tmvar, x, y, f, m, n variables

covar, c coercion variables

 $\begin{array}{c} datacon,\ K\\ const,\ T\\ tyfam,\ F\\ index,\ i \end{array}$

index, i indices

```
Role
role, R
                                           ::=
                                                    \mathbf{Nom}
                                                    Rep
                                                    R_1 \cap R_2
                                                                                    S
relflag, \ \rho
                                                                                                          relevance flag
constraint, \phi
                                                                                                          props
                                                    a \sim_{A/R} b
                                                                                    S
S
                                                    (\phi)
                                                    \phi\{b/x\}
                                                                                    S
                                                    |\phi|
tm, a, b, v, w, A, B
                                                                                                          types and kinds
                                                    \lambda^{\rho}x:A/R.b
                                                                                    \mathsf{bind}\;x\;\mathsf{in}\;b
                                                    \lambda^{R,\rho}x.b
                                                                                    \mathsf{bind}\;x\;\mathsf{in}\;b
                                                    a b^{R,\rho}
                                                     T
                                                    \Pi^{\rho}x:A/R\to B
                                                                                    \mathsf{bind}\ x\ \mathsf{in}\ B
                                                     a \triangleright_R \gamma
                                                    \forall c : \phi.B
                                                                                    bind c in B
                                                    \Lambda c : \phi . b
                                                                                    \mathsf{bind}\ c\ \mathsf{in}\ b
                                                    \Lambda c.b
                                                                                    \mathsf{bind}\ c\ \mathsf{in}\ b
                                                     a[\gamma]
                                                    K
                                                    {f match}~a~{f with}~brs
                                                    \operatorname{\mathbf{sub}} R a
                                                                                    S
                                                     a\{b/x\}
                                                                                    S
                                                                                    S
                                                     a\{\gamma/c\}
                                                                                    S
                                                     a
                                                                                    S
                                                     (a)
                                                                                    S
                                                                                                              parsing precedence is hard
                                                                                    S
                                                    |a|R
                                                                                    S
                                                    \mathbf{Int}
                                                                                    S
                                                    Bool
                                                                                    S
                                                    Nat
                                                                                    S
                                                    Vec
                                                                                    S
                                                    0
                                                                                    S
                                                    S
                                                                                    S
                                                    True
```

```
S
                                       \mathbf{Fix}
                                                                            S
                                       a \rightarrow b
                                      \phi \Rightarrow A
                                                                            S
                                       ab^{R,+}
                                                                            S
                                       \lambda^R x.a
                                                                            S
                                                                            S
                                       \lambda x : A.a
                                      \forall\,x:A/R\to B\quad \mathsf{S}
brs
                                                                                                          case branches
                           ::=
                                       none
                                       K \Rightarrow a; brs
                                                                            S
                                       brs\{a/x\}
                                                                            S
                                       brs\{\gamma/c\}
                                                                            S
                                       (brs)
co, \gamma
                           ::=
                                                                                                          explicit coercions
                                       c
                                       \operatorname{\mathbf{red}} a\ b
                                       \mathbf{refl}\;a
                                       (a \models \mid_{\gamma} b)
                                       \mathbf{sym}\,\gamma
                                      \gamma_1; \gamma_2
                                       \mathbf{sub}\,\gamma
                                      \Pi^{R,\rho}x\!:\!\gamma_1.\gamma_2
                                                                            bind x in \gamma_2
                                      \lambda^{R,\rho} x : \gamma_1 \cdot \gamma_2
\gamma_1 \ \gamma_2^{R,\rho}
                                                                            bind x in \gamma_2
                                       \mathbf{piFst}\, \gamma
                                       \mathbf{cpiFst}\,\gamma
                                       \mathbf{isoSnd}\,\gamma
                                       \gamma_1@\gamma_2
                                       \forall c: \gamma_1.\gamma_3
                                                                            bind c in \gamma_3
                                                                            bind c in \gamma_3
                                       \lambda c: \gamma_1.\gamma_3@\gamma_4
                                       \gamma(\gamma_1,\gamma_2)
                                      \gamma@(\gamma_1 \sim \gamma_2)
                                       \gamma_1 \triangleright_R \gamma_2
                                       \gamma_1 \sim_A \gamma_2
                                       conv \phi_1 \sim_{\gamma} \phi_2
                                       \mathbf{eta}\,a
                                       left \gamma \gamma'
                                       \mathbf{right} \, \gamma \, \gamma'
                                                                            S
                                       (\gamma)
                                                                            S
                                       \gamma\{a/x\}
                                                                            S
                                                                                                          signature classifier
sig\_sort
                                       \mathbf{Cs}\,A
```

```
\mathbf{Ax}\ a\ A\ R
sort
                                      ::=
                                                                                          binding classifier
                                               \mathbf{Tm}\,A\,R
                                               \mathbf{Co}\,\phi
context, \ \Gamma
                                                                                          contexts
                                               Ø
                                               \Gamma, x : A/R
                                               \Gamma, c: \phi
                                               \Gamma\{b/x\}
                                                                                  Μ
                                               \Gamma\{\gamma/c\}
                                                                                  Μ
                                               \Gamma, \Gamma'
                                                                                  Μ
                                               |\Gamma|
                                                                                  Μ
                                               (\Gamma)
                                                                                  Μ
                                                                                  Μ
role\_context, \ \Omega
                                                                                          {\rm role}_contexts
                                      ::=
                                               Ø
                                               \Omega, x:R,\Omega'
                                               \Omega, c: R, \Omega'
                                               \Omega, \Omega'
                                                                                  Μ
                                               (\Omega)
                                                                                  Μ
                                                                                  Μ
available\_props,\ \Delta
                                               Ø
                                               \Delta, c
                                               \widetilde{\Gamma}
                                                                                  Μ
                                                                                  Μ
sig, \Sigma
                                                                                          signatures
                                               Ø
                                               \Sigma \cup \{T: A/R\}
                                               \Sigma \cup \{F \sim a : A/R\}
                                               \Sigma_0
                                                                                  Μ
                                               \Sigma_1
                                                                                  Μ
                                               |\Sigma|
                                                                                  Μ
terminals
                                               \leftrightarrow
                                               \Leftrightarrow
                                               min
```

 \in

```
\not\in
\Leftarrow
\Lambda
F
\neq
\triangleright
 ok
Ø
fv
dom
\asymp
\mathbf{fst}
\operatorname{snd}
|\Rightarrow|
\vdash_=
\mathbf{refl_2}
++
judgement
x: A/R \in \Gamma
c: \phi \in \Gamma
T:A/R\in\Sigma
F \sim a : A/R \in \Sigma
K:T\Gamma\in\Sigma
x \in \Delta
c\,\in\,\Delta
c \, \mathbf{not} \, \mathbf{relevant} \, \in \, \gamma
x \not\in \mathsf{fv} a
```

 $x\not\in\operatorname{dom}\Gamma$

 $formula, \psi$

```
c \not\in \operatorname{dom} \Gamma
                              T \not\in \mathsf{dom}\, \Sigma
                              F \not\in \mathsf{dom}\, \Sigma
                              a = b
                             \phi_1 = \phi_2
                             \Gamma_1 = \Gamma_2
                              \gamma_1 = \gamma_2
                              \neg \psi
                             \psi_1 \wedge \psi_2
                             \psi_1 \vee \psi_2
                             \psi_1 \Rightarrow \psi_2
                              (\psi)
                              c:(a:A\sim b:B)\in\Gamma
                                                                       suppress lc hypothesis generated by Ott
JSubRole
                             R_1 \leq R_2
                                                                       Subroling judgement
JValue
                      ::=
                             \mathbf{CoercedValue}\,R\,A
                                                                       Values with at most one coercion at the top
                             Value_R A
                                                                       values
                              Value Type RA
                                                                       Types with head forms (erased language)
Jconsistent
                             consistent a b
                                                                       (erased) types do not differ in their heads
Jerased
                      ::=
                             \Omega \vdash erased\_tma\:R
JChk
                      ::=
                              (\rho = +) \lor (x \not\in \mathsf{fv}\ A)
                                                                       irrelevant argument check
Jpar
                             \Omega \vDash a \Rightarrow_R b
                                                                       parallel reduction (implicit language)
                             \Omega \vdash a \Rightarrow_R^* b
                                                                       multistep parallel reduction
                             \Omega \vdash a \Leftrightarrow_R b
                                                                       parallel reduction to a common term
Jbeta
                      ::=
                             \vDash a > b/R
                                                                       primitive reductions on erased terms
                             \models a \leadsto \dot{b}/R
                                                                       single-step head reduction for implicit language
                             \models a \leadsto^* b/R
                                                                       multistep reduction
Jett
                      ::=
                             \Gamma \vDash \phi ok
                                                                       Prop wellformedness
                             \Gamma \vDash a : A/R
                                                                       typing
                             \Gamma; \Delta \vDash \phi_1 \equiv \phi_2
                                                                       prop equality
```

```
\Gamma; \Delta \vDash a \equiv b : A/R
                                \models \Gamma
Jsig
                        ::=
                                \models \Sigma
Jann
                         ::=
                                 \Gamma \vdash \phi ok
                                \Gamma \vdash a : A/R
                                \Gamma; \Delta \vdash \gamma : \phi_1 \sim \phi_2
                                \Gamma; \Delta \vdash \gamma : A \sim_R B
                                \vdash \Gamma
                                \vdash \Sigma
Jred
                        ::=
                                \Gamma \vdash a \leadsto b/R
judgement
                         ::=
                                 JSubRole
                                 JValue
                                 J consistent
                                 Jerased
                                 JChk
                                 Jpar
                                 Jbeta
                                 Jett
                                 Jsig
                                 Jann
                                 Jred
user\_syntax
                         ::=
                                 tmvar
                                 covar
                                 data con
                                 const
                                 tyfam
                                 index
                                 role
                                 relflag
                                 constraint
                                 tm
                                 brs
                                 co
                                 sig\_sort
                                 sort
                                 context
```

 $role_context$

definitional equality

 $context\ well formedness$

signature wellformedness

prop wellformedness

coercion between props

coercion between types

 $context\ well formedness$

signature wellformedness

single-step, weak head reduction to values for annotated lang

typing

| available_props | sig | terminals | formula

$R_1 \leq R_2$ Subroling judgement

$$\label{eq:correction} \begin{split} \frac{\mathsf{Value}_R\ a}{\mathbf{CoercedValue}\ R\ a} & \quad \mathrm{CV} \\ \\ \frac{\mathsf{Value}_R\ a}{\neg (R_1 \leq R)} & \quad \mathrm{CC} \\ \\ \overline{\mathbf{CoercedValue}\ R\ (a \rhd_{R_1} \gamma)} & \quad \mathrm{CC} \end{split}$$

 $Value_R A$ values

$$\overline{\text{Value}_R} \star \quad \text{Value_STAR}$$

$$\overline{\text{Value}_R} \ \overline{\Pi^\rho x \colon A/R_1 \to B} \quad \text{Value_CPI}$$

$$\overline{\text{Value}_R} \ \forall c \colon \phi \ldotp B \quad \text{Value_CPI}$$

$$\overline{\text{Value}_R} \ \lambda^+ x \colon A/R_1 \ldotp a \quad \text{Value_AbsRel}$$

$$\overline{\text{Value}_R} \ \lambda^{R_1,+} x \ldotp a \quad \text{Value_UAbsRel}$$

$$\overline{\text{Value}_R} \ \lambda^{R_1,+} x \ldotp a \quad \text{Value_UAbsIrrel}$$

$$\overline{\text{Value}_R} \ \lambda^{R_1,-} x \ldotp a \quad \text{Value_AbsIrrel}$$

$$\overline{\text{Value}_R} \ \lambda^- x \colon A/R_1 \ldotp a \quad \text{Value_AbsIrrel}$$

$$\overline{\text{Value}_R} \ \Lambda c \colon \phi \ldotp a \quad \text{Value_CAbs}$$

$$\overline{\text{Value}_R} \ \Lambda c \colon a \quad \text{Value_UCAbs}$$

$$F \sim a \colon A/R_1 \in \Sigma_0$$

$$\neg (R_1 \le R) \quad \text{Value_Ax}$$

 $\overline{\mathbf{ValueType}\,R\,\star}$ VALUE_TYPE_STAR

```
value_type_Pi
                                   \overline{\mathbf{ValueType}\,R\,\Pi^{\rho}x\!:\!A/R_1	o B}
                                        \overline{\mathbf{ValueType}\,R\,orall c\!:\!\phi.B} VALUE_TYPE_CPI
                                           F \sim a : A/R_1 \in \Sigma_0
                                          \frac{\neg (R_1 \le R)}{\mathbf{ValueType} \, R \, F} \quad \text{VALUE\_TYPE\_AX}
                             (erased) types do not differ in their heads
consistent a b
                                           \frac{}{\text{consistent} \star \star} Consistent_a_Star
                                                                                                    CONSISTENT_A_PI
                 \overline{\mathbf{consistent}\left(\Pi^{\rho}x_{1}:A_{1}/R\to B_{1}\right)\left(\Pi^{\rho}x_{2}:A_{2}/R\to B_{2}\right)}
                                                                                     CONSISTENT_A_CPI
                            \overline{\mathbf{consistent} (\forall c_1 : \phi_1.A_1) (\forall c_2 : \phi_2.A_2)}
                                       ¬ValueType R b CONSISTENT_A_STEP_R
                                       \negValueType R a Consistent_A_Step_L
\Omega \vdash erased\_tma\ R
                                                                         ERASED_A_BULLET
                                          \overline{\Omega \vdash erased\_tm \Box R}
                                                                            ERASED_A_STAR
                                           \overline{\Omega \vdash erased\_tm \star R}
                                                                                   ERASED_A_VAR
                                      \overline{\Omega, x: R, \Omega' \vdash erased\_tmx R}
                                     \Omega, x: R_1, \Omega' \vdash erased\_tma R
                                                                                      ERASED_A_ABS
                                    \overline{\Omega, \Omega' \vdash erased\_tm(\lambda^{R_1, \rho}x.a) R}
                                             \Omega \vdash erased\_tmaR
                                             \Omega \vdash erased\_tmb\ R
                                                                                 ERASED_A_APP
                                       \overline{\Omega \vdash erased\_tm(a\ b^{R_1,\rho})\,R}
                                      \Omega \vdash erased\_tmAR_2
                                      \Omega, x: R_1, \Omega' \vdash erased\_tmBR
                                      R_1 \leq R_2
                                                                                             erased_a_Pi
                               \overline{\Omega, \Omega' \vdash erased\_tm(\Pi^{\rho}x : A/R_1 \to B) R}
                                     \Omega \vdash erased\_tma\ R_1
                                     \Omega \vdash erased\_tmb R_1
                                     \Omega \vdash erased\_tmA R_2
                                     \Omega, c: R_1, \Omega' \vdash erased\_tmB\:R
                                     R_1 \leq R_2
                                                                                            ERASED_A_CPI
                              \frac{1 - \frac{1}{2}}{\Omega, \Omega' \vdash erased\_tm(\forall c : a \sim_{A/R_1} b.B) R}
                                    \frac{\Omega, c: R_1, \Omega' \vdash erased\_tmb R}{\Omega, \Omega' \vdash erased\_tm(\Lambda c.b) R} \quad \text{ERASED\_A\_CABS}
                                            \Omega \vdash erased\_tma\ R
                                                                               ERASED_A_CAPP
                                        \overline{\Omega \vdash erased\_tm(a[\bullet]) R}
                                                                            ERASED_A_FAM
                                             \overline{\Omega \vdash erased\_tmFR}
```

$$\frac{\Omega \vDash a_1 \Rightarrow_{R_1} a_2}{\Omega \vDash a_1 \triangleright_{R} \bullet \Rightarrow_{R_1} a_2 \triangleright_{R} \bullet} \quad \text{Par_Cong}$$

$$\frac{\Omega \vDash a_1 \Rightarrow_{R_1} (a_2 \triangleright_{R} \bullet)}{\Omega \vDash (a_1 \triangleright_{R} \bullet) \Rightarrow_{R_1} (a_2 \triangleright_{R} \bullet)} \quad \text{Par_Combine}$$

$$\frac{\Omega \vDash a_1 \Rightarrow_{R_1} (a_2 \triangleright_{R} \bullet)}{\Omega \vDash a_1 \Rightarrow_{R_1} (a_2 \triangleright_{R} \bullet)} \quad \text{Par_Combine}$$

$$\frac{\Omega \vDash a_1 \Rightarrow_{R_1} (a_2 \triangleright_{R} \bullet)}{\Omega \vDash a_1 b_1^{R_2,+} \Rightarrow_{R_1} (a_2 (b_2 \triangleright_{R} \bullet)^{R_2,+}) \triangleright_{R} \bullet} \quad \text{Par_Push}$$

$$\frac{\Omega \vDash a_1 \Rightarrow_{R_1} (a_2 \triangleright_{R} \bullet)}{\Omega \vDash a_1 [\bullet] \Rightarrow_{R_1} (a_2 [\bullet]) \triangleright_{R} \bullet} \quad \text{Par_CPush}$$

 $\Omega \vdash a \Rightarrow_R^* b$ multistep parallel reduction

$$\frac{\Omega \vdash a \Rightarrow_{R}^{*} a}{\Omega \vdash a \Rightarrow_{R}^{*} b} \text{MP-Refl}$$

$$\frac{\Omega \vdash a \Rightarrow_{R} b}{\Omega \vdash b \Rightarrow_{R}^{*} a'}$$

$$\frac{\Omega \vdash a \Rightarrow_{R}^{*} a'}{\Omega \vdash a \Rightarrow_{R}^{*} a'}$$
MP-Step

 $\Omega \vdash a \Leftrightarrow_R b$ parallel reduction to a common term

$$\frac{\Omega \vdash a_1 \Rightarrow_R^* b}{\Omega \vdash a_2 \Rightarrow_R^* b} \\
\frac{\Omega \vdash a_1 \Leftrightarrow_R a_2}{\Omega \vdash a_1 \Leftrightarrow_R a_2} \quad \text{JOIN}$$

 $\models a > b/R$ primitive reductions on erased terms

$$\frac{\mathsf{Value}_{R_1} \ (\lambda^{R,\rho} x.v)}{\vDash (\lambda^{R,\rho} x.v) \ b^{R,\rho} > v\{b/x\}/R_1} \quad \mathsf{BETA_APPABS}$$

$$\frac{\vdash (\Lambda c.a')[\bullet] > a'\{\bullet/c\}/R}{\vdash (\Lambda c.a')[\bullet] > a'\{\bullet/c\}/R} \quad \mathsf{BETA_CAPPCABS}$$

$$\frac{F \sim a : A/R \in \Sigma_0}{\vDash F > a/R} \quad \mathsf{BETA_AXIOM}$$

 $\models a \leadsto b/R$ single-step head reduction for implicit language

$$\begin{array}{c} \models a \leadsto a'/R_1 \\ \hline \models \lambda^{R,-}x.a \leadsto \lambda^{R,-}x.a'/R_1 \\ \hline \models a \leadsto a'/R_1 \\ \hline \models a \ b^{R,\rho} \leadsto a' \ b^{R,\rho}/R_1 \\ \hline \hline \models a \ b^{R,\rho} \leadsto a' \ b^{R,\rho}/R_1 \\ \hline \hline \vdash a \ b^{R,\rho} \leadsto a' \ b^{R,\rho}/R_1 \\ \hline \hline \vdash a \ b^{R,\rho} \leadsto a' \ b^{R,\rho}/R_1 \\ \hline \hline \vdash a \ b^{R,\rho} \leadsto a' \ b^{R,\rho}/R_1 \\ \hline \hline \hline \begin{array}{c} \vdash a \leadsto a'/R \\ \hline \vdash a \ b^{R,\rho} \leadsto a' \ b^{R,\rho}/R_1 \\ \hline \hline \hline \hline \\ \hline \hline \hline \\ \hline \hline \hline \\ \hline \\ \hline \hline \\ \hline$$

$$\begin{array}{c}
\vDash a \leadsto a'/R_1 \\
\neg(R \le R_1) \\
\vDash a \triangleright_R \bullet \leadsto a' \triangleright_R \bullet/R_1
\end{array}
\quad \text{E_CONG}$$

$$\frac{}{\vDash (a \triangleright_R \bullet) \triangleright_R \bullet \leadsto a \triangleright_R \bullet/R_1} \quad \text{E_COMBINE}$$

$$\frac{}{\vDash (v_1 \triangleright_R \bullet) v_2^{R_1,+} \leadsto (v_1 (v_2 \triangleright_R \bullet)^{R_1,+}) \triangleright_R \bullet/R_2} \quad \text{E_PUSH}$$

$$\frac{}{\vDash (v_1 \triangleright_R \bullet) [\bullet] \leadsto (v_1 [\bullet]) \triangleright_R \bullet/R_1} \quad \text{E_CPUSH}$$

 $\models a \leadsto^* b/R$ multistep reduction

 $\Gamma \vDash \phi$ ok Prop wellformedness

$$\begin{array}{l} \Gamma \vDash a : A/R \\ \Gamma \vDash b : A/R \\ \hline \Gamma \vDash A : \star/R \\ \hline \Gamma \vDash a \sim_{A/R} b \text{ ok} \end{array} \quad \text{E-Wff}$$

 $\Gamma \vDash a : A/R$ typing

$$R_{1} \leq R_{2}$$

$$\Gamma \vDash a : A/R_{1}$$

$$\Gamma \vDash a : A/R_{2}$$

$$E_{-}SUBROLE$$

$$E \Gamma$$

$$\Gamma \vDash x : A/R \qquad E_{-}STAR$$

$$E \Gamma$$

$$X : A/R \vDash \Gamma$$

$$\Gamma \vDash x : A/R \qquad E_{-}VAR$$

$$\Gamma, x : A/R \vDash B : */R'$$

$$\Gamma \vDash A : */R$$

$$R \leq R'$$

$$\Gamma \vDash \Pi^{\rho}x : A/R \Rightarrow B : */R'$$

$$\Gamma \vDash A : */R$$

$$(\rho = +) \lor (x \not\in \text{fv } a)$$

$$R \leq R'$$

$$\Gamma \vDash b : \Pi^{+}x : A/R \Rightarrow B/R'$$

$$\Gamma \vDash a : A/R$$

$$\Gamma \vDash b : \Pi^{+}x : A/R \Rightarrow B/R'$$

$$\Gamma \vDash a : A/R$$

$$\Gamma \vDash b : \Pi^{-}x : A/R \Rightarrow B/R'$$

$$\Gamma \vDash b : \Pi^{-}x : A/R \Rightarrow B/R'$$

$$\Gamma \vDash b : \Pi^{-}x : A/R \Rightarrow B/R'$$

$$\Gamma \vDash b : \Pi^{-}x : A/R \Rightarrow B/R'$$

$$\Gamma \vDash b : \Pi^{-}x : A/R \Rightarrow B/R'$$

$$\Gamma \vDash a : A/R$$

$$\Gamma \vDash b : \Pi^{-}x : A/R \Rightarrow B/R'$$

$$\Gamma \vDash a : A/R$$

$$\Gamma \vDash b : \Pi^{-}x : A/R \Rightarrow B/R'$$

$$\Gamma \vDash a : A/R$$

$$\Gamma \vDash b : \Pi^{-}x : A/R \Rightarrow B/R'$$

$$\Gamma \vDash a : A/R$$

$$\Gamma \vDash b : \Pi^{-}x : A/R \Rightarrow B/R'$$

$$\Gamma \vDash a : A/R$$

$$\Gamma \vDash b : \Pi^{-}x : A/R \Rightarrow B/R'$$

$$\Gamma \vDash a : A/R$$

$$\Gamma \vDash b : \Pi^{-}x : A/R \Rightarrow B/R'$$

$$\Gamma \vDash a : A/R$$

$$\Gamma \vDash b : \Pi^{-}x : A/R \Rightarrow B/R'$$

$$\Gamma \vDash a : A/R$$

$$\Gamma \vDash b : \Pi^{-}x : A/R \Rightarrow B/R'$$

$$\Gamma \vDash a : A/R$$

$$\Gamma \vDash b : \Pi^{-}x : A/R \Rightarrow B/R'$$

$$\Gamma \vDash a : A/R$$

$$\Gamma \vDash b : \Pi^{-}x : A/R \Rightarrow B/R'$$

$$\Gamma \vDash a : A/R$$

$$\Gamma \vDash b : \Pi^{-}x : A/R \Rightarrow B/R'$$

$$\Gamma \vDash a : A/R$$

$$\Gamma \vDash b : \Pi^{-}x : A/R \Rightarrow B/R'$$

$$\Gamma \vDash a : A/R$$

$$\Gamma \vDash b : \Pi^{-}x : A/R \Rightarrow B/R'$$

$$\Gamma \vDash a : A/R$$

$$\Gamma \vDash b : \Pi^{-}x : A/R \Rightarrow B/R'$$

$$\Gamma \vDash a : A/R$$

$$\Gamma \vDash b : \Pi^{-}x : A/R \Rightarrow B/R'$$

$$\Gamma \vDash a : A/R$$

$$\Gamma \vDash b : \Pi^{-}x : A/R \Rightarrow B/R'$$

$$\Gamma \vDash a : A/R$$

$$\Gamma \vDash b : \Pi^{-}x : A/R \Rightarrow B/R'$$

$$\Gamma \vDash a : A/R$$

$$\Gamma \vDash a : A/R$$

$$\Gamma \vDash b : \Pi^{-}x : A/R \Rightarrow B/R'$$

$$\Gamma \vDash a : A/R$$

$$\Gamma \vDash b : \Pi^{-}x : A/R \Rightarrow B/R'$$

$$\Gamma \vDash a : A/R$$

$$\Gamma \vDash b : \Pi^{-}x : A/R \Rightarrow B/R'$$

$$\Gamma \vDash a : A/R$$

$$\Gamma \vDash b : \Pi^{-}x : A/R \Rightarrow B/R'$$

$$\Gamma \vDash a : A/R$$

$$\Gamma \vDash b : \Pi^{-}x : A/R \Rightarrow B/R'$$

$$\Gamma \vDash a : A/R$$

$$\Gamma \vDash b : \Pi^{-}x : A/R \Rightarrow B/R'$$

$$\Gamma \vDash a : A/R$$

$$\Gamma \vDash b : \Pi^{-}x : A/R \Rightarrow B/R'$$

$$\Gamma \vDash a : A/R$$

$$\Gamma$$

$$\Gamma \vDash a : A/R$$

$$\Gamma; \widetilde{\Gamma} \vDash A \equiv B : \star/R$$

$$\Gamma \vDash a : B/R$$

$$\Gamma \vDash a : B/R$$

$$\Gamma \vDash a : B/R$$

$$\Gamma \vDash \phi \text{ ok}$$

$$\Gamma \vDash \forall c : \phi : B : \star/R$$

$$\Gamma \vDash \phi \text{ ok}$$

$$\Gamma \vDash \forall c : \phi : B : \star/R$$

$$\Gamma \vDash \phi \text{ ok}$$

$$\Gamma \vDash Ac \cdot a : \forall c : \phi \cdot B/R$$

$$\Gamma \vDash a_1 : \forall c : (a \sim_{A/R} b) \cdot B_1/R'$$

$$\Gamma; \widetilde{\Gamma} \vDash a \equiv b : A/R$$

$$\Gamma \vDash a_1 : \forall c : (a \sim_{A/R} b) \cdot B_1/R'$$

$$\Gamma; \widetilde{\Gamma} \vDash a \equiv b : A/R$$

$$\Gamma \vDash a : A_1/R$$

$$\Gamma \vDash a : A_1/R \in \Sigma_0$$

$$\emptyset \vDash A : \star/R$$

$$\Gamma \vDash a : A_1/R_1$$

$$\Gamma; \widetilde{\Gamma} \vDash A_1 \equiv A_2 : \star/R_2$$

$$\neg(R_2 \le R_1)$$

$$\Gamma \vDash A_2 : \star/R_2$$

$$\Gamma \vDash a \triangleright_{R_2} \bullet : A_2/R_1$$
E-TyCast equality
$$\Gamma; \Delta \vDash A_1 \equiv A_2 : A/R$$

$$\Gamma; \Delta \vDash B_1 \equiv B_2 : A/R$$

$$\Gamma; \Delta \vDash B_1 \equiv B_2 : A/R$$

$$\Gamma; \Delta \vDash A_1 \sim_{A/R} B_1 \equiv A_2 \sim_{A/R} B_2$$

$$\Sigma = \text{PropCon}$$

 $\Gamma; \Delta \vDash \phi_1 \equiv \phi_2$ prop equality

$$\begin{array}{c} \Gamma; \Delta \vDash A_1 \equiv A_2 : A/R \\ \Gamma; \Delta \vDash B_1 \equiv B_2 : A/R \\ \hline \Gamma; \Delta \vDash A_1 \sim_{A/R} B_1 \equiv A_2 \sim_{A/R} B_2 \end{array} \quad \text{E_PropCong} \\ \Gamma; \Delta \vDash A_1 \sim_{A/R} B_1 \equiv A_2 \sim_{A/R} B_2 \\ \Gamma; \Delta \vDash A_1 \approx_{A/R} A_2 \text{ ok} \\ \Gamma \vDash A_1 \sim_{A/R} A_2 \text{ ok} \\ \hline \Gamma; \Delta \vDash A_1 \sim_{A/R} A_2 \equiv A_1 \sim_{B/R} A_2 \\ \hline \Gamma; \Delta \vDash \forall c : \phi_1.B_1 \equiv \forall c : \phi_2.B_2 : \star/R \\ \hline \Gamma; \Delta \vDash \phi_1 \equiv \phi_2 \end{array} \quad \text{E_CPiFst} \\ \hline \end{array}$$

 $\Gamma; \Delta \vDash a \equiv b : A/R$ definitional equality

$$\begin{array}{l} \vDash \Gamma \\ c: (a \sim_{A/R} b) \in \Gamma \\ \hline c \in \Delta \\ \hline \Gamma; \Delta \vDash a \equiv b: A/R \end{array} \quad \text{E_ASSN} \\ \hline \frac{\Gamma \vDash a: A/R}{\Gamma; \Delta \vDash a \equiv a: A/R} \quad \text{E_REFL} \\ \hline \frac{\Gamma; \Delta \vDash b \equiv a: A/R}{\Gamma; \Delta \vDash a \equiv b: A/R} \quad \text{E_SYM} \\ \hline \end{array}$$

```
\Gamma; \Delta \vDash a \equiv a_1 : A/R
                                   \Gamma; \Delta \vDash a_1 \equiv b : A/R
                                                                                 E_Trans
                                    \Gamma; \Delta \vDash a \equiv b : A/R
                                      \Gamma; \Delta \vDash a \equiv b : A/R_1
                                      R_1 \leq R_2
                                                                                   E_Sub
                                    \overline{\Gamma: \Delta \vDash a \equiv b: A/R_2}
                                           \Gamma \vDash a_1 : B/R
                                           \Gamma \vDash a_2 : B/R
                                           \models a_1 > a_2/R
                                                                                 E_Beta
                                    \overline{\Gamma; \Delta \vDash a_1 \equiv a_2 : B/R}
                       \Gamma; \Delta \vDash A_1 \equiv A_2 : \star / R
                       \Gamma, x: A_1/R; \Delta \vDash B_1 \equiv B_2: \star/R'
                       \Gamma \vDash A_1 : \star / R
                       \Gamma \vDash \Pi^{\rho} x : A_1/R \to B_1 : \star/R'
                       \Gamma \vDash \Pi^{\rho} x : A_2/R \rightarrow B_2 : \star/R'
                       R \leq R'
                                                                                                              E_PiCong
 \overline{\Gamma; \Delta \vDash (\Pi^{\rho}x : A_1/R \to B_1) \equiv (\Pi^{\rho}x : A_2/R \to B_2) : \star/R'}
                     \Gamma, x: A_1/R; \Delta \vDash b_1 \equiv b_2: B/R'
                     \Gamma \vDash A_1 : \star / R
                      R \leq R'
                     (\rho = +) \lor (x \not\in \mathsf{fv}\ b_1)
                     (\rho = +) \lor (x \not\in \mathsf{fv}\ b_2)
                                                                                                           E_ABSCONG
\overline{\Gamma;\Delta\vDash(\lambda^{R,\rho}x.b_1)\equiv(\lambda^{R,\rho}x.b_2):(\Pi^{\rho}x\!:\!A_1/R\to B)/R'}
                \Gamma; \Delta \vDash a_1 \equiv b_1 : (\Pi^+ x : A/R \to B)/R'
               \Gamma; \Delta \vDash a_2 \equiv b_2 : A/R
                                                                                                E_AppCong
          \Gamma; \Delta \vDash a_1 \ a_2^{R,+} \equiv b_1 \ b_2^{R,+} : (B\{a_2/x\})/R'
              \Gamma; \Delta \vDash a_1 \equiv b_1 : (\Pi^- x : A/R \rightarrow B)/R'
              \Gamma \vDash a : A/R
          \overline{\Gamma; \Delta \vDash a_1 \square^{R,-} \equiv b_1 \square^{R,-} : (B\{a/x\})/R'} E_IAPPCONG
      \frac{\Gamma; \Delta \vDash \Pi^{\rho} x : A_1/R \to B_1 \equiv \Pi^{\rho} x : A_2/R \to B_2 : \star/R'}{\Gamma; \Delta \vDash A_1 \equiv A_2 : \star/R}
      \Gamma; \Delta \vDash \Pi^{\rho} x : A_1/R \to B_1 \equiv \Pi^{\rho} x : A_2/R \to B_2 : \star/R'
      \Gamma; \Delta \vDash a_1 \equiv a_2 : A_1/R
                                                                                                              E_PiSnd
                    \Gamma; \Delta \vDash B_1\{a_1/x\} \equiv B_2\{a_2/x\} : \star/R'
                 \Gamma; \Delta \vDash a_1 \sim_{A_1/R} b_1 \equiv a_2 \sim_{A_2/R} b_2
                 \Gamma, c: a_1 \sim_{A_1/R} b_1; \Delta \vDash A \equiv B: \star/R'
                 \Gamma \vDash a_1 \sim_{A_1/R} b_1 ok
                 \Gamma \vDash \forall c : a_1 \sim_{A_1/R} b_1 . A : \star/R'
                \Gamma \vDash \forall c : a_2 \sim_{A_2/R} b_2 . B : \star / R'
                                                                                                          E_CPICONG
 \overline{\Gamma; \Delta \vDash \forall c : a_1 \sim_{A_1/R} b_1 . A \equiv \forall c : a_2 \sim_{A_2/R} b_2 . B : \star / R'}
                         \Gamma, c: \phi_1; \Delta \vDash a \equiv b: B/R
                                                                                         E_CABSCONG
               \overline{\Gamma; \Delta \vDash (\Lambda c.a) \equiv (\Lambda c.b) : \forall c : \phi_1.B/R}
             \Gamma; \Delta \vDash a_1 \equiv b_1 : (\forall c : (a \sim_{A/R} b).B)/R'
            \Gamma; \widetilde{\Gamma} \vDash a \equiv b : A/R
                \Gamma; \Delta \vDash a_1[\bullet] \equiv b_1[\bullet] : (B\{\bullet/c\})/R' E_CAPPCONG
```

$$\Gamma; \Delta \vDash \forall c : (a_1 \sim_{A/R} a_2).B_1 \equiv \forall c : (a_1' \sim_{A'/R'} a_2').B_2 : \star/R_0$$

$$\Gamma; \widetilde{\Gamma} \vDash a_1 \equiv a_2 : A/R$$

$$\Gamma; \widetilde{\Gamma} \vDash a_1' \equiv a_2' : A'/R'$$

$$\Gamma; \Delta \vDash B_1 \{ \bullet/c \} \equiv B_2 \{ \bullet/c \} : \star/R_0$$

$$\Gamma; \Delta \vDash a \equiv b : A/R$$

$$\Gamma; \Delta \vDash a \sim_{A/R} b \equiv a' \sim_{A'/R'} b'$$

$$\Gamma; \Delta \vDash a' \equiv b' : A'/R'$$

$$\Gamma; \Delta \vDash a \equiv b : A/R_1$$

$$\Gamma; \widetilde{\Gamma} \vDash A \equiv B : \star/R_2$$

$$R_1 \leq R_2$$

$$\Gamma; \Delta \vDash a \equiv b : B/R_2$$

$$\Gamma; \Delta \vDash a \equiv b : B/R_2$$

$$E_CAST$$

$$\Gamma; \Delta \vDash a \equiv b : A/R_1$$

$$\Gamma; \widetilde{\Gamma} \vDash A \equiv B : \star/R_2$$

$$\Gamma; \Delta \vDash a \equiv b : A/R_1$$

$$\Gamma; \Delta \vDash a \equiv a \implies A/R_1$$

$$\Gamma; \Delta \vDash a \implies A/R_$$

$\models \Gamma$ context wellformedness

$$\overline{\models \varnothing} \quad \text{E-EMPTY}$$

$$\vDash \Gamma$$

$$\Gamma \vDash A : \star / R$$

$$x \not\in \text{dom } \Gamma$$

$$\vDash \Gamma, x : A / R$$

$$\vDash \Gamma$$

$$\Gamma \vDash \phi \text{ ok}$$

$$\frac{c \not\in \text{dom } \Gamma}{\vDash \Gamma, c : \phi}$$

$$\text{E-ConsCo}$$

$\models \Sigma$ signature wellformedness

 $\Gamma \vdash \phi$ ok prop wellformedness

$$\begin{split} & \Gamma \vdash a : A/R \\ & \Gamma \vdash b : B/R \\ & \frac{|A|R = |B|R}{\Gamma \vdash a \sim_{A/R} b \text{ ok}} \quad \text{An_WFF} \end{split}$$

$$\Gamma \vdash a : A/R$$
 typing

$$\frac{\vdash \Gamma}{\Gamma \vdash \star : \star / R} \quad \text{An_Star}$$

$$\vdash \Gamma$$

$$\frac{x : A/R \in \Gamma}{\Gamma \vdash x : A/R} \quad \text{An_Var}$$

$$\frac{\Gamma, x : A/R \vdash B : \star / R'}{\Gamma \vdash A : \star / R} \quad \text{An_Pl}$$

$$\frac{\Gamma, x : A/R \vdash B : \star / R'}{\Gamma \vdash \Pi^p x : A/R \rightarrow B : \star / R'} \quad \text{An_Pl}$$

$$\frac{\Gamma \vdash A : \star / R}{\Gamma \vdash \Pi^p x : A/R \rightarrow B : \star / R'} \quad \text{An_Pl}$$

$$\frac{\Gamma \vdash A : \star / R}{\Gamma \vdash A : \star / R} \quad \text{End} \quad \text{An_Abs}$$

$$\frac{\Gamma \vdash A : \star / R}{\Gamma \vdash A : \star / R} \quad \text{An_Abs}$$

$$\frac{\Gamma \vdash A : \star / R}{\Gamma \vdash A : \star / R} \quad \text{An_Abs}$$

$$\frac{\Gamma \vdash b : (\Pi^p x : A/R \rightarrow B)/R'}{\Gamma \vdash b : a^R \cdot p : (B\{a/x\})/R'} \quad \text{An_App}$$

$$\frac{\Gamma \vdash a : A/R}{\Gamma \vdash B : \star / R} \quad \text{An_Conv}$$

$$\frac{\Gamma \vdash \phi \text{ ok}}{\Gamma \vdash A : b \Rightarrow \star / R} \quad \text{An_Conv}$$

$$\frac{\Gamma \vdash \phi \text{ ok}}{\Gamma \vdash A : b \Rightarrow \star / R} \quad \text{An_Conv}$$

$$\frac{\Gamma \vdash \phi \text{ ok}}{\Gamma \vdash A : b \Rightarrow \star / R} \quad \text{An_Chi}$$

$$\frac{\Gamma \vdash \alpha : (\forall c : a \sim A_1/R \ b \cdot B)/R'}{\Gamma \vdash A : (\forall c : a \sim A_1/R \ b \cdot B)/R'} \quad \text{An_CAbs}$$

$$\frac{\Gamma \vdash \alpha_1 : (\forall c : a \sim A_1/R \ b \cdot B)/R'}{\Gamma \vdash \alpha_1 : (\forall c : a \sim A_1/R \ b \cdot B)/R'} \quad \text{An_CApp}$$

$$\frac{\Gamma \vdash \alpha_1 : (\forall c : a \sim A_1/R \ b \cdot B)/R'}{\Gamma \vdash \alpha_1 : (\forall c : a \sim A_1/R \ b \cdot B)/R'} \quad \text{An_CApp}$$

$$\frac{\Gamma \vdash \alpha_1 : (\forall c : a \sim A_1/R \ b \cdot B)/R'}{\Gamma \vdash \alpha_1 : (\forall c : a \sim A_1/R \ b \cdot B)/R'} \quad \text{An_CApp}$$

$$\frac{\Gamma \vdash \alpha_1 : A/R}{\Gamma \vdash F : A/R} \quad \text{An_FAM}$$

$$\frac{R_1 \le R_2}{\Gamma \vdash a : A/R_1} \quad \text{An_SubRole}$$

$$\frac{\Gamma \vdash \Delta \vdash A_1 : A_1/R}{\Gamma \vdash a : A/R_2} \quad \text{An_SubRole}$$

$$\frac{\Gamma \vdash \Delta \vdash A_1 : A_1/R}{\Gamma \vdash A_1 \sim A_1/R} \quad \text{Bi ok}$$

$$\Gamma \vdash \Delta \vdash A_1 \sim A_1/R} \quad \text{Bi ok}$$

$$\Gamma \vdash \Delta \vdash (\gamma_1 \sim A_1/2) : (A_1 \sim A_1/R} \quad B_1) \sim (A_2 \sim A_1/R} \quad B_2)$$

$$\text{An_PropCong}$$

```
\frac{\Gamma; \Delta \vdash \gamma : \forall c : \phi_1.A_2 \sim_R \forall c : \phi_2.B_2}{\Gamma; \Delta \vdash \mathbf{cpiFst} \ \gamma : \phi_1 \sim \phi_2} \quad \text{An\_CPiFst}
                                                                  \frac{\Gamma; \Delta \vdash \gamma : \phi_1 \sim \phi_2}{\Gamma; \Delta \vdash \mathbf{sym} \ \gamma : \phi_2 \sim \phi_1} \quad \text{An_IsoSym}
                                                                     \Gamma; \Delta \vdash \gamma : A \sim_R B
                                                                      \Gamma \vdash a_1 \sim_{A/R} a_2 ok
                                                                      \Gamma \vdash a_1' \sim_{B/R} a_2' ok
                                                                      |a_1|R = |a_1'|R
                                                                      |a_2|R = |a_2'|R
      \overline{\Gamma; \Delta \vdash \mathbf{conv} \ (a_1 \sim_{A/R} a_2) \sim_{\gamma} (a_1' \sim_{B/R} a_2') : (a_1 \sim_{A/R} a_2) \sim (a_1' \sim_{B/R} a_2')}
\Gamma; \Delta \vdash \gamma : A \sim_R B
                                                   coercion between types
                                                                            \vdash \Gamma
                                                                            c: a \sim_{A/R} b \in \Gamma
                                                                            \frac{c \in \Delta}{\Gamma; \Delta \vdash c : a \sim_R b} \quad \text{An\_Assn}
                                                                      \frac{\Gamma \vdash a : A/R}{\Gamma ; \Delta \vdash \mathbf{refl} \; a : a \sim_R a} \quad \text{An\_Refl}
                                                                      \Gamma \vdash a : A/R
                                                                      \Gamma \vdash b : B/R
                                                                      |a|R = |b|R
                                                             \frac{\Gamma; \widetilde{\Gamma} \vdash \gamma : A \sim_R B}{\Gamma; \Delta \vdash (a \mid = \mid_{\gamma} b) : a \sim_R b} \quad \text{An\_eraseeq}
                                                                           \Gamma \vdash b : B/R
                                                                           \Gamma \vdash a : A/R
                                                                            \Gamma; \widetilde{\Gamma} \vdash \gamma_1 : B \sim_R A
                                                                           \Gamma; \Delta \vdash \gamma : b \sim_R a
                                                                                                                              An_Sym
                                                                       \overline{\Gamma; \Delta \vdash \mathbf{sym}\, \gamma : a \sim_R b}
                                                                        \Gamma; \Delta \vdash \gamma_1 : a \sim_R a_1
                                                                        \Gamma; \Delta \vdash \gamma_2 : a_1 \sim_R b
                                                                        \Gamma \vdash a : A/R
                                                                        \Gamma \vdash a_1 : A_1/R
                                                                        \Gamma; \widetilde{\Gamma} \vdash \gamma_3 : A \sim_R A_1
                                                                                                                           An_Trans
                                                                    \overline{\Gamma; \Delta \vdash (\gamma_1; \gamma_2) : a \sim_R b}
                                                                          \Gamma \vdash a_1 : B_0/R
                                                                          \Gamma \vdash a_2 : B_1/R
                                                                          |B_0|R = |B_1|R
                                                                 \frac{ \vdash |a_1|R > |a_2|R/R}{\Gamma; \Delta \vdash \mathbf{red} \ a_1 \ a_2 : a_1 \sim_R \ a_2}
                                                                                                                                 An_Beta
                                                         \Gamma; \Delta \vdash \gamma_1 : A_1 \sim_{R'} A_2
                                                         \Gamma, x: A_1/R; \Delta \vdash \gamma_2: B_1 \sim_{R'} B_2
                                                         B_3 = B_2\{x \triangleright_{R'} \operatorname{sym} \gamma_1/x\}
                                                         \Gamma \vdash \Pi^{\rho} x : A_1/R \to B_1 : \star/R'
                                                         \Gamma \vdash \Pi^{\rho} x : A_1/R \rightarrow B_2 : \star/R'
                                                         \Gamma \vdash \Pi^{\rho} x : A_2/R \rightarrow B_3 : \star/R'
                                                         R \leq R'
                                                                                                                                                                            An_PiCong
                   \overline{\Gamma;\Delta \vdash \Pi^{R,\rho}x \colon\! \gamma_1.\gamma_2 : (\Pi^{\rho}x \colon\! A_1/R \to B_1) \sim_{R'} (\Pi^{\rho}x \colon\! A_2/R \to B_3)}
```

```
\Gamma; \Delta \vdash \gamma_1 : A_1 \sim_R A_2
                                            \Gamma, x: A_1/R; \Delta \vdash \gamma_2: b_1 \sim_{R'} b_2
                                            b_3 = b_2\{x \triangleright_{R'} \operatorname{sym} \gamma_1/x\}
                                            \Gamma \vdash A_1 : \star / R
                                            \Gamma \vdash A_2 : \star / R
                                            (\rho = +) \lor (x \not\in \mathsf{fv} \mid b_1 \mid R')
                                            (\rho = +) \lor (x \not\in \mathsf{fv} \mid b_3 \mid R')
                                            \Gamma \vdash (\lambda^{\rho} x : A_1/R.b_2) : B/R'
                                            R \leq R'
                                                                                                                                                 An_AbsCong
              \overline{\Gamma; \Delta \vdash (\lambda^{R,\rho}x : \gamma_1.\gamma_2) : (\lambda^{\rho}x : A_1/R.b_1) \sim_{R'} (\lambda^{\rho}x : A_2/R.b_3)}
                                                      \Gamma; \Delta \vdash \gamma_1 : a_1 \sim_{R'} b_1
                                                      \Gamma; \Delta \vdash \gamma_2 : a_2 \sim_R b_2
                                                      \Gamma \vdash a_1 \ a_2^{R,\rho} : A/R'
                                                      \Gamma \vdash b_1 \ b_2^{R,\rho} : B/R'
                                                      \Gamma; \widetilde{\Gamma} \vdash \gamma_3 : A \sim_{R'} B
                                    \frac{1}{\Gamma; \Delta \vdash \gamma_1 \ \gamma_2^{R,\rho}: a_1 \ a_2^{R,\rho} \sim_{R'} b_1 \ b_2^{R,\rho}} \quad \text{An\_AppCong}
                          \Gamma; \Delta \vdash \gamma: \Pi^{\rho}x \colon A_1/R \to \underline{B_1 \sim_{R'} \Pi^{\rho}x \colon A_2/R \to B_2}
                                                    \Gamma; \Delta \vdash \mathbf{piFst} \ \gamma : A_1 \sim_R A_2
                          \Gamma; \Delta \vdash \gamma_1 : \Pi^{\rho} x : A_1/R \to B_1 \sim_{R'} \Pi^{\rho} x : A_2/R \to B_2
                          \Gamma; \Delta \vdash \gamma_2 : a_1 \sim_R a_2
                          \Gamma \vdash a_1 : A_1/R
                          \Gamma \vdash a_2 : A_2/R
                                                                                                                                                An_PiSnd
                                      \Gamma; \Delta \vdash \gamma_1 @ \gamma_2 : B_1 \{a_1/x\} \sim_{R'} B_2 \{a_2/x\}
                                   \Gamma; \Delta \vdash \gamma_1 : a_1 \sim_{A_1/R} b_1 \sim a_2 \sim_{A_2/R} b_2
                                   \Gamma, c: a_1 \sim_{A_1/R} b_1; \Delta \vdash \gamma_3: B_1 \sim_{R'} B_2
                                    B_3 = B_2\{c \triangleright_{R'} \operatorname{\mathbf{sym}} \gamma_1/c\}
                                   \Gamma \vdash \forall c : a_1 \sim_{A_1/R} b_1.B_1 : \star/R'
                                   \Gamma \vdash \forall c : a_2 \sim_{A_2/R} b_2 . B_3 : \star / R'
                                   \Gamma \vdash \forall c : a_1 \sim_{A_1/R} b_1.B_2 : \star/R'
                                                                                                                                                         An_CPiCong
      \overline{\Gamma; \Delta \vdash (\forall c : \gamma_1.\gamma_3) : (\forall c : a_1 \sim_{A_1/R} b_1.B_1) \sim_R (\forall c : a_2 \sim_{A_2/R} b_2.B_3)}
                      \Gamma; \Delta \vdash \gamma_1 : b_0 \sim_{A_1/R} b_1 \sim b_2 \sim_{A_2/R} b_3
                      \Gamma, c: b_0 \sim_{A_1/R} b_1; \Delta \vdash \gamma_3: a_1 \sim_{R'} a_2
                      a_3 = a_2 \{c \triangleright_{R'} \operatorname{\mathbf{sym}} \gamma_1/c\}
                      \Gamma \vdash (\Lambda c : b_0 \sim_{A_1/R} b_1.a_1) : \forall c : b_0 \sim_{A_1/R} b_1.B_1/R'
                      \Gamma \vdash (\Lambda c : b_0 \sim_{A_1/R} b_1.a_2) : B/R'
                      \Gamma \vdash (\Lambda c : b_2 \sim_{A_2/R} b_3.a_3) : \forall c : b_2 \sim_{A_2/R} b_3.B_2/R'
                      \Gamma; \widetilde{\Gamma} \vdash \gamma_4 : \forall c : b_0 \sim_{A_1/R} b_1.B_1 \sim_{R'} \forall c : \phi_2.B_2
                                                                                                                                                            An_CABSCONG
\overline{\Gamma; \Delta \vdash (\lambda c : \gamma_1.\gamma_3@\gamma_4) : (\Lambda c : b_0 \sim_{A_1/R} b_1.a_1) \sim_{R'} (\Lambda c : b_2 \sim_{A_2/R} b_3.a_3)}
                                                     \Gamma; \Delta \vdash \gamma_1 : a_1 \sim_R b_1
                                                     \Gamma; \widetilde{\Gamma} \vdash \gamma_2 : a_2 \sim_{R'} b_2
                                                     \Gamma; \widetilde{\Gamma} \vdash \gamma_3 : a_3 \sim_{R'} b_3
                                                     \Gamma \vdash a_1[\gamma_2] : A/R
                                                     \Gamma \vdash b_1[\gamma_3] : B/R
                                                     \Gamma; \Gamma \vdash \gamma_4 : A \sim_R B
                                                                                                                 An_CAppCong
                                     \overline{\Gamma; \Delta \vdash \gamma_1(\gamma_2, \gamma_3) : a_1[\gamma_2] \sim_R b_1[\gamma_3]}
```

$$\begin{array}{l} \Gamma; \Delta \vdash \gamma_{1} : (\forall c_{1} : a \sim_{A/R} a'.B_{1}) \sim_{R_{0}} (\forall c_{2} : b \sim_{B/R'} b'.B_{2}) \\ \Gamma; \widetilde{\Gamma} \vdash \gamma_{2} : a \sim_{R} a' \\ \Gamma; \widetilde{\Gamma} \vdash \gamma_{3} : b \sim_{R'} b' \\ \hline \Gamma; \Delta \vdash \gamma_{1} @ (\gamma_{2} \sim \gamma_{3}) : B_{1}\{\gamma_{2}/c_{1}\} \sim_{R_{0}} B_{2}\{\gamma_{3}/c_{2}\} \\ \hline \frac{\Gamma; \Delta \vdash \gamma_{1} : a \sim_{R_{1}} a'}{\Gamma; \Delta \vdash \gamma_{2} : a \sim_{A/R_{1}} a' \sim b \sim_{B/R_{2}} b'} \quad \text{An_CAST} \\ \hline \frac{\Gamma; \Delta \vdash \gamma_{1} \triangleright_{R_{2}} \gamma_{2} : b \sim_{R_{2}} b'}{\Gamma; \Delta \vdash \gamma_{1} \triangleright_{R_{2}} \gamma_{2} : b \sim_{R_{2}} b'} \quad \text{An_LSoSnD} \\ \hline \frac{\Gamma; \Delta \vdash \gamma : (a \sim_{A/R} a') \sim (b \sim_{B/R} b')}{\Gamma; \Delta \vdash \mathbf{isoSnd} \gamma : A \sim_{R} B} \quad \text{An_IsoSnD} \\ \hline \frac{\Gamma; \Delta \vdash \gamma : a \sim_{R_{1}} b}{\Gamma; \Delta \vdash \mathbf{sub} \gamma : a \sim_{R_{2}} b} \quad \text{An_SuB} \end{array}$$

$\vdash \Gamma$ context wellformedness

$\vdash \Sigma$ signature wellformedness

$$\begin{array}{ccc} & & & & \\ & & \vdash \varSigma \\ & \varnothing \vdash A : \star / R \\ & \varnothing \vdash a : \star / R \\ & \varnothing \vdash a : A / R \\ & & \vdash \Xi \cup \{F \sim a : A / R\} \end{array} \quad \text{An_Sig_ConsAx}$$

 $\Gamma \vdash a \leadsto b/R$ single-step, weak head reduction to values for annotated language

$$\frac{\Gamma \vdash a \leadsto a'/R_1}{\Gamma \vdash a \ b^{R,\rho} \leadsto a' \ b^{R,\rho}/R_1} \quad \text{An_AppLeft}$$

$$\frac{\text{Value}_R \ (\lambda^\rho x \colon A/R.w)}{\Gamma \vdash (\lambda^\rho x \colon A/R.w) \ a^{R,\rho} \leadsto w \{a/x\}/R} \quad \text{An_AppAbs}$$

$$\frac{\Gamma \vdash a \leadsto a'/R}{\Gamma \vdash a[\gamma] \leadsto a'[\gamma]/R} \quad \text{An_CAppLeft}$$

$$\frac{\Gamma \vdash (\Lambda c \colon \phi.b)[\gamma] \leadsto b\{\gamma/c\}/R}{\Gamma \vdash A \colon \star/R} \quad \text{An_CAppCAbs}$$

$$\frac{\Gamma \vdash A \colon \star/R}{\Gamma, x \colon A/R \vdash b \leadsto b'/R_1} \quad \text{An_AbsTerm}$$

$$\frac{\Gamma \vdash (\lambda^- x \colon A/R.b) \leadsto (\lambda^- x \colon A/R.b')/R_1}{\Gamma \vdash (\lambda^- x \colon A/R.b) \leadsto (\lambda^- x \colon A/R.b')/R_1} \quad \text{An_AbsTerm}$$

$$\frac{F \sim a : A/R \in \Sigma_{1}}{\Gamma \vdash F \leadsto a/R} \quad \text{An_AXIOM}$$

$$\frac{\Gamma \vdash a \leadsto a'/R}{\Gamma \vdash a \rhd_{R_{1}} \gamma \leadsto a' \rhd_{R_{1}} \gamma/R} \quad \text{An_ConvTerm}$$

$$\frac{\text{Value}_{R} \ v}{\Gamma \vdash (v \rhd_{R_{2}} \gamma_{1}) \rhd_{R_{2}} \gamma_{2} \leadsto v \rhd_{R_{2}} (\gamma_{1}; \gamma_{2})/R} \quad \text{An_Combine}$$

$$\text{Value}_{R} \ v$$

$$\Gamma; \widetilde{\Gamma} \vdash \gamma : \Pi^{\rho} x_{1} : A_{1}/R \to B_{1} \sim_{R'} \Pi^{\rho} x_{2} : A_{2}/R \to B_{2}$$

$$b' = b \rhd_{R'} \text{ sym } (\text{piFst } \gamma)$$

$$\gamma' = \gamma@(b') = |_{(\text{piFst } \gamma)} b)$$

$$\Gamma \vdash (v \rhd_{R'} \gamma) \ b^{R,\rho} \leadsto ((v \ b'^{R,\rho}) \rhd_{R'} \gamma')/R} \quad \text{An_Push}$$

$$\text{Value}_{R} \ v$$

$$\Gamma; \widetilde{\Gamma} \vdash \gamma : \forall c_{1} : a_{1} \sim_{B_{1}/R} b_{1}.A_{1} \sim_{R'} \forall c_{2} : a_{2} \sim_{B_{2}/R} b_{2}.A_{2}$$

$$\gamma'_{1} = \gamma_{1} \rhd_{R'} \text{ sym } (\text{cpiFst } \gamma)$$

$$\gamma' = \gamma@(\gamma'_{1} \sim \gamma_{1})$$

$$\Gamma \vdash (v \rhd_{R'} \gamma) [\gamma_{1}] \leadsto ((v [\gamma'_{1}]) \rhd_{R'} \gamma')/R} \quad \text{An_CPush}$$

Definition rules: 161 good 0 bad Definition rule clauses: 475 good 0 bad