$tmvar$, x, y, f, m, n	variables
$covar$, c	coercion variables
$datacon$, K	
$const$, T, F	
$index$, i	indices

relflag, ρ	::=			relevance flag
	|	$+$		
	|	$-$		
appflag, ν	::=			applicative flag
	|	R		
	|	ρ		
role, R	::=			Role
	|	**Nom**		
	|	**Rep**		
	|	**Phm**		
	|	$R_1 \cap R_2$	S	
	|	**param** R_1 R_2	S	
	|	$app\_role\nu$	S	
	|	(R)	S	
constraint, ϕ	::=			props
	|	$a \sim_{A/R} b$		
	|	(ϕ)	S	
	|	$\phi\{b/x\}$	S	
	|	$\lvert\phi\rvert$	S	
	|	$a \sim_R b$	S	
tm, a, b, v, w, A, B	::=			types and kinds
	|	$\star$		
	|	x		
	|	$\lambda^\rho x\!:\!A.b$	bind x in b	
	|	$\lambda^\rho x.b$	bind x in b	
	|	$a\ b^\nu$		
	|	$\Pi^\rho x\!:\!A \to B$	bind x in B	
	|	$\Lambda c\!:\!\phi.b$	bind c in b	
	|	$\Lambda c.b$	bind c in b	
	|	$a[\gamma]$		
	|	$\forall c\!:\!\phi.B$	bind c in B	
	|	$a \triangleright_R \gamma$		
	|	F		
	|	$\square$		
	|	$case_R\ a\ of\ a' \to b_1 \Vert_{\_} \to b_2$		
	|	$caseRa_1ofa_2 \to b_1 \Vert_{\_} \to b_2$		
	|	K		
	|	**match** a **with** brs		
	|	**sub** R a		
	|	$a\{b/x\}$	S	
	|	$a\{\gamma/c\}$	S	
	|	a	S	
	|	a	S	

	|	(a)	S	
	|	a	S	parsing precedence is hard
	|	$\lvert a\rvert_R$	S	
	|	**Int**	S	
	|	**Bool**	S	
	|	**Nat**	S	
	|	**Vec**	S	
	|	0	S	
	|	S	S	
	|	**True**	S	
	|	**Fix**	S	
	|	**Age**	S	
	|	$a \to b$	S	
	|	$\phi \Rightarrow A$	S	
	|	$a\ b$	S	
	|	$\lambda x.a$	S	
	|	$\lambda x\!:\!A.a$	S	
	|	$\forall\, x : A \to B$	S	
	|	**if** ϕ **then** a **else** b	S	
brs	::=			case branches
	|	**none**		
	|	$K \Rightarrow a; brs$		
	|	$brs\{a/x\}$	S	
	|	$brs\{\gamma/c\}$	S	
	|	(brs)	S	
$co,\ \gamma$	::=			explicit coercions
	|	$\bullet$		
	|	c		
	|	**red** $a\ b$		
	|	**refl** a		
	|	$(a\ \lvert{=}\rvert_{\gamma}\ b)$		
	|	**sym** γ		
	|	$\gamma_1; \gamma_2$		
	|	**sub** γ		
	|	$\Pi^{R,\rho} x\!:\!\gamma_1.\gamma_2$	bind x in γ_2	
	|	$\lambda^{R,\rho} x\!:\!\gamma_1.\gamma_2$	bind x in γ_2	
	|	$\gamma_1\ \gamma_2^{R,\rho}$		
	|	**piFst** γ		
	|	**cpiFst** γ		
	|	**isoSnd** γ		
	|	$\gamma_1 @ \gamma_2$		
	|	$\forall c\!:\!\gamma_1.\gamma_3$	bind c in γ_3	
	|	$\lambda c\!:\!\gamma_1.\gamma_3 @ \gamma_4$	bind c in γ_3	
	|	$\gamma(\gamma_1, \gamma_2)$		

	|	$\gamma @ (\gamma_1 \sim \gamma_2)$		
	|	$\gamma_1 \triangleright_R \gamma_2$		
	|	$\gamma_1 \sim_A \gamma_2$		
	|	$\mathbf{conv}\ \phi_1 \sim_\gamma \phi_2$		
	|	$\mathbf{eta}\, a$		
	|	$\mathbf{left}\, \gamma\, \gamma'$		
	|	$\mathbf{right}\, \gamma\, \gamma'$		
	|	(γ)	S	
	|	γ	S	
	|	$\gamma\{a/x\}$	S	
$\mathit{role\_context},\ \Omega$	::=			role$_c$*ontexts*
	|	$\varnothing$		
	|	$\Omega, x : R$		
	|	(Ω)	M	
	|	Ω	M	
$\mathit{roles},\ \mathit{Rs}$	::=			
	|	$\mathbf{nilR}$		
	|	R, Rs		
$\mathit{sig\_sort}$	::=			signature classifier
	|	$: A @ \mathit{Rs}$		
	|	$\sim a : A / R @ \mathit{Rs}$		
sort	::=			binding classifier
	|	$\mathbf{Tm}\, A$		
	|	$\mathbf{Co}\, \phi$		
$\mathit{context},\ \Gamma$	::=			contexts
	|	$\varnothing$		
	|	$\Gamma, x : A$		
	|	$\Gamma, c : \phi$		
	|	$\Gamma\{b/x\}$	M	
	|	$\Gamma\{\gamma/c\}$	M	
	|	Γ, Γ'	M	
	|	$\lvert\Gamma\rvert$	M	
	|	(Γ)	M	
	|	Γ	M	
$\mathit{sig},\ \Sigma$	::=			signatures
	|	$\varnothing$		
	|	$\Sigma \cup \{F\, \mathit{sig\_sort}\}$		
	|	Σ_0	M	
	|	Σ_1	M	
	|	$\lvert\Sigma\rvert$	M	

$available\_props,\ \Delta$	$::=$		
	$\|$	$\varnothing$	
	$\|$	Δ, c	
	$\|$	$\widetilde{\Gamma}$	M
	$\|$	(Δ)	M

$terminals$	$::=$	
	$\|$	$\leftrightarrow$
	$\|$	$\Leftrightarrow$
	$\|$	$\longrightarrow$
	$\|$	min
	$\|$	$\equiv$
	$\|$	$\forall$
	$\|$	$\in$
	$\|$	$\not\in$
	$\|$	$\Leftarrow$
	$\|$	$\Rightarrow$
	$\|$	$\Rightarrow^{*}$
	$\|$	$\rightarrow$
	$\|$	Λ
	$\|$	
	$\|$	$\Box$
	$\|$	$\vdash$
	$\|$	$\dashv$
	$\|$	$\models$
	$\|$	$\models$
	$\|$	$\neq$
	$\|$	$\triangleright$
	$\|$	ok
	$\|$	$\_$
	$\|$	$\leadsto$
	$\|$	$\leadsto^{*}$
	$\|$	$\leadsto$
	$\|$	$\varnothing$
	$\|$	$\circ$
	$\|$	fv
	$\|$	dom
	$\|$	$\sim$
	$\|$	$\asymp$
	$\|$	$\|$
	$\|$	$\bullet$
	$\|$	$\mathbf{fst}$
	$\|$	$\mathbf{snd}$
	$\|$	$\| \Rightarrow \|$
	$\|$	$\vdash_{=}$
	$\|$	$\mathbf{refl_2}$

	$\mid$	$++$	
formula, ψ	$::=$		
	$\mid$	*judgement*	
	$\mid$	$x : A \in \Gamma$	
	$\mid$	$x : R \in \Omega$	
	$\mid$	$c : \phi \in \Gamma$	
	$\mid$	$F\ sig\_sort \in \Sigma$	
	$\mid$	$x \in \Delta$	
	$\mid$	$c \in \Delta$	
	$\mid$	$c\ \mathbf{not\ relevant} \in \gamma$	
	$\mid$	$x \notin \mathsf{fv}\, a$	
	$\mid$	$x \notin \mathsf{dom}\, \Gamma$	
	$\mid$	$uniq(\Omega)$	
	$\mid$	$c \notin \mathsf{dom}\, \Gamma$	
	$\mid$	$T \notin \mathsf{dom}\, \Sigma$	
	$\mid$	$F \notin \mathsf{dom}\, \Sigma$	
	$\mid$	$R_1 = R_2$	
	$\mid$	$a = b$	
	$\mid$	$\phi_1 = \phi_2$	
	$\mid$	$\Gamma_1 = \Gamma_2$	
	$\mid$	$\gamma_1 = \gamma_2$	
	$\mid$	$\neg\psi$	
	$\mid$	$\psi_1 \wedge \psi_2$	
	$\mid$	$\psi_1 \vee \psi_2$	
	$\mid$	$\psi_1 \Rightarrow \psi_2$	
	$\mid$	(ψ)	
	$\mid$	ψ	
	$\mid$	$c : (a : A \sim b : B) \in \Gamma$	
	$\mid$		suppress lc hypothesis generated by Ott
	$\mid$	$\{y/x\}B = B_1$	
	$\mid$	$\{c_1/c_2\}B = B_1$	
JSubRole	$::=$		
	$\mid$	$R_1 \leq R_2$	Subroling judgement
JPath	$::=$		
	$\mid$	$\mathsf{Path}_R\ a = F @ Rs$	Type headed by constant (partial function)
JPat	$::=$		
	$\mid$	$\Gamma \vDash a : A\ \mathbf{pat}\ @Rs$	Pattern judgment
JValue	$::=$		
	$\mid$	$\mathsf{Value}_R\ A$	values
JValueType	$::=$		
	$\mid$	$\mathsf{ValueType}_R\ A$	Types with head forms (erased language)

$Jconsistent$	$::=$		
	$\mid$	$\mathsf{consistent}_R\ a b$	(erased) types do not differ in their heads
$Jroleing$	$::=$		
	$\mid$	$\Omega \vDash a : R$	
$JChk$	$::=$		
	$\mid$	$(\rho = +) \vee (x \notin \mathsf{fv}\ A)$	irrelevant argument check
$Jpar$	$::=$		
	$\mid$	$\Omega \vDash a \Rightarrow_R b$	parallel reduction (implicit language)
	$\mid$	$\Omega \vdash a \Rightarrow^*_R b$	multistep parallel reduction
	$\mid$	$\Omega \vdash a \Leftrightarrow_R b$	parallel reduction to a common term
$Jbeta$	$::=$		
	$\mid$	$\vDash a > b/R$	primitive reductions on erased terms
	$\mid$	$\vDash a \leadsto b/R$	single-step head reduction for implicit language
	$\mid$	$\vDash a \leadsto^* b/R$	multistep reduction
$Jett$	$::=$		
	$\mid$	$\Gamma \vDash \phi\ \mathsf{ok}$	Prop wellformedness
	$\mid$	$\Gamma \vDash a : A$	typing
	$\mid$	$\Gamma; \Delta \vDash \phi_1 \equiv \phi_2$	prop equality
	$\mid$	$\Gamma; \Delta \vDash a \equiv b : A/R$	definitional equality
	$\mid$	$\vDash \Gamma$	context wellformedness
$Jsig$	$::=$		
	$\mid$	$\vDash \Sigma$	signature wellformedness
$judgement$	$::=$		
	$\mid$	$JSubRole$	
	$\mid$	$JPath$	
	$\mid$	$JPat$	
	$\mid$	$JValue$	
	$\mid$	$JValueType$	
	$\mid$	$Jconsistent$	
	$\mid$	$Jroleing$	
	$\mid$	$JChk$	
	$\mid$	$Jpar$	
	$\mid$	$Jbeta$	
	$\mid$	$Jett$	
	$\mid$	$Jsig$	
$user\_syntax$	$::=$		
	$\mid$	$tmvar$	
	$\mid$	$covar$	
	$\mid$	$datacon$	

| *const*
| *index*
| *relflag*
| *appflag*
| *role*
| *constraint*
| *tm*
| *brs*
| *co*
| *role_context*
| *roles*
| *sig_sort*
| *sort*
| *context*
| *sig*
| *available_props*
| *terminals*
| *formula*

$\boxed{R_1 \leq R_2}$ Subroling judgement

$$\frac{}{\mathbf{Nom} \leq R} \quad \textsc{NomBot}$$

$$\frac{}{R \leq \mathbf{Rep}} \quad \textsc{RepTop}$$

$$\frac{}{R \leq R} \quad \textsc{Refl}$$

$$\frac{R_1 \leq R_2 \quad R_2 \leq R_3}{R_1 \leq R_3} \quad \textsc{Trans}$$

$\boxed{\mathsf{Path}_R \ a = F@Rs}$ Type headed by constant (partial function)

$$\frac{F \ : A@Rs \ \in \ \Sigma_0}{\mathsf{Path}_R \ F = F@Rs} \quad \textsc{Path\_AbsConst}$$

$$\frac{F \ \sim a : A/R_1@Rs \ \in \ \Sigma_0 \quad \neg(R_1 \leq R)}{\mathsf{Path}_R \ F = F@R, Rs} \quad \textsc{Path\_Const}$$

$$\frac{\mathsf{Path}_R \ a = F@R_1, Rs \quad app\_role \nu = R_1}{\mathsf{Path}_R \ (a \ b'^{\nu}) = F@Rs} \quad \textsc{Path\_App}$$

$$\frac{\mathsf{Path}_R \ a = F@Rs}{\mathsf{Path}_R \ (a[\bullet]) = F@Rs} \quad \textsc{Path\_CApp}$$

$\boxed{\Gamma \vDash a : A \, \mathbf{pat} \, @Rs}$ Pattern judgment

$$\frac{F \ : A@Rs \ \in \ \Sigma_0}{\varnothing \vDash F : A \, \mathbf{pat} \, @Rs} \quad \textsc{Pat\_AbsConst}$$

$$\frac{F \ \sim a : A/R_1@Rs \ \in \ \Sigma_0 \quad \neg(R_1 \leq R)}{\varnothing \vDash F : A \, \mathbf{pat} \, @R, Rs} \quad \textsc{Pat\_Const}$$

$$\frac{\begin{array}{l}\Gamma \vDash a : \Pi^\rho y{:}A_1 \to B_1\ \mathbf{pat}\ @R_1, Rs\\ \{y/x\}B = B_1\\ \mathit{app\_role}\nu = R_1\end{array}}{\Gamma, x : A_1 \vDash (a\ x^\nu) : B\ \mathbf{pat}\ @Rs}\ \textsc{Pat\_App}$$

$$\frac{\begin{array}{l}\Gamma \vDash a : \forall c_1{:}\phi.B_1\ \mathbf{pat}\ @Rs\\ \{c_1/c\}B = B_1\end{array}}{\Gamma, c : \phi \vDash (a[\bullet]) : B\ \mathbf{pat}\ @Rs}\ \textsc{Pat\_CApp}$$

$\boxed{\mathsf{Value}_R\ A}$ values

$$\frac{}{\mathsf{Value}_R\ \star}\ \textsc{Value\_Star}$$

$$\frac{}{\mathsf{Value}_R\ \Pi^\rho x{:}A \to B}\ \textsc{Value\_Pi}$$

$$\frac{}{\mathsf{Value}_R\ \forall c{:}\phi.B}\ \textsc{Value\_CPi}$$

$$\frac{}{\mathsf{Value}_R\ \lambda^+ x{:}A.a}\ \textsc{Value\_AbsRel}$$

$$\frac{}{\mathsf{Value}_R\ \lambda^+ x.a}\ \textsc{Value\_UAbsRel}$$

$$\frac{\mathsf{Value}_R\ a}{\mathsf{Value}_R\ \lambda^- x.a}\ \textsc{Value\_UAbsIrrel}$$

$$\frac{}{\mathsf{Value}_R\ \Lambda c{:}\phi.a}\ \textsc{Value\_CAbs}$$

$$\frac{}{\mathsf{Value}_R\ \Lambda c.a}\ \textsc{Value\_UCAbs}$$

$$\frac{\mathsf{Path}_R\ a = F@Rs}{\mathsf{Value}_R\ a}\ \textsc{Value\_Path}$$

$\boxed{\mathsf{ValueType}_R\ A}$ Types with head forms (erased language)

$$\frac{}{\mathsf{ValueType}_R\ \star}\ \textsc{value\_type\_Star}$$

$$\frac{}{\mathsf{ValueType}_R\ \Pi^\rho x{:}A \to B}\ \textsc{value\_type\_Pi}$$

$$\frac{}{\mathsf{ValueType}_R\ \forall c{:}\phi.B}\ \textsc{value\_type\_CPi}$$

$$\frac{\mathsf{Path}_R\ a = F@Rs}{\mathsf{ValueType}_R\ a}\ \textsc{value\_type\_Path}$$

$\boxed{\mathsf{consistent}_R\ ab}$ (erased) types do not differ in their heads

$$\frac{}{\mathsf{consistent}_R\ \star\star}\ \textsc{consistent\_a\_Star}$$

$$\frac{}{\mathsf{consistent}_{R'}\ (\Pi^\rho x_1{:}A_1 \to B_1)(\Pi^\rho x_2{:}A_2 \to B_2)}\ \textsc{consistent\_a\_Pi}$$

$$\frac{}{\mathsf{consistent}_R\ (\forall c_1{:}\phi_1.A_1)(\forall c_2{:}\phi_2.A_2)}\ \textsc{consistent\_a\_CPi}$$

$$\frac{\begin{array}{l}\mathsf{Path}_R\ a_1 = F@Rs\\ \mathsf{Path}_R\ a_2 = F@Rs\end{array}}{\mathsf{consistent}_R\ a_1\,a_2}\ \textsc{consistent\_a\_Path}$$

$$\frac{\neg\mathsf{ValueType}_R\ b}{\mathsf{consistent}_R\ ab}\quad \textsc{consistent\_a\_Step\_R}$$

$$\frac{\neg\mathsf{ValueType}_R\ a}{\mathsf{consistent}_R\ ab}\quad \textsc{consistent\_a\_Step\_L}$$

$\boxed{\Omega \vDash a : R}$

$$\frac{uniq(\Omega)}{\Omega \vDash \Box : R}\quad \textsc{role\_a\_Bullet}$$

$$\frac{uniq(\Omega)}{\Omega \vDash \star : R}\quad \textsc{role\_a\_Star}$$

$$\frac{\begin{array}{l} uniq(\Omega) \\ x : R \in \Omega \\ R \leq R_1 \end{array}}{\Omega \vDash x : R_1}\quad \textsc{role\_a\_Var}$$

$$\frac{\Omega, x : \mathbf{Nom} \vDash a : R}{\Omega \vDash (\lambda^\rho x.a) : R}\quad \textsc{role\_a\_Abs}$$

$$\frac{\begin{array}{l} \Omega \vDash a : R \\ \Omega \vDash b : app\_role \nu \end{array}}{\Omega \vDash (a\ b^\nu) : R}\quad \textsc{role\_a\_App}$$

$$\frac{\begin{array}{l} \Omega \vDash A : R \\ \Omega, x : \mathbf{Nom} \vDash B : R \end{array}}{\Omega \vDash (\Pi^\rho x\!:\!A \to B) : R}\quad \textsc{role\_a\_Pi}$$

$$\frac{\begin{array}{l} \Omega \vDash a : R_1 \\ \Omega \vDash b : R_1 \\ \Omega \vDash A : R_0 \\ \Omega \vDash B : R \end{array}}{\Omega \vDash (\forall c\!:\!a \sim_{A/R_1} b.B) : R}\quad \textsc{role\_a\_CPi}$$

$$\frac{\Omega \vDash b : R}{\Omega \vDash (\Lambda c.b) : R}\quad \textsc{role\_a\_CAbs}$$

$$\frac{\Omega \vDash a : R}{\Omega \vDash (a[\bullet]) : R}\quad \textsc{role\_a\_CApp}$$

$$\frac{\begin{array}{l} uniq(\Omega) \\ F\ : A@Rs \in \Sigma_0 \end{array}}{\Omega \vDash F : R}\quad \textsc{role\_a\_Const}$$

$$\frac{\begin{array}{l} uniq(\Omega) \\ F\ \sim a : A/R@Rs \in \Sigma_0 \end{array}}{\Omega \vDash F : R_1}\quad \textsc{role\_a\_Fam}$$

$$\frac{\begin{array}{l} F\ sig\_sort \in \Sigma_0 \\ \Omega \vDash a : R \\ \Omega \vDash b_1 : R_1 \\ \Omega \vDash b_2 : R_1 \end{array}}{\Omega \vDash (case_R\ a\ of\ F \to b_1 \| \_ \to b_2) : R_1}\quad \textsc{role\_a\_Pattern}$$

$\boxed{(\rho = +) \lor (x \notin \mathsf{fv}\ A)}$ irrelevant argument check

$$\frac{}{(+ = +) \vee (x \notin \mathsf{fv}\ A)} \quad \textsc{Rho\_Rel}$$

$$\frac{x \ \notin \ \mathsf{fv} A}{(- = +) \vee (x \notin \mathsf{fv}\ A)} \quad \textsc{Rho\_IrrRel}$$

$\boxed{\Omega \vDash a \Rightarrow_R b}$ parallel reduction (implicit language)

$$\frac{\Omega \vDash a : R}{\Omega \vDash a \Rightarrow_R a} \quad \textsc{Par\_Refl}$$

$$\frac{\begin{array}{c}\Omega \vDash a \Rightarrow_R (\lambda^\rho x.a') \\ \Omega \vDash b \Rightarrow_{\mathit{app\_role}\nu} b'\end{array}}{\Omega \vDash a\ b^\nu \Rightarrow_R a'\{b'/x\}} \quad \textsc{Par\_Beta}$$

$$\frac{\begin{array}{c}\Omega \vDash a \Rightarrow_R a' \\ \Omega \vDash b \Rightarrow_{\mathit{app\_role}\nu} b'\end{array}}{\Omega \vDash a\ b^\nu \Rightarrow_R a'\ b'^\nu} \quad \textsc{Par\_App}$$

$$\frac{\Omega \vDash a \Rightarrow_R (\Lambda c.a')}{\Omega \vDash a[\bullet] \Rightarrow_R a'\{\bullet/c\}} \quad \textsc{Par\_CBeta}$$

$$\frac{\Omega \vDash a \Rightarrow_R a'}{\Omega \vDash a[\bullet] \Rightarrow_R a'[\bullet]} \quad \textsc{Par\_CApp}$$

$$\frac{\Omega, x : \mathbf{Nom} \vDash a \Rightarrow_R a'}{\Omega \vDash \lambda^\rho x.a \Rightarrow_R \lambda^\rho x.a'} \quad \textsc{Par\_Abs}$$

$$\frac{\begin{array}{c}\Omega \vDash A \Rightarrow_R A' \\ \Omega, x : \mathbf{Nom} \vDash B \Rightarrow_R B'\end{array}}{\Omega \vDash \Pi^\rho x\!:\!A \to B \Rightarrow_R \Pi^\rho x\!:\!A' \to B'} \quad \textsc{Par\_Pi}$$

$$\frac{\Omega \vDash a \Rightarrow_R a'}{\Omega \vDash \Lambda c.a \Rightarrow_R \Lambda c.a'} \quad \textsc{Par\_CAbs}$$

$$\frac{\begin{array}{c}\Omega \vDash A \Rightarrow_{R_0} A' \\ \Omega \vDash a \Rightarrow_{R_1} a' \\ \Omega \vDash b \Rightarrow_{R_1} b' \\ \Omega \vDash B \Rightarrow_R B'\end{array}}{\Omega \vDash \forall c\!:\!a \sim_{A/R_1} b.B \Rightarrow_R \forall c\!:\!a' \sim_{A'/R_1} b'.B'} \quad \textsc{Par\_CPi}$$

$$\frac{\begin{array}{l}F \ \sim a : A/R_1 @Rs \ \in \ \Sigma_0 \\ R_1 \leq R \\ \mathit{uniq}(\Omega)\end{array}}{\Omega \vDash F \Rightarrow_R a} \quad \textsc{Par\_Axiom}$$

$$\frac{\begin{array}{l}F\ \mathit{sig\_sort} \ \in \ \Sigma_0 \\ \Omega \vDash a \Rightarrow_R a' \\ \Omega \vDash b_1 \Rightarrow_{R_0} b_1' \\ \Omega \vDash b_2 \Rightarrow_{R_0} b_2'\end{array}}{\Omega \vDash \mathit{case}_R\ a\ \mathit{of}\ F \to b_1 \|_\_ \to b_2 \Rightarrow_{R_0} \mathit{case}_R\ a'\ \mathit{of}\ F \to b_1' \|_\_ \to b_2'} \quad \textsc{Par\_Pattern}$$

$$\frac{\begin{array}{l}\Omega \vDash a \Rightarrow_R a' \\ \Omega \vDash b_1 \Rightarrow_{R_0} b_1' \\ \Omega \vDash b_2 \Rightarrow_{R_0} b_2' \\ \mathsf{Path}_R\ a' = F@Rs\end{array}}{\Omega \vDash \mathit{case}_R\ a\ \mathit{of}\ F \to b_1 \|_\_ \to b_2 \Rightarrow_{R_0} b_1'} \quad \textsc{Par\_PatternTrue}$$

$$\frac{\begin{array}{l} F\ sig\_sort\ \in\ \Sigma_0 \\ \Omega \vDash a \Rightarrow_R a' \\ \Omega \vDash b_1 \Rightarrow_{R_0} b_1' \\ \Omega \vDash b_2 \Rightarrow_{R_0} b_2' \\ \mathsf{Value}_R\ a' \\ \neg(\mathsf{Path}_R\ a' = F@Rs) \end{array}}{\Omega \vDash case_R\ a\ of\ F \to b_1 \| \_ \to b_2 \Rightarrow_{R_0} b_2'} \quad \text{PAR\_PATTERNFALSE}$$

$\boxed{\Omega \vdash a \Rightarrow^*_R b}$ multistep parallel reduction

$$\frac{}{\Omega \vdash a \Rightarrow^*_R a} \quad \text{MP\_REFL}$$

$$\frac{\begin{array}{l} \Omega \vDash a \Rightarrow_R b \\ \Omega \vdash b \Rightarrow^*_R a' \end{array}}{\Omega \vdash a \Rightarrow^*_R a'} \quad \text{MP\_STEP}$$

$\boxed{\Omega \vdash a \Leftrightarrow_R b}$ parallel reduction to a common term

$$\frac{\begin{array}{l} \Omega \vdash a_1 \Rightarrow^*_R b \\ \Omega \vdash a_2 \Rightarrow^*_R b \end{array}}{\Omega \vdash a_1 \Leftrightarrow_R a_2} \quad \text{JOIN}$$

$\boxed{\vDash a > b/R}$ primitive reductions on erased terms

$$\frac{\mathsf{Value}_{R_1}\ (\lambda^\rho x.v)}{\vDash (\lambda^\rho x.v)\ b^\nu > v\{b/x\}/R_1} \quad \text{BETA\_APPABS}$$

$$\frac{}{\vDash (\Lambda c.a')[\bullet] > a'\{\bullet/c\}/R} \quad \text{BETA\_CAPPCABS}$$

$$\frac{\begin{array}{l} F\ \sim a : A/R@Rs\ \in\ \Sigma_0 \\ R \leq R_1 \end{array}}{\vDash F > a/R_1} \quad \text{BETA\_AXIOM}$$

$$\frac{\mathsf{Path}_R\ a = F@Rs}{\vDash case_R\ a\ of\ F \to b_1 \| \_ \to b_2 > b_1/R_0} \quad \text{BETA\_PATTERNTRUE}$$

$$\frac{\begin{array}{l} F\ sig\_sort\ \in\ \Sigma_0 \\ \mathsf{Value}_R\ a \\ \neg(\mathsf{Path}_R\ a = F@Rs) \end{array}}{\vDash case_R\ a\ of\ F \to b_1 \| \_ \to b_2 > b_2/R_0} \quad \text{BETA\_PATTERNFALSE}$$

$\boxed{\vDash a \rightsquigarrow b/R}$ single-step head reduction for implicit language

$$\frac{\vDash a \rightsquigarrow a'/R_1}{\vDash \lambda^- x.a \rightsquigarrow \lambda^- x.a'/R_1} \quad \text{E\_ABSTERM}$$

$$\frac{\vDash a \rightsquigarrow a'/R_1}{\vDash a\ b^\nu \rightsquigarrow a'\ b^\nu/R_1} \quad \text{E\_APPLEFT}$$

$$\frac{\vDash a \rightsquigarrow a'/R}{\vDash a[\bullet] \rightsquigarrow a'[\bullet]/R} \quad \text{E\_CAPPLEFT}$$

$$\frac{\vDash a \rightsquigarrow a'/R}{\vDash case_R\ a\ of\ F \to b_1 \| \_ \to b_2 \rightsquigarrow case_R\ a'\ of\ F \to b_1 \| \_ \to b_2/R_0} \quad \text{E\_PATTERN}$$

$$\frac{\vDash a > b/R}{\vDash a \rightsquigarrow b/R} \quad \text{E\_PRIM}$$

$\boxed{\vDash a \leadsto^* b/R}$ multistep reduction

$$\frac{}{\vDash a \leadsto^* a/R}\ \textsc{Equal}$$

$$\frac{\begin{array}{c}\vDash a \leadsto b/R\\ \vDash b \leadsto^* a'/R\end{array}}{\vDash a \leadsto^* a'/R}\ \textsc{Step}$$

$\boxed{\Gamma \vDash \phi\ \mathsf{ok}}$ Prop wellformedness

$$\frac{\begin{array}{c}\Gamma \vDash a : A\\ \Gamma \vDash b : A\\ \Gamma \vDash A : \star\end{array}}{\Gamma \vDash a \sim_{A/R} b\ \mathsf{ok}}\ \text{E\_}\textsc{Wff}$$

$\boxed{\Gamma \vDash a : A}$ typing

$$\frac{\vDash \Gamma}{\Gamma \vDash \star : \star}\ \text{E\_}\textsc{Star}$$

$$\frac{\begin{array}{c}\vDash \Gamma\\ x : A \in \Gamma\end{array}}{\Gamma \vDash x : A}\ \text{E\_}\textsc{Var}$$

$$\frac{\begin{array}{c}\Gamma, x : A \vDash B : \star\\ \Gamma \vDash A : \star\end{array}}{\Gamma \vDash \Pi^\rho x\!:\!A \to B : \star}\ \text{E\_}\textsc{Pi}$$

$$\frac{\begin{array}{l}\Gamma, x : A \vDash a : B\\ \Gamma \vDash A : \star\\ (\rho = +) \vee (x \notin \mathsf{fv}\ a)\end{array}}{\Gamma \vDash \lambda^\rho x.a : (\Pi^\rho x\!:\!A \to B)}\ \text{E\_}\textsc{Abs}$$

$$\frac{\begin{array}{l}\Gamma \vDash b : \Pi^+ x\!:\!A \to B\\ \Gamma \vDash a : A\end{array}}{\Gamma \vDash b\ a^+ : B\{a/x\}}\ \text{E\_}\textsc{App}$$

$$\frac{\begin{array}{l}\Gamma \vDash b : \Pi^+ x\!:\!A \to B\\ \Gamma \vDash a : A\\ \mathsf{Path}_R\ a = F@Rs\end{array}}{\Gamma \vDash b\ a^R : B\{a/x\}}\ \text{E\_T}\textsc{App}$$

$$\frac{\begin{array}{l}\Gamma \vDash b : \Pi^- x\!:\!A \to B\\ \Gamma \vDash a : A\end{array}}{\Gamma \vDash b\ \square^- : B\{a/x\}}\ \text{E\_I}\textsc{App}$$

$$\frac{\begin{array}{l}\Gamma \vDash a : A\\ \Gamma; \widetilde{\Gamma} \vDash A \equiv B : \star/\mathbf{Rep}\\ \Gamma \vDash B : \star\end{array}}{\Gamma \vDash a : B}\ \text{E\_}\textsc{Conv}$$

$$\frac{\begin{array}{l}\Gamma, c : \phi \vDash B : \star\\ \Gamma \vDash \phi\ \mathsf{ok}\end{array}}{\Gamma \vDash \forall c\!:\!\phi.B : \star}\ \text{E\_CP}\textsc{i}$$

$$\frac{\begin{array}{l}\Gamma, c : \phi \vDash a : B\\ \Gamma \vDash \phi\ \mathsf{ok}\end{array}}{\Gamma \vDash \Lambda c.a : \forall c\!:\!\phi.B}\ \text{E\_CA}\textsc{bs}$$

$$\frac{\begin{array}{l}\Gamma \vDash a_1 : \forall c\!:\!(a \sim_{A/R} b).B_1 \\ \Gamma; \widetilde{\Gamma} \vDash a \equiv b : A/R\end{array}}{\Gamma \vDash a_1[\bullet] : B_1\{\bullet/c\}} \quad \text{E\_CApp}$$

$$\frac{\begin{array}{l}\vDash \Gamma \\ F\ : A@Rs\ \in\ \Sigma_0 \\ \varnothing \vDash A : \star\end{array}}{\Gamma \vDash F : A} \quad \text{E\_Const}$$

$$\frac{\begin{array}{l}\vDash \Gamma \\ F\ \sim a : A/R_1@Rs\ \in\ \Sigma_0 \\ \varnothing \vDash A : \star\end{array}}{\Gamma \vDash F : A} \quad \text{E\_Fam}$$

$$\frac{\begin{array}{l}F\ sig\_sort\ \in\ \Sigma_0 \\ \Gamma \vDash a : A \\ \Gamma \vDash b_1 : B \\ \Gamma \vDash b_2 : B\end{array}}{\Gamma \vDash case_R\ a\ of\ F \to b_1\|_- \to b_2 : B} \quad \text{E\_Pat}$$

$$\frac{\begin{array}{l}\Gamma \vDash a_1 : A \\ \Gamma' \vDash a_2 : A\,\mathbf{pat}\,@R, Rs \\ \Gamma, (\Gamma', c : \phi_1) \vDash b_1 : B \\ \Gamma \vDash b_2 : B \\ \phi_1 = (a_1 \sim_{A/R} a_2)\end{array}}{\Gamma \vDash caseRa_1 of a_2 \to b_1\|_- \to b_2 : B} \quad \text{E\_case}$$

$\boxed{\Gamma; \Delta \vDash \phi_1 \equiv \phi_2}$ prop equality

$$\frac{\begin{array}{l}\Gamma; \Delta \vDash A_1 \equiv A_2 : A/R \\ \Gamma; \Delta \vDash B_1 \equiv B_2 : A/R\end{array}}{\Gamma; \Delta \vDash A_1 \sim_{A/R} B_1 \equiv A_2 \sim_{A/R} B_2} \quad \text{E\_PropCong}$$

$$\frac{\begin{array}{l}\Gamma; \Delta \vDash A \equiv B : \star/R_0 \\ \Gamma \vDash A_1 \sim_{A/R} A_2\ \mathsf{ok} \\ \Gamma \vDash A_1 \sim_{B/R} A_2\ \mathsf{ok}\end{array}}{\Gamma; \Delta \vDash A_1 \sim_{A/R} A_2 \equiv A_1 \sim_{B/R} A_2} \quad \text{E\_IsoConv}$$

$$\frac{\Gamma; \Delta \vDash \forall c\!:\!(a_1 \sim_{A/R_1} a_2).B_1 \equiv \forall c\!:\!(b_1 \sim_{B/R_2} b_2).B_2 : \star/R'}{\Gamma; \Delta \vDash a_1 \sim_{A/R_1} a_2 \equiv b_1 \sim_{B/R_2} b_2} \quad \text{E\_CPiFst}$$

$\boxed{\Gamma; \Delta \vDash a \equiv b : A/R}$ definitional equality

$$\frac{\begin{array}{l}\vDash \Gamma \\ c : (a \sim_{A/R} b)\ \in\ \Gamma \\ c\ \in\ \Delta\end{array}}{\Gamma; \Delta \vDash a \equiv b : A/R} \quad \text{E\_Assn}$$

$$\frac{\Gamma \vDash a : A}{\Gamma; \Delta \vDash a \equiv a : A/\mathbf{Nom}} \quad \text{E\_Refl}$$

$$\frac{\Gamma; \Delta \vDash b \equiv a : A/R}{\Gamma; \Delta \vDash a \equiv b : A/R} \quad \text{E\_Sym}$$

$$\frac{\begin{array}{l}\Gamma; \Delta \vDash a \equiv a_1 : A/R \\ \Gamma; \Delta \vDash a_1 \equiv b : A/R\end{array}}{\Gamma; \Delta \vDash a \equiv b : A/R} \quad \text{E\_Trans}$$

$$\frac{\begin{array}{l}\Gamma;\Delta \vDash a \equiv b : A/R_1\\ R_1 \leq R_2\end{array}}{\Gamma;\Delta \vDash a \equiv b : A/R_2}\ \textsc{E\_Sub}$$

$$\frac{\begin{array}{l}\Gamma \vDash a_1 : B\\ \Gamma \vDash a_2 : B\\ \vDash a_1 > a_2/R\end{array}}{\Gamma;\Delta \vDash a_1 \equiv a_2 : B/R}\ \textsc{E\_Beta}$$

$$\frac{\begin{array}{l}\Gamma;\Delta \vDash A_1 \equiv A_2 : \star/R'\\ \Gamma, x : A_1;\Delta \vDash B_1 \equiv B_2 : \star/R'\\ \Gamma \vDash A_1 : \star\\ \Gamma \vDash \Pi^\rho x\!:\!A_1 \to B_1 : \star\\ \Gamma \vDash \Pi^\rho x\!:\!A_2 \to B_2 : \star\end{array}}{\Gamma;\Delta \vDash (\Pi^\rho x\!:\!A_1 \to B_1) \equiv (\Pi^\rho x\!:\!A_2 \to B_2) : \star/R'}\ \textsc{E\_PiCong}$$

$$\frac{\begin{array}{l}\Gamma, x : A_1;\Delta \vDash b_1 \equiv b_2 : B/R'\\ \Gamma \vDash A_1 : \star\\ (\rho = +) \lor (x \notin \mathsf{fv}\ b_1)\\ (\rho = +) \lor (x \notin \mathsf{fv}\ b_2)\end{array}}{\Gamma;\Delta \vDash (\lambda^\rho x.b_1) \equiv (\lambda^\rho x.b_2) : (\Pi^\rho x\!:\!A_1 \to B)/R'}\ \textsc{E\_AbsCong}$$

$$\frac{\begin{array}{l}\Gamma;\Delta \vDash a_1 \equiv b_1 : (\Pi^+ x\!:\!A \to B)/R'\\ \Gamma;\Delta \vDash a_2 \equiv b_2 : A/R'\end{array}}{\Gamma;\Delta \vDash a_1\ {a_2}^+ \equiv b_1\ {b_2}^+ : (B\{a_2/x\})/R'}\ \textsc{E\_AppCong}$$

$$\frac{\begin{array}{l}\Gamma;\Delta \vDash a_1 \equiv b_1 : (\Pi^- x\!:\!A \to B)/R'\\ \Gamma \vDash a : A\end{array}}{\Gamma;\Delta \vDash a_1\ \square^- \equiv b_1\ \square^- : (B\{a/x\})/R'}\ \textsc{E\_IAppCong}$$

$$\frac{\Gamma;\Delta \vDash \Pi^\rho x\!:\!A_1 \to B_1 \equiv \Pi^\rho x\!:\!A_2 \to B_2 : \star/R'}{\Gamma;\Delta \vDash A_1 \equiv A_2 : \star/R'}\ \textsc{E\_PiFst}$$

$$\frac{\begin{array}{l}\Gamma;\Delta \vDash \Pi^\rho x\!:\!A_1 \to B_1 \equiv \Pi^\rho x\!:\!A_2 \to B_2 : \star/R'\\ \Gamma;\Delta \vDash a_1 \equiv a_2 : A_1/R'\end{array}}{\Gamma;\Delta \vDash B_1\{a_1/x\} \equiv B_2\{a_2/x\} : \star/R'}\ \textsc{E\_PiSnd}$$

$$\frac{\begin{array}{l}\Gamma;\Delta \vDash a_1 \sim_{A_1/R} b_1 \equiv a_2 \sim_{A_2/R} b_2\\ \Gamma, c : a_1 \sim_{A_1/R} b_1;\Delta \vDash A \equiv B : \star/R'\\ \Gamma \vDash a_1 \sim_{A_1/R} b_1\ \mathsf{ok}\\ \Gamma \vDash \forall c\!:\!a_1 \sim_{A_1/R} b_1.A : \star\\ \Gamma \vDash \forall c\!:\!a_2 \sim_{A_2/R} b_2.B : \star\end{array}}{\Gamma;\Delta \vDash \forall c\!:\!a_1 \sim_{A_1/R} b_1.A \equiv \forall c\!:\!a_2 \sim_{A_2/R} b_2.B : \star/R'}\ \textsc{E\_CPiCong}$$

$$\frac{\begin{array}{l}\Gamma, c : \phi_1;\Delta \vDash a \equiv b : B/R\\ \Gamma \vDash \phi_1\ \mathsf{ok}\end{array}}{\Gamma;\Delta \vDash (\Lambda c.a) \equiv (\Lambda c.b) : \forall c\!:\!\phi_1.B/R}\ \textsc{E\_CAbsCong}$$

$$\frac{\begin{array}{l}\Gamma;\Delta \vDash a_1 \equiv b_1 : (\forall c\!:\!(a \sim_{A/R} b).B)/R'\\ \Gamma;\widetilde{\Gamma} \vDash a \equiv b : A/\mathbf{param}\, R\, R'\end{array}}{\Gamma;\Delta \vDash a_1[\bullet] \equiv b_1[\bullet] : (B\{\bullet/c\})/R'}\ \textsc{E\_CAppCong}$$

$$\frac{\begin{array}{l}\Gamma;\Delta \vDash \forall c\!:\!(a_1 \sim_{A/R} a_2).B_1 \equiv \forall c\!:\!(a_1' \sim_{A'/R'} a_2').B_2 : \star/R_0\\ \Gamma;\widetilde{\Gamma} \vDash a_1 \equiv a_2 : A/\mathbf{param}\, R\, R_0\\ \Gamma;\widetilde{\Gamma} \vDash a_1' \equiv a_2' : A'/\mathbf{param}\, R'\, R_0\end{array}}{\Gamma;\Delta \vDash B_1\{\bullet/c\} \equiv B_2\{\bullet/c\} : \star/R_0}\ \textsc{E\_CPiSnd}$$

$$\frac{\begin{array}{l}\Gamma;\Delta \vDash a \equiv b : A/R\\ \Gamma;\Delta \vDash a \sim_{A/R} b \equiv a' \sim_{A'/R'} b'\end{array}}{\Gamma;\Delta \vDash a' \equiv b' : A'/R'} \quad \text{E\_CAST}$$

$$\frac{\begin{array}{l}\Gamma;\Delta \vDash a \equiv b : A/R\\ \Gamma;\widetilde{\Gamma} \vDash A \equiv B : \star/\mathbf{Rep}\\ \Gamma \vDash B : \star\end{array}}{\Gamma;\Delta \vDash a \equiv b : B/R} \quad \text{E\_EQCONV}$$

$$\frac{\Gamma;\Delta \vDash a \sim_{A/R_1} b \equiv a' \sim_{A'/R_1} b'}{\Gamma;\Delta \vDash A \equiv A' : \star/\mathbf{Rep}} \quad \text{E\_ISOSND}$$

$$\frac{\begin{array}{l}F\ sig\_sort\ \in\ \Sigma_0\\ \Gamma;\Delta \vDash a \equiv a' : A/R\\ \Gamma;\Delta \vDash b_1 \equiv b_1' : B/R_0\\ \Gamma;\Delta \vDash b_2 \equiv b_2' : B/R_0\end{array}}{\Gamma;\Delta \vDash case_R\ a\ of\ F \to b_1 \|\_ \to b_2 \equiv case_R\ a'\ of\ F \to b_1' \|\_ \to b_2' : B/R_0} \quad \text{E\_PATCONG}$$

$$\frac{\begin{array}{l}\mathsf{Path}_{R'}\ a = F@R, Rs\\ \mathsf{Path}_{R'}\ a' = F@R, Rs\\ \Gamma \vDash a : \Pi^+ x\!:\!A \to B\\ \Gamma \vDash b : A\\ \Gamma \vDash a' : \Pi^+ x\!:\!A \to B\\ \Gamma \vDash b' : A\\ \Gamma;\Delta \vDash a\ b^{R_1} \equiv a'\ b'^{R_1} : B\{b/x\}/R'\\ \Gamma;\widetilde{\Gamma} \vDash B\{b/x\} \equiv B\{b'/x\} : \star/R'\end{array}}{\Gamma;\Delta \vDash a \equiv a' : \Pi^+ x\!:\!A \to B/R'} \quad \text{E\_LEFTREL}$$

$$\frac{\begin{array}{l}\mathsf{Path}_{R'}\ a = F@R, Rs\\ \mathsf{Path}_{R'}\ a' = F@R, Rs\\ \Gamma \vDash a : \Pi^- x\!:\!A \to B\\ \Gamma \vDash b : A\\ \Gamma \vDash a' : \Pi^- x\!:\!A \to B\\ \Gamma \vDash b' : A\\ \Gamma;\Delta \vDash a\ \Box^- \equiv a'\ \Box^- : B\{b/x\}/R'\\ \Gamma;\widetilde{\Gamma} \vDash B\{b/x\} \equiv B\{b'/x\} : \star/R_0\end{array}}{\Gamma;\Delta \vDash a \equiv a' : \Pi^- x\!:\!A \to B/R'} \quad \text{E\_LEFTIRREL}$$

$$\frac{\begin{array}{l}\mathsf{Path}_{R'}\ a = F@R, Rs\\ \mathsf{Path}_{R'}\ a' = F@R, Rs\\ \Gamma \vDash a : \Pi^+ x\!:\!A \to B\\ \Gamma \vDash b : A\\ \Gamma \vDash a' : \Pi^+ x\!:\!A \to B\\ \Gamma \vDash b' : A\\ \Gamma;\Delta \vDash a\ b^{+} \equiv a'\ b'^{+} : B\{b/x\}/R'\\ \Gamma;\widetilde{\Gamma} \vDash B\{b/x\} \equiv B\{b'/x\} : \star/R_0\end{array}}{\Gamma;\Delta \vDash b \equiv b' : A/\mathbf{param}\, R_1\ R'} \quad \text{E\_RIGHT}$$

$$\frac{\begin{array}{l}\mathsf{Path}_{R'}\ a = F@R, Rs\\ \mathsf{Path}_{R'}\ a' = F@R, Rs\\ \Gamma \vDash a : \forall c\!:\!(a_1 \sim_{A/R_1} a_2).B\\ \Gamma \vDash a' : \forall c\!:\!(a_1 \sim_{A/R_1} a_2).B\\ \Gamma;\widetilde{\Gamma} \vDash a_1 \equiv a_2 : A/R'\\ \Gamma;\Delta \vDash a[\bullet] \equiv a'[\bullet] : B\{\bullet/c\}/R'\end{array}}{\Gamma;\Delta \vDash a \equiv a' : \forall c\!:\!(a_1 \sim_{A/R_1} a_2).B/R'} \quad \text{E\_CLEFT}$$

$\boxed{\vDash \Gamma}$ context wellformedness

$$\frac{}{\vDash \varnothing} \quad \textsc{E\_Empty}$$

$$\frac{\begin{array}{l} \vDash \Gamma \\ \Gamma \vDash A : \star \\ x \notin \mathsf{dom}\,\Gamma \end{array}}{\vDash \Gamma, x : A} \quad \textsc{E\_ConsTm}$$

$$\frac{\begin{array}{l} \vDash \Gamma \\ \Gamma \vDash \phi \;\; \mathsf{ok} \\ c \notin \mathsf{dom}\,\Gamma \end{array}}{\vDash \Gamma, c : \phi} \quad \textsc{E\_ConsCo}$$

$\boxed{\vDash \Sigma}$ signature wellformedness

$$\frac{}{\vDash \varnothing} \quad \textsc{Sig\_Empty}$$

$$\frac{\begin{array}{l} \vDash \Sigma \\ \varnothing \vDash A : \star \\ F \notin \mathsf{dom}\,\Sigma \end{array}}{\vDash \Sigma \cup \{F : A@Rs\}} \quad \textsc{Sig\_ConsConst}$$

$$\frac{\begin{array}{l} \vDash \Sigma \\ \varnothing \vDash a : A \\ F \notin \mathsf{dom}\,\Sigma \end{array}}{\vDash \Sigma \cup \{F \sim a : A/R@Rs\}} \quad \textsc{Sig\_ConsAx}$$

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Definition rules:        122 good    0 bad
Definition rule clauses: 361 good    0 bad
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