

# Data Structures

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# Outline

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- 5 (Optional) Binary Tree

# Today's Goal

- Understand how to create dynamically allocated data structures e.g. *linked list* and *binary tree*.
- Practice allocating and deallocating memory in the data structures.

# Data Structures

- Which data type would you use if you have to store information of students who registered to a course?
  - Array? You may need more...
- What kind of things you need to think about when choosing a data structure?
  - Performance? Memory usage?

# Linked list

- Linked list is a linear data structure, where each node contains the address of (i.e. pointer to) the next node.
- Many variations exist, including doubly linked list, or circular list.

```
struct Node {  
    int val;  
    Node* next;  
}
```

# Linked List Find

- Find node  $N$  with value  $v$  in the linked list  $\mathcal{L}$ .
  - Let  $iter$  be the address of  $\mathcal{L}$ .
  - Advance  $iter$  (i.e.  $iter = iter \rightarrow next$ ) until a node with value  $v$  is found.

# Linked List Insertion

- Insert a node next to a given node  $N$ .
  - ① Store address of the next node of  $N$  into  $Old$ .
  - ② Create a new object  $New$  and its next node becomes  $Old$ .
  - ③ Substitute  $New$  for  $Old$  as the next node of  $N$ .

```
void insertNext(Node* node , int val)
{
    auto oldNext = node->next;
    auto newNext = new Node{val , oldNext};
    node->next = newNext;
}
```

# Linked List Deletion

- Delete a given node  $N$  from the linked list  $\mathcal{L}$ .
  - ① Find  $P$  in the  $\mathcal{L}$ , whose next node is  $N$ .
  - ② Set the next node of  $P$  to be the next node of  $N$ .
  - ③ Deallocate memory pointed by  $N$ .

```
bool deleteNode(Node* head , Node* node)
{
    auto p = prevNode(head , node);
    if (!p)
        return false;
    p->next = node->next;
    delete node;
    return true;
}
```

- We may do better if we keep the pointer to the previous node!



# Stack

- *Stack* has limited interfaces `push(stack, value)` and `pop(stack)`.
- `push()` stores value to the stack.
- `pop()` removes a value from the stack.
- Latest value added to the stack will be popped first. (LIFO)
- The underlying implementation may use either an array or a linked list.

```
struct Stack { ... };  
void push(Stack* stack, int val);  
int pop(Stack* stack);
```

# Queue

- *Queue* has limited interfaces `enqueue(q, v)` and `dequeue(q)`
- `enqueue()` stores value to the queue.
- `dequeue()` removes a value from the queue.
- Oldest value added to the queue will be dequeued first. (FIFO)
- The underlying implementation may use either an array, a linked list or two stacks. (How?)

```
struct Queue { ... };  
void enqueue(Queue* queue, int val);  
int dequeue(Queue* queue);
```

# (Optional) Binary Tree

- Linked list has a linear structure. In many cases we should perform linear search.
- Linear search is not bad. Can we do better? Use *Binary Trees*!
- Each node in the binary tree has two children nodes.

```
struct Node {
    int val;
    Node *left, *right;
};
```

- A binary called a *Binary Search Tree*, if it satisfies this property:
  - $\forall N \in \mathcal{T}, N \rightarrow \text{left} \rightarrow \text{value} \leq N \rightarrow \text{value}$  and  $N \rightarrow \text{value} \leq N \rightarrow \text{right} \rightarrow \text{value}$

# (Optional) Binary Search Tree Insertion

```

void insertLeaf(Node* root, int val)
{
    auto iter = root;
    auto leaf = new Node{val, nullptr, nullptr};
    while(true) {
        if (val < iter->val) {
            if (!iter->left)
                iter->left = leaf; break;
            iter = iter->left;
        } else {
            if (!iter->right)
                iter->right = leaf; break;
            iter = iter->right;
        }
    }
}

```