tmvar, x, y, f, m, n variables

covar, c coercion variables

 $\begin{array}{c} datacon,\ K\\ const,\ T\\ tyfam,\ F\\ index,\ i \end{array}$

index, i indices

```
Role
role, R
                                           ::=
                                                    \mathbf{Nom}
                                                    Rep
                                                    R_1 \cap R_2
                                                                                    S
relflag, \ \rho
                                                                                                          relevance flag
constraint, \phi
                                                                                                          props
                                                    a \sim_{A/R} b
                                                                                    S
S
                                                    (\phi)
                                                    \phi\{b/x\}
                                                                                    S
                                                    |\phi|
tm, a, b, v, w, A, B
                                                                                                          types and kinds
                                                    \lambda^{\rho}x:A/R.b
                                                                                    \mathsf{bind}\;x\;\mathsf{in}\;b
                                                    \lambda^{R,\rho}x.b
                                                                                    \mathsf{bind}\;x\;\mathsf{in}\;b
                                                    a b^{R,\rho}
                                                     T
                                                    \Pi^{\rho}x:A/R\to B
                                                                                    \mathsf{bind}\ x\ \mathsf{in}\ B
                                                     a \triangleright_R \gamma
                                                    \forall c : \phi.B
                                                                                    bind c in B
                                                    \Lambda c : \phi . b
                                                                                    \mathsf{bind}\ c\ \mathsf{in}\ b
                                                    \Lambda c.b
                                                                                    \mathsf{bind}\ c\ \mathsf{in}\ b
                                                     a[\gamma]
                                                    K
                                                    {f match}~a~{f with}~brs
                                                    \operatorname{\mathbf{sub}} R a
                                                                                    S
                                                     a\{b/x\}
                                                                                    S
                                                                                    S
                                                     a\{\gamma/c\}
                                                                                    S
                                                     a
                                                                                    S
                                                     (a)
                                                                                    S
                                                                                                              parsing precedence is hard
                                                                                    S
                                                    |a|R
                                                                                    S
                                                    \mathbf{Int}
                                                                                    S
                                                    Bool
                                                                                    S
                                                    Nat
                                                                                    S
                                                    Vec
                                                                                    S
                                                    0
                                                                                    S
                                                    S
                                                                                    S
                                                    True
```

```
Fix
                                                                          S
                                                                         S
                                   a \rightarrow b
                                                                         S
                                   \phi \Rightarrow A
                                   ab^{R,+}
                                                                         S
                                   \lambda^R x.a
                                                                          S
                                                                          S
                                   \lambda x : A.a
                                   \forall\,x:A/R\to B\quad \mathsf{S}
brs
                      ::=
                                                                                                        case branches
                                   none
                                   K \Rightarrow a; brs
                                                                          S
                                   brs\{a/x\}
                                                                          S
                                   brs\{\gamma/c\}
                                   (brs)
co, \gamma
                                                                                                        explicit coercions
                                   \mathbf{red} \ a \ b
                                   \mathbf{refl}\;a
                                   (a \models \mid_{\gamma} b)
                                   \operatorname{\mathbf{sym}} \gamma
                                   \gamma_1; \gamma_2
                                   \mathbf{sub}\,\gamma
                                  \Pi^{R,\rho} \dot{x} : \gamma_1.\gamma_2
                                                                          \text{bind } x \text{ in } \gamma_2
                                  \lambda^{R,\rho} x : \gamma_1 \cdot \gamma_2
\gamma_1 \ \gamma_2^{R,\rho}
                                                                          \text{bind }x\text{ in }\gamma_2
                                   \mathbf{piFst}\,\gamma
                                   \mathbf{cpiFst}\,\gamma
                                   \mathbf{isoSnd}\,\gamma
                                   \gamma_1@\gamma_2
                                   \forall c: \gamma_1.\gamma_3
                                                                          bind c in \gamma_3
                                   \lambda c: \gamma_1.\gamma_3@\gamma_4
                                                                          \text{bind } c \text{ in } \gamma_3
                                   \gamma(\gamma_1,\gamma_2)
                                   \gamma@(\gamma_1 \sim \gamma_2)
                                   \gamma_1 \triangleright_R \gamma_2
                                   \gamma_1 \sim_A \gamma_2
                                   conv \phi_1 \sim_{\gamma} \phi_2
                                   \mathbf{eta}\,a
                                   left \gamma \gamma'
                                   right \gamma \gamma'
                                   (\gamma)
                                                                          S
                                   \gamma
                                   \gamma\{a/x\}
                                                                                                        binding classifier
sort
                                   \mathbf{Tm}\,A\,R
```

```
\mathbf{Co}\,\phi
sig\_sort
                                        ::=
                                                                                              signature classifier
                                                 \operatorname{\mathbf{Cs}} A
                                                 \mathbf{Ax} \ a \ A \ R
context, \ \Gamma
                                                                                              contexts
                                                 Ø
                                                 \Gamma, x : A/R
                                                 \Gamma, c: \phi
                                                 \Gamma\{b/x\}
                                                                                      Μ
                                                 \Gamma\{\gamma/c\}
                                                                                      Μ
                                                 \Gamma, \Gamma'
                                                                                      Μ
                                                 |\Gamma|
                                                                                      Μ
                                                 (\Gamma)
                                                                                      Μ
                                                                                      Μ
sig,~\Sigma
                                                                                              signatures
                                        ::=
                                                 Ø
                                                 \Sigma \cup \{\, T : A/R\}
                                                 \Sigma \cup \{F \sim a : A/R\}
                                                 \Sigma_0 \\ \Sigma_1
                                                                                      Μ
                                                                                      Μ
                                                 |\Sigma|
                                                                                      Μ
available\_props,\ \Delta
                                                 Ø
                                                 \Delta, c
                                                 \widetilde{\Gamma}
                                                                                      Μ
                                                 (\Delta)
                                                                                      Μ
role\_context, \Omega
                                                                                              role_contexts
                                                 Ø
                                                 \Omega, x:R
                                                 (\Omega)
                                                                                      Μ
                                                 \Omega
                                                                                      Μ
terminals
                                                 \leftrightarrow
                                                 \Leftrightarrow
                                                 \min
                                                 \not\in
```

```
F
                                        \neq
                                         ok
                                        Ø
                                        0
                                        fv
                                        \mathsf{dom} \\
                                        \asymp
                                        \mathbf{fst}
                                        \operatorname{snd}
                                        |\Rightarrow|
                                       \vdash_{=}
                                        \mathbf{refl_2}
                                        ++
formula, \psi
                              ::=
                                        judgement
                                        x:A/R\in\Gamma
                                        x:R\,\in\,\Omega
                                        c:\phi\,\in\,\Gamma
                                        T: A/R \, \in \, \Sigma
                                        F \sim a : A/R \in \Sigma
                                        K:T\Gamma \in \Sigma
                                        x\,\in\,\Delta
                                        c\,\in\,\Delta
                                        c \, \mathbf{not} \, \mathbf{relevant} \, \in \, \gamma
                                        x \not\in \mathsf{fv} a
                                        x\not\in\operatorname{dom}\Gamma
                                        rctx\_uniq\Omega
```

```
c \not\in \operatorname{dom} \Gamma
                              T \not\in \mathsf{dom}\, \Sigma
                              F \not\in \mathsf{dom}\, \Sigma
                              a = b
                              \phi_1 = \phi_2
                              \Gamma_1 = \Gamma_2
                              \gamma_1 = \gamma_2
                              \neg \psi
                              \psi_1 \wedge \psi_2
                              \psi_1 \vee \psi_2
                              \psi_1 \Rightarrow \psi_2
                              (\psi)
                              c:(a:A\sim b:B)\in\Gamma
                                                                       suppress lc hypothesis generated by Ott
JSubRole
                              R_1 \leq R_2
                                                                       Subroling judgement
JValue
                      ::=
                              \mathbf{CoercedValue}\,R\,A
                                                                       Values with at most one coercion at the top
                              Value_R A
                                                                       values
                              Value Type RA
                                                                       Types with head forms (erased language)
Jconsistent
                              consistent a \ b \ R
                                                                       (erased) types do not differ in their heads
Jerased
                      ::=
                             \Omega \vDash erased\_tm \; a \; R
JChk
                      ::=
                              (\rho = +) \lor (x \not\in \mathsf{fv}\ A)
                                                                       irrelevant argument check
Jpar
                             \Omega \vDash a \Rightarrow_R b
                                                                       parallel reduction (implicit language)
                             \Omega \vdash a \Rightarrow_R^* b
                                                                       multistep parallel reduction
                             \Omega \vdash a \Leftrightarrow_R b
                                                                       parallel reduction to a common term
Jbeta
                      ::=
                             \vDash a > b/R
                                                                       primitive reductions on erased terms
                             \models a \leadsto \dot{b}/R
                                                                       single-step head reduction for implicit language
                             \models a \leadsto^* b/R
                                                                       multistep reduction
Jett
                      ::=
                             \Gamma \vDash \phi ok
                                                                       Prop wellformedness
                             \Gamma \vDash a : A/R
                                                                       typing
                             \Gamma; \Delta \vDash \phi_1 \equiv \phi_2
                                                                       prop equality
```

sig

definitional equality $context\ well formedness$ signature wellformedness prop wellformedness typing coercion between props coercion between types $context\ well formedness$ signature wellformedness single-step, weak head reduction to values for annotated lang available_props role_context terminals formula

$R_1 \leq R_2$ Subroling judgement

$$\label{eq:correction} \begin{split} \frac{\mathsf{Value}_R\ a}{\mathbf{CoercedValue}\ R\ a} & \quad \mathrm{CV} \\ \\ \frac{\mathsf{Value}_R\ a}{\neg (R_1 \leq R)} & \quad \mathrm{CC} \\ \\ \overline{\mathbf{CoercedValue}\ R\ (a \rhd_{R_1} \gamma)} & \quad \mathrm{CC} \end{split}$$

 $Value_R A$ values

$$\overline{\text{Value}_R} \star \begin{array}{c} \text{Value_STAR} \\ \hline \\ \overline{\text{Value}_R} \ \overline{\Pi^\rho x \colon A/R_1 \to B} \end{array} \begin{array}{c} \text{Value_PI} \\ \hline \\ \overline{\text{Value}_R} \ \overline{Vc \colon \phi.B} \end{array} \begin{array}{c} \text{Value_CPI} \\ \hline \\ \overline{\text{Value}_R} \ \lambda^+ x \colon A/R_1.a \end{array} \begin{array}{c} \text{Value_AbsRel} \\ \hline \\ \overline{\text{Value}_R} \ \lambda^{R_1,+} x.a \end{array} \begin{array}{c} \text{Value_UAbsRel} \\ \hline \\ \overline{\text{Value}_R} \ \lambda^{R_1,+} x.a \end{array} \begin{array}{c} \text{Value_UAbsIrrel} \\ \hline \\ \overline{\text{Value}_R} \ \lambda^{R_1,-} x.a \end{array} \begin{array}{c} \text{Value_UAbsIrrel} \\ \hline \\ \overline{\text{Value}_R} \ \lambda^{R_1,-} x.a \end{array} \begin{array}{c} \text{Value_AbsIrrel} \\ \hline \\ \overline{\text{Value}_R} \ \lambda^{-} x \colon A/R_1.a \end{array} \begin{array}{c} \text{Value_CAbs} \\ \hline \\ \overline{\text{Value}_R} \ \Lambda c \colon \phi.a \end{array} \begin{array}{c} \text{Value_CAbs} \\ \hline \\ \overline{\text{Value}_R} \ \Lambda c.a \end{array} \begin{array}{c} \text{Value_UCAbs} \\ \hline \\ \overline{\text{Value}_R} \ \Lambda c.a \end{array} \begin{array}{c} \text{Value_UCAbs} \\ \hline \\ \overline{\text{Value}_R} \ \Lambda c.a \end{array} \begin{array}{c} \text{Value_Ax} \\ \hline \\ \overline{\text{Value}_R} \ F \end{array} \begin{array}{c} \text{Value_Ax} \\ \hline \\ \hline \\ \hline \\ \hline \end{array} \begin{array}{c} \text{Value_Ax} \end{array}$$

$$\overline{\text{ValueType } R \star}$$
 VALUE_TYPE_STAR

```
\overline{\mathbf{ValueType}\,R\,\Pi^{\rho}x\!:\!A/R_1	o B}
                                                                             VALUE_TYPE_CPI
                                        \overline{\mathbf{ValueType}\,R\,orall c\!:\!\phi.B}
                                          F \sim a : A/R_1 \in \Sigma_0
                                          \neg (R_1 \le R)
ValueType R F
                                                                            VALUE_TYPE_AX
consistent a b R
                                (erased) types do not differ in their heads
                                                                       CONSISTENT_A_STAR
                                         \overline{\mathbf{consistent}\,\star\star R}
               \overline{\mathbf{consistent} \left( \Pi^{\rho} x_1 : A_1/R \to B_1 \right) \left( \Pi^{\rho} x_2 : A_2/R \to B_2 \right) R'}
                                                                                         CONSISTENT_A_CPI
                          \overline{\mathbf{consistent} (\forall c_1 : \phi_1.A_1) (\forall c_2 : \phi_2.A_2) R}
                                                                         CONSISTENT_A_FAM
                                         consistent FFR'
                                       \negValueType R b
                                                                       CONSISTENT_A_STEP_R

\overline{\text{consistent } a \ b \ R}

                                       \negValueType R a
                                                                        CONSISTENT_A_STEP_L
                                        consistent a \ b \ R
\Omega \vDash erased\_tm \ a \ R
                                               rctx\_uniq\Omega
                                                                          ERASED_A_BULLET
                                         \overline{\Omega \vDash erased\_tm \square R}
                                                 rctx\_uniq\Omega
                                                                             ERASED_A_STAR
                                           \overline{\Omega \vDash erased\_tm \, \star \, R}
                                                  rctx\_uniq\Omega
                                                  x:R\in\Omega
                                                  R \leq R_1
                                                                              ERASED_A_VAR
                                           \overline{\Omega \vDash erased\_tm \ x \ R_1}
                                       \Omega, x : R_1 \vDash erased\_tm \ a \ R
                                                                                    ERASED_A_ABS
                                      \overline{\Omega \vDash erased\_tm(\lambda^{R_1,\rho}x.a)R}
                                            \Omega \vDash erased\_tm \ a \ R
                                            \Omega \vDash erased\_tm \ b \ R_1
                                                                                   ERASED_A_APP
                                      \Omega \vDash erased\_tm\left(a\ b^{R_{1},\rho}\right)R
                                        \Omega \vDash erased\_tm \ A \ R_1
                                        \Omega, x: R_1 \vDash erased\_tm \ B \ R
                                                                                            ERASED_A_PI
                                 \overline{\Omega \vDash erased\_tm\left(\Pi^{\rho}x : A/R_1 \to B\right)R}
                                           \Omega \vDash erased\_tm \ a \ R_1
                                           \Omega \vDash erased\_tm \ b \ R_1
                                           \Omega \vDash erased\_tm \ A \ R_1
                                           \Omega \vDash erased\_tm \ B \ R
                                                                                          ERASED_A_CPI
                               \overline{\Omega \vDash erased\_tm\left(\forall c : a \sim_{A/R_1} b.B\right)R}
                                           \Omega \vDash erased\_tm \ b \ R
                                                                                ERASED_A_CABS
                                       \Omega \vDash erased\_tm(\Lambda c.b)R
```

$$\begin{array}{c} \Omega \vDash erased.tm\ a\ R \\ \hline \Omega \vDash erased.tm\ (a | \bullet |)\ R \\ \hline R \vdash erased.tm\ (a | \bullet |)\ R \\ \hline R \vdash erased.tm\ (a \mid \bullet | \bullet |)\ R \\ \hline R \vdash erased.tm\ F\ R_1 \\ \hline \Omega \vDash erased.tm\ T\ R \\ \hline \Omega \vDash erased.tm\ T\ R \\ \hline \Omega \vDash erased.tm\ (a \triangleright_{R_1} \bullet)\ R \\ \hline \Omega \vDash erased.tm\ (a \triangleright_{R_1} \bullet)\ R \\ \hline \Theta \vDash erased.tm\ (a \triangleright_{R_1} \bullet)\ R \\ \hline (\rho = +) \lor (x \not\in \text{fv}\ A) \\ \hline (\rho = +) \lor (x \not\in \text{fv}\$$

 $\Omega \vdash a \Rightarrow_R^* b$ multistep parallel reduction

$$\frac{\Omega \vdash a \Rightarrow_R^* a}{\Omega \vdash a \Rightarrow_R^* a} \quad \text{MP_Refl}$$

$$\frac{\Omega \vdash a \Rightarrow_R b}{\Omega \vdash b \Rightarrow_R^* a'}$$

$$\frac{\Omega \vdash a \Rightarrow_R^* a'}{\Omega \vdash a \Rightarrow_R^* a'} \quad \text{MP_STEP}$$

 $\Omega \vdash a \Leftrightarrow_R b$ parallel reduction to a common term

$$\begin{array}{c}
\Omega \vdash a_1 \Rightarrow_R^* b \\
\Omega \vdash a_2 \Rightarrow_R^* b \\
\Omega \vdash a_1 \Leftrightarrow_R a_2
\end{array}$$
JOIN

 $\models a > b/R$ primitive reductions on erased terms

$$\frac{\mathsf{Value}_{R_1} \ (\lambda^{R,\rho} x.v)}{\vDash (\lambda^{R,\rho} x.v) \ b^{R,\rho} > v\{b/x\}/R_1} \quad \text{Beta_AppAbs}$$

$$\frac{\vdash (\Lambda c.a')[\bullet] > a'\{\bullet/c\}/R}{\vdash (\Lambda c.a')[\bullet] > a'\{\bullet/c\}/R} \quad \text{Beta_CAppCAbs}$$

$$\frac{F \sim a : A/R \in \Sigma_0}{\vDash F > a/R} \quad \text{Beta_Axiom}$$

 $\models a \leadsto b/R$ single-step head reduction for implicit language

$$\frac{\models a \leadsto a'/R_1}{\models \lambda^{R,-}x.a \leadsto \lambda^{R,-}x.a'/R_1} \quad \text{E_ABSTERM}$$

$$\frac{\models a \leadsto a'/R_1}{\models a \ b^{R,\rho} \leadsto a' \ b^{R,\rho}/R_1} \quad \text{E_APPLEFT}$$

$$\frac{\models a \leadsto a'/R}{\models a \ | \implies a'/R} \quad \text{E_CAPPLEFT}$$

$$\frac{\operatorname{Value}_{R_1} \ (\lambda^{R,\rho}x.v)}{\models (\lambda^{R,\rho}x.v) \ a^{R,\rho} \leadsto v \{a/x\}/R_1} \quad \text{E_APPABS}$$

$$\frac{\vdash (\lambda^{R,\rho}x.v) \ a^{R,\rho} \leadsto v \{a/x\}/R_1}{\models (\Lambda c.b)[\bullet] \leadsto b \{\bullet/c\}/R} \quad \text{E_CAPPCABS}$$

$$F \sim a : A/R \in \Sigma_0$$

$$\frac{R \leq R_1}{\models F \leadsto a/R_1} \quad \text{E_AXIOM}$$

$$\stackrel{\vdash}{\models a \leadsto a'/R_1}$$

$$\frac{\neg (R \leq R_1)}{\models a \bowtie_R \bullet \leadsto a' \bowtie_R \bullet/R_1} \quad \text{E_CONG}$$

$$\frac{\vdash (a \bowtie_R \bullet) \bowtie_R \bullet \leadsto a \bowtie_R \bullet/R_1}{\models (v_1 \bowtie_R \bullet) v_2^{R_1,+} \leadsto (v_1(v_2 \bowtie_R \bullet)^{R_1,+}) \bowtie_R \bullet/R_2} \quad \text{E_PUSH}$$

$$\stackrel{\vdash}{\models} (v_1 \bowtie_R \bullet)[\bullet] \leadsto (v_1[\bullet]) \bowtie_R \bullet/R_1} \quad \text{E_CPUSH}$$

$$\frac{\Leftrightarrow^* b/R}{\bowtie} \quad \text{multistep reduction}$$

 $\models a \leadsto^* b/R$

 $\Gamma \vDash \phi$ ok Prop wellformedness

$$\begin{array}{l} \Gamma \vDash a : A/R \\ \Gamma \vDash b : A/R \\ \hline \Gamma \vDash A : \star/R \\ \hline \Gamma \vDash a \sim_{A/R} b \text{ ok} \end{array} \quad \text{E-Wff}$$

 $\Gamma \vDash a : A/R$ typing

$$R_{1} \leq R_{2}$$

$$\Gamma \vDash a : A/R_{1}$$

$$\Gamma \vDash a : A/R_{2}$$

$$E_SUBROLE$$

$$\frac{\vDash \Gamma}{\Gamma \vDash \star : \star/R} \quad E_STAR$$

$$\vDash \Gamma$$

$$\frac{x : A/R \in \Gamma}{\Gamma \vDash x : A/R} \quad E_VAR$$

$$\Gamma, x : A/R \vDash B : \star/R'$$

$$\Gamma \vDash A : \star/R$$

$$\frac{R \leq R'}{\Gamma \vDash \Pi^{\rho}x : A/R \to B : \star/R'} \quad E_PI$$

$$\Gamma, x : A/R \vDash a : B/R'$$

$$\Gamma \vDash A : \star/R$$

$$(\rho = +) \lor (x \not\in \text{fv } a)$$

$$R \leq R'$$

$$\Gamma \vDash \lambda^{R,\rho}x . a : (\Pi^{\rho}x : A/R \to B)/R'$$

$$E_ABS$$

$$\begin{array}{c} \Gamma \vDash b : \Pi^+x : A/R \to B/R' \\ \Gamma \vDash a : A/R \\ \hline \Gamma \vDash b : a^{R,+} : B\{a/x\}/R' \\ \hline \Gamma \vDash b : \Pi^-x : A/R \to B/R' \\ \hline \Gamma \vDash a : A/R \\ \hline \Gamma \vDash b : \Pi^-x : A/R \to B/R' \\ \hline \Gamma \vDash a : A/R \\ \hline \Gamma \vDash b : D^{R,-} : B\{a/x\}/R' \\ \hline \Gamma \vDash a : A/R \\ \hline \Gamma \vDash a : A/R \\ \hline \Gamma \vDash a : A/R \\ \hline \Gamma \vDash a : B/R \\ \hline \Gamma \vDash \phi \text{ ok} \\ \hline \Gamma \vDash \phi \text{ ok} \\ \hline \Gamma \vDash \phi \text{ ok} \\ \hline \Gamma \vDash a : b/R \\ \hline \Gamma \vDash a \Rightarrow b : A/R \\ \hline \Gamma \vDash a : b/R \\ \hline \Gamma \vDash a : b/R \\ \hline \Gamma \vDash a : b/R \\ \hline \Gamma \vDash a : A/R \\ \hline \Gamma \vDash A/R$$

 $\Gamma; \Delta \vDash \phi_1 \equiv \phi_2$

prop equality

$$\begin{array}{c} \Gamma \vDash A: \star/R \\ \Gamma; \Delta \vDash A_1 \equiv A_2: A/R \\ \Gamma; \Delta \vDash B_1 \equiv B_2: A/R \\ \hline \Gamma; \Delta \vDash A_1 \sim_{A/R} B_1 \equiv A_2 \sim_{A/R} B_2 \end{array} \quad \text{E_PropCong} \\ \Gamma; \Delta \vDash A \equiv B: \star/R \\ \Gamma \vDash A_1 \sim_{A/R} A_2 \text{ ok} \\ \hline \Gamma \vDash A_1 \sim_{A/R} A_2 \text{ ok} \\ \hline \Gamma; \Delta \vDash A_1 \sim_{A/R} A_2 \equiv A_1 \sim_{B/R} A_2 \end{array} \quad \text{E_IsoConv} \\ \frac{\Gamma; \Delta \vDash A_1 \sim_{A/R} A_2 \equiv A_1 \sim_{B/R} A_2}{\Gamma; \Delta \vDash \forall c: \phi_1.B_1 \equiv \forall c: \phi_2.B_2: \star/R} \quad \text{E_CPiFst} \\ \hline \Gamma; \Delta \vDash \phi_1 \equiv \phi_2 \end{array}$$

 $\Gamma; \Delta \vDash a \equiv b : A/R$

definitional equality

$$c: (a \sim_{A/R} b) \in \Gamma$$

$$c \in \Delta$$

$$\Gamma; \Delta \vDash a \equiv b : A/R$$

$$E_{-ASSN}$$

```
\frac{\Gamma \vDash a : A/R}{\Gamma : \Delta \vDash a \equiv a : A/R} \quad \text{E\_Refl}
                                       \Gamma; \Delta \vDash b \equiv a : A/R
                                                                                   E_Sym
                                       \Gamma; \Delta \vDash a \equiv b : A/R
                                    \Gamma; \Delta \vDash a \equiv a_1 : A/R
                                    \Gamma; \Delta \vDash a_1 \equiv b : A/R
                                                                                 E_Trans
                                    \Gamma; \Delta \vDash a \equiv b : A/R
                                      \Gamma; \Delta \vDash a \equiv b : A/R_1
                                      R_1 \leq R_2
                                                                                    E_Sub
                                     \Gamma; \Delta \vDash a \equiv b : A/R_2
                                            \Gamma \vDash a_1 : B/R
                                            \Gamma \vDash a_2 : B/R
                                            \models a_1 > a_2/R
                                                                                 E_BETA
                                    \overline{\Gamma; \Delta \vDash a_1 \equiv a_2 : B/R}
                       \Gamma; \Delta \vDash A_1 \equiv A_2 : \star / R
                       \Gamma, x: A_1/R; \Delta \vDash B_1 \equiv B_2: \star/R'
                       \Gamma \vDash A_1 : \star / R
                       \Gamma \vDash \Pi^{\rho} x : A_1/R \to B_1 : \star/R'
                       \Gamma \vDash \Pi^{\rho} x : A_2/R \to B_2 : \star/R'
                       R \leq R'
                                                                                                               E_PiCong
 \overline{\Gamma; \Delta \vDash (\Pi^{\rho}x : A_1/R \to B_1)} \equiv (\Pi^{\rho}x : A_2/R \to B_2) : \star/R'
                      \Gamma, x: A_1/R; \Delta \vDash b_1 \equiv b_2: B/R'
                      \Gamma \vDash A_1 : \star / R
                      R \leq R'
                     (\rho = +) \lor (x \not\in \mathsf{fv}\ b_1)
\frac{(\rho = +) \vee (x \not\in \mathsf{fv} \ b_2)}{\Gamma; \Delta \vDash (\lambda^{R,\rho} x. b_1) \equiv (\lambda^{R,\rho} x. b_2) : (\Pi^{\rho} x : A_1/R \to B)/R'}
                                                                                                            E_ABSCONG
               \Gamma; \Delta \vDash a_1 \equiv b_1 : (\Pi^+ x : A/R \to B)/R'
               \Gamma; \Delta \vDash a_2 \equiv b_2 : A/R
                                                                                                  E_AppCong
          \Gamma; \Delta \vDash a_1 \ a_2^{R,+} \equiv b_1 \ b_2^{R,+} : (B\{a_2/x\})/R'
              \Gamma; \Delta \vDash a_1 \equiv b_1 : (\Pi^- x : A/R \rightarrow B)/R'
              \Gamma \vDash a : A/R
                                                                                               E_IAPPCONG
           \Gamma: \Delta \vDash a_1 \square^{R,-} \equiv b_1 \square^{R,-} : (B\{a/x\})/R'
      \Gamma; \Delta \vDash \Pi^{\rho} x : A_1/R \xrightarrow{} B_1 \equiv \Pi^{\rho} x : A_2/R \xrightarrow{} B_2 : \star/R'
                                   \Gamma: \Delta \vDash A_1 \equiv A_2 : \star / R
      \Gamma; \Delta \vDash \Pi^{\rho} x : A_1/R \to B_1 \equiv \Pi^{\rho} x : A_2/R \to B_2 : \star/R'
      \Gamma; \Delta \vDash a_1 \equiv a_2 : A_1/R
                                                                                                               E_PiSnd
                    \Gamma; \Delta \vDash B_1\{a_1/x\} \equiv B_2\{a_2/x\} : \star/R'
                 \Gamma; \Delta \vDash a_1 \sim_{A_1/R} b_1 \equiv a_2 \sim_{A_2/R} b_2
                 \Gamma, c: a_1 \sim_{A_1/R} b_1; \Delta \vDash A \equiv B: \star/R'
                 \Gamma \vDash a_1 \sim_{A_1/R} b_1 ok
                 \Gamma \vDash \forall c : a_1 \sim_{A_1/R} b_1 . A : \star/R'
                 \Gamma \vDash \forall c : a_2 \sim_{A_2/R} b_2 . B : \star / R'
                                                                                                           E_CPiCong
 \overline{\Gamma;\Delta \vDash \forall c\!:\! a_1 \sim_{A_1/R} b_1.A \equiv \forall c\!:\! a_2 \sim_{A_2/R} b_2.B : \star/R'}
```

 $\models \Gamma$ context wellformedness

$$\overline{\models \varnothing} \quad \text{E-EMPTY}$$

$$\vDash \Gamma$$

$$\Gamma \vDash A : \star / R$$

$$x \not\in \text{dom } \Gamma$$

$$\vDash \Gamma, x : A / R$$

$$\vDash \Gamma$$

$$\Gamma \vDash \phi \text{ ok}$$

$$c \not\in \text{dom } \Gamma$$

$$\vDash \Gamma, c : \phi$$

$$\text{E-ConsCo}$$

 $\models \Sigma$ | signature wellformedness

 $\Gamma \vdash \phi$ ok prop wellformedness

$$\begin{split} & \Gamma \vdash a : A/R \\ & \Gamma \vdash b : B/R \\ & \frac{|A|R = |B|R}{\Gamma \vdash a \sim_{A/R} b \text{ ok}} \quad \text{AN_WFF} \end{split}$$

 $\Gamma \vdash a : A/R$ typing

$$\frac{\vdash \Gamma}{\Gamma \vdash \star : \star / R} \quad \text{An_Star}$$

$$\vdash \Gamma$$

$$\frac{x : A/R \in \Gamma}{\Gamma \vdash x : A/R} \quad \text{An_Var}$$

$$\frac{\Gamma, x : A/R \vdash B : \star / R'}{\Gamma \vdash A : \star / R} \quad \text{An_Pi}$$

$$\frac{\Gamma \vdash A : \star / R}{\Gamma \vdash \Pi^{\rho} x : A/R \rightarrow B : \star / R'} \quad \text{An_Pi}$$

$$\frac{\Gamma \vdash A : \star / R}{\Gamma \vdash \Pi^{\rho} x : A/R \rightarrow B : \star / R'} \quad \text{An_Abs}$$

$$\frac{\Gamma \vdash A : \star / R}{\Gamma \vdash \lambda^{\rho} x : A/R \vdash a : B/R'} \quad \text{An_Abs}$$

$$\frac{\Gamma \vdash \lambda \circ (\Pi^{\rho} x : A/R \rightarrow B)/R'}{\Gamma \vdash \lambda \circ (\Pi^{\rho} x : A/R \rightarrow B)/R'} \quad \text{An_App}$$

$$\frac{\Gamma \vdash b : (\Pi^{\rho} x : A/R \rightarrow B)/R'}{\Gamma \vdash a : A/R} \quad \text{An_App}$$

$$\frac{\Gamma \vdash b : A/R}{\Gamma \vdash b : a^{R,\rho} : (B\{a/x\})/R'} \quad \text{An_App}$$

$$\frac{\Gamma \vdash a : A/R}{\Gamma \vdash B : \star / R} \quad \text{An_Conv}$$

$$\frac{\Gamma \vdash \phi \text{ ok}}{\Gamma \vdash B : \star / R} \quad \text{An_Conv}$$

$$\frac{\Gamma \vdash \phi \text{ ok}}{\Gamma \vdash Ac : \phi \cdot B : \star / R} \quad \text{An_Conv}$$

$$\frac{\Gamma \vdash \phi \text{ ok}}{\Gamma \vdash Ac : \phi \cdot B : \star / R} \quad \text{An_CAbs}$$

$$\frac{\Gamma \vdash a : (\forall c : a \sim_{A_1/R} b \cdot B)/R'}{\Gamma \vdash A : (\phi c : a \sim_{A_1/R} b \cdot B)/R'} \quad \text{An_CAbs}$$

$$\frac{\Gamma \vdash a : (\forall c : a \sim_{A_1/R} b \cdot B)/R'}{\Gamma \vdash a : A/R} \quad \text{An_CApp}$$

$$\frac{\vdash \Gamma}{\Gamma \vdash A : \star / R} \quad \text{An_CApp}$$

$$\frac{\vdash \Gamma}{\Gamma \vdash A : \star / R} \quad \text{An_CApp}$$

$$\frac{\vdash \Gamma}{\Gamma \vdash A : \star / R} \quad \text{An_FAM}$$

$$\frac{R_1 \leq R_2}{\Gamma \vdash a : A/R_1} \quad \text{An_FAM}$$

$$\frac{R_1 \leq R_2}{\Gamma \vdash a : A/R_2} \quad \text{An_SubRole}$$

```
\Gamma; \Delta \vdash \gamma : \phi_1 \sim \phi_2
                                                   coercion between props
                                                                   \Gamma; \Delta \vdash \gamma_1 : A_1 \sim_R A_2
                                                                   \Gamma; \Delta \vdash \gamma_2 : B_1 \sim_R B_2
                                                                   \Gamma \vdash A_1 \sim_{A/R} B_1 ok
                                                                  \Gamma \vdash A_2 \sim_{A/R}^{\square} B_2 ok
                                                                                                                                                              An_PropCong

\overrightarrow{\Gamma; \Delta \vdash (\gamma_1 \sim_A \gamma_2) : (A_1 \sim_{A/R} B_1) \sim (A_2 \sim_{A/R} B_2)}

                                                       \frac{\Gamma; \Delta \vdash \gamma : \forall c : \phi_1. A_2 \sim_R \forall c : \phi_2. B_2}{\text{An\_CPiFst}} \quad \text{An\_CPiFst}
                                                              \Gamma; \Delta \vdash \mathbf{cpiFst} \ \gamma : \phi_1 \sim \phi_2
                                                                     \frac{\Gamma; \Delta \vdash \gamma : \phi_1 \sim \phi_2}{\Gamma; \Delta \vdash \mathbf{sym} \, \gamma : \phi_2 \sim \phi_1} \quad \text{An_IsoSym}
                                                                        \Gamma; \Delta \vdash \gamma : A \sim_R B
                                                                        \Gamma \vdash a_1 \sim_{A/R} a_2 ok
                                                                        \Gamma \vdash a_1' \sim_{B/R} a_2' ok
                                                                         |a_1|R = |a_1'|R
      \frac{|a_{2}|R = |a_{2}^{'}|R}{\Gamma; \Delta \vdash \mathbf{conv} \ (a_{1} \sim_{A/R} a_{2}) \sim_{\gamma} (a_{1}^{'} \sim_{B/R} a_{2}^{'}) : (a_{1} \sim_{A/R} a_{2}) \sim (a_{1}^{'} \sim_{B/R} a_{2}^{'})} \quad \text{An\_IsoConv}
                                                     coercion between types
\Gamma; \Delta \vdash \gamma : A \sim_R B
                                                                               c: a \sim_{A/R} b \in \Gamma
                                                                               c \in \Delta
                                                                              \frac{c \in \Delta}{\Gamma; \Delta \vdash c : a \sim_R b} \quad \text{An\_Assn}
                                                                         \frac{\Gamma \vdash a : A/R}{\Gamma ; \Delta \vdash \mathbf{refl} \; a : a \sim_R a} \quad \text{An\_Refl}
                                                                         \Gamma \vdash a : A/R
                                                                         \Gamma \vdash b : B/R
                                                                         |a|R = |b|R
                                                               \frac{\Gamma; \widetilde{\Gamma} \vdash \gamma : A \sim_R B}{\Gamma; \Delta \vdash (a \mid =\mid_{\gamma} b) : a \sim_R b} \quad \text{An\_eraseeq}
                                                                              \Gamma \vdash b : B/R
                                                                              \Gamma \vdash a : A/R
                                                                              \Gamma; \widetilde{\Gamma} \vdash \gamma_1 : B \sim_R A
                                                                          \frac{\Gamma; \Delta \vdash \gamma: b \sim_R a}{\Gamma; \Delta \vdash \mathbf{sym} \ \gamma: a \sim_R b} \quad \text{An\_Sym}
                                                                          \Gamma; \Delta \vdash \gamma_1 : a \sim_R a_1
                                                                          \Gamma; \Delta \vdash \gamma_2 : a_1 \sim_R b
                                                                          \Gamma \vdash a : A/R
                                                                          \Gamma \vdash a_1 : A_1/R
                                                                      \frac{\Gamma; \widetilde{\Gamma} \vdash \gamma_3 : A \sim_R A_1}{\Gamma; \Delta \vdash (\gamma_1; \gamma_2) : a \sim_R b} \quad \text{An\_Trans}
                                                                             \Gamma \vdash a_1 : B_0/R
                                                                             \Gamma \vdash a_2 : B_1/R
                                                                              |B_0|R = |B_1|R
                                                                   \frac{ \vdash |a_1|R > |a_2|R/R}{\Gamma; \Delta \vdash \mathbf{red} \ a_1 \ a_2 : a_1 \sim_R a_2} \quad \text{An\_Beta}
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```
\Gamma; \Delta \vdash \gamma_1 : A_1 \sim_{R'} A_2
                                            \Gamma, x: A_1/R; \Delta \vdash \gamma_2: B_1 \sim_{R'} B_2
                                            B_3 = B_2\{x \triangleright_{R'} \operatorname{sym} \gamma_1/x\}
                                            \Gamma \vdash \Pi^{\rho} x : A_1/R \rightarrow B_1 : \star/R'
                                            \Gamma \vdash \Pi^{\rho} x : A_1/R \rightarrow B_2 : \star/R'
                                            \Gamma \vdash \Pi^{\rho} x : A_2/R \rightarrow B_3 : \star/R'
                                            R \leq R'
                                                                                                                                                       An_PiCong
         \overline{\Gamma; \Delta \vdash \Pi^{R,\rho} x : \gamma_1.\gamma_2 : (\Pi^{\rho} x : A_1/R \to B_1) \sim_{R'} (\Pi^{\rho} x : A_2/R \to B_3)}
                                           \Gamma; \Delta \vdash \gamma_1 : A_1 \sim_R A_2
                                           \Gamma, x: A_1/R; \Delta \vdash \gamma_2: b_1 \sim_{R'} b_2
                                           b_3 = b_2\{x \triangleright_{R'} \operatorname{sym} \gamma_1/x\}
                                           \Gamma \vdash A_1 : \star / R
                                           \Gamma \vdash A_2 : \star / R
                                           (\rho = +) \lor (x \not\in \mathsf{fv} \mid b_1 \mid R')
                                           (\rho = +) \lor (x \not\in \mathsf{fv} \mid b_3 \mid R')
                                           \Gamma \vdash (\lambda^{\rho} x : A_1/R.b_2) : B/R'
                                           R \leq R'
                                                                                                                                              An_AbsCong
              \overline{\Gamma; \Delta \vdash (\lambda^{R,\rho}x : \gamma_1.\gamma_2) : (\lambda^{\rho}x : A_1/R.b_1) \sim_{R'} (\lambda^{\rho}x : A_2/R.b_3)}
                                                     \Gamma; \Delta \vdash \gamma_1 : a_1 \sim_{R'} b_1
                                                     \Gamma; \Delta \vdash \gamma_2 : a_2 \sim_R b_2
                                                     \Gamma \vdash a_1 \ a_2^{R,\rho} : A/R'
                                                     \Gamma \vdash b_1 \ b_2^{R,\rho} : B/R'
                                   \frac{\Gamma; \widetilde{\Gamma} \vdash \gamma_3 : A \sim_{R'} B}{\Gamma; \Delta \vdash \gamma_1 \ \gamma_2^{R,\rho} : a_1 \ a_2^{R,\rho} \sim_{R'} b_1 \ b_2^{R,\rho}} \quad \text{An\_AppCong}
                          \Gamma; \Delta \vdash \gamma : \Pi^{\rho}x : A_1/R \to B_1 \sim_{R'} \Pi^{\rho}x : A_2/R \to B_2
                                                                                                                                            An_PiFst
                                                   \Gamma; \Delta \vdash \mathbf{piFst} \gamma : A_1 \sim_R A_2
                          \Gamma; \Delta \vdash \gamma_1 : \Pi^{\rho} x : A_1/R \to B_1 \sim_{R'} \Pi^{\rho} x : A_2/R \to B_2
                          \Gamma; \Delta \vdash \gamma_2 : a_1 \sim_R a_2
                          \Gamma \vdash a_1 : A_1/R
                          \Gamma \vdash a_2 : A_2/R
                                                                                                                                            An_PiSnd
                                     \Gamma; \Delta \vdash \gamma_1@\gamma_2 : \overline{B_1\{a_1/x\} \sim_{R'} B_2\{a_2/x\}}
                                   \Gamma; \Delta \vdash \gamma_1 : a_1 \sim_{A_1/R} b_1 \sim a_2 \sim_{A_2/R} b_2
                                   \Gamma, c: a_1 \sim_{A_1/R} b_1; \Delta \vdash \gamma_3: B_1 \sim_{R'} B_2
                                   B_3 = B_2\{c \triangleright_{R'} \operatorname{\mathbf{sym}} \gamma_1/c\}
                                   \Gamma \vdash \forall c : a_1 \sim_{A_1/R} b_1 . B_1 : \star / R'
                                   \Gamma \vdash \forall c : a_2 \sim_{A_2/R} b_2.B_3 : \star/R'
                                   \Gamma \vdash \forall c : a_1 \sim_{A_1/R} b_1.B_2 : \star/R'
                                                                                                                                                      An_CPiCong
      \overline{\Gamma; \Delta \vdash (\forall c : \gamma_1.\gamma_3) : (\forall c : a_1 \sim_{A_1/R} b_1.B_1) \sim_R (\forall c : a_2 \sim_{A_2/R} b_2.B_3)}
                     \Gamma; \Delta \vdash \gamma_1 : b_0 \sim_{A_1/R} b_1 \sim b_2 \sim_{A_2/R} b_3
                     \Gamma, c: b_0 \sim_{A_1/R} b_1; \Delta \vdash \gamma_3: a_1 \sim_{R'} a_2
                      a_3 = a_2 \{c \triangleright_{R'} \operatorname{\mathbf{sym}} \gamma_1/c\}
                     \Gamma \vdash (\Lambda c : b_0 \sim_{A_1/R} b_1.a_1) : \forall c : b_0 \sim_{A_1/R} b_1.B_1/R'
                     \Gamma \vdash (\Lambda c : b_0 \sim_{A_1/R} b_1.a_2) : B/R'
                     \Gamma \vdash (\Lambda c : b_2 \sim_{A_2/R} b_3.a_3) : \forall c : b_2 \sim_{A_2/R} b_3.B_2/R'
                     \Gamma; \Gamma \vdash \gamma_4 : \forall c : b_0 \sim_{A_1/R} b_1.B_1 \sim_{R'} \forall c : \phi_2.B_2
                                                                                                                                                         An_CABSCONG
\Gamma; \Delta \vdash (\lambda c : \gamma_1.\gamma_3@\gamma_4) : (\Lambda c : b_0 \sim_{A_1/R} b_1.a_1) \sim_{R'} (\Lambda c : b_2 \sim_{A_2/R} b_3.a_3)
```

$$\begin{array}{c} \Gamma; \Delta \vdash \gamma_{1} : a_{1} \sim_{R} b_{1} \\ \Gamma; \widetilde{\Gamma} \vdash \gamma_{2} : a_{2} \sim_{R'} b_{2} \\ \Gamma; \widetilde{\Gamma} \vdash \gamma_{3} : a_{3} \sim_{R'} b_{3} \\ \Gamma \vdash a_{1}[\gamma_{2}] : A/R \\ \Gamma \vdash b_{1}[\gamma_{3}] : B/R \\ \Gamma; \widetilde{\Gamma} \vdash \gamma_{4} : A \sim_{R} B \\ \hline \Gamma; \Delta \vdash \gamma_{1}(\gamma_{2}, \gamma_{3}) : a_{1}[\gamma_{2}] \sim_{R} b_{1}[\gamma_{3}] \end{array} \quad \text{An_CAPPCong} \\ \Gamma; \Delta \vdash \gamma_{1} : (\forall c_{1} : a \sim_{A/R} a'.B_{1}) \sim_{R_{0}} (\forall c_{2} : b \sim_{B/R'} b'.B_{2}) \\ \Gamma; \widetilde{\Gamma} \vdash \gamma_{2} : a \sim_{R} a' \\ \Gamma; \widetilde{\Gamma} \vdash \gamma_{3} : b \sim_{R'} b' \\ \hline \Gamma; \Delta \vdash \gamma_{1} @ (\gamma_{2} \sim \gamma_{3}) : B_{1}\{\gamma_{2}/c_{1}\} \sim_{R_{0}} B_{2}\{\gamma_{3}/c_{2}\} \\ \hline \Gamma; \Delta \vdash \gamma_{1} : a \sim_{R_{1}} a' \\ \hline \Gamma; \Delta \vdash \gamma_{2} : a \sim_{A/R_{1}} a' \sim b \sim_{B/R_{2}} b' \\ \hline \Gamma; \Delta \vdash \gamma_{1} \rhd_{R_{2}} \gamma_{2} : b \sim_{R_{2}} b' \\ \hline \Gamma; \Delta \vdash \gamma : (a \sim_{A/R} a') \sim (b \sim_{B/R} b') \\ \hline \Gamma; \Delta \vdash \mathbf{isoSnd} \gamma : A \sim_{R} B \\ \hline \Gamma; \Delta \vdash \gamma : a \sim_{R_{1}} b \\ \hline R_{1} \leq R_{2} \\ \hline \Gamma; \Delta \vdash \mathbf{sub} \gamma : a \sim_{R_{2}} b \end{array} \quad \text{An_Sub}$$

 $\vdash \Gamma$ context wellformedness

 $\vdash \Sigma$ signature wellformedness

 $\Gamma \vdash a \leadsto b/R$ single-step, weak head reduction to values for annotated language

$$\frac{\Gamma \vdash a \leadsto a'/R_1}{\Gamma \vdash a \ b^{R,\rho} \leadsto a' \ b^{R,\rho}/R_1} \quad \text{An_Appleft}$$

$$\frac{\text{Value}_R \ (\lambda^\rho x \colon A/R.w)}{\Gamma \vdash (\lambda^\rho x \colon A/R.w) \ a^{R,\rho} \leadsto w \{a/x\}/R} \quad \text{An_Appabs}$$

$$\frac{\Gamma \vdash a \leadsto a'/R}{\Gamma \vdash a[\gamma] \leadsto a'[\gamma]/R} \quad \text{An_CAPPLEFT}$$

$$\overline{\Gamma \vdash (\Lambda c : \phi, b)[\gamma] \leadsto b\{\gamma/c\}/R} \quad \text{An_CAPPCABS}$$

$$\Gamma \vdash A : \star/R$$

$$\Gamma, x : A/R \vdash b \leadsto b'/R_1$$

$$\overline{\Gamma \vdash (\lambda^- x : A/R . b) \leadsto (\lambda^- x : A/R . b')/R_1} \quad \text{An_ABSTERM}$$

$$\frac{F \sim a : A/R \in \Sigma_1}{\Gamma \vdash F \leadsto a/R} \quad \text{An_AXIOM}$$

$$\frac{\Gamma \vdash a \leadsto a'/R}{\Gamma \vdash a \bowtie_{R_1} \gamma \leadsto a' \bowtie_{R_1} \gamma/R} \quad \text{An_CONVTERM}$$

$$\frac{\text{Value}_R \ v}{\Gamma \vdash (v \bowtie_{R_2} \gamma_1) \bowtie_{R_2} \gamma_2 \leadsto v \bowtie_{R_2} (\gamma_1; \gamma_2)/R} \quad \text{An_COMBINE}$$

$$\text{Value}_R \ v$$

$$\Gamma; \widetilde{\Gamma} \vdash \gamma : \Pi^\rho x_1 : A_1/R \to B_1 \leadsto_{R'} \Pi^\rho x_2 : A_2/R \to B_2$$

$$b' = b \bowtie_{R'} \text{sym} (\text{piFst} \gamma)$$

$$\gamma' = \gamma @ (b' \models |_{(\text{piFst} \gamma)} \ b)$$

$$\overline{\Gamma \vdash (v \bowtie_{R'} \gamma) \ b^{R,\rho} \leadsto ((v \ b'^{R,\rho}) \bowtie_{R'} \gamma')/R}} \quad \text{An_PUSH}$$

$$\text{Value}_R \ v$$

$$\Gamma; \widetilde{\Gamma} \vdash \gamma : \forall c_1 : a_1 \sim_{B_1/R} b_1 . A_1 \sim_{R'} \forall c_2 : a_2 \sim_{B_2/R} b_2 . A_2$$

$$\gamma'_1 = \gamma_1 \bowtie_{R'} \text{sym} (\text{cpiFst} \gamma)$$

$$\gamma' = \gamma @ (\gamma'_1 \sim \gamma_1)$$

$$\Gamma \vdash (v \bowtie_{R'} \gamma) [\gamma_1] \leadsto ((v [\gamma'_1]) \bowtie_{R'} \gamma')/R} \quad \text{An_CPUSH}$$

Definition rules: 163 good 0 bad Definition rule clauses: 484 good 0 bad