Assignment #2

Name:	ID:	
This	assignment has ${\bf 4}$ questions, for a total of ${\bf 25}$ marks.	
_	nyms: SOS (structural operational semantics), COS (BG (big step), CBV (call by value), CBN (call by name	
	Γ ion	
• $t_1 = f(3+5) : \mathbb{N}$		[3]

•
$$t_2 = f((\lambda x : \mathbb{N}. x + 2)5) : \mathbb{N}$$

[3]

Question 2: Encoding	nave as
• sequencing: $t := \cdots \mid t; t'$. Semantics: t is evaluated first, then t' is evaluated.	[1]
• let-in: $t := \cdots \mid let \ x = t \ in \ t'$. Semantics: t is evaluated into a value v and then t' is evaluate v in place of x .	ted for [1]
• arrays of length 4: $t ::= \cdots \mid [t,t,t,t]$. Values include arrays of values: $v ::= \cdots \mid [v,v,v,v]$. (no semantics for this case)	[1]
• array field access: $t := \cdots \mid t.i \ (i \in 03)$. Semantics: for $i \in 03$ we have that $[v_0, v_1, v_1, v_2]$ returns v_i (show the encodings for at least two cases of i).	[2]
• array update: $t ::= \cdots \mid t.i = t (i \in 03)$. Semantics: for $i \in 03$ we have that $[v_0, v_1, v_2, v_3]$ returns $[v_0, v_1, v, v_3]$ (show the encodings for at least two cases of i).].2 = v $[2]$

Consider extending STLC with named function, call this language STLCN. A program is no longer a term, but a collection of named functions $(P := \varnothing \mid P; F)$. A named function defines a function name, a parameter of a certain type and the function body $(F := f(x : \tau) : \tau \mapsto t)$. A function body is a term. Terms now must include new constructs to call other functions.

Define the COS judgements for STLCN. Define the primitive reduction rules for STLCN as well as evaluation contexts. Define the typing judgements for STLCN, starting from how to determine when a program is well typed. Define the typing rules for STLCN.

Some primitive reductions and typing rules from STLCN will be similar to those of STLC. To avoid duplicates, name all the STLC rules that have analogous ones and show how to convert two of them to the new forms only.

STLCN should still be safe, i.e., it should not get stuck trying to call a function with a parameter of the wrong type, or calling a function that does not exist. However, STLCN is not normalising: there may be functions that mutually call each other and thus diverge.

Define the linking two programs (P_1+P_2) , which yields the program obtained by joining the two. Linking is only possible when both programs are well-typed and when they fulfill each other's exports with their own imports.