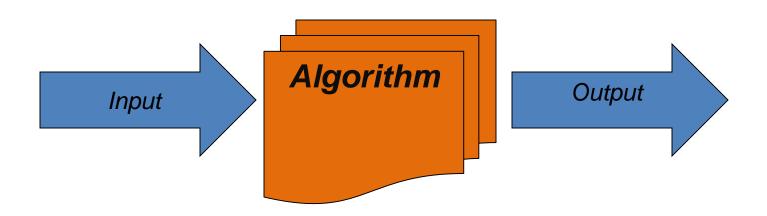
Algorithms Design and Analysis

What is an algorithm?

- A computational procedure that takes some value, or set of values, as *input* and produces some value, or set of values, as *output*.
- A sequence of computational steps that transform the input into the output.



What is an algorithm?

- A computational problem is a mathematical problem, specified by an input/output relation.
- An algorithm is a computational procedure for solving a computational problem.
- Example: Sorting
 - Input: A sequence of N numbers a₁...a_n
 - **Output**: the permutation (reordering) of the input sequence such that $a_1 \le a_2 \le ... \le a_n$

What will we study?

- Expressing algorithms
 - Define a problem precisely and abstractly
 - Presenting algorithms using pseudocode
- Algorithm validation
 - Prove that an algorithm is correct
- Algorithm analysis
 - Time and space complexity
 - What problems are so hard that efficient algorithms are unlikely to exist
- Designing algorithms
 - Algorithms for classical problems
 - Meta algorithms (classes of algorithms) and when you should use which

Pseudocode

- High-level description of an algorithm
- More structured than English prose
- Less detailed than a program
- Preferred notation for describing algorithms
- Hides program design issues

Example: find max element of an array

Algorithm arrayMax(A, n)Input array A of n integers
Output maximum element of A

 $\begin{array}{l} \textit{currentMax} \leftarrow A[0] \\ \textit{for } \textit{i} \leftarrow 1 \textit{ to } \textit{n} - 1 \textit{ do} \\ \textit{if } A[\textit{i}] > \textit{currentMax} \textit{ then} \\ \textit{currentMax} \leftarrow A[\textit{i}] \\ \textit{return } \textit{currentMax} \end{array}$

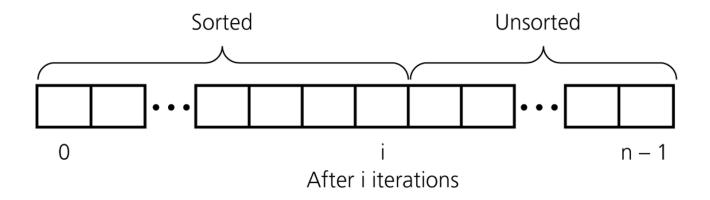
Algorithm Design

- Learn general approaches to algorithm design
 - Divide and conquer
 - Greedy method
 - Dynamic Programming
 - Basic Search and Traversal Technique
 - Graph Theory
 - Linear Programming
 - Approximation Algorithm
 - NP Problem

Sorting Algorithms

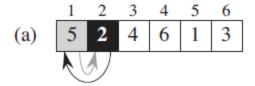
Insertion Sort

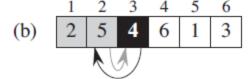
- Divide the list into two portions: sorted and unsorted.
- while some elements unsorted:
 - Using linear search, find the location in the sorted portion where the 1st element of the unsorted portion should be inserted
 - Move all the elements after the insertion location up one position to make space for the new element

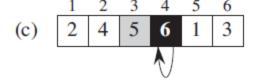


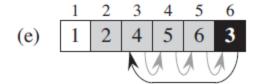
Insertion Sort

```
for j = 2 to A. length
key = A[j]
# Insert A[j] into the sorted sequence A[1...j-1].
i = j-1
while i > 0 and A[i] > key
A[i+1] = A[i]
i = i-1
A[i+1] = key
```









Analysis of Insertion Sort

```
INSERTION-SORT (A)
                                                     times
                                             cost
   for j = 2 to A. length
                                             c_1
                                                     n
                                             c_2 \qquad n-1
2 	 key = A[j]
   // Insert A[j] into the sorted
           sequence A[1..j-1].
                                                     n-1
                                                    n-1
4 	 i = j - 1
                                             C_{\Delta}
                                             c_5 \qquad \sum_{i=2}^n t_i
5 while i > 0 and A[i] > key
                                             c_6 \qquad \sum_{i=2}^{n} (t_i - 1)
          A[i + 1] = A[i]
6
                                             c_7 \qquad \sum_{i=2}^{n} (t_i - 1)
   i = i - 1
      A[i+1] = key
                                                     n-1
                                             C_8
```

The running time T(n) =

$$c_1 n + c_2 (n-1) + c_4 (n-1) + c_5 \sum_{j=2}^{n} t_j + c_6 \sum_{j=2}^{n} (t_j - 1) + c_7 \sum_{j=2}^{n} (t_j - 1) + c_8 (n-1)$$

Analysis of Insertion Sort

• For the best case: $t_j = 1$ for j = 2,3,....,n $T(n) = c_1 n + c_2 (n-1) + c_4 (n-1) + c_5 (n-1) + c_8 (n-1)$ $= (c_1 + c_2 + c_4 + c_5 + c_8)n - (c_2 + c_4 + c_5 + c_8)$

it is thus a *linear function* of n.

• For the worst case: $t_j = j$ for j = 2,3,...,n

$$\sum_{j=2}^{n} j = \frac{n(n+1)}{2} - 1$$
 and

$$\sum_{j=2}^{n} (j-1) = \frac{n(n-1)}{2}$$

Analysis of Insertion Sort

$$T(n) = c_1 n + c_2 (n-1) + c_4 (n-1) + c_5 \left(\frac{n(n+1)}{2} - 1\right)$$

$$+ c_6 \left(\frac{n(n-1)}{2}\right) + c_7 \left(\frac{n(n-1)}{2}\right) + c_8 (n-1)$$

$$= \left(\frac{c_5}{2} + \frac{c_6}{2} + \frac{c_7}{2}\right) n^2 + \left(c_1 + c_2 + c_4 + \frac{c_5}{2} - \frac{c_6}{2} - \frac{c_7}{2} + c_8\right) n$$

$$- \left(c_2 + c_4 + c_5 + c_8\right).$$

It is thus a *quadratic function* of n