### Administrivia

HW2 solutions released: see HW2 repo

#### Thursday

- Midterm I
- HW3 release
- Project I Part 3 release

## Midterm I Review

## **Topics**

#### Data Models

- ER models
- Relational model
- Pros/Cons of hierarchical/network models
- ER → Relational Model translation

Relational Algebra

## **Topics**

#### DDL statements

- CREATE TABLE statements
- Integrity Constraints

#### SQL

- Select, From, Where clauses
- Expressions
- Group by, Joins
- Conceptual evaluation
- WITH, Views, Tables

#### Relational Constraints

- Gives the DBMS a list of consistency checks
- It is run with respect to whatever data that exists in the database
  - DBMS doesn't understand anything about what application wants or intends or "should" have
  - It simply goes through each constraint one by one and checks them against the data in the database

### Relational Constraints

Domain constraints

Foreign key constraints

Unique constraints

Primary key constraints

**NOT NULL constraints** 

CHECK constraints

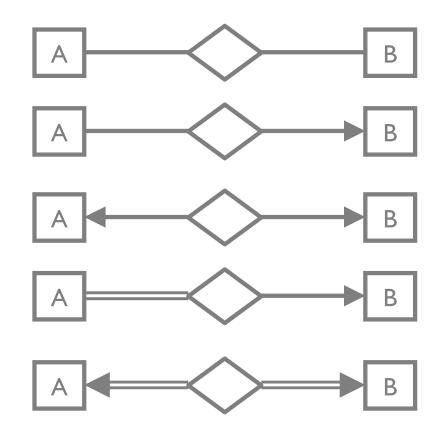
### ER -> Relational translation

Translate entities and relationships into relational tables
Translate ER constraints into relational constraints
The translation is correct if for *any* database instance:

Constraints violated in ER are violated in relational Constraints violated in relational are violated in ER If ER doesn't violate, neither should relational If relational doesn't violate, neither should ER

Note: some translations are not possible

## The following can be cleanly translated



(notice there's an "at most I" if there's an "at least I")

### How to adhoc check a translation?

- Come up with data that satisfies the ER constraints and check that they don't violate relational version
- 2. For each ER constraint
  - I. Come up with data that violated the ER constraint. Does relational version identify the violation given the same data?
- 3. Vice Versa

# Non-exhaustive examples

(entity sets don't include attrs for simplicity)

And how to check them



```
ab(
a references A(a)
b references B(b)
primary key (a, b)
)
```

B( b int primary key )

a	name
1	eve
2	bob

a	b
1	3
1	4
2	3

	b	
3		
4		



a	name
1	eve
2	bob

ab(	
a references	s A(a)
b reference	s B(b)
primary key	(a, b)
unique(a)	
)	

a	Ь
1	3
2	3

B(			
b	int	primary	key
)			

	b
3	
4	

This database satisfies ER constraints.
Thus it should satisfy the relational constraints



a	name
1	eve
2	bob

ab(
a references A(a)
b references B(b)
primary key (a, b)
unique(a)
)

a	b
1	3
2	null

B( b int primary key )

b	
3	
4	

This database violates ER constraints.

Thus, it should violate the relational constraints



a	name
1	eve
2	bob

ab(
a references A(a)
b references B(b)
primary key (a, b)
unique(a)
)

a	b
1	3
1	4
2	3

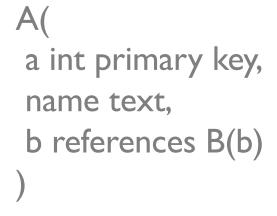
B(			
b	int	primary	key
)			

	b
3	
4	

This database violates ER constraints.

Thus, it should violate the relational constraints





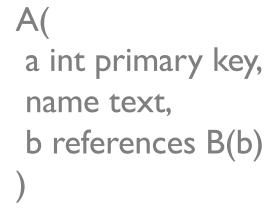
a	name	Ь
1	eve	3
2	bob	3

B(			
b	int	primary	key
)			

	b
3	
4	

This database satisfies ER constraints.
Thus should satisfy relational constraints





a	name	b
1	eve	3
2	bob	3
1	eve	4

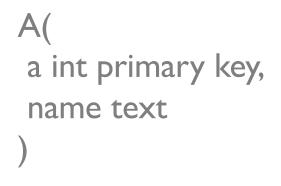
B(			
b	int	primary	key
)			

	b
3	
4	

This database violates ER constraints.

Thus should violate relational constraints





a	name
1	eve
2	bob

ab(
a references A(a)
b references B(b)
primary key (a, b)
unique(b)
)

a	Ь
1	3
2	3

B(			
b	int	primary	key
)			

	b
3	
4	

This database satisfies ER constraints.
BUT it doesn't satisfy the relational constraints.





#### Can this be expressed relationally?

A(
a int primary key,
name text
)

a	name
1	eve
2	bob

ab(
a references A(a)
b references B(b)
primary key (a, b)
unique(a)
)

b	int	primary	key
)			

ì	Ь		b
	3		3
		'	4

R/

This database violates ER constraints.
BUT it does satisfy the relational constraints.

2





```
A(
a int primary key,
b int primary key)
name text,
b references B(b) NOT NULL
)
```

a	name	b
1	eve	3
2	bob	3

	b
3	
4	

This database satisfies ER constraints. Also satisfies relational constraints. OK



```
A(
a int primary key,
b int primary key)
name text,
b references B(b) NOT NULL
)
```

a	name	b
1	eve	null
2	bob	3

	b	
3		
4		

This database violates ER constraints. Also violates relational constraints. OK



```
A(
a int primary key,
b int primary key)
name text,
b references B(b) NOT NULL
)
```

a	name	b
2	bob	3

b34

This database satisfies ER constraints. Also satisfies relational constraints. OK

# Some tips

There are not that many ways to express a relationship between A and B for any combination of constraints

```
A(...), AB(...), B(...) // all three

A_AB(...), B(...) // A and AB are merged

A(...), B_AB(...) // AB and B are merged

AB(...) // all three merged
```

### You should understand...

Why the following cannot be expressed or cannot be expressed without redundancy

