# SP-51 C-10: Code Library



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## 1基本配置

#### 1.1 快速讀寫

```
struct FastIS {
 static const int LEN = 1<<21 | 1;
 char buf[LEN], *l, *r, c, *t; bool neof, neg;
 FastIS(): l(buf), r(buf), neof(1) {}
 inline operator bool() { return neof; }
 inline char gc() {
   return
      !neof || (l=r & (r = (l=buf) + fread(buf, 1, LEN, stdin), l=r)) ?
      neof=0. -1 : *l++:
 FastIS& operator>>(char &c) {
    for (c=gc(); isspace(c); c=gc()) {} return *this;
 FastIS& operator>>(char *s) {
    for (c=gc(); isspace(c); c=gc()) {}
    for (t=s; ~c and !isspace(c); c=gc()) *t++=c;
    *t='\0'; return *this;
 template<typename T> FastIS& operator>>>(T &x) {
    for (c=gc(); \sim c \text{ and } ! isdigit(c); c=gc()) \text{ neg } ^= c = '-';
    for (x=0; isdigit(c); c=gc()) x = 10*x+c-'0';
    if (neg) { x=-x, neg=0; } return *this;
} fast in;
struct FastOS {
 static const int LEN=1<<21 | 1, W=21;
 char buf[LEN], *l, *r, *t, stk[W]; int top;
 FastOS() : l(buf), r(buf+LEN) {}
 inline void flush() { fwrite(buf, 1, l-buf, stdout); l=buf; }
 FastOS& operator<<(char c) { *l++=c; if (l=r) flush(); return *this; }
 FastOS& operator << (char *s) {
    for (t=s; *t≠'\0'; ++t) { *this<<*t; } return *this;
 FastOS& operator<<(const char *s) { return *this<<(char*)s; }
 template<typename T> FastOS& operator<<(T x) {</pre>
   if (!x) *this<<'0'; else if (x<0) *this<<'-', x=-x;
    for (top=0; x; x \ne 10) stk[++top] = x\%10+'0';
    while (top) { *this<<stk[top--]; } return *this;</pre>
```

```
~FastOS() { flush(); }
} fast_out;
1.2 Montogomery 乘法
using u64 = unsigned long long;
using u32 = unsigned int;
using i64 = long long;
struct u128 {
  u64 h, l;
  static u128 mul(const u64 &x, const u64 &y) {
    u32 xh = x>>32, yh = y>>32, xl = x, yl = y;
    u64 hh = static cast<u64>(xh) * yh, ll = static cast<u64>(xl) * yl;
    u64 hl = static_cast<u64>(xh) * yl, lh = static_cast<u64>(xl) * yh;
    u64 c = (h1 << 32 >> 32) + (lh << 32 >> 32) + (ll >> 32);
    u64 h = hh + (hl>>32) + (lh>>32) + (c>>32);
    u64 l = (hl << 32) + (lh << 32) + ll;
    return {h, l};
};
struct Montgomery {
  u64 mod, inv, r2;
  u64 reduce(u128 x) {
    u64 q = x.l * inv;
                                       // q \equiv x / n \pmod{r}
    i64 a = x.h - u128::mul(q, mod).h; // x - qn \equiv 0 (mod r)
    return a<0 ? a+=mod : a;</pre>
  u64 mul(u64 x, u64 y) { return reduce(u128::mul(x, y)); }
  Montgomery(u64 mod) : mod(mod), inv(1), r2(-mod % mod) {
    // 2 ^ (2 ^ 6) = 2 ^ 64
    for (int i=0; i<6; ++i) inv *= 2 - mod * inv;
    // 2 ^ (4 * 2 ^ 4) = 2 ^ 64
    for (int i=0; i<4; ++i) if ((r2 \Leftarrow 1) \geqslant mod) r2 -= mod;
    for (int i=0; i<4; ++i) r2 = mul(r2, r2);
  inline u64 init(u64 x) { return mul(x, r2); }
  inline u64 reduce(u64 x) { return reduce({0, x}); }
};
1.3 暫用 Treap 隨機數生成器
struct rot xor {
  int seed = 0×95abcfad;
  inline unsigned int operator()() {
```

return seed = (seed  $\ll$  1)  $^{\circ}$  ((seed  $\gg$  31) & 0×a53a9be9);

```
}
};
```

## 2字符串

#### 2.1 Z 算法

```
void z_func(char s[], int n, int z[]) {
  int l = 0, r = 0; z[0] = 0;
  for (int i=1; i<n; ++i) {
    z[i] = i < r ? min(z[i-l], r-i) : 0;
    while (i+z[i] < n & s[i+z[i]] = s[z[i]]) ++z[i];
    if (chkmx(r, i+z[i])) l = i;
  }
}</pre>
```

#### 2.2 Manacher

```
int manacher_init(char in[], int n, char out[]) {
   int m = (n+1)<<1;
   out[0] = out[1] = '#'; out[m+1] = '@';
   for (int i=1; i ≤ n; ++i) out[i<<1] = in[i-1], out[i<<1|1] = '#';
   return m;
}

void manacher_work(char s[], int n, int p[]) {
   int mid = 0, r = 0; p[0] = 1;
   for (int i=1; i<n; ++i) {
     p[i] = i < r ? min(p[2*pos-i], r-i) : 1;
     while (s[i-p[i]] = s[i+p[i]]) ++p[i];
     if (chkmx(r, i+p[i])) mid = i;
}</pre>
```

#### 2.3 迴文自動機

slink 指向沿着 fail 邊長度不等差的第一個後綴。 trans 指向長度小於等於當前迴文串一半的最長後綴。

```
namespace PAM {
  int len[MAX_N], ch[MAX_N][26], fail[MAX_N];
  int diff[MAX_N], slink[MAX_N];
  int trans[MAX_N];
  int dep[MAX_N];
```

```
int indx. last:
  char s[MAX_N];
  int new node(int l) {
   len[indx] = l;
    memset(ch[indx], 0, sizeof(ch[indx])); return indx++;
  void init() {
   indx = 0; new_node(0); new_node(-1);
    s[0] = '#'; fail[0] = 1; dep[0] = dep[1] = 0; diff[0] = 0; last = 0;
  int getfail(int p, int i) {
    while (s[i-len[p]-1] \neq s[i]) p = fail[p];
    return p;
  void insert(int i) {
    int p = getfail(last, i);
    if (!ch[p][s[i]-'a']) {
     int q = new node(len[p]+2);
      fail[q] = ch[getfail(fail[p], i)][s[i]-'a'];
      dep[q] = dep[fail[q]] + 1;
      diff[q] = len[q]-len[fail[q]];
      slink[q] = diff[q]=diff[fail[q]] ? slink[fail[q]] : fail[q];
      if (len[q] \leq 2) {
       trans[q] = fail[q];
      } else {
        int r = trans[p];
        while (s[i-len[r]-1] \neq s[i] \mid (len[r]+2)*2 > len[q]) r = fail[r];
        trans[q] = ch[r][s[i]-'a'];
      ch[p][s[i]-'a'] = q;
    last = ch[p][s[i]-'a'];
} // namespace PAM
2.4 後綴自動機
namespace SAM {
int len[MAX_2N], ch[MAX_2N][26], fail[MAX_2N];
int cnt[MAX 2N];
int indx, last;
int new node(int l) {
 len[indx] = l; cnt[indx] = 1;
 memset(ch[indx], -1, sizeof(ch[indx])); return indx++;
int copy_node(int q, int l) {
```

```
len[indx] = l; cnt[indx] = 0; fail[indx] = fail[q];
 memcpy(ch[indx], ch[q], sizeof(ch[q])); return indx++;
void init() {
 indx = 0; new node(0);
 fail[0] = -1; last = 0;
void insert(int w) {
 int cur = new_node(len[last] + 1);
 int p = last;
 for (; \sim p \& ch[p][w] = -1; p = fail[p]) ch[p][w] = cur;
 if (p = -1) {
   fail[cur] = 0;
 else if (int q = ch[p][w]; len[q] = len[p] + 1) {
   fail[cur] = q;
 } else {
    int r = copy_node(q, len[p] + 1);
    for (; \sim p \ 66 \ ch[p][w] = q; p = fail[p]) \ ch[p][w] = r;
    fail[cur] = fail[q] = r;
 last = cur;
} // namespace SAM
```

#### 2.5 Lyndon 分解

#### pms 前綴的最小後綴。

```
void duval(const char s[], int n, int pms[]) {
  int i = 0;
  while (i < n) {
    pms[i] = i;
  int j = i, k = i+1;
  while (k < n &6 s[j] ≤ s[k]) {
    if (s[j] < s[k]) {
       pms[k] = i;
       j = i, ++k;
    } else {
       pms[k] = pms[j] + (k - j);
       ++j, ++k;
    }
  }
  while (i ≤ j) {
       // handle [i, i + (k - j))
       i += k - j;
  }
}</pre>
```

}

### 3 圖論

#### 3.1 樹上啓發式合併

```
void dsu(int u) {
  for (auto &v: e[u]) if (v ≠ fa[u] && v ≠ hson[u]) {
    dsu(v);
    for (int i=dfn[v]; i≤rdfn[v]; ++i) del(dfn[i]);
  }
  if (hson[u]) dfs(hson[u]);
  for (auto &v: e[u]) if (v ≠ fa[u] && v ≠ hson[u]) {
    for (int i=dfn[v]; i≤rdfn[v]; ++i) add(dfn[i]);
  }
}
```

## 4 數學

#### 4.1 線性基

qry\_kth 包括零,下標從零開始的第 k 大。只關注 k 的後 dim 位。 qry\_rnk 返回小於 x 的向量個數(包括零),eq 表示能否表示出 x。 涉及刪除元素時可以考慮兩種方案:

- · 記錄每個基的過期時間 exp。
- 記錄每個向量用基的表示方法。

```
template<int w> struct LinearBasis {
  bitset<w> base[w]; int dim;
  LinearBasis() : dim(0) { for (int i=0; i<w; ++i) base[i].reset(); }
  int insert(bitset<w> x) {
   for (int i=w-1; i ≥ 0; --i) if (x[i]) {
      if (base[i][i]) {
        x ^= base[i];
      } else {
        // qry_kth requires this:
```

```
for (int j=i-1; j \ge 0; --j) if (x[j] \& base[j][j]) \times ^= base[j];
        for (int j=i+1; j < w; ++j) if (base[j][i]) base[j] ^= x;
        base[i] = x;
        ++dim; return i;
    return -1;
  bitset<w> qry_kth(bitset<w> k) {
    bitset<w> res;
    for (int i=0; i<w; ++i) if (base[i][i]) {</pre>
     if (k[0]) res ^= base[i];
     k >≔ 1;
    } return res;
  bitset<w> qry rnk(bitset<w> x, bool &eq) {
    int df = dim;
    bitset<w> res, cur;
    for (int i=w-1; i \ge 0; --i) {
     if (base[i][i]) --df;
      if (x[i] & (!cur[i] | base[i][i])) res.set(df);
      if (x[i] \neq cur[i]) {
       if (base[i][i]) cur ^= base[i];
        else break;
    eq = cur = x;
    return res;
};
```

## 5 數據結構

#### 5.1 左偏堆

二路分治進行合併操作可實現線性建堆。

```
namespace LeftistHeap {
int val[MAX_N], dist[MAX_N];
int lc[MAX_N], rc[MAX_N];

void push_up(int x) {
  if (dist[lc[x]] < dist[rc[x]]) swap(lc[x], rc[x]);</pre>
```

```
dist[x] = dist[rc[x]] + 1;
int merge(int x, int y) {
 if (!x || !y) return x | y;
 if (val[x] > val[y]) swap(x, y);
 rc[x] = merge(rc[x], y);
  push_up(x); return x;
} // namespace LeftistHeap
5.2 笛卡爾樹
static int stk[MAX N];
int top = 0:
for (int i=1; i \le n; ++i) {
 int x = i, y = 0;
 while (top & h[i] < h[stk[top]]) y = stk[top--];</pre>
 lc[x] = v;
 rc[stk[top]] = x;
 stk[++top] = x;
return rc[0];
```

#### 5.3 FHQ Treap

可持續化時只需要在 split 時複製結點。 △ 未測試 △

```
namespace FHQTreap {
int rnd[MAX N], val[MAX N];
int siz[MAX N], fa[MAX N];
int lc[MAX_N], rc[MAX_N];
int indx, root;
int new node(int v);
void push_up(int p);
void split(int p, int k, int &x, int &y) {
 if (!p) \{ x = 0; y = 0; \}
 if (int t = siz[p] - siz[rc[p]]; t \leq k) {
   x = p; split(rc[x], k-t, rc[x], y);
   fa[rc[x]] = x;
 } else {
   y = p; split(lc[y], k, x, lc[y]);
   fa[lc[y]] = y;
 } push up(p);
```

```
int merge(int x, int y) {
   if (!x || !y) return x | y;
   if (rnd[x] < rnd[y]) {
     rc[x] = merge(rc[x], y);
     push_up(x); return fa[rc[x]] = x;
   } else {
     lc[y] = merge(x, lc[y]);
     push_up(y); return fa[lc[y]] = y;
   }
}
// namespace FHQTreap</pre>
```

#### 5.4 Splay

#### △ 未測試 △

```
namespace Splay {
#define lc(x) (ch[(x)][0])
#define rc(x) (ch[(x)][1])
#define isrc(x) (rc(fa[(x)]) = (x))
int val[MAX_N], cnt[MAX_N], siz[MAX_N];
int fa[MAX_N], ch[MAX_N][2];
int indx, root;
int new_node(int v);
void push up(int x);
void rotate(int x) {
 int y = fa[x], z = fa[y], k = isrc(x), w = ch[x][!k];
 if (y) \{ ch[z][isrc(y)] = x; \} ch[x][!k] = y; ch[y][k] = w;
 if (w) { fa[w] = y; } fa[y] = x; fa[x] = z;
 push_up(y);
void splay(int x) {
 for (int y; y = fa[x], x; rotate(x))
   if (y) rotate(isrc(x)^isrc(y) ? x : y);
 push_up(x);
 root = x;
void upd_ins(int v) {
 if (!root) return root = new node(v);
 int p = root;
 while (1) {
    if (val[p] = v) {
     ++cnt[p]; return splay(p);
    } else {
     int k = v > val[p];
```

```
if (!ch[p][k]) {
        int q = new_node(v);
       ch[p][k] = q; fa[q] = p;
       return splay(q);
     } else p = ch[p][k];
 }
int qry_find(int v);
int qry pre(int x);
void upd del(int v) {
 int p = qry_find(v);
 if (!p) return;
 if (cnt[p] > 1) --cnt[p], --siz[p];
 else if (!lc(p) \& !rc(p)) root = 0;
 else if (!lc(p) || !rc(p)) fa[root = lc(p) | rc(p)] = 0;
 else {
   fa[lc(p)] = 0;
   qry pre(p);
   fa[rc(root) = rc(p)] = root;
   push up(root);
#undef lc
#undef rc
#undef isrc
} // namespace Splay
```

#### 5.5 Link Cut Tree

#### ⚠ 未測試 △

```
namespace LCT {
#define lc(x) (ch[(x)][0])
#define rc(x) (ch[(x)][1])
#define isrc(x) (rc(fa[(x)]) = (x))
#define nroot(x) (lc(fa[(x)]) = (x) || rc(fa[(x)]) = (x))
int fa[MAX_N], ch[MAX_N][2];
bool rev_tag[MAX_N];

void push_rev(int x);
void push_dn(int x);
void push_up(int x);
void rotate(int x) {
   int y = fa[x], z = fa[y], k = isrc(x), w = ch[x][!k];
   if (nroot(y)) { ch[z][isrc(y)] = x; } ch[x][!k] = y; ch[y][k] = w;
   if (w) { fa[w] = y; } fa[y] = x; fa[x] = z;
```

```
push_up(y);
void splay(int x) {
 static int stk[MAX N], top;
 for (int y = stk[top = 1] = x; nroot(y); stk[++top] = (y = fa[y]));
 while (top) push_dn(stk[top--]);
 for (int y; y = fa[x], nroot(x); rotate(x))
   if (nroot(y)) rotate(isrc(x)^isrc(y) ? x : y);
 push_up(x);
int access(int x) {
 int y = 0;
 for(; x; x = fa[y=x]) splay(x), rc(x) = y, push_up(x);
 return y;
void evert(int x) { push_rev(access(x)); }
int split(int x,int y) { evert(x); return access(y); }
int find(int x) {
 for (x = access(x); lc(x); x = lc(x)) push_dn(x);
 return splay(x), x;
void link(int x,int y) { evert(x), splay(x); fa[x] = y; }
void cut(int x, int y) { split(x,y), splay(x); rc(x) = fa[y] = 0; push_up(x); }
bool chk_link(int x,int y) { evert(x); return find(y) \neq x ? fa[x] = y : 0; }
bool chk_cut(int x,int y){
 evert(x);
 if (find(y)\neq x || fa[y]\neq x || ch[y][0]) return 0;
 rc(x) = fa[y]=0; push_up(x); return 1;
#undef lc
#undef rc
#undef isrc
#undef nroot
} // namespace LCT
```