

Exercise class: propositional logic (solutions)

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Week 5

1. Logical equivalence

We consider the following truth table:

a	b	$a \wedge \neg b$	$a \vee (\neg a \wedge b)$	$a \wedge (\neg a \vee \neg b)$	$\neg(a \rightarrow b)$	$a \vee b \rightarrow \neg a$
<i>true</i>	<i>true</i>	<i>false</i>	<i>true</i>	<i>false</i>	<i>false</i>	<i>false</i>
<i>true</i>	<i>false</i>	<i>true</i>	<i>true</i>	<i>true</i>	<i>true</i>	<i>false</i>
<i>false</i>	<i>true</i>	<i>false</i>	<i>true</i>	<i>false</i>	<i>false</i>	<i>true</i>
<i>false</i>	<i>false</i>	<i>false</i>	<i>false</i>	<i>false</i>	<i>false</i>	<i>true</i>

It follows that $a \wedge \neg b$ is equivalent to $a \wedge (\neg a \vee \neg b)$ and $\neg(a \rightarrow b)$ but not to $a \vee (\neg a \wedge b)$ or $a \vee b \rightarrow \neg a$.

2. Conjunctive normal form

We find

$$\begin{aligned}\neg(a \vee b) \vee (c \rightarrow \neg d) \\ &\equiv \neg(a \vee b) \vee (\neg c \vee \neg d) \\ &\equiv (\neg a \wedge \neg b) \vee (\neg c \vee \neg d) \\ &\equiv (\neg a \vee \neg c \vee \neg d) \wedge (\neg b \vee \neg c \vee \neg d)\end{aligned}$$

$$\begin{aligned}(a \rightarrow (b \rightarrow c)) \rightarrow d \\ &\equiv \neg(\neg a \vee (\neg b \vee c)) \vee d \\ &\equiv (\neg\neg a \wedge \neg\neg b \wedge \neg c) \vee d \\ &\equiv (a \wedge b \wedge \neg c) \vee d \\ &\equiv (a \vee d) \wedge (b \vee d) \wedge (\neg c \vee d)\end{aligned}$$

$$\begin{aligned}a \wedge \neg(b \wedge (c \vee \neg(d \vee e))) \\ &\equiv a \wedge (\neg b \vee (\neg c \wedge \neg\neg(d \vee e))) \\ &\equiv a \wedge (\neg b \vee (\neg c \wedge (d \vee e))) \\ &\equiv a \wedge (\neg b \vee \neg c) \wedge (\neg b \vee d \vee e)\end{aligned}$$

3. Refutation

Let us first convert $\neg(\neg(a \rightarrow b \vee c) \rightarrow \neg(d \wedge e))$ into CNF:

$$\begin{aligned} & \neg(\neg(a \rightarrow b \vee c) \rightarrow \neg(d \wedge e)) \\ & \equiv \neg(\neg\neg(\neg a \vee b \vee c) \vee \neg(d \wedge e)) \\ & \equiv \neg(\neg a \vee b \vee c \vee \neg(d \wedge e)) \\ & \equiv (\neg\neg a \wedge \neg b \wedge \neg c \wedge \neg\neg(d \wedge e)) \\ & \equiv a \wedge \neg b \wedge \neg c \wedge d \wedge e \end{aligned}$$

- We need to show that $\{a, \neg b, \neg c, d, e, \neg a\}$ is unsatisfiable, by showing that the empty clause can be derived using resolution. From a and $\neg a$ we immediately obtain the empty clause using resolution.
- We need to show that $\{a, \neg b, \neg c, d, e, \neg a \vee b \vee \neg e\}$ is unsatisfiable. From $\neg a \vee b \vee \neg e$ and a we derive

$$b \vee \neg e$$

which combined with e gives

$$b$$

which together with $\neg b$ allows us to derive the empty clause.

4. Refutation II

$K \cup \{\neg\alpha\}$ contains the following clauses:

- (1) $a \vee e$
- (2) $a \vee \neg f$
- (3) $b \vee d$
- (4) c
- (5) $\neg d$
- (6) $\neg e \vee f$
- (7) $\neg a \vee \neg b$
- (8) $\neg a \vee \neg c$

Using resolution we find:

- (9) $\neg b \vee e$ (1 + 7)
- (10) $d \vee e$ (3 + 9)
- (11) e (5 + 10)
- (12) $\neg c \vee \neg f$ (2 + 8)
- (13) $\neg f$ (4 + 12)
- (14) $\neg e$ (6 + 13)

Combining (11) and (14) we derive the empty clause.