Relational Databases

Design Theory

8 February 2024

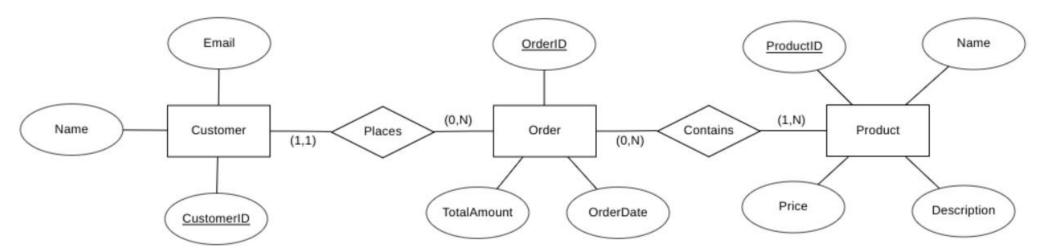
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Quick Recap: Entity-Relationship Diagram

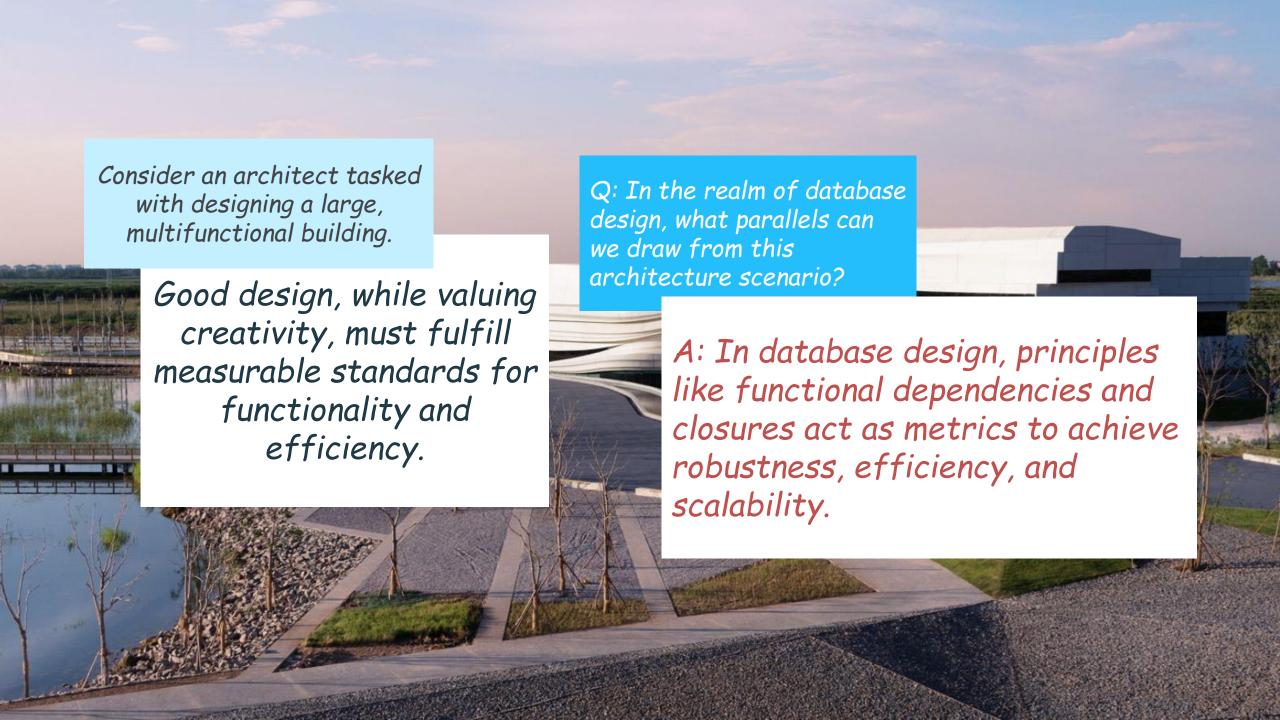
Key Concepts:

- ERDs are foundational in visualizing and structuring the data relationships in a database, facilitating clear communication and planning in the design process. They highlight entities, their attributes, and the relationships between them, crucial for understanding the database's conceptual layout.
- One-to-One, One-to-Many, Many-to-Many, Self-Referencing, and M-way are critical in ERDs. These relationships help in accurately modeling complex real-world interactions within the database.

ERD Example



Scenario: Customer's Order History.

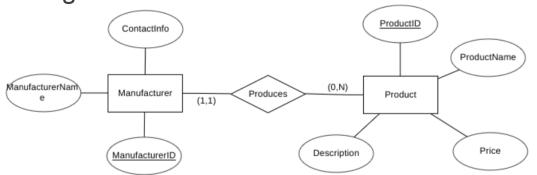


Transfroming ER Models to Relational Schemas

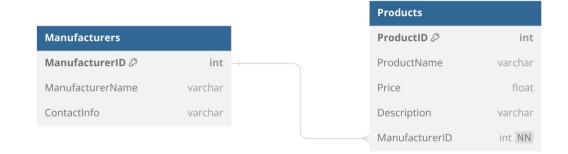
Key Points:

- Entity Conversion: Each entity in the ER model becomes a table in the schema.
- Relationship Handling: Relationships in ER models are translated into foreign keys and join tables.
- Attribute Translation: Attributes of entities in the ER model become the columns of the tables in the schema.

ER Diagram:



Relational Schema:



Best Practices:

- Consistency: Keep names and types consistent during the transition.
- Integrity Constraints: Enforce data integrity through constraints derived from the ER model (e.g., not null, unique).

Normalisation Consideration

The schema often needs normalisation to ensure efficiency and data integrity.

What Makes a Good Data Model?

Efficient Representation of Data Relationships

 A good data model accurately represents the relationships between different data entities, ensuring clarity and logical structure.

Minimised Redundancy

• Reduces unnecessary repetition of data across the database, enhancing efficiency and storage utilization.

Minimised Modification Anomalies

 Effectively manages and minimizes potential errors in data modification, such as update, insertion, and deletion anomalies.

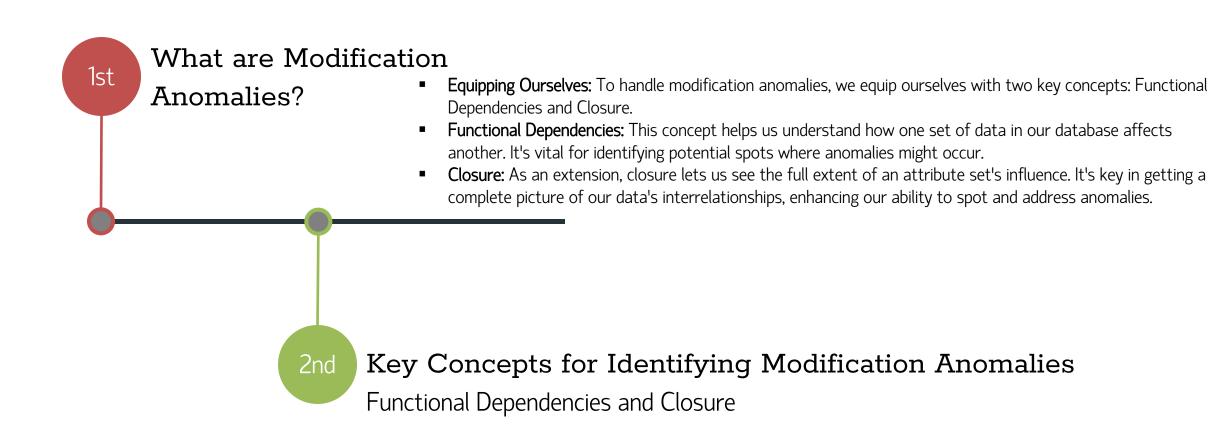
Data Integrity and Accessibility

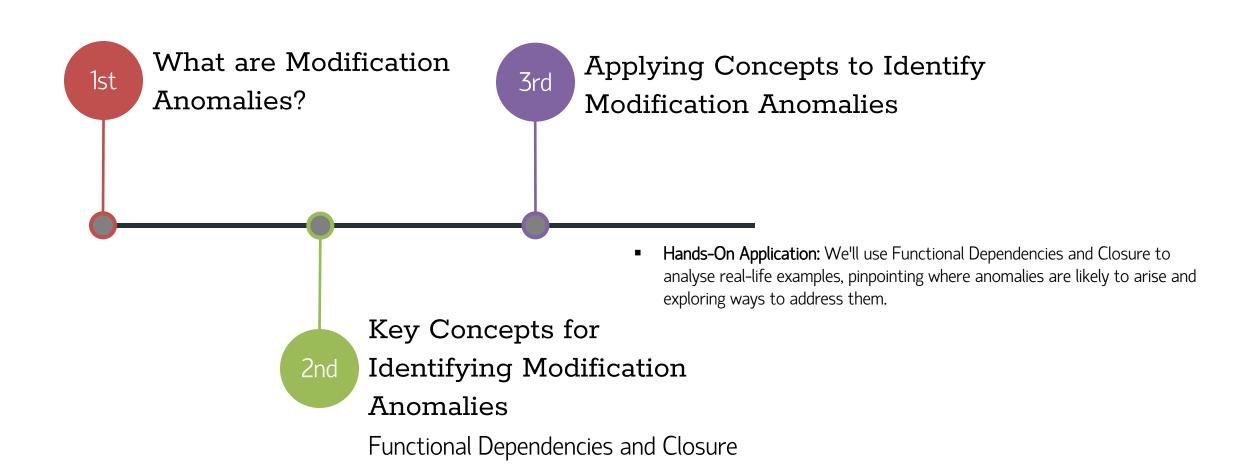
 Maintains high data integrity and provides ease of access and use, making the database reliable and user-friendly.

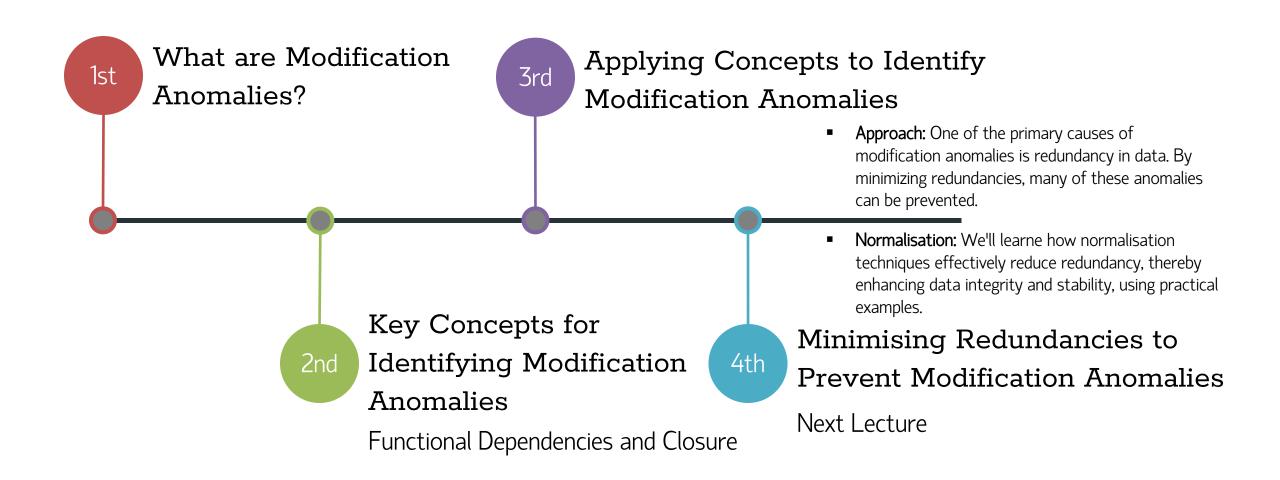
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What are Modification Anomalies?

- The Foundation: Before managing modification anomalies, it's essential to know what they are. These are issues in databases that lead to inconsistent data during Insertions, Deletions, and Updates.
- Why It Matters: Identifying these anomalies is the first critical step. It's about understanding the challenges in database operations to develop effective strategies for handling them.







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What are Modification Anomalies?

- The Foundation: Before managing modification anomalies, it's essential to know what they are. These are issues in databases that lead to inconsistent data during Insertions, Deletions, and Updates.
- Why It Matters: Identifying these anomalies is the first critical step. It's about understanding the challenges in database operations to develop effective strategies for handling them.

Modification Anomalies

Key Points:

- Modification Anomalies are unintended side effects in the database system, often arising from inadequate or flaw design.
- These lead to inconsistent data, redundant or duplicated entries and impact in overall efficiency and performance of the database

Types of Anomalies

Insertion Anomaly, Deletion Anomaly and Update Anomaly

Insertion Anomalies

Definition

• An insertion anomaly occurs in a database when the addition of new data is hindered or results in incomplete records due to the reliance on other related data being present.

Example Scenario

• Context: In an e-commerce database with combined product and supplier details, adding a new supplier without associated products leads to records with NULL values for product columns.

Product ID Produc	ct Name S	Supplier ID	Supplier Name	Supplier Contact	
1001 Solar NULL NULL	String Lights S S			+1234567890 +1122334455	

Explanation

In this scenario, adding 'NewEco Inc.' as a new supplier without a corresponding product creates entries with NULL values for product fields. This occurrence, typical in denormalized databases, exemplifies an insertion anomaly: you can add supplier details, but the absence of related product data leads to incomplete records.

Deletion Anomalies

Definition

 A deletion anomaly occurs when removing a record from a database inadvertently leads to the loss of additional, valuable data.

Example Scenario

• Context: With product details and their supplier information are stored in the same table. When the last product from a specific supplier is deleted, the supplier's details are also lost from the database

Product ID	Product Name	Supplier ID	Supplier Name	Supplier Contact
1001	Solar String Lights	S01	EcoLights Co.	
1002	LED Lantern	S02	BrightLumen Inc.	

Deletion Action

Delete 'Solar String Lights' (the only product supplied by 'EcoLights Co.')

Explanation

• By deleting 'Solar String Lights', all information about its supplier 'EcoLights Co.' is also removed from the table. This deletion anomaly leads to the loss of important supplier details that might be needed for future reference or orders.

Update Anomalies

Definition

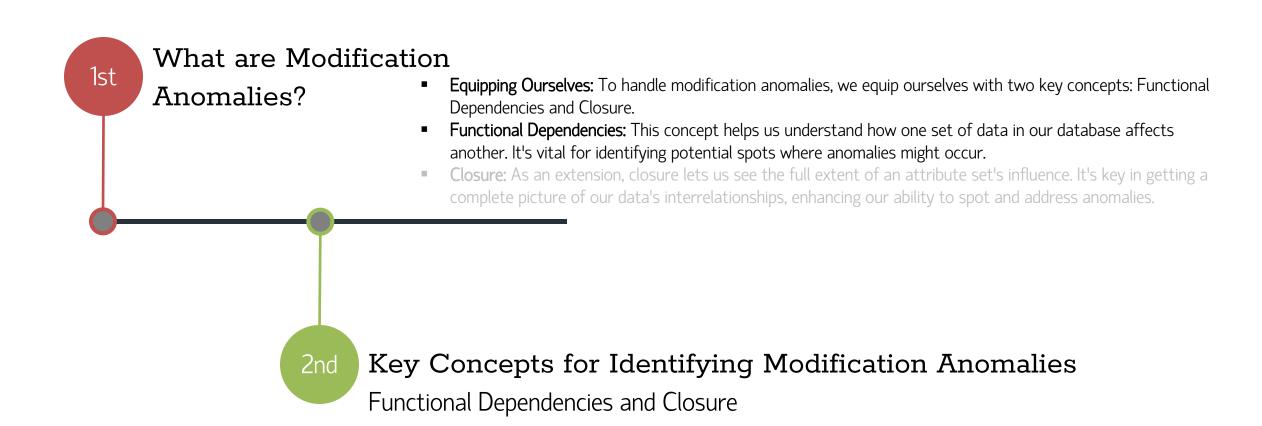
 Update anomalies occur when changes made to data in one part of a denormalized database are not consistently replicated in other related parts, leading to data inconsistencies or redundancies.

Example Scenario

• Context: Consider an database where each product's record includes supplier details. An update anomaly happens when changing a supplier's contact information in one product record does not automatically update the same information in other products supplied by the same supplier.

Explanation

In this example, updating the contact information for 'EcoLights Co.' in the record for 'Solar String Lights' doesn't update the same information for 'Eco Desk Lamp'. This inconsistency in supplier contact details across different product records is a classic update anomaly.



Functional Dependencies

Definition

- Functional Dependencies are relationships where one set of data in a database uniquely determines another set. These relationships are key to structuring databases effectively.
- Importance: Understanding FDs helps in creating logically organized and efficient databases.

Example Scenario

Customer Table:

List of Dependencies:

```
CustomerID -> CustomerName, Address
```

E-Commerce Example

Scenario

Product Table:

List of Dependencies:

```
SKU Number -> Name, Price, Description
```

Explanation

- A single SKU Number uniquely determines a product's Price, Name, Description, and Category.
- This dependency ensures consistent and reliable data across the database, which is essential for inventory management, customer experience, and analytics.

Identifying Functional Dependencies

The Process of Identifying Functional Dependencies

- Identifying FDs often involves looking for patterns and consistent relationships within your data.
- Example Scenario: (Sales Table)

```
Total Amount |
                                                        Product ID
Order ID
           Order Date l
                         Customer ID |
01001
           2021-07-01
                         C123
                                        150.00
                                                        P001
01002
           2021-07-02
                                        200.00
                                                        P002
                         C124
01003
          | 2021-07-02 |
                         C125
                                        250.00
                                                        P003
```

- Practical Tips:
 - Unique Identifiers: Start with primary keys or unique identifiers as they often determine other information.
 - One-to-One Relationship: Look for cases where one data element always corresponds to a specific value of another.
 - Real-World Logic: Apply your understanding of the real-world context. For instance, in a sales database, an 'Order ID' will
 determine 'Order Date' and 'Total Amount'."
- List of Dependencies:

```
Order ID -> Order Date, Customer ID, Total Amount, Product ID
```

Identifying Functional Dependencies (cont.)

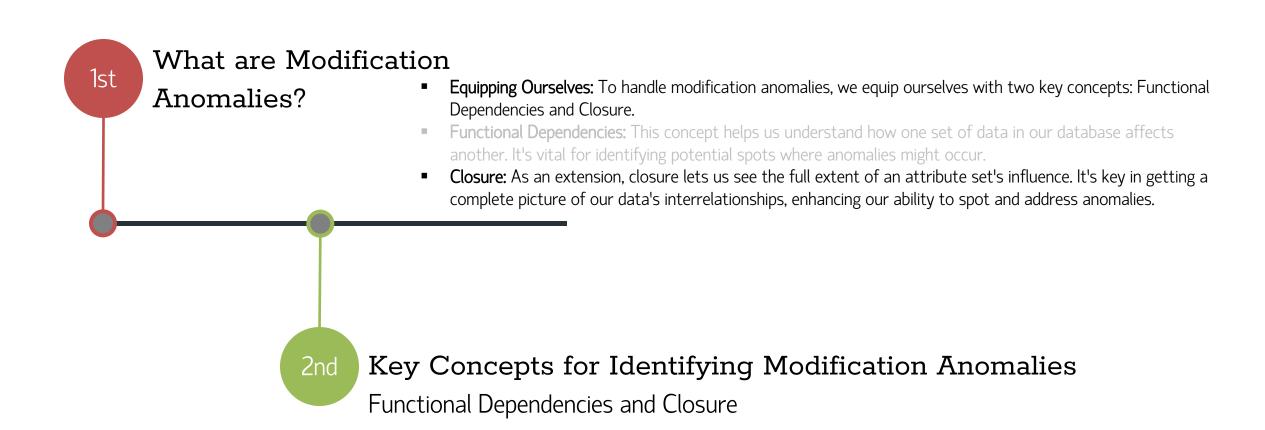
Example Scenario: (Sales Table)

Order ID Product ID	Product Name	Customer ID	Customer Name	Order Total
		-		
01001 P001	Wireless Mouse	C123	Alice	300.00
01001 P002	Bluetooth Keyboard	C123	Alice	300.00
01002 P002	Bluetooth Keyboard	C124	Bob	200.00
01003 P003	USB-C Hub	C125	Carol	450.00
01003 P001	Wireless Mouse	C125	Carol	450.00

List of Dependencies:

```
Order ID -> Order Total, Customer ID, Customer Name
Product ID -> Product Name
Customer ID -> Customer Name
```

- The current table's structure, where Order ID and Product ID are in the same table without a clear functional dependency, can lead to data redundancy and complications in maintaining the database, as seen with repeated customer and order total information.
- In a well-structured database, a separate junction table or line items table would typically be used to appropriately represent this many-to-many relationship.



Understanding Closure

What is Closure in Functional Dependencies?

- Closure in database theory, denoted as X+, refers to the complete set of attributes functionally determined by a given set of attributes X. It represents all data inferred from X.
- Point: Closure expands functional dependencies, revealing the full impact of a set of attributes within a table.

Example Scenario

Closure:

```
| Set of Attributes | Closure (X+) |
|-----|
| CustomerID | Name, Address |
```

- This means that knowing the CustomerID of any entry allows us to determine the corresponding customer's Name and Address uniquely.
- This closure indicates a direct and exclusive relationship between CustomerID and the two other attributes, highlighting the role of CustomerID as a potential primary key for this table.

Identifying Superkeys and Keys

Closure and Superkeys

 A superkey is a set of one or more attributes that can uniquely identify each record in a table. Closures help us find these superkeys.

Example Scenario

Customer ID	Name	Address	Email
C001	Alice Smith	123 Maple St	alice@example.com
C002	Bob Johnson	456 Oak Ave	bob@example.com
C003	Carol Davis	789 Pine Rd	carol@example.com

Closure:

```
| Attribute Set | Closure (All Determined Attributes) | Is Superkey? |
|-----|
| CustomerID | CustomerID, Email, Name, Address | Yes |
```

• In a Customer table with CustomerID, Email, Name, and Address, if the closure of CustomerID includes all these attributes, then CustomerID is a superkey.

Identifying Superkeys and Keys (cont.)

Example Scenario

Customer ID	Name	Address	Email
C001	Alice Smith	123 Maple St	alice@example.com
C002	Bob Johnson	456 Oak Ave	bob@example.com
C003	Carol Davis	789 Pine Rd	carol@example.com

Refining Superkeys to Candidate Keys

- A candidate key is the smallest possible superkey it cannot have any redundant attributes.
- Closure:

- If 'CustomerlD' alone can uniquely identify every record (closure of 'CustomerlD' is the entire set of table attributes), then 'CustomerlD' is not just a superkey but also a candidate key. It's the minimal attribute set needed for unique identification."
- Additional Consideration: If adding Email to CustomerID doesn't expand the closure beyond the table's attributes, then (CustomerID, Email) is not a candidate key due to its non-minimality.

Utilising Closures on Complex Example

Order Details Example

OrderID	CustomerID	ProductID	OrderDate	Quantity	Price
	 C001	 P001	 2021-01-01	•	 20.00
0102 0103	C002 C003	P002 P003	2021-01-02 2021-01-03	•	15.00 10.00

Closure:

```
Closure (All Determined Attributes)
Attribute Set
                                                           Is Superkey?
OrderID
                     OrderID, CustomerID, ProductID,
                                                            Yes
                     OrderDate, Quantity, Price
CustomerID
                     CustomerID
                                                            No
ProductID
                     ProductID
                                                            No
OrderID, ProductID |
                     OrderID, CustomerID, ProductID,
                                                            Yes
                     OrderDate, Quantity, Price
```

Utilising Closures on Complex Example

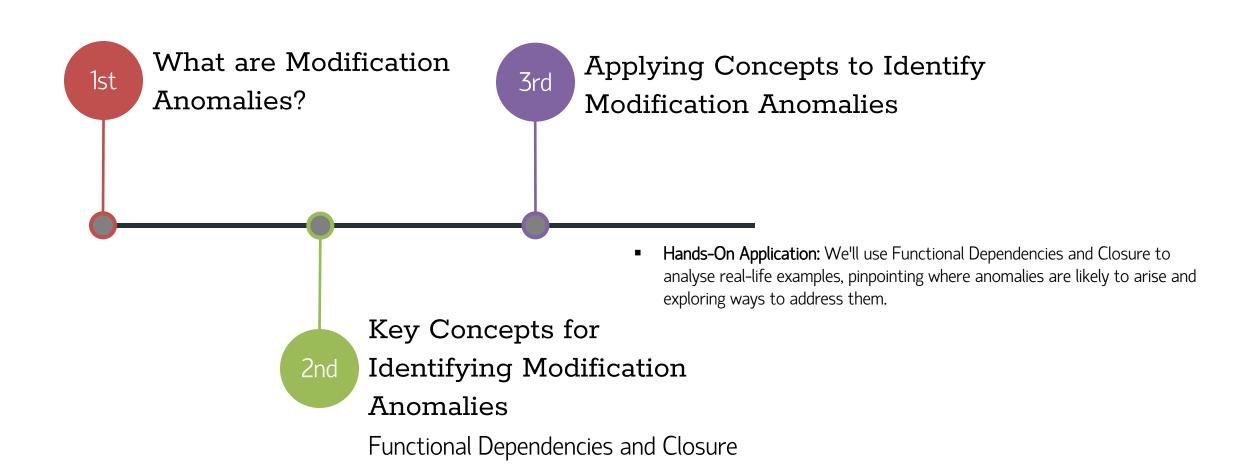
Order Details Example

OrderID Custo	merID Product	ID OrderDate Quan [.]	tity Price
0101 C001	P001	2021-01-01 2	20.00
0102 C002	P002	2021-01-02 1	15.00
0103 C003	P003	2021-01-03 5	10.00

Closure:

```
Attribute Set
                     Closure (All Determined Attributes)
                                                            Is Superkey?
                                                                          | Is Key?
OrderID
                     OrderID, CustomerID, ProductID,
                                                                            No
                     OrderDate, Quantity, Price
                     CustomerID
CustomerID
                                                            No
                                                                            No
ProductID
                     ProductID
                                                                            No
                                                            No
                                                                            Yes (Candidate Key)
OrderID, ProductID |
                     OrderID, CustomerID, ProductID,
                                                            Yes
                     OrderDate, Quantity, Price
```

- OrderID alone is a superkey but not a candidate key as it's not minimal.
- CustomerID and ProductID alone do not determine all attributes, hence not superkeys.
- The combination (OrderID, Product ID) is both a superkey and a candidate key, being minimal yet determining all table attributes.



Recap and Beyond: Functional Dependencies and Closures

Functional Depedencies

- Focus: This is about the relationship between sets of attributes in a database. Specifically, it indicates that if you know the value of one set of attributes, you can determine the value of another set.
- Modification Anomalies: FDs identify potential update anomalies. By structuring tables based on FDs, update
 anomalies can be minimized, as each piece of information is stored only once and updated consistently.

Closure

- Focus: Closure involves identifying all the attributes that a given set of attributes can functionally determine. It extends the concept of functional dependencies to a broader set of attributes.
- Modification Anomalies: Closures help prevent insertion and deletion anomalies. By understanding which attributes determine others, database designers can organize tables to allow independent data entry (minimizing insertion anomalies) and prevent unintended data loss (minimizing deletion anomalies).

Functional Dependencies: Identifying Update Anomalies

Order Table Example

```
CustomerID
             CustomerName
                            Address
                                                OrderID
                                                          OrderDate
                                                                        OrderAmount
                            123 Apple Street
C001
             Alice Smith
                                                0100
                                                          2021-07-15
                                                                        $150
C001
             Alice Smith
                            123 Apple Street
                                                0101
                                                          2021-07-20
                                                                        $200
C002
             Bob Johnson
                            456 Berry Avenue
                                                0102
                                                          2021-07-18
                                                                        $250
```

Functional Dependency and Anomaly Analysis

- Anomaly Detection Pointer: In summary, seeing a supposed-to-be unique identifier like CustomerID and OrderID repeating in a dataset is a flag for a potential anomaly that should be investigated further.
- Scenario: Imagine a scenario where, in future data entries, 'CustomerName' varies for the same 'CustomerlD', which would be an update anomaly. This inconsistency would indicate a violation of the functional dependency ('CustomerlD' → 'CustomerName', 'Address').

Closure and Keys: Identifying Insertion Anomalies

Product Table Example

ProductID	ProductName	SupplierID	SupplierName
P001	Laptop	S001	TechCorp
P002	Phone	S002	MobileInc

Closure and Key Analysis

Attribute Set 	Closure (All Determined Attributes)	Is Superkey? 	Is Key?	Anomaly Indication
ProductID 	ProductID, ProductName	No I	No	Potential for Insertion Anomaly (Incomplete data without `SupplierID` and `SupplierName`)
SupplierID ProductID, SupplierID 	SupplierID, SupplierName ProductID, ProductName, SupplierID, SupplierName	No Yes 	No Yes (Candidate Key)	No apparent anomaly

- Anomaly Detection Pointer: Watch for missing data in entries. If key attributes like 'SupplierID' and 'SupplierName', expected from the closure of 'ProductID', are absent in new records, it suggests an insertion anomaly.
- Scenario: Imagine adding a new product (P003) without 'SupplierID' and 'SupplierName' exemplifies an insertion anomaly. This case violates the closure rule where 'ProductID' and 'SupplierID' together should determine complete product details, including 'ProductName' and 'SupplierName'.

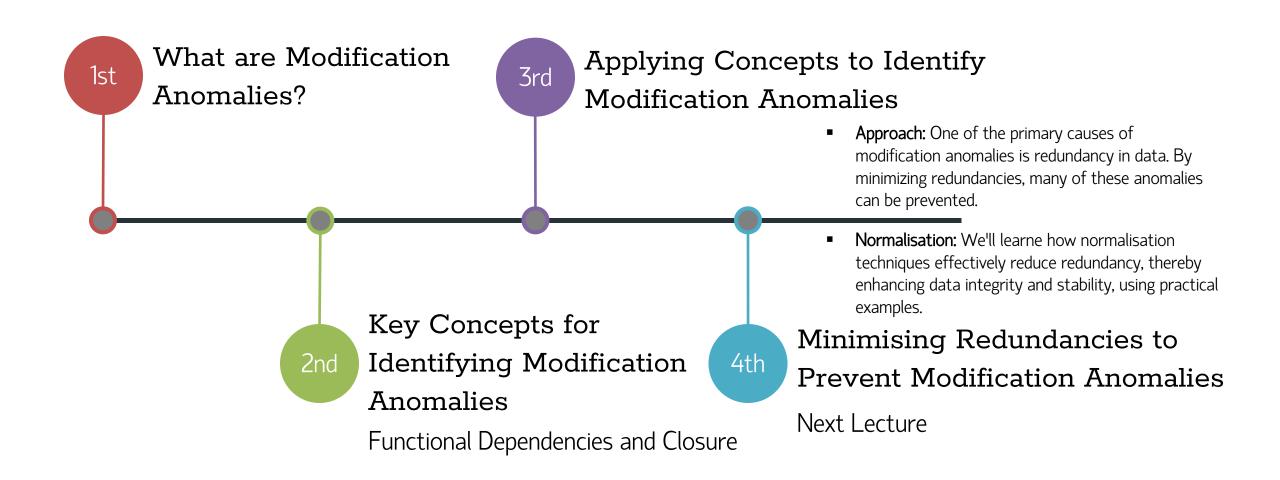
Closure and Keys: Identifying Deletion Anomalies

Product Table Example

SupplierID	SupplierName	ProductID	ProductName
S001	TechCorp	P001	Laptop
S001	TechCorp	P002	Phone
S002	MobileInc	P003	Tablet

Closure and Key Analysis

- Anomaly Detection Pointer: Pay attention to how interconnected attributes are based on closures. Deleting a
 record that's a key part of a closure can lead to unintended loss of crucial related information, representing a
 deletion anomaly.
- Scenario: Imagine deleting a supplier (like S001) could lead to the deletion of all related products (P001, P002). This would be a deletion anomaly as the removal of a supplier record inadvertently leads to the loss of important product data.



Recap and Key Takeaways

- A good data model is characterized by its efficiency in representing data relationships and its effectiveness in minimizing modification anomalies, ensuring data consistency and integrity across the database.
- Functional dependencies and closure serve as crucial metrics for identifying potential anomalies. These insights guide database design, helping to structure data in a way that enhances reliability and scalability while minimizing redundancy and data anomalies.

CustomerID	CustomerName	Address	1
 C001 C002	 Alice Smith Bob Johnson	123 Apple Street 456 Berry Avenue	
C001	Alice Smith	789 Cherry Lane	Anomaly

An anomaly is observed with CustomerID C001. It is associated with two different addresses ('123 Apple Street' and '789 Cherry Lane'), which violates the functional dependency. This inconsistency suggests an update anomaly.

Preparing for Normalisation:

- Next Lecture Preview:
 - Understand key methods for organizing database structures to enhance efficiency and minimise modification anomalies.
 - Explore various normalization forms and their significance in optimizing database design and functionality.

In what ways do you think the concepts of superkeys and candidate keys are essential for maintaining the integrity of a relational database?