$tmvar$, x, y, f, m, n	variables
$covar$, c	coercion variables
$datacon$, K	
$const$, T, F	
$index$, i	indices

relflag, ρ	$::=$			relevance flag
	$\mid$	$+$		
	$\mid$	$-$		
appflag, ν	$::=$			applicative flag
	$\mid$	R		
	$\mid$	ρ		
role, R	$::=$			Role
	$\mid$	**Nom**		
	$\mid$	**Rep**		
	$\mid$	**Phm**		
	$\mid$	$R_1 \cap R_2$	S	
	$\mid$	$\mathbf{param}\, R_1\; R_2$	S	
	$\mid$	$app\_role\nu$	S	
	$\mid$	(R)	S	
constraint, ϕ	$::=$			props
	$\mid$	$a \sim_{A/R} b$		
	$\mid$	(ϕ)	S	
	$\mid$	$\phi\{b/x\}$	S	
	$\mid$	$\lvert\phi\rvert$	S	
	$\mid$	$a \sim_R b$	S	
tm, a, b, v, w, A, B	$::=$			types and kinds
	$\mid$	$\star$		
	$\mid$	x		
	$\mid$	$\lambda^{\rho} x : A.b$	bind x in b	
	$\mid$	$\lambda^{\rho} x.b$	bind x in b	
	$\mid$	$a\ b^{\nu}$		
	$\mid$	$\Pi^{\rho} x : A \to B$	bind x in B	
	$\mid$	$\Lambda c : \phi.b$	bind c in b	
	$\mid$	$\Lambda c.b$	bind c in b	
	$\mid$	$a[\gamma]$		
	$\mid$	$\forall c : \phi.B$	bind c in B	
	$\mid$	$a \triangleright_R \gamma$		
	$\mid$	F		
	$\mid$	$\square$		
	$\mid$	$caseRa_1 of\, a_2 \to b_1 \Vert_{\_} \to b_2$		
	$\mid$	K		
	$\mid$	$\mathbf{match}\, a\, \mathbf{with}\, brs$		
	$\mid$	$\mathbf{sub}\, R\; a$		
	$\mid$	$a\{b/x\}$	S	
	$\mid$	$a\{\gamma/c\}$	S	
	$\mid$	a	S	
	$\mid$	a	S	
	$\mid$	(a)	S	

	|	a	S	parsing precedence is hard
	|	$\|a\|_R$	S	
	|	**Int**	S	
	|	**Bool**	S	
	|	**Nat**	S	
	|	**Vec**	S	
	|	0	S	
	|	S	S	
	|	**True**	S	
	|	**Fix**	S	
	|	**Age**	S	
	|	$a \to b$	S	
	|	$\phi \Rightarrow A$	S	
	|	$a\ b$	S	
	|	$\lambda x.a$	S	
	|	$\lambda x : A.a$	S	
	|	$\forall\, x : A \to B$	S	
	|	**if** ϕ **then** a **else** b	S	
brs	::=			case branches
	|	**none**		
	|	$K \Rightarrow a; brs$		
	|	$brs\{a/x\}$	S	
	|	$brs\{\gamma/c\}$	S	
	|	(brs)	S	
$co,\ \gamma$	::=			explicit coercions
	|	$\bullet$		
	|	c		
	|	**red** $a\ b$		
	|	**refl** a		
	|	$(a \mid=\mid_{\gamma}\ b)$		
	|	**sym** γ		
	|	$\gamma_1; \gamma_2$		
	|	**sub** γ		
	|	$\Pi^{R,\rho} x : \gamma_1.\gamma_2$	bind x in γ_2	
	|	$\lambda^{R,\rho} x : \gamma_1.\gamma_2$	bind x in γ_2	
	|	$\gamma_1\ \gamma_2^{R,\rho}$		
	|	**piFst** γ		
	|	**cpiFst** γ		
	|	**isoSnd** γ		
	|	$\gamma_1 @ \gamma_2$		
	|	$\forall c : \gamma_1.\gamma_3$	bind c in γ_3	
	|	$\lambda c : \gamma_1.\gamma_3 @ \gamma_4$	bind c in γ_3	
	|	$\gamma(\gamma_1, \gamma_2)$		
	|	$\gamma @ (\gamma_1 \sim \gamma_2)$		

	$\|$	$\gamma_1 \triangleright_R \gamma_2$		
	$\|$	$\gamma_1 \sim_A \gamma_2$		
	$\|$	$\mathbf{conv}\ \phi_1 \sim_\gamma \phi_2$		
	$\|$	$\mathbf{eta}\, a$		
	$\|$	$\mathbf{left}\, \gamma\, \gamma'$		
	$\|$	$\mathbf{right}\, \gamma\, \gamma'$		
	$\|$	(γ)	S	
	$\|$	γ	S	
	$\|$	$\gamma\{a/x\}$	S	
$role\_context,\ \Omega$	::=			role$_c ontexts$
	$\|$	$\varnothing$		
	$\|$	$\Omega, x : R$		
	$\|$	(Ω)	M	
	$\|$	Ω	M	
$roles,\ Rs$	::=			
	$\|$	$\mathbf{nilR}$		
	$\|$	R, Rs		
sig_sort	::=			signature classifier
	$\|$	$: A@Rs$		
	$\|$	$\sim a : A/R@Rs$		
$sort$	::=			binding classifier
	$\|$	$\mathbf{Tm}\, A$		
	$\|$	$\mathbf{Co}\, \phi$		
$context,\ \Gamma$	::=			contexts
	$\|$	$\varnothing$		
	$\|$	$\Gamma, x : A$		
	$\|$	$\Gamma, c : \phi$		
	$\|$	$\Gamma\{b/x\}$	M	
	$\|$	$\Gamma\{\gamma/c\}$	M	
	$\|$	Γ, Γ'	M	
	$\|$	$\|\Gamma\|$	M	
	$\|$	(Γ)	M	
	$\|$	Γ	M	
$sig,\ \Sigma$	::=			signatures
	$\|$	$\varnothing$		
	$\|$	$\Sigma \cup \{F\, sig\_sort\}$		
	$\|$	Σ_0	M	
	$\|$	Σ_1	M	
	$\|$	$\|\Sigma\|$	M	
$available\_props,\ \Delta$	::=			

$|$ $\varnothing$
$|$ Δ, c
$|$ $\widetilde{\Gamma}$ M
$|$ (Δ) M

terminals $::=$
$|$ $\leftrightarrow$
$|$ $\Leftrightarrow$
$|$ $\longrightarrow$
$|$ min
$|$ $\equiv$
$|$ $\forall$
$|$ $\in$
$|$ $\not\in$
$|$ $\Leftarrow$
$|$ $\Rightarrow$
$|$ $\Rightarrow^*$
$|$ $\rightarrow$
$|$ Λ
$|$
$|$ $\Box$
$|$ $\vdash$
$|$ $\dashv$
$|$ $\models$
$|$ $\models$
$|$ $\neq$
$|$ $\triangleright$
$|$ ok
$|$ $\_$
$|$ $\leadsto$
$|$ $\leadsto^*$
$|$ $\leadsto$
$|$ $\varnothing$
$|$ $\circ$
$|$ fv
$|$ dom
$|$ $\sim$
$|$ $\asymp$
$|$ $|$
$|$ $\bullet$
$|$ $\mathbf{fst}$
$|$ $\mathbf{snd}$
$|$ $|\Rightarrow|$
$|$ $\vdash_=$
$|$ $\mathbf{refl_2}$
$|$ $++$

formula, ψ	$::=$		
	$\|$	*judgement*	
	$\|$	$x : A \in \Gamma$	
	$\|$	$x : R \in \Omega$	
	$\|$	$c : \phi \in \Gamma$	
	$\|$	$F\ sig\_sort \in \Sigma$	
	$\|$	$x \in \Delta$	
	$\|$	$c \in \Delta$	
	$\|$	$c\,\textbf{not relevant} \in \gamma$	
	$\|$	$x \notin \mathsf{fv}\,a$	
	$\|$	$x \notin \mathsf{dom}\,\Gamma$	
	$\|$	$uniq(\Omega)$	
	$\|$	$c \notin \mathsf{dom}\,\Gamma$	
	$\|$	$T \notin \mathsf{dom}\,\Sigma$	
	$\|$	$F \notin \mathsf{dom}\,\Sigma$	
	$\|$	$R_1 = R_2$	
	$\|$	$a = b$	
	$\|$	$\phi_1 = \phi_2$	
	$\|$	$\Gamma_1 = \Gamma_2$	
	$\|$	$\gamma_1 = \gamma_2$	
	$\|$	$\neg\psi$	
	$\|$	$\psi_1 \wedge \psi_2$	
	$\|$	$\psi_1 \vee \psi_2$	
	$\|$	$\psi_1 \Rightarrow \psi_2$	
	$\|$	(ψ)	
	$\|$	ψ	
	$\|$	$c : (a : A \sim b : B) \in \Gamma$	
	$\|$		suppress lc hypothesis generated by Ott
	$\|$	$\{y/x\}B = B_1$	
	$\|$	$\{c_1/c_2\}B = B_1$	
JSubRole	$::=$		
	$\|$	$R_1 \leq R_2$	Subroling judgement
JPath	$::=$		
	$\|$	$\mathsf{Path}_R\ a = F@Rs$	Type headed by constant (partial function)
JPat	$::=$		
	$\|$	$\Gamma \vDash a : Apat/R$	Pattern judgment
JTypeRoleList	$::=$		
	$\|$	$\mathsf{Roles}(A) = Rs$	type role list
JMatchSubst	$::=$		
	$\|$	$\mathsf{match}_R\,a_1\mathsf{with}\,a_2 \to b_1 = b_2$	match and substitute
JValue	$::=$		

	$\mid$	$\mathsf{Value}_R\ A$	values
JValueType	::=		
	$\mid$	$\mathsf{ValueType}_R\ A$	Types with head forms (erased language)
Jconsistent	::=		
	$\mid$	$\mathsf{consistent}_R\ ab$	(erased) types do not differ in their heads
Jroleing	::=		
	$\mid$	$\Omega \vDash a : R$	
JChk	::=		
	$\mid$	$(\rho = +) \vee (x \notin \mathsf{fv}\ A)$	irrelevant argument check
Jpar	::=		
	$\mid$	$\Omega \vDash a \Rightarrow_R b$	parallel reduction (implicit language)
	$\mid$	$\Omega \vdash a \Rightarrow^*_R b$	multistep parallel reduction
	$\mid$	$\Omega \vdash a \Leftrightarrow_R b$	parallel reduction to a common term
Jbeta	::=		
	$\mid$	$\vDash a > b/R$	primitive reductions on erased terms
	$\mid$	$\vDash a \leadsto b/R$	single-step head reduction for implicit language
	$\mid$	$\vDash a \leadsto^* b/R$	multistep reduction
Jett	::=		
	$\mid$	$\Gamma \vDash \phi\ \mathsf{ok}$	Prop wellformedness
	$\mid$	$\Gamma \vDash a : A$	typing
	$\mid$	$\Gamma; \Delta \vDash \phi_1 \equiv \phi_2$	prop equality
	$\mid$	$\Gamma; \Delta \vDash a \equiv b : A/R$	definitional equality
	$\mid$	$\vDash \Gamma$	context wellformedness
Jsig	::=		
	$\mid$	$\vDash \Sigma$	signature wellformedness
judgement	::=		
	$\mid$	*JSubRole*	
	$\mid$	*JPath*	
	$\mid$	*JPat*	
	$\mid$	*JTypeRoleList*	
	$\mid$	*JMatchSubst*	
	$\mid$	*JValue*	
	$\mid$	*JValueType*	
	$\mid$	*Jconsistent*	
	$\mid$	*Jroleing*	
	$\mid$	*JChk*	
	$\mid$	*Jpar*	
	$\mid$	*Jbeta*	

	$\mid$	*Jett*
	$\mid$	*Jsig*
user_syntax	::=	
	$\mid$	*tmvar*
	$\mid$	*covar*
	$\mid$	*datacon*
	$\mid$	*const*
	$\mid$	*index*
	$\mid$	*relflag*
	$\mid$	*appflag*
	$\mid$	*role*
	$\mid$	*constraint*
	$\mid$	*tm*
	$\mid$	*brs*
	$\mid$	*co*
	$\mid$	*role_context*
	$\mid$	*roles*
	$\mid$	*sig_sort*
	$\mid$	*sort*
	$\mid$	*context*
	$\mid$	*sig*
	$\mid$	*available_props*
	$\mid$	*terminals*
	$\mid$	*formula*

$\boxed{R_1 \leq R_2}$ Subroling judgement

$$\frac{}{\mathbf{Nom} \leq R} \quad \textsc{NomBot}$$

$$\frac{}{R \leq \mathbf{Rep}} \quad \textsc{RepTop}$$

$$\frac{}{R \leq R} \quad \textsc{Refl}$$

$$\frac{\begin{array}{l} R_1 \leq R_2 \\ R_2 \leq R_3 \end{array}}{R_1 \leq R_3} \quad \textsc{Trans}$$

$\boxed{\mathsf{Path}_R\ a = F@Rs}$ Type headed by constant (partial function)

$$\frac{F\ : A@Rs\ \in \Sigma_0}{\mathsf{Path}_R\ F = F@Rs} \quad \textsc{Path\_AbsConst}$$

$$\frac{\begin{array}{l} F\ \sim a : A/R_1@Rs\ \in \Sigma_0 \\ \neg(R_1 \leq R) \end{array}}{\mathsf{Path}_R\ F = F@Rs} \quad \textsc{Path\_Const}$$

$$\frac{\begin{array}{l} \mathsf{Path}_R\ a = F@R_1, Rs \\ app\_role \nu = R_1 \end{array}}{\mathsf{Path}_R\ (a\ b'^{\nu}) = F@Rs} \quad \textsc{Path\_App}$$

$$\frac{\mathsf{Path}_R\ a = F@Rs}{\mathsf{Path}_R\ (a[\bullet]) = F@Rs} \quad \textsc{Path\_CApp}$$

$\boxed{\Gamma \vDash a : A pat/R}$ Pattern judgment

$$\frac{F\ :A@Rs\ \in\ \Sigma_0}{\varnothing \vDash F : A pat/R}\quad \textsc{Pat\_AbsConst}$$

$$\frac{\begin{array}{l} F\ \sim a : A/R_1@Rs\ \in\ \Sigma_0 \\ \neg(R_1 \leq R) \end{array}}{\varnothing \vDash F : A pat/R}\quad \textsc{Pat\_Const}$$

$$\frac{\begin{array}{l} \Gamma \vDash a : \Pi^\rho y\!:\!A_1 \to B_1 pat/R \\ x\ \notin\ \mathsf{dom}\,\Gamma \\ \{y/x\}B = B_1 \end{array}}{\Gamma, x : A_1 \vDash (a\ x^\rho) : B pat/R}\quad \textsc{Pat\_App}$$

$$\frac{\begin{array}{l} \Gamma \vDash a : \forall c_1\!:\!\phi.B_1 pat/R \\ c\ \notin\ \mathsf{dom}\,\Gamma \\ \{c_1/c\}B = B_1 \end{array}}{\Gamma, c : \phi \vDash (a[\bullet]) : B pat/R}\quad \textsc{Pat\_CApp}$$

$\boxed{\mathsf{Roles}(A) = Rs}$ type role list

$$\frac{}{\mathsf{Roles}(\star) = \mathbf{nilR}}\quad \textsc{TypeRoleList\_Star}$$

$$\frac{\mathsf{Roles}(B) = Rs}{\mathsf{Roles}((\Pi^\rho x\!:\!A \to B)) = R_1, Rs}\quad \textsc{TypeRoleList\_Pi}$$

$$\frac{\mathsf{Roles}(B) = Rs}{\mathsf{Roles}((\forall c\!:\!\phi.B)) = Rs}\quad \textsc{TypeRoleList\_CPi}$$

$\boxed{\mathsf{match}_R\, a_1 \mathsf{with}\, a_2 \to b_1 = b_2}$ match and substitute

$$\frac{F\ :A@Rs\ \in\ \Sigma_0}{\mathsf{match}_R F \mathsf{with} F \to b = b}\quad \textsc{MatchSubst\_AbsConst}$$

$$\frac{\begin{array}{l} F\ \sim a : A/R_1@Rs\ \in\ \Sigma_0 \\ \neg(R_1 \leq R) \end{array}}{\mathsf{match}_R F \mathsf{with} F \to b = b}\quad \textsc{MatchSubst\_Const}$$

$$\frac{\mathsf{match}_R\, a_1 \mathsf{with}\, a_2 \to b_1 = b_2}{\mathsf{match}_R (a_1\ a^{R'}) \mathsf{with} (a_2\ x^+) \to b_1 = (b_2\{a/x\})}\quad \textsc{MatchSubst\_AppRelR}$$

$$\frac{\mathsf{match}_R\, a_1 \mathsf{with}\, a_2 \to b_1 = b_2}{\mathsf{match}_R (a_1\ a^{+}) \mathsf{with} (a_2\ x^+) \to b_1 = (b_2\{a/x\})}\quad \textsc{MatchSubst\_AppRel}$$

$$\frac{\mathsf{match}_R\, a_1 \mathsf{with}\, a_2 \to b_1 = b_2}{\mathsf{match}_R (a_1\ \square^-) \mathsf{with} (a_2\ \square^-) \to b_1 = b_2}\quad \textsc{MatchSubst\_AppIrrel}$$

$$\frac{\mathsf{match}_R\, a_1 \mathsf{with}\, a_2 \to b_1 = b_2}{\mathsf{match}_R (a_1[\bullet]) \mathsf{with} (a_2[\bullet]) \to b_1 = b_2}\quad \textsc{MatchSubst\_CApp}$$

$\boxed{\mathsf{Value}_R\ A}$ values

$$\frac{}{\mathsf{Value}_R\ \star}\quad \textsc{Value\_Star}$$

$$\frac{}{\mathsf{Value}_R\ \Pi^\rho x\!:\!A \to B}\quad \textsc{Value\_Pi}$$

$$\frac{}{\mathsf{Value}_R\ \forall c\!:\!\phi.B}\quad \textsc{Value\_CPi}$$

$$\frac{}{\mathsf{Value}_R\ \lambda^+ x\!:\!A.a} \quad \textsc{Value\_AbsRel}$$

$$\frac{}{\mathsf{Value}_R\ \lambda^+ x.a} \quad \textsc{Value\_UAbsRel}$$

$$\frac{\mathsf{Value}_R\ a}{\mathsf{Value}_R\ \lambda^- x.a} \quad \textsc{Value\_UAbsIrrel}$$

$$\frac{}{\mathsf{Value}_R\ \Lambda c\!:\!\phi.a} \quad \textsc{Value\_CAbs}$$

$$\frac{}{\mathsf{Value}_R\ \Lambda c.a} \quad \textsc{Value\_UCAbs}$$

$$\frac{\mathsf{Path}_R\ a = F@Rs}{\mathsf{Value}_R\ a} \quad \textsc{Value\_Path}$$

$\boxed{\mathsf{ValueType}_R\ A}$ Types with head forms (erased language)

$$\frac{}{\mathsf{ValueType}_R\ \star} \quad \textsc{value\_type\_Star}$$

$$\frac{}{\mathsf{ValueType}_R\ \Pi^\rho x\!:\!A \to B} \quad \textsc{value\_type\_Pi}$$

$$\frac{}{\mathsf{ValueType}_R\ \forall c\!:\!\phi.B} \quad \textsc{value\_type\_CPi}$$

$$\frac{\mathsf{Path}_R\ a = F@Rs}{\mathsf{ValueType}_R\ a} \quad \textsc{value\_type\_Path}$$

$\boxed{\mathsf{consistent}_R\ ab}$ (erased) types do not differ in their heads

$$\frac{}{\mathsf{consistent}_R\ \star\star} \quad \textsc{consistent\_a\_Star}$$

$$\frac{}{\mathsf{consistent}_{R'}\ (\Pi^\rho x_1\!:\!A_1 \to B_1)(\Pi^\rho x_2\!:\!A_2 \to B_2)} \quad \textsc{consistent\_a\_Pi}$$

$$\frac{}{\mathsf{consistent}_R\ (\forall c_1\!:\!\phi_1.A_1)(\forall c_2\!:\!\phi_2.A_2)} \quad \textsc{consistent\_a\_CPi}$$

$$\frac{\begin{array}{c}\mathsf{Path}_R\ a_1 = F@Rs\\ \mathsf{Path}_R\ a_2 = F@Rs\end{array}}{\mathsf{consistent}_R\ a_1 a_2} \quad \textsc{consistent\_a\_Path}$$

$$\frac{\neg\mathsf{ValueType}_R\ b}{\mathsf{consistent}_R\ ab} \quad \textsc{consistent\_a\_Step\_R}$$

$$\frac{\neg\mathsf{ValueType}_R\ a}{\mathsf{consistent}_R\ ab} \quad \textsc{consistent\_a\_Step\_L}$$

$\boxed{\Omega \vDash a : R}$

$$\frac{\mathit{uniq}(\Omega)}{\Omega \vDash \square : R} \quad \textsc{role\_a\_Bullet}$$

$$\frac{\mathit{uniq}(\Omega)}{\Omega \vDash \star : R} \quad \textsc{role\_a\_Star}$$

$$\frac{\begin{array}{l}\mathit{uniq}(\Omega)\\ x : R \in \Omega\\ R \leq R_1\end{array}}{\Omega \vDash x : R_1} \quad \textsc{role\_a\_Var}$$

$$\frac{\Omega, x : \mathbf{Nom} \vDash a : R}{\Omega \vDash (\lambda^{\rho} x.a) : R} \quad \textsc{role\_a\_Abs}$$

$$\frac{\begin{array}{l}\Omega \vDash a : R \\ \Omega \vDash b : \mathit{app\_role}\nu\end{array}}{\Omega \vDash (a\ b^{\nu}) : R} \quad \textsc{role\_a\_App}$$

$$\frac{\begin{array}{l}\Omega \vDash A : R \\ \Omega, x : \mathbf{Nom} \vDash B : R\end{array}}{\Omega \vDash (\Pi^{\rho} x\!:\!A \to B) : R} \quad \textsc{role\_a\_Pi}$$

$$\frac{\begin{array}{l}\Omega \vDash a : R_1 \\ \Omega \vDash b : R_1 \\ \Omega \vDash A : R_0 \\ \Omega \vDash B : R\end{array}}{\Omega \vDash (\forall c\!:\!a \sim_{A/R_1} b.B) : R} \quad \textsc{role\_a\_CPi}$$

$$\frac{\Omega \vDash b : R}{\Omega \vDash (\Lambda c.b) : R} \quad \textsc{role\_a\_CAbs}$$

$$\frac{\Omega \vDash a : R}{\Omega \vDash (a[\bullet]) : R} \quad \textsc{role\_a\_CApp}$$

$$\frac{\begin{array}{l}\mathit{uniq}(\Omega) \\ F\ : A@Rs\ \in\ \Sigma_0\end{array}}{\Omega \vDash F : R} \quad \textsc{role\_a\_Const}$$

$$\frac{\begin{array}{l}\mathit{uniq}(\Omega) \\ F\ \sim a : A/R@Rs\ \in\ \Sigma_0\end{array}}{\Omega \vDash F : R_1} \quad \textsc{role\_a\_Fam}$$

$$\frac{\begin{array}{l}\Omega \vDash a_1 : R \\ \Omega \vDash b_1 : R_1 \\ \Omega \vDash b_2 : R_1\end{array}}{\Omega \vDash \mathit{caseR}a_1 \mathit{of}\, a_2 \to b_1 \|_{\_} \to b_2 : R_1} \quad \textsc{role\_a\_Pattern}$$

$\boxed{(\rho = +) \vee (x \notin \mathsf{fv}\ A)}$ irrelevant argument check

$$\frac{}{(+ = +) \vee (x \notin \mathsf{fv}\ A)} \quad \textsc{Rho\_Rel}$$

$$\frac{x\ \notin\ \mathsf{fv} A}{(- = +) \vee (x \notin \mathsf{fv}\ A)} \quad \textsc{Rho\_IrrRel}$$

$\boxed{\Omega \vDash a \Rightarrow_R b}$ parallel reduction (implicit language)

$$\frac{\Omega \vDash a : R}{\Omega \vDash a \Rightarrow_R a} \quad \textsc{Par\_Refl}$$

$$\frac{\begin{array}{l}\Omega \vDash a \Rightarrow_R (\lambda^{\rho} x.a') \\ \Omega \vDash b \Rightarrow_{\mathit{app\_role}\nu} b'\end{array}}{\Omega \vDash a\ b^{\nu} \Rightarrow_R a'\{b'/x\}} \quad \textsc{Par\_Beta}$$

$$\frac{\begin{array}{l}\Omega \vDash a \Rightarrow_R a' \\ \Omega \vDash b \Rightarrow_{\mathit{app\_role}\nu} b'\end{array}}{\Omega \vDash a\ b^{\nu} \Rightarrow_R a'\ b'^{\nu}} \quad \textsc{Par\_App}$$

$$\frac{\Omega \vDash a \Rightarrow_R (\Lambda c.a')}{\Omega \vDash a[\bullet] \Rightarrow_R a'\{\bullet/c\}} \quad \textsc{Par\_CBeta}$$

$$\frac{\Omega \vDash a \Rightarrow_R a'}{\Omega \vDash a[\bullet] \Rightarrow_R a'[\bullet]} \quad \textsc{Par\_CApp}$$

$$\frac{\Omega, x : \mathbf{Nom} \vDash a \Rightarrow_R a'}{\Omega \vDash \lambda^\rho x.a \Rightarrow_R \lambda^\rho x.a'} \quad \textsc{Par\_Abs}$$

$$\frac{\begin{array}{l}\Omega \vDash A \Rightarrow_R A' \\ \Omega, x : \mathbf{Nom} \vDash B \Rightarrow_R B'\end{array}}{\Omega \vDash \Pi^\rho x\!:\!A \to B \Rightarrow_R \Pi^\rho x\!:\!A' \to B'} \quad \textsc{Par\_Pi}$$

$$\frac{\Omega \vDash a \Rightarrow_R a'}{\Omega \vDash \Lambda c.a \Rightarrow_R \Lambda c.a'} \quad \textsc{Par\_CAbs}$$

$$\frac{\begin{array}{l}\Omega \vDash A \Rightarrow_{R_0} A' \\ \Omega \vDash a \Rightarrow_{R_1} a' \\ \Omega \vDash b \Rightarrow_{R_1} b' \\ \Omega \vDash B \Rightarrow_R B'\end{array}}{\Omega \vDash \forall c\!:\!a \sim_{A/R_1} b.B \Rightarrow_R \forall c\!:\!a' \sim_{A'/R_1} b'.B'} \quad \textsc{Par\_CPi}$$

$$\frac{\begin{array}{l}F \sim a : A/R_1 @Rs \in \Sigma_0 \\ R_1 \leq R \\ \mathit{uniq}(\Omega)\end{array}}{\Omega \vDash F \Rightarrow_R a} \quad \textsc{Par\_Axiom}$$

$$\frac{\begin{array}{l}\Omega \vDash a_1 \Rightarrow_R a_1' \\ \Omega \vDash b_1 \Rightarrow_{R_0} b_1' \\ \Omega \vDash b_2 \Rightarrow_{R_0} b_2'\end{array}}{\Omega \vDash (\mathit{case} R a_1 \mathit{of} a_2 \to b_1 \| \_ \to b_2) \Rightarrow_{R_0} (\mathit{case} R a_1' \mathit{of} a_2 \to b_1' \| \_ \to b_2')} \quad \textsc{Par\_Pattern}$$

$$\frac{\begin{array}{l}\Omega \vDash a_1 \Rightarrow_R a_1' \\ \Omega \vDash b_1 \Rightarrow_{R_0} b_1' \\ \Omega \vDash b_2 \Rightarrow_{R_0} b_2' \\ \mathsf{match}_R\, a_1' \mathsf{with}\, a_2 \to b_1' = b\end{array}}{\Omega \vDash (\mathit{case} R a_1 \mathit{of} a_2 \to b_1 \| \_ \to b_2) \Rightarrow_{R_0} b} \quad \textsc{Par\_PatternTrue}$$

$$\frac{\begin{array}{l}\Omega \vDash a_1 \Rightarrow_R a_1' \\ \Omega \vDash b_1 \Rightarrow_{R_0} b_1' \\ \Omega \vDash b_2 \Rightarrow_{R_0} b_2' \\ \mathsf{Value}_R\ a_1' \\ \neg(\mathsf{match}_R\, a_1' \mathsf{with}\, a_2 \to b_1' = b)\end{array}}{\Omega \vDash (\mathit{case} R a_1 \mathit{of} a_2 \to b_1 \| \_ \to b_2) \Rightarrow_{R_0} b_2'} \quad \textsc{Par\_PatternFalse}$$

$\boxed{\Omega \vdash a \Rightarrow_R^* b}$ multistep parallel reduction

$$\frac{}{\Omega \vdash a \Rightarrow_R^* a} \quad \textsc{MP\_Refl}$$

$$\frac{\begin{array}{l}\Omega \vDash a \Rightarrow_R b \\ \Omega \vdash b \Rightarrow_R^* a'\end{array}}{\Omega \vdash a \Rightarrow_R^* a'} \quad \textsc{MP\_Step}$$

$\boxed{\Omega \vdash a \Leftrightarrow_R b}$ parallel reduction to a common term

$$\frac{\Omega \vdash a_1 \Rightarrow^*_R b \qquad \Omega \vdash a_2 \Rightarrow^*_R b}{\Omega \vdash a_1 \Leftrightarrow_R a_2} \quad \textsc{join}$$

$\boxed{\vDash a > b/R}$ primitive reductions on erased terms

$$\frac{\mathsf{Value}_{R_1}\ (\lambda^\rho x.v)}{\vDash (\lambda^\rho x.v)\ b^\nu > v\{b/x\}/R_1} \quad \textsc{Beta\_AppAbs}$$

$$\frac{}{\vDash (\Lambda c.a')[\bullet] > a'\{\bullet/c\}/R} \quad \textsc{Beta\_CAppCAbs}$$

$$\frac{F\ \sim a : A/R@Rs\ \in\ \Sigma_0 \qquad R \leq R_1}{\vDash F > a/R_1} \quad \textsc{Beta\_Axiom}$$

$$\frac{\mathsf{match}_R\, a_1 \mathsf{with}\, a_2 \to b_1 = b}{\vDash case R a_1 of\, a_2 \to b_1 \|_\_ \to b_2 > b/R_0} \quad \textsc{Beta\_PatternTrue}$$

$$\frac{\mathsf{Value}_R\ a_1 \qquad \neg(\mathsf{match}_R\, a_1 \mathsf{with}\, a_2 \to b_1 = b)}{\vDash case R a_1 of\, a_2 \to b_1 \|_\_ \to b_2 > b_2/R_0} \quad \textsc{Beta\_PatternFalse}$$

$\boxed{\vDash a \rightsquigarrow b/R}$ single-step head reduction for implicit language

$$\frac{\vDash a \rightsquigarrow a'/R_1}{\vDash \lambda^- x.a \rightsquigarrow \lambda^- x.a'/R_1} \quad \textsc{E\_AbsTerm}$$

$$\frac{\vDash a \rightsquigarrow a'/R_1}{\vDash a\ b^\nu \rightsquigarrow a'\ b^\nu/R_1} \quad \textsc{E\_AppLeft}$$

$$\frac{\vDash a \rightsquigarrow a'/R}{\vDash a[\bullet] \rightsquigarrow a'[\bullet]/R} \quad \textsc{E\_CAppLeft}$$

$$\frac{\vDash a \rightsquigarrow a_1'/R}{\vDash case R a_1 of\, a_2 \to b_1 \|_\_ \to b_2 \rightsquigarrow case R a_1' of\, a_2 \to b_1 \|_\_ \to b_2/R_0} \quad \textsc{E\_Pattern}$$

$$\frac{\vDash a > b/R}{\vDash a \rightsquigarrow b/R} \quad \textsc{E\_Prim}$$

$\boxed{\vDash a \rightsquigarrow^* b/R}$ multistep reduction

$$\frac{}{\vDash a \rightsquigarrow^* a/R} \quad \textsc{Equal}$$

$$\frac{\vDash a \rightsquigarrow b/R \qquad \vDash b \rightsquigarrow^* a'/R}{\vDash a \rightsquigarrow^* a'/R} \quad \textsc{Step}$$

$\boxed{\Gamma \vDash \phi\ \mathsf{ok}}$ Prop wellformedness

$$\frac{\Gamma \vDash a : A \qquad \Gamma \vDash b : A \qquad \Gamma \vDash A : \star}{\Gamma \vDash a \sim_{A/R} b\ \ \mathsf{ok}} \quad \textsc{E\_Wff}$$

$\boxed{\Gamma \vDash a : A}$ typing

$$\frac{\vDash \Gamma}{\Gamma \vDash \star : \star} \quad \text{E\_Star}$$

$$\frac{\begin{array}{l}\vDash \Gamma \\ x : A \in \Gamma\end{array}}{\Gamma \vDash x : A} \quad \text{E\_Var}$$

$$\frac{\begin{array}{l}\Gamma, x : A \vDash B : \star \\ \Gamma \vDash A : \star\end{array}}{\Gamma \vDash \Pi^\rho x{:}A \to B : \star} \quad \text{E\_Pi}$$

$$\frac{\begin{array}{l}\Gamma, x : A \vDash a : B \\ \Gamma \vDash A : \star \\ (\rho = +) \vee (x \notin \mathsf{fv}\ a)\end{array}}{\Gamma \vDash \lambda^\rho x.a : (\Pi^\rho x{:}A \to B)} \quad \text{E\_Abs}$$

$$\frac{\begin{array}{l}\Gamma \vDash b : \Pi^+ x{:}A \to B \\ \Gamma \vDash a : A\end{array}}{\Gamma \vDash b\ a^+ : B\{a/x\}} \quad \text{E\_App}$$

$$\frac{\begin{array}{l}\Gamma \vDash b : \Pi^+ x{:}A \to B \\ \Gamma \vDash a : A \\ \mathsf{Path}_{R'}\ a = F@R, Rs\end{array}}{\Gamma \vDash b\ a^R : B\{a/x\}} \quad \text{E\_TApp}$$

$$\frac{\begin{array}{l}\Gamma \vDash b : \Pi^- x{:}A \to B \\ \Gamma \vDash a : A\end{array}}{\Gamma \vDash b\ \Box^- : B\{a/x\}} \quad \text{E\_IApp}$$

$$\frac{\begin{array}{l}\Gamma \vDash a : A \\ \Gamma; \widetilde{\Gamma} \vDash A \equiv B : \star/\mathbf{Rep} \\ \Gamma \vDash B : \star\end{array}}{\Gamma \vDash a : B} \quad \text{E\_Conv}$$

$$\frac{\begin{array}{l}\Gamma, c : \phi \vDash B : \star \\ \Gamma \vDash \phi\ \mathsf{ok}\end{array}}{\Gamma \vDash \forall c{:}\phi.B : \star} \quad \text{E\_CPi}$$

$$\frac{\begin{array}{l}\Gamma, c : \phi \vDash a : B \\ \Gamma \vDash \phi\ \mathsf{ok}\end{array}}{\Gamma \vDash \Lambda c.a : \forall c{:}\phi.B} \quad \text{E\_CAbs}$$

$$\frac{\begin{array}{l}\Gamma \vDash a_1 : \forall c{:}(a \sim_{A/R} b).B_1 \\ \Gamma; \widetilde{\Gamma} \vDash a \equiv b : A/R\end{array}}{\Gamma \vDash a_1[\bullet] : B_1\{\bullet/c\}} \quad \text{E\_CApp}$$

$$\frac{\begin{array}{l}\vDash \Gamma \\ F\ : A@Rs \in \Sigma_0 \\ \varnothing \vDash A : \star\end{array}}{\Gamma \vDash F : A} \quad \text{E\_Const}$$

$$\frac{\begin{array}{l}\vDash \Gamma \\ F\ \sim a : A/R_1@Rs \in \Sigma_0 \\ \varnothing \vDash A : \star\end{array}}{\Gamma \vDash F : A} \quad \text{E\_Fam}$$

$$\frac{\begin{array}{l}\Gamma \vDash a_1 : A \\ \Gamma' \vDash a_2 : Apat/R \\ \Gamma, (\Gamma', c : \phi_1) \vDash b_1 : B \\ \Gamma \vDash b_2 : B \\ \phi_1 = (a_1 \sim_{A/R} a_2)\end{array}}{\Gamma \vDash caseR a_1 of a_2 \to b_1 \| \_ \to b_2 : B} \quad \text{E\_CASE}$$

$\boxed{\Gamma; \Delta \vDash \phi_1 \equiv \phi_2}$ prop equality

$$\frac{\begin{array}{c}\Gamma; \Delta \vDash A_1 \equiv A_2 : A/R \\ \Gamma; \Delta \vDash B_1 \equiv B_2 : A/R\end{array}}{\Gamma; \Delta \vDash A_1 \sim_{A/R} B_1 \equiv A_2 \sim_{A/R} B_2} \quad \text{E\_PROPCONG}$$

$$\frac{\begin{array}{l}\Gamma; \Delta \vDash A \equiv B : \star/R_0 \\ \Gamma \vDash A_1 \sim_{A/R} A_2 \ \mathsf{ok} \\ \Gamma \vDash A_1 \sim_{B/R} A_2 \ \mathsf{ok}\end{array}}{\Gamma; \Delta \vDash A_1 \sim_{A/R} A_2 \equiv A_1 \sim_{B/R} A_2} \quad \text{E\_ISOCONV}$$

$$\frac{\Gamma; \Delta \vDash \forall c\!:\!(a_1 \sim_{A/R_1} a_2).B_1 \equiv \forall c\!:\!(b_1 \sim_{B/R_2} b_2).B_2 : \star/R'}{\Gamma; \Delta \vDash a_1 \sim_{A/R_1} a_2 \equiv b_1 \sim_{B/R_2} b_2} \quad \text{E\_CPIFST}$$

$\boxed{\Gamma; \Delta \vDash a \equiv b : A/R}$ definitional equality

$$\frac{\begin{array}{l}\vDash \Gamma \\ c : (a \sim_{A/R} b) \in \Gamma \\ c \in \Delta\end{array}}{\Gamma; \Delta \vDash a \equiv b : A/R} \quad \text{E\_ASSN}$$

$$\frac{\Gamma \vDash a : A}{\Gamma; \Delta \vDash a \equiv a : A/\mathbf{Nom}} \quad \text{E\_REFL}$$

$$\frac{\Gamma; \Delta \vDash b \equiv a : A/R}{\Gamma; \Delta \vDash a \equiv b : A/R} \quad \text{E\_SYM}$$

$$\frac{\begin{array}{l}\Gamma; \Delta \vDash a \equiv a_1 : A/R \\ \Gamma; \Delta \vDash a_1 \equiv b : A/R\end{array}}{\Gamma; \Delta \vDash a \equiv b : A/R} \quad \text{E\_TRANS}$$

$$\frac{\begin{array}{l}\Gamma; \Delta \vDash a \equiv b : A/R_1 \\ R_1 \leq R_2\end{array}}{\Gamma; \Delta \vDash a \equiv b : A/R_2} \quad \text{E\_SUB}$$

$$\frac{\begin{array}{l}\Gamma \vDash a_1 : B \\ \Gamma \vDash a_2 : B \\ \vDash a_1 > a_2/R\end{array}}{\Gamma; \Delta \vDash a_1 \equiv a_2 : B/R} \quad \text{E\_BETA}$$

$$\frac{\begin{array}{l}\Gamma; \Delta \vDash A_1 \equiv A_2 : \star/R' \\ \Gamma, x : A_1; \Delta \vDash B_1 \equiv B_2 : \star/R' \\ \Gamma \vDash A_1 : \star \\ \Gamma \vDash \Pi^\rho x\!:\!A_1 \to B_1 : \star \\ \Gamma \vDash \Pi^\rho x\!:\!A_2 \to B_2 : \star\end{array}}{\Gamma; \Delta \vDash (\Pi^\rho x\!:\!A_1 \to B_1) \equiv (\Pi^\rho x\!:\!A_2 \to B_2) : \star/R'} \quad \text{E\_PICONG}$$

$$\frac{\begin{array}{l}\Gamma, x : A_1; \Delta \vDash b_1 \equiv b_2 : B/R' \\ \Gamma \vDash A_1 : \star \\ (\rho = +) \vee (x \notin \mathsf{fv}\ b_1) \\ (\rho = +) \vee (x \notin \mathsf{fv}\ b_2)\end{array}}{\Gamma; \Delta \vDash (\lambda^\rho x.b_1) \equiv (\lambda^\rho x.b_2) : (\Pi^\rho x\!:\!A_1 \to B)/R'} \quad \text{E\_AbsCong}$$

$$\frac{\begin{array}{l}\Gamma; \Delta \vDash a_1 \equiv b_1 : (\Pi^+ x\!:\!A \to B)/R' \\ \Gamma; \Delta \vDash a_2 \equiv b_2 : A/R'\end{array}}{\Gamma; \Delta \vDash a_1\ {a_2}^+ \equiv b_1\ {b_2}^+ : (B\{a_2/x\})/R'} \quad \text{E\_AppCong}$$

$$\frac{\begin{array}{l}\Gamma; \Delta \vDash a_1 \equiv b_1 : (\Pi^+ x\!:\!A \to B)/R' \\ \mathsf{Path}_{R'}\ a_1 = F@R, Rs \\ \Gamma; \Delta \vDash a_2 \equiv b_2 : A/\mathbf{param}\, R\, R'\end{array}}{\Gamma; \Delta \vDash a_1\ {a_2}^R \equiv b_1\ {b_2}^R : (B\{a_2/x\})/R'} \quad \text{E\_TAppCong}$$

$$\frac{\begin{array}{l}\Gamma; \Delta \vDash a_1 \equiv b_1 : (\Pi^- x\!:\!A \to B)/R' \\ \Gamma \vDash a : A\end{array}}{\Gamma; \Delta \vDash a_1\ \Box^- \equiv b_1\ \Box^- : (B\{a/x\})/R'} \quad \text{E\_IAppCong}$$

$$\frac{\Gamma; \Delta \vDash \Pi^\rho x\!:\!A_1 \to B_1 \equiv \Pi^\rho x\!:\!A_2 \to B_2 : \star/R'}{\Gamma; \Delta \vDash A_1 \equiv A_2 : \star/R'} \quad \text{E\_PiFst}$$

$$\frac{\begin{array}{l}\Gamma; \Delta \vDash \Pi^\rho x\!:\!A_1 \to B_1 \equiv \Pi^\rho x\!:\!A_2 \to B_2 : \star/R' \\ \Gamma; \Delta \vDash a_1 \equiv a_2 : A_1/R'\end{array}}{\Gamma; \Delta \vDash B_1\{a_1/x\} \equiv B_2\{a_2/x\} : \star/R'} \quad \text{E\_PiSnd}$$

$$\frac{\begin{array}{l}\Gamma; \Delta \vDash a_1 \sim_{A_1/R} b_1 \equiv a_2 \sim_{A_2/R} b_2 \\ \Gamma, c : a_1 \sim_{A_1/R} b_1; \Delta \vDash A \equiv B : \star/R' \\ \Gamma \vDash a_1 \sim_{A_1/R} b_1\ \ \mathsf{ok} \\ \Gamma \vDash \forall c\!:\!a_1 \sim_{A_1/R} b_1.A : \star \\ \Gamma \vDash \forall c\!:\!a_2 \sim_{A_2/R} b_2.B : \star\end{array}}{\Gamma; \Delta \vDash \forall c\!:\!a_1 \sim_{A_1/R} b_1.A \equiv \forall c\!:\!a_2 \sim_{A_2/R} b_2.B : \star/R'} \quad \text{E\_CPiCong}$$

$$\frac{\begin{array}{l}\Gamma, c : \phi_1; \Delta \vDash a \equiv b : B/R \\ \Gamma \vDash \phi_1\ \ \mathsf{ok}\end{array}}{\Gamma; \Delta \vDash (\Lambda c.a) \equiv (\Lambda c.b) : \forall c\!:\!\phi_1.B/R} \quad \text{E\_CAbsCong}$$

$$\frac{\begin{array}{l}\Gamma; \Delta \vDash a_1 \equiv b_1 : (\forall c\!:\!(a \sim_{A/R} b).B)/R' \\ \Gamma; \widetilde{\Gamma} \vDash a \equiv b : A/\mathbf{param}\, R\, R'\end{array}}{\Gamma; \Delta \vDash a_1[\bullet] \equiv b_1[\bullet] : (B\{\bullet/c\})/R'} \quad \text{E\_CAppCong}$$

$$\frac{\begin{array}{l}\Gamma; \Delta \vDash \forall c\!:\!(a_1 \sim_{A/R} a_2).B_1 \equiv \forall c\!:\!(a_1' \sim_{A'/R'} a_2').B_2 : \star/R_0 \\ \Gamma; \widetilde{\Gamma} \vDash a_1 \equiv a_2 : A/\mathbf{param}\, R\, R_0 \\ \Gamma; \widetilde{\Gamma} \vDash a_1' \equiv a_2' : A'/\mathbf{param}\, R'\, R_0\end{array}}{\Gamma; \Delta \vDash B_1\{\bullet/c\} \equiv B_2\{\bullet/c\} : \star/R_0} \quad \text{E\_CPiSnd}$$

$$\frac{\begin{array}{l}\Gamma; \Delta \vDash a \equiv b : A/R \\ \Gamma; \Delta \vDash a \sim_{A/R} b \equiv a' \sim_{A'/R'} b'\end{array}}{\Gamma; \Delta \vDash a' \equiv b' : A'/R'} \quad \text{E\_Cast}$$

$$\frac{\begin{array}{l}\Gamma; \Delta \vDash a \equiv b : A/R \\ \Gamma; \widetilde{\Gamma} \vDash A \equiv B : \star/\mathbf{Rep} \\ \Gamma \vDash B : \star\end{array}}{\Gamma; \Delta \vDash a \equiv b : B/R} \quad \text{E\_EqConv}$$

$$\frac{\Gamma; \Delta \vDash a \sim_{A/R_1} b \equiv a' \sim_{A'/R_1} b'}{\Gamma; \Delta \vDash A \equiv A' : \star/\mathbf{Rep}} \quad \text{E\_IsoSnd}$$

$$\frac{\begin{array}{l} \Gamma; \Delta \vDash a_1 \equiv a_1' : A/R \\ \Gamma; \Delta \vDash b_1 \equiv b_1' : B/R_0 \\ \Gamma; \Delta \vDash b_2 \equiv b_2' : B/R_0 \end{array}}{\Gamma; \Delta \vDash caseR a_1 of a_2 \to b_1 \| \_ \to b_2 \equiv caseR a_1' of a_2 \to b_1' \| \_ \to b_2' : B/R_0} \quad \text{E\_PatCong}$$

$$\frac{\begin{array}{l} \mathsf{Path}_{R'}\ a = F@R, Rs \\ \mathsf{Path}_{R'}\ a' = F@R, Rs \\ \Gamma \vDash a : \Pi^{+} x\!:\!A \to B \\ \Gamma \vDash b : A \\ \Gamma \vDash a' : \Pi^{+} x\!:\!A \to B \\ \Gamma \vDash b' : A \\ \Gamma; \Delta \vDash a\ b^{R_1} \equiv a'\ b'^{R_1} : B\{b/x\}/R' \\ \Gamma; \widetilde{\Gamma} \vDash B\{b/x\} \equiv B\{b'/x\} : \star/R' \end{array}}{\Gamma; \Delta \vDash a \equiv a' : \Pi^{+} x\!:\!A \to B/R'} \quad \text{E\_LeftRel}$$

$$\frac{\begin{array}{l} \mathsf{Path}_{R'}\ a = F@R, Rs \\ \mathsf{Path}_{R'}\ a' = F@R, Rs \\ \Gamma \vDash a : \Pi^{-} x\!:\!A \to B \\ \Gamma \vDash b : A \\ \Gamma \vDash a' : \Pi^{-} x\!:\!A \to B \\ \Gamma \vDash b' : A \\ \Gamma; \Delta \vDash a\ \square^{-} \equiv a'\ \square^{-} : B\{b/x\}/R' \\ \Gamma; \widetilde{\Gamma} \vDash B\{b/x\} \equiv B\{b'/x\} : \star/R_0 \end{array}}{\Gamma; \Delta \vDash a \equiv a' : \Pi^{-} x\!:\!A \to B/R'} \quad \text{E\_LeftIrrel}$$

$$\frac{\begin{array}{l} \mathsf{Path}_{R'}\ a = F@R, Rs \\ \mathsf{Path}_{R'}\ a' = F@R, Rs \\ \Gamma \vDash a : \Pi^{+} x\!:\!A \to B \\ \Gamma \vDash b : A \\ \Gamma \vDash a' : \Pi^{+} x\!:\!A \to B \\ \Gamma \vDash b' : A \\ \Gamma; \Delta \vDash a\ b^{+} \equiv a'\ b'^{+} : B\{b/x\}/R' \\ \Gamma; \widetilde{\Gamma} \vDash B\{b/x\} \equiv B\{b'/x\} : \star/R_0 \end{array}}{\Gamma; \Delta \vDash b \equiv b' : A/\mathbf{param}\, R_1\, R'} \quad \text{E\_Right}$$

$$\frac{\begin{array}{l} \mathsf{Path}_{R'}\ a = F@R, Rs \\ \mathsf{Path}_{R'}\ a' = F@R, Rs \\ \Gamma \vDash a : \forall c\!:\!(a_1 \sim_{A/R_1} a_2).B \\ \Gamma \vDash a' : \forall c\!:\!(a_1 \sim_{A/R_1} a_2).B \\ \Gamma; \widetilde{\Gamma} \vDash a_1 \equiv a_2 : A/R' \\ \Gamma; \Delta \vDash a[\bullet] \equiv a'[\bullet] : B\{\bullet/c\}/R' \end{array}}{\Gamma; \Delta \vDash a \equiv a' : \forall c\!:\!(a_1 \sim_{A/R_1} a_2).B/R'} \quad \text{E\_CLeft}$$

$\boxed{\vDash \Gamma}$ context wellformedness

$$\frac{}{\vDash \varnothing} \quad \text{E\_Empty}$$

$$\frac{\begin{array}{l} \vDash \Gamma \\ \Gamma \vDash A : \star \\ x \notin \mathsf{dom}\,\Gamma \end{array}}{\vDash \Gamma, x : A} \quad \text{E\_ConsTm}$$

$$\frac{\models \Gamma \quad \Gamma \models \phi \ \mathsf{ok} \quad c \notin \mathsf{dom}\,\Gamma}{\models \Gamma, c : \phi} \quad \text{E\_ConsCo}$$

$\boxed{\models \Sigma}$ signature wellformedness

$$\frac{}{\models \varnothing} \quad \text{Sig\_Empty}$$

$$\frac{\models \Sigma \quad \varnothing \models A : \star \quad F \notin \mathsf{dom}\,\Sigma \quad \mathsf{Roles}(A) = Rs}{\models \Sigma \cup \{F : A@Rs\}} \quad \text{Sig\_ConsConst}$$

$$\frac{\models \Sigma \quad \varnothing \models a : A \quad F \notin \mathsf{dom}\,\Sigma \quad \mathsf{Roles}(A) = Rs}{\models \Sigma \cup \{F \sim a : A/R@Rs\}} \quad \text{Sig\_ConsAx}$$

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Definition rules:        131 good    0 bad
Definition rule clauses: 376 good    0 bad
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