tmvar, x, y, f, m, n variables

covar, c coercion variables

 $\begin{array}{c} datacon,\ K\\ const,\ T\\ tyfam,\ F\\ index,\ i \end{array}$ 

index, i indices

```
Role
role, R
                                           ::=
                                                    \mathbf{Nom}
                                                    Rep
                                                    R_1 \cap R_2
                                                                                    S
relflag, \ \rho
                                                                                                          relevance flag
constraint, \phi
                                                                                                          props
                                                    a \sim_{A/R} b
                                                                                    S
S
                                                    (\phi)
                                                    \phi\{b/x\}
                                                                                    S
                                                    |\phi|
tm, a, b, v, w, A, B
                                                                                                          types and kinds
                                                    \lambda^{\rho}x:A/R.b
                                                                                    \mathsf{bind}\;x\;\mathsf{in}\;b
                                                    \lambda^{R,\rho}x.b
                                                                                    \mathsf{bind}\;x\;\mathsf{in}\;b
                                                    a b^{R,\rho}
                                                     T
                                                    \Pi^{\rho}x:A/R\to B
                                                                                    \mathsf{bind}\ x\ \mathsf{in}\ B
                                                     a \triangleright_R \gamma
                                                    \forall c : \phi.B
                                                                                    bind c in B
                                                    \Lambda c : \phi . b
                                                                                    \mathsf{bind}\ c\ \mathsf{in}\ b
                                                    \Lambda c.b
                                                                                    \mathsf{bind}\ c\ \mathsf{in}\ b
                                                     a[\gamma]
                                                    K
                                                    {f match}~a~{f with}~brs
                                                    \operatorname{\mathbf{sub}} R a
                                                                                    S
                                                     a\{b/x\}
                                                                                    S
                                                                                    S
                                                     a\{\gamma/c\}
                                                                                    S
                                                     a
                                                                                    S
                                                     (a)
                                                                                    S
                                                                                                              parsing precedence is hard
                                                                                    S
                                                    |a|R
                                                                                    S
                                                    \mathbf{Int}
                                                                                    S
                                                    Bool
                                                                                    S
                                                    Nat
                                                                                    S
                                                    Vec
                                                                                    S
                                                    0
                                                                                    S
                                                    S
                                                                                    S
                                                    True
```

```
Fix
                                                                         S
                                                                        S
                                  a \rightarrow b
                                                                        S
                                  \phi \Rightarrow A
                                  ab^{R,+}
                                                                        S
                                  \lambda^R x.a
                                                                         S
                                                                        S
                                  \lambda x : A.a
                                  \forall\,x:A/R\to B\quad \mathsf{S}
brs
                      ::=
                                                                                                       case branches
                                  none
                                  K \Rightarrow a; brs
                                                                        S
                                  brs\{a/x\}
                                                                        S
                                  brs\{\gamma/c\}
                                  (brs)
co, \gamma
                                                                                                       explicit coercions
                                  \mathbf{red} \ a \ b
                                  \mathbf{refl}\;a
                                  (a \models \mid_{\gamma} b)
                                  \operatorname{\mathbf{sym}} \gamma
                                  \gamma_1; \gamma_2
                                  \mathbf{sub}\,\gamma
                                  \Pi^{R,\rho} \dot{x} : \gamma_1.\gamma_2
                                                                        \text{bind } x \text{ in } \gamma_2
                                  \lambda^{R,\rho} x : \gamma_1 \cdot \gamma_2
\gamma_1 \ \gamma_2^{R,\rho}
                                                                        \text{bind }x\text{ in }\gamma_2
                                  \mathbf{piFst}\,\gamma
                                  \mathbf{cpiFst}\,\gamma
                                  \mathbf{isoSnd}\,\gamma
                                  \gamma_1@\gamma_2
                                  \forall c: \gamma_1.\gamma_3
                                                                        bind c in \gamma_3
                                  \lambda c: \gamma_1.\gamma_3@\gamma_4
                                                                        bind c in \gamma_3
                                  \gamma(\gamma_1,\gamma_2)
                                  \gamma@(\gamma_1 \sim \gamma_2)
                                  \gamma_1 \triangleright_R \gamma_2
                                  \gamma_1 \sim_A \gamma_2
                                  conv \phi_1 \sim_{\gamma} \phi_2
                                  \mathbf{eta}\,a
                                  left \gamma \gamma'
                                  right \gamma \gamma'
                                  (\gamma)
                                                                        S
                                  \gamma
                                  \gamma\{a/x\}
                                                                                                       binding classifier
sort
                                  \mathbf{Tm}\,A\,R
```

```
\mathbf{Co}\,\phi
sig\_sort
                                        ::=
                                                                                              signature classifier
                                                 \operatorname{\mathbf{Cs}} A
                                                 \mathbf{Ax} \ a \ A \ R
context, \ \Gamma
                                                                                              contexts
                                                 Ø
                                                 \Gamma, x : A/R
                                                 \Gamma, c: \phi
                                                 \Gamma\{b/x\}
                                                                                      Μ
                                                 \Gamma\{\gamma/c\}
                                                                                      Μ
                                                 \Gamma, \Gamma'
                                                                                      Μ
                                                 |\Gamma|
                                                                                      Μ
                                                 (\Gamma)
                                                                                      Μ
                                                                                      Μ
sig,~\Sigma
                                                                                              signatures
                                        ::=
                                                 Ø
                                                 \Sigma \cup \{\, T : A/R\}
                                                 \Sigma \cup \{F \sim a : A/R\}
                                                 \Sigma_0 \\ \Sigma_1
                                                                                      Μ
                                                                                      Μ
                                                 |\Sigma|
                                                                                      Μ
available\_props,\ \Delta
                                                 Ø
                                                 \Delta, c
                                                 \widetilde{\Gamma}
                                                                                      Μ
                                                 (\Delta)
                                                                                      Μ
role\_context, \Omega
                                                                                              role_contexts
                                                 Ø
                                                 \Omega, x:R
                                                 (\Omega)
                                                                                      Μ
                                                 \Omega
                                                                                      Μ
terminals
                                                 \leftrightarrow
                                                 \Leftrightarrow
                                                 \min
                                                 \not\in
```

```
F
                                        \neq
                                         ok
                                        Ø
                                        0
                                        fv
                                        \mathsf{dom} \\
                                        \asymp
                                        \mathbf{fst}
                                        \operatorname{snd}
                                        |\Rightarrow|
                                        \vdash_{=}
                                        \mathbf{refl_2}
                                        ++
formula, \psi
                              ::=
                                        judgement
                                        x:A/R\in\Gamma
                                        x:R\,\in\,\Omega
                                        c:\phi\,\in\,\Gamma
                                         T: A/R \, \in \, \Sigma
                                        F \sim a : A/R \in \Sigma
                                        K:T\Gamma \in \Sigma
                                        x\,\in\,\Delta
                                        c\,\in\,\Delta
                                        c \, \mathbf{not} \, \mathbf{relevant} \, \in \, \gamma
                                        x \not\in \mathsf{fv} a
                                        x \not\in \operatorname{dom} \Gamma
                                        rctx\_uniq\Omega
```

```
c \not\in \operatorname{dom} \Gamma
                              T \not\in \mathsf{dom}\, \Sigma
                              F \not\in \mathsf{dom}\, \Sigma
                              a = b
                              \phi_1 = \phi_2
                              \Gamma_1 = \Gamma_2
                              \gamma_1 = \gamma_2
                              \neg \psi
                              \psi_1 \wedge \psi_2
                              \psi_1 \vee \psi_2
                              \psi_1 \Rightarrow \psi_2
                              (\psi)
                              c:(a:A\sim b:B)\in\Gamma
                                                                       suppress lc hypothesis generated by Ott
JSubRole
                              R_1 \leq R_2
                                                                       Subroling judgement
JValue
                      ::=
                              \mathbf{CoercedValue}\,R\,A
                                                                       Values with at most one coercion at the top
                              Value_R A
                                                                       values
                              Value Type RA
                                                                       Types with head forms (erased language)
Jconsistent
                              consistent a \ b \ R
                                                                       (erased) types do not differ in their heads
Jerased
                      ::=
                             \Omega \vDash erased\_tm \; a \; R
JChk
                      ::=
                              (\rho = +) \lor (x \not\in \mathsf{fv}\ A)
                                                                       irrelevant argument check
Jpar
                             \Omega \vDash a \Rightarrow_R b
                                                                       parallel reduction (implicit language)
                             \Omega \vdash a \Rightarrow_R^* b
                                                                       multistep parallel reduction
                             \Omega \vdash a \Leftrightarrow_R b
                                                                       parallel reduction to a common term
Jbeta
                      ::=
                             \vDash a > b/R
                                                                       primitive reductions on erased terms
                             \models a \leadsto \dot{b}/R
                                                                       single-step head reduction for implicit language
                             \models a \leadsto^* b/R
                                                                       multistep reduction
Jett
                      ::=
                             \Gamma \vDash \phi ok
                                                                       Prop wellformedness
                             \Gamma \vDash a : A/R
                                                                       typing
                             \Gamma; \Delta \vDash \phi_1 \equiv \phi_2
                                                                       prop equality
```

sig

definitional equality  $context\ well formedness$ signature wellformedness prop wellformedness typing coercion between props coercion between types  $context\ well formedness$ signature wellformedness single-step, weak head reduction to values for annotated lang available\_props role\_context terminals formula

## $R_1 \leq R_2$ Subroling judgement

$$\label{eq:correction} \begin{split} \frac{\mathsf{Value}_R\ a}{\mathbf{CoercedValue}\ R\ a} & \quad \mathrm{CV} \\ & \quad \mathsf{Value}_R\ a \\ & \quad \neg (R_1 \leq R) \\ & \quad \mathbf{CoercedValue}\ R\ (a \rhd_{R_1} \gamma) \end{split} \quad \mathbf{CC} \end{split}$$

 $Value_R A$  values

$$\overline{\text{Value}_R} \star \begin{array}{c} \text{Value\_STAR} \\ \hline \\ \overline{\text{Value}_R} \ \overline{\Pi^\rho x \colon A/R_1 \to B} \end{array} \begin{array}{c} \text{Value\_PI} \\ \hline \\ \overline{\text{Value}_R} \ \overline{Vc \colon \phi.B} \end{array} \begin{array}{c} \text{Value\_CPI} \\ \hline \\ \overline{\text{Value}_R} \ \lambda^+ x \colon A/R_1.a \end{array} \begin{array}{c} \text{Value\_AbsRel} \\ \hline \\ \overline{\text{Value}_R} \ \lambda^{R_1,+} x.a \end{array} \begin{array}{c} \text{Value\_UAbsRel} \\ \hline \\ \overline{\text{Value}_R} \ \lambda^{R_1,+} x.a \end{array} \begin{array}{c} \text{Value\_UAbsIrrel} \\ \hline \\ \overline{\text{Value}_R} \ \lambda^{R_1,-} x.a \end{array} \begin{array}{c} \text{Value\_UAbsIrrel} \\ \hline \\ \overline{\text{Value}_R} \ \lambda^{R_1,-} x.a \end{array} \begin{array}{c} \text{Value\_AbsIrrel} \\ \hline \\ \overline{\text{Value}_R} \ \lambda^{-} x \colon A/R_1.a \end{array} \begin{array}{c} \text{Value\_CAbs} \\ \hline \\ \overline{\text{Value}_R} \ \Lambda c \colon \phi.a \end{array} \begin{array}{c} \text{Value\_CAbs} \\ \hline \\ \overline{\text{Value}_R} \ \Lambda c.a \end{array} \begin{array}{c} \text{Value\_UCAbs} \\ \hline \\ \overline{\text{Value}_R} \ \Lambda c.a \end{array} \begin{array}{c} \text{Value\_UCAbs} \\ \hline \\ \overline{\text{Value}_R} \ \Lambda c.a \end{array} \begin{array}{c} \text{Value\_Ax} \\ \hline \\ \overline{\text{Value}_R} \ F \end{array} \begin{array}{c} \text{Value\_Ax} \\ \hline \\ \hline \\ \hline \\ \hline \end{array} \begin{array}{c} \text{Value\_Ax} \end{array}$$

$$\overline{\text{ValueType } R \star}$$
 VALUE\_TYPE\_STAR

```
\overline{\mathbf{ValueType}\,R\,\Pi^{\rho}x\!:\!A/R_1	o B}
                                                                                VALUE_TYPE_CPI
                                          \overline{\mathbf{ValueType}\,R\,orall c\!:\!\phi.B}
                                             F \sim a : A/R_1 \in \Sigma_0
                                            \frac{\neg (R_1 \le R)}{\mathbf{ValueType} \, R \, F} \quad \text{VALUE\_TYPE\_AX}
consistent a \ b \ R
                                  (erased) types do not differ in their heads
                                                                         CONSISTENT_A_STAR
                                           \overline{\mathbf{consistent}\,\star\star R}
                \overline{\mathbf{consistent} \left( \Pi^{\rho} x_1 : A_1/R \to B_1 \right) \left( \Pi^{\rho} x_2 : A_2/R \to B_2 \right) R'}
                                                                                             CONSISTENT_A_CPI
                           \overline{\mathbf{consistent} (\forall c_1 : \phi_1.A_1) (\forall c_2 : \phi_2.A_2) R}
                                         \negValueType R b
                                                                          CONSISTENT_A_STEP_R
                                         \overline{\text{consistent } a \ b \ R}
                                         \negValueType R a
                                                                           CONSISTENT_A_STEP_L
                                          consistent a \ b \ R
\Omega \vDash erased\_tm \ a \ R
                                                 rctx\_uniq\Omega
                                                                              ERASED_A_BULLET
                                           \overline{\Omega \vDash erased\_tm \square R}
                                                   rctx\_uniq\Omega
                                                                                 ERASED_A_STAR
                                             \overline{\Omega \vDash erased\_tm \, \star \, R}
                                                    rctx\_uniq\Omega
                                                     x:R\,\in\,\Omega
                                              \frac{R \leq R_1}{\Omega \vDash erased\_tm \ x \ R_1}
                                                                                  ERASED_A_VAR
                                         \Omega, x : R_1 \vDash erased\_tm \ a \ R
                                                                                        ERASED_A_ABS
                                        \overline{\Omega \vDash erased\_tm(\lambda^{R_1,\rho}x.a)R}
                                              \Omega \vDash erased\_tm \ a \ R
                                              \Omega \vDash erased\_tm \ b \ R_1
                                                                                       ERASED_A_APP
                                         \Omega \vDash erased\_tm(a\ b^{R_1,\rho})R
                                          \Omega \vDash erased\_tm \ A \ R_1
                                          \Omega, x : R_1 \vDash erased\_tm \ B \ R
                                                                                                ERASED_A_PI
                                   \overline{\Omega \vDash erased\_tm\left(\Pi^{\rho}x : A/R_1 \to B\right)R}
                                              \Omega \vDash erased\_tm \ a \ R_1
                                              \Omega \vDash erased\_tm \ b \ R_1
                                              \Omega \vDash erased\_tm \ A \ R_1
                                              \Omega \vDash erased\_tm \ B \ R
                                                                                              ERASED_A_CPI
                                 \overline{\Omega \vDash erased\_tm\left(\forall c : a \sim_{A/R_1} b.B\right)R}
                                             \Omega \vDash erased\_tm\ b\ R
                                                                                    ERASED_A_CABS
                                         \overline{\Omega \vDash erased\_tm\left(\Lambda c.b\right)R}
                                             \Omega \vDash erased\_tm \ a \ R
                                                                                   ERASED_A_CAPP
                                          \Omega \vDash erased\_tm\left(a[\bullet]\right)R
```

$$rctx.uniq\Omega\\F \sim a: A/R \in \Sigma_0\\R \leq R_1\\\Omega \vDash erased.tm\ F\ R_1$$

$$rctx.uniq\Omega\\\Omega \vDash erased.tm\ T\ R$$

$$\Omega \vDash erased.tm\ T\ R$$

$$\Omega \vDash erased.tm\ (a \rhd_{R_1} \bullet)\ R$$

$$RASED_A.CONV$$

$$(\rho = +) \lor (x \not\in \text{fv}\ A)$$

$$retx.uniq\Omega\\\Omega \vDash erased.tm\ (a \rhd_{R_1} \bullet)\ R$$

$$RASED_A.CONV$$

$$(\rho = +) \lor (x \not\in \text{fv}\ A)$$

$$retx.uniq\Omega\\\Omega \vDash erased.tm\ a\ R$$

$$(+ = +) \lor (x \not\in \text{fv}\ A)$$

$$retx.uniq\Omega\\Rightarrow Rased.tm\ a\ R = RASED_A.CONV$$

$$RHO_REL$$

$$x \not\in \text{fv}\ A$$

$$(- = +) \lor (x \not\in \text{fv}\ A)$$

$$RHO_LRREL$$

$$x \not\in \text{fv}\ A$$

$$(- = +) \lor (x \not\in \text{fv}\ A)$$

$$RHO_LRREL$$

$$\alpha \vDash a \Rightarrow_R a$$

$$\alpha \vDash a \Rightarrow_R a$$

$$\alpha \vDash a \Rightarrow_R a$$

$$\alpha \vDash b \Rightarrow_{R_1} b'$$

$$\alpha \vDash a \Rightarrow_R a'$$

$$\alpha \vDash b \Rightarrow_{R_1} b'$$

$$\alpha \vDash a \Rightarrow_R a'$$

$$\alpha \vDash b \Rightarrow_{R_1} b'$$

$$\alpha \vDash a \Rightarrow_R a'$$

$$\alpha \vDash$$

$$\frac{\Omega \vDash a_1 \Rightarrow_{R_1} a_2}{\Omega \vDash a_1 \triangleright_{R} \bullet \Rightarrow_{R_1} a_2 \triangleright_{R} \bullet} \quad \text{Par\_Cong}$$

$$\frac{\Omega \vDash a_1 \Rightarrow_{R_1} (a_2 \triangleright_{R} \bullet)}{\Omega \vDash (a_1 \triangleright_{R} \bullet) \Rightarrow_{R_1} (a_2 \triangleright_{R} \bullet)} \quad \text{Par\_Combine}$$

$$\frac{\Omega \vDash a_1 \Rightarrow_{R_1} (a_2 \triangleright_{R} \bullet)}{\Omega \vDash b_1 \Rightarrow_{R_2} b_2} \quad \text{Par\_Push}$$

$$\frac{\Omega \vDash a_1 b_1^{R_2,+} \Rightarrow_{R_1} (a_2 (b_2 \triangleright_{R} \bullet)^{R_2,+}) \triangleright_{R} \bullet}{\Omega \vDash a_1 b_1^{R_2,+} \Rightarrow_{R_1} (a_2 (b_2 \triangleright_{R} \bullet)} \quad \text{Par\_PushCombine}$$

$$\frac{\Omega \vDash a_1 b_1^{R_2,+} \Rightarrow_{R_1} (a_2 (b_2 \triangleright_{R} \bullet)}{\Omega \vDash a_1 b_1^{R_2,+} \Rightarrow_{R_1} (a_2 (b_2 \triangleright_{R} \bullet)^{R_2,+}) \triangleright_{R} \bullet} \quad \text{Par\_PushCombine}$$

$$\frac{\Omega \vDash a_1 \Rightarrow_{R_1} (a_2 \triangleright_{R} \bullet)}{\Omega \vDash a_1 [\bullet] \Rightarrow_{R_1} (a_2 [\bullet]) \triangleright_{R} \bullet} \quad \text{Par\_CPush}$$

 $\Omega \vdash a \Rightarrow_R^* b$  multistep parallel reduction

$$\begin{array}{ll} \overline{\Omega \vdash a \Rightarrow_R^* a} & \text{MP-Refl} \\ \Omega \vdash a \Rightarrow_R b \\ \underline{\Omega \vdash b \Rightarrow_R^* a'} \\ \overline{\Omega \vdash a \Rightarrow_R^* a'} & \text{MP-Step} \end{array}$$

 $\boxed{\Omega \vdash a \Leftrightarrow_R b}$  parallel reduction to a common term

$$\frac{\Omega \vdash a_1 \Rightarrow_R^* b}{\Omega \vdash a_2 \Rightarrow_R^* b} \\
\frac{\Omega \vdash a_1 \Leftrightarrow_R a_2}{\Omega \vdash a_1 \Leftrightarrow_R a_2}$$
JOIN

 $| \vdash a > b/R |$  primitive reductions on erased terms

$$\frac{\mathsf{Value}_{R_1} \ (\lambda^{R,\rho} x.v)}{\vDash (\lambda^{R,\rho} x.v) \ b^{R,\rho} > v\{b/x\}/R_1} \quad \text{Beta\_AppAbs}$$

$$\frac{\vdash (\Lambda c.a')[\bullet] > a'\{\bullet/c\}/R}{\vdash (\Lambda c.a')[\bullet] > a'\{\bullet/c\}/R} \quad \text{Beta\_CAppCAbs}$$

$$\frac{F \sim a : A/R \in \Sigma_0}{\vDash F > a/R} \quad \text{Beta\_Axiom}$$

 $\models a \leadsto b/R$  single-step head reduction for implicit language

$$\frac{\models a \leadsto a'/R_1}{\models \lambda^{R,-}x.a \leadsto \lambda^{R,-}x.a'/R_1} \quad \text{E\_ABSTERM}$$

$$\frac{\models a \leadsto a'/R_1}{\models a \ b^{R,\rho} \leadsto a' \ b^{R,\rho}/R_1} \quad \text{E\_APPLEFT}$$

$$\frac{\models a \leadsto a'/R}{\models a[\bullet] \leadsto a'[\bullet]/R} \quad \text{E\_CAPPLEFT}$$

$$\frac{\text{Value}_{R_1} (\lambda^{R,\rho}x.v)}{\models (\lambda^{R,\rho}x.v) \ a^{R,\rho} \leadsto v\{a/x\}/R_1} \quad \text{E\_APPABS}$$

$$\frac{\models (\Lambda c.b)[\bullet] \leadsto b\{\bullet/c\}/R} \quad \text{E\_CAPPCABS}$$

$$F \sim a : A/R \in \Sigma_{0}$$

$$R \leq R_{1}$$

$$\models F \rightsquigarrow a/R_{1}$$

$$\vdash a \rightsquigarrow a'/R_{1}$$

$$\neg (R \leq R_{1})$$

$$\models a \triangleright_{R} \bullet \rightsquigarrow a' \triangleright_{R} \bullet/R_{1}$$

$$E\_CONG$$

$$\vdash (a \triangleright_{R} \bullet) \triangleright_{R} \bullet \rightsquigarrow a \triangleright_{R} \bullet/R_{1}$$

$$\vdash (v_{1} \triangleright_{R} \bullet) v_{2}^{R_{1},+} \rightsquigarrow (v_{1}(v_{2} \triangleright_{R} \bullet)^{R_{1},+}) \triangleright_{R} \bullet/R_{2}$$

$$\vdash (v_{1} \triangleright_{R} \bullet)[\bullet] \rightsquigarrow (v_{1}[\bullet]) \triangleright_{R} \bullet/R_{1}$$

 $\models a \leadsto^* b/R$  multistep reduction

 $\Gamma \vDash \phi$  ok Prop wellformedness

$$\begin{array}{l} \Gamma \vDash a : A/R \\ \Gamma \vDash b : A/R \\ \hline \Gamma \vDash A : \star/R \\ \hline \Gamma \vDash a \sim_{A/R} b \text{ ok} \end{array} \quad \text{E-Wff}$$

 $\Gamma \vDash a : A/R$  typing

$$\begin{array}{c} R_1 \leq R_2 \\ \Gamma \vDash a : A/R_1 \\ \hline \Gamma \vDash a : A/R_2 \end{array} \quad \text{E\_SubRole} \\ \\ \frac{\vDash \Gamma}{\Gamma \vDash \star : \star/R} \quad \text{E\_STAR} \\ \\ \vDash \Gamma \\ \frac{x : A/R \in \Gamma}{\Gamma \vDash \star : \star/R} \quad \text{E\_VAR} \\ \\ \Gamma, x : A/R \vDash B : \star/R' \\ \Gamma \vDash A : \star/R \\ R \leq R' \\ \hline \Gamma \vDash \Pi^{\rho}x : A/R \to B : \star/R' \quad \text{E\_PI} \\ \\ \Gamma, x : A/R \vDash a : B/R' \\ \hline \Gamma \vDash A : \star/R \\ (\rho = +) \lor (x \not\in \text{fv } a) \\ R \leq R' \\ \hline \Gamma \vDash \lambda^{R,\rho}x . a : (\Pi^{\rho}x : A/R \to B)/R' \quad \text{E\_Abs} \\ \\ \Gamma \vDash b : \Pi^{+}x : A/R \to B/R' \\ \hline \Gamma \vDash a : A/R \\ \hline \Gamma \vDash b : A^{R,+} : B\{a/x\}/R' \quad \text{E\_App} \\ \hline \end{array}$$

 $\frac{\Gamma \vDash a : A/R}{\Gamma ; \Delta \vDash a \equiv a : A/R} \quad \text{E\_Refl}$ 

```
\frac{\Gamma; \Delta \vDash b \equiv a : A/R}{\Gamma: \Delta \vDash a \equiv b : A/R} \quad \text{E\_SYM}
                                      \Gamma; \Delta \vDash a \equiv a_1 : A/R
                                      \Gamma; \Delta \vDash a_1 \equiv b : A/R
                                                                                       E_Trans
                                       \Gamma : \Delta \vDash a \equiv b : A/R
                                         \Gamma; \Delta \vDash a \equiv b : A/R_1
                                        R_1 \leq R_2
                                        \frac{R_1 \le R_2}{\Gamma; \Delta \vDash a \equiv b : A/R_2}
                                                                                          E_Sub
                                              \Gamma \vDash a_1 : B/R
                                              \Gamma \vDash a_2 : B/R
                                      \frac{\models a_1 > a_2/R}{\Gamma; \Delta \models a_1 \equiv a_2 : B/R} \quad \text{E\_Beta}
                        \Gamma; \Delta \vDash A_1 \equiv A_2 : \star / R
                        \Gamma, x: A_1/R; \Delta \vDash B_1 \equiv B_2: \star/R'
                        \Gamma \vDash A_1 : \star / R
                        \Gamma \vDash \Pi^{\rho} x : A_1/R \to B_1 : \star/R'
                        \Gamma \vDash \Pi^{\rho} x : A_2/R \rightarrow B_2 : \star/R'
                        R \leq R'
                                                                                                                       E_PiCong
\overline{\Gamma; \Delta \vDash (\Pi^{\rho}x : A_1/R \to B_1) \equiv (\Pi^{\rho}x : A_2/R \to B_2) : \star/R'}
                       \Gamma, x: A_1/R; \Delta \vDash b_1 \equiv b_2: B/R'
                       \Gamma \vDash A_1 : \star / R
                       R \leq R'
                       (\rho = +) \lor (x \not\in \mathsf{fv}\ b_1)
                       (\rho = +) \lor (x \not\in \mathsf{fv}\ b_2)
                                                                                                                    E_ABSCONG
\overline{\Gamma; \Delta \vDash (\lambda^{R,\rho} x. b_1) \equiv (\lambda^{R,\rho} x. b_2) : (\Pi^{\rho} x: A_1/R \to B)/R'}
                \Gamma; \Delta \vDash a_1 \equiv b_1 : (\Pi^+ x : A/R \to B)/R'
          \frac{\Gamma; \Delta \vDash a_2 \equiv b_2 : A/R}{\Gamma; \Delta \vDash a_1 \ a_2^{R,+} \equiv b_1 \ b_2^{R,+} : (B\{a_2/x\})/R'} \quad \text{E\_AppCong}
               \Gamma; \Delta \vDash a_1 \equiv b_1 : (\Pi^- x : A/R \rightarrow B)/R'
               \Gamma \vDash a : A/R
           \frac{1}{\Gamma; \Delta \vDash a_1 \square^{R,-} \equiv b_1 \square^{R,-} : (B\{a/x\})/R'} \quad \text{E_IAPPCong}
      \frac{\Gamma; \Delta \vDash \Pi^{\rho} x : A_1/R \to B_1 \equiv \Pi^{\rho} x : A_2/R \to B_2 : \star/R'}{\Gamma; \Delta \vDash A_1 \equiv A_2 : \star/R}
      \Gamma; \Delta \vDash \Pi^{\rho} x : A_1/R \to B_1 \equiv \Pi^{\rho} x : A_2/R \to B_2 : \star/R'
      \Gamma; \Delta \vDash a_1 \equiv a_2 : A_1/R
                      a_1 \equiv a_2 : A_1/R
\Gamma; \Delta \vDash B_1\{a_1/x\} \equiv B_2\{a_2/x\} : \star/R'
                                                                                                                       E_PiSnd
                  \Gamma; \Delta \vDash a_1 \sim_{A_1/R} b_1 \equiv a_2 \sim_{A_2/R} b_2
                  \Gamma, c: a_1 \sim_{A_1/R} b_1; \Delta \vDash A \equiv B: \star/R'
                  \Gamma \vDash a_1 \sim_{A_1/R} b_1 ok
                  \Gamma \vDash \forall c : a_1 \sim_{A_1/R} b_1.A : \star/R'
                 \Gamma \vDash \forall c : a_2 \sim_{A_2/R} b_2.B : \star/R'
\frac{\ldots \ldots \widetilde{c}}{\Gamma; \Delta \vDash \forall c \colon a_1 \sim_{A_1/R} b_1.A \equiv \forall c \colon a_2 \sim_{A_2/R} b_2.B \colon \star/R'}
                                                                                                                   E_CPICONG
                          \Gamma, c: \phi_1; \Delta \vDash a \equiv b: B/R
                          \Gamma \vDash \phi_1 ok
                                                                                                E_CABSCONG
                \overline{\Gamma; \Delta \vDash (\Lambda c.a) \equiv (\Lambda c.b) : \forall c : \phi_1.B/R}
```

$$\begin{array}{ll} \Gamma; \Delta \vDash a_1 \equiv b_1 : (\forall c \colon (a \sim_{A/R} b) . B) / R' \\ \hline \Gamma; \widetilde{\Gamma} \vDash a \equiv b : A / R \\ \hline \Gamma; \Delta \vDash a_1 [\bullet] \equiv b_1 [\bullet] : (B \{\bullet/c\}) / R' \end{array} \qquad \text{E\_CAPPCONG} \\ \hline \Gamma; \Delta \vDash \forall c \colon (a_1 \sim_{A/R} a_2) . B_1 \equiv \forall c \colon (a_1' \sim_{A'/R'} a_2') . B_2 \colon \star / R_0 \\ \Gamma; \widetilde{\Gamma} \vDash a_1 \equiv a_2 : A / R \\ \hline \Gamma; \widetilde{\Gamma} \vDash a_1' \equiv a_2' : A' / R' \\ \hline \hline \Gamma; \Delta \vDash B_1 \{\bullet/c\} \equiv B_2 \{\bullet/c\} \colon \star / R_0 \\ \hline \hline \Gamma; \Delta \vDash a \equiv b \colon A / R \\ \hline \Gamma; \Delta \vDash a \sim b \colon A / R \\ \hline \Gamma; \Delta \vDash a' \equiv b' \colon A' / R' \end{array} \qquad \text{E\_CAST} \\ \hline \hline \Gamma; \Delta \vDash a' \equiv b \colon A / R \\ \hline \Gamma; \Delta \vDash a' \equiv b' \colon A' / R' \end{array} \qquad \text{E\_CAST} \\ \hline \hline \Gamma; \Delta \vDash a \equiv b \colon A / R_1 \\ \hline \Gamma; \Delta \vDash a \equiv b \colon A / R_1 \\ \hline \Gamma; \Delta \vDash a \equiv b \colon A / R_2 \\ \hline \Gamma; \Delta \vDash a \equiv b \colon A / R_2 \\ \hline \Gamma; \Delta \vDash a \equiv b \colon A / R_2 \\ \hline \Gamma; \Delta \vDash a \equiv b \colon A / R_2 \\ \hline \Gamma; \Delta \vDash a \equiv b \colon A / R_2 \\ \hline \Gamma; \Delta \vDash a \equiv b \colon A / R_2 \end{array} \qquad \text{E\_EQCONV} \\ \hline \hline \Gamma; \Delta \vDash a \Rightarrow A / R \land A \equiv A' \colon \star / R \end{cases} \qquad \text{E\_ISOSND} \\ \hline \Gamma; \Delta \vDash a_1 \equiv a_2 \colon A / R_1 \\ \hline \Gamma; \Delta \vDash A \equiv B \colon \star / R_2 \\ \hline \neg (R_2 \le R_1) \\ \hline \Gamma \vDash B \colon \star / R_1 \\ \hline \Gamma; \Delta \vDash a_1 \rhd_{R_2} \bullet \equiv a_2 \rhd_{R_2} \bullet \colon B / R_1 \end{cases} \qquad \text{E\_CASTCONG}$$

## $\models \Gamma$ context wellformedness

## $\models \Sigma$ signature wellformedness

 $\Gamma \vdash \phi$  ok prop wellformedness

$$\begin{array}{l} \Gamma \vdash a : A/R \\ \Gamma \vdash b : B/R \\ \frac{|A|R = |B|R}{\Gamma \vdash a \sim_{A/R} b \text{ ok}} \end{array} \quad \text{An\_Wff}$$

 $\Gamma \vdash a : A/R$  typing

$$\frac{\vdash \Gamma}{\Gamma \vdash \star : \star / R} \quad \text{An\_STAR}$$

$$\vdash \Gamma$$

$$\frac{x : A/R \in \Gamma}{\Gamma \vdash x : A/R} \quad \text{An\_VAR}$$

$$\frac{\Gamma, x : A/R \vdash B : \star / R'}{\Gamma \vdash A : \star / R} \quad \text{An\_PI}$$

$$\frac{\Gamma \vdash A : \star / R}{\Gamma \vdash \Pi^{\rho} x : A/R \vdash a : B/R'} \quad \text{An\_PI}$$

$$\Gamma \vdash A : \star / R$$

$$\Gamma, x : A/R \vdash a : B/R'$$

$$(\rho = +) \lor (x \notin \text{fv} \mid a \mid R')$$

$$R \leq R'$$

$$\Gamma \vdash b : (\Pi^{\rho} x : A/R \rightarrow B)/R'$$

$$\Gamma \vdash b : (H^{\rho} x : A/R \rightarrow B)/R'$$

$$\Gamma \vdash a : A/R$$

$$\Gamma \vdash b : a^{R,\rho} : (B\{a/x\})/R'$$

$$AN\_APP$$

$$\Gamma \vdash a : A/R$$

$$\Gamma; \widetilde{\Gamma} \vdash \gamma : A \sim_R B$$

$$\Gamma \vdash B : \star / R$$

$$\Gamma \vdash B : \star / R$$

$$\Gamma \vdash A \circ k$$

$$\Gamma \vdash \varphi \circ k$$

$$\Gamma, c : \varphi \vdash B : \star / R$$

$$\Gamma \vdash \varphi \circ k$$

$$\Gamma, c : \varphi \vdash a : B/R$$

$$\Gamma \vdash \varphi \circ k$$

$$\Gamma, c : \varphi \vdash a : B/R$$

$$\Gamma \vdash Ac : \varphi. a : (\forall c : \varphi. B)/R$$

$$\Lambda \land CABS$$

$$\Gamma \vdash a_1 : (\forall c : a \sim_{A_1/R} b.B)/R'$$

$$\Gamma \vdash \alpha : (\forall c : a \sim_{A_1/R} b.B)/R'$$

$$\Gamma \vdash \alpha : (\forall c : a \sim_{A_1/R} b.B)/R'$$

$$\Gamma \vdash \alpha : A/R \in \Sigma_1$$

$$\varnothing \vdash A : \star / R$$

$$\Gamma \vdash F : A/R$$

$$R_1 \leq R_2$$

$$\Gamma \vdash a : A/R_1$$

$$\Gamma \vdash a : A/R_2$$

$$\Lambda \land SUBROLE$$

 $\Gamma; \Delta \vdash \gamma : \phi_1 \sim \phi_2$ 

coercion between props

```
\Gamma; \Delta \vdash \gamma_1 : A_1 \sim_R A_2
                                                                     \Gamma; \Delta \vdash \gamma_2 : B_1 \sim_R B_2
                                                                     \Gamma \vdash A_1 \sim_{A/R} B_1 ok
                                \frac{\Gamma \vdash A_2 \sim_{A/R} B_2 \text{ ok}}{\Gamma; \Delta \vdash (\gamma_1 \sim_A \gamma_2) : (A_1 \sim_{A/R} B_1) \sim (A_2 \sim_{A/R} B_2)} \quad \text{An\_PropCong}
                                                       \frac{\Gamma; \Delta \vdash \gamma : \forall c : \phi_1.A_2 \sim_R \forall c : \phi_2.B_2}{\Gamma; \Delta \vdash \mathbf{cpiFst} \ \gamma : \phi_1 \sim \phi_2} \quad \text{An\_CPiFst}
                                                                       \frac{\Gamma; \Delta \vdash \gamma : \phi_1 \sim \phi_2}{\Gamma; \Delta \vdash \mathbf{sym} \ \gamma : \phi_2 \sim \phi_1} \quad \text{An_IsoSym}
                                                                           \Gamma; \Delta \vdash \gamma : A \sim_R B
                                                                           \Gamma \vdash a_1 \sim_{A/R} a_2 ok
                                                                           \Gamma \vdash a_1' \sim_{B/R} a_2' ok
                                                                           |a_1|R = |a_1'|R
                                                                           |a_2|R = |a_2'|R
       \overline{\Gamma; \Delta \vdash \mathbf{conv} \ (a_1 \sim_{A/R} a_2) \sim_{\gamma} (a'_1 \sim_{B/R} a'_2) : (a_1 \sim_{A/R} a_2) \sim (a'_1 \sim_{B/R} a'_2)} \quad \text{An\_IsoConv}
\Gamma; \Delta \vdash \gamma : A \sim_R B
                                                       coercion between types
                                                                                  \vdash \Gamma
                                                                                  c: a \sim_{A/R} b \in \Gamma
                                                                                 \frac{c \in \Delta}{\Gamma; \Delta \vdash c : a \sim_R b} \quad \text{An\_Assn}
                                                                            \frac{\Gamma \vdash a : A/R}{\Gamma ; \Delta \vdash \mathbf{refl} \; a : a \sim_R a} \quad \text{An\_Refl}
                                                                           \Gamma \vdash a : A/R
                                                                           \Gamma \vdash b : B/R
                                                                           |a|R = |b|R
                                                                  \frac{\Gamma; \widetilde{\Gamma} \vdash \gamma : A \sim_R B}{\Gamma; \Delta \vdash (a \mid = \mid_{\gamma} b) : a \sim_R b} \quad \text{An\_eraseeq}
                                                                                 \Gamma \vdash b : B/R
                                                                                 \Gamma \vdash a : A/R
                                                                                 \Gamma; \widetilde{\Gamma} \vdash \gamma_1 : B \sim_R A
                                                                             \frac{\Gamma; \Delta \vdash \gamma : b \sim_R a}{\Gamma; \Delta \vdash \mathbf{sym} \ \gamma : a \sim_R b} \quad \text{An\_Sym}
                                                                             \Gamma; \Delta \vdash \gamma_1 : a \sim_R a_1
                                                                             \Gamma; \Delta \vdash \gamma_2 : a_1 \sim_R b
                                                                             \Gamma \vdash a : A/R
                                                                             \Gamma \vdash a_1 : A_1/R
                                                                        \frac{\Gamma; \widetilde{\Gamma} \vdash \gamma_3 : A \sim_R A_1}{\Gamma; \Delta \vdash (\gamma_1; \gamma_2) : a \sim_R b} \quad \text{An\_Trans}
                                                                                \Gamma \vdash a_1 : B_0/R
                                                                                \Gamma \vdash a_2 : B_1/R
                                                                                |B_0|R = |B_1|R
                                                                      \frac{ \vdash |a_1|R > |a_2|R/R}{\Gamma; \Delta \vdash \mathbf{red} \ a_1 \ a_2 : a_1 \sim_R \ a_2} \quad \text{An\_Beta}
```

```
\Gamma; \Delta \vdash \gamma_1 : A_1 \sim_{R'} A_2
                                            \Gamma, x: A_1/R; \Delta \vdash \gamma_2: B_1 \sim_{R'} B_2
                                            B_3 = B_2\{x \triangleright_{R'} \operatorname{sym} \gamma_1/x\}
                                            \Gamma \vdash \Pi^{\rho} x : A_1/R \rightarrow B_1 : \star/R'
                                            \Gamma \vdash \Pi^{\rho} x : A_1/R \rightarrow B_2 : \star/R'
                                            \Gamma \vdash \Pi^{\rho} x : A_2/R \rightarrow B_3 : \star/R'
                                            R \leq R'
                                                                                                                                                       An_PiCong
         \overline{\Gamma; \Delta \vdash \Pi^{R,\rho} x : \gamma_1.\gamma_2 : (\Pi^{\rho} x : A_1/R \to B_1) \sim_{R'} (\Pi^{\rho} x : A_2/R \to B_3)}
                                           \Gamma; \Delta \vdash \gamma_1 : A_1 \sim_R A_2
                                           \Gamma, x: A_1/R; \Delta \vdash \gamma_2: b_1 \sim_{R'} b_2
                                           b_3 = b_2\{x \triangleright_{R'} \operatorname{sym} \gamma_1/x\}
                                           \Gamma \vdash A_1 : \star / R
                                           \Gamma \vdash A_2 : \star / R
                                           (\rho = +) \lor (x \not\in \mathsf{fv} \mid b_1 \mid R')
                                           (\rho = +) \lor (x \not\in \mathsf{fv} \mid b_3 \mid R')
                                           \Gamma \vdash (\lambda^{\rho} x : A_1/R.b_2) : B/R'
                                           R \leq R'
                                                                                                                                              An_AbsCong
              \overline{\Gamma; \Delta \vdash (\lambda^{R,\rho}x : \gamma_1.\gamma_2) : (\lambda^{\rho}x : A_1/R.b_1) \sim_{R'} (\lambda^{\rho}x : A_2/R.b_3)}
                                                     \Gamma; \Delta \vdash \gamma_1 : a_1 \sim_{R'} b_1
                                                     \Gamma; \Delta \vdash \gamma_2 : a_2 \sim_R b_2
                                                     \Gamma \vdash a_1 \ a_2^{R,\rho} : A/R'
                                                     \Gamma \vdash b_1 \ b_2^{R,\rho} : B/R'
                                   \frac{\Gamma; \widetilde{\Gamma} \vdash \gamma_3 : A \sim_{R'} B}{\Gamma; \Delta \vdash \gamma_1 \ \gamma_2^{R,\rho} : a_1 \ a_2^{R,\rho} \sim_{R'} b_1 \ b_2^{R,\rho}} \quad \text{An\_AppCong}
                          \Gamma; \Delta \vdash \gamma : \Pi^{\rho}x : A_1/R \to B_1 \sim_{R'} \Pi^{\rho}x : A_2/R \to B_2
                                                                                                                                            An_PiFst
                                                   \Gamma; \Delta \vdash \mathbf{piFst} \gamma : A_1 \sim_R A_2
                          \Gamma; \Delta \vdash \gamma_1 : \Pi^{\rho} x : A_1/R \to B_1 \sim_{R'} \Pi^{\rho} x : A_2/R \to B_2
                          \Gamma; \Delta \vdash \gamma_2 : a_1 \sim_R a_2
                          \Gamma \vdash a_1 : A_1/R
                          \Gamma \vdash a_2 : A_2/R
                                                                                                                                            An_PiSnd
                                     \Gamma; \Delta \vdash \gamma_1@\gamma_2 : \overline{B_1\{a_1/x\} \sim_{R'} B_2\{a_2/x\}}
                                   \Gamma; \Delta \vdash \gamma_1 : a_1 \sim_{A_1/R} b_1 \sim a_2 \sim_{A_2/R} b_2
                                   \Gamma, c: a_1 \sim_{A_1/R} b_1; \Delta \vdash \gamma_3: B_1 \sim_{R'} B_2
                                   B_3 = B_2\{c \triangleright_{R'} \operatorname{\mathbf{sym}} \gamma_1/c\}
                                   \Gamma \vdash \forall c : a_1 \sim_{A_1/R} b_1 . B_1 : \star / R'
                                   \Gamma \vdash \forall c : a_2 \sim_{A_2/R} b_2.B_3 : \star/R'
                                   \Gamma \vdash \forall c : a_1 \sim_{A_1/R} b_1.B_2 : \star/R'
                                                                                                                                                      An_CPiCong
      \overline{\Gamma; \Delta \vdash (\forall c : \gamma_1.\gamma_3) : (\forall c : a_1 \sim_{A_1/R} b_1.B_1) \sim_R (\forall c : a_2 \sim_{A_2/R} b_2.B_3)}
                     \Gamma; \Delta \vdash \gamma_1 : b_0 \sim_{A_1/R} b_1 \sim b_2 \sim_{A_2/R} b_3
                     \Gamma, c: b_0 \sim_{A_1/R} b_1; \Delta \vdash \gamma_3: a_1 \sim_{R'} a_2
                      a_3 = a_2 \{c \triangleright_{R'} \operatorname{\mathbf{sym}} \gamma_1/c\}
                     \Gamma \vdash (\Lambda c : b_0 \sim_{A_1/R} b_1.a_1) : \forall c : b_0 \sim_{A_1/R} b_1.B_1/R'
                     \Gamma \vdash (\Lambda c : b_0 \sim_{A_1/R} b_1.a_2) : B/R'
                     \Gamma \vdash (\Lambda c : b_2 \sim_{A_2/R} b_3.a_3) : \forall c : b_2 \sim_{A_2/R} b_3.B_2/R'
                     \Gamma; \Gamma \vdash \gamma_4 : \forall c : b_0 \sim_{A_1/R} b_1.B_1 \sim_{R'} \forall c : \phi_2.B_2
                                                                                                                                                         An_CABSCONG
\Gamma; \Delta \vdash (\lambda c : \gamma_1.\gamma_3@\gamma_4) : (\Lambda c : b_0 \sim_{A_1/R} b_1.a_1) \sim_{R'} (\Lambda c : b_2 \sim_{A_2/R} b_3.a_3)
```

$$\begin{array}{c} \Gamma; \Delta \vdash \gamma_{1} : a_{1} \sim_{R} b_{1} \\ \Gamma; \widetilde{\Gamma} \vdash \gamma_{2} : a_{2} \sim_{R'} b_{2} \\ \Gamma; \widetilde{\Gamma} \vdash \gamma_{3} : a_{3} \sim_{R'} b_{3} \\ \Gamma \vdash a_{1}[\gamma_{2}] : A/R \\ \Gamma \vdash b_{1}[\gamma_{3}] : B/R \\ \Gamma; \widetilde{\Gamma} \vdash \gamma_{4} : A \sim_{R} B \\ \hline \Gamma; \Delta \vdash \gamma_{1}(\gamma_{2}, \gamma_{3}) : a_{1}[\gamma_{2}] \sim_{R} b_{1}[\gamma_{3}] \end{array} \quad \text{An\_CAPPCong} \\ \Gamma; \Delta \vdash \gamma_{1} : (\forall c_{1} : a \sim_{A/R} a'.B_{1}) \sim_{R_{0}} (\forall c_{2} : b \sim_{B/R'} b'.B_{2}) \\ \Gamma; \widetilde{\Gamma} \vdash \gamma_{2} : a \sim_{R} a' \\ \Gamma; \widetilde{\Gamma} \vdash \gamma_{3} : b \sim_{R'} b' \\ \hline \Gamma; \Delta \vdash \gamma_{1} @ (\gamma_{2} \sim \gamma_{3}) : B_{1}\{\gamma_{2}/c_{1}\} \sim_{R_{0}} B_{2}\{\gamma_{3}/c_{2}\} \\ \hline \Gamma; \Delta \vdash \gamma_{1} : a \sim_{R_{1}} a' \\ \hline \Gamma; \Delta \vdash \gamma_{2} : a \sim_{A/R_{1}} a' \sim b \sim_{B/R_{2}} b' \\ \hline \Gamma; \Delta \vdash \gamma_{1} \rhd_{R_{2}} \gamma_{2} : b \sim_{R_{2}} b' \\ \hline \Gamma; \Delta \vdash \gamma : (a \sim_{A/R} a') \sim (b \sim_{B/R} b') \\ \hline \Gamma; \Delta \vdash \mathbf{isoSnd} \gamma : A \sim_{R} B \\ \hline \Gamma; \Delta \vdash \gamma : a \sim_{R_{1}} b \\ \hline R_{1} \leq R_{2} \\ \hline \Gamma; \Delta \vdash \mathbf{sub} \gamma : a \sim_{R_{2}} b \end{array} \quad \text{An\_Sub}$$

 $\vdash \Gamma$  context wellformedness

 $\vdash \Sigma$  signature wellformedness

 $\Gamma \vdash a \leadsto b/R$  single-step, weak head reduction to values for annotated language

$$\frac{\Gamma \vdash a \leadsto a'/R_1}{\Gamma \vdash a \ b^{R,\rho} \leadsto a' \ b^{R,\rho}/R_1} \quad \text{An_Appleft}$$
 
$$\frac{\text{Value}_R \ (\lambda^\rho x \colon A/R.w)}{\Gamma \vdash (\lambda^\rho x \colon A/R.w) \ a^{R,\rho} \leadsto w\{a/x\}/R} \quad \text{An_Appabs}$$

$$\frac{\Gamma \vdash a \leadsto a'/R}{\Gamma \vdash a[\gamma] \leadsto a'[\gamma]/R} \quad \text{An\_CAPPLEFT}$$

$$\overline{\Gamma \vdash (\Lambda c : \phi, b)[\gamma] \leadsto b\{\gamma/c\}/R} \quad \text{An\_CAPPCABS}$$

$$\Gamma \vdash A : \star/R$$

$$\Gamma, x : A/R \vdash b \leadsto b'/R_1$$

$$\overline{\Gamma \vdash (\lambda^- x : A/R.b) \leadsto (\lambda^- x : A/R.b')/R_1} \quad \text{An\_ABSTERM}$$

$$\frac{F \sim a : A/R \in \Sigma_1}{\Gamma \vdash F \leadsto a/R} \quad \text{An\_AXIOM}$$

$$\frac{\Gamma \vdash a \leadsto a'/R}{\Gamma \vdash a \bowtie_{R_1} \gamma \leadsto a' \bowtie_{R_1} \gamma/R} \quad \text{An\_CONVTERM}$$

$$\frac{Value_R \ v}{\Gamma \vdash (v \bowtie_{R_2} \gamma_1) \bowtie_{R_2} \gamma_2 \leadsto v \bowtie_{R_2} (\gamma_1; \gamma_2)/R} \quad \text{An\_COMBINE}$$

$$Value_R \ v$$

$$\Gamma; \widetilde{\Gamma} \vdash \gamma : \Pi^\rho x_1 : A_1/R \to B_1 \leadsto_{R'} \Pi^\rho x_2 : A_2/R \to B_2$$

$$b' = b \bowtie_{R'} \text{sym} (\text{piFst} \gamma)$$

$$\gamma' = \gamma@(b') = |_{(\text{piFst} \gamma)} \ b)$$

$$\Gamma \vdash (v \bowtie_{R'} \gamma) \ b^{R,\rho} \leadsto ((v \ b'^{R,\rho}) \bowtie_{R'} \gamma')/R} \quad \text{An\_PUSH}$$

$$Value_R \ v$$

$$\Gamma; \widetilde{\Gamma} \vdash \gamma : \forall c_1 : a_1 \leadsto_{B_1/R} b_1.A_1 \leadsto_{R'} \forall c_2 : a_2 \leadsto_{B_2/R} b_2.A_2$$

$$\gamma_1' = \gamma_1 \bowtie_{R'} \text{sym} (\text{cpiFst} \gamma)$$

$$\gamma' = \gamma@(\gamma'_1 \leadsto_{1})$$

$$\Gamma \vdash (v \bowtie_{R'} \gamma) [\gamma_1] \leadsto ((v[\gamma'_1]) \bowtie_{R'} \gamma')/R \quad \text{An\_CPUSH}$$

Definition rules: 162 good 0 bad Definition rule clauses: 482 good 0 bad