Université Libre de Bruxelles Faculté des Sciences

INFO-F-403 – Introduction to language theory and compiling

First session examination

January, 17th, 2014

Instructions

- This is a closed book test. You are not allowed to use any kind of reference.
- You can anwser in french or in english.
- Write your first and last names on each sheet that you hand in.
- Write clearly: you can use a pencil or a ballpen or even a quill as long as your answers are readable!
- Always provide full and rigourous justifications along with your answers.
- This test is worth 12 points out of 20. The weight of each questions is given as a reference.
- In your answers (diagrams representing automata, grammars,...), you can always use the conventions adopted in the course, without recalling them explicitely. If you deviate from these conventions, be sure to make it clear.

Question 1 — 4 points

Define formally the following notions:

- 1. regular language
- 2. regular expression (RE)
- 3. deterministic finite automaton (DFA)
- 4. non-deterministic finite automaton with ε -transitions (ε -NFA)

and prove *formally* that the classes of REs, DFAs and ε -NFAs recognise *exactly* the class of regular languages. That is, prove that all regular languages can be defined by an RE (resp. DFA, ε -NFA); and that the language defined by an RE (DFA, ε -NFA) is always regular.

Question 2 — 2 points

Define the notion of ambiguous context-free grammar. Is the following grammar ambiguous (where the set of terminals is $\{+, \times, id, (,)\}$? Prove your answer. If the grammar is ambiguous, give a non-ambiguous context-free grammar that defines the same language.

- \rightarrow E + E
- (3) $T \rightarrow T \times T$

- $\begin{array}{ccc}
 (5) & F & \rightarrow & id \\
 (6) & & (E)
 \end{array}$

Finally, explain why ambiguous grammars are an issue when building a compiler (you can use the above grammar as an example to support your answer).

Question 3 — 2 points

Explain the role of the *scanner* and the *parser* during compiling. Explain their respective place in the compiling chain, and how they interact.

Question 4 — 1 point

Give the diagram of a deterministic pushdown automaton, on the alphabet $\Sigma = \{a, b, c, d, e\}$, that accepts the language $L = \{(ab)^n c(de)^n \mid n \geq 0\}$ using the empty stack acceptance condition.

Question 5 — 3 points

Give the LR(0) canonical finite state machine (CFSM) of the following grammar (where the set of variables is $\{S', S, A, B, C\}$ and the set of terminals is $\{a, b, c, d\}$):

- (1) $S' \rightarrow S$
- (2) $S \rightarrow aAb$
- (3) aB
- (4) c
- (5) $A \rightarrow dA$
- (7) $B \rightarrow bC$
- (8) $C \rightarrow \varepsilon$