# CVX Implementation Notes

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## 1 Regularization

The quadratic approximation of the objective function f(x) at iterate  $x_k$  is

$$\tilde{f}(x_k + \Delta x) = f(x_k) + \nabla f(x_k)' \Delta x + \frac{1}{2} \Delta x' H_k \Delta x,$$

where  $H_k$  is the Hessian  $\nabla^2 f(x_k)$  or an approximation thereof. The search direction  $\Delta x$  from iterate  $x_k$  is computed by minimizing this function over the variable  $\Delta x$  resulting in the equation

$$H_k \Delta x = -\nabla f(x_k) := -y_k \tag{1}$$

for the search direction  $\Delta x$ . Clearly we can assume that  $y_k = \nabla f(x_k) \neq 0$  since the search terminates at a zero gradient.

In our algorithms we can also assume that  $H_k$  is positive semidefinite, but not necessarily nonsingular. If  $H_k$  is nonsingular (i.e. positive definite), then the solution  $\Delta x$  is a descent direction:

$$\Delta x' \nabla f(x_k) = -\Delta x' H_k \Delta x < 0$$

and this is all we care about. We will solve (1) by Cholesky factorization  $H_k = LL'$  with lower triangular L, by solving the triangular systems

$$Lv = -y_k, \quad L'\Delta x = v.$$
 (2)

The Cholesky factorization fails if  $H_k$  is singular. In that case we replace  $H_k$  with  $H_k(\delta) := H_k + \delta I$ , for some positive constant  $\delta$ . This matrix is now positive definite, in fact

$$(H_k(\delta)u, u) = u'H_k(\delta)u = u'H_ku + \delta u'Iu \ge \delta \|u\|^2$$

so that (1) now yields a descent direction. Moreover it improves the conditioning of (1): if  $H_k + \delta I = L(\delta)L(\delta)'$  is the Cholesky factorization of  $H_k(\delta)$ , then we have

$$|L(\delta)_{ii}| \geq \delta$$
.

Indeed, the diagonal element  $L(\delta)_{ii}$  is an eigenvalue of the triangular matrix  $L(\delta)'$ . Let u be a corresponding eigenvector. Then we have

$$|L(\delta)_{ii}|^2 ||u||^2 = ||L(\delta)'u||^2 = (L(\delta)L(\delta)'u, u) = ((H_k(\delta)u, u) \ge \delta ||u||^2$$

from which it follows that

$$|L(\delta)_{ii}| \ge \sqrt{\delta}$$

with obvious implications for the numerical stability of the triangular systems (2). Moreover this suggests that we should replace  $H_k$  with  $H_k + \delta I$  not only if the Cholesky factorization fails but rahter as soon as the minimal diagonal element (in absolute value) of the Cholesky factor L is below the threshold  $\sqrt{\delta}$ .

Trust region interpretation. The passage from the matrix  $H_k$  to the regularization  $H_k + \delta I$  has an interpretation in terms of *trust regions*: the solution  $\Delta x^*$  of

$$H_k(\delta)\Delta x = -y_k$$
, where  $y_k = \nabla f(x_k)$ ,

is the minimizer of the quadratic function

$$\phi(\Delta x) = f(x_k) + y_k' \Delta x + \Delta x' H_k(\delta) \Delta x$$
  
=  $f(x_k) + y_k' \Delta x + \Delta x' H_k \Delta x + \delta \|\Delta x\|^2$   
=  $\tilde{f}(x + \Delta x) + \delta \|\Delta x\|^2$ .

Now note that this minimizer  $\Delta x^*$  is automatically also the minimizer of the quadratic approximation  $\tilde{f}(x_k + \Delta x)$  on the ball  $B(x_k, r_k)$  with radius  $r_k = \|\Delta x^*\|$ . Indeed, if this ball contained a point u with  $\tilde{f}(x_k + u) < \tilde{f}(x_k + \Delta x^*)$ , then, since also  $\|u\| \leq \|\Delta x^*\|$  it follows that

$$\phi(u) = \tilde{f}(x_k + u) + \delta \|u\|^2 < \tilde{f}(x_k + \Delta x^*) + \delta \|\Delta x^*\|^2 = \phi(\Delta x^*).$$

In other words: passing from  $H_k$  to  $H_k + \delta I$  we compute the search direction  $\Delta x$  by minimizing the quadratic approximation  $\tilde{f}(x_k + \Delta x)$  not globally but instead on the ball  $B(x_k, r_k)$  (the region in which we trust the approximation) where the trust radius  $r_k$  is defined implicitly as  $r_k = ||\Delta x^*||$ .

This indicates that the regularization  $H_k \to H_k(\delta)$  is not unreasonable and in any case it solves the problem of nonsingularity of  $H_k$  for us, improves the conditioning and results in a descent direction  $\Delta x$  at iterate  $x_k$ .

### 2 Hessian

Let  $f: \mathbb{R}^n \to \mathbb{R}$  be a twice continuously differentiable function. The second order Taylor expansion of f centered at x has the form

$$f(x+h) = f(x) + L(h) + \frac{1}{2}B(h,h) + R(h),$$

where L is a linear function of  $h \in \mathbb{R}^n$ , B(u,v) is a bilinear function of  $(u,v) \in \mathbb{R}^n \times \mathbb{R}^n$  and the remainder R(h) satisfies  $R(h) = o(\|h\|^2)$ . This condition on the remainder ensures that L and B are uniquely determined as

$$L(h) = \nabla f(x)'h$$
 and  $B(u, v) = u'Hv$ ,

where  $H := \nabla^2 f(x) \in Mat_{n \times n}(\mathbb{R})$  is the matrix with entries

$$H_{ij} = B(e_i, e_j) = \frac{\partial^2 f}{\partial x_i \partial x_j}(x),$$

i.e. the Hessian matrix of f at x.

To compute the gradient and Hessian of a  $C^2$ -function f we make use of the fact that the remaider condition  $R(h) = o(\|h\|^2)$  in a quadratic expansion

$$f(x+h) = f(x) + \Delta' h + h' H h + R(h)$$

uniquely determines the "coefficients"  $\Delta$  and H as  $\Delta = \nabla f(x)$  and  $H = \nabla^2 f(x)$ . We only need to find such an expansion for f(x+h) and check that the remainder satisfies  $R(h) = o(\|h\|^2)$ . This is how we will derive our formulas below.

**Hessian of affine transformation.** Let  $f: \mathbb{R}^n \to \mathbb{R}$  be a  $C^2$ -function and  $\overline{f}(u) = f(x_0 + Fu)$ , where  $x_0 \in \mathbb{R}^n$  and  $F: \mathbb{R}^m \to \mathbb{R}^n$  is a linear map, i.e.  $F \in Mat_{n \times m}$ .

We want to compute the gradient and Hessian of h at any point u from those of f. To get these do a second order Taylor expansion of f about x:

$$f(x+h) = f(x) + h^{T}g + h^{T}Hh + o(\|h\|^{2}),$$

where  $g = \nabla f(x)$  and  $H = \nabla^2 f(x)$  are uniquely determined by the fact that the residual is  $o(\|h\|^2)$ . Applying this to the point  $x = x_0 + Fu$  this implies that

$$\overline{f}(u+h) = f(x_0 + Fu + Fh) = f(x_0 + Fu) + (Fh)^T g + (Fh)^T HFh + o(||Fh||^2).$$

Since  $o(\|Fh\|^2)$  is  $o(\|h\|^2)$  we conclude from this that

$$\nabla \overline{f}(u) = F^T g$$
 and  $\nabla^2 \overline{f}(u) = F^T H F$ ,

or, more explicitly

$$\nabla \overline{f}(u) = F^T \nabla f(x_0 + Fu) \quad \text{and} \quad \nabla^2 \overline{f}(u) = F^T \nabla^2 f(x_0 + Fu) F.$$
 (3)

**Hessian of composition.** With a similar approach we can compute the Hessian of a composition g(f(x)), where here  $g: \mathbb{R} \to \mathbb{R}$  is a scalar function of one variable (more general g are much harder to handle and we do not need them). Indeed, set

$$y = f(x), \quad \nabla = \nabla f(x) \quad H = \nabla^2 f(x) \quad \text{and} \quad k = h^T \nabla + \frac{1}{2} h^T H h$$

and use second order Taylor approximations on f at the point x and g at the point y = f(x) to obtain:

$$g(f(x+h))) = g\left(f(x) + h^T \nabla + \frac{1}{2}h^T H h\right)$$
  
=  $g(y+k) = g(y) + g'(y) + \frac{1}{2}g''(y)k^2 + o(k^2)$   
=  $g(y) + g'(y)\left(h^T \nabla + \frac{1}{2}h^T H h\right) + \frac{1}{2}g''(y)\left(h^T \nabla + \frac{1}{2}h^T H h\right)^2 + o(\|h\|^2).$ 

Here we have used that  $o(k^2) = o(\|h\|^2)$ . Collect terms of first and second order in h together and sticking all terms of higher order into the residual  $o(\|h\|^2)$ . Note that the squared term contributes no first order terms and only one second order term, this being the term

$$(\boldsymbol{h}^T \nabla)^2 = (\boldsymbol{h}^T \nabla)(\boldsymbol{h}^T \nabla) = (\boldsymbol{h}^T \nabla)(\boldsymbol{h}^T \nabla)^T = \boldsymbol{h}^T (\nabla \nabla^T) \boldsymbol{h}.$$

We obtain

$$g(f(x+h))) = g(y) + h^{T}[g'(y)\nabla] + \frac{1}{2}h^{T}[g'(y)H + g''(y)\nabla\nabla^{T}]h + o(\|h\|^{2}).$$

and from this we can read off that

$$\nabla(g \circ f)(x) = g'(f(x))\nabla f(x), \quad \text{and}$$
(4)

$$H(g \circ f)(x) = g'(f(x))H + g''(f(x))\nabla f(x)\nabla f(x)^{T}$$
(5)

Note that here  $d = \nabla f(x)$  is viewed as a column vector and so  $\nabla f(x) \nabla f(x)^T$  is the *outer product* 

$$\nabla f(x)\nabla f(x)^T = dd^T = (d_i d_j)_{ij}.$$

We will need the following example to construct test functions for unconstrained minimization:

**Example 2.1.** Let  $\phi: G \subseteq \mathbb{R} \to \mathbb{R}$  be a function of one variable defined on an open subset  $G \subseteq \mathbb{R}$ , fix  $a \in \mathbb{R}^n$  and set  $g(x) = \phi(a \cdot x)$ , for all  $x \in \mathbb{R}^n$  such that  $a \cdot x \in G$ .

Since  $f(x) = a \cdot x$  satisfies  $\nabla f(x) = a$  and  $H = \nabla^2 f(x) = 0$ , for all  $x \in \mathbb{R}^n$ , the formulas (4) and (5) yield

$$\nabla g(x) = \phi'(a \cdot x)a$$
 and  $\nabla^2 g(x) = \phi''(a \cdot x)aa'$ .

The idea is to construct test functions of the form

$$objF(x) = \sum_{j} \phi_{j}(a_{j} \cdot x)$$

where all the  $\phi_j = \phi_j(u)$  have a unique minimum at u = 0. Then objF(x) assumes its global minimum at all points x satisfying  $a_j \cdot x = 0$ , for all j, equivalently Ax = 0, where A is the matrix with rows  $a_j$ , provided such a solution exists.

If the functions  $\phi_j$  are all convex, the same is true of our objective function objF (sums and compositions of convex functions are again convex) With this we can construct examples with well conditioned, poorly conditioned and even singular Hessians  $(ker(A) \neq \{0\})$  with known minimizers. The conditioning of  $\nabla^2 f(x)$  is closely related to that of the matrix A.

# 3 Nullspace

Here we discuss methods to solve an underdetermined set of linear equations Ax = b where the number n of variables is larger than the number m of equations. In other words A is an  $m \times n$  matrix with m < n.

The set of solutions is then the hyperplane  $x_0 + ker(A)$  where ker(A) denotes the nullspace of A and  $x_0$  is any particular solution of Ax = b.

We will represent the nullspace of A as the columnspace colspace(F) of a matrix F such that

$$AF = 0.$$

If F satisfies this equation, then the column space of F is a subspace of the null space ker(A). Recall also that the column space of F is the range Im(F) of the linear map defined by the matrix F.

We will generally assume that A has full rank, i.e. rank(A) = m. The nullspace ker(A) then has dimension n - m and the columns of F span the entire nullspace of A if and only if rank(F) = n - m.

If we have found such a matrix F and  $x_0$  with  $Ax_0 = b$ , then the hyperplane of all solutions to Ax = b can be represented as

$$x_0 + ker(A) = x_0 + Im(F).$$

We want to apply this to the minimization problem

minimize 
$$f: \mathbb{R}^n \to \mathbb{R}$$
 under the constraint  $Ax = b$ . (6)

Since the solutions x to the constraint are exactly the vectors x of the form  $x = x_0 + Fu$ ,  $u \in \mathbb{R}^{n-m}$ , we can turn this into the unconstrained problem

minimize 
$$g(u) = f(x_0 + Fu)$$
 on all of  $\mathbb{R}^{n-m}$ . (7)

Thereby we reduce the dimension of the problem from n to n-m. In large scale problems one would not do this since this operation can destroy the sparseness of the Hessian H(f) which is critical for computation in very large dimensional problems.

In dense small problem (up to say n = 3000 variables) the above elimination of the constraint Ax = b is a reasonable approach. We will discuss two algorithms to finding F and  $x_0$ . In both cases the columns of F will be orthonormal and hence an orthonormal basis for the null space ker(A).

### 3.1 Null space with QR-decomposition of A

This approach relies on the relation  $ker(A) = Im(A')^{\perp}$ , where the prime denotes the matrix transpose and  $V^{\perp}$  is the orthogonal complement of a subspace V. Recall that

$$n = \dim(\ker(A)) + \dim(\ker(A)^{\perp}) = \dim(\ker(A)) + \dim(\operatorname{Im}(A'))$$
$$= (n - m) + \dim(\operatorname{Im}(A')).$$

Thus dim(Im(A')) = m. We will find the range Im(A') and its orthogonal complement using the QR-decomposition of the transpose A' (which has dimension  $n \times m$ , i.e. maps from  $\mathbb{R}^m$  to  $\mathbb{R}^n$ ):

$$A' = QR, (8)$$

where Q is orthogonal and R upper triangular with nonzero diagonal, where R is  $p \times m$  (i.e. maps from  $\mathbb{R}^m \to \mathbb{R}^p$  and Q is  $m \times p$ , i.e maps from  $\mathbb{R}^p \to \mathbb{R}^m$ .

In such a factorization Q and R must be  $n \times p$  and  $p \times m$  respectively, since A' is  $n \times m$ . Moreover  $rank(Q) \ge rank(A') = rank(A) = m$  and so we must have  $p \ge m$ . Because the columns of Q are linearly independent and of dimension n we also must have  $p \le n$ .

Indeed such factorizations exist for all  $m \leq p \leq n$  but in practice (i.e. available in libraries) there are only two flavours:

(A) The reduced (minimal) form gives us R as an  $m \times m$  matrix and  $Q = [q_1, \ldots, q_m]$  as an  $n \times m$  matrix with m-columns  $q_j \in \mathbb{R}^n$ . It follows that R maps onto  $\mathbb{R}^m = dom(Q)$  and hence

$$Im(A') = Im(Q) = span(q_1, \dots, q_m).$$

From this we see that  $ker(A) = Im(A')^{\perp} = span(q_1, \ldots, q_m)^{\perp}$  but if we use the reduced form we have to compute this orthogonal complement ourselves. In other words we have to extend  $\{q_1, \ldots, q_m\}$  to an orthonormal basis  $\{q_1, \ldots, q_m, q_{m+1}, \ldots, q_n\}$  of  $\mathbb{R}^n$  ourselves.

In R the minimal form is obtained as follows

$$qrA \leftarrow qr(t(A)); Q \leftarrow qr.Q(qrA); R \leftarrow qr.R(qrA)$$

Here we first compute a QR-decomposition object qrA and from this object extract the matrices Q and R using the helper functions qr.Q and qr.R.

The intermediate forms of the decomposition (8) with  $m \leq p \leq n$  extend this matrix Q by adding (arbitrarily) orthonormal columns on the right

$$Q \to Q_+ = [q_1, \dots, q_m, q_{m+1}, \dots q_n]$$

and adjusting the matrix R by adding zero rows at the bottom (so that R becomes  $p \times m$ . This of course does not change the matrix product QR as we can see from block multiplication

$$Q_{+} = [Q, Q_{1}], \ R_{+} = \begin{pmatrix} R \\ 0 \end{pmatrix} \implies Q_{+}R_{+} = QR + Q1 * 0 = QR.$$

(B) The full (maximal) form of the decomposition has p = n, that is, the columns of Q have been extended to a full orthonormal basis of  $\mathbb{R}^n$  and so clearly

$$ker(A) = Im(A')^{\perp} = span(q_1, \dots, q_m)^{\perp} = span(q_{m+1}, \dots, q_n).$$
 (9)

In other words the matrix F can be chosen to be the matrix with columns  $q_{m+1}, \ldots, q_n$ :

$$F = [q_{m+1}, \dots, q_n] \in Mat_{n \times (n-m)}(\mathbb{R}). \tag{10}$$

To get a special solution  $x_0$  of Ax = b note that A = R'Q' and solve the factored form R'Q'x = b. Set y = Q'x. Then forward solve the lower triangular system R'y = b and get x from Q'x = y as x = Qy.

To get the complete factorization  $A' = Q_+ R_+$  in R we do

and then extract the matrix Q via  $\mathbb{Q} \leftarrow \mathbb{Qp[,1:m]}$ . R is already in incomplete form (no zero rows at bottom) since we have not specified complete=TRUE in the function  $\mathbb{qr.R}$  (the default is complete=FALSE).

With this the factorization of A' becomes A' = QR (and not A' = QpR) and this is what we need in the computation of the special solution  $x_0$  above.

The same holds true in the scala breeze library where we get the QR decomposition as

val 
$$qr.QR(q,r)=qr(A)$$
.

This yields Q = q as  $n \times n$ -matrix but yields R = r as  $m \times m$ -matrix, that is, without the n - m bottom zero rows. With this the decomposition of A' becomes  $A' = Q[\cdot, 1:m] * R$  and so A = R'P' where the matrix

$$P = Q[\cdot, 1:m] = [col_1(Q), \dots, col_m(Q)]$$

has orthogonal columns and we must solve for  $x_0$  as R'y = b, P'x = y, that is, x = Py. Let us note that this  $x_0$  is the minimal norm solution of the system Ax = b. Indeed  $x_0 = Qy \in span(col_1(Q), \ldots, col_m(Q))$  while  $Fu \in span(col_{m+1}(Q), \ldots, col_n(Q))$  and so  $Fu \perp x_0$ , for all  $u \in \mathbb{R}^{n-m}$ , from which the claim follows.

The QR factorization is computed by repeatedly applying (orthogonal) Householder updates and is thus a very stable algorithm.

#### 3.2 Nullspace via SVD decomposition of A

We can also compute the nullspace ker(A) and special solution  $x_0$  of Ax = b using the more involved SVD-decomposition of A:

$$A = U\Sigma V'$$

where  $U \in Mat_{m \times m}(\mathbb{R})$  and  $V \in Mat_{n \times m}(\mathbb{R})$  are orthonormal matrices and  $\Sigma$  is an  $m \times n$  diagonal matrix with entries

$$\sigma_1 \ge \sigma_2 \cdots \ge \sigma_m \ge 0$$

(the singular values of A). Since rank(A) = m we must have  $\sigma_i > 0$  and so the matrix  $\Sigma$  is invertible. Write  $V = [v_1, \dots, v_m]$ , where the  $v_j \in \mathbb{R}^n$  are the columns of V. Thus, for all  $x \in \mathbb{R}^n$ ,

$$Ax = 0 \iff U\Sigma V'x = 0$$

$$\iff \Sigma V'x = 0$$

$$\iff 0 = V'x = (v_1 \cdot x, v_2 \cdot x, \dots, v_m \cdot x)' = 0$$

$$\iff x \perp \{v_1, \dots, v_m\}$$

$$\iff x \in span(v_{m+1}, \dots, v_n)$$

In other words

$$ker(A) = span(v_{m+1}, \dots, v_n). \tag{11}$$

where  $v_{m+1}, \ldots, v_n \in \mathbb{R}^n$  are vectors which extend  $v_1, \ldots, v_m$  to an ON-basis of  $\mathbb{R}^n$ . To get the full  $n \times n$  matrix

$$V_+ = [V, v_{m+1}, \dots, v_n]$$

we must compute the SVD in R via

Now a particular solution of  $U\Sigma V'x=Ax=b$  can be found by solving  $\Sigma V'x=U'b:=c$  which is equivalent to

$$V'x = (c_1/\sigma_1, \dots, c_m/\sigma_m)',$$

A particular solution of this is given by

$$x_0 = Vw$$
, where  $w = (c_1/\sigma_1, \dots, c_m/\sigma_m)' \in \mathbb{R}^m$ . (12)

since the columns of V are orthonormal.