Practice Exercises

8.1 Suppose that we decompose the schema r(A, B, C, D, E) into

$$r_1(A, B, C)$$

 $r_2(A, D, E)$

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Show that this decomposition is a lossless decomposition if the following set *F* of functional dependencies holds:

$$A \to BC$$

$$CD \to E$$

$$B \to D$$

$$E \to A$$

- **8.3** Explain how functional dependencies can be used to indicate the following:
 - A one-to-one relationship set exists between entity sets *student* and *instructor*.
 - A many-to-one relationship set exists between entity sets *student* and *instructor*.
- 8.4 Use Armstrong's axioms to prove the soundness of the union rule. (*Hint*: Use the augmentation rule to show that, if $\alpha \to \beta$, then $\alpha \to \alpha\beta$. Apply the augmentation rule again, using $\alpha \to \gamma$, and then apply the transitivity rule.)
- 8.5 Use Armstrong's axioms to prove the soundness of the pseudotransitivity 5 rule.
- **8.6** Compute the closure of the following set F of functional dependencies for relation schema r (A, B, C, D, E).

$$A \rightarrow BC$$

$$CD \rightarrow E$$

$$B \rightarrow D$$

$$E \rightarrow A$$

List the candidate keys for *R*.

8.7 Using the functional dependencies of Practice Exercise 8.6, compute the canonical cover F_c .

| A | В | С |
|-------|-------|-------|
| a_1 | b_1 | c_1 |
| a_1 | b_1 | c_2 |
| a_2 | b_1 | c_1 |
| a_2 | b_1 | C3 |

Figure 8.17 Relation of Practice Exercise 8.2.

- **8.9** Given the database schema R(a, b, c), and a relation r on the schema R, write an SQL query to test whether the functional dependency $b \rightarrow c$ holds on relation r. Also write an SQL assertion that enforces the functional dependency; assume that no null values are present. (Although part of the SQL standard, such assertions are not supported by any database imple-
- mentation currently.)

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- Our discussion of lossless-join decomposition implicitly assumed that attributes on the left-hand side of a functional dependency cannot take on null values. What could go wrong on decomposition, if this property is violated?
- In the BCNF decomposition algorithm, suppose you use a functional dependency $\alpha \to \beta$ to decompose a relation schema $r(\alpha, \beta, \gamma)$ into $r_1(\alpha, \beta)$ and $r_2(\alpha, \gamma)$.
 - What primary and foreign-key constraint do you expect to hold on the decomposed relations?
 - Give an example of an inconsistency that can arise due to an erroneous update, if the foreign-key constraint were not enforced on the decomposed relations above.
 - When a relation is decomposed into 3NF using the algorithm in Section 8.5.2, what primary and foreign key dependencies would you expect will hold on the decomposed schema?
- Let R_1, R_2, \ldots, R_n be a decomposition of schema U. Let u(U) be a relation, and let $r_i = \Pi_{R_i}(u)$. Show that

$$u \subseteq r_1 \bowtie r_2 \bowtie \cdots \bowtie r_n$$

Show that it is possible to ensure that a dependency-preserving decomposition into 3NF is a lossless decomposition by guaranteeing that at least one schema contains a candidate key for the schema being decomposed. (Hint: Show that the join of all the projections onto the schemas of the decomposition cannot have more tuples than the original relation.)

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- **8.16** Let a **prime** attribute be one that appears in at least one candidate key. Let α and β be sets of attributes such that $\alpha \to \beta$ holds, but $\beta \to \alpha$ does not hold. Let A be an attribute that is not in α , is not in β , and for which $\beta \to A$ holds. We say that A is **transitively dependent** on α . We can restate our definition of 3NF as follows: A relation schema R is in 3NF with respect to a set F of functional dependencies if there are no nonprime attributes A in R for which A is transitively dependent on a key for R. Show that this new definition is equivalent to the original one.
- **8.17** A functional dependency $\alpha \to \beta$ is called a **partial dependency** if there is a proper subset γ of α such that $\gamma \to \beta$. We say that β is *partially dependent* on α . A relation schema R is in **second normal form** (2NF) if each attribute A in R meets one of the following criteria:
 - It appears in a candidate key.
 - It is not partially dependent on a candidate key.

Show that every 3NF schema is in 2NF. (*Hint*: Show that every partial dependency is a transitive dependency.)

8.18 Give an example of a relation schema *R* and a set of dependencies such that *R* is in BCNF but is not in 4NF. 10

Exercises

8.21 Normalize the following schema, with given constraints, to 4NF.

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books(accessionno, isbn, title, author, publisher) users(userid, name, deptid, deptname) accessionno \rightarrow isbn isbn \rightarrow title isbn \rightarrow publisher isbn \rightarrow author userid \rightarrow name userid \rightarrow deptid deptid \rightarrow deptname

8.22 Explain what is meant by *repetition of information* and *inability to represent information*. Explain why each of these properties may indicate a bad relational database design.

- 8.23 Why are certain functional dependencies called *trivial* functional dependencies? 10
- Use the definition of functional dependency to argue that each of Armstrong's axioms (reflexivity, augmentation, and transitivity) is sound.
- 8.25 Consider the following proposed rule for functional dependencies: If $\alpha \to \beta$ and $\gamma \to \beta$, then $\alpha \to \gamma$. Prove that this rule is *not* sound by showing a relation *r* that satisfies $\alpha \to \beta$ and $\gamma \to \beta$, but does not satisfy $\alpha \to \gamma$.
- 8.26 Use Armstrong's axioms to prove the soundness of the decomposition rule. 5

8.29 Consider the following set F of functional dependencies on the relation schema r(A, B, C, D, E, F):

$$\begin{array}{l} A \rightarrow BCD \\ BC \rightarrow DE \\ B \rightarrow D \\ D \rightarrow A \end{array}$$

- a. Compute B^+ .
- b. Prove (using Armstrong's axioms) that *AF* is a superkey.
- c. Compute a canonical cover for the above set of functional dependencies F; give each step of your derivation with an explanation.
- d. Give a 3NF decomposition of r based on the canonical cover.
- e. Give a BCNF decomposition of r using the original set of functional dependencies.
- f. Can you get the same BCNF decomposition of *r* as above, using the canonical cover?
- **8.30** List the three design goals for relational databases, and explain why each is desirable. **5**

- **8.31** In designing a relational database, why might we choose a non-BCNF design? **5**
- **8.33** Given a relational schema r(A, B, C, D), does $A \rightarrow BC$ logically imply $A \rightarrow B$ and $A \rightarrow C$? If yes prove it, else give a counter example.

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8.34 Explain why 4NF is a normal form more desirable than BCNF. 10

QUESTIONS FROM STUDY MATERIALS

Module 3

| Define redundancy? | 2 |
|--|---|
| Define functional dependency? | 2 |
| Discuss normalization? | 2 |
| Illustrate functional dependency with example? | 2 |
| Illustrate fully functional dependency with example? | 2 |
| Demonstrate transitive dependency? Give an example? | 2 |
| Discuss Domain-Key Normal Form? | 2 |
| Define Armstrong axioms for FD's? | 2 |
| Define First Normal Form? | 2 |
| Define second Normal Form? | 2 |
| Define third Normal Form? | 2 |
| Define Fourth Normal Form? | 2 |
| List out the Problems related to decompositions? | 2 |
| Explain about Loss less-join dependency? | 2 |
| Explain about BCNF? | 2 |

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| Explain about multi-valued dependencies? | 2 |
| Define join dependency and fifth normal form? | 2 |
| Explain the concept scheme refinement in database design? | 2 |
| Define dependency preserving decomposition? | 2 |
| Explain about inclusion dependency? | 2 |
| Define schema? | 2 |
| Define the terms Entity type | 2 |
| Define strong entity sets? | 2 |
| Explain about stored attributes? | 2 |
| Define a Transaction? | 2 |
| Define normalization? Explain 1NF, 2NF, 3NF Normal forms? | 5 |
| Compare and contrast BCNF with 3NF? | 5 |
| Describe properties of decompositions? | 5 |
| Explain about Schema refinement in Database design? | 5 |
| Illustrate Multivalued dependencies and Fourth normal form with example? | 5 |

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| Discuss about Join dependencies and Fifth normal form? | 5 |
| Illustrate Inclusion dependencies with example? | 5 |
| Explain Atomicity | 5 |
| Discuss How do you implement Durability? | 5 |
| Discuss How do you implement Atomicity ? | 5 |
| Discuss Serializability in detail? | 5 |
| Explain Durability | 5 |
| Explain consistancy | 5 |
| Explain Isolation | 5 |
| Discuss in detail Multiple Granularity? | 5 |
| Given a relation R(A, B, C, D) and Functional Dependency set FD = { AB \rightarrow CD, B \rightarrow C}, determine whether the given R is in 2NF? If not convert it into 2 NF. | 10 |
| Given a relation R(X, Y, Z, W, P) and Functional Dependency set FD = { $X \rightarrow Y, Y \rightarrow P$, and $Z \rightarrow W$ }, determine whether the given R is in 3NF? If not convert it into 3 NF. | 10 |
| Given a relational schema R(X, Y, Z) set of functional dependencies P and Q such that: $P = \{X \to Y, Y \to Z, Z \to X\} \text{ and } Q = \{X \to YZ, Y \to X, Z \to X\} \text{ using FD sets P and Q proof that P=Q}$ | 10 |
| Discuss about Lossless Decomposition with example | 10 |
| Discuss about Dependency Perverse with example | 10 |
| Discuss Multivalued dependency with example | 10 |

| Discuss Join dependency with example | 10 |
|--|----|
| Discuss Inclusion dependence with example | 10 |
| Given a relational Schema R(V, W, X, Y, Z) and set of Function Dependency FD = { $V \rightarrow W$, $VW \rightarrow X$, $Y \rightarrow VXZ$ }. Find the canonical cover? | 10 |
| Discuss Reflexsive, augmentation and transitive rule | 10 |
| Discuss 4NF and 5NF normalization with example | 10 |
| Discuss 2NF and 3NF normalization with example | 10 |
| Discuss the different type of normal form and their description | 10 |
| Discuss the advantage and disadvantage of normaliztion | 10 |
| Discuss Insertion anomaly | 2 |
| Show that: if $\alpha \to \beta$ and $\alpha \to \gamma$ then $\alpha \to \beta \gamma$ | 10 |
| Consider the following relational schemes for a library database: Book (Title, Author, Catalog_no, Publisher, Year, Price) Collection (Title, Author, Catalog_no) the following are functional dependencies: a. Title Author> Catalog_no b. Catalog_no> Title Author Publisher Year c. Publisher Title Year> Price d. Assume {Author, Title} is the key for both schemes. Apply the appropriate normal form for Book and Cancellation? | 10 |
| Consider a schema R (A, B, C, D) and functional dependencies A -> B and C -> D. Solve and find whether the decomposition of R into R1 (A, B) and R2(C, D) belongs to which one or both (dependency preserving and loss less join)? | 10 |