Verified Multisig Contract

Raphaël Cauderlier

Nomadic Labs

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Introduction

- Raphaël Cauderlier
- Nomadic Labs
- Verification of Tezos Smart Contracts in Coq

Outline

- Introduction
- Pormal Methods for smart contracts
- Multisig Contracts
- 4 Mi-Cho-Coq
- Conclusion

Motivation

Introduction

- Smart contracts manipulate money (sometimes a lot)
- They are here to stay: in case of bug, they are hard to update
- Security: bugs may become exploits

Before uploading them, we want to be sure there is no bug in them!

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Before uploading them, we want to be sure there is no bug in them!

 infinitely-many possible input values so testing cannot be exhaustive

Formal methods: methods for mathematically reasoning about programs

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- Semantics: Description of the meaning of all instructions of the programming language
- Specification: Formula in some logic describing the expected behaviour of the program
- Goal: verify that the program satisfies the specification
 - Mathematical proof, more or less automatized

Approaches

- Model Checking
 Abstract the program into a state automaton called the model that can be checked on all inputs.
- Abstract Interpretation
 Abstract the values as domains (for example intervals). Refine the abstraction when needed.
- Deductive Verification
 Reduce to the theorem proving problem.

Model Checking

Introduction

Abstract the program into a state automaton called the *model* that can be checked on all inputs.

- Specifications:
 - Safety No bad state can be reached
 - Liveness Good states are reached infinitely often
 - Temporal properties
- Problem:
 - Finding the model
 - Linking it to the concrete program

Abstract Interpretation

Introduction

Abstract the values as *domains* (for example intervals). Refine the abstraction when needed.

- Specifications:
 - Safety
 - Arithmetic
- Problem
 - False alarms

Deductive Verification

Introduction

Reduce to the theorem proving problem.

- Specifications:
 - Functional properties { precondition } Program { postcondition }
 - Very rich logics
- Problem
 - Requires a lot of user interaction

Introduction

Michelson has been designed to ease formal methods

- Static typing
- Explicit failure
- No overflow nor division by zero
- Clear documented semantics

Michelson contracts are necessarily small and simple

Formal methods for Michelson

• Model Checking:

- Example: auction
- Spec: Anybody either win the auction or lose no money
- Tool: Cubicle Model-Checker
- Abstract Interpretation:
 - Bound on gas
 - Token freeze
- Deductive Verification:
 - Example: multisig
 - Spec: multisig succeeds IFF enough valid signatures
 - Tool: Mi-Cho-Cog

The multisig contract

- *n* persons share the ownership of the contract.
- they agree on a threshold t (an integer).
- to do anything with the contract, at least t owners must agree.
- possible actions:
 - list of operations (to be run atomically)
 - changing the list of owners and the threshold

The multisig contract

Introduction

Michelson implementation:

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- Integrated in the tezos-client
 - tezos-client deploy multisig ...
 - tezos-client sign multisig transaction ...
 - tezos-client from multisig contract ... transfer ...

Multisig anti-replay

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 - authenticate the owners
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- Signed data = (action, counter, multisig address)
 Counter incremented at each successful run

Multisig storage

Introduction

We need to store

- the keys
- the threshold
- the anti-replay counter

Multisig parameter

Introduction

Two entrypoints

- Default: take my tokens (by anybody)
- Main: Perform an action (requires enough signatures)

Multisig code

https://github.com/murbard/smart-contracts/blob/master/multisig/michelson/generic.tz

Introduction

The Coq interactive theorem prover

- Developped for more than 30 years
- Non-trivial mathematical theorems: 4-color, odd-order
- CompCert: certified C compiler

Parameter A : Type.

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Parameter B : A -> Prop.

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Parameter f : forall x : A, B x.

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Parameter A : Type.
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Parameter a : A.

Parameter B : A -> Prop.

Parameter f : forall x : A, B x.

Check f a. (* Answer: f a : B a *)

Coq: Implicit arguments

Explicit polymorphism: Types are regular terms

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Lemma id_2 : id 2 = 2.

Proof. reflexivity. Qed.
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Coq: Inductive Types

Generalisation of ADT to dependent types:

```
Inductive vector (A : Type) : nat -> Type :=
   | Nil : vector A 0
   | Cons n : A -> vector A n -> vector A (1 + n).
```

Mi-Cho-Coq

Coq formalisation of Michelson

- syntax
- semantics
- typing

Functional verification of a few contracts

Verification of the Multisig

Proven:

- Characterisation of the relation between input and output: eval multisig input = Success output <-> R input output
- Not proven: Security property:
 - signatures cannot be forged
 - signatures sent to a multisig cannot be replayed

Future work

- Prove more contracts
- Improve automation
- Formalise more of Michelson (gas, contract life, security)
- Higher-level languages with certified compilers
- Extract the Coq interpreter

Conclusion

- Formal methods are complementary.
- Language design matters.