Al-Baath University ICPC Team Notebook (2018)

1 Combinatorial optimization

1.1 Sparse max-flow

```
// Adjacency list implementation of Dinic's blocking flow algorithm.
// This is very fast in practice, and only loses to push-relabel flow.
// Running time:
       O(|V|^2 |E|)
// INPUT:
       - graph, constructed using AddEdge()
       - source and sink
      - maximum flow value
       - To obtain actual flow values, look at edges with capacity > 0
         (zero capacity edges are residual edges).
#include < cstdio >
#include < vector >
#include < queue >
using namespace std;
typedef long long LL;
struct Edge {
 int u, v;
  LL cap, flow;
 Edge() {}
 Edge(int u, int v, LL cap): u(u), v(v), cap(cap), flow(0) {}
struct Dinic {
 int N;
  vector<Edge> E;
 vector<vector<int>> g;
 vector<int> d, pt;
 Dinic(int N): N(N), E(0), g(N), d(N), pt(N) {}
  void AddEdge(int u, int v, LL cap) {
    if (u != v) {
      E.emplace_back(Edge(u, v, cap));
      g[u].emplace_back(E.size() - 1);
      E.emplace_back(Edge(v, u, 0));
      g[v].emplace_back(E.size() - 1);
 bool BFS(int S, int T) {
    queue<int> q({S});
    fill(d.begin(), d.end(), N + 1);
    d[S] = 0;
    while(!q.empty()) {
      int u = q.front(); q.pop();
      if (u == T) break;
      for (int k: g[u]) {
        Edge &e = E[k];
        if (e.flow < e.cap && d[e.v] > d[e.u] + 1) {
          d[e.v] = d[e.u] + 1;
          q.emplace(e.v);
```

```
return d[T] != N + 1;
  LL DFS (int u, int T, LL flow = -1) {
    if (u == T || flow == 0) return flow;
    for (int &i = pt[u]; i < g[u].size(); ++i) {</pre>
      Edge &e = E[g[u][i]];
      Edge &oe = E[g[u][i]^1];
      if (d[e.v] == d[e.u] + 1) {
        LL amt = e.cap - e.flow;
        if (flow != -1 && amt > flow) amt = flow;
        if (LL pushed = DFS(e.v, T, amt)) {
          e.flow += pushed;
          oe.flow -= pushed;
          return pushed;
    return 0;
  LL MaxFlow(int S, int T) {
    LL total = 0;
    while (BFS(S, T)) {
      fill(pt.begin(), pt.end(), 0);
      while (LL flow = DFS(S, T))
        total += flow;
    return total;
};
// The following code solves SPOJ problem #4110: Fast Maximum Flow (
    FASTFLOW)
int main()
  int N, E;
  scanf("%d%d", &N, &E);
 Dinic dinic(N);
  for (int i = 0; i < E; i++)
    int u, v;
    LL cap;
    scanf("%d%d%lld", &u, &v, &cap);
    dinic.AddEdge(u - 1, v - 1, cap);
   dinic.AddEdge(v - 1, u - 1, cap);
  printf("%lld\n", dinic.MaxFlow(0, N - 1));
  return 0;
// END CUT
```

1.2 Min-cost max-flow

```
// Implementation of min cost max flow algorithm using adjacency // matrix (Edmonds and Karp 1972). This implementation keeps track of // forward and reverse edges separately (so you can set cap[i][j] != // cap[j][i]). For a regular max flow, set all edge costs to 0. // Running time, O(|V|^2) cost per augmentation // max flow: O(|V|^3) augmentations
```

```
min cost max flow: O(|V|^4 * MAX_EDGE_COST) augmentations
// INPUT:
11
      - graph, constructed using AddEdge()
       - source
       - sink
// OUTPUT:
      - (maximum flow value, minimum cost value)
       - To obtain the actual flow, look at positive values only.
#include <cmath>
#include <vector>
#include <iostream>
using namespace std;
typedef vector<int> VI;
typedef vector<VI> VVI;
typedef long long L;
typedef vector<L> VL;
typedef vector<VL> VVL;
typedef pair<int, int> PII;
typedef vector<PII> VPII;
const L INF = numeric_limits<L>::max() / 4;
struct MinCostMaxFlow {
 int N;
  VVL cap, flow, cost;
  VI found:
 VL dist, pi, width;
 VPII dad;
 MinCostMaxFlow(int N) :
    N(N), cap(N, VL(N)), flow(N, VL(N)), cost(N, VL(N)),
    found(N), dist(N), pi(N), width(N), dad(N) {}
  void AddEdge(int from, int to, L cap, L cost) {
    this->cap[from][to] = cap;
    this->cost[from][to] = cost;
  void Relax(int s, int k, L cap, L cost, int dir) {
   L val = dist[s] + pi[s] - pi[k] + cost;
    if (cap && val < dist[k]) {</pre>
      dist[k] = val;
     dad[k] = make_pair(s, dir);
     width[k] = min(cap, width[s]);
  L Dijkstra(int s, int t) {
    fill(found.begin(), found.end(), false);
    fill(dist.begin(), dist.end(), INF);
    fill(width.begin(), width.end(), 0);
    dist[s] = 0;
    width[s] = INF;
    while (s != -1) {
     int best = -1;
      found[s] = true;
      for (int k = 0; k < N; k++) {
        if (found[k]) continue;
        Relax(s, k, cap[s][k] - flow[s][k], cost[s][k], 1);
        Relax(s, k, flow[k][s], -cost[k][s], -1);
        if (best == -1 \mid \mid dist[k] < dist[best]) best = k;
      s = best;
```

```
for (int k = 0; k < N; k++)
     pi[k] = min(pi[k] + dist[k], INF);
    return width[t];
  pair<L, L> GetMaxFlow(int s, int t) {
   L \text{ totflow} = 0, totcost = 0;
    while (L amt = Dijkstra(s, t)) {
      totflow += amt;
      for (int x = t; x != s; x = dad[x].first) {
        if (dad[x].second == 1) {
          flow[dad[x].first][x] += amt;
          totcost += amt * cost[dad[x].first][x];
        } else {
          flow[x][dad[x].first] -= amt;
          totcost -= amt * cost[x][dad[x].first];
    return make_pair(totflow, totcost);
};
// BEGIN CUT
// The following code solves UVA problem #10594: Data Flow
int main() {
 int N, M;
  while (scanf("%d%d", &N, &M) == 2) {
   VVL v(M, VL(3));
    for (int i = 0; i < M; i++)
     scanf("%Ld%Ld%Ld", &v[i][0], &v[i][1], &v[i][2]);
    L D, K;
    scanf("%Ld%Ld", &D, &K);
    MinCostMaxFlow mcmf(N+1);
    for (int i = 0; i < M; i++) {
     mcmf.AddEdge(int(v[i][0]), int(v[i][1]), K, v[i][2]);
     mcmf.AddEdge(int(v[i][1]), int(v[i][0]), K, v[i][2]);
    mcmf.AddEdge(0, 1, D, 0);
    pair<L, L> res = mcmf.GetMaxFlow(0, N);
    if (res.first == D) {
     printf("%Ld\n", res.second);
    } else {
     printf("Impossible.\n");
  return 0;
// END CUT
```

1.3 Push-relabel max-flow

```
// Adjacency list implementation of FIFO push relabel maximum flow // with the gap relabeling heuristic. This implementation is // significantly faster than straight Ford-Fulkerson. It solves // random problems with 10000 vertices and 1000000 edges in a few // seconds, though it is possible to construct test cases that // achieve the worst-case.
```

```
// Running time:
       0(|V|^3)
// INPUT:
       - graph, constructed using AddEdge()
       - source
       - sink
// OUTPUT:
      - maximum flow value
       - To obtain the actual flow values, look at all edges with
         capacity > 0 (zero capacity edges are residual edges).
#include <cmath>
#include <vector>
#include <iostream>
#include <queue>
using namespace std;
typedef long long LL;
struct Edge {
  int from, to, cap, flow, index;
  Edge(int from, int to, int cap, int flow, int index) :
    from(from), to(to), cap(cap), flow(flow), index(index) {}
};
struct PushRelabel {
  int N;
  vector<vector<Edge> > G;
  vector<LL> excess;
 vector<int> dist, active, count;
  queue<int> Q;
  PushRelabel(int N): N(N), G(N), excess(N), dist(N), active(N), count(2*N)
  void AddEdge(int from, int to, int cap) {
    G[from].push_back(Edge(from, to, cap, 0, G[to].size()));
    if (from == to) G[from].back().index++;
    G[to].push_back(Edge(to, from, 0, 0, G[from].size() - 1));
  void Enqueue(int v) {
    if (!active[v] && excess[v] > 0) { active[v] = true; Q.push(v); }
  void Push(Edge &e) {
    int amt = int(min(excess[e.from], LL(e.cap - e.flow)));
    if (dist[e.from] <= dist[e.to] || amt == 0) return;</pre>
    e.flow += amt;
    G[e.to][e.index].flow -= amt;
    excess[e.to] += amt;
    excess[e.from] -= amt;
    Enqueue(e.to);
  void Gap(int k) {
    for (int v = 0; v < N; v++) {
      if (dist[v] < k) continue;</pre>
      count[dist[v]]--;
      dist[v] = max(dist[v], N+1);
      count[dist[v]]++;
      Enqueue (v);
```

```
void Relabel(int v) {
    count[dist[v]]--;
    dist[v] = 2*N;
    for (int i = 0; i < G[v].size(); i++)
     if (G[v][i].cap - G[v][i].flow > 0)
        dist[v] = min(dist[v], dist[G[v][i].to] + 1);
    count[dist[v]]++;
    Enqueue (v);
  void Discharge(int v) {
    for (int i = 0; excess[v] > 0 && i < G[v].size(); i++) Push(G[v][i]);
    if (excess[v] > 0) {
      if (count[dist[v]] == 1)
        Gap(dist[v]);
      else
        Relabel(v);
  LL GetMaxFlow(int s, int t) {
    count[0] = N-1;
    count[N] = 1;
    dist[s] = N;
    active[s] = active[t] = true;
    for (int i = 0; i < G[s].size(); i++) {</pre>
      excess[s] += G[s][i].cap;
      Push(G[s][i]);
    while (!Q.empty()) {
      int v = Q.front();
      Q.pop();
      active[v] = false;
      Discharge(v);
    LL totflow = 0;
    for (int i = 0; i < G[s].size(); i++) totflow += G[s][i].flow;</pre>
    return totflow;
};
// The following code solves SPOJ problem #4110: Fast Maximum Flow (
int main() {
  int n, m;
  scanf("%d%d", &n, &m);
 PushRelabel pr(n);
  for (int i = 0; i < m; i++) {</pre>
   int a, b, c;
    scanf("%d%d%d", &a, &b, &c);
    if (a == b) continue;
    pr.AddEdge(a-1, b-1, c);
    pr.AddEdge(b-1, a-1, c);
 printf("%Ld\n", pr.GetMaxFlow(0, n-1));
  return 0;
// END CUT
```

```
// Min cost bipartite matching via shortest augmenting paths
// This is an O(n^3) implementation of a shortest augmenting path
// algorithm for finding min cost perfect matchings in dense
// graphs. In practice, it solves 1000x1000 problems in around 1
// second.
    cost[i][j] = cost for pairing left node i with right node j
    Lmate[i] = index of right node that left node i pairs with
    Rmate[j] = index of left node that right node j pairs with
// The values in cost[i][j] may be positive or negative. To perform
// maximization, simply negate the cost[][] matrix.
#include <algorithm>
#include <cstdio>
#include <cmath>
#include <vector>
using namespace std;
typedef vector<double> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
double MinCostMatching(const VVD &cost, VI &Lmate, VI &Rmate) {
 int n = int(cost.size());
  // construct dual feasible solution
 VD u(n);
 VD v(n);
  for (int i = 0; i < n; i++) {
   u[i] = cost[i][0];
    for (int j = 1; j < n; j++) u[i] = min(u[i], cost[i][j]);</pre>
  for (int j = 0; j < n; j++) {
   v[j] = cost[0][j] - u[0];
    for (int i = 1; i < n; i++) v[j] = min(v[j], cost[i][j] - u[i]);</pre>
  // construct primal solution satisfying complementary slackness
  Lmate = VI(n, -1);
  Rmate = VI(n, -1);
  int mated = 0;
  for (int i = 0; i < n; i++) {
   for (int j = 0; j < n; j++) {
      if (Rmate[j] != -1) continue;
      if (fabs(cost[i][j] - u[i] - v[j]) < 1e-10) {</pre>
       Lmate[i] = j;
       Rmate[j] = i;
       mated++;
       break;
 VD dist(n);
 VI dad(n);
 VI seen(n);
  // repeat until primal solution is feasible
  while (mated < n) {</pre>
    // find an unmatched left node
    int s = 0;
    while (Lmate[s] != -1) s++;
```

```
// initialize Dijkstra
  fill(dad.begin(), dad.end(), -1);
  fill(seen.begin(), seen.end(), 0);
  for (int k = 0; k < n; k++)
   dist[k] = cost[s][k] - u[s] - v[k];
  int i = 0;
  while (true) {
    // find closest
    \frac{1}{1} = -1;
    for (int k = 0; k < n; k++) {
      if (seen[k]) continue;
      if (j == -1 || dist[k] < dist[j]) j = k;</pre>
    seen[j] = 1;
    // termination condition
    if (Rmate[j] == -1) break;
    // relax neighbors
    const int i = Rmate[j];
    for (int k = 0; k < n; k++) {
      if (seen[k]) continue;
      const double new_dist = dist[j] + cost[i][k] - u[i] - v[k];
      if (dist[k] > new_dist) {
       dist[k] = new_dist;
        dad[k] = j;
  // update dual variables
  for (int k = 0; k < n; k++) {
    if (k == j || !seen[k]) continue;
    const int i = Rmate[k];
   v[k] += dist[k] - dist[j];
   u[i] -= dist[k] - dist[j];
  u[s] += dist[j];
  // augment along path
  while (dad[i] >= 0) {
   const int d = dad[i];
   Rmate[j] = Rmate[d];
   Lmate[Rmate[j]] = j;
    j = d;
  Rmate[j] = s;
  Lmate[s] = j;
  mated++;
double value = 0;
for (int i = 0; i < n; i++)
  value += cost[i][Lmate[i]];
return value;
```

1.5 Max bipartite matchine

```
INPUT: w[i][j] = edge between row node i and column node j
     OUTPUT: mr[i] = assignment for row node i, -1 if unassigned
             mc[i] = assignment for column node j, -1 if unassigned
             function returns number of matches made
#include <vector>
using namespace std;
typedef vector<int> VI;
typedef vector<VI> VVI;
bool FindMatch (int i, const VVI &w, VI &mr, VI &mc, VI &seen) {
  for (int j = 0; j < w[i].size(); j++) {</pre>
    if (w[i][j] && !seen[j]) {
      seen[j] = true;
      if (mc[j] < 0 || FindMatch(mc[j], w, mr, mc, seen)) {</pre>
       mr[i] = j;
        mc[j] = i;
        return true;
  return false:
int BipartiteMatching(const VVI &w, VI &mr, VI &mc) {
 mr = VI(w.size(), -1);
 mc = VI(w[0].size(), -1);
  int ct = 0;
  for (int i = 0; i < w.size(); i++) {</pre>
   VI seen(w[0].size());
    if (FindMatch(i, w, mr, mc, seen)) ct++;
  return ct;
```

1.6 Global min-cut

```
// Adjacency matrix implementation of Stoer-Wagner min cut algorithm.
// Running time:
      0(|V|^3)
// INPUT:
      - graph, constructed using AddEdge()
// OUTPUT:
      - (min cut value, nodes in half of min cut)
#include <cmath>
#include <vector>
#include <iostream>
using namespace std;
typedef vector<int> VI;
typedef vector<VI> VVI;
const int INF = 1000000000;
pair<int, VI> GetMinCut(VVI &weights) {
  int N = weights.size();
  VI used(N), cut, best_cut;
```

```
int best_weight = -1;
  for (int phase = N-1; phase >= 0; phase--) {
   VI w = weights[0];
    VI added = used;
    int prev, last = 0;
    for (int i = 0; i < phase; i++) {</pre>
     prev = last;
      last = -1;
      for (int j = 1; j < N; j++)
        if (!added[j] && (last == -1 \mid \mid w[j] > w[last])) last = j;
      if (i == phase-1) {
        for (int j = 0; j < N; j++) weights[prev][j] += weights[last][j];</pre>
        for (int j = 0; j < N; j++) weights[j][prev] = weights[prev][j];</pre>
        used[last] = true;
        cut.push_back(last);
        if (best_weight == -1 || w[last] < best_weight) {</pre>
          best_cut = cut;
          best_weight = w[last];
      } else {
        for (int j = 0; j < N; j++)
         w[j] += weights[last][j];
        added[last] = true;
 return make_pair(best_weight, best_cut);
// The following code solves UVA problem #10989: Bomb, Divide and Conquer
int main() {
 int N:
  cin >> N;
  for (int i = 0; i < N; i++) {
    int n, m;
    cin >> n >> m;
    VVI weights(n, VI(n));
    for (int j = 0; j < m; j++) {
     int a, b, c;
     cin >> a >> b >> c;
     weights[a-1][b-1] = weights[b-1][a-1] = c;
    pair<int, VI> res = GetMinCut(weights);
    cout << "Case #" << i+1 << ": " << res.first << endl;
// END CUT
```

2 Geometry

2.1 Convex hull

```
// Compute the 2D convex hull of a set of points using the monotone chain
// algorithm. Eliminate redundant points from the hull if REMOVE_REDUNDANT
    is
// #defined.
//
// Running time: O(n log n)
//
// INPUT: a vector of input points, unordered.
// OUTPUT: a vector of points in the convex hull, counterclockwise,
    starting
// with bottommost/leftmost point
```

```
#include <cstdio>
#include <cassert>
#include <vector>
#include <algorithm>
#include <cmath>
// BEGIN CUT
#include <map>
// END CUT
using namespace std;
#define REMOVE_REDUNDANT
typedef double T;
const T EPS = 1e-7;
struct PT {
 T x, y;
 PT() {}
 PT(T x, T y) : x(x), y(y) \{ \}
 bool operator<(const PT &rhs) const { return make_pair(y,x) < make_pair(</pre>
      rhs.v.rhs.x): }
 bool operator == (const PT &rhs) const { return make pair(y,x) == make pair
       (rhs.y,rhs.x); }
T cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
T area2(PT a, PT b, PT c) { return cross(a,b) + cross(b,c) + cross(c,a); }
#ifdef REMOVE REDUNDANT
bool between (const PT &a, const PT &b, const PT &c) {
  return (fabs(area2(a,b,c)) < EPS && (a.x-b.x)*(c.x-b.x) <= 0 && (a.y-b.y)
      *(c.y-b.y) <= 0);
#endif
void ConvexHull(vector<PT> &pts) {
 sort(pts.begin(), pts.end());
  pts.erase(unique(pts.begin(), pts.end()), pts.end());
 vector<PT> up, dn;
  for (int i = 0; i < pts.size(); i++) {</pre>
    while (up.size() > 1 \&\& area2(up[up.size()-2], up.back(), pts[i]) >= 0)
         up.pop_back();
    while (dn.size() > 1 \&\& area2(dn[dn.size()-2], dn.back(), pts[i]) <= 0)
         dn.pop back();
    up.push_back(pts[i]);
    dn.push_back(pts[i]);
  for (int i = (int) up.size() - 2; i >= 1; i--) pts.push_back(up[i]);
#ifdef REMOVE REDUNDANT
  if (pts.size() <= 2) return;</pre>
  dn.clear();
  dn.push back(pts[0]);
  dn.push_back(pts[1]);
  for (int i = 2; i < pts.size(); i++) {</pre>
    if (between(dn[dn.size()-2], dn[dn.size()-1], pts[i])) dn.pop_back();
    dn.push_back(pts[i]);
  if (dn.size() >= 3 && between(dn.back(), dn[0], dn[1])) {
    dn[0] = dn.back();
    dn.pop_back();
 pts = dn;
#endif
// BEGIN CUT
```

```
// The following code solves SPOJ problem #26: Build the Fence (BSHEEP)
int main() {
  int t;
  scanf("%d", &t);
  for (int caseno = 0; caseno < t; caseno++) {</pre>
    int n;
    scanf("%d", &n);
    vector<PT> v(n);
    for (int i = 0; i < n; i++) scanf("%lf%lf", &v[i].x, &v[i].y);</pre>
    vector<PT> h(v);
    map<PT, int> index;
    for (int i = n-1; i >= 0; i--) index[v[i]] = i+1;
    ConvexHull(h);
    double len = 0;
    for (int i = 0; i < h.size(); i++) {</pre>
      double dx = h[i].x - h[(i+1)%h.size()].x;
      double dy = h[i].y - h[(i+1)%h.size()].y;
      len += sqrt(dx*dx+dy*dy);
    if (caseno > 0) printf("\n");
    printf("%.2f\n", len);
    for (int i = 0; i < h.size(); i++) {</pre>
      if (i > 0) printf(" ");
      printf("%d", index[h[i]]);
    printf("\n");
// END CUT
```

2.2 Miscellaneous geometry

```
// C++ routines for computational geometry.
#include <iostream>
#include <vector>
#include <cmath>
#include <cassert>
using namespace std;
double INF = 1e100;
double EPS = 1e-12;
struct PT {
  double x, y;
  PT() {}
  PT(double x, double y) : x(x), y(y) {}
  PT(const PT &p) : x(p.x), y(p.y)
  PT operator + (const PT &p) const { return PT(x+p.x, y+p.y); }
  PT operator - (const PT &p) const {
                                      return PT(x-p.x, y-p.y);
  PT operator * (double c)
                               const {
                                      return PT(x*c, y*c);
  PT operator / (double c)
                               const { return PT(x/c, y/c);
                          { return p.x*q.x+p.y*q.y; }
double dot (PT p, PT q)
double dist2(PT p, PT q) { return dot(p-q,p-q); }
double cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
ostream & operator << (ostream &os, const PT &p) {
  return os << "(" << p.x << "," << p.y << ")";
// rotate a point CCW or CW around the origin
```

```
PT RotateCCW90 (PT p) { return PT (-p.y,p.x); }
PT RotateCW90 (PT p)
                     { return PT(p.y,-p.x); }
PT RotateCCW(PT p, double t) {
  return PT(p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.y*cos(t));
// project point c onto line through a and b
// assuming a != b
PT ProjectPointLine(PT a, PT b, PT c) {
  return a + (b-a) *dot(c-a, b-a) /dot(b-a, b-a);
// project point c onto line segment through a and b
PT ProjectPointSegment (PT a, PT b, PT c) {
  double r = dot(b-a,b-a);
  if (fabs(r) < EPS) return a;</pre>
  r = dot(c-a, b-a)/r;
  if (r < 0) return a;</pre>
 if (r > 1) return b;
 return a + (b-a) *r;
// compute distance from c to segment between a and b
double DistancePointSegment(PT a, PT b, PT c) {
  return sqrt(dist2(c, ProjectPointSegment(a, b, c)));
// compute distance between point (x,y,z) and plane ax+by+cz=d
double DistancePointPlane(double x, double y, double z,
                          double a, double b, double c, double d)
 return fabs(a*x+b*y+c*z-d)/sqrt(a*a+b*b+c*c);
// determine if lines from a to b and c to d are parallel or collinear
bool LinesParallel(PT a, PT b, PT c, PT d) {
  return fabs(cross(b-a, c-d)) < EPS;</pre>
bool LinesCollinear(PT a, PT b, PT c, PT d) {
  return LinesParallel(a, b, c, d)
      && fabs(cross(a-b, a-c)) < EPS
      && fabs(cross(c-d, c-a)) < EPS;
// determine if line segment from a to b intersects with
// line segment from c to d
bool SegmentsIntersect(PT a, PT b, PT c, PT d) {
  if (LinesCollinear(a, b, c, d)) {
    if (dist2(a, c) < EPS || dist2(a, d) < EPS ||</pre>
      dist2(b, c) < EPS || dist2(b, d) < EPS) return true;
    if (dot(c-a, c-b) > 0 && dot(d-a, d-b) > 0 && dot(c-b, d-b) > 0)
      return false:
    return true;
  if (cross(d-a, b-a) * cross(c-a, b-a) > 0) return false;
  if (cross(a-c, d-c) * cross(b-c, d-c) > 0) return false;
  return true;
// compute intersection of line passing through a and b
// with line passing through c and d, assuming that unique
// intersection exists; for segment intersection, check if
// seaments intersect first
PT ComputeLineIntersection(PT a, PT b, PT c, PT d) {
 b=b-a; d=c-d; c=c-a;
  assert (dot (b, b) > EPS && dot (d, d) > EPS);
  return a + b*cross(c, d)/cross(b, d);
```

```
// compute center of circle given three points
PT ComputeCircleCenter(PT a, PT b, PT c) {
  b = (a+b)/2;
  c = (a + c) / 2;
  return ComputeLineIntersection(b, b+RotateCW90(a-b), c, c+RotateCW90(a-c)
// determine if point is in a possibly non-convex polygon (by William
// Randolph Franklin); returns 1 for strictly interior points, 0 for
// strictly exterior points, and 0 or 1 for the remaining points.
// Note that it is possible to convert this into an *exact* test using
// integer arithmetic by taking care of the division appropriately
// (making sure to deal with signs properly) and then by writing exact
// tests for checking point on polygon boundary
bool PointInPolygon(const vector<PT> &p, PT q) {
  bool c = 0;
  for (int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1)%p.size();
    if ((p[i].y \le q.y \&\& q.y < p[j].y ||
      p[j].y \le q.y \&\& q.y < p[i].y) \&\&
      q.x < p[i].x + (p[j].x - p[i].x) * (q.y - p[i].y) / (p[j].y - p[i].y)
      c = !c;
  return c;
// determine if point is on the boundary of a polygon
bool PointOnPolygon(const vector<PT> &p, PT q) {
  for (int i = 0; i < p.size(); i++)</pre>
    if (dist2(ProjectPointSegment(p[i], p[(i+1)*p.size()], q), q) < EPS)</pre>
      return true:
    return false;
// compute intersection of line through points a and b with
// circle centered at c with radius r > 0
vector<PT> CircleLineIntersection(PT a, PT b, PT c, double r) {
 vector<PT> ret;
 b = b-a;
 a = a-c;
  double A = dot(b, b);
  double B = dot(a, b);
  double C = dot(a, a) - r*r;
  double D = B*B - A*C;
  if (D < -EPS) return ret;</pre>
  ret.push_back(c+a+b*(-B+sqrt(D+EPS))/A);
  if (D > EPS)
    ret.push_back(c+a+b*(-B-sqrt(D))/A);
  return ret;
// compute intersection of circle centered at a with radius r
// with circle centered at b with radius R
vector<PT> CircleCircleIntersection(PT a, PT b, double r, double R) {
  vector<PT> ret:
  double d = sqrt(dist2(a, b));
  if (d > r+R \mid | d+min(r, R) < max(r, R)) return ret;
  double x = (d*d-R*R+r*r)/(2*d);
  double y = sqrt(r*r-x*x);
 PT v = (b-a)/d;
 ret.push_back(a+v*x + RotateCCW90(v)*y);
 if (\mathbf{v} > 0)
   ret.push_back(a+v*x - RotateCCW90(v)*y);
  return ret:
```

```
// This code computes the area or centroid of a (possibly nonconvex)
// polygon, assuming that the coordinates are listed in a clockwise or
// counterclockwise fashion. Note that the centroid is often known as
// the "center of gravity" or "center of mass".
double ComputeSignedArea(const vector<PT> &p) {
  double area = 0;
  for(int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) % p.size();
    area += p[i].x*p[j].y - p[j].x*p[i].y;
  return area / 2.0;
double ComputeArea(const vector<PT> &p) {
  return fabs(ComputeSignedArea(p));
PT ComputeCentroid(const vector<PT> &p) {
  PT c(0,0);
  double scale = 6.0 * ComputeSignedArea(p);
  for (int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) % p.size();
    c = c + (p[i]+p[j])*(p[i].x*p[j].y - p[j].x*p[i].y);
  return c / scale;
// tests whether or not a given polygon (in CW or CCW order) is simple
bool IsSimple(const vector<PT> &p) {
  for (int i = 0; i < p.size(); i++) {</pre>
    for (int k = i+1; k < p.size(); k++) {</pre>
      int j = (i+1) % p.size();
      int 1 = (k+1) % p.size();
      if (i == 1 \mid | j == k) continue;
      if (SegmentsIntersect(p[i], p[j], p[k], p[l]))
        return false;
  return true;
int main() {
  // expected: (-5,2)
  cerr << RotateCCW90(PT(2,5)) << endl;</pre>
  // expected: (5,-2)
  cerr << RotateCW90(PT(2,5)) << endl;</pre>
  // expected: (-5,2)
  cerr << RotateCCW(PT(2,5),M_PI/2) << endl;</pre>
  // expected: (5,2)
  cerr << ProjectPointLine(PT(-5,-2), PT(10,4), PT(3,7)) << endl;
  // expected: (5,2) (7.5,3) (2.5,1)
  cerr << ProjectPointSegment(PT(-5,-2), PT(10,4), PT(3,7)) << ""
       << ProjectPointSegment(PT(7.5,3), PT(10,4), PT(3,7)) << " "
       << ProjectPointSegment(PT(-5,-2), PT(2.5,1), PT(3,7)) << endl;</pre>
  // expected: 6.78903
  cerr << DistancePointPlane(4,-4,3,2,-2,5,-8) << endl;</pre>
  // expected: 1 0 1
  cerr << LinesParallel(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << " "
       << LinesParallel(PT(1,1), PT(3,5), PT(2,0), PT(4,5)) << " "
       << LinesParallel(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;
  // expected: 0 0 1
```

```
cerr << LinesCollinear(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << " "</pre>
     << LinesCollinear(PT(1,1), PT(3,5), PT(2,0), PT(4,5)) << " "
     << LinesCollinear(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;
// expected: 1 1 1 0
cerr << SegmentsIntersect(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) << " "
     << SegmentsIntersect(PT(0,0), PT(2,4), PT(4,3), PT(0,5)) << " "
     << SegmentsIntersect(PT(0,0), PT(2,4), PT(2,-1), PT(-2,1)) << ""
     << SegmentsIntersect(PT(0,0), PT(2,4), PT(5,5), PT(1,7)) << endl;
// expected: (1,2)
cerr << ComputeLineIntersection(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) <<
    endl;
// expected: (1,1)
cerr << ComputeCircleCenter(PT(-3,4), PT(6,1), PT(4,5)) << endl;</pre>
vector<PT> v;
v.push_back(PT(0,0));
v.push_back(PT(5,0));
v.push_back(PT(5,5));
v.push_back(PT(0,5));
// expected: 1 1 1 0 0
cerr << PointInPolygon(v, PT(2,2)) << " "</pre>
     << PointInPolygon(v, PT(2,0)) << " "
     << PointInPolygon(v, PT(0,2)) << " "</pre>
     << PointInPolygon(v, PT(5,2)) << " "
     << PointInPolygon(v, PT(2,5)) << endl;</pre>
// expected: 0 1 1 1 1
cerr << PointOnPolygon(v, PT(2,2)) << " "</pre>
     << PointOnPolygon(v, PT(2,0)) << " "</pre>
     << PointOnPolygon(v, PT(0,2)) << " "</pre>
     << PointOnPolygon(v, PT(5,2)) << " "
     << PointOnPolygon(v, PT(2,5)) << endl;
// expected: (1,6)
             (5,4) (4,5)
             blank line
             (4,5) (5,4)
             blank line
             (4,5) (5,4)
vector<PT> u = CircleLineIntersection(PT(0,6), PT(2,6), PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
u = CircleLineIntersection(PT(0,9), PT(9,0), PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
u = CircleCircleIntersection(PT(1,1), PT(10,10), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
u = CircleCircleIntersection(PT(1,1), PT(8,8), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 10, sqrt(2.0)/2.0);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 5, sqrt(2.0)/2.0);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
// area should be 5.0
// centroid should be (1.1666666, 1.166666)
PT pa[] = { PT(0,0), PT(5,0), PT(1,1), PT(0,5) };
vector<PT> p(pa, pa+4);
PT c = ComputeCentroid(p);
cerr << "Area: " << ComputeArea(p) << endl;</pre>
cerr << "Centroid: " << c << endl;</pre>
return 0;
```

3 Numerical algorithms

3.1 Number theory (modular, Chinese remainder, linear Diophantine)

```
// This is a collection of useful code for solving problems that
// involve modular linear equations. Note that all of the
// algorithms described here work on nonnegative integers.
#include <iostream>
#include <vector>
#include <algorithm>
using namespace std;
typedef vector<int> VI;
typedef pair<int, int> PII;
// return a % b (positive value)
int mod(int a, int b) {
        return ((a%b) + b) % b;
// computes gcd(a,b)
int gcd(int a, int b) {
        while (b) { int t = a%b; a = b; b = t; }
        return a:
// computes lcm(a,b)
int lcm(int a, int b) {
        return a / gcd(a, b) *b;
// (a^b) mod m via successive squaring
int powermod(int a, int b, int m)
        int ret = 1;
        while (b)
                if (b & 1) ret = mod(ret*a, m);
                a = mod(a*a, m);
                b >>= 1;
        return ret;
// returns g = gcd(a, b); finds x, y such that d = ax + by
int extended_euclid(int a, int b, int &x, int &y) {
        int xx = y = 0;
        int yy = x = 1;
        while (b) {
                int q = a / b;
                int t = b; b = a%b; a = t;
                t = xx; xx = x - q*xx; x = t;
                t = yy; yy = y - q*yy; y = t;
        return a;
// finds all solutions to ax = b \pmod{n}
VI modular_linear_equation_solver(int a, int b, int n) {
        int x, y;
        VI ret;
```

```
int g = extended_euclid(a, n, x, y);
        if (!(b%q)) {
                x = mod(x*(b / q), n);
                for (int i = 0; i < g; i++)
                        ret.push_back(mod(x + i*(n / q), n));
        return ret;
// computes b such that ab = 1 \pmod{n}, returns -1 on failure
int mod_inverse(int a, int n) {
        int x, y;
        int g = extended_euclid(a, n, x, y);
        if (q > 1) return -1;
        return mod(x, n);
// Chinese remainder theorem (special case): find z such that
// z % m1 = r1, z % m2 = r2. Here, z is unique modulo M = lcm(m1, m2).
// Return (z, M). On failure, M = -1.
PII chinese_remainder_theorem(int m1, int r1, int m2, int r2) {
        int s, t;
        int g = extended euclid(m1, m2, s, t);
        if (r1%g != r2%g) return make_pair(0, -1);
        return make_pair(mod(s*r2*m1 + t*r1*m2, m1*m2) / g, m1*m2 / g);
// Chinese remainder theorem: find z such that
// z % m[i] = r[i] for all i. Note that the solution is
// unique modulo M = lcm_i (m[i]). Return (z, M). On
// failure, M = -1. Note that we do not require the a[i]'s
// to be relatively prime.
PII chinese_remainder_theorem(const VI &m, const VI &r) {
        PII ret = make pair(r[0], m[0]);
        for (int i = 1; i < m.size(); i++) {</pre>
                ret = chinese_remainder_theorem(ret.second, ret.first, m[i
                    ], r[i]);
                if (ret.second == -1) break;
        return ret;
// computes x and y such that ax + by = c
// returns whether the solution exists
bool linear diophantine (int a, int b, int c, int &x, int &y) {
        if (!a && !b)
                if (c) return false;
                x = 0; y = 0;
                return true;
        if (!a)
                if (c % b) return false;
                x = 0; v = c / b;
                return true;
        if (!b)
                if (c % a) return false;
                x = c / a; y = 0;
               return true;
        int g = gcd(a, b);
        if (c % q) return false;
        x = c / q * mod inverse(a / q, b / q);
        y = (c - a*x) / b;
        return true;
```

```
int main() {
        // expected: 2
        cout << gcd(14, 30) << endl;
        // expected: 2 -2 1
        int x, y;
        int g = extended_euclid(14, 30, x, y);
        cout << g << " " << x << " " << y << endl;
        // expected: 95 451
        VI sols = modular_linear_equation_solver(14, 30, 100);
        for (int i = 0; i < sols.size(); i++) cout << sols[i] << " ";</pre>
        // expected: 8
        cout << mod_inverse(8, 9) << endl;</pre>
        // expected: 23 105
                   11 12
        PII ret = chinese_remainder_theorem(VI({ 3, 5, 7 }), VI({ 2, 3, 2
        cout << ret.first << " " << ret.second << endl;</pre>
        ret = chinese_remainder_theorem(VI({ 4, 6 }), VI({ 3, 5 }));
        cout << ret.first << " " << ret.second << endl;</pre>
        // expected: 5 -15
        if (!linear_diophantine(7, 2, 5, x, y)) cout << "ERROR" << endl;</pre>
        cout << x << " " << y << endl;
        return 0;
```

3.2 Systems of linear equations, matrix inverse, determinant

```
// Gauss-Jordan elimination with full pivoting.
// (1) solving systems of linear equations (AX=B)
// (2) inverting matrices (AX=I)
    (3) computing determinants of square matrices
// Running time: O(n^3)
             a[][] = an nxn matrix
             b[][] = an nxm matrix
// OUTPUT:
                   = an nxm matrix (stored in b[][])
             A^{-1} = an nxn matrix (stored in a[][])
             returns determinant of a[][]
#include <iostream>
#include <vector>
#include <cmath>
using namespace std;
const double EPS = 1e-10;
typedef vector<int> VI;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
T GaussJordan (VVT &a, VVT &b) {
  const int n = a.size();
```

```
const int m = b[0].size();
 VI irow(n), icol(n), ipiv(n);
 T \det = 1;
  for (int i = 0; i < n; i++) {
    int pj = -1, pk = -1;
    for (int j = 0; j < n; j++) if (!ipiv[j])</pre>
     for (int k = 0; k < n; k++) if (!ipiv[k])</pre>
        if (pj == -1 || fabs(a[j][k]) > fabs(a[pj][pk])) { <math>pj = j; pk = k;
    if (fabs(a[pj][pk]) < EPS) { cerr << "Matrix is singular." << endl;</pre>
        exit(0); }
    ipiv[pk]++;
    swap(a[pj], a[pk]);
    swap(b[pj], b[pk]);
    if (pj != pk) det *= -1;
    irow[i] = pj;
    icol[i] = pk;
    T c = 1.0 / a[pk][pk];
    det *= a[pk][pk];
    a[pk][pk] = 1.0;
    for (int p = 0; p < n; p++) a[pk][p] *= c;</pre>
    for (int p = 0; p < m; p++) b[pk][p] *= c;</pre>
    for (int p = 0; p < n; p++) if (p != pk) {
     c = a[p][pk];
     a[p][pk] = 0;
     for (int q = 0; q < n; q++) a[p][q] -= a[pk][q] * c;
     for (int q = 0; q < m; q++) b[p][q] -= b[pk][q] * c;
  for (int p = n-1; p >= 0; p--) if (irow[p] != icol[p]) {
    for (int k = 0; k < n; k++) swap(a[k][irow[p]], a[k][icol[p]]);</pre>
  return det;
int main() {
  const int n = 4;
  const int m = 2;
  double A[n][n] = \{ \{1,2,3,4\}, \{1,0,1,0\}, \{5,3,2,4\}, \{6,1,4,6\} \};
  double B[n][m] = \{ \{1,2\}, \{4,3\}, \{5,6\}, \{8,7\} \};
 VVT a(n), b(n);
  for (int i = 0; i < n; i++) {
   a[i] = VT(A[i], A[i] + n);
   b[i] = VT(B[i], B[i] + m);
  double det = GaussJordan(a, b);
  // expected: 60
 cout << "Determinant: " << det << endl;</pre>
  // expected: -0.233333 0.166667 0.133333 0.0666667
               0.166667 0.166667 0.333333 -0.333333
               0.05 -0.75 -0.1 0.2
  cout << "Inverse: " << endl;</pre>
  for (int i = 0; i < n; i++) {
    for (int j = 0; j < n; j++)
      cout << a[i][j] << ' ';
    cout << endl;</pre>
  // expected: 1.63333 1.3
               -0.166667 0.5
               2.36667 1.7
```

```
// -1.85 -1.35
cout << "Solution: " << endl;
for (int i = 0; i < n; i++) {
   for (int j = 0; j < m; j++)
      cout << b[i][j] << ' ';
   cout << endl;
}</pre>
```

3.3 Reduced row echelon form, matrix rank

```
// Reduced row echelon form via Gauss-Jordan elimination
// with partial pivoting. This can be used for computing
// the rank of a matrix.
// Running time: O(n^3)
// INPUT:
             a[][] = an nxm matrix
// OUTPUT: rref[][] = an nxm matrix (stored in a[][])
             returns rank of a[][]
#include <iostream>
#include <vector>
#include <cmath>
using namespace std;
const double EPSILON = 1e-10;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
int rref(VVT &a) {
  int n = a.size();
  int m = a[0].size();
  int r = 0;
  for (int c = 0; c < m && r < n; c++) {
    int j = r;
    for (int i = r + 1; i < n; i++)
     if (fabs(a[i][c]) > fabs(a[j][c])) j = i;
    if (fabs(a[j][c]) < EPSILON) continue;</pre>
    swap(a[j], a[r]);
    T s = 1.0 / a[r][c];
    for (int j = 0; j < m; j++) a[r][j] *= s;
    for (int i = 0; i < n; i++) if (i != r) {
     T t = a[i][c];
     for (int j = 0; j < m; j++) a[i][j] -= t * a[r][j];
    r++;
  return r;
int main() {
  const int n = 5, m = 4;
  double A[n][m] = {
   {16, 2, 3, 13},
    { 5, 11, 10, 8},
   { 9, 7, 6, 12},
    { 4, 14, 15, 1},
    {13, 21, 21, 13}};
  VVT a(n);
  for (int i = 0; i < n; i++)
    a[i] = VT(A[i], A[i] + m);
```

```
int rank = rref(a);

// expected: 3
cout << "Rank: " << rank << endl;

// expected: 1 0 0 1
// 0 1 0 3
// 0 0 1 -3
// 0 0 0 3.10862e-15
cout << "rref: " << endl;

for (int i = 0; i < 5; i++) {
   for (int j = 0; j < 4; j++)
      cout << endl;
}

cout << endl;
}</pre>
```

4 Graph algorithms

4.1 Eulerian path

```
struct Edge;
typedef list<Edge>::iterator iter;
struct Edge
        int next vertex;
        iter reverse_edge;
        Edge(int next_vertex)
                :next_vertex(next_vertex)
                { }
};
const int max vertices = ;
int num vertices;
list<Edge> adj[max_vertices];
                                       // adjacency list
vector<int> path;
void find_path(int v)
        while(adj[v].size() > 0)
                int vn = adj[v].front().next_vertex;
                adj[vn].erase(adj[v].front().reverse_edge);
                adj[v].pop_front();
                find_path(vn);
        path.push_back(v);
void add_edge(int a, int b)
        adj[a].push_front(Edge(b));
        iter ita = adj[a].begin();
        adj[b].push_front(Edge(a));
        iter itb = adj[b].begin();
        ita->reverse_edge = itb;
        itb->reverse_edge = ita;
```

5 Data structures

5.1 Binary Indexed Tree

```
#include <iostream>
using namespace std;
#define LOGSZ 17
int tree[(1<<LOGSZ)+1];</pre>
int N = (1 << LOGSZ);
// add v to value at x
void set(int x, int v) {
 while (x \le N) {
   tree[x] += v;
    x += (x \& -x);
// get cumulative sum up to and including x
int get(int x) {
 int res = 0;
 while(x) {
   res += tree[x];
    x = (x \& -x);
 return res;
// get largest value with cumulative sum less than or equal to x;
// for smallest, pass x-1 and add 1 to result
int getind(int x) {
 int idx = 0, mask = N;
  while(mask && idx < N) {</pre>
    int t = idx + mask;
    if(x >= tree[t]) {
     idx = t;
     x -= tree[t];
    mask >>= 1;
  return idx;
```

6 Miscellaneous

6.1 Dates

```
// Routines for performing computations on dates. In these routines,
// months are expressed as integers from 1 to 12, days are expressed
// as integers from 1 to 31, and years are expressed as 4-digit
// integers.

#include <iostream>
#include <string>
using namespace std;
string dayOfWeek[] = {"Mon", "Tue", "Wed", "Thu", "Fri", "Sat", "Sun"};
```

```
// converts Gregorian date to integer (Julian day number)
int dateToInt (int m, int d, int v) {
  return
    1461 * (y + 4800 + (m - 14) / 12) / 4 +
    367 * (m - 2 - (m - 14) / 12 * 12) / 12 -
    3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +
    d - 32075;
// converts integer (Julian day number) to Gregorian date: month/day/year
void intToDate (int jd, int &m, int &d, int &y) {
  int x, n, i, j;
 x = jd + 68569;
 n = 4 * x / 146097;
  x = (146097 * n + 3) / 4;
  i = (4000 * (x + 1)) / 1461001;
  x = 1461 * i / 4 - 31;
  \dot{j} = 80 * x / 2447;
 d = x - 2447 * j / 80;
 x = \frac{1}{7} / 11;
 m = 1 + 2 - 12 * x;
 v = 100 * (n - 49) + i + x;
// converts integer (Julian day number) to day of week
string intToDay (int jd) {
  return dayOfWeek[jd % 7];
int main (int argc, char **argv) {
  int jd = dateToInt (3, 24, 2004);
  int m, d, y;
  intToDate (jd, m, d, y);
  string day = intToDay (jd);
  // expected output:
       2453089
        3/24/2004
       Wed
  cout << jd << endl
    << m << "/" << d << "/" << v << endl
    << day << endl;
```

6.2 Prime numbers

```
// O(sqrt(x)) Exhaustive Primality Test
#include <cmath>
#define EPS 1e-7
typedef long long LL;
bool IsPrimeSlow (LL x)
  if(x<=1) return false;</pre>
  if(x<=3) return true;</pre>
  if (!(x%2) || !(x%3)) return false;
  LL s=(LL) (sqrt((double)(x))+EPS);
  for (LL i=5; i<=s; i+=6)</pre>
    if (!(x%i) || !(x%(i+2))) return false;
  return true;
// Primes less than 1000:
            3
                   5
             43
                   47
                                 59
                          53
```

```
107
                           109
                                      127
                                                 137
                                                            149
                                                            211
                                                                  223
     157
           163
                167
                      173
                           179
                                 181
                                      191
                                            193
                                                 197
                                                       199
     227
           229
                                            263
                                                 269
                                                       271
                                                            277
                                                                  281
                233
                      239
                           241
                                 251
                                      257
     283
           293
                307
                      311
                           313
                                 317
                                      331
                                            337
                                                 347
                                                       349
                                                            353
                                                                  359
          373
     367
                379
                      383
                           389
                                 397
                                      401
                                            409
                                                 419
                                                      421
                                                            431
                                                                  433
     439
           443
                449
                     457
                           461
                                 463
                                      467
                                            479
                                                 487
                                                      491
                                                            499
                                                                  503
     509
           521
                523
                     541
                           547
                                 557
                                      563
                                            569
                                                 571
                                                      577
                                                            587
                                                                  593
           601
                607
                     613
                           617
                                 619
                                      631
                                            641
                                                 643
                                                      647
                                                            653
                                                                  659
     661
           673
                677
                      683
                           691
                                      709
                                            719
                                                 727
                                                       733
                                                            739
                                                                  743
     751
           757
                761
                      769
                           773
                                 787
                                      797
                                                 811
                                                       821
                                                            823
                                                                  827
           839
                853
                     857
                                 863
                                                            907
     919
           929
                937
                      941
                           947
                                 953
                                                      983
                                                            991
                                                                  997
// Other primes:
     The largest prime smaller than 10 is 7.
     The largest prime smaller than 100 is 97.
     The largest prime smaller than 1000 is 997.
     The largest prime smaller than 10000 is 9973.
     The largest prime smaller than 100000 is 99991.
     The largest prime smaller than 1000000 is 999983.
     The largest prime smaller than 10000000 is 9999991.
     The largest prime smaller than 100000000 is 99999989.
     The largest prime smaller than 1000000000 is 999999937.
     The largest prime smaller than 10000000000 is 9999999967.
     The largest prime smaller than 10000000000 is 9999999977.
     The largest prime smaller than 100000000000 is 999999999999.
     The largest prime smaller than 1000000000000 is 999999999971.
     The largest prime smaller than 1000000000000 is 9999999999973.
     The largest prime smaller than 100000000000000 is 99999999999989.
     The largest prime smaller than 100000000000000 is 99999999999937.
```

6.3 C++ input/output

```
#include <iostream>
#include <iomanip>
using namespace std;
int main()
    // Ouput a specific number of digits past the decimal point,
    // in this case 5
    cout.setf(ios::fixed); cout << setprecision(5);</pre>
    cout << 100.0/7.0 << endl;
    cout.unsetf(ios::fixed);
    // Output the decimal point and trailing zeros
    cout.setf(ios::showpoint);
    cout << 100.0 << endl;
    cout.unsetf(ios::showpoint);
    // Output a '+' before positive values
    cout.setf(ios::showpos);
    cout << 100 << " " << -100 << endl;
    cout.unsetf(ios::showpos);
    // Output numerical values in hexadecimal
    cout << hex << 100 << " " << 1000 << " " << 10000 << dec << endl;
```

6.4 Knuth-Morris-Pratt

```
Finds all occurrences of the pattern string p within the
text string t. Running time is O(n + m), where n and m
are the lengths of p and t, respecitvely.
#include <iostream>
#include <string>
#include <vector>
using namespace std;
typedef vector<int> VI;
void buildPi(string& p, VI& pi)
  pi = VI(p.length());
  int k = -2;
  for(int i = 0; i < p.length(); i++) {</pre>
    while (k >= -1 \&\& p[k+1] != p[i])
     k = (k == -1) ? -2 : pi[k];
    pi[i] = ++k;
int KMP (string& t, string& p)
  VI pi;
 buildPi(p, pi);
  int k = -1;
  for(int i = 0; i < t.length(); i++) {</pre>
    while (k >= -1 \&\& p[k+1] != t[i])
      k = (k == -1) ? -2 : pi[k];
    if(k == p.length() - 1) {
      // p matches t[i-m+1, ..., i]
      cout << "matched at index " << i-k << ": ";</pre>
      cout << t.substr(i-k, p.length()) << endl;</pre>
      k = (k == -1) ? -2 : pi[k];
  return 0;
int main()
  string a = "AABAACAADAABAABA", b = "AABA";
  KMP(a, b); // expected matches at: 0, 9, 12
  return 0;
```

6.5 Latitude/longitude

```
/*
Converts from rectangular coordinates to latitude/longitude and vice
versa. Uses degrees (not radians).
*/
#include <iostream>
#include <cmath>
using namespace std;
```

```
struct 11
  double r, lat, lon;
};
struct rect
 double x, y, z;
11 convert (rect& P)
 11 Q;
 Q.r = sqrt(P.x*P.x+P.y*P.y+P.z*P.z);
 Q.lat = 180/M_PI*asin(P.z/Q.r);
 Q.lon = 180/M_PI*acos(P.x/sqrt(P.x*P.x+P.y*P.y));
 return 0;
rect convert (11& Q)
 rect P;
 P.x = Q.r*cos(Q.lon*M_PI/180)*cos(Q.lat*M_PI/180);
 P.y = Q.r*sin(Q.lon*M_PI/180)*cos(Q.lat*M_PI/180);
 P.z = Q.r*sin(Q.lat*M_PI/180);
 return P;
int main()
 rect A;
 11 B;
 A.x = -1.0; A.y = 2.0; A.z = -3.0;
 B = convert(A);
 cout << B.r << " " << B.lat << " " << B.lon << endl;
 A = convert(B);
 cout << A.x << " " << A.y << " " << A.z << endl;
```

6.6 Dates (Java)

```
// Example of using Java's built-in date calculation routines
import java.text.SimpleDateFormat;
import java.util.*;
public class Dates {
    public static void main(String[] args) {
        Scanner s = new Scanner(System.in);
        SimpleDateFormat sdf = new SimpleDateFormat("M/d/yyyy");
        while (true) {
            int n = s.nextInt();
            if (n == 0) break;
            GregorianCalendar c = new GregorianCalendar(n, Calendar.JANUARY
            while (c.get(Calendar.DAY_OF_WEEK) != Calendar.SATURDAY)
                c.add(Calendar.DAY_OF_YEAR, 1);
            for (int i = 0; i < 12; i++) {
                System.out.println(sdf.format(c.getTime()));
                while (c.get(Calendar.MONTH) == i) c.add(Calendar.
                    DAY_OF_YEAR, 7);
```

}

7 Hussain

7.1 Adaptive Simpson

```
// Adaptive Simpson's Rule (Wikipedia Article)
dbl adaptiveSimpsons(dbl (*f)(dbl), // ptr to function
                           dbl a, dbl b, // interval [a,b]
                           dbl epsilon, // error tolerance
                           int maxRecursionDepth) { // recursion cap
  db1 c = (a + b)/2, h = b - a;
  dbl fa = f(a), fb = f(b), fc = f(c);
  dbl S = (h/6) * (fa + 4*fc + fb);
  return adaptiveSimpsonsAux(f, a, b, epsilon, S, fa, fb, fc,
      maxRecursionDepth);
// Recursive auxiliary function for adaptiveSimpsons() function below
dbl adaptiveSimpsonsAux(dbl (*f)(dbl), dbl a, dbl b, dbl epsilon,
                         dbl S, dbl fa, dbl fb, dbl fc, int bottom) {
  dbl c = (a + b)/2, h = b - a;
  db1 d = (a + c)/2, e = (c + b)/2;
  dbl fd = f(d), fe = f(e);
  dbl Sleft = (h/12)*(fa + 4*fd + fc);
  dbl Sright = (h/12)*(fc + 4*fe + fb);
  dbl S2 = Sleft + Sright;
  if (bottom <= 0 || fabs(S2 - S) <= 15*epsilon) // magic 15 comes from</pre>
      error analysis
    return S2 + (S2 - S)/15;
  return adaptiveSimpsonsAux(f, a, c, epsilon/2, Sleft, fa, fc, fd, bottom
      -1) +
         adaptiveSimpsonsAux(f, c, b, epsilon/2, Sright, fc, fb, fe, bottom
             -1);
int main(){
float I = adaptiveSimpsons(sin, 0, 2, 0.001, 100);
                                                           // compute
     integral of sin(x)
                                                          // from 0 to 2
                                                              and store it
                                                          // the new
                                                               variable I
 printf("I = %lf\n",I); // print the result
 return 0;
```

7.2 Binomial Coeff (constant N)

```
C[0] = 1

for (int k = 0; k < n; ++ k)

C[k+1] = (C[k] * (n-k)) / (k+1)

// C[i] = C(n,i)
```

7.3 Generate (x,y) pairs s.t. x AND y=y

for(int y = x; y; y = (y-1)&x)
 cout<<y<<endl;</pre>

for(**int** x = 1; x <= n; x++)