Al-Baath University ICPC Team Notebook (2018)

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1 Combinatorial optimization

1.1 Sparse max-flow

```
// Adjacency list implementation of Dinic's
    blocking flow algorithm.
// This is very fast in practice, and only loses
     to push-relabel flow.
// Running time:
      O(|V|^2 |E|)
// INPUT:
      - graph, constructed using AddEdge()
       - source and sink
// OUTPUT:
      - maximum flow value
      - To obtain actual flow values, look at
    edges with capacity > 0
         (zero capacity edges are residual edges
#include<cstdio>
#include < vector >
#include < queue >
using namespace std;
typedef long long LL;
struct Edge {
 int u, v;
 LL cap, flow;
  Edge() {}
 Edge(int u, int v, LL cap): u(u), v(v), cap(
      cap), flow(0) {}
struct Dinic {
 int N;
 vector<Edge> E;
 vector<vector<int>> q;
  vector<int> d, pt;
  Dinic(int N): N(N), E(0), g(N), d(N), pt(N) {}
  void AddEdge(int u, int v, LL cap) {
   if (u != v) {
     E.emplace_back(Edge(u, v, cap));
     g[u].emplace_back(E.size() - 1);
      E.emplace_back(Edge(v, u, 0));
      g[v].emplace_back(E.size() - 1);
  bool BFS(int S, int T) {
    queue<int> q({S});
    fill(d.begin(), d.end(), N + 1);
   d[S] = 0;
    while(!q.empty()) {
     int u = q.front(); q.pop();
     if (u == T) break;
      for (int k: q[u]) {
        Edge &e = E[k];
        if (e.flow < e.cap && d[e.v] > d[e.u] +
          d[e.v] = d[e.u] + 1;
          q.emplace(e.v);
    return d[T] != N + 1;
```

```
LL DFS (int u, int T, LL flow = -1) {
   if (u == T || flow == 0) return flow;
    for (int &i = pt[u]; i < g[u].size(); ++i) {</pre>
      Edge &e = E[g[u][i]];
      Edge &oe = E[g[u][i]^1];
      if (d[e.v] == d[e.u] + 1) {
       LL amt = e.cap - e.flow;
        if (flow != -1 && amt > flow) amt = flow
        if (LL pushed = DFS(e.v, T, amt)) {
          e.flow += pushed;
         oe.flow -= pushed;
         return pushed;
   return 0;
 LL MaxFlow(int S, int T) {
   LL total = 0;
   while (BFS(S, T)) {
     fill(pt.begin(), pt.end(), 0);
     while (LL flow = DFS(S, T))
       total += flow:
   return total;
// The following code solves SPOJ problem #4110:
     Fast Maximum Flow (FASTFLOW)
int main()
 int N, E;
 scanf("%d%d", &N, &E);
 Dinic dinic(N);
  for (int i = 0; i < E; i++)
   int u, v;
   scanf("%d%d%lld", &u, &v, &cap);
   dinic.AddEdge(u - 1, v - 1, cap);
   dinic.AddEdge(v - 1, u - 1, cap);
 printf("%lld\n", dinic.MaxFlow(0, N - 1));
 return 0;
// END CUT
```

1.2 Min-cost max-flow

```
// Implementation of min cost max flow algorithm
    using adjacency
// matrix (Edmonds and Karp 1972). This
    implementation keeps track of
// forward and reverse edges separately (so you
    can set cap[i][j] !=
// cap[j][i]). For a regular max flow, set all
    edge costs to 0.
//
// Running time, O(|V|^2) cost per augmentation
// max flow: O(|V|^3)
```

```
augmentations
      min cost max flow: O(|V|^4 *
    MAX_EDGE_COST) augmentations
// INPUT:
      - graph, constructed using AddEdge()
      - source
      - sink
// OUTPUT:
11
     - (maximum flow value, minimum cost value
    - To obtain the actual flow, look at
    positive values only.
#include <cmath>
#include <vector>
#include <iostream>
using namespace std;
typedef vector<int> VI;
typedef vector<VI> VVI;
typedef long long L;
typedef vector<L> VL;
typedef vector<VL> VVL;
typedef pair<int, int> PII;
typedef vector<PII> VPII;
const L INF = numeric_limits<L>::max() / 4;
struct MinCostMaxFlow {
 int N;
 VVL cap, flow, cost;
 VI found:
 VL dist, pi, width;
 VPII dad;
 MinCostMaxFlow(int N) :
    N(N), cap(N, VL(N)), flow(N, VL(N)), cost(N,
         VL(N)),
    found(N), dist(N), pi(N), width(N), dad(N)
  void AddEdge(int from, int to, L cap, L cost)
    this->cap[from][to] = cap;
   this->cost[from][to] = cost;
  void Relax(int s, int k, L cap, L cost, int
    L \text{ val} = \text{dist}[s] + \text{pi}[s] - \text{pi}[k] + \text{cost};
    if (cap && val < dist[k]) {</pre>
     dist[k] = val;
     dad[k] = make_pair(s, dir);
     width[k] = min(cap, width[s]);
  L Dijkstra(int s, int t) {
    fill(found.begin(), found.end(), false);
    fill(dist.begin(), dist.end(), INF);
    fill(width.begin(), width.end(), 0);
    dist[s] = 0;
    width[s] = INF;
    while (s != -1) {
     int best = -1;
```

```
found[s] = true;
      for (int k = 0; k < N; k++) {
        if (found[k]) continue;
        Relax(s, k, cap[s][k] - flow[s][k], cost
            [s][k], 1);
        Relax(s, k, flow[k][s], -cost[k][s], -1)
        if (best == -1 || dist[k] < dist[best])</pre>
            best = k;
      s = best;
    for (int k = 0; k < N; k++)
     pi[k] = min(pi[k] + dist[k], INF);
    return width[t];
  pair<L, L> GetMaxFlow(int s, int t) {
   L \text{ totflow} = 0, \text{ totcost} = 0;
    while (L amt = Dijkstra(s, t)) {
     totflow += amt;
     for (int x = t; x != s; x = dad[x].first)
        if (dad[x].second == 1) {
          flow[dad[x].first][x] += amt;
          totcost += amt * cost[dad[x].first][x
        } else {
          flow[x][dad[x].first] -= amt;
          totcost -= amt * cost[x][dad[x].first
     }
    return make_pair(totflow, totcost);
};
// BEGIN CUT
// The following code solves UVA problem #10594:
     Data Flow
int main() {
  int N, M;
  while (scanf("%d%d", &N, &M) == 2) {
    VVL v(M, VL(3));
    for (int i = 0; i < M; i++)
      scanf("%Ld%Ld%Ld", &v[i][0], &v[i][1], &v[
          i][2]);
    L D, K;
    scanf("%Ld%Ld", &D, &K);
    MinCostMaxFlow mcmf(N+1);
    for (int i = 0; i < M; i++) {
     mcmf.AddEdge(int(v[i][0]), int(v[i][1]), K
          , v[i][2]);
     mcmf.AddEdge(int(v[i][1]), int(v[i][0]), K
          , v[i][2]);
    mcmf.AddEdge(0, 1, D, 0);
    pair<L, L> res = mcmf.GetMaxFlow(0, N);
    if (res.first == D) {
     printf("%Ld\n", res.second);
    } else {
     printf("Impossible.\n");
```

1.3 Push-relabel max-flow

```
// Adjacency list implementation of FIFO push
    relabel maximum flow
// with the gap relabeling heuristic. This
    implementation is
// significantly faster than straight Ford-
    Fulkerson. It solves
// random problems with 10000 vertices and
    1000000 edges in a few
// seconds, though it is possible to construct
    test cases that
// achieve the worst-case.
// Running time:
    0(|V|^3)
// INPUT:
// - graph, constructed using AddEdge()
      - source
      - sink
// OUTPUT:
      - maximum flow value
       - To obtain the actual flow values, look
    at all edges with
      capacity > 0 (zero capacity edges are
    residual edges).
#include <cmath>
#include <vector>
#include <iostream>
#include <queue>
using namespace std;
typedef long long LL;
struct Edge {
  int from, to, cap, flow, index;
  Edge (int from, int to, int cap, int flow, int
    from (from), to (to), cap (cap), flow (flow),
        index(index) {}
struct PushRelabel {
  int N;
  vector<vector<Edge> > G;
  vector<LL> excess;
  vector<int> dist, active, count;
  queue<int> Q;
  PushRelabel(int N) : N(N), G(N), excess(N),
      dist(N), active(N), count(2*N) {}
  void AddEdge(int from, int to, int cap) {
    G[from].push_back(Edge(from, to, cap, 0, G[
        to].size()));
```

```
if (from == to) G[from].back().index++;
  G[to].push back(Edge(to, from, 0, 0, G[from
       ].size() - 1));
void Enqueue(int v) {
  if (!active[v] && excess[v] > 0) { active[v]
       = true; Q.push(v); }
void Push(Edge &e) {
  int amt = int(min(excess[e.from], LL(e.cap -
       e.flow)));
  if (dist[e.from] <= dist[e.to] || amt == 0)</pre>
      return;
  e.flow += amt;
  G[e.to][e.index].flow -= amt;
  excess[e.to] += amt;
  excess[e.from] -= amt;
  Enqueue(e.to);
void Gap(int k) {
  for (int v = 0; v < N; v++) {
    if (dist[v] < k) continue;</pre>
    count[dist[v]]--;
    dist[v] = max(dist[v], N+1);
    count[dist[v]]++;
    Enqueue (v);
void Relabel(int v) {
  count[dist[v]]--;
  dist[v] = 2*N;
  for (int i = 0; i < G[v].size(); i++)</pre>
    if (G[v][i].cap - G[v][i].flow > 0)
      dist[v] = min(dist[v], dist[G[v][i].to]
          + 1);
  count[dist[v]]++;
  Enqueue (v);
void Discharge(int v) {
  for (int i = 0; excess[v] > 0 && i < G[v].
      size(); i++) Push(G[v][i]);
  if (excess[v] > 0) {
    if (count[dist[v]] == 1)
      Gap(dist[v]);
    else
      Relabel(v);
LL GetMaxFlow(int s, int t) {
  count[0] = N-1;
  count[N] = 1;
  dist[s] = N;
  active[s] = active[t] = true;
  for (int i = 0; i < G[s].size(); i++) {</pre>
   excess[s] += G[s][i].cap;
   Push(G[s][i]);
  while (!Q.empty()) {
    int v = Q.front();
    Q.pop();
    active[v] = false;
    Discharge(v);
```

```
LL totflow = 0;
    for (int i = 0; i < G[s].size(); i++)</pre>
        totflow += G[s][i].flow;
    return totflow;
};
// BEGIN CUT
// The following code solves SPOJ problem #4110:
     Fast Maximum Flow (FASTFLOW)
int main() {
  int n, m;
  scanf("%d%d", &n, &m);
  PushRelabel pr(n);
  for (int i = 0; i < m; i++) {
   int a, b, c;
    scanf("%d%d%d", &a, &b, &c);
    if (a == b) continue;
    pr.AddEdge(a-1, b-1, c);
    pr.AddEdge(b-1, a-1, c);
  printf("%Ld\n", pr.GetMaxFlow(0, n-1));
  return 0;
// END CUT
```

1.4 Min-cost matching

```
// Min cost bipartite matching via shortest
    augmenting paths
// This is an O(n^3) implementation of a
    shortest augmenting path
// algorithm for finding min cost perfect
    matchings in dense
// graphs. In practice, it solves 1000x1000
    problems in around 1
// second.
    cost[i][j] = cost for pairing left node i
    with right node j
// Lmate[i] = index of right node that left
    node i pairs with
// Rmate[j] = index of left node that right
    node j pairs with
// The values in cost[i][j] may be positive or
    negative. To perform
// maximization, simply negate the cost[][]
#include <algorithm>
#include <cstdio>
#include <cmath>
#include <vector>
using namespace std;
typedef vector<double> VD;
typedef vector<VD> VVD;
```

```
typedef vector<int> VI;
double MinCostMatching(const VVD &cost, VI &
    Lmate, VI &Rmate) {
 int n = int(cost.size());
 // construct dual feasible solution
 VD u(n):
 VD v(n);
 for (int i = 0; i < n; i++) {
   u[i] = cost[i][0];
    for (int j = 1; j < n; j++) u[i] = min(u[i],
         cost[i][j]);
  for (int j = 0; j < n; j++) {
   v[j] = cost[0][j] - u[0];
   for (int i = 1; i < n; i++) v[j] = min(v[j]),
         cost[i][j] - u[i]);
 // construct primal solution satisfying
      complementary slackness
 Lmate = VI(n, -1);
 Rmate = VI(n, -1);
 int mated = 0;
 for (int i = 0; i < n; i++) {</pre>
   for (int j = 0; j < n; j++) {
     if (Rmate[j] != -1) continue;
     if (fabs(cost[i][j] - u[i] - v[j]) < 1e
          -10) {
        Lmate[i] = j;
        Rmate[j] = i;
        mated++;
       break;
 }
 VD dist(n);
 VI dad(n);
 VI seen(n);
  // repeat until primal solution is feasible
  while (mated < n) {</pre>
    // find an unmatched left node
   int s = 0;
   while (Lmate[s] != -1) s++;
    // initialize Dijkstra
    fill(dad.begin(), dad.end(), -1);
    fill(seen.begin(), seen.end(), 0);
    for (int k = 0; k < n; k++)
     dist[k] = cost[s][k] - u[s] - v[k];
   int j = 0;
   while (true) {
      // find closest
      j = -1;
      for (int k = 0; k < n; k++) {
       if (seen[k]) continue;
       if (j == -1 || dist[k] < dist[j]) j = k;</pre>
     seen[j] = 1;
      // termination condition
      if (Rmate[j] == -1) break;
```

```
// relax neighbors
    const int i = Rmate[i];
    for (int k = 0; k < n; k++) {
     if (seen[k]) continue;
     const double new_dist = dist[j] + cost[i
          [k] - u[i] - v[k];
     if (dist[k] > new_dist) {
       dist[k] = new_dist;
       dad[k] = j;
   }
  // update dual variables
  for (int k = 0; k < n; k++) {
   if (k == j || !seen[k]) continue;
    const int i = Rmate[k];
   v[k] += dist[k] - dist[j];
   u[i] -= dist[k] - dist[j];
  u[s] += dist[j];
  // augment along path
  while (dad[j] >= 0) {
   const int d = dad[j];
   Rmate[j] = Rmate[d];
   Lmate[Rmate[j]] = j;
    j = d;
  Rmate[j] = s;
  Lmate[s] = j;
  mated++:
double value = 0;
for (int i = 0; i < n; i++)
 value += cost[i][Lmate[i]];
return value;
```

1.5 Max bipartite matchine

```
// This code performs maximum bipartite matching
    .
//
// Running time: O(|E| |V|) -- often much faster
    in practice
//
// INPUT: w[i][j] = edge between row node i
    and column node j
// OUTPUT: mr[i] = assignment for row node i,
    -1 if unassigned
// mc[j] = assignment for column node
    j, -1 if unassigned
// function returns number of matches
    made

#include <vector>
using namespace std;

typedef vector<int> VI;
typedef vector<VI> VVI;
```

```
bool FindMatch (int i, const VVI &w, VI &mr, VI &
    mc, VI &seen) {
  for (int j = 0; j < w[i].size(); j++) {
    if (w[i][j] && !seen[j]) {
     seen[j] = true;
      if (mc[j] < 0 || FindMatch(mc[j], w, mr,</pre>
          mc, seen)) {
        mr[i] = j;
        mc[j] = i;
        return true;
  return false:
int BipartiteMatching(const VVI &w, VI &mr, VI &
  mr = VI(w.size(), -1);
  mc = VI(w[0].size(), -1);
  int ct = 0;
  for (int i = 0; i < w.size(); i++) {</pre>
   VI seen(w[0].size());
   if (FindMatch(i, w, mr, mc, seen)) ct++;
  return ct;
```

1.6 Global min-cut

```
// Adjacency matrix implementation of Stoer-
    Wagner min cut algorithm.
// Running time:
// O(|V|^3)
// INPUT:

    graph, constructed using AddEdge()

// OUTPUT:
   - (min cut value, nodes in half of min
#include <cmath>
#include <vector>
#include <iostream>
using namespace std;
typedef vector<int> VI;
typedef vector<VI> VVI;
const int INF = 1000000000;
pair<int, VI> GetMinCut(VVI &weights) {
 int N = weights.size();
 VI used(N), cut, best_cut;
 int best_weight = -1;
 for (int phase = N-1; phase >= 0; phase--) {
   VI w = weights[0];
   VI added = used;
   int prev, last = 0;
   for (int i = 0; i < phase; i++) {</pre>
     prev = last;
```

```
last = -1;
      for (int j = 1; j < N; j++)
       if (!added[j] && (last == -1 \mid \mid w[j] > w
            [last]) last = j;
      if (i == phase-1) {
        for (int j = 0; j < N; j++) weights[prev</pre>
            [j] += weights[last][j];
        for (int j = 0; j < N; j++) weights[j][</pre>
            prev] = weights[prev][j];
       used[last] = true;
       cut.push_back(last);
        if (best weight == -1 || w[last] <</pre>
            best_weight) {
          best_cut = cut;
         best_weight = w[last];
      } else {
        for (int j = 0; j < N; j++)
         w[j] += weights[last][j];
        added[last] = true;
 return make_pair(best_weight, best_cut);
// BEGIN CUT
// The following code solves UVA problem #10989:
     Bomb, Divide and Conquer
int main() {
 int N;
 cin >> N:
  for (int i = 0; i < N; i++) {
   int n, m;
   cin >> n >> m;
   VVI weights(n, VI(n));
   for (int j = 0; j < m; j++) {
     int a, b, c;
     cin >> a >> b >> c;
     weights[a-1][b-1] = weights[b-1][a-1] = c;
   pair<int, VI> res = GetMinCut(weights);
   cout << "Case #" << i+1 << ": " << res.first
// END CUT
```

2 Geometry

2.1 Convex hull

```
// Compute the 2D convex hull of a set of points
    using the monotone chain
// algorithm. Eliminate redundant points from
    the hull if REMOVE_REDUNDANT is
// #defined.
//
// Running time: O(n log n)
//
// INPUT: a vector of input points,
    unordered.
// OUTPUT: a vector of points in the convex
    hull, counterclockwise, starting
// with bottommost/leftmost point
```

```
#include <cstdio>
#include <cassert>
#include <vector>
#include <algorithm>
#include <cmath>
// BEGIN CUT
#include <map>
// END CUT
using namespace std;
#define REMOVE_REDUNDANT
typedef double T;
const T EPS = 1e-7;
struct PT {
 T x, y;
  PT() {}
  PT(T x, T y) : x(x), y(y) {}
 bool operator<(const PT &rhs) const { return</pre>
      make_pair(y,x) < make_pair(rhs.y,rhs.x);</pre>
 bool operator==(const PT &rhs) const { return
      make_pair(y,x) == make_pair(rhs.y,rhs.x);
};
T cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
T area2(PT a, PT b, PT c) { return cross(a,b) +
    cross(b,c) + cross(c,a); }
#ifdef REMOVE REDUNDANT
bool between (const PT &a, const PT &b, const PT
  return (fabs(area2(a,b,c)) < EPS && (a.x-b.x)</pre>
      *(c.x-b.x) \le 0 \&\& (a.y-b.y) *(c.y-b.y) \le
       0);
#endif
void ConvexHull(vector<PT> &pts) {
  sort(pts.begin(), pts.end());
  pts.erase(unique(pts.begin(), pts.end()), pts.
      end());
  vector<PT> up, dn;
  for (int i = 0; i < pts.size(); i++) {</pre>
    while (up.size() > 1 && area2(up[up.size()
         -2], up.back(), pts[i]) >= 0) up.
         pop_back();
    while (dn.size() > 1 && area2(dn[dn.size()
        -2], dn.back(), pts[i]) <= 0) dn.
        pop back();
    up.push_back(pts[i]);
    dn.push back(pts[i]);
  for (int i = (int) up.size() - 2; i >= 1; i--)
       pts.push_back(up[i]);
#ifdef REMOVE REDUNDANT
  if (pts.size() <= 2) return;</pre>
  dn.clear();
  dn.push_back(pts[0]);
  dn.push_back(pts[1]);
  for (int i = 2; i < pts.size(); i++) {</pre>
    if (between(dn[dn.size()-2], dn[dn.size()
        -1], pts[i])) dn.pop_back();
    dn.push_back(pts[i]);
```

```
if (dn.size() >= 3 && between(dn.back(), dn
       [0], dn[1])) {
    dn[0] = dn.back();
    dn.pop_back();
  pts = dn;
#endif
// BEGIN CUT
// The following code solves SPOJ problem #26:
    Build the Fence (BSHEEP)
int main() {
 int t;
  scanf("%d", &t);
  for (int caseno = 0; caseno < t; caseno++) {</pre>
    int n;
    scanf("%d", &n);
    vector<PT> v(n);
    for (int i = 0; i < n; i++) scanf("%lf%lf",</pre>
        &v[i].x, &v[i].v);
    vector<PT> h(v);
    map<PT, int> index;
    for (int i = n-1; i >= 0; i--) index[v[i]] =
         i+1:
    ConvexHull(h);
    double len = 0;
    for (int i = 0; i < h.size(); i++) {</pre>
      double dx = h[i].x - h[(i+1)\%h.size()].x;
      double dy = h[i].y - h[(i+1)%h.size()].y;
      len += sqrt (dx*dx+dy*dy);
    if (caseno > 0) printf("\n");
    printf("%.2f\n", len);
    for (int i = 0; i < h.size(); i++) {</pre>
      if (i > 0) printf(" ");
      printf("%d", index[h[i]]);
    printf("\n");
// END CUT
```

2.2 Miscellaneous geometry

```
#include <iostream>
#include <vector>
#include <cmath>
#include <cassert>

using namespace std;

double INF = 1e100;
double EPS = 1e-12;

struct PT {
   double x, y;
   PT() {}
   PT(double x, double y) : x(x), y(y) {}
   PT(const PT &p) : x(p.x), y(p.y) {}
```

```
PT operator - (const PT &p) const { return PT
 (x-p.x, y-p.y); }
PT operator * (double c)
                                const { return PT
       (x*c, y*c ); }
  PT operator / (double c)
                                const { return PT
      (x/c, y/c);
double dot (PT p, PT q)
                           { return p.x*q.x+p.y*
    q.v; }
double dist2(PT p, PT q)
                           { return dot(p-q,p-q)
double cross(PT p, PT q)
                          { return p.x*q.y-p.y*
ostream & operator << (ostream & os, const PT &p) {
 return os << "(" << p.x << "," << p.y << ")";
// rotate a point CCW or CW around the origin
PT RotateCCW90(PT p) { return PT(-p.y,p.x); }
PT RotateCW90 (PT p) { return PT (p.v,-p.x); }
PT RotateCCW(PT p, double t) {
  return PT(p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.
      v*cos(t));
// project point c onto line through a and b
// assuming a != b
PT ProjectPointLine(PT a, PT b, PT c) {
  return a + (b-a) *dot(c-a, b-a)/dot(b-a, b-a);
// project point c onto line segment through a
PT ProjectPointSegment (PT a, PT b, PT c) {
  double r = dot(b-a,b-a);
  if (fabs(r) < EPS) return a;</pre>
 r = dot(c-a, b-a)/r;
  if (r < 0) return a;</pre>
  if (r > 1) return b;
  return a + (b-a) *r;
// compute distance from c to segment between a
double DistancePointSegment(PT a, PT b, PT c) {
  return sqrt(dist2(c, ProjectPointSegment(a, b,
       c)));
// compute distance between point (x, y, z) and
    plane ax+by+cz=d
double DistancePointPlane (double x, double v,
    double z,
                           double a, double b,
                               double c, double
  return fabs(a*x+b*y+c*z-d)/sqrt(a*a+b*b+c*c);
// determine if lines from a to b and c to d are
     parallel or collinear
bool LinesParallel (PT a, PT b, PT c, PT d) {
  return fabs(cross(b-a, c-d)) < EPS;</pre>
```

PT operator + (const PT &p) const { return PT

(x+p.x, y+p.y); }

```
bool LinesCollinear(PT a, PT b, PT c, PT d) {
  return LinesParallel(a, b, c, d)
      && fabs(cross(a-b, a-c)) < EPS
      && fabs(cross(c-d, c-a)) < EPS;
// determine if line segment from a to b
    intersects with
// line segment from c to d
bool SegmentsIntersect(PT a, PT b, PT c, PT d) {
  if (LinesCollinear(a, b, c, d)) {
    if (dist2(a, c) < EPS || dist2(a, d) < EPS</pre>
      dist2(b, c) < EPS \mid | dist2(b, d) < EPS)
          return true;
    if (dot(c-a, c-b) > 0 \&\& dot(d-a, d-b) > 0
        && dot(c-b, d-b) > 0
      return false;
    return true;
  if (cross(d-a, b-a) * cross(c-a, b-a) > 0)
      return false;
  if (cross(a-c, d-c) * cross(b-c, d-c) > 0)
      return false:
  return true;
// compute intersection of line passing through
    a and b
// with line passing through c and d, assuming
    that unique
// intersection exists; for segment intersection
    , check if
// segments intersect first
PT ComputeLineIntersection(PT a, PT b, PT c, PT
    d) {
  b=b-a; d=c-d; c=c-a;
  assert (dot (b, b) > EPS && dot (d, d) > EPS);
  return a + b*cross(c, d)/cross(b, d);
// compute center of circle given three points
PT ComputeCircleCenter(PT a, PT b, PT c) {
 b = (a+b)/2;
  c = (a + c) / 2;
  return ComputeLineIntersection(b, b+RotateCW90
      (a-b), c, c+RotateCW90(a-c));
// determine if point is in a possibly non-
    convex polygon (by William
// Randolph Franklin); returns 1 for strictly
    interior points, 0 for
// strictly exterior points, and 0 or 1 for the
    remaining points.
// Note that it is possible to convert this into
     an *exact* test using
// integer arithmetic by taking care of the
    division appropriately
// (making sure to deal with signs properly) and
     then by writing exact
// tests for checking point on polygon boundary
bool PointInPolygon(const vector<PT> &p, PT q) {
 bool c = 0;
  for (int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) p.size();
    if ((p[i].y \le q.y \& q.y < p[j].y | |
      p[j].y \le q.y \&\& q.y < p[i].y) \&\&
```

```
q.x < p[i].x + (p[j].x - p[i].x) * (q.y -
          p[i].y) / (p[j].y - p[i].y))
      c = !c;
  return c;
// determine if point is on the boundary of a
    polvaon
bool PointOnPolygon(const vector<PT> &p, PT q) {
  for (int i = 0; i < p.size(); i++)
    if (dist2(ProjectPointSegment(p[i], p[(i+1)%
        p.size()], q), q) < EPS)
      return true;
    return false;
// compute intersection of line through points a
     and b with
// circle centered at c with radius r > 0
vector<PT> CircleLineIntersection(PT a, PT b, PT
     c, double r) {
  vector<PT> ret;
  b = b-a;
  a = a-c;
  double A = dot(b, b);
  double B = dot(a, b);
  double C = dot(a, a) - r * r;
  double D = B*B - A*C;
  if (D < -EPS) return ret;</pre>
  ret.push_back(c+a+b*(-B+sqrt(D+EPS))/A);
  if (D > EPS)
    ret.push_back(c+a+b*(-B-sqrt(D))/A);
  return ret;
// compute intersection of circle centered at a
    with radius r
// with circle centered at b with radius R
vector<PT> CircleCircleIntersection(PT a, PT b,
    double r, double R) {
  vector<PT> ret;
  double d = sqrt(dist2(a, b));
  if (d > r+R \mid \mid d+min(r, R) < max(r, R)) return
       ret:
  double x = (d*d-R*R+r*r)/(2*d);
  double y = sqrt(r*r-x*x);
  PT v = (b-a)/d;
  ret.push_back(a+v*x + RotateCCW90(v)*y);
  if (y > 0)
    ret.push_back(a+v*x - RotateCCW90(v)*y);
  return ret;
// This code computes the area or centroid of a
     (possibly nonconvex)
// polygon, assuming that the coordinates are
    listed in a clockwise or
// counterclockwise fashion. Note that the
    centroid is often known as
// the "center of gravity" or "center of mass".
double ComputeSignedArea(const vector<PT> &p) {
  double area = 0;
  for(int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) % p.size();
    area += p[i].x*p[j].y - p[j].x*p[i].y;
  return area / 2.0;
```

```
double ComputeArea(const vector<PT> &p) {
  return fabs(ComputeSignedArea(p));
PT ComputeCentroid(const vector<PT> &p) {
  PT c(0,0);
  double scale = 6.0 * ComputeSignedArea(p);
  for (int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) % p.size();
    c = c + (p[i]+p[j])*(p[i].x*p[j].y - p[j].x*
         p[i].y);
  return c / scale;
// tests whether or not a given polygon (in CW
    or CCW order) is simple
bool IsSimple(const vector<PT> &p) {
  for (int i = 0; i < p.size(); i++) {</pre>
    for (int k = i+1; k < p.size(); k++) {
      int j = (i+1) % p.size();
      int \bar{1} = (k+1) % p.size();
      if (i == 1 \mid \mid j == k) continue;
      if (SegmentsIntersect(p[i], p[j], p[k], p[
           11))
        return false;
  return true;
int main() {
  // expected: (-5,2)
  cerr << RotateCCW90(PT(2,5)) << endl;</pre>
  // expected: (5,-2)
  cerr << RotateCW90(PT(2,5)) << endl;</pre>
  // expected: (-5,2)
  cerr << RotateCCW(PT(2,5),M_PI/2) << endl;</pre>
  // expected: (5,2)
  cerr << ProjectPointLine(PT(-5,-2), PT(10,4),</pre>
       PT(3,7)) << endl;
  // expected: (5,2) (7.5,3) (2.5,1)
  cerr << ProjectPointSegment(PT(-5,-2), PT</pre>
       (10,4), PT(3,7)) << " "
       << ProjectPointSegment(PT(7.5,3), PT</pre>
             (10,4), PT(3,7)) << " "
       << ProjectPointSegment(PT(-5,-2), PT
             (2.5,1), PT(3,7)) << endl;
  // expected: 6.78903
  cerr << DistancePointPlane (4, -4, 3, 2, -2, 5, -8)
       << endl:
  // expected: 1 0 1
  cerr << LinesParallel(PT(1,1), PT(3,5), PT</pre>
       (2,1), PT(4,5)) << " "
       << LinesParallel(PT(1,1), PT(3,5), PT
            (2,0), PT(4,5)) << " "
       << LinesParallel(PT(1,1), PT(3,5), PT
            (5,9), PT(7,13)) << endl;
  // expected: 0 0 1
  cerr << LinesCollinear(PT(1,1), PT(3,5), PT</pre>
```

```
(2,1), PT(4,5)) << " "
     << LinesCollinear(PT(1,1), PT(3,5), PT
          (2,0), PT(4,5)) << " "
     << LinesCollinear(PT(1,1), PT(3,5), PT
          (5,9), PT(7,13)) << endl;
// expected: 1 1 1 0
cerr << SegmentsIntersect(PT(0,0), PT(2,4), PT</pre>
    (3,1), PT(-1,3)) << " "
     << SegmentsIntersect(PT(0,0), PT(2,4), PT
          (4,3), PT(0,5)) << " "
     << SegmentsIntersect(PT(0,0), PT(2,4), PT
          (2,-1), PT(-2,1)) << ""
     << SegmentsIntersect(PT(0,0), PT(2,4), PT
          (5,5), PT(1,7)) << endl;
// expected: (1,2)
cerr << ComputeLineIntersection(PT(0,0), PT</pre>
    (2,4), PT(3,1), PT(-1,3)) << endl;
// expected: (1,1)
cerr << ComputeCircleCenter(PT(-3,4), PT(6,1),</pre>
     PT(4,5)) << end1;
vector<PT> v;
v.push_back(PT(0,0));
v.push_back(PT(5,0));
v.push_back(PT(5,5));
v.push_back(PT(0,5));
// expected: 1 1 1 0 0
cerr << PointInPolygon(v, PT(2,2)) << " "</pre>
     << PointInPolygon(v, PT(2,0)) << " "
     << PointInPolygon(v, PT(0,2)) << " "
     << PointInPolygon(v, PT(5,2)) << " "
     << PointInPolygon(v, PT(2,5)) << endl;</pre>
// expected: 0 1 1 1 1
cerr << PointOnPolygon(v, PT(2,2)) << " "</pre>
     << PointOnPolygon(v, PT(2,0)) << " "
     << PointOnPolygon(v, PT(0,2)) << " "
     << PointOnPolygon(v, PT(5,2)) << " "
     << PointOnPolygon(v, PT(2,5)) << endl;
// expected: (1,6)
             (5,4) (4,5)
             blank line
             (4,5) (5,4)
             blank line
             (4,5) (5,4)
vector<PT> u = CircleLineIntersection(PT(0,6),
     PT(2,6), PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i</pre>
    ] << " "; cerr << endl;
u = CircleLineIntersection(PT(0,9), PT(9,0),
    PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i</pre>
    ] << " "; cerr << endl;
u = CircleCircleIntersection(PT(1,1), PT
    (10,10), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i</pre>
    ] << " "; cerr << endl;
u = CircleCircleIntersection(PT(1,1), PT(8,8),
     5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i</pre>
    | << " "; cerr << endl;</pre>
u = CircleCircleIntersection(PT(1,1), PT
    (4.5, 4.5), 10, sqrt(2.0)/2.0);
for (int i = 0; i < u.size(); i++) cerr << u[i</pre>
```

2.3 3D geometry

public class Geom3D {

```
// distance from point (x, y, z) to plane aX +
     bY + cZ + d = 0
public static double ptPlaneDist(double x,
    double y, double z,
    double a, double b, double c, double d) {
  return Math.abs(a*x + b*y + c*z + d) / Math.
      sqrt(a*a + b*b + c*c);
// distance between parallel planes aX + bY +
    cZ + d1 = 0 and
// aX + bY + cZ + d2 = 0
public static double planePlaneDist(double a,
    double b, double c,
    double d1, double d2) {
  return Math.abs(d1 - d2) / Math.sqrt(a*a + b
      *b + c*c);
// distance from point (px, py, pz) to line (
    x1, y1, z1) – (x2, y2, z2)
// (or ray, or segment; in the case of the ray
    , the endpoint is the
// first point)
public static final int LINE = 0;
public static final int SEGMENT = 1;
public static final int RAY = 2;
public static double ptLineDistSq(double x1,
    double y1, double z1,
    double x2, double y2, double z2, double px
        , double py, double pz,
    int type) {
  double pd2 = (x1-x2)*(x1-x2) + (y1-y2)*(y1-
      y2) + (z1-z2) * (z1-z2);
  double x, y, z;
  if (pd2 == 0) {
   x = x1;
   y = y1;
   z = z1;
  } else {
    double u = ((px-x1)*(x2-x1) + (py-y1)*(y2-x1)
        y1) + (pz-z1)*(z2-z1)) / pd2;
   x = x1 + u * (x2 - x1);
   y = y1 + u * (y2 - y1);
   z = z1 + u * (z2 - z1);
```

```
if (type != LINE && u < 0) {</pre>
     x = x1;
     y = y1;
      z = z1;
    if (type == SEGMENT && u > 1.0) {
     x = x2;
     y = y2;
      z = z2;
  return (x-px)*(x-px) + (y-py)*(y-py) + (z-pz)
      ) * (z-pz);
public static double ptLineDist(double x1,
    double y1, double z1,
    double x2, double y2, double z2, double px
        , double py, double pz,
    int type) {
  return Math.sqrt(ptLineDistSq(x1, y1, z1, x2
      , y2, z2, px, py, pz, type));
```

2.4 3D geometry - Ashley

```
struct point3 {
double x, y, z;
point3 (double x=0, double y=0, double z=0):x(x),
    y(y), z(z) \{ \}
point3 operator+(point3 p)const ?{ return point3
    (x + p.x, y)
+ p.y, z + p.z); }
point3 operator*(double k)const { return point3(
    k*x, k*y,
k * z);
point3 operator-(point3 p)const ?{ return *this
    + (p*-1.0);
point3 operator/(double k)const { return *this
    *(1.0/k);}
double norm() { return x*x + y*y + z*z; }
double abs() { return sqrt(norm()); }
point3 normalize() { return *this/this->abs(); }
};
// dot product
double dot(point3 a, point3 b) {
return a.x*b.x + a.y*b.y + a.z*b.z;
// cross product
point3 cross(point3 a, point3 b) {
return point3(a.y*b.z - b.y*a.z, b.x*a.z - a.x*b
    .z, a.x*b.y
- b.x*a.y);
struct line {
point3 a, b;
line(point3 A=point3(), point3 B=point3()) : a(A
    ), b(B) {}
// Direction unit vector a -> b
point3 dir() { return (b - a).normalize(); }
// Returns closest point on an infinite line u
    to the point p
point3 cpoint_iline(line u, point3 p) {
point3 ud = u.dir();
```

```
return u.a - ud*dot(u.a - p, ud);
// Returns Shortest distance between two
    infinite lines u and v
double dist_ilines(line u, line v) {
return dot(v.a - u.a, cross(u.dir(), v.dir()).
    normalize());
// Finds the closest point on infinite line u to
     infinite line v
// Note: if (uv*uv - uu*vv) is zero then the
    lines are parallel
// and such a single closest point does not
    exist. Check for
// this if needed.
point3 cpoint_ilines(line u, line v) {
point3 ud = u.dir(); point3 vd = v.dir();
double uu = dot(ud, ud), vv = dot(vd, vd), uv =
    dot(ud, vd);
double t = dot(u.a, ud) - dot(v.a, ud); t *= vv;
t = uv*(dot(u.a, vd) - dot(v.a, vd));
t /= (uv*uv - uu*vv);
return u.a + ud*t;
// Closest point on a line segment u to a given
    point p
point3 cpoint_lineseq(line u, point3 p) {
point3 ud = u.b - u.a; double s = dot(u.a - p)
ud) /ud.norm();
if (s < -1.0) return u.b;
if (s > ?0.0) return u.a;
return u.a - ud*s;
struct plane {
point3 n, p;
plane(point3 ni = point3(), point3 pi = point3()
    ) : n(ni),
p(pi) {}
plane(point3 a, point3 b, point3 c): n(cross(b-
    a, ca).normalize()), p(a) {}
//Value of d for the equation ax + by + cz + d =
double d() { return -dot(n, p); }
};
// Closest point on a plane u to a given point p
point3 cpoint_plane(plane u, point3 p) {
return p - u.n*(dot(u.n, p) + u.d());
// Point of intersection of an infinite line v
    and a plane u.
// Note: if dot(u.n, vd) == 0 then the line and
    plane do not
// intersect at a single point. Check for this
    if needed.
point3 iline isect plane(plane u, line v) {
point3 vd = v.dir();
return v.a - vd*((dot(u.n, v.a) + u.d())/dot(u.n
    , vd));
// Infinite line of intersection between two
    planes u and v.
// Note: if dot(v.n, uvu) == 0 then the planes
    do not intersect
// at a line. Check for this case if it is
    needed.
line isect planes (plane u, plane v) {
point3 o = u.n*-u.d(), uv = cross(u.n, v.n);
point3 uvu = cross(uv, u.n);
point3 a = o - uvu*((dot(v.n, o) + v.d())/(dot(v.n, o)))
```

```
n.
uvu) *uvu.norm()));
return line(a, a + uv);
// Returns great circle distance (lat[-90,90],
    long[-180,180])
double greatcircle (double lt1, double lo1,
    double 1t2, double
lo2, double r) {
double a = M_PI*(lt1/180.0), b = M_PI*(lt2
    /180.0);
double c = M_PI*((102-101)/180.0);
return r*acos(sin(a)*sin(b) + cos(a)*cos(b)*cos(
// Rotates point p around directed line a->b
    with angle 'theta'
point3 rotate(point3 a, point3 b, point3 p,
    double theta) {
point3 o = cpoint_iline(line(a,b),p);
point3 perp = cross(b-a,p-o);
return o+perp*sin(theta)+(p-o)*cos(theta);
```

2.5 Circles

```
// Returns whether they form a circle or not.
// 'center' and 'r' contain the circle if there
    is one
bool get_circle(point p1, point p2, point p3,
    point &center,
double &r) {
double g = 2 \times imag(conj(p2-p1) \times (p3-p2));
if (abs(g) < eps) return false;</pre>
center = p1*(norm(p3)-norm(p2));
center += p2*(norm(p1)-norm(p3));
center += p3*(norm(p2)-norm(p1));
center /= point(0, g); r = abs(p1-center);
return true;
// Returns number of circles that are tangent to
     all three lines
// 'cirs' has all possible circles with radius >
// It has zero circles when two of them are
    coincide
// It has two circles when only two of them are
    parallel
// It has four circles when they form a triangle
     . In this case
// first circle is incircle. Next circles are ex
     -circles tangent
// to edge a,b,c of triangle respectively.
int get_circle(point a1, point a2, point b1,
    point b2, point c1,
point c2, vector<circle> &cirs) {
point a,b,c;
int sa=line_line_inter(a1,a2,b1,b2,c);
int sb=line_line_inter(b1,b2,c1,c2,a);
int sc=line_line_inter(c1, c2, a1, a2, b);
if(sa==-1 || sb==-1 || sc==-1)
return 0:
if(sa+sb+sc==0)
```

```
return 0;
if(sb==0) {
swap(a1,c1);
swap (a2, c2);
if(sc==0) {
swap (b1, c1);
swap (b2, c2);
sa=line line inter(a1,a2,b1,b2,c);
line_line_inter(b1,b2,c1,c2,a);
line_line_inter(c1, c2, a1, a2, b);
if(sa==0) {
point v1 = polar(1.0, (arg(a2-a1) + arg(a-b))/2) + b;
point v2 = polar(1.0, (arg(a1-a2) + arg(a-b))/2) + b;
point v3 = polar(1.0, (arg(b2-b1) + arg(a-b))/2) + a;
point v4 = polar(1.0, (arg(b1-b2) + arg(a-b))/2) + a;
point p;
if(line\_line\_inter(b, v1, a, v3, p) == 0)
swap(v3, v4);
line_line_inter(b, v1, a, v3, p);
circle c1.c2;
c1.c = p;
line_line_inter(b, v2, a, v4, p);
c2.c = p;
c1.r = c2.r = abs(((a1-b1)/(b2-b1)).imag()*abs(
    h2-
b1))/2;
cirs.push_back(c1);
cirs.push back(c2);
} else {
if (abs(a-b) <eps)</pre>
return 0;
point bisec1[4][2];
point bisec2[4][2];
bisec1[0][0]=polar(1.0, (arg(c-a)+arg(b-a))/2);
bisec1[0][1]=a;
bisec2[0][0]=polar(1.0, (arg(c-b)+arg(a-b))/2);
bisec2[0][1]=b;
bisec1[1][0]=polar(1.0, (arg(c-a)+arg(b-a))/2);
bisec1[1][1]=a;
bisec2[1][0]=polar(1.0, (arg(c-b)+arg(b-a))/2);
bisec2[1][1]=b;
bisec1[2][0]=polar(1.0, (arg(a-b)+arg(c-b))/2);
bisec1[2][1]=b;
bisec2[2][0]=polar(1.0, (arg(a-c)+arg(c-b))/2);
bisec2[2][1]=c;
bisec1[3][0]=polar(1.0, (arg(b-c)+arg(a-c))/2);
bisec1[3][1]=b;
bisec2[3][0]=polar(1.0, (arg(b-a)+arg(a-c))/2);
bisec2[3][1]=c;
for(int i=0;i<4;i++) {</pre>
point p;
line_line_inter(bisec1[i][1],bisec1[i][1]+bisec1
[0],bisec2[i][1],bisec2[i][1]+bisec2[i][0],p);
circle c1;
c1.c = p;
c1.r = abs(((p-a)/(b-a)).imag())*abs(b-a);
cirs.push_back(c1);
return cirs.size();
// Returns number of circles that pass through
```

```
point a and b and
// are tangent to the line c-d
// 'ans' has all possible circles with radius >
int get_circle(point a, point b, point c, point
vector<circle> &ans) {
point pa = (a+b)/2.0;
point pb = (b-a) * point(0,1) + pa;
vector<point> ta;
parabola_line_inter(a,c,d,pa,pb,ta);
for(int i=0;i<ta.size();i++)</pre>
ans.push_back(circle(ta[i],abs(a-ta[i])));
return ans.size();
// Returns number of circles that pass through
    point p and are
// tangent to the lines a-b and c-d
// 'ans' has all possible circles with radius
    greater than zero
int get_circle(point p, point a, point b, point
    c, point d,
vector<circle> &ans) {
point inter;
int st = line_line_inter(a,b,c,d,inter);
if(st==-1) return 0;
d-=c;
b-=a;
vector<point> ta;
if(st==0) {
point pa = point (0, imag((a-c)/d)/2)*d+c;
point pb = b+pa;
parabola_line_inter(p,a,a+b,pa,pb,ta);
} else {
if(abs(inter-p)>eps) {
point bi;
bi = polar(1.0, (arg(b) + arg(d))/2) + inter;
vector<point> temp;
parabola_line_inter(p,a,a+b,inter,bi,temp);
ta.insert(ta.end(),temp.begin(),temp.end());
temp.clear();
bi = polar(1.0, (arg(b) + arg(d) + M_PI)/2) + inter;
parabola line inter(p,a,a+b,inter,bi,temp);
ta.insert(ta.end(),temp.begin(),temp.end());
for(int i=0;i<ta.size();i++)</pre>
ans.push_back(circle(ta[i],abs(p-ta[i])));
return ans.size();
```

2.6 Parabola Circle Intersection

```
// Find intersection of the line d-e and the
    parabola that
// is defined by point 'p' and line a-b
// Returns the number of intersections
// 'ans' has intersection points
int parabola_line_inter(point p, point a, point b, point d,
point e, vector<point> &ans) {
b = b-a;
p/=b; a/=b; d/=b; e/=b;
a-=p; d-=p; e-=p;
point n = (e-d) *point(0,1);
```

```
double c = -dot(n,e);
if(abs(n.imag()) < eps) {</pre>
if(abs(a.imag())>eps) {
double x = -c/n.real();
ans.push_back(point(x,a.imag()/2-x*x/(2*a.imag())
    )));
} else {
double aa = 1;
double bb = -2*a.imag()*n.real()/n.imag();
double cc = -2*a.imag()*c/n.imag()-a.imag()*a.
double delta = bb*bb-4*aa*cc;
if(delta>-eps) {
if(delta<0)</pre>
delta = 0;
delta = sqrt(delta);
double x = (-bb+delta)/(2*aa);
ans.push_back(point(x, (-c-n.real()*x)/n.imag()))
if(delta>eps) {
double x = (-bb-delta)/(2*aa);
ans.push back(point(x,(-cn.real()*x)/n.imag()));
for(int i=0;i<ans.size();i++)</pre>
ans[i] = (ans[i]+p)*b;
return ans.size();
```

3 Numerical algorithms

3.1 Number theory (modular, Chinese remainder, linear Diophantine)

```
// This is a collection of useful code for
    solving problems that
// involve modular linear equations. Note that
    all of the
// algorithms described here work on nonnegative
     integers.
#include <iostream>
#include <vector>
#include <algorithm>
using namespace std;
typedef vector<int> VI;
typedef pair<int, int> PII;
// return a % b (positive value)
int mod(int a, int b) {
        return ((a%b) + b) % b;
// computes gcd(a,b)
int gcd(int a, int b) {
        while (b) { int t = a%b; a = b; b = t; }
        return a;
// computes lcm(a,b)
```

```
int lcm(int a, int b) {
        return a / gcd(a, b) *b;
// (a^b) mod m via successive squaring
int powermod(int a, int b, int m)
        int ret = 1:
        while (b)
                if (b & 1) ret = mod(ret*a, m);
                a = mod(a*a, m);
                b >>= 1;
        return ret;
// returns g = gcd(a, b); finds x, y such that d
     = ax + by
int extended euclid(int a, int b, int &x, int &y
    ) {
        int xx = y = 0;
        int yy = x = 1;
        while (b) {
                int q = a / b;
                int t = b; b = a%b; a = t;
                t = xx; xx = x - q*xx; x = t;
                t = yy; yy = y - q*yy; y = t;
        return a:
// finds all solutions to ax = b \pmod{n}
VI modular linear equation solver(int a, int b,
    int n) {
        int x, y;
        VI ret;
        int g = extended_euclid(a, n, x, y);
        if (!(b%g)) {
                x = mod(x*(b / g), n);
                for (int i = 0; i < g; i++)</pre>
                        ret.push\_back(mod(x + i
                             *(n / q), n));
        return ret;
// computes b such that ab = 1 \pmod{n}, returns
    -1 on failure
int mod_inverse(int a, int n) {
        int x, y;
        int g = extended_euclid(a, n, x, y);
        if (q > 1) return -1;
        return mod(x, n);
// Chinese remainder theorem (special case):
    find z such that
// z % m1 = r1, z % m2 = r2. Here, z is unique
    modulo\ M = lcm(m1, m2).
// Return (z, M). On failure, M = -1.
PII chinese_remainder_theorem(int ml, int rl,
    int m2, int r2) {
        int s, t;
        int g = extended_euclid(m1, m2, s, t);
        if (r1%g != r2%g) return make pair(0,
            -1);
        return make_pair(mod(s*r2*m1 + t*r1*m2,
            m1*m2) / g, m1*m2 / g);
```

```
// Chinese remainder theorem: find z such that
// z % m[i] = r[i] for all i. Note that the
    solution is
// unique modulo M = lcm_i (m[i]). Return (z, M
// failure, M = -1. Note that we do not require
    the a[i]'s
// to be relatively prime.
PII chinese_remainder_theorem (const VI &m, const
        PII ret = make_pair(r[0], m[0]);
        for (int i = 1; i < m.size(); i++) {</pre>
                ret = chinese_remainder_theorem(
                    ret.second, ret.first, m[i
                     ], r[i]);
                if (ret.second == -1) break;
        return ret;
// computes x and y such that ax + by = c
// returns whether the solution exists
bool linear_diophantine(int a, int b, int c, int
     &x, int &y) {
        if (!a && !b)
                if (c) return false;
                x = 0; y = 0;
                return true;
        if (!a)
                if (c % b) return false;
                x = 0; y = c / b;
                return true;
        if (!b)
                if (c % a) return false;
                x = c / a; y = 0;
                return true;
        int q = qcd(a, b);
        if (c % g) return false;
        x = c / g * mod_inverse(a / g, b / g);
        y = (c - a * x) / b;
        return true;
int main() {
        // expected: 2
        cout << gcd(14, 30) << endl;
        // expected: 2 -2 1
        int g = extended_euclid(14, 30, x, y);
        cout << g << " " << x << " " << y <<
            endl;
        // expected: 95 451
        VI sols = modular_linear_equation_solver
             (14, 30, 100);
        for (int i = 0; i < sols.size(); i++)</pre>
            cout << sols[i] << " ";
        cout << endl;</pre>
        // expected: 8
```

```
cout << mod_inverse(8, 9) << endl;

// expected: 23 105
// 11 12
PII ret = chinese_remainder_theorem(VI({
      3, 5, 7 }), VI({ 2, 3, 2 }));
cout << ret.first << " " << ret.second
      << endl;
ret = chinese_remainder_theorem(VI({ 4,
      6 }), VI({ 3, 5 }));
cout << ret.first << " " << ret.second
      << endl;
// expected: 5 -15
if (!linear_diophantine(7, 2, 5, x, y))
      cout << "ERROR" << endl;
cout << x << " " << y << endl;
return 0;</pre>
```

3.2 Systems of linear equations, matrix inverse, determinant

```
// Gauss-Jordan elimination with full pivoting.
// Uses:
    (1) solving systems of linear equations (AX
    =B)
     (2) inverting matrices (AX=I)
    (3) computing determinants of square
    matrices
// Running time: O(n^3)
// INPUT:
             a[][] = an nxn matrix
             b[][] = an nxm matrix
// OUTPUT:
                    = an nxm matrix (stored in b
    [][])
             A^{-1} = an nxn matrix (stored in a
    [][])
             returns determinant of a[][]
#include <iostream>
#include <vector>
#include <cmath>
using namespace std;
const double EPS = 1e-10;
typedef vector<int> VI;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
T GaussJordan (VVT &a, VVT &b) {
 const int n = a.size();
  const int m = b[0].size();
  VI irow(n), icol(n), ipiv(n);
  T \det = 1;
  for (int i = 0; i < n; i++) {
    int p_{1} = -1, p_{k} = -1;
    for (int j = 0; j < n; j++) if (!ipiv[j])</pre>
      for (int k = 0; k < n; k++) if (!ipiv[k])
```

```
if (pj == -1 \mid | fabs(a[j][k]) > fabs(a[
             pj][pk])) { pj = j; pk = k; }
    if (fabs(a[pj][pk]) < EPS) { cerr << "Matrix</pre>
         is singular." << endl; exit(0); }
    ipiv[pk]++;
    swap(a[pj], a[pk]);
    swap(b[pj], b[pk]);
    if (pj != pk) det *= -1;
    irow[i] = pj;
    icol[i] = pk;
    T c = 1.0 / a[pk][pk];
    det *= a[pk][pk];
    a[pk][pk] = 1.0;
    for (int p = 0; p < n; p++) a[pk][p] *= c;</pre>
    for (int p = 0; p < m; p++) b[pk][p] *= c;
    for (int p = 0; p < n; p++) if (p != pk) {
      c = a[p][pk];
      a[p][pk] = 0;
      for (int q = 0; q < n; q++) a[p][q] -= a[
          pk][q] * c;
      for (int q = 0; q < m; q++) b[p][q] -= b[
          pk][q] * c;
  for (int p = n-1; p >= 0; p--) if (irow[p] !=
      icol[p]) {
    for (int k = 0; k < n; k++) swap(a[k][irow[p</pre>
        ]], a[k][icol[p]]);
  return det;
int main() {
 const int n = 4;
  const int m = 2;
 double A[n][n] = {
      \{1,2,3,4\},\{1,0,1,0\},\{5,3,2,4\},\{6,1,4,6\}
  double B[n][m] = \{ \{1,2\}, \{4,3\}, \{5,6\}, \{8,7\} \};
  VVT a(n), b(n);
  for (int i = 0; i < n; i++) {
   a[i] = VT(A[i], A[i] + n);
   b[i] = VT(B[i], B[i] + m);
 double det = GaussJordan(a, b);
 // expected: 60
 cout << "Determinant: " << det << endl;</pre>
  // expected: -0.233333 0.166667 0.133333
      0.0666667
               0.166667 0.166667 0.3333333
      -0.333333
               0.233333 0.833333 -0.133333
      -0.0666667
               0.05 -0.75 -0.1 0.2
 cout << "Inverse: " << endl;</pre>
 for (int i = 0; i < n; i++) {</pre>
   for (int j = 0; j < n; j++)
     cout << a[i][j] << ' ';
   cout << endl;</pre>
  // expected: 1.63333 1.3
               -0.166667 0.5
```

```
// 2.36667 1.7
// -1.85 -1.35
cout << "Solution: " << endl;
for (int i = 0; i < n; i++) {
   for (int j = 0; j < m; j++)
      cout << b[i][j] << ' ';
   cout << endl;
}</pre>
```

3.3 Reduced row echelon form, matrix rank

```
// Reduced row echelon form via Gauss-Jordan
    elimination
// with partial pivoting. This can be used for
    computing
// the rank of a matrix.
// Running time: O(n^3)
            a[][] = an nxm matrix
// OUTPUT:
            rref[][] = an nxm matrix (stored in
            returns rank of a[][]
#include <iostream>
#include <vector>
#include <cmath>
using namespace std;
const double EPSILON = 1e-10;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
int rref(VVT &a) {
 int n = a.size();
  int m = a[0].size();
  for (int c = 0; c < m && r < n; c++) {
   int j = r;
    for (int i = r + 1; i < n; i++)
     if (fabs(a[i][c]) > fabs(a[j][c])) j = i;
    if (fabs(a[j][c]) < EPSILON) continue;</pre>
    swap(a[j], a[r]);
    T s = 1.0 / a[r][c];
    for (int j = 0; j < m; j++) a[r][j] *= s;
    for (int i = 0; i < n; i++) if (i != r) {
     T t = a[i][c];
     for (int j = 0; j < m; j++) a[i][j] -= t *
           a[r][j];
    r++;
 return r;
int main() {
  const int n = 5, m = 4;
  double A[n][m] = {
    {16, 2, 3, 13},
```

```
{ 5, 11, 10, 8},
 { 9, 7, 6, 12},
 { 4, 14, 15, 1},
 {13, 21, 21, 13}};
VVT a(n);
for (int i = 0; i < n; i++)
 a[i] = VT(A[i], A[i] + m);
int rank = rref(a);
// expected: 3
cout << "Rank: " << rank << endl;</pre>
// expected: 1 0 0 1
          0 1 0 3
             0 0 1 -3
             0 0 0 3.10862e-15
            0 0 0 2.22045e-15
cout << "rref: " << endl;</pre>
for (int i = 0; i < 5; i++) {
  for (int j = 0; j < 4; j++)
   cout << a[i][j] << ' ';
  cout << endl;</pre>
```

3.4 Number Theory Essentials

// Sieve

```
ll sieve size;
bitset<10000010> bs:
vi primes;
void sieve(ll upperbound) {
    // create list of primes in [0..upperbound]
    _sieve_size = upperbound + 1;
    bs.set();bs[0] = bs[1] = 0;
    for (11 i = 2; i <= _sieve_size; i++) if (bs</pre>
        for (ll j = i * i; j <= _sieve_size; j</pre>
             += i) bs[j] = 0;
        primes.push_back((int)i);
bool isPrime(11 N) {
    if (N <= _sieve_size) return bs[N];</pre>
    for (int i = 0; i < (int)primes.size(); i++)
        if (N % primes[i] == 0) return false;
    return true;
// Prime Factors
vi primeFactors(ll N) {
    vi factors;
    11 PF_idx = 0, PF = primes[PF_idx];
    while (PF * PF <= N) {</pre>
        while (N % PF == 0) { N \neq PF; factors.
             push_back(PF); }
        PF = primes[++PF_idx];
    if (N != 1) factors.push_back(N);
    return factors;
// NumDiv
11 numDiv(ll N) {
    11 \text{ PF\_idx} = 0, PF = primes[PF\_idx], ans = 1;
    while (PF * PF <= N) {</pre>
```

```
11 power = 0;
        while (N % PF == 0) { N /= PF; power++;
        ans \star = (power + 1);
        PF = primes[++PF_idx];
    if (N != 1) ans \star= 2;
    return ans:
// SumDiv
ll sumDiv(ll N) {
    11 \text{ PF\_idx} = 0, PF = primes[PF\_idx], ans = 1;
    while (PF * PF <= N) {</pre>
        11 power = 0;
        while (N % PF == 0) { N /= PF; power++;
        ans \star= ((11)pow((double)PF, power + 1.0)
             -1) / (PF - 1);
        PF = primes[++PF_idx];
    if (N != 1) ans \star= ((11)pow((double)N, 2.0)
        -1) / (N - 1);
    return ans;
// EulerPhi
ll EulerPhi(ll N) {
   11 \text{ PF\_idx} = 0, PF = primes[PF\_idx], ans = N;
    while (PF * PF <= N) {
        if (N % PF == 0) ans -= ans / PF;
        while (N % PF == 0) N /= PF;
        PF = primes[++PF_idx];
    if (N != 1) ans -= ans / N;
    return ans;
```

4 Graph algorithms

4.1 Eulerian path

5 Data structures

5.1 Binary Indexed Tree

```
#include <iostream>
using namespace std;
#define LOGSZ 17
int tree[(1<<LOGSZ)+1];</pre>
int N = (1 << LOGSZ);
// add v to value at x
void set(int x, int v) {
  while (x \le N)
    tree[x] += v;
    x += (x & -x);
// get cumulative sum up to and including x
int get(int x) {
  int res = 0;
  while(x) {
    res += tree[x];
    x = (x \& -x);
  return res:
// get largest value with cumulative sum less
    than or equal to x;
// for smallest, pass x-1 and add 1 to result
int getind(int x) {
  int idx = 0, mask = N;
  while(mask && idx < N) {</pre>
    int t = idx + mask;
    if(x >= tree[t]) {
     idx = t;
      x -= tree[t];
    mask >>= 1;
  return idx;
```

6 Miscellaneous

6.1 Dates

```
// Routines for performing computations on dates
    . In these routines,
// months are expressed as integers from 1 to
    12, days are expressed
// as integers from 1 to 31, and years are
    expressed as 4-digit
// integers.
#include <iostream>
#include <string>
using namespace std:
string dayOfWeek[] = {"Mon", "Tue", "Wed", "Thu"
    , "Fri", "Sat", "Sun"};
// converts Gregorian date to integer (Julian
    day number)
int dateToInt (int m, int d, int y) {
  return
    1461 * (y + 4800 + (m - 14) / 12) / 4 +
    367 * (m - 2 - (m - 14) / 12 * 12) / 12 -
    3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +
    d = 32075:
// converts integer (Julian day number) to
    Gregorian date: month/day/year
void intToDate (int jd, int &m, int &d, int &y) {
 int x, n, i, j;
  x = id + 68569;
  n = 4 * x / 146097;
  x = (146097 * n + 3) / 4;
  i = (4000 * (x + 1)) / 1461001;
  x = 1461 * i / 4 - 31;
  \dot{1} = 80 * x / 2447;
  d = x - 2447 * j / 80;
  x = \frac{1}{2} / 11;
  m = j + 2 - 12 * x;
  y = 100 * (n - 49) + i + x;
// converts integer (Julian day number) to day
    of week
string intToDay (int jd) {
  return dayOfWeek[jd % 7];
int main (int argc, char **argv) {
  int jd = dateToInt (3, 24, 2004);
  int m, d, y;
  intToDate (jd, m, d, y);
  string day = intToDay (jd);
  // expected output:
        2453089
        3/24/2004
        Wed
```

6.2 Prime numbers

```
// O(sgrt(x)) Exhaustive Primality Test
#include <cmath>
#define EPS 1e-7
typedef long long LL;
bool IsPrimeSlow (LL x)
 if(x<=1) return false;</pre>
 if(x<=3) return true;</pre>
 if (!(x%2) || !(x%3)) return false;
 LL s=(LL) (sqrt ((double)(x))+EPS);
 for (LL i=5; i<=s; i+=6)
   if (!(x\%i) \mid | !(x\%(i+2))) return false;
 return true;
// Primes less than 1000:
       2
             3
                  5
                             11
                                   13
                                         17
            23
                  29
       19
                        31
                              37
            43
                  47
                        53
                             59
      41
                                         67
       71
            73
                  79
                        83
                              89
           101
                103
                      107
                            109
                                  113
                                        127
    131 137 139 149 151
     157
           163
                167
                      173
                            179
                                        191
    193
         197
              199
                     211
                           223
                233
           229
                      239
                            241
                                        257
          269
               271
                           281
          293
                307
                      311
                            313
                                  317
                                        331
     283
    337 347
               349
                     353
                           359
     367
          373
                379
                      383
                            389
                                  397
                                        401
         419
               421
                     431 433
     439
         443
                449
                      4.57
                            461
                                  463
                                        467
    479 487
               491 499
                           503
     509
         521
                 523
                      541
                            547
                                        563
         571
                577
                     587
                           593
     599
         601
                607
                      613
                            617
                                        631
         643
                647
                     653
                           659
                                  701
     661
         673
                677
                      683
                            691
                                        709
         727
                     739
    719
                733
                           743
                761
                            77.3
     751
          7.5.7
                      769
                                  787
                                        797
    809
         811
                     823
                           827
                821
     829
                853
                      857
                            859
          839
                                        877
         883
                     907
                           911
               937
         929
                      941
                            947
    971 977 983
     The largest prime smaller than 10 is 7.
     The largest prime smaller than 100 is 97.
     The largest prime smaller than 1000 is
    997.
    The largest prime smaller than 10000 is
     The largest prime smaller than 100000 is
     The largest prime smaller than 1000000 is
     The largest prime smaller than 10000000 is
     9999991.
```

```
The largest prime smaller than 100000000
is 99999989.
 The largest prime smaller than 1000000000
is 999999937.
The largest prime smaller than 10000000000
is 9999999967.
The largest prime smaller than
100000000000 is 9999999977.
The largest prime smaller than
10000000000000 is 9999999999999.
 The largest prime smaller than
100000000000000 is 999999999971.
 The largest prime smaller than
1000000000000000 is 9999999999973.
 The largest prime smaller than
The largest prime smaller than
100000000000000000 is 999999999999937.
The largest prime smaller than
The largest prime smaller than
```

6.3 C++ input/output

```
#include <iostream>
#include <iomanip>
using namespace std;
int main()
    // Ouput a specific number of digits past
        the decimal point,
    // in this case 5
    cout.setf(ios::fixed); cout << setprecision</pre>
    cout << 100.0/7.0 << endl;
    cout.unsetf(ios::fixed);
    // Output the decimal point and trailing
    cout.setf(ios::showpoint);
    cout << 100.0 << endl;
    cout.unsetf(ios::showpoint);
    // Output a '+' before positive values
    cout.setf(ios::showpos);
    cout << 100 << " " << -100 << endl;
    cout.unsetf(ios::showpos);
    // Output numerical values in hexadecimal
    cout << hex << 100 << " " << 1000 << " " <<
        10000 << dec << endl;
```

6.4 Knuth-Morris-Pratt

```
/*
Finds all occurrences of the pattern string p
   within the
text string t. Running time is O(n + m), where n
   and m
are the lengths of p and t, respecitively.
```

```
#include <iostream>
#include <string>
#include <vector>
using namespace std;
typedef vector<int> VI;
void buildPi(string& p, VI& pi)
 pi = VI(p.length());
 int k = -2;
  for(int i = 0; i < p.length(); i++) {</pre>
    while (k \ge -1 \&\& p[k+1] != p[i])
    k = (k == -1) ? -2 : pi[k];
    pi[i] = ++k;
int KMP(string& t, string& p)
  VI pi;
 buildPi(p, pi);
  int k = -1;
  for(int i = 0; i < t.length(); i++) {</pre>
    while (k >= -1 \&\& p[k+1] != t[i])
     k = (k == -1) ? -2 : pi[k];
    k++;
    if(k == p.length() - 1) {
      // p matches t[i-m+1, ..., i]
      cout << "matched at index " << i-k << ": "
      cout << t.substr(i-k, p.length()) << endl;</pre>
      k = (k == -1) ? -2 : pi[k];
  return 0:
int main()
  string a = "AABAACAADAABAABA", b = "AABA";
  KMP(a, b); // expected matches at: 0, 9, 12
  return 0;
```

6.5 Latitude/longitude

```
/*
Converts from rectangular coordinates to
    latitude/longitude and vice
versa. Uses degrees (not radians).
*/
#include <iostream>
#include <cmath>

using namespace std;
struct l1
{
    double r, lat, lon;
};
```

```
struct rect
  double x, y, z;
11 convert(rect& P)
 11 Q;
 Q.r = sqrt(P.x*P.x+P.y*P.y+P.z*P.z);
 O.lat = 180/M PI*asin(P.z/O.r);
 Q.lon = 180/M_PI*acos(P.x/sqrt(P.x*P.x+P.y*P.y)
  return Q;
rect convert(ll& Q)
 rect P:
 P.x = Q.r*cos(Q.lon*M_PI/180)*cos(Q.lat*M_PI
      /180);
 P.y = Q.r*sin(Q.lon*M_PI/180)*cos(Q.lat*M_PI
  P.z = Q.r*sin(Q.lat*M_PI/180);
  return P;
int main()
 rect A;
 11 B;
 A.x = -1.0; A.y = 2.0; A.z = -3.0;
 B = convert(A);
 cout << B.r << " " << B.lat << " " << B.lon <<
 A = convert(B);
 cout << A.x << " " << A.y << " " << A.z <<
```

6.6 Dates (Java)

```
// Example of using Java's built-in date
    calculation routines
import java.text.SimpleDateFormat;
import java.util.*;
public class Dates {
    public static void main(String[] args) {
        Scanner s = new Scanner(System.in);
        SimpleDateFormat sdf = new
            SimpleDateFormat("M/d/yyyy");
        while (true) {
            int n = s.nextInt();
            if (n == 0) break;
            GregorianCalendar c = new
                GregorianCalendar(n, Calendar.
                 JANUARY, 1);
            while (c.get(Calendar.DAY_OF_WEEK)
                 != Calendar.SATURDAY)
                c.add(Calendar.DAY_OF_YEAR, 1);
```

7 Hussain

7.1 Adaptive Simpson

```
// Adaptive Simpson's Rule (Wikipedia Article)
dbl adaptiveSimpsons(dbl (*f)(dbl), // ptr to
   dbl a, dbl b, // interval [a,b]
   dbl epsilon, // error tolerance
   int maxRecursionDepth) { // recursion cap
  db1 c = (a + b)/2, h = b - a;
  dbl fa = f(a), fb = f(b), fc = f(c);
  dbl S = (h/6) * (fa + 4*fc + fb);
  return adaptiveSimpsonsAux(f, a, b, epsilon, S
      , fa, fb, fc, maxRecursionDepth);
// Recursive auxiliary function for
    adaptiveSimpsons() function below
dbl adaptiveSimpsonsAux(dbl (*f)(dbl), dbl a,
    dbl b, dbl epsilon,
    dbl S, dbl fa, dbl fb, dbl fc, int bottom) {
  db1 c = (a + b)/2, h = b - a;
  db1 d = (a + c)/2, e = (c + b)/2;
  dbl fd = f(d), fe = f(e);
  dbl Sleft = (h/12)*(fa + 4*fd + fc);
  dbl Sright = (h/12)*(fc + 4*fe + fb);
  dbl S2 = Sleft + Sright;
  if (bottom \leq 0 \mid \mid fabs(S2 - S) \leq 15 * epsilon)
         // magic 15 comes from error analysis
    return S2 + (S2 - S)/15;
  return adaptiveSimpsonsAux(f, a, c, epsilon/2,
       Sleft, fa, fc, fd, bottom-1) +
         adaptiveSimpsonsAux(f, c, b, epsilon/2,
              Sright, fc, fb, fe, bottom-1);
int main(){
// compute integral of sin(x)
// from 0 to 2 and store it in
// the new variable I
 float I = adaptiveSimpsons(sin, 0, 2, 0.001,
 printf("I = %lf\n",I); // print the result
 return 0;
```

7.2 Binomial Coeff (constant N)

```
C[0] = 1

for (int k = 0; k < n; ++ k)

C[k+1] = (C[k] * (n-k)) / (k+1)

// C[i] = C(n,i)
```

7.3 Generate (x,y) pairs s.t. x AND y=y

7.4 Index of LSB

```
int msb(unsigned x) {
union { double a; int b[2]; };
a = x;
return (b[1] >> 20) - 1023;
}
```

8 Malek

8.1 Finding bridges in graph

```
int dfslow[N];
int dfsnum[N];
int dfscnt = 1;
vector<int> adj[N];
void dfs(int u, int p) {
 dfslow[u] = dfsnum[u] = dfscnt++;
  for (int i = 0; i < adj[u].size(); i++) {</pre>
    int v = adj[u][i];
    if (!dfsnum[v]) {
      dfs(v, u);
      if(dfslow[v]>dfsnum[u]){
        //it's a bridge
      dfslow[u]=min(dfslow[u],dfslow[v]);
    else if(v!=p)
      //back edge
      dfslow[u]=min(dfslow[u],dfsnum[v]);
```

8.2 LCA(Sparse Table) and Centroid Decomposition

```
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
typedef vector<int> vi;
#define lp(i,n) for(int i=0;i<(int)n;i++)
#define lp1(i,n) for (int i=1; i <= (int)n; i++)
const int N = 1e5 + 5;
const int LOGN = 20;
vi adj[N];
int dp[LOGN][N]; //sparse
int level[N];
int cs; //composition size
int sub[N];
bool cen[N];
int par[N];
int ans[N];
void dfs1(int u) {
 lp(i,adj[u].size())
    int v = adi[u][i];
    if (v != dp[0][u]) {
      level[v] = level[u] + 1;
      dp[0][v] = u;
      dfs1(v);
void sparse() {
  dfs1(0);
  lp1(i, LOGN-1)
    lp(j,n)
      dp[i][j] = dp[i - 1][dp[i - 1][j]];
int lca(int a, int b) {
  if (level[a] > level[b])
    swap(a, b);
  int dif = level[b] - level[a];
  lp(i,LOGN)
    if (dif & (1 << i))</pre>
      b = dp[i][b];
  if (a == b)
    return b;
  for (int i = LOGN - 1; i >= 0; i--) {
    if (dp[i][a] != dp[i][b])
      a = dp[i][a], b=dp[i][b];
  return dp[0][a];
int dist(int a, int b) {
  return level[a] + level[b] - 2 * level[lca(a,
      b)];
/*---rot&decay----*/
void dfssub(int u, int p) {
  sub[u] = 1;
  lp(i,adj[u].size())
    int v = adj[u][i];
```

```
if (!cen[v] && v != p)
      dfssub(v, u), sub[u] += sub[v];
int dfscen(int u, int p) {
 lp(i,adj[u].size())
    int v = adj[u][i];
    if (!cen[v] && v != p&&sub[v]>cs/2)
     return dfscen(v, u);
 return u;
void decomp(int root, int p) {
 cs = 0;
  dfssub(root, p);
 int centroid = dfscen(root, p);
// cout<<centroid<<endl;</pre>
 cen[centroid] = 1;
  par[centroid] = p;
  lp(i,adj[centroid].size())
    if (!cen[adj[centroid][i]])
      decomp(adj[centroid][i], centroid);
void update(int u) {
 int x = u;
  while (x != -1) {
    ans[x] = min(ans[x], dist(u, x));
    x = par[x];
int query(int u) {
 int x = u;
 int mn = ans[u];
 while (x != -1) {
   mn = min(mn, ans[x] + dist(u, x));
    x = par[x];
 return mn;
int main() {
 int m;
  sii(n, m);
    int x, y;
    lp(i, n-1)
     sii(x,y);
     x--; y--;
     adj[x].push_back(y);
     adj[y].push_back(x);
  sparse();
  decomp(0, -1);
  lp(i,n)ans[i]=1e9;
  update(0);
 while (m--) {
// int x, v;
     cin>>x>>y;
     cout << dist(x, y) << endl;</pre>
    int t,u;
    sii(t,u);
    if(t==1) update(u);
```

```
else printf("%d\n",query(u));
}
```

9 Marsil

9.1 2D geomtry using Complex

```
src: http://codeforces.com/blog/entry/22175
Functions using std::complex
1) Vector addition: a + b
2) Scalar multiplication: r * a
3) Dot product: (conj(a) * b).x
4) Cross product: (conj(a) * b).v
5) notice: conj(a) * b = (ax*bx + ay*by) + i
              av*bx)
6) Squared distance: norm(a - b)
7) Euclidean distance: abs(a - b)
8) Angle of elevation: arg(b - a)
9) Slope of line (a, b): tan(arg(b - a))
10) Polar to cartesian: polar(r, theta)
11) Cartesian to polar: point(abs(p), arg(p))
12) Rotation about the origin: a * polar(1.0,
    theta)
13) Rotation about pivot p: (a-p) * polar(1.0,
    theta) + p
14) Angle ABC: abs(remainder(arg(a-b) - arg(c-b)
    , 2.0 * M_PI))
      remainder normalizes the angle to be
           between [-PI, PI]. Thus, we can get
           the positive non-reflex angle by
           taking its abs value.
15) Project p onto vector v: v * dot(p, v) /
    norm(v):
16) Project p onto line (a, b): a + (b - a) *
    dot(p - a, b - a) / norm(b - a)
17) Reflect p across line (a, b): a + conj((p -
    a) / (b - a) ) * (b - a)
18) Intersection of line (a, b) and (p, q):
point intersection (point a, point b, point p,
  double c1 = cross(p - a, b - a), c2 = cross(q
      -a, b-a);
  return (c1 * q - c2 * p) / (c1 - c2); //
      undefined if parallel
Drawbacks:
Using std::complex is very advantageous, but it
    has one disadvantage: you can't use std::
    cin or scanf. Also, if we macro x and y, we
     can't use them as variables. But that's
    rather minor, don't you think?
EDIT: Credits to Zlobober for pointing out that
    std::complex has issues with integral data
    types. The library will work for simple
    arithmetic like vector addition and such.
    but not for polar or abs. It will compile
    but there will be some errors in
    correctness! The tip then is to rely on the
     library only if you're using floating
    point data all throughout.
```

9.2 bottom up lasy segment tree

```
template<typename T, typename U> struct
    seq_tree_lazy {
    int S, H;
   T zero;
    vector<T> value;
   U noop;
    vector<bool> dirty;
   vector<U> prop;
    seg_tree_lazy<T, U>(int _S, T _zero = T(), U
         _{noop} = U()) {
        zero = _zero, noop = _noop;
        for (S = 1, H = 1; S < \_S; ) S *= 2, H
        value.resize(2*S, zero);
       dirty.resize(2*S, false);
       prop.resize(2*S, noop);
   void set_leaves(vector<T> &leaves) {
        copy(leaves.begin(), leaves.end(), value
            .begin() + S);
        for (int i = S - 1; i > 0; i--)
           value[i] = value[2 * i] + value[2 *
                i + 11;
    void apply(int i, U &update) {
       value[i] = update(value[i]);
       if(i < S) {
           prop[i] = prop[i] + update;
            dirty[i] = true;
   void rebuild(int i) {
        for (int 1 = i/2; 1; 1 /= 2) {
            T combined = value[2*1] + value[2*1]
            value[1] = prop[1](combined);
    void propagate(int i) {
        for (int h = H; h > 0; h--) {
            int 1 = i \gg h:
           if (dirty[1]) {
                apply(2*1, prop[1]);
                apply(2*l+1, prop[1]);
                prop[1] = noop;
                dirty[1] = false;
    void upd(int i, int j, U update) {
        i += S, i += S;
        propagate(i), propagate(j);
```

```
for (int 1 = i, r = j; 1 \le r; 1 \ne 2, r
             /= 2) {
            if((1&1) == 1) apply(1++, update);
            if((r\&1) == 0) apply(r--, update);
        rebuild(i), rebuild(j);
    T query(int i, int j) {
        i += S, j += S;
        propagate(i), propagate(j);
        T res_left = zero, res_right = zero;
        for(; i <= j; i /= 2, j /= 2) {
            if((i&1) == 1) res_left = res_left +
                 value[i++];
            if((j&1) == 0) res_right = value[j
                --] + res_right;
        return res_left + res_right;
};
As an example, let's see how to use it to
    support the follow operations:
    Type 1: Add amount V to the values in range
        [L, R].
    Type 2: Reset the values in range [L, R] to
        value V.
    Type 3: Query for the sum of the values in
        range [L, R].
//The T struct would look like this:
```

```
struct node {
   int sum, width;
   node operator+(const node &n) {
        return { sum + n.sum, width + n.width };
};
//And the U struct would look like this:
struct update {
   bool type; // 0 for add, 1 for reset
   int value;
   node operator()(const node &n) {
        if (type) return { n.width * value, n.
            width };
        else return { n.sum + n.width * value, n
            .width };
   update operator+(const update &u) {
        if (u.type) return u;
        return { type, value + u.value };
};
```

9.3 Ordered Statistics Tree

```
#define lp(i,n) for (int i = 0; i < (int) (n); ++
i)</pre>
```

```
#include <ext/pb ds/assoc container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __qnu_pbds;
typedef tree<</pre>
double,
int,
less<double>,
rb_tree_tag,
tree_order_statistics_node_update> map_t;
typedef tree<</pre>
int,
null_type,
less<int>,
rb_tree_tag,
tree_order_statistics_node_update> ordered_set;
int main(){
    map_t a;
    a[0] = 1;
    a[2] = 1;
    a[5] = 1;
    cout << a.find_by_order(1) -> first << endl;</pre>
    cout << a.order_of_key(-5) << endl;</pre>
```