Lecture Assignment 2 Name: Sneyny THA SID: 12113053

19. 0 n

factor of

(a) n' > [m) = 4 n' increase by 4

 $0 \quad \stackrel{?}{n} \rightarrow (n+1) = \stackrel{?}{n+2} + 2n+1 \text{ in crease by } 2n+1$ 

(b). n

a.  $\vec{n} \rightarrow (2n)^3 = 8\vec{n}$  increase by factor of 8

6. P → (n+) = n+3+3+3n+1 increased by 3++3n+1

c)-1001

a). (001) -> (4001) = 2(1001) increased by factor of 4 b) (001) -> 100(n+1) = 1001] +20011+100 increased by

2007 +100.

d). nlogn

a).  $nlogn \rightarrow (2nlog2n) = nlog(2n) = nlog(2n) = nlog(2n)$  = n(log n + log 4n) = nlogn + nlog4n

= n/ogn+ n(logn+ loga)

= nlogn + nlogn= 2nlogn + 2n ( $logn = log_2n$ ) increased by factor of 2 plus 2n.

b.  $nlogn \rightarrow (n+1) log(n+1) = nlog(n+1) + log(n+1)$ Let remove nlogn

= nlog(n+1) +nlog (n+1)-nlogn=nlogn+1)+

= log(n+1)+n[log(n+1)-logn]

increased by log(n+1)+n[log(n+1)-logn]

(e). 2

a. 2" -> 2" = (z") increased by Square factor

5. given 
$$f(n) = O(g(n))$$
,  $f(n) \leq cg(n)$ 

a). False,  $f(n) = 2$ ,  $g(n) = 1$ ,  $log(g(n)) = 0 \Rightarrow log f(n) > log(g(n))$ 

b). False,  $f(n) = 2n$ 

g(n) = n

 $f(n) \leq c \cdot g(n)$  ( $f(n) = O(g(n))$ )

c  $g(n)$ 
 $f(n) \leq c \cdot g(n)$ 

C. True

 $f(n) \leq c \cdot g(n)$ 
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