排序

希尔排序: Shell序列 + 输入敏感性

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Shell's Sequence

- ❖ Shell 1959: $\mathcal{H}_{shell} = \{1, 2, 4, 8, 16, 32, 64, ..., 2^k, ...\}$
- ❖ 实际上,采用 $\mathscr{A}_{\mathsf{shell}}$,在最坏情况下需要运行 $\Omega(n^2)$ 时间...
- **❖** 考查由子序列 A = unsort[0, 2^{N-1}) 和 B = unsort[2^{N-1}, 2^N) <mark>交错而成的序列</mark>

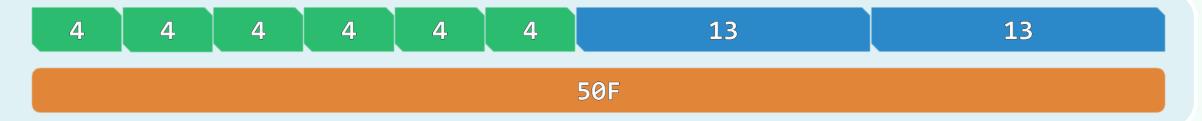


❖ 在做2-sorting时, A、B各成一列; 故此后必然各自有序

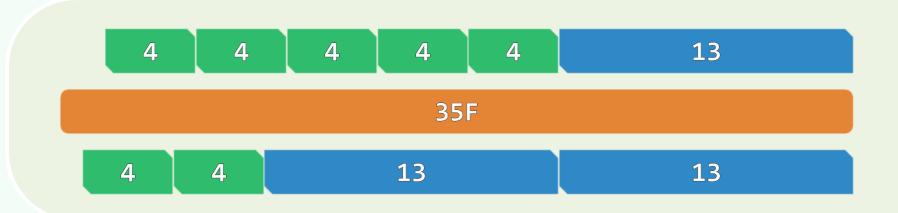
- * 然而其中的逆序对依然很多,最后的1-sorting仍需 $1+2+3+\cdots+2^{N-1}=\Omega(n^2/4)$ 时间
- ❖ 根源在于,淨_{shell}中各项并不互素,甚至相邻项也非互素

Postage Problem

❖ The postage for a letter is 50F, and a postcard 35F
But there are only stamps of 4F and 13F available



Possible to stamp
the letter and
the postcard
EXACTLY?



� Given a postage P, determine whether $P \in \{ n \cdot 4 + m \cdot 13 \mid n, m \in \mathcal{N} \}$

Linear Combination

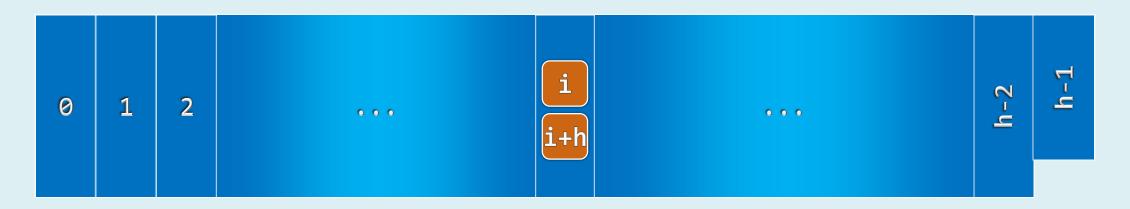
- \clubsuit Let $g,h\in\mathcal{N}$
- lacktriangledown For any $n,m\in\mathcal{N}$, $n\cdot g+m\cdot h$ is called a linear combination of g and h
- ❖ Denote $\mathbf{C}(g,h) = \{ ng + mh \mid n,m \in \mathcal{N} \}$ $\mathbf{N}(g,h) = \mathcal{N} \setminus \mathbf{C}(g,h) \quad \text{//numbers that are NOT combinations of g and h}$ $\mathbf{x}(g,h) = \max\{ \mathbf{N}(g,h) \} \quad \text{//always exists?}$
- \clubsuit Theorem: when g and h are relatively prime, we have

$$\mathbf{x}(g,h) = (g-1) \cdot (h-1) - 1 = gh - g - h$$

e.g.
$$\mathbf{x}(3,7) = 11$$
, $\mathbf{x}(4,9) = 23$, $\mathbf{x}(\boxed{4},\boxed{13}) = \boxed{35}$, $\mathbf{x}(5,14) = 51$

h-sorting & h-ordered

- **�** A sequence S[0,n) is called h-ordered if $S[i] \leq S[i+h], \; \forall \; 0 \leq i < n-h$
- ❖ A 1-ordered sequence is sorted
- ❖ h-sorting: an h-ordered sequence is obtained by
 - arranging S into a 2D matrix with h columns and
 - sorting each column respectively

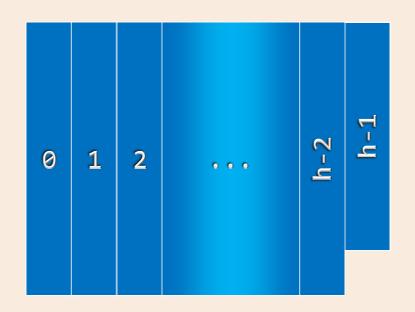


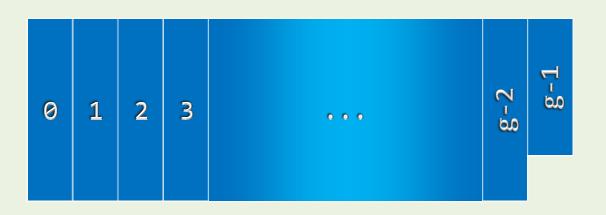
Theorem K

❖ [Knuth, ACP Vol.3 p.90]

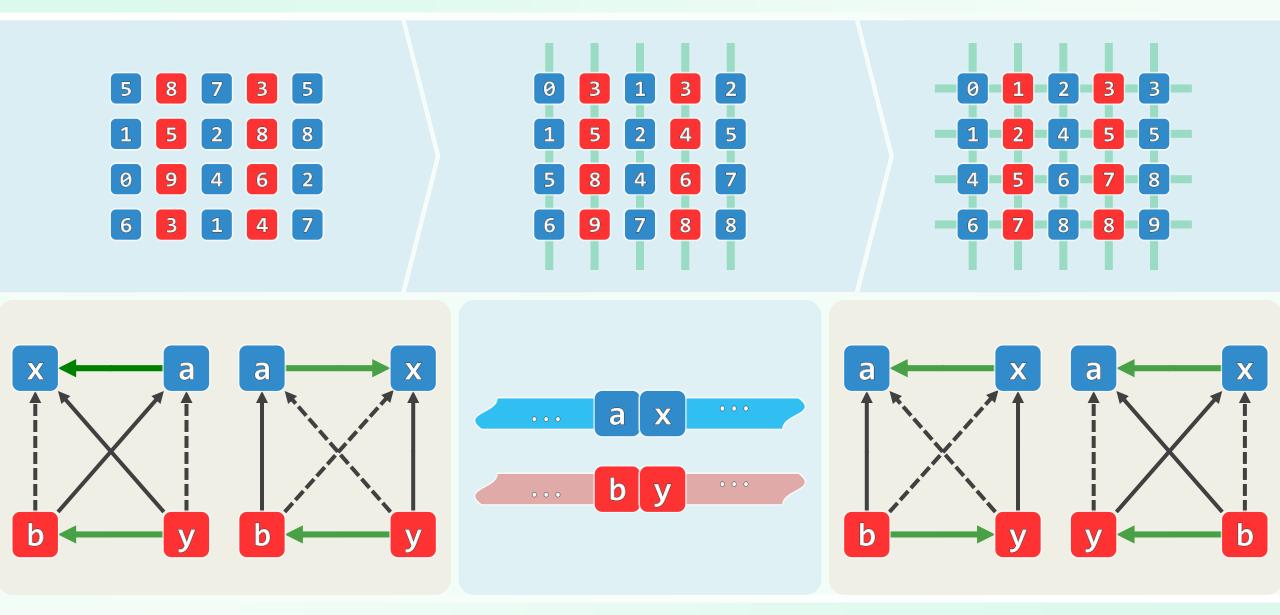
//习题解析[12-12, 12-13]

A g-ordered sequence REMAINS g-ordered after being h-sorted.

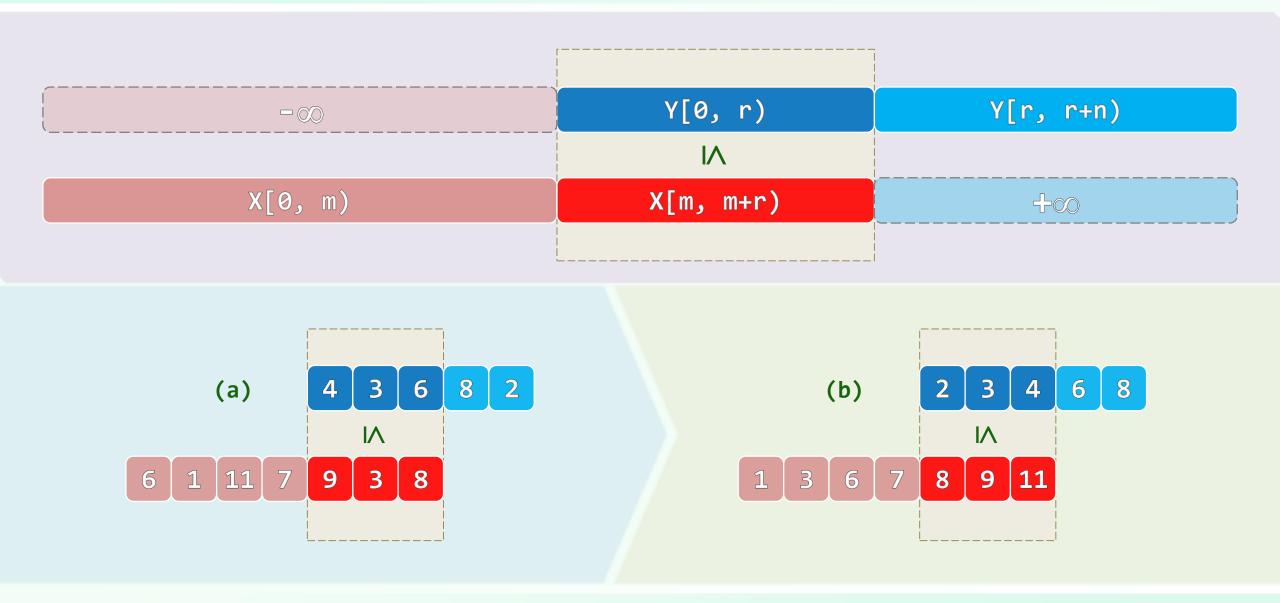




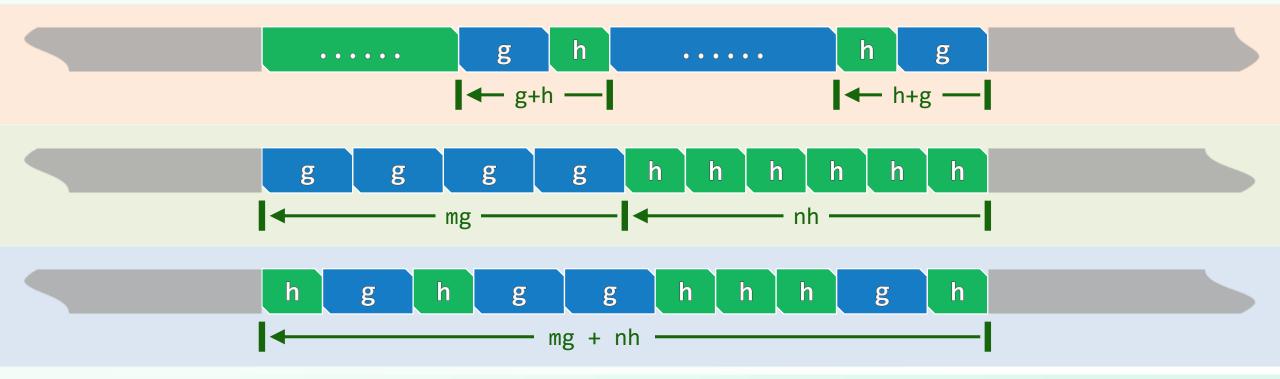
Order Preservation



Lemma L



Linear Combination



Inversion

- \clubsuit Let S[0,n) be a (g,h)-ordered sequence, where g and h are relatively prime
- ❖ Then for all elements S[j] and S[i], we have

$$i - j > \mathbf{x}(g, h)$$
 only if $S[j] \leq S[i]$

lacktriangledown This implies that to the LEFT of each element, only the previous $\mathbf{x}(g,h)$ elements could be GREATER

inversion free could be greater than S[i] [i] gh - g - h ———

lacktriangledown There would be no more than $n \cdot \mathbf{x}(g,h)$ INVERSIONs altogether