Polyalphabetic Ciphers

- □ A polyalphabetic cipher is any cipher based on substitution, using multiple substitution alphabets.
- polyalphabetic cipher techniques have the following features in common:
 - A set of related monoalphabetic substitution rules is used.
 - □ A key determines which particular rule is chosen for a given transformation.

Polyalphabetic Ciphers Encryption

☐ Assume

a	b	c	d	e	f	g	h	i	j	k	1	m
0	1	2	3	4	5	6	7	8	9	10	11	12

n	O	p	q	r	S	t	u	V	W	X	y	Z
13	14	15	16	17	18	19	20	21	22	23	24	25

☐ Then We set these Rules:

Polyalphabetic Ciphers Encryption

- ☐ Then We set Key as Rules:
- 1) Shift the first letter three position to the right
- 2) Shift the second letter five position to the right
- 3) Shift the third letter seven position to the right

☐ Given Plaintext = security

Polyalphabetic Ciphers Encryption

☐ Given Plaintext = security

1) Divide Plaintext to three words

a	b	c	d	e	f	g	h	i	j	k	1	m
0	1	2	3	4	5	6	7	8	9	10	11	12

n	O	p	q	r	S	t	u	V	W	X	y	Z
13	14	15	16	17	18	19	20	21	22	23	24	25

- 2) P= sec uri ty
- 3) C= VJJ XWP WD

Polyalphabetic Ciphers Decryption

- ☐ Then We set Key as Rules: (reverse)
- 1) Shift the first letter three position to the left
- 2) Shift the second letter five position to the left
- 3) Shift the third letter seven position to the left

☐ Given Ciphertext = VJJXWPWD

Polyalphabetic Ciphers Decryption

☐ Given C= VJJXWPWD

1) Divide Plaintext to three words as your rules number

a	b	С	d	e	f	g	h	i	j	k	1	m
0	1	2	3	4	5	6	7	8	9	10	11	12

n	O	p	q	r	S	t	u	V	W	X	y	Z
13	14	15	16	17	18	19	20	21	22	23	24	25

- 2) C= VJJ XWP WD
- 3) C= SEC URI TY

Vigenère Cipher

☐ The Vigenère cipher, was invented by a Frenchman, Blaise de Vigenère in the 16th century.

□ Vigenère cipher is a simple polyalphabetic cipher

Vigenère Cipher

$$\square C_i = (P_i + K) \mod 26$$

$$\square P_i = (C_i - K) \mod 26$$

□ Repeating key

 \square K=deceptive

 \square P=we are discovered save yourself

key:

deceptivedeceptive

plaintext:

wearediscoveredsaveyourself

a	a	b	С	d	e	f	g	h	i	j	k	1	m
()	1	2	3	4	5	6	7	8	9	10	11	12

n	O	p	q	r	S	t	u	V	W	X	y	Z
13	14	15	16	17	18	19	20	21	22	23	24	25

key:

deceptivedeceptivedeceptive

plaintext:

wearediscoveredsaveyourself

key	3	4	2	4	15	19	8	21	4	3	4	2	4	15
plaintext	22	4	0	17	4	3	8	18	2	14	21	4	17	4
ciphertext	25	8	2	21	19	22	16	13	6	17	25	6	21	19

key	19	8	21	4	3	4	2	4	15	19	8	21	4
plaintext	3	18	0	21	4	24	14	20	17	18	4	11	5
ciphertext	22	0	21	25	7	2	16	24	6	11	12	6	9

☐ Result

key:

plaintext:

ciphertext:

deceptivedeceptive

wearediscoveredsaveyourself

ZICVTWQNGRZGVTWAVZHCQYGLMGJ

The strength of Vigenère Cipher is that there are multiple ciphertext letters for each plaintext letter

- decryption simply works in reverse
- $\square P_i = (C_i K) \mod 26$

ciphertext:

key:

plaintext:

ZICVTWQNGRZGVTWAVZHCQYGLMGJ

deceptivedeceptivedeceptive

wearediscoveredsaveyourself

Autokey Cipher

An autokey cipher (also known as the autoclave cipher) is a cipher which incorporates the message (the plaintext) into the key.

$$\square P = \{p_1, p_2, p_3, ..., p_n\}$$

$$\square K = \{k_1, p_1, p_2, p_3, ..., p_{n-1}\}$$

$$\Box C = \{c_1, c_2, c_3, ..., c_n\}$$

Autokey Cipher

$$\square C_i = (P_i + K_i) mod 26$$

$$\square P_i = (C_i - K_i) \mod 26$$

a	b	С	d	e	f	g	h	i	j	k	1	m
0	1	2	3	4	5	6	7	8	9	10	11	12

n	О	p	q	r	S	t	u	V	W	X	y	Z
13	14	15	16	17	18	19	20	21	22	23	24	25

Autokey Cipher Encryption

 \square K=m

 \square P=attack is today

1													
Plaintext	a	t	t	a	C	k	i	S	t	o	d	a	у
P Value	0	19	19	0	2	10	8	18	19	14	3	0	24
Key	12	0	19	19	0	2	10	8	18	19	14	3	0
C Value	12	19	12	19	2	12	18	0	11	7	17	3	24
Ciphertext	m	t	m	t	C	m	S	а	ı	h	r	d	у

Classical Encryption Techniques 2

Autokey Cipher Decryption

 \square K=m

 \Box C=mtmtcmsalhrdy

m	t	m	t	С	m	S	а	I	h	r	d	у
12	19	12	19	2	12	18	0	11	7	17	3	24
12	0	19	19	0	2	10	8	18	19	14	3	0
0	19	19	0	2	10	8	18	19	14	3	0	24
а	t	t	а	С	k	i	S	t	o	d	а	У
	12 12 0	12 19 12 0 0 19	12 19 12 12 0 19 0 19 19	12 19 12 19 12 0 19 19 0 19 19 0	12 19 12 19 2 12 0 19 19 0 0 19 19 0 2	12 19 12 19 2 12 12 0 19 19 0 2 0 19 19 0 2 10	12 19 12 19 2 12 18 12 0 19 19 0 2 10 0 19 19 0 2 10 8	12 19 12 19 2 12 18 0 12 0 19 19 0 2 10 8 0 19 19 0 2 10 8 18	12 19 12 19 2 12 18 0 11 12 0 19 19 0 2 10 8 18 0 19 19 0 2 10 8 18 19	12 19 12 19 2 12 18 0 11 7 12 0 19 19 0 2 10 8 18 19 0 19 19 0 2 10 8 18 19 14	12 19 12 19 2 12 18 0 11 7 17 12 0 19 19 0 2 10 8 18 19 14 0 19 19 0 2 10 8 18 19 14 3	m t m t c m s a l h r d 12 19 12 19 2 12 18 0 11 7 17 3 12 0 19 19 0 2 10 8 18 19 14 3 0 19 19 0 2 10 8 18 19 14 3 0 a t t a c k i s t o d a

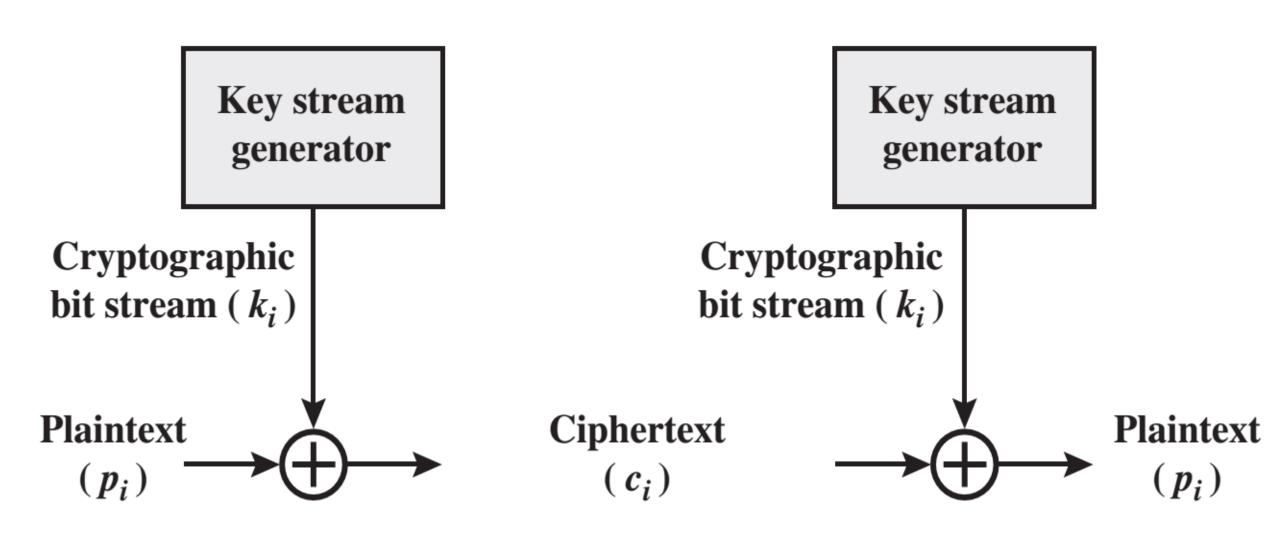
Dr Mohamed Loey

Vernam Cipher

☐ Vernam Cipher was introduced by an AT&T engineer named Gilbert Vernam in 1918.

☐ The ultimate defense against such a cryptanalysis is to choose a keyword that is as long as the plaintext and has no statistical relationship to it.

Vernam Cipher



Vernam Cipher

☐ Encryption

$$\succ C = P XOR K$$

☐ Decryption

$$\triangleright P = C XOR K$$

Vernam Cipher Encryption

- □ P=11100011101010101101
- □ K=1001010101

- □ P=11100011101010101101
- □ K=10010101011001010101
- □ C=011101101100111111000

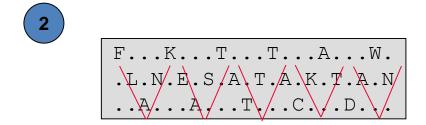
Vernam Cipher Decryption

- □ C=011101101100111111000
- □ K=10010101011001010101
- □ P=11100011101010101101

Transposition Ciphers



The clear text message would be encoded using a key of 3.



Use a rail fence cipher and a key of 3.



The clear text message would appear as follows.

Transposition Techniques

□ Transposition Techniques performing some sort of permutation on the plaintext letters (reorder the position of letters in plaintext).

- ☐ Types:
 - Rail Fence Cipher
 - Row Transposition Cipher

Rail Fence Cipher Encryption

- ☐ P= meet me after the toga party
- □ K=2

1)
$$p = \begin{bmatrix} m & e & ma & t & r & h & t & g & p & r & y \\ e & t & e & f & e & t & e & o & a & a & t \end{bmatrix}$$

2) C=mematrhtgpryetefeteoaat

Rail Fence Cipher Decryption

- ☐ C=mematrhtgpryetefeteoaat
- □ K=2

1)
$$C = \begin{bmatrix} m & e & ma & t & r & h & t & g & p & r & y \\ e & t & e & f & e & t & e & o & a & a & t \end{bmatrix}$$

2) P= meetmeafterthetogaparty

Row Transposition Cipher Encryption

☐ P= attack postponed until two am

$$\Box \ C = \begin{pmatrix} 4 & 3 & 1 & 2 & 5 & 6 & 7 \\ a & t & t & a & c & k & p \\ o & s & t & p & o & n & e \\ d & u & n & t & i & l & t \\ w & o & a & m & x & x & x \end{pmatrix} = ttnaaptmtsuoaodwcoixknlxpetx$$

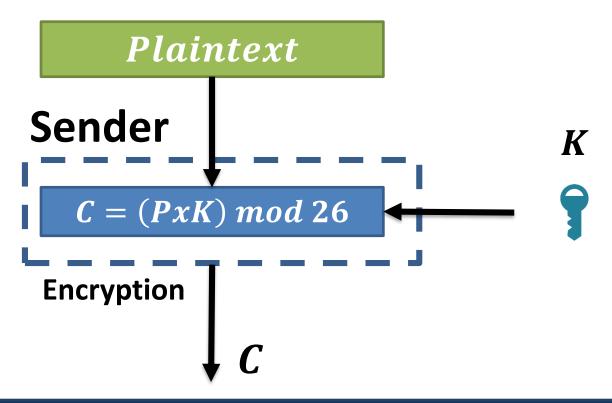
Row Transposition Cipher Decryption

- \Box C= ttnaaptmtsuoaodwcoixknlxpetx = Len(C)=28
- □ K= 4312567
- ☐ Each Column have 28/7= 4 letter

$$\Box P = \begin{pmatrix} 4 & 3 & 1 & 2 & 5 & 6 & 7 \\ a & t & t & a & c & k & p \\ o & s & t & p & o & n & e \\ d & u & n & t & i & l & t \\ w & o & a & m & x & x & x \end{pmatrix} = attackpostponeduntiltwoamxxx$$

Task 1

 \Box As shown in Figure below, use Multiplicative Cipher to encrypt "enemy attack tonight" with key = 4.



Task2

☐ As shown in Figure below, use Affine Cipher to encrypt "enemy attack tonight" with key pair (4,3).

