

Robust Hyperproperty Preservation for Secure Compilation

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FOR SOFTWARE SYSTEMS





Special Thanks to:



Contents

Robust Compilation Criteria

Proof Techniques

Where is FAC?

Background & Motivation

- many criteria imply secure compilation
preserving memory safety, CFI, non interference, program
equivalence

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equivalence

Is that all?

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Goal: study criteria that

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properties formally

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Goal: study criteria that

- are **security-driven** and preserve security
properties **formally**
- are **robust** (hold for all adversarial context)

Background & Motivation

- many criteria imply secure compilation
preserving memory safety, CFI, non interference, program
equivalence

Relate **backtranslation** techniques
and property preservation

Goal

- are **security-driven** and preserve security
properties **formally**
- are **robust** (hold for all adversarial context)

HyperProperties

Clarkson & Schneider '08

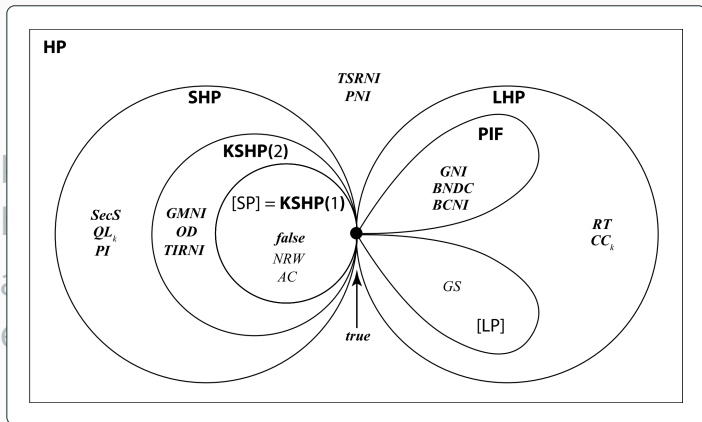
- properties = sets of traces

HyperProperties Clarkson & Schneider '08

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- are organised in subclasses for expressiveness

HyperProperties Clarkson & Schneider '08



Robust Compilation Criteria

Robust Compilation Partial Order

In the partial order:

- higher notions are **stronger**

Robust Compilation Partial Order

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 - and **trickier** to achieve

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Robust Compilation Partial Order

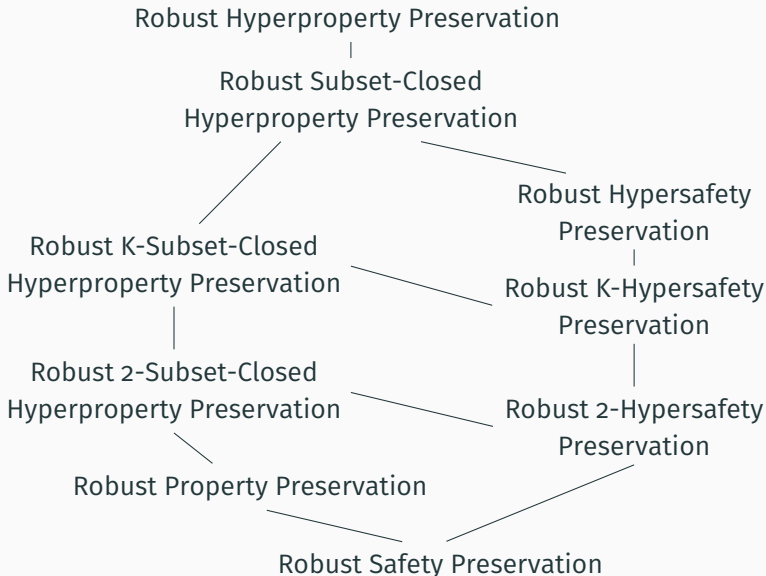
In the partial order:

- higher notions are **stronger**
 - and **trickier** to achieve
- each notion comes in **two flavours**
 - one with **clear HP correspondence**
 - one for **simpler proofs**

Notation

- P_s, P_t : components of S and T
- C_s, C_t : contexts
- $C_s[P_s], C_t[P_t]$: whole programs
- $[\![\cdot]\!] : P_s \rightarrow P_t$: compiler from S to T
- β : traces (possibly infinite), I/O with an environment
- $\text{Behav}(P_s)$: set of traces of P_s
- π : prefix (finite)
- $<$: prefixing

Robust Compilation Criteria



RPP: Robust Property Preservation

Definition (RPP)

$\llbracket \cdot \rrbracket \in \text{RPP} \stackrel{\text{def}}{=} \forall P_s, P.$

if $(\forall C_s. \text{Behav}(C_s[P_s]) \subseteq P)$

then $(\forall C_t. \text{Behav}(C_t[\llbracket P_s \rrbracket]) \subseteq P)$

RC: Robust Compilation

Definition: (RC)

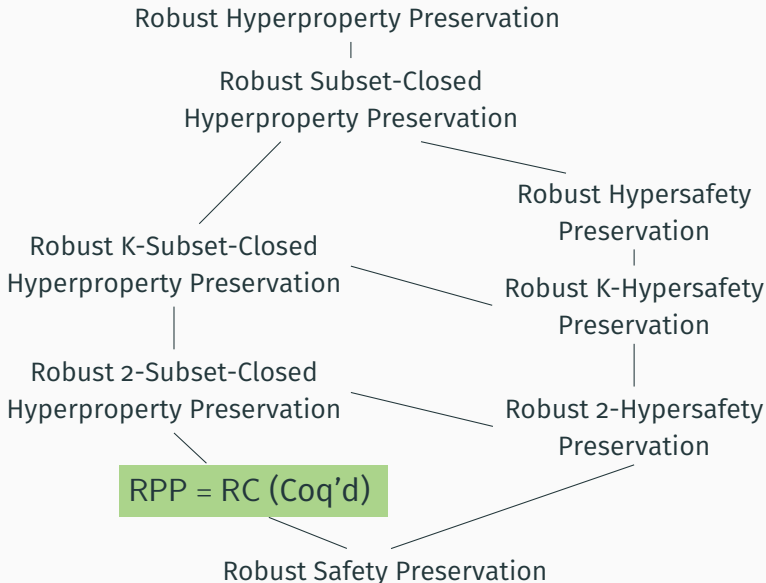
$$\begin{aligned} \llbracket \cdot \rrbracket \in \text{RC} &\stackrel{\text{def}}{=} \forall \mathbb{C}_t, P_s, \beta. \exists \mathbb{C}_s. \\ &\quad \text{if } \beta \in \text{Behav}(\mathbb{C}_t[\llbracket P_s \rrbracket]) \\ &\quad \text{then } \beta \in \text{Behav}(\mathbb{C}_s[P_s]) \end{aligned}$$

RC: Robust Compilation

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Robust Compilation Criteria



Robust Safety Property Preservation

Definition (RSPP)

$$[[\cdot]] \in \text{RSPP} \stackrel{\text{def}}{=} \forall P_s, P \in SP.$$

$$\text{if } (\forall C_s. \text{Behav}(C_s[P_s]) \subseteq P)$$

$$\text{then } (\forall C_t. \text{Behav}(C_t[[P_s]]) \subseteq P)$$

Robust Safety Compilation

Definition: (RC)

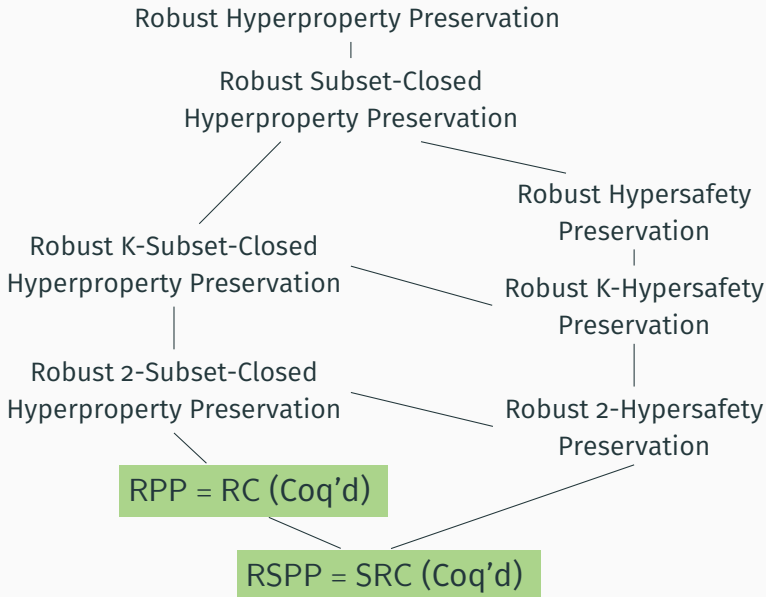
$$\begin{aligned} \llbracket \cdot \rrbracket \in \text{RC} &\stackrel{\text{def}}{=} \forall \mathbb{C}_t, \mathbb{P}_s, \beta. \exists \mathbb{C}_s. \\ &\quad \text{if } \beta \in \text{Behav}(\mathbb{C}_t[\llbracket \mathbb{P}_s \rrbracket]) \\ &\quad \text{then } \beta \in \text{Behav}(\mathbb{C}_s[\mathbb{P}_s]) \end{aligned}$$

Robust Safety Compilation

Definition: (SRC)

$$\begin{aligned} \llbracket \cdot \rrbracket \in \text{SRC} &\stackrel{\text{def}}{=} \forall \mathbb{C}_t, P_s, \pi. \exists \mathbb{C}_s. \\ &\quad \text{if } \pi < \text{Behav}(\mathbb{C}_t[\llbracket P_s \rrbracket]) \\ &\quad \text{then } \pi < \text{Behav}(\mathbb{C}_s[P_s]) \end{aligned}$$

Robust Compilation Criteria



RSHP: Robust Hypersafety Preservation

Definition (RPP)

$$[[\cdot]] \in \text{RPP} \stackrel{\text{def}}{=} \forall P_s, P.$$

$$\text{if } (\forall C_s. \text{Behav}(C_s[P_s]) \subseteq P)$$

$$\text{then } (\forall C_t. \text{Behav}(C_t[[P_s]]) \subseteq P)$$

RSHP: Robust Hypersafety Preservation

Definition (RSHP)

$$\begin{aligned} \llbracket \cdot \rrbracket \in \text{RSHP} &\stackrel{\text{def}}{=} \forall P_s, H \in \text{SHP}. \\ &\quad \text{if } (\forall C_s. \text{Behav}(C_s[P_s]) \in H) \\ &\quad \text{then } (\forall C_t. \text{Behav}(C_t[\llbracket P_s \rrbracket]) \in H) \end{aligned}$$

SHRC: Hypersafety Robust Compilation

Definition: (RC)

$$\begin{aligned} \llbracket \cdot \rrbracket \in \text{RC} &\stackrel{\text{def}}{=} \forall P_s, C_t, \pi. \exists C_s. \\ &\text{if } \pi < \text{Behav}(C_t[\llbracket P_s \rrbracket]) \\ &\text{then } \pi < \text{Behav}(C_s[P_s]) \end{aligned}$$

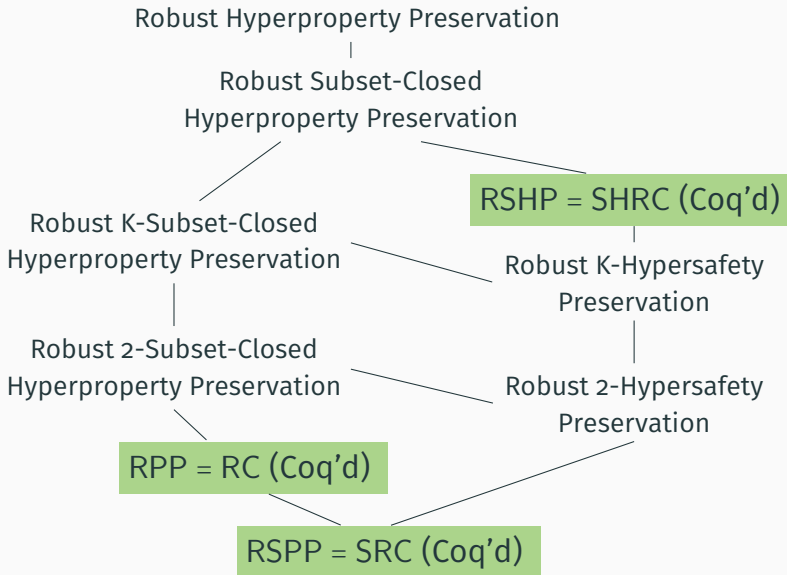
SHRC: Hypersafety Robust Compilation

Definition: (SHRC)

$$\begin{aligned} \llbracket \cdot \rrbracket \in \text{SHRC} &\stackrel{\text{def}}{=} \forall P_s, C_t, \hat{\pi}. \exists C_s. \\ &\text{if } \hat{\pi} \preceq \text{Behav}(C_t[\llbracket P_s \rrbracket]) \\ &\text{then } \hat{\pi} \preceq \text{Behav}(C_s[P_s]) \end{aligned}$$

$\hat{\pi}$: finite set of prefixes

Robust Compilation Criteria



Robust Hyperproperty Preservation

Definition (RHP)

$\llbracket \cdot \rrbracket \in \text{RHP} \stackrel{\text{def}}{=} \forall P_s, H.$

if $(\forall C_s. \text{Behav}(C_s[P_s]) \in H)$

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Hyperproperty Robust Compilation

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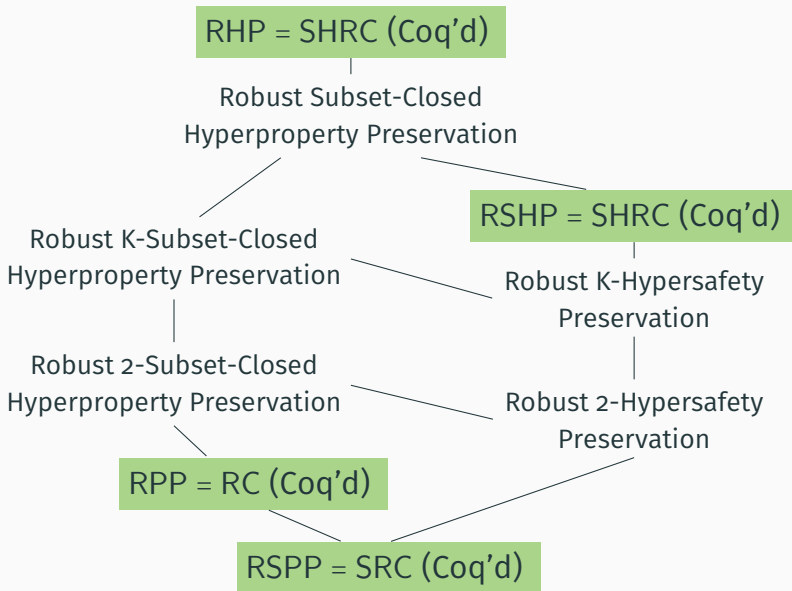
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Hyperproperty Robust Compilation

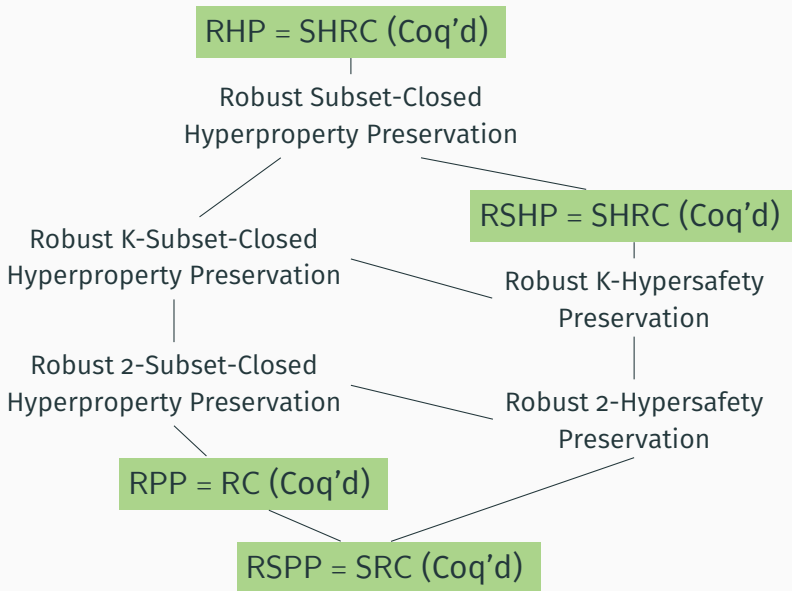
Definition: (HRC)

$$\begin{aligned} \llbracket \cdot \rrbracket \in \text{HRC} &\stackrel{\text{def}}{=} \forall \mathbb{C}_t, P_s, \exists \mathbb{C}_s. \forall \beta. \\ &\quad \beta \in \text{Behav}(\mathbb{C}_t[\llbracket P_s \rrbracket]) \\ &\quad \iff \beta \in \text{Behav}(\mathbb{C}_s[P_s]) \end{aligned}$$

Robust Compilation Criteria

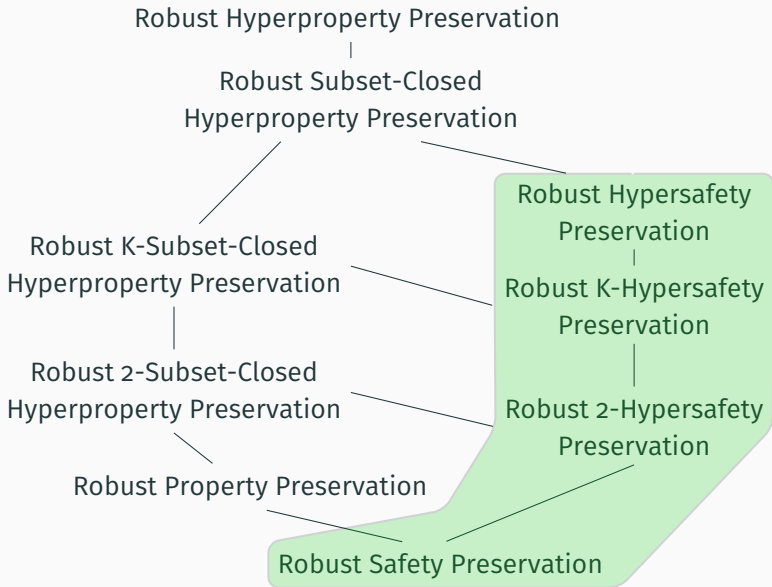


Robust Compilation Criteria

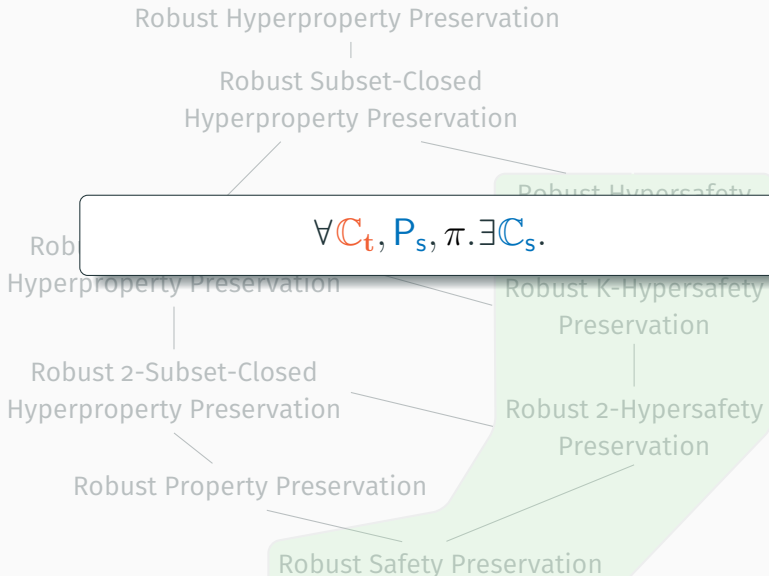


Proof Techniques

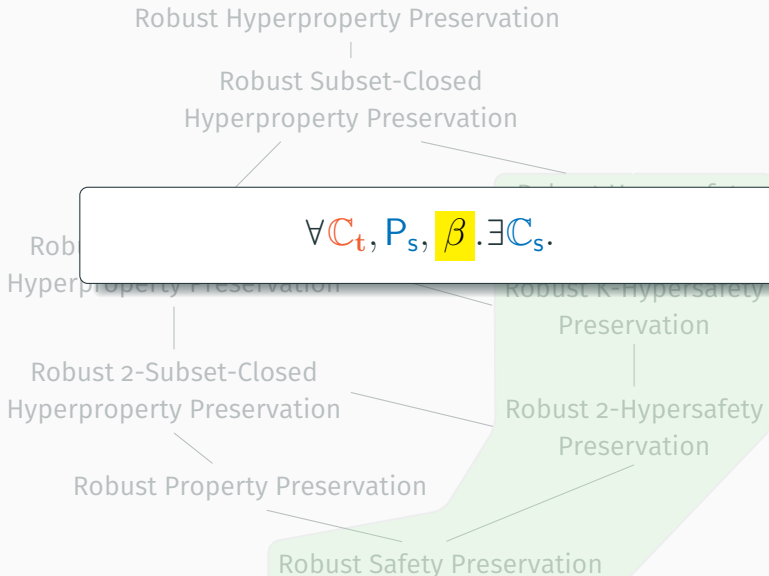
Proof Techniques



Proof Techniques

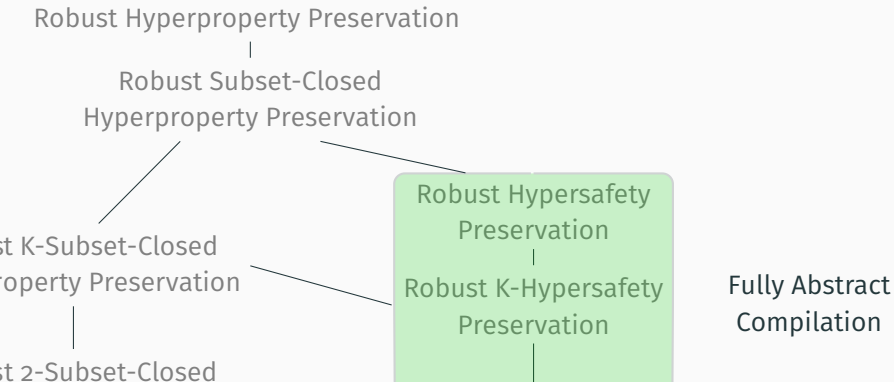


Proof Techniques

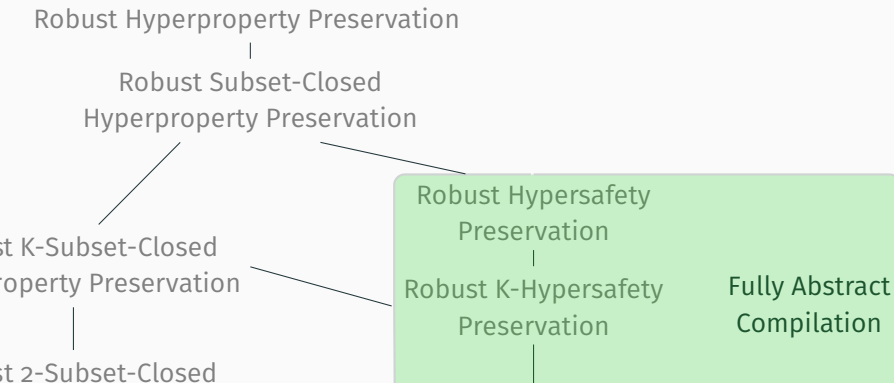


Where is FAC?

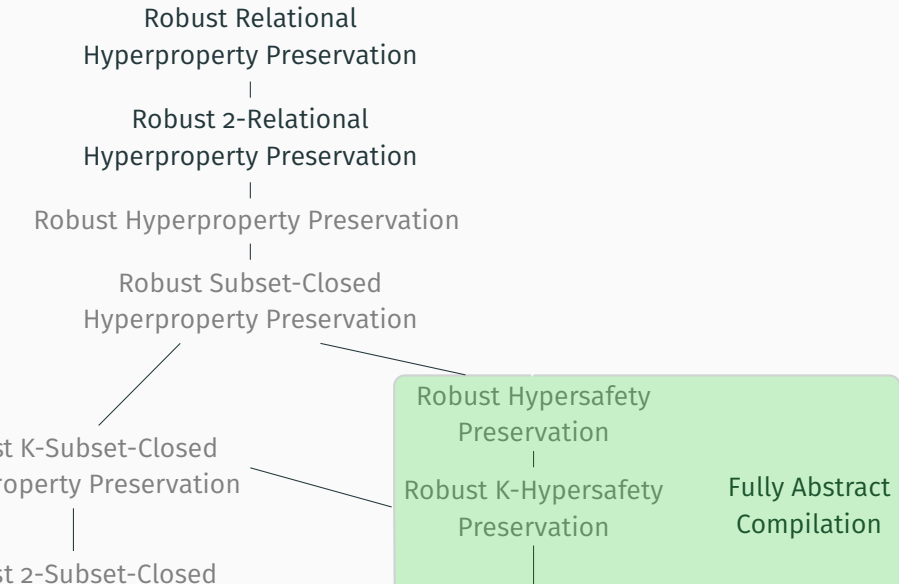
There is FAC



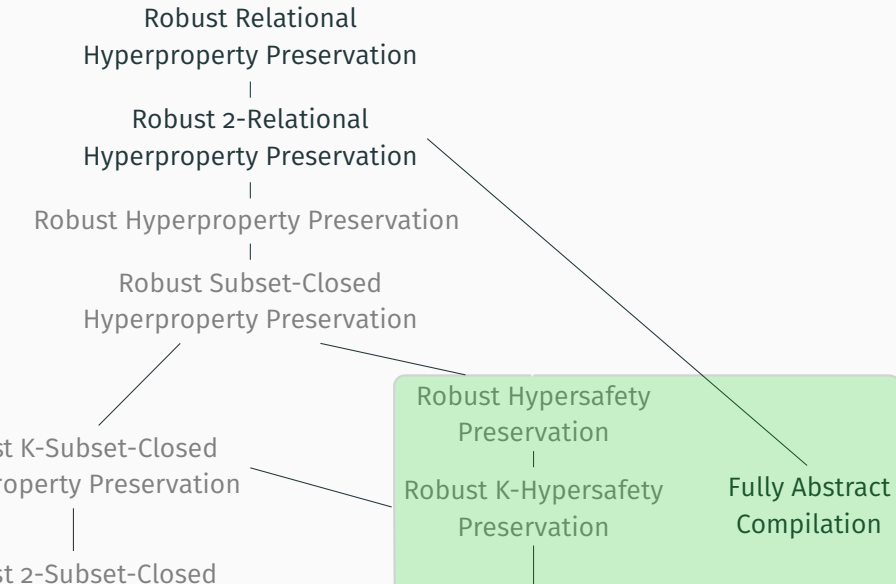
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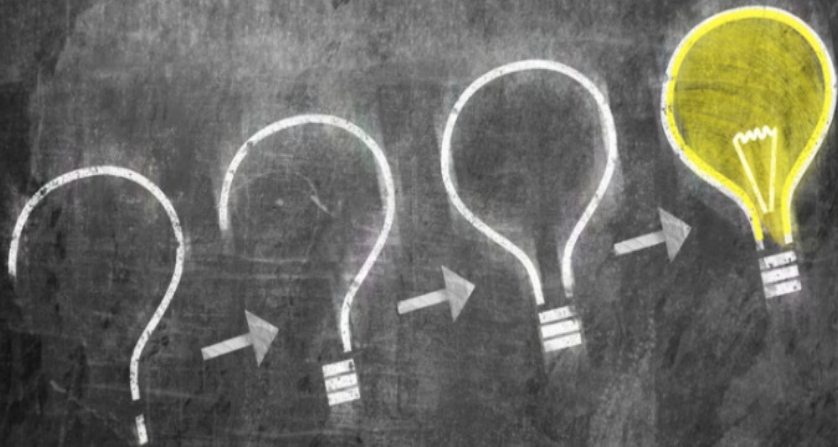
Conclusion

- motivated the Robust Compilation Partial Order
- discussed some of these criteria
- analysed proof techniques for some criteria

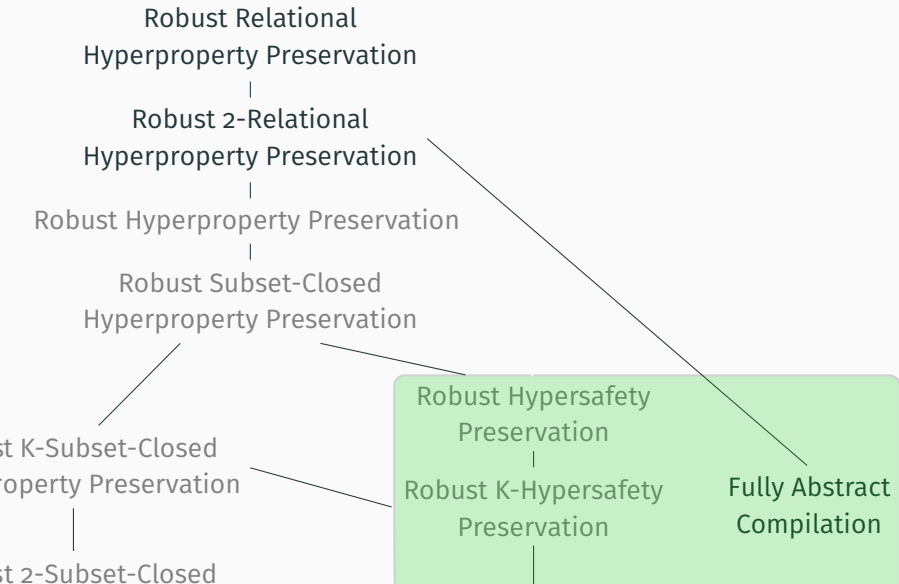
Conclusion

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Conclusion



Robust Compilation Criteria



Robust Relational Hyperproperty Preservation

Definition: (HRC)

$$\begin{aligned} \llbracket \cdot \rrbracket \in \text{HRC} &\stackrel{\text{def}}{=} \forall \mathbb{C}_t, \mathbb{P}_s. \exists \mathbb{C}_s. \beta \\ &\quad \beta \in \text{Behav}(\mathbb{C}_t[\llbracket \mathbb{P}_s \rrbracket]) \\ &\iff \beta \in \text{Behav}(\mathbb{C}_s[\mathbb{P}_s]) \end{aligned}$$

Robust Relational Hyperproperty Preservation

Definition: (RRHP)

$$\begin{aligned} \llbracket \cdot \rrbracket \in \text{RRHP} &\stackrel{\text{def}}{=} \forall \mathbb{C}_t, \exists \mathbb{C}_s. \forall P_s, \beta \\ &\quad \beta \in \text{Behav}(\mathbb{C}_t[\llbracket P_s \rrbracket]) \\ &\iff \beta \in \text{Behav}(\mathbb{C}_s[P_s]) \end{aligned}$$

Robust 2-Relational Hyperproperty Preservation

Definition: (RRHP)

$$\llbracket \cdot \rrbracket \in \text{RRHP} \stackrel{\text{def}}{=} \forall \mathbb{C}_t, \exists \mathbb{C}_s. \forall \mathbb{P}_s. \beta \left(\begin{array}{l} \beta \in \text{Behav}(\mathbb{C}_t[\llbracket \mathbb{P}_s \rrbracket]) \\ \iff \beta \in \text{Behav}(\mathbb{C}_s[\mathbb{P}_s]) \end{array} \right)$$

Robust 2-Relational Hyperproperty Preservation

Definition: (R2RHP)

$$\llbracket \cdot \rrbracket \in \text{R2RHP} \stackrel{\text{def}}{=} \forall \mathbb{C}_t, \exists \mathbb{C}_s. \forall P_s, P'_s. \beta$$
$$\left(\begin{array}{l} \beta \in \text{Behav}(\mathbb{C}_t[\llbracket P_s \rrbracket]) \\ \iff \beta \in \text{Behav}(\mathbb{C}_s[P_s]) \end{array} \right)$$

and

$$\left(\begin{array}{l} \beta \in \text{Behav}(\mathbb{C}_t[\llbracket P'_s \rrbracket]) \\ \iff \beta \in \text{Behav}(\mathbb{C}_s[P'_s]) \end{array} \right)$$