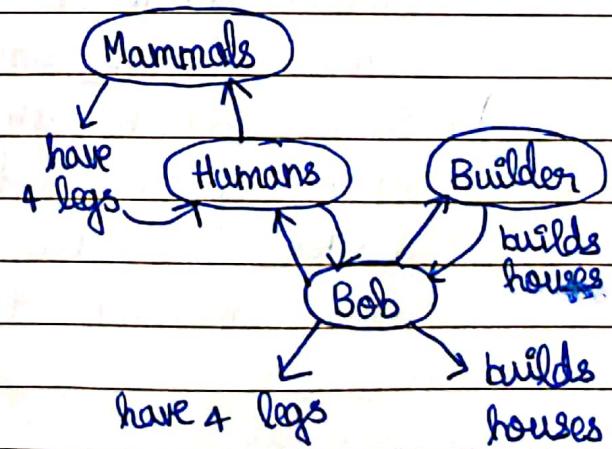
has (dog, legs) :- is (dog, mammals), has (mammal, legs)

* Frames, slots, slot values

- Aggregation & association
- * → Overriding symbol. (Add it before slot name to indicate that the slot value can be changed)
- We can express slot as a frame.

<u>Frame Name</u>	<u>Slot</u>	<u>Slot value</u>
No. of legs	Min	0
	Max	4

Humans are Mammals
 Mammals have 4 legs
 Bob is a human
 Bob is a builder
 Builder builds houses.



Admissibility :
c) Invocability : search methods are not allowed

Procedures

Procedures are set of instructions associated with frames which are called on request.

Eg:

Slot reader : is called when it wants to read the value of a slot.

Slot writer : value to be inserted into a slot.

Instance constructor : creates new instance for a class.

→ WHEN NEEDED PROCEDURES

Demons

- It is a special type of procedure that is run automatically whenever a particular value changes or a particular event occurs.
- Some demons are called when a particular value is read or they are called automatically when the user of the system wants to know what value it plays in a particular slot.
Such demons are called WHEN READ PROCEDURES.
- Some demons are automatically called when value of a slot is automatically changed. Such demons are called WHEN CHANGED PROCEDURES & OR WHEN WRITTEN PROCEDURES.

Implementation of Frames

```
Function find_slot_value(s, F)
{
    if F[s] == v
        then return v
    else if instance(F, F')
        then return find_slot_value(s, F')
    else if subclass(F, F_i)
        then return find_slot_value(s, F_i)
    else return fail
```

AI & ML Lab

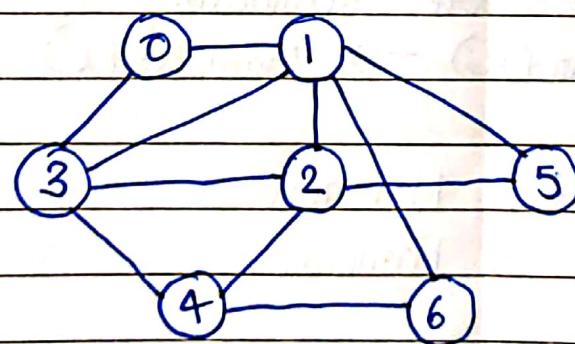
Graph Traversals

DFS

- Depth First Search.
- Uses stack (FILO).

BFS

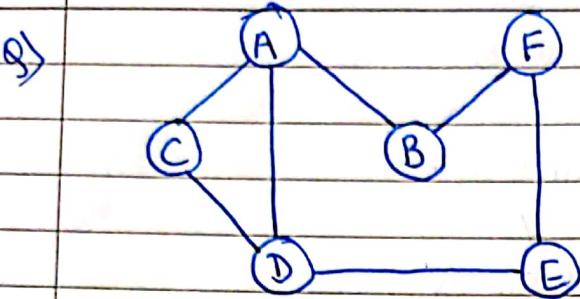
- Breadth First Search.
- Uses queue (FIFO)
- Level order traversal.



tg: 6 4 3 5 2 1 0

fg: 0 1 3 2 5 6 4

• Application : LinkedIn



Taking A as root

DFS : C D E F B A

BFS : A B C D F E

C
D
E
F
B
A

A	B	C	D	F	E
---	---	---	---	---	---

Representational Adequacy

i) All Humans are Mammals

$$\rightarrow \forall x \text{ Human}(x) \rightarrow \text{Mammal}(x)$$

ii) All Dogs are Mammals

$$\rightarrow \forall x \text{ Dog}(x) \rightarrow \text{Mammal}$$

iii) Some Dogs have 3 legs

$$\rightarrow \exists x \text{ Dog}(x) \rightarrow \text{Legs}(x, 3)$$

iv) Not all like both Mathematics & Science

$$\rightarrow \neg \forall x \ x \rightarrow \text{like}(x, \text{Mathematics}) \wedge \text{like}(x, \text{Science})$$

\neg → Negation sign

\wedge → And

- v) Some dogs chase cats & eat cheese
 $\rightarrow \exists x \text{ Dogs}(x) \rightarrow \text{chase}(x, \text{cats}) \wedge \text{eat}(x, \text{cheese})$
- vi) Some dogs chase cats & do not eat cheese
 $\rightarrow \exists x \text{ Dogs}(x) \rightarrow \text{chase}(x, \text{cats}) \wedge \neg[\text{eat}(x, \text{cheese})]$

Search Space

- It is a representation of ~~the~~ set of all possible choices in a given problem.
 - One or more choices can be the soln to the problem.
- * One method to find a solution is Heuristic approach. (But this method does not ensure an optimistic soln).

Robot, Block example → Imp.

Semantic Tree

Semantic net

- 1) One predecessor (parent node) & can have more than one successors (children).
leaf node
- 2) Root node which do not have predecessor in the starting pt.
- 3) The nodes do not have successors is called leaf nodes.

The leaf nodes which lead to a successful search are called goals.

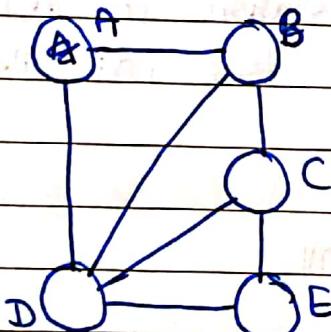
4) Path is the root through the semantic tree.

A path which leads from root node to the goal node is called complete path.

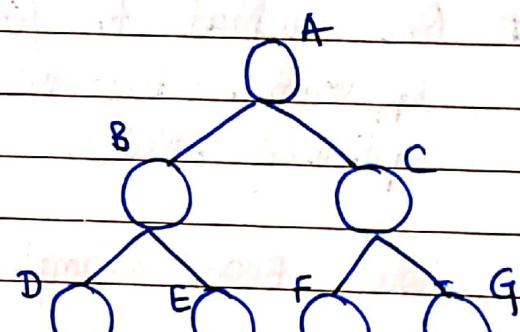
A path which leads from root node to the leaf node which are not goal nodes is called partial path.

5) An ancestor of a node is a node further up the tree in some path.

Descendant comes after a node in a path in a tree.

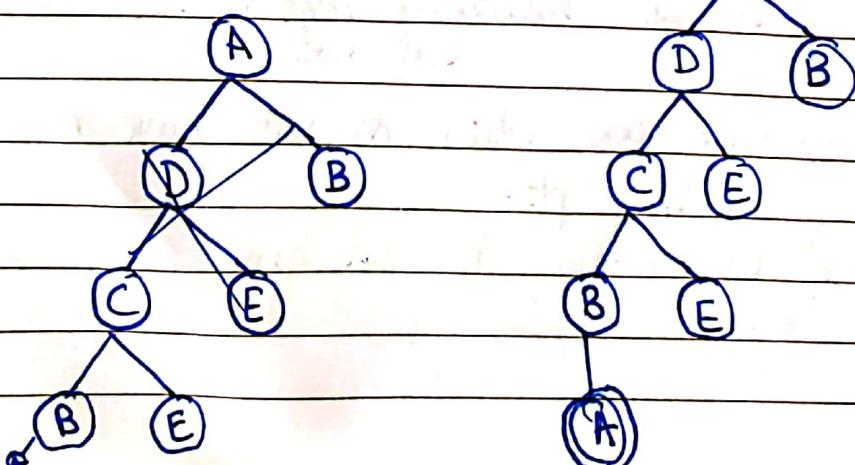


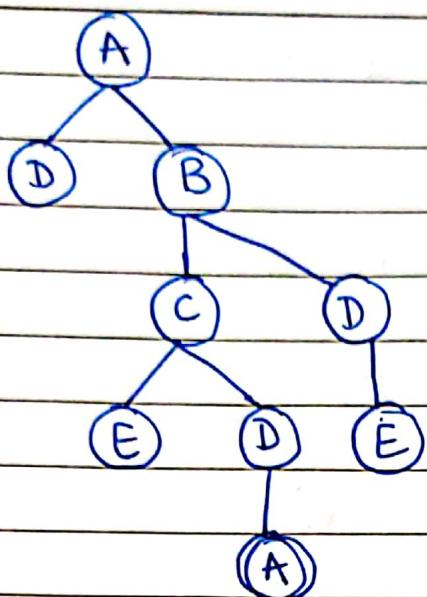
i) Semantic net



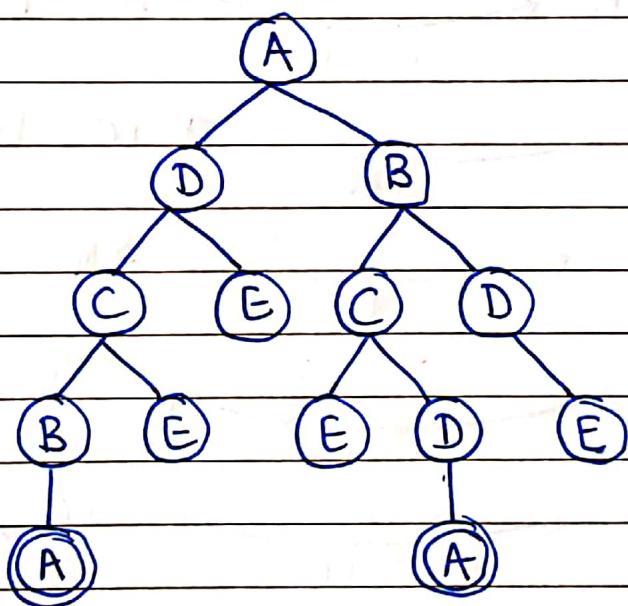
ii) Semantic tree

Search trees





Final solution (All paths)



Missionaries & Cannibals

1) Move 1C to the other side

2) Move 2C " " "

3) " 1M " " "

4) 2M " " " "

5) " 1C & 1M " " "

{
1)
2)
3)
4)
5)
operators

$3C, 3M$ $1C$

$2C$

$0, 0$

$1C, 3M$

$1C$

$1C, 0$

$3M$

$2C$

$1C$

$3M$

$1C$

$2C$

$1C, 1M$

$2M$

$2C$

$1C, 1M$

$1C, 1M$

$1C, 1M$

$2C$

$2M$

$1C, 1M$

$2C$

$3M$

$1C$

$2C$

$3M$

$1C$

$1C$

$1C, 3M$

$2C$

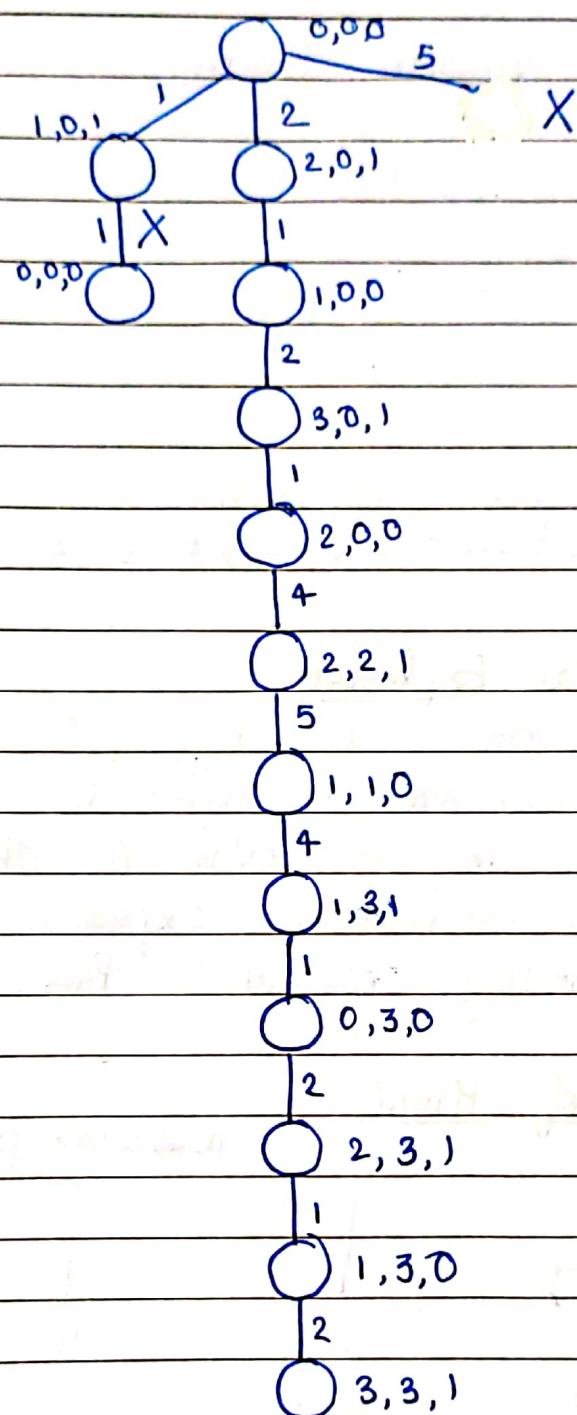
$1C, 3M$

$0, 0, 0$

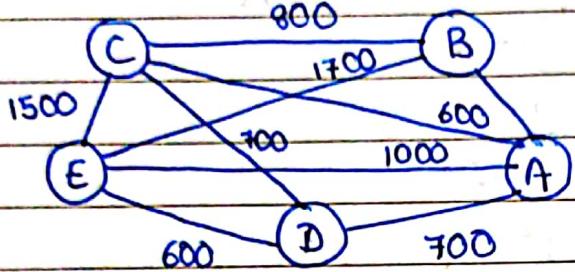
cannibals

cannoy

missinories



Travelling Salesman Problem



Nearest neighbor approach

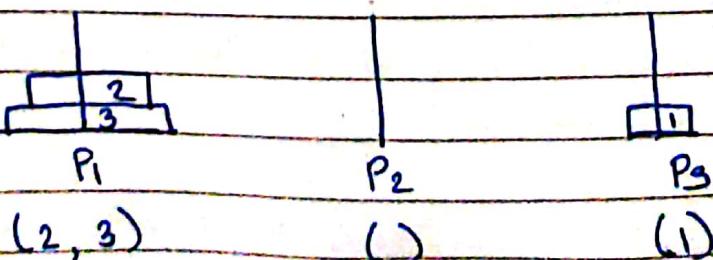
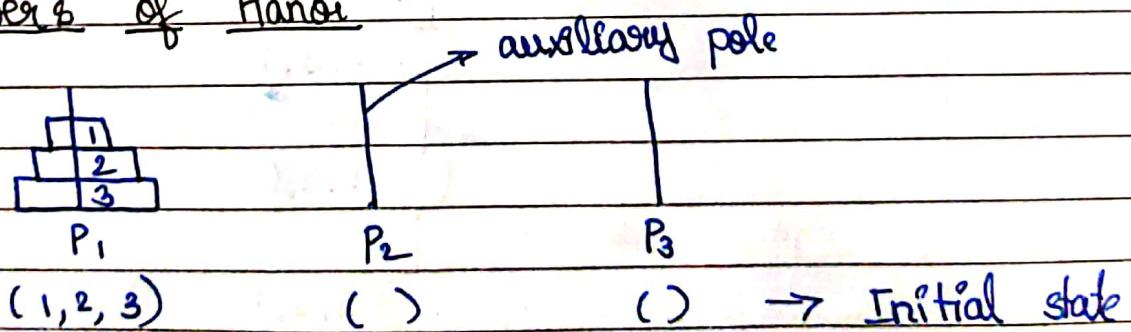
$$A \rightarrow C \rightarrow D \rightarrow E \rightarrow B \rightarrow A = 4500$$

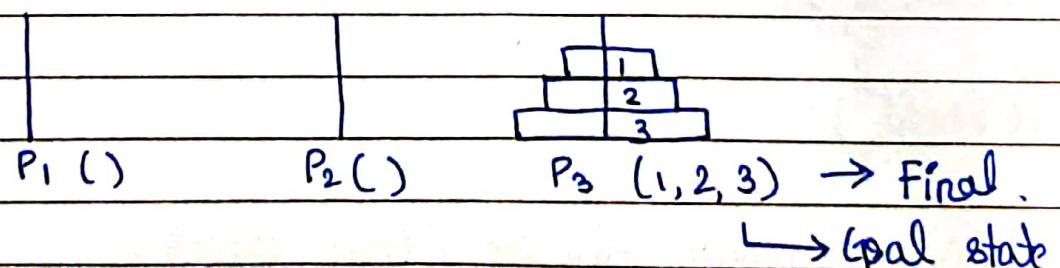
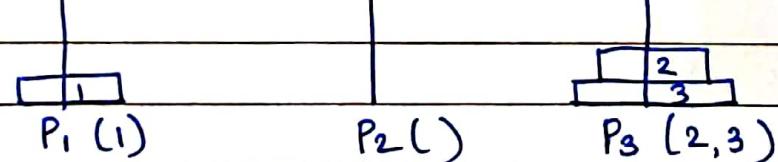
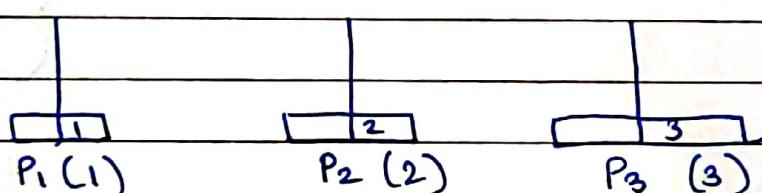
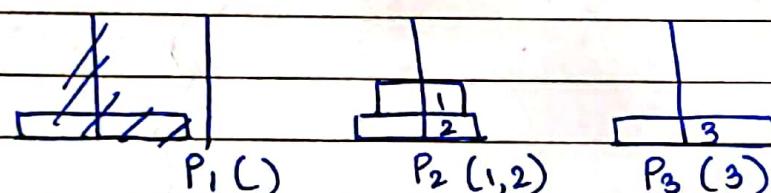
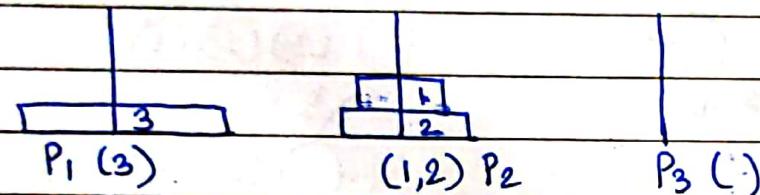
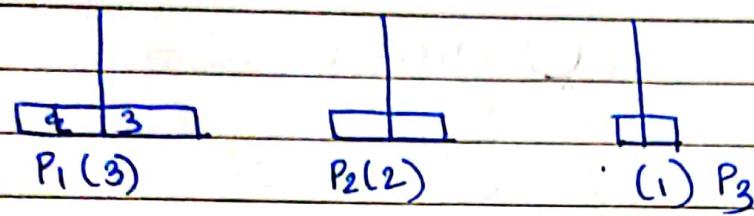
Combinatorial Explosion

As the problem gets bigger, it becomes more & more unreasonable to expect the computer to be pgm to be able to solve it. This problem is known as combinatorial explosion.

Eg : Travelling Salesman's Problem.

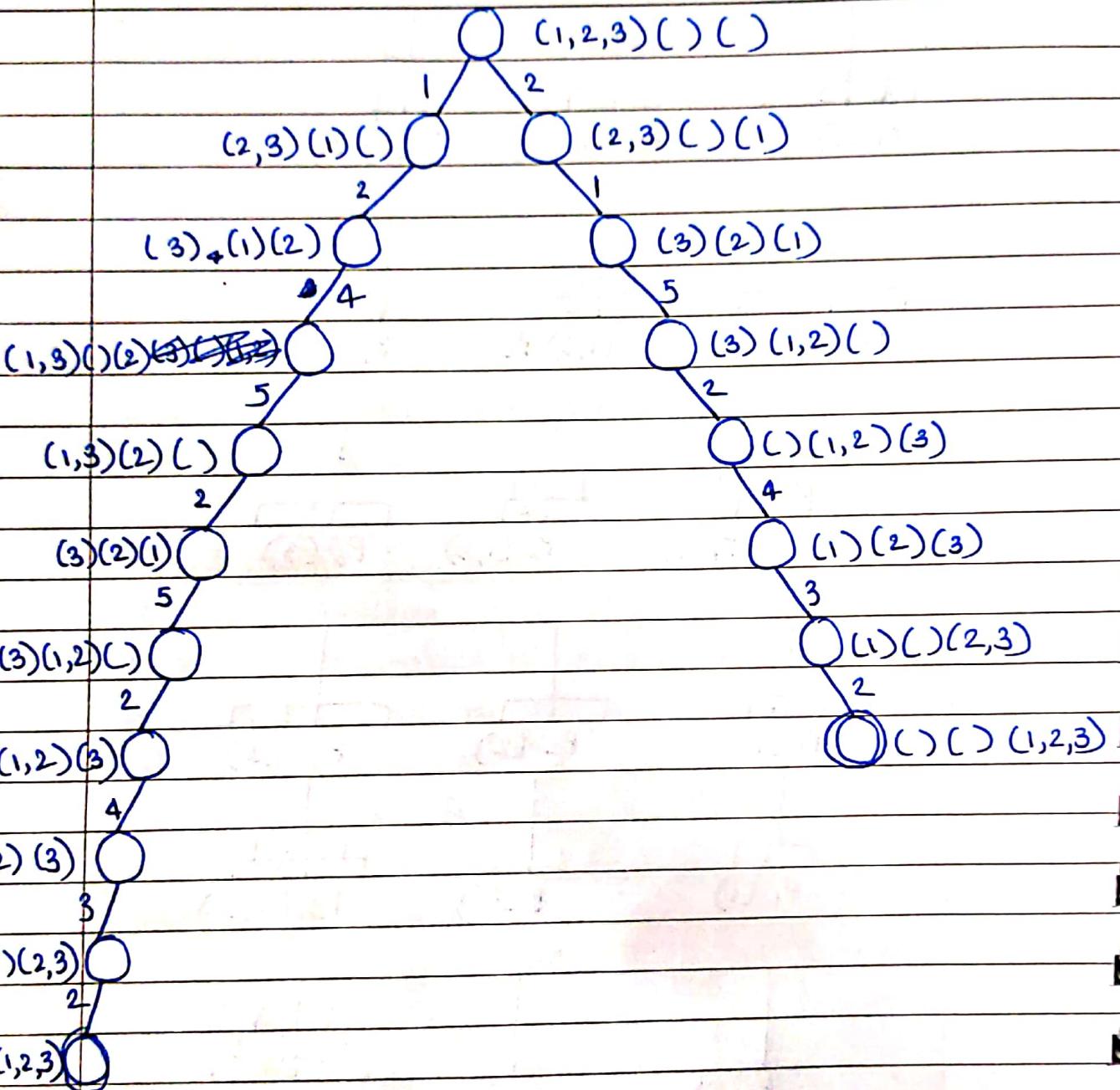
Towers of Hanoi





Operators

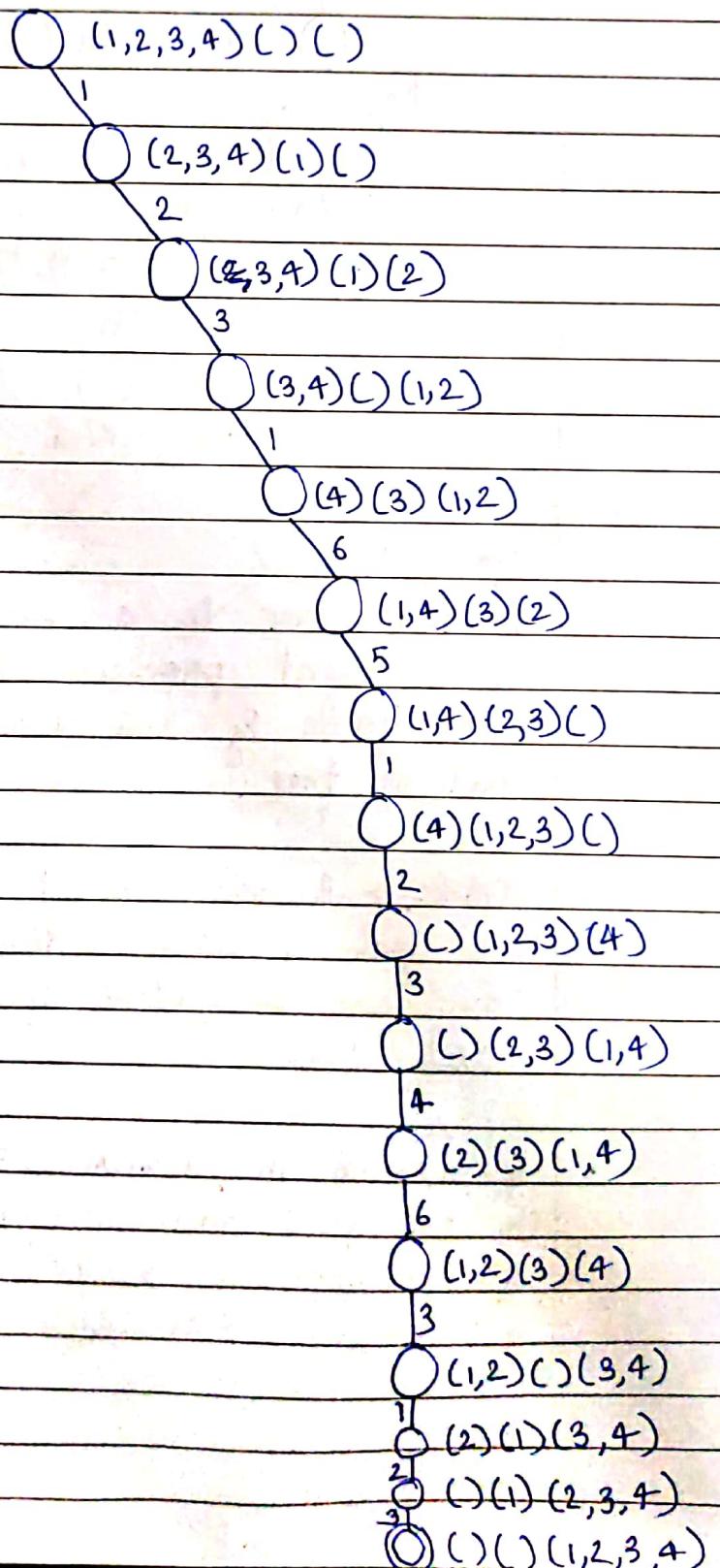
- 1) Move disc from P₁ to P₂
- 2) " " " P₁ " P₃
- 3) " " " P₂ " P₃
- 4) " " " P₂ " P₁
- 5) " " " P₃ " P₂
- 6) " " " P₃ " P₁



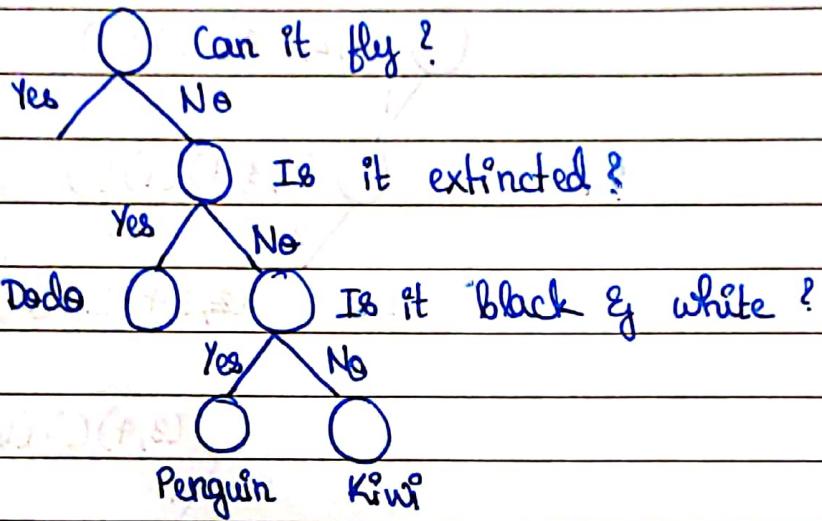
Problem reduction / Goal reduction

A complex problem can be solved by breaking it down into several smaller problems. If we solve all those smaller sub problems then we have solved the main problem.

Where no. of discs are 4 :



Describe & Match

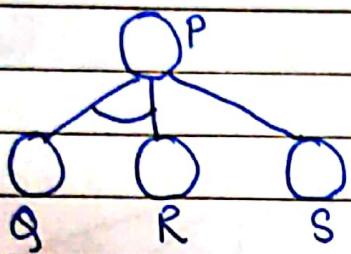


Goal tree

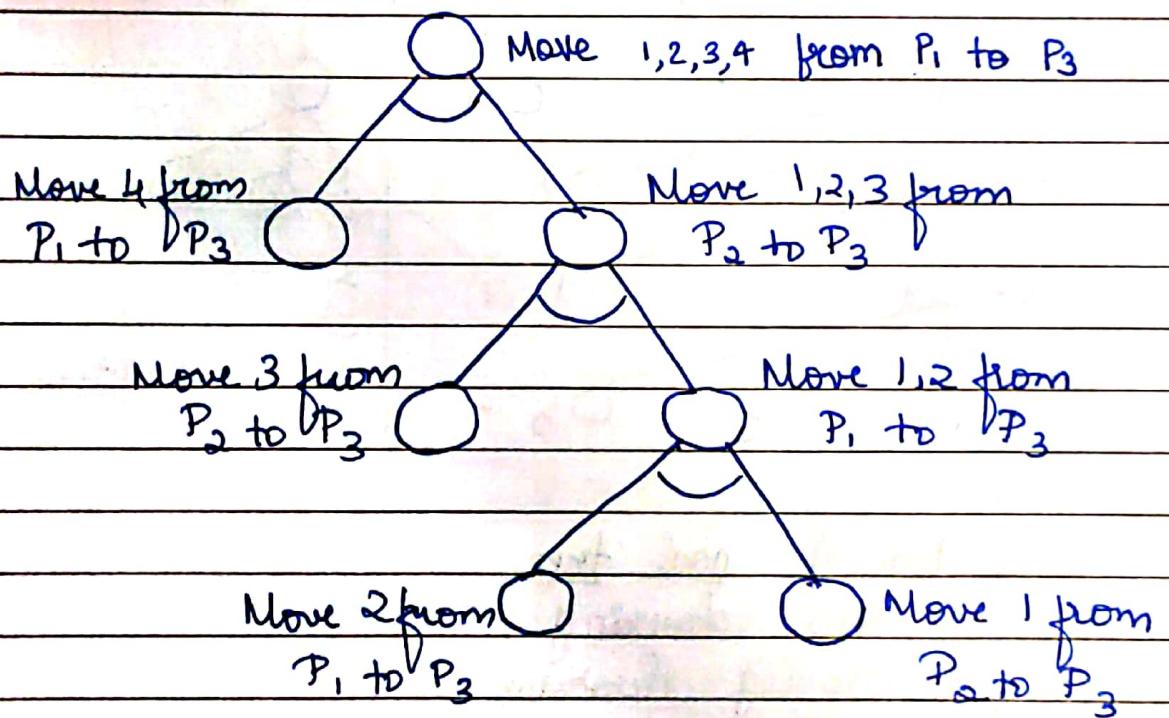
A goal tree is a form of semantic tree used to represent problems than can be broken down into goals & sub goals. These are also called And-or tree.

~~Goal~~ → A "solution" to the problem is referred as a goal & each individual step performed along the way to achieve the goal. is known as a sub goal.

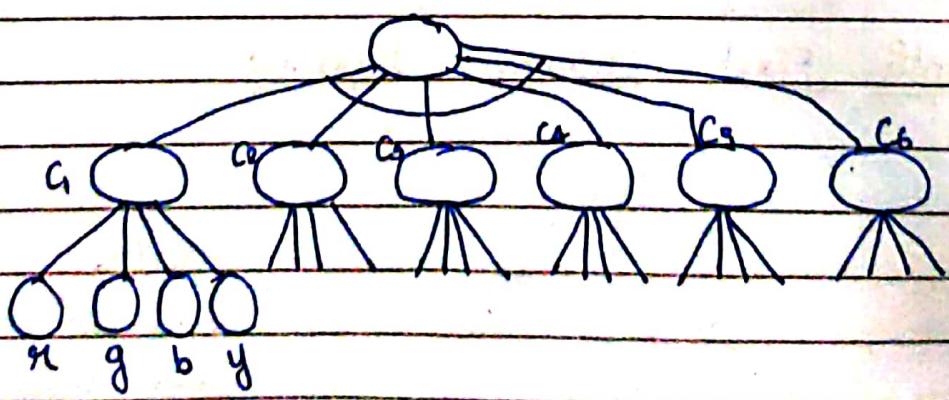
- * If a goal is achieved only by achieving all of its sub goals then such nodes are called and nodes.
- * If a goal is achieved by solving only one of its subgoals then such nodes are called or nodes.

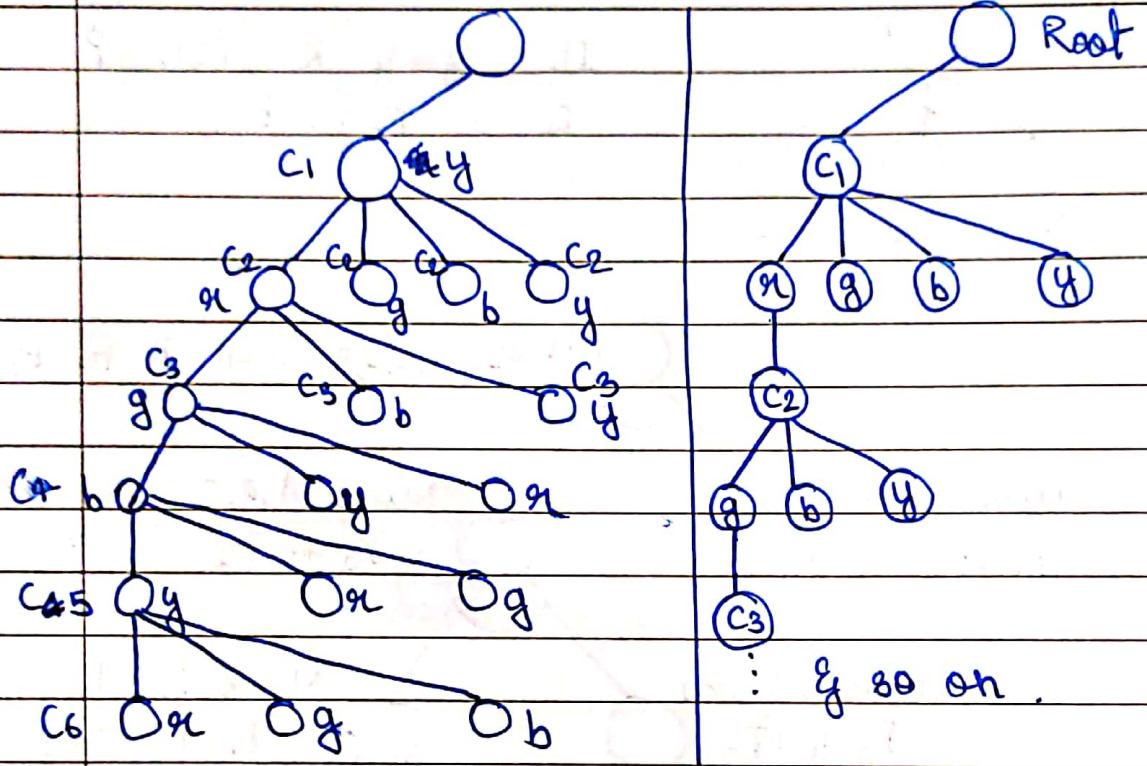


If $b \leq g \leq R$ then P
If $s \leq g$ then P .



Map colouring



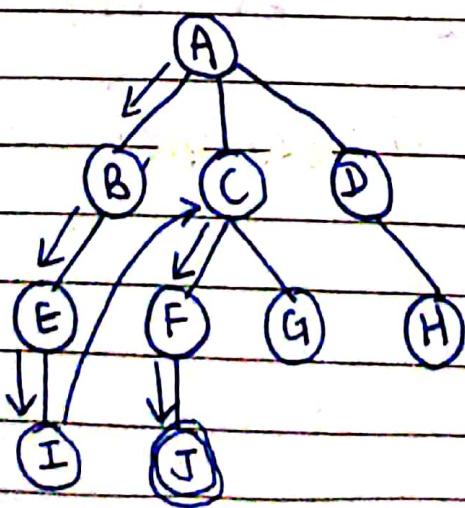


Uses of goal tree

- i) Map colouring
- ii) Proving theorems

LAB

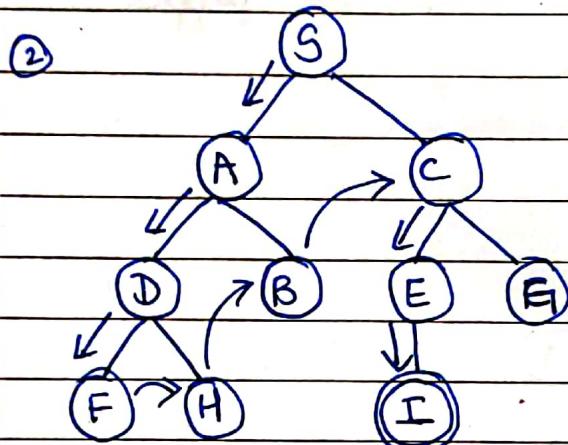
Depth First Iterative Depening



ABEE

Both DFS & BFS are used here.

$d=0$ A
 $d=1$ A B C D
 $d=2$ A B E C F G D H
 $d=3$ A B E I C F J



$d=0$ \$S
 $d=1$ S A C
 $d=2$ S A D B C E G
 $d=3$ S A D F H B C E I

Taskwork - 1

Initial stack

~~task~~ ['a', 'b', 'c']

Elements popped from stack

~~c~~

Stack

stack = []

stack.append('a')

stack.append('b')

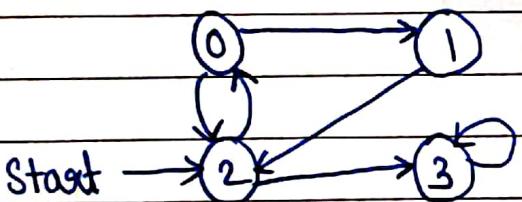
stack.append('c')

print("Initial stack")

print(stack)

print("In Elements popped from stack")
print(stack.pop())
print(stack.pop())
print(stack.pop())
print("Stack after elements are popped")
print(stack)

Graph



BFS

```
from collections import defaultdict
```

```
class Graph:
```

```
    def __init__(self):  
        self.graph = defaultdict(list)
```

```
    def addEdge(self, u, v):  
        self.graph[u].append(v)
```

```
    def BFS(self, s):
```

```
        visited = [False] * (max(self.graph) + 1)
```

```
        queue = []
```

```
        queue.append(s)
```

```
        visited[s] = True
```

while queue:

$s = \text{queue.pop(0)}$

```
print(s, end=" ")
```

for i in self.graph[s]:

if visited[i] == False :

queue.append(i)

`visited[i] = True`

```
g = Graph()
g.addEdge(0, 1)
g.addEdge(0, 2)
g.addEdge(1, 2)
g.addEdge(2, 0)
g.addEdge(2, 3)
g.addEdge(3, 3)
```

print ("Following is Breadth First Traversal (starting from vertex 2) ")

g. BFS (2)

Output

BFS: 2031

$$\overleftrightarrow{2} | 03 | 1$$

2031

1

Ex DFS:

3
1
0
2

31002

2013

DFS

```
from collections import defaultdict
```

```
class Graph:
```

```
    def __init__(self):
```

```
        self.graph = defaultdict(list)
```

```
    def addEdge(self, u, v):
```

```
        self.graph[u].append(v)
```

```
    def DFSUtil(self, v, visited):
```

```
        visited.add(v)
```

```
        print(v, end='')
```

```
        for neighbour in self.graph[v]:
```

```
            if neighbour not in visited:
```

```
                self.DFSUtil(neighbour, visited)
```

```
    def DFS(self, v):
```

```
        visited = set()
```

```
        self.DFSUtil(v, visited)
```

```
g = Graph()
```

```
g.addEdge(0, 1)
```

```
g.addEdge(0, 2)
```

```
g.addEdge(1, 2)
```

```
g.addEdge(2, 0)
```

```
g.addEdge(2, 3)
```

```
g.addEdge(3, 3)
```

point ("Following is DFS from (starting from vertex 2))
g. DFS(2).

Search Methodologies

Search is a method that can be used by computers to examine a problem space in order to find a goal.



Data driven Approach

- Also called forward chaining.
- Top down
- Goal state not known.
- Eg : Colouring of map

Goal driven Approach

- Also called backward chaining.
- Bottom up
- Goal state known (start from goal & reach initial)
- Eg: Maze prob (start from exit will be easier)

* Goal driven search is particularly useful in situations in which the goal can be clearly specified.

Eg: A theorem that is to be proved.

* Data driven search is most useful when the initial data is provided & it is not clear what the goal is.

Eg: Analyzing astronomical data.

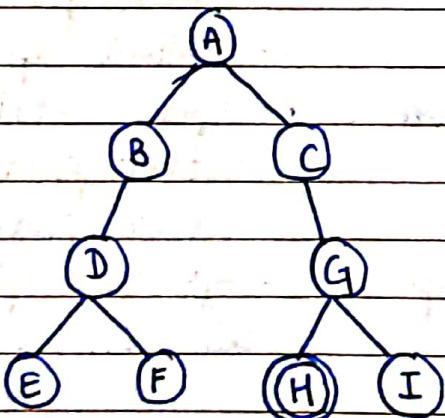
No technique, simply search (Try & error)

Generate & Test → Blind Search → Brute Force Search
Generator

3 properties

- It should be complete (all possible solns must be present &/produced)
- It should be non redundant.
- It must be well informed. (it should not advice any solⁿ that does not match the pbm st).

DFS & BFS



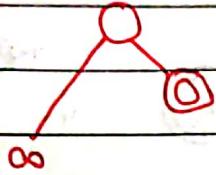
DFS: A B D E F C G H

DFS: E D F B C G H A

Properties of Search Methods

- 1) Complexity - Time & Space
- 2) Completeness
- 3) Optimality
- 4) Admissibility
- 5) Irreversability - explore only one path (no backtracking)

BFS is considered as complete search method but
DFS is not " " incomplete " " /tree.



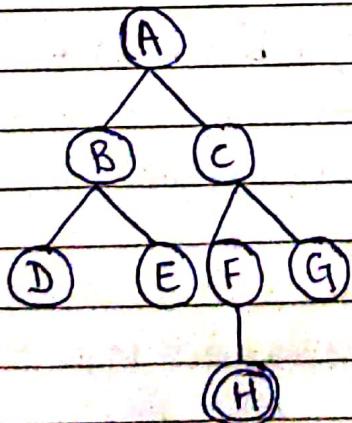
→ In this case DFS will never reach goal state.

BFS : Branch factor is the no. of children for a node.

The search method that finds optimal solⁿ in quickest time is called admissible search time.

Search methods for which backtracking is allowed is called tentative methods.

DFID → Depth First Iterative Depening

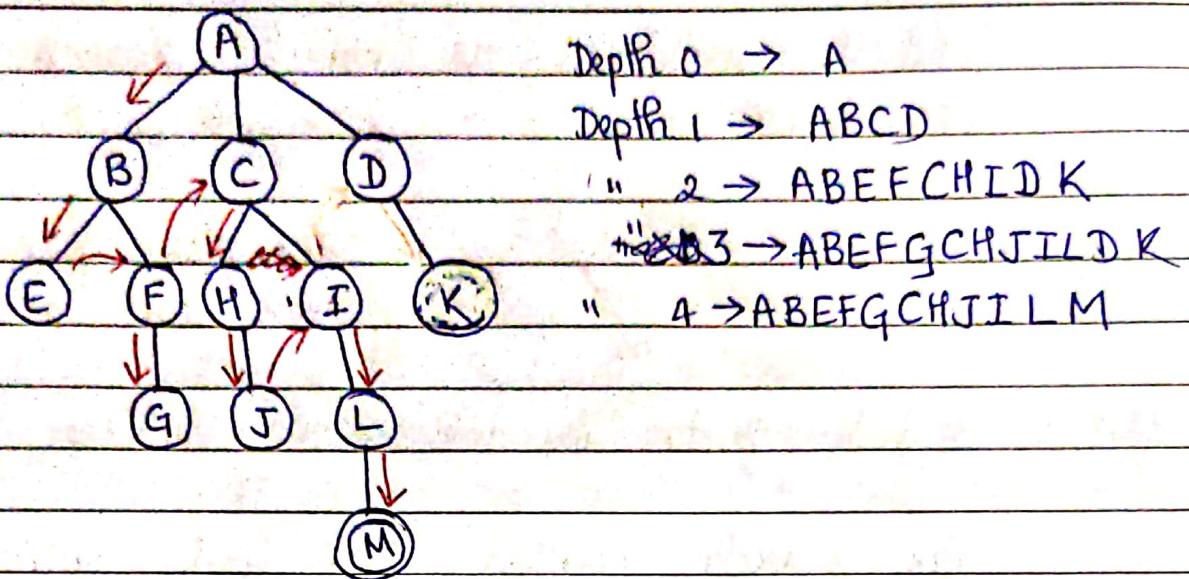


Depth 0 → A

Depth 1 → ABC

Depth 2 → ABDECFG

Depth 3 → ABDECFH



Depth d & branching factor b

$$1 + b + b^2 + b^3 + \dots + b^d = \frac{1 - b^{d+1}}{1 - b} \text{ nodes}$$

$$\text{Eg: } d=2, b=2 \therefore \frac{1 - 2^3}{1 - 2} = \frac{-7}{-1} = 7 \text{ nodes}$$

DFID ↗

$$(d+1)1 + b(d) + b^2(d-1) + b^3(d-2) + \dots + b^d$$

↗ No. of times a node is visited.

In DFID, leaf nodes are visited only once, & other nodes will be explored multiple times.

DFS & BFS

$$1 + b + b^2 + \dots + b^d$$

For smaller trees DFID is inefficient.

g) $d = 5 \quad b = 10$

\rightarrow DFS & BFS:

$$\text{Nodes} = \frac{1 - 10^6}{1 - 10} = \frac{1 - 1000000}{1 - 10} = \frac{999999}{9} = 111111$$

DFID

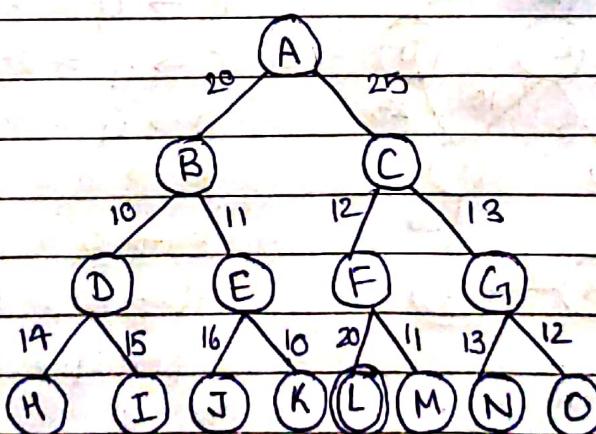
$$(5+1) + 10(5) + 10^2(4) + 10^3(3) + 10^4(2) + 10^5 \\ = 6 + 50 + 400 + 3000 + 20000 + 100000 \\ = 123456$$

IMP

Time complexity of DFID & BFS = $O(b^d)$

Space " " DFID & DFS = $O(bd)$

Best First Search



A \rightarrow L \rightarrow 50

B \rightarrow L \rightarrow 40

C \rightarrow L \rightarrow 30

D \rightarrow L \rightarrow 20

E \rightarrow L \rightarrow 18

F \rightarrow L \rightarrow 15

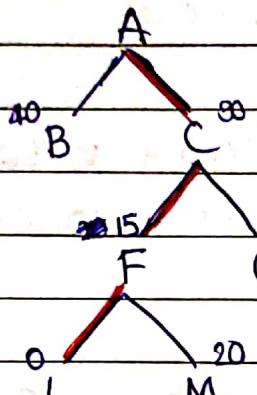
Heuristic

G \rightarrow L \rightarrow 20

L \rightarrow L \rightarrow 0

M \rightarrow L \rightarrow 20

We need 2 queues (open queue, closed queue)
 Priority queue based queue



Open queue

A B C

CB

FGB

FGB

LMGB

KMGB

MGB

Closed queue

A

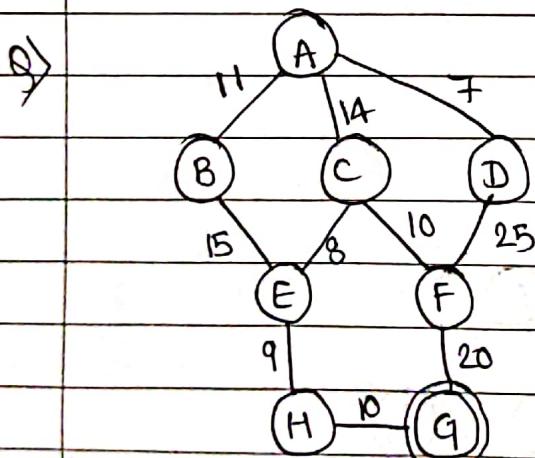
AC

AC

AC F

ACFL

ACFL



$$A \rightarrow G = 40$$

$$B \rightarrow G = 32$$

$$C \rightarrow G = 25$$

$$D \rightarrow G = 35$$

$$F \rightarrow G = 17$$

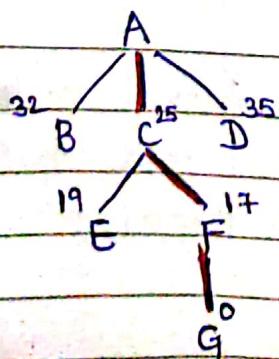
$$E \rightarrow G = 19$$

$$H \rightarrow G = 10$$

$$G \rightarrow G = 0$$

BDCBA

CBDA



Open queue

A

ABCD

EBD

FEBD

GEBD

Closed queue

A

AC

ACF

ACFG

Implementation of open queue

Function best()

{

queue = [];

state = root-node;

while (true)

{ if is-goal (state)

then return SUCCESS;

use

{

add-to-front-queue (successor (state));

sort (queue);

}

if queue == []

then return FAILURE;

state = queue[0];

remove-from-front (queue);

}

}

Tracing

State = A

queue → B C D

sort → C B D

for

State = C

queue → E F .

sort → F E .

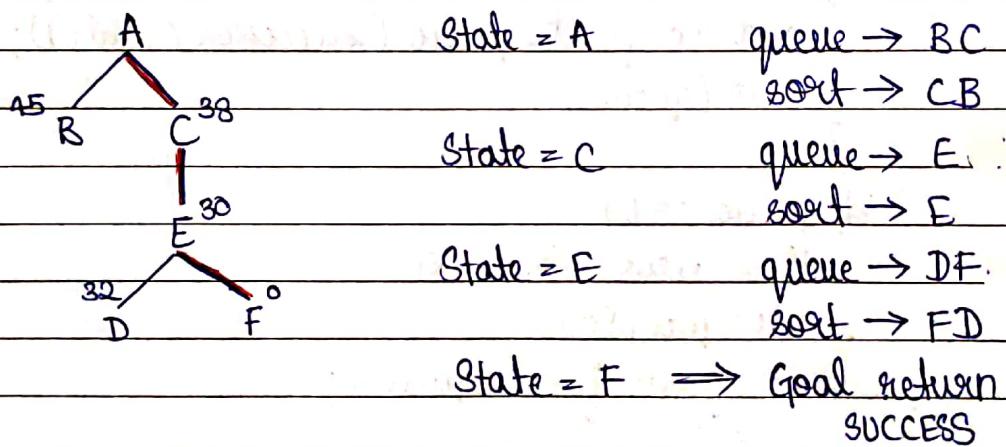
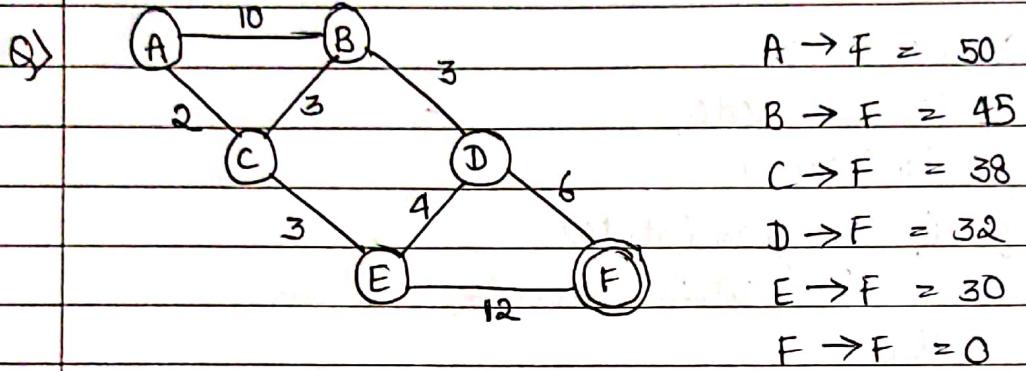
State = F

queue → G .

sort → G .

State = G → Goal. return SUCCESS

Text book pbm



The node with highest Heuristic value is the start node.

Advantages

Gives solⁿ quicker

Disadvantage

Does not give best soln

TC of Best First Search < TC of BFS, DFS, DFID