International University of Business Agriculture and Technology

May 6, 2025

			4	D.ľ	Number Theory	5		7.9 Notes	15
				4.1	1 - N Divisor	5		7.10 Topological Sort	15
				4.2	Catalan Numbers	6		7.11 dijkstra simple	15
	$\operatorname{Contents}$			4.3	Eular Totient Of Every Number 1-N	6			
				4.4	GCD LCM Equations	6	8	G.Tree	16
1	A.Settings and Script	1		4.5	LOG with any base	7		8.1 2D Fenwick Tree	
	1.1 Generator	1		4.6	Large-number-fibonacci	7		8.2 MST Kruskal's	
	1.2 Main Script	1		4.7	Sieve of Eratosthenes	7		8.3 MST To find Kth Sortest Number in Rang	
	1.3 Settings	$\frac{1}{2}$		4.8	Total Digit of N factorial	7		8.4 Merger Sort Tree Using Segment Tree .	17
	1.4 Template	2						8.5 Minimum Spanning Tree – MST using	
	1.4 Template	2	5	Dat	ta Structures	7		Prim's Algo	17
2	B.Combinatorics	2		5.1	MO's algo coordinate compression	7	•	TT Ct .	
_	2.1 Binomial coefficient	2		5.2	Segment Tree Lazy	8	9	H.Strings	17
				5.3	Segment Tree	9		9.1 KMP Algorithm	
	2.2 Extended Euclidean Algorithm	2		5.4	ordered set	10		9.2 LCS	18
	2.3 General Equation	3						9.3 Lexiographically smallest string by cycle	
	2.4 How Many Digit X^Y	3	6	$\mathbf{E}.\mathbf{D}$		10		shift	
	2.5 How many digit in N!	3		6.1	8	10		9.4 Longest Palindromic Subsequence	
	2.6 Last Non Zero Digit of Factorial	3		6.2	Knapsack 1	11		9.5 Manacher's Algorithm	
	2.7 Matrix Exponentitation	3		6.3	Longest Incresing Subsequence	11		9.6 Trie	
	2.8 Modular Inverse	3	_	П.С	N 1			9.7 Z function	
	2.9 combination1	3	7	F.G	Graph	11		9.8 string hashing	19
	2.10 nCr with Big Mod	4		7.1	Bellman Ford Negetive Cycle Detection	11	10	T D'	0.0
	2.11 star and bar	4		7.2	DFS Articulation Point	11	10	I.Bit	20
	2.11 Star and Dar	4		7.3	DFS Bridge Graph	12		10.1 Equation	
Q	C.Geometry	4		7.4	DFS LCA	12		10.2 Maximum And Pair	
J	v			7.5	DSU	13		10.3 Maximum XOR of all subsequence	
	3.1 Line And Point	4		7.6	Dijkstra by GM	13		10.4 Minimum XOR Operation	
	3.2 Points Inside Polygon	4		7.7	Eular Path on Undirected Graph	14		10.5 Number of Subarrays with XOR 0	
	3.3 Solid-Equations	5		7.8	Minimum Spanning Tree from Each Egde	14		10.6 Sum of XOR of All subset in Array	20

	10.7 Sum of all and of all subset	20							
11	J.Misc								
	11.1 Base conversion	21							
	11.2 Equation Solve	21							
	11.3 LCS - longest common subsequences	21							
	11.4 LIS - Longest Increasing Subsequence .	21							
	11.5 Mex of all Subarray	21							
	11.6 Mo's Algorithm	22							
	11.7 Nim Game 2d	23							
	11.8 Number of Subsegment Equal to K	23							
	11.9 Ternary Search	24							
	11.10Tower of Hanoi	24							

1 A.Settings and Script

1.1 Generator

```
#include <bits/stdc++.h>
using namespace std;
#define int long long
#define accuracy
chrono::steady_clock::now().
time_since_epoch().count()
#define rep(i, a, n) for (int i = a; i <= n; ++i)
#define nl
const int N = 1e6 + 4;
int32_t permutation[N];
mt19937 rng(accuracy);
int rand(int 1, int r)
   uniform_int_distribution<int> ludo(1, r);
   return ludo(rng);
}
const int inf = 1LL << 31;</pre>
using pii = pair<int, int>;
namespace generator
string gen_string(int len = 0, bool upperCase =
    false, int l = 1, int r = 26)
{
```

```
assert(len >= 0 && len <= 5e6);
   string str(len, (upperCase ? 'A' : 'a'));
   for (char &ch : str)
       ch += rand(1, r) - 1;
   return str;
vector<int> gen_array(int len = 0, int minRange =
    0, int maxRange = inf)
   assert(len >= 0 and len <= 5e6);
   vector<int> a(len);
   for (int &x : a)
       x = rand(minRange, maxRange);
   return a:
using namespace generator;
template <typename T = int> ostream
    &operator<<(ostream &other, const vector<T>
    &v)
   for (const T &x : v)
       other << x << ' ';
   other << '\n':
   return other;
#define SINGLE_TEST
const int max_tests = 1;
void generate_test()
signed main()
   srand(accuracy);
   int t = 1;
   #ifndef SINGLE_TEST
       t = rand(1, max_tests), cout << t << '\n';
   #endif
   while (t--)
       generate_test();
```

```
}
}
```

1.2 Main Script

```
g++ code.cpp -o code
g++ gen.cpp -o gen
g++ brute.cpp -o brute
for((i = 1; ; ++i)); do
   ./gen $i > input_file
   ./code < input_file > myAnswer
   ./brute < input_file > correctAnswer
   diff -Z myAnswer correctAnswer > /dev/null ||
        break
   echo "Passed test: " $i
done
echo "WA on the following test:"
cat input_file
echo "Your answer is:"
cat myAnswer
echo "Correct answer is:"
cat correctAnswer
```

1.3 Settings

```
// Layout -> Two column split row
// Wrap -> On, wrapping indent -> true
// Autosave -> After Delay
// Sticky scroll -> off
// gedit ~/.bashrc -> ulimit -s 2000123

// Command lines:
// generator (bash):
// chmod +x test.sh
// ./test.sh

// C++:
// g++ -std=c++17 -02 -Wno-unused-result
-static a.cpp -o x
```

```
// ./x
// g++ a.cpp -o x
// ./x
```

1.4 Template

```
#include "bits/stdc++.h"
#pragma GCC target ("avx2")
#pragma GCC optimization ("03")
#pragma GCC optimization ("unroll-loops")
using namespace std;
typedef long long int 11;
typedef unsigned long long 11;
typedef double db;
#define pi acos(-1)
#define prDouble(x, y) fixed << setprecision(y)</pre>
    << x
#define lcm(a, b) (a / \_gcd(a, b)) * b
#define fast io
    ios_base::sync_with_stdio(false); \
    cin.tie(NULL);
    cout.tie(NULL)
#define nl << "\n"
#define debug(x) cout << #x << " = " << x nl
#define casePrint(ans, cn) cout << "Case " << cn</pre>
    << ": " << ans nl
int getbit(long long n, long long bit) {
    return (1LL << bit) & n:
}
void setbit(long long &n, long long bit) {
   n |= (1LL << bit);</pre>
}
void unsetbit(long long &n, long long bit) {
    n &= ~(1LL << bit);
}
void solve(int caseNumber) {
}
```

```
void init_code() {
#ifndef ONLINE_JUDGE
   freopen("input.txt", "r", stdin);
   freopen("output.txt", "w", stdout);
   freopen("error.txt", "w", stderr);
#endif
}
signed main() {
   init_code();
   fast_io;
   int testcase = 1;
   // cin >> testcase;
   for (int caseNumber = 1: caseNumber <=</pre>
        testcase; ++caseNumber) {
       solve(caseNumber):
   }
   return 0;
```

2 B.Combinatorics

2.1 Binomial coefficient

2.2 Extended Euclidean Algorithm

```
ll exgcd(ll a, ll b, ll &x, ll &y)
{
    if(b == 0)
    {
        x = 1, y = 0;
        return a;
    }
    ll t = exgcd(b, a%b, y, x);
    y -= a / b * x;
    return t;
}
```

2.3 General Equation

- $\sum_{0 \le k \le n} {n-k \choose k} = Fib_{n+1}$
- \bullet $\binom{n}{k} = \binom{n}{n-k}$
- $\bullet \binom{n}{k} + \binom{n}{k+1} = \binom{n+1}{k+1}$
- $4 \cdot k \binom{n}{k} = n \binom{n-1}{k-1}$
- \bullet $\binom{n}{k} = \frac{n}{k} \binom{n-1}{k-1}$
- $\sum_{i=0}^{n} \binom{n}{i} = 2^n$
- $\sum_{i>0} \binom{n}{2i} = 2^{n-1}$
- $\sum_{i>0} \binom{n}{2i+1} = 2^{n-1}$
- $\sum_{i=0}^{k} (-1)^{i} {n \choose i} = (-1)^{k} {n-1 \choose k}$
- $\sum_{i=0}^{k} {n+i \choose i} = {n+k+1 \choose k}$
- $1\binom{n}{1} + 2\binom{n}{2} + 3\binom{n}{3} + \dots + n\binom{n}{n} = n2^{n-1}$
- $\sum_{i=0}^{n} C(n,i) = 2^{n}$
- $\sum_{i>0} C(n,2i) = 2^{n-1}$
- $\sum_{i>0} C(n,2i+1) = 2^{n-1}$

2.4 How Many Digit X^Y

```
Let, No. of Digits D. D
=floor [log_{10}(X^Y)] + 1
=floor [Y \times log_{10}(X)] + 1
```

2.5 How many digit in N!

```
 \begin{array}{c} \text{Let, Number of Digits D} \\ \text{D=floor}[log_{10}(N!)] + 1 \\ = \text{floor}[log_{10}(1 \times 2 \times 3 \times 4. \dots (N-1) \times N)] + 1 \\ = \text{floor}[log_{10}(1) + log_{10}(2) + \dots + \\ log_{10}(N)] + 1 \end{array}
```

2.6 Last Non Zero Digit of Factorial

2.7 Matrix Exponentitation

```
REP(i, N) REP(j, N) B[i][j] = R[i][j];
   return:
}
void MatPow(LL R[N][N],LL M[N][N],LL P) {
   while(P) {
       if(P & 1) MatMul(M,R);
       MatMul(M,M);
       P = P >> 1:
   }
}
int main() {
   LL n,M[N][N],R[N][N]; // M is Co-efficient
        Matrix, R is Base case Matrix
   //Take input values of M and R matrix
   //Input n, We have to find f(n)
   MatPow(R,M,n-2); // Here n-2 may changes in
        diffrent problems
   //value of f(n) is in R[0][0] position
   return 0:
```

2.8 Modular Inverse

```
int binExpIter(int a, int b)
{
    int ans = 1;
    while (b)
    {
        if (b & 1)
            ans = (ans * 1LL * a) % MOD;
        a = (a * 1LL * a) % MOD;
        b >>= 1;
    }
    return ans;
}

int MMI(int n)
{
    return binExpIter(n, MOD - 2);
}
```

2.9 combination1

```
#include<bits/stdc++.h>
using namespace std;
const int N = 1e6, mod = 1e9 + 7;
int power(long long n, long long k) {
  int ans = 1 % mod; n %= mod; if (n < 0) n +=</pre>
      mod;
  while (k) {
   if (k \& 1) ans = (long long) ans * n % mod;
   n = (long long) n * n % mod;
   k >>= 1:
 return ans:
int f[N], invf[N];
int nCr(int n, int r) {
 if (n < r \text{ or } n < 0) return 0;
 return 1LL * f[n] * invf[r] % mod * invf[n - r]
      % mod:
}
int nPr(int n, int r) {
 if (n < r \text{ or } n < 0) return 0;
 return 1LL * f[n] * invf[n - r] % mod;
int32_t main() {
 ios_base::sync_with_stdio(0);
 cin.tie(0);
 f[0] = 1;
 for (int i = 1; i < N; i++) {</pre>
   f[i] = 1LL * i * f[i - 1] % mod;
  invf[N-1] = power(f[N-1], mod - 2);
  for (int i = N - 2; i \ge 0; i--) {
   invf[i] = 1LL * invf[i + 1] * (i + 1) % mod:
  cout << nCr(6, 2) << '\n';
 cout << nPr(6, 2) << '\n';
 return 0;
```

2.10 nCr with Big Mod

```
11 Big_Mod(ll a, ll b)
    if(!b) return 1;
    ll ans = Big_Mod(a, b/2);
    ans %= mod;ans *= ans;ans %= mod;
    if(b&1)
       a %= mod; ans *= a; ans %= mod;
    return ans;
ll nCr(ll n, ll k)
    return fact(n)*111*Big_Mod(fact(k),
        mod-2)%mod*1ll*Big_Mod(fact(n-k),
        mod-2)%mod;
```

2.11 star and bar

```
1. x_1 + x_2 + \cdots + x_r = n where x_i >
   0. Solution: \binom{n-1}{n-1}
2. x_1 + x_2 + \cdots + x_r = n where x_i \geq 0. Solution: \binom{n+r-1}{r-1}
3. x_1 + x_2 + \cdots + x_r = n where x_i \geq
   b. Solution: \binom{n-rb+r-1}{r-1}
```

C.Geometry

3.1 Line And Point

```
struct pt {
   double x, y;
struct line {
   double a, b, c;
```

```
};
const double EPS = 1e-9;
// determinant of a 2x2 matrix
double det(double a, double b, double c, double
   return a*d - b*c:
bool intersect(line m, line n, pt & res) {
   double zn = det(m.a, m.b, n.a, n.b);
   if (abs(zn) < EPS)
       return false;
   res.x = -det(m.c, m.b, n.c, n.b) / zn;
   res.v = -det(m.a, m.c, n.a, n.c) / zn;
   return true:
bool parallel(line m, line n) {
   return abs(det(m.a, m.b, n.a, n.b)) < EPS;</pre>
bool equivalent(line m, line n) {
   return abs(det(m.a, m.b, n.a, n.b)) < EPS</pre>
          && abs(det(m.a, m.c, n.a, n.c)) < EPS
          && abs(det(m.b. m.c. n.b. n.c)) < EPS:
```

3.2 Points Inside Polygon

```
//points inside convex polygon O(logn)
const 11 N = 100009;
struct point {
   11 x, y;
}a[N];
double cross(const point& p1,const point&
    p2, const point& org) {
   return
        ((p1.x-org.x)*1.0)*(p2.y-org.y)-((p2.x-org.x)*1.0)*(p1.y-org.y);
inline bool comp(const point& x,const point& y) {
   return cross(x,y,a[0]) >= 0;
bool inside(point& p) {
   if (cross(a[0], a[n-1], p)>=0) return false;
```

```
if (cross(a[0], a[1], p) <=0) return false;</pre>
   11 1 =1;
   11 r = n-1;
   while(l<r) {</pre>
       11 m = 1 + (r-1)/2;
       if(cross(a[m],p,a[0]) >=0)
           1=m+1;
       else
           r=m;
   if(1 == 0)
       return false;
   return cross(a[1-1],a[1],p) >0;
sort(a+1,a+n,comp);
```

3.3 Solid-Equations

1. Rectangular Parallelepiped

- Area: 2(ab + bc + ca)
- Volume: abc
- Diagonal: $\sqrt{a^2+b^2+c^2}$

2. Cube

- Total surface area: $6a^2$
- Volume: a^3
- Diagonal: $\sqrt{3} \cdot a$

• Total surface area: 2(base area) + $(perimeter of base) \times (height)$

3. Prism

• Volume: (base area) \times (height)

4. Pyramid

- Total surface area: (base area) + $\frac{1}{2}$ × (perimeter of base) × (slant height) (for a right regular pyramid)
- Slant height: $\sqrt{h^2 + r^2}$ (for a right regular pyramid with base inscribed in a circle of radius r)
- Volume: $\frac{1}{3} \times (\text{base area}) \times (\text{height})$

5. Right Circular Cone

- Surface area: $\pi r(l+r)$
- Volume: $\frac{1}{3} \times (\text{base area}) \times (\text{height})$

6. Sphere

- Surface area: $4\pi r^2$
- Volume: $\frac{4}{3}\pi r^3$
- Radius of the circle formed by a section at a distance h from the center: $r_h = \sqrt{r^2 h^2}$

7. Cylinder

- Total surface area: $2\pi r(r+h)$
- Curved surface area: $2\pi rh$
- Volume: $\pi r^2 h$

8. Cone

- \bullet Curved surface area: $\pi r l$
- Total surface area: $\pi r(r+l)$
- Volume: $\frac{1}{3}\pi r^2 h$

9. Frustum of a Cone

- Curved surface area: $\pi(r_1 + r_2)l$
- Total surface area: $\pi(r_1 + r_2)l + \pi r_1^2 + \pi r_2^2$
- Volume: $\frac{1}{3}\pi h(r_1^2 + r_2^2 + r_1r_2)$

10. Hemisphere

- Curved surface area: $2\pi r^2$
- Total surface area: $3\pi r^2$
- Volume: $\frac{2}{3}\pi r^3$

11. Torus

- Surface area: $4\pi^2 Rr$
- Volume: $2\pi^2 Rr^2$

12. Ellipsoid

- Surface area: (Complex formula, approximation: $4\pi \left(\frac{a^pb^p+a^pc^p+b^pc^p}{3}\right)^{1/p}$, where $p\approx 1.6075$)
- Volume: $\frac{4}{3}\pi abc$

4 D.Number Theory

4.1 1 - N Divisor

```
// Number of Divisor
for (int i = 1; i * i <= n; i++) {
   for (int j = i * i; j <= n; j += i) {
      if ( i * i == j) a[j]++;
      else a[j] += 2;
   }
}
// Note: When we find the remainder of a % m, the
   answer will their GCD (Greatest Common
   Divisor). So, a % m = GCD (a, m)</pre>
```

```
// Sum of Divisor
for ( int i = 1; i * i <= n; i++) {
   for (int j = i*i; j < n; j += i) {</pre>
       if(j == i*i) a[j] += i;
       else a[i] += i + (i / i);
   }
}
long long SumOfDivisors(long long num) {
   long long total = 1;
   for (int i = 2; (long long)i * i \le num; i++)
       if (num % i == 0) {
           int e = 0;
           do {
               num /= i;
           } while (num % i == 0);
           long long sum = 0, pow = 1;
           do {
               sum += pow;
              pow *= i;
           } while (e-- > 0);
           total *= sum;
       }
   if (num > 1) {
       total *= (1 + num);
   return total;
}
long long numberOfDivisors(long long num) {
   long long total = 1;
   for (int i = 2; (long long)i * i <= num; i++)</pre>
       if (num % i == 0) {
           int e = 0;
           do {
               e++:
              num /= i;
           } while (num % i == 0);
```

```
total *= e + 1;
}

if (num > 1) {
    total *= 2;
}
return total;
}
```

4.2 Catalan Numbers

The Catalan sequence is the sequence $C_0, C_1, C_2, ...,$ where $C_0 = 1, C_1 = 1$, and

$$C_n = C_0 C_{n-1} + C_1 C_{n-2} + \ldots + C_{n-1} C_0, \quad n \ge 2.$$

Therefore,

$$C_n = \frac{C(2n, n)}{n+1}, \quad n = 0, 1, 2, \dots;$$

$$C_n = \frac{4n-2}{n+1} \times C_{n-1}, \quad n > 1.$$

The Catalan sequence is a frequent counting sequence. For example:

- C_n is the number of stack-sortable permutations of $\{1, \ldots, n\}$.
- C_n is the number of different ways that a convex polygon with n+2 sides can be cut into triangles by connecting vertices with non-crossing line segments.
- C_n is the number of rooted binary trees with n nodes.

4.3 Eular Totient Of Every Number 1-N

```
int euler_phi [mxm];
void Euler_Totient (int n) {
    // Euler Totient of number 1 to n
        euler_phi[1] = 1;
    for(int i = 2; i <= n; i++) {
        if(euler_phi[i] > 0) continue;
        for(int j = i + i; j <= n; j += i) {
            if(euler_phi[j] == 0) euler_phi[j] = j;
            euler_ph[j] -= (euler_phi[j] / i);
        }
    }
}
// Note: Normally euler totient returns the
    amount of number which are co-prime with n.</pre>
```

4.4 GCD LCM Equations

- gcd(a,0) = a
- $gcd(a, b) = gcd(b, a \mod b)$
- Every common divisor of a and b is a divisor of gcd(a, b).
- If m is any integer, then gcd(a+mb, b) = gcd(a, b)
- The gcd is a multiplicative function in the following sense: if a_1 and a_2 are relatively prime, then $gcd(a_1a_2, b) = gcd(a_1, b) gcd(a_2, b)$.
- $gcd(a, b) \cdot lcm[a, b] = |ab|$
- gcd(a, lcm[b, c]) = lcm[gcd(a, b), gcd(a, c)]
- $\operatorname{lcm}[a, \gcd(b, c)] = \gcd(\operatorname{lcm}[a, b], \operatorname{lcm}[a, c)]$
- For non-negative integers a and b, where a and b are not both zero,

$$\gcd(n1, n1) = n\gcd(a, b)$$

 $\bullet \ \gcd(a,b) = \frac{kla \cdot kib}{k}$

- $gcd(i, n) = k] = \frac{\varphi(n)}{k} \sum_{i=1}^{n} \alpha(i)$
- $gcd(k,n) = \sum_{d|n} \alpha(d)\varphi(\frac{n}{d})$
- $\sum_{k=1}^{n} \frac{1}{\gcd(k,n)} = \sum_{d|n} \frac{1}{d} \cdot \phi(\frac{n}{d}) = \frac{1}{n} \sum_{d|n} d \cdot \phi(d)$
- $\bullet \sum_{k=1}^{n} \frac{k}{\gcd(k,n)} = \frac{n}{2} \cdot \sum_{d|n} \frac{1}{d} \cdot \phi(\frac{n}{d}) = \frac{n}{2} \cdot \frac{1}{n} \cdot \sum_{d|n} d \cdot \phi(d)$
- $\sum_{k=1}^{n} \frac{n}{\gcd(k,n)} = 2 * \sum_{k=1}^{n} \frac{k}{\gcd(k,n)} 1$, for n > 1
- $\sum_{i=1}^{n} \sum_{j=1}^{n} [gcd(i,j) = 1] = \sum_{d=1}^{n} \mu(d) \lfloor \frac{n}{d} \rfloor^2$
- $\sum_{i=1}^{n} \sum_{j=1}^{n} gcd(i,j) = \sum_{d=1}^{n} \phi(d) \lfloor \frac{n}{d} \rfloor^2$
- $\sum_{i=1}^{n} \sum_{j=1}^{n} i \cdot j[gcd(i,j) = 1] = \sum_{i=1}^{n} \phi(i)i^2$
- $F(n) = \sum_{i=1}^{n} \sum_{j=1}^{n} lcm(i,j) = \sum_{l=1}^{n} \left(\frac{(1+\lfloor \frac{n}{l} \rfloor)(\lfloor \frac{n}{l} \rfloor)}{2}\right)^{2} \sum_{d|l} \mu(d)ld$
- gcd(lcm(a,b), lcm(b,c), lcm(a,c)) = lcm(gcd(a,b), gcd(b,c), gcd(a,c))
- $gcd(A_L, A_{L+1}, ..., A_R) = gcd(A_L, A_{L+1} A_L, ..., A_R A_{R-1})$
- Givenn, IfSUM = LCM(1, n) + LCM(2, n) + ... + LCM(n, n) thenSUM = n2((d)xd) + 1
- xed(k) = x.()

4.5 LOG with any base

```
// find log with any base
double LOG(double base, double num)
{
    return log(num) / log(base);
}
```

4.6 Large-number-fibonacci

```
const 11 M = 1e9 + 7;
#define vpll vector<pair<11, 11>>
#define pll pair<11, 11>
vector<pair<11, 11>> base;
vpll mul(vpll a, vpll b) {
 vpll c;
 ll a1 = (a[0].first * b[0].first + a[0].second
      * b[1].first) % M:
 11 \ a2 = (a[0].first * b[0].second + a[0].second
      * b[1].second) % M:
 ll b1 = (a[1].first * b[0].first + a[1].second
      * b[1].first) % M;
 11 b2 = (a[1].first * b[0].second + a[1].second
      * b[1].second) % M;
 pll x = \{a1, a2\};
 pll y = \{b1, b2\};
 c.push_back(x);
 c.push_back(y);
 return c;
vpll power(vpll a, ll b) {
 if (b == 1) return base;
 vpll r = power(a, b / 2);
 r = mul(r, r);
 if (b \% 2) r = mul(r, a);
 return r;
}
void fib(vpll vp) { cout << vp[0].first +</pre>
    vp[0].second << endl; }</pre>
int main() {
 11 n:
 cin >> n;
 base.push_back(\{1, 1\});
 base.push_back({1, 0});
```

```
vpll ans = power(base, n - 2);
fib(ans);
return 0;
}
```

4.7 Sieve of Eratosthenes

```
// remember to call the function
vector<bool> prime(1e7+10, true);

void SieveOfEratosthenes(){
   prime[0]=prime[1]=false;
   for (int p = 2; p * p <= 1e7+10; p++) {
      if (prime[p] == true) {
        for (int i = p * p; i <= 1e7+10; i += p)
            prime[i] = false;
      }
   }
}</pre>
```

4.8 Total Digit of N factorial

5 Data Structures

5.1 MO's algo coordinate compression

```
vector<int>compressed, a;
const int N = 2e5 + 8;
vector<int>cnt(N, 0);
int cur_result = 0;
struct query {
   int 1, r, idx;
};
int block;
void add(int i) {
```

```
cnt[a[i]]++;
   if(cnt[a[i]] == 1) {
       cur_result++;
void Remove(int i) {
   cnt[a[i]]--;
   if(cnt[a[i]] == 0) {
       cur_result--;
bool comp1(query p , query q) {
   if (p.1 / block != q.1 / block) {
       if (p.1 == q.1) return p.r < q.r;</pre>
       return p.l < q.l;</pre>
   return (p.l / block & 1) ? (p.r < q.r) : (p.r
        > q.r);
void mos_algorithm(int n, vector<query>&queries) {
   vector<int>answers(queries.size());
   block = (int)sqrt(n);
   sort(queries.begin(), queries.end(), comp1);
   int cur_1 = 0;
   int cur_r = -1;
   for (query q : queries) {
       while (cur_1 > q.1) {
           cur 1--:
           add(cur_1);
       while (cur_r < q.r) {</pre>
           cur_r++;
           add(cur_r);
       while (cur_1 < q.1) {</pre>
           Remove(cur_1);
           cur_1++;
       while (cur_r > q.r) {
           Remove(cur_r);
           cur_r--;
       answers[q.idx] = cur_result;
```

```
for (int i : answers) {
       cout << i << nl:
   }
void solve() {
   int n, q; cin >> n >> q;
   a.assign(n, 0);
   for (int i = 0; i < n; i++) {</pre>
       cin >> a[i]:
   vector<query>queries(q);
   for(int i = 0; i < q; i++) {</pre>
       cin >> queries[i].1;
       queries[i].1--;
       cin >> queries[i].r;
       queries[i].r--;
       queries[i].idx = i;
   }
   compressed = a;
   sort(compressed.begin(), compressed.end());
   compressed.resize(unique(compressed.begin(),
        compressed.end()) - compressed.begin());
   for (int i = 0; i < n; i++) {</pre>
       a[i] = lower_bound(compressed.begin(),
            compressed.end(), a[i]) -
            compressed.begin() + 1;
   }
   mos_algorithm(n, queries);
}
int main() {
   fast_IO;
   int t = 1;
   // cin >> t;
   while (t--) {
       solve();
   }
   return 0;
```

5.2 Segment Tree Lazy

```
struct node
{
 ll sum;
 11 setAll;
 11 increment;
 bool setAllValid;
 node()
   sum = 0:
   setAllValid = 0;
   increment = 0;
 void Reset()
   setAllValid = increment = 0;
 }
};
// 0-based indexing
class segtree
public:
 int range;
 vector<node> tree;
 void build(vector<int> &v)
   range = v.size();
   tree.assign(4 * range, node());
   build_recur(v, 0, range - 1, 0);
 void build_recur(vector<int> &v, int 1, int r,
      int node no)
   if (1 == r)
     if (1 < v.size())</pre>
       tree[node_no].sum = v[1];
       tree[node_no].sum = 0;
     return;
   }
   int mid = (1 + r) / 2:
   build_recur(v, 1, mid, 2 * node_no + 1);
   build_recur(v, mid + 1, r, 2 * node_no + 2);
```

```
tree[node no].sum = tree[2 * node no + 1].sum
      + tree[2 * node_no + 2].sum;
11 range_query(int L, int R)
 return range_query_recur(0, 0, range - 1, L,
     R);
void incUpdate_recur(int node, int 1, int r,
    int &L, int &R, int &X)
 if (r < L || R < 1 || 1 >= range)
   return;
 if (L <= 1 && R >= r)
   tree[node].increment += X;
   return;
 applyAggr(node, 1, r);
 int mid = (1 + r) / 2;
 incUpdate_recur(2 * node + 1, 1, mid, L, R,
 incUpdate_recur(2 * node + 2, mid + 1, r, L,
     R, X);
 applyAggr(2 * node + 1, 1, mid);
 applyAggr(2 * node + 2, mid + 1, r);
 tree[node].sum = tree[2 * node + 1].sum +
      tree[2 * node + 2].sum;
void incUpdate(int L, int R, int X)
 incUpdate_recur(0, 0, range - 1, L, R, X);
void setUpdate_recur(int node, int 1, int r,
    int &L. int &R. int &X)
 if (r < L || R < 1 || 1 >= range)
   return;
 if (L <= 1 && R >= r)
   tree[node].setAllValid = 1;
   tree[node].setAll = X;
   tree[node].increment = 0;
   return;
```

```
applyAggr(node, 1, r);
 int mid = (1 + r) / 2;
 setUpdate_recur(2 * node + 1, 1, mid, L, R,
 setUpdate_recur(2 * node + 2, mid + 1, r, L,
      R, X);
 applyAggr(2 * node + 1, 1, mid);
 applyAggr(2 * node + 2, mid + 1, r);
 tree[node].sum = tree[2 * node + 1].sum +
      tree[2 * node + 2].sum;
}
void setUpdate(int L, int R, int X)
 setUpdate_recur(0, 0, range - 1, L, R, X);
void compose(int par, int child)
 if (tree[par].setAllValid)
   tree[child].setAllValid = 1;
   tree[child].setAll = tree[par].setAll;
   tree[child].increment = tree[par].increment;
 else
   tree[child].increment += tree[par].increment;
void applyAggr(int node, int 1, int r)
 if (tree[node].setAllValid)
   tree[node].sum = (r - l + 1) *
       tree[node].setAll:
 tree[node].sum += (r - l + 1) *
      tree[node].increment;
 if (1 != r)
 {
   compose(node, 2 * node + 1);
   compose(node, 2 * node + 2);
 tree[node].Reset();
ll range_query_recur(int node, int 1, int r,
    int &L, int &R)
{
```

```
if (r < L || R < 1 || 1 >= range)
    return 0;
applyAggr(node, 1, r);
if (L <= 1 && R >= r)
    return tree[node].sum;
int mid = (1 + r) / 2;
return range_query_recur(2 * node + 1, 1,
    mid, L, R) + range_query_recur(2 * node +
    2, mid + 1, r, L, R);
};
```

5.3 Segment Tree

```
class SegmentTree
public:
   vector<vector<int>> tree; // 0-mn, 1-mx, 2-sum
   SegmentTree(int n)
   {
       tree.assign(3, vector<int>(4 * n + 1, 0));
   void build(vector<int> &a, int 1, int r, int
       node)
   {
       if (1 > r)
          return:
       if (1 == r)
          tree[0][node] = a[1]:
          tree[1][node] = a[1];
          tree[2][node] = a[1];
          return;
       int mid = (1 + r) / 2;
       build(a, 1, mid, 2 * node);
       build(a, mid + 1, r, 2 * node + 1);
       tree[0][node] = min(tree[0][2 * node],
           tree[0][2 * node + 1]);
       tree[1][node] = max(tree[1][2 * node],
           tree[1][2 * node + 1]);
```

```
tree[2][node] = tree[2][2 * node] +
        tree[2][2 * node + 1];
}
void update(int node, int 1, int r, int idx,
    int value)
   if (1 > r)
       return:
    if (1 == r)
       tree[0][node] = value;
       tree[1][node] = value;
       tree[2][node] = value:
       return;
   }
    int mid = (1 + r) / 2;
    if (idx >= 1 and idx <= mid)</pre>
       update(2 * node, 1, mid, idx, value);
    else
       update(2 * node + 1, mid + 1, r, idx,
            value);
    tree[0][node] = min(tree[0][2 * node],
        tree[0][2 * node + 1]);
   tree[1][node] = max(tree[1][2 * node].
        tree[1][2 * node + 1]);
    tree[2][node] = tree[2][2 * node] +
        tree[2][2 * node + 1];
}
// b=begin, e=end
int query_mn(int node, int 1, int r, int b,
    int e)
   if (1 \ge b \text{ and } r \le e)
       return tree[0][node];
   if (1 > e \text{ or } r < b)
       return 1e18; // return invalid value;
    int mid = (1 + r) / 2;
```

```
int left = query_mn(2 * node, 1, mid, b,
       int right = query_mn(2 * node + 1, mid +
            1, r, b, e);
       return min(left, right);
   }
   int query_mx(int node, int 1, int r, int b,
        int e)
   {
       if (1 \ge b \text{ and } r \le e)
           return tree[1][node];
       if (1 > e \text{ or } r < b)
           return -1e18; // return invalid value;
       int mid = (1 + r) / 2:
       int left = query_mx(2 * node, 1, mid, b,
       int right = query_mx(2 * node + 1, mid +
            1, r, b, e);
       return max(left, right);
   int query_sum(int node, int 1, int r, int b,
        int e)
   {
       if (1 \ge b \text{ and } r \le e)
           return tree[2][node];
       if (1 > e \text{ or } r < b)
           return 0; // return invalid value;
       int mid = (1 + r) / 2;
       int left = query_sum(2 * node, 1, mid, b,
            e);
       int right = query_sum(2 * node + 1, mid +
            1, r, b, e);
       return left + right;
   }
};
```

5.4 ordered set

```
// Ordered Set
#include<bits/stdc++.h>
#include<ext/pb_ds/assoc_container.hpp>
#include<ext/pb_ds/tree_policy.hpp>
```

```
using namespace __gnu_pbds;
using namespace std;
template <typename T> using o_set = tree<T,</pre>
    null_type, less<T>,rb_tree_tag,
    tree_order_statistics_node_update>;
int32_t main() {
   ios_base::sync_with_stdio(0);
   cin.tie(0);
   o set<int> se:
   cout << se.order_of_key(5) << '\n'; // number</pre>
        of elements < 5
   se.erase(3); // erase by value
   cout << (*se.find_by_order(1)) << '\n'; // if</pre>
        you imagine this as a 0-indexed vector,
        what is se[1]?
   return 0;
```

6 E.DP

6.1 Digit DP

```
int len,id,inp[10];
LL dp[10][10][2][2];
int vis[10][10][2][2];
LL call(int pos, bool isSmall, bool isStart, int
    total) {
   if(pos==len) return total;
    if(vis[pos][total][isSmall][isStart]==id)
        return dp[pos][total][isSmall][isStart];
   vis[pos][total][isSmall][isStart]=id;
    int last=9;
   if(!isSmall) last=inp[pos];
   LL ret=0;
   if(isStart) {
       nfor(int i=0; i<=last; i++) {</pre>
           ret+=call(pos+1,isSmall |
                i<inp[pos],1,(i==0)+total);</pre>
       }
   }
```

```
else {
       for(int i=1; i<=last; i++) {</pre>
           ret+=call(pos+1,isSmall |
                i<inp[pos],1,(i==0)+total);</pre>
       ret+=call(pos+1,1,0,0);
    return dp[pos][total][isSmall][isStart]=ret;
LL solve(LL x) {
    if(x<0) return 0;</pre>
    len=0;
    while(x) {
       inp[len++]=x%10;
       x/=10;
    reverse(inp,inp+len);
    return call(0,0,0,0)+1;
int main() {
    int t;
    scanf("%d",&t);
    FOR(tc,1,t) {
       LL lo,hi;
       scanf("%lld %lld",&lo,&hi);
       printf("Case %d:
            %1ld\n",tc,solve(hi)-solve(lo-1));
   }
```

6.2 Knapsack 1

```
int n, w; cin>>n>w;
vector<int> dp(w+1, 0);
for(int i=0; i<n; i++){
   int weight, val; cin>>weight>> val;
   for(int prev_weight=w-weight; prev_weight>=0;
        prev_weight--){
        dp[prev_weight+weight] =
            max(dp[prev_weight+weight],
            dp[prev_weight]+val);
}
```

```
cout<<dp[w]<<nl;</pre>
```

int longest(int n) {

6.3 Longest Incresing Subsequence

```
vector<int>v:
   for(int i=1; i<=n; i++) {</pre>
       int pos = lower_bound(v.begin(), v.end(),
           a[i]) - v.begin();
       if(pos == v.size()) v.push_back(a[i]);
       else v[pos] = a[i];
   return v.size();
}
// longest common subsequence
#include<bits/stdc++.h>
using namespace std;
const int N = 3030;
string a, b;
int dp[N][N];
int lcs(int i, int j) {
 if (i >= a.size() or j >= b.size()) return 0;
 if (dp[i][j] != -1) return dp[i][j];
 int ans = lcs(i + 1, j);
 ans = max(ans, lcs(i, j + 1));
 if (a[i] == b[i]) {
   ans = \max(ans, lcs(i + 1, j + 1) + 1);
 return dp[i][j] = ans;
void print(int i, int j) {
 if (i >= a.size() or j >= b.size()) return;
 if (a[i] == b[i]) {
   cout << a[i];
   print(i + 1, j + 1);
   return:
 int x = lcs(i + 1, j);
  int y = lcs(i, j + 1);
```

```
if (x >= y) {
    print(i + 1, j);
}
else {
    print(i, j + 1);
}
int32_t main() {
    ios_base::sync_with_stdio(0);
    cin.tie(0);
    cin >> a >> b;
    memset(dp, -1, sizeof dp);
    // cout << lcs(0, 0) << '\n';
    print(0, 0);
    return 0;
}</pre>
```

7 F.Graph

7.1 Bellman Ford Negetive Cycle Detection

```
struct node {
   int u;
   int v;
   int wt:
   node(int first, int second, int weight)
       u = first:
       v = second;
       wt = weight;
   }
const int inf = 10000000;
const int SZ = 1e6+5;
int dist[SZ];
vector<node> edges;
void bellmenFord(int n) {
   for(int i = 0; i<=n-1; i++) {</pre>
       for(auto it: edges) {
           if(dist[it.u] + it.wt < dist[it.v]) {</pre>
```

7.2 DFS Articulation Point

```
const 11 SZ = 1e6 + 5:
11 tin[SZ], low[SZ], vis[SZ], isArticulation[SZ];
vector<ll> adj[SZ];
ll timer:
void dfs(ll node, ll parent, ll timer) {
   vis[node] = 1;
   tin[node] = low[node] = timer++;
   11 childCnt = 0;
   for (auto child : adj[node]) {
       if (child == parent) {
          continue;
       if (vis[child] == 0) {
          dfs(child, node, timer);
          low[node] = min(low[node], low[child]);
          if (low[child] >= tin[node] && parent
               != -1) {
              isArticulation[node] = 1;
       }
       else {
          low[node] = min(low[node], tin[child]);
   if (parent == -1 && childCnt > 1) {
       isArticulation[node] = 1;
```

```
}
}
void init(ll n) {
    mem(vis, 0);
    mem(isArticulation, 0);
    mem(tin, -1);
    mem(low, -1);
    for (int i = 1; i <= n; i++) {
        adj[i].clear();
    }
}
//dfs call -> dfs(i, -1, timer);
```

7.3 DFS Bridge Graph

```
const 11 SZ = 1e6 + 5;
11 tin[SZ], low[SZ], vis[SZ];
vector<11> adj[SZ];
11 timer;
void dfs(ll node, ll parent, ll timer) {
   vis[node] = 1;
   tin[node] = low[node] = timer++;
   for (auto child : adj[node]) {
       if (child == parent) continue;
       if (vis[child] == 0) {
           dfs(child, node, timer);
           low[node] = min(low[node], low[child]);
           if (low[child] > tin[node])
               vp.push_back({min(child, node),
               max(node, child)});
       }
       else low[node] = min(low[node],
           tin[child]);
void init()
   memset(vis, 0, sizeof(vis));
   memset(tin, -1, sizeof(tin));
   memset(low, -1, sizeof(low));
   for (int i = 0; i < n; i++) adj[i].clear();</pre>
```

```
}
// dfs call -> dfs(i, -1, timer);
```

7.4 DFS LCA

```
const int N = 3e5 + 9, LG = 18;
vector<int> g[N]:
int par[N][LG + 1], dep[N], sz[N];
// Call at first
void dfs(int u, int p = 0) {
 par[u][0] = p;
 dep[u] = dep[p] + 1;
 sz[u] = 1;
 for (int i = 1; i <= LG; i++) par[u][i] =</pre>
      par[par[u][i - 1]][i - 1];
 for (auto v: g[u]) if (v != p) {
   dfs(v. u):
   sz[u] += sz[v];
 }
// Lowest Common Ancestor
int lca(int u, int v) {
 if (dep[u] < dep[v]) swap(u, v);
 for (int k = LG; k \ge 0; k--) if
      (dep[par[u][k]] >= dep[v]) u = par[u][k];
 if (u == v) return u;
 for (int k = LG; k >= 0; k--) if (par[u][k] !=
      par[v][k]) u = par[u][k], v = par[v][k];
 return par[u][0];
}
// K-th root from u
int kth(int u, int k) {
 assert(k >= 0);
 for (int i = 0; i <= LG; i++) if (k & (1 << i))
      u = par[u][i];
 return u;
// Calculate distance between u and v
int dist(int u, int v) {
 int 1 = lca(u, v);
 return dep[u] + dep[v] - (dep[1] << 1);</pre>
```

```
//kth node from u to v. Oth node is u
int go(int u, int v, int k) {
 int 1 = lca(u, v);
 int d = dep[u] + dep[v] - (dep[1] << 1);
 assert(k <= d);
 if (dep[1] + k <= dep[u]) return kth(u, k);</pre>
 k \rightarrow dep[u] - dep[1];
 return kth(v, dep[v] - dep[1] - k);
int32 t main() {
 int n; cin >> n;
 for (int i = 1; i < n; i++) {</pre>
   int u, v; cin >> u >> v;
   g[u].push_back(v);
   g[v].push_back(u);
  dfs(1);
  int q; cin >> q;
  while (q--) {
   int u, v; cin >> u >> v;
   cout << dist(u, v) << '\n';
 return 0;
```

7.5 DSU

```
// 1-based indexing
class DSU
{
public:
    vector<int> parent, rating;
    int cc_count;
    DSU(int n)
    {
        cc_count = n;
        parent.assign(n + 10, 0);
        rating.assign(n + 1, 1);
        for (int i = 1; i <= n; i++)
        {
            parent[i] = i;
            rating[i] = 1;
        }
}</pre>
```

```
}
}
int find_parent(int node)
   return parent[node] = (parent[node] ==
       node ? node :
       find_parent(parent[node]));
}
bool is_in_same(int u, int v)
   int par1 = find_parent(u);
   int par2 = find_parent(v);
   if (par1 == par2)
       return true;
   }
   return false;
// Returns size of component
int get_size(int u) { return
    rating[find_parent(u)]; }
// False means already connected
bool merge(int u, int v)
   if (is_in_same(u, v))
       return false;
   int par1 = find_parent(u);
   int par2 = find_parent(v);
   if (rating[par1] >= rating[par2])
       parent[par2] = par1;
       rating[par1] += rating[par2];
   }
   else
       parent[par1] = par2;
       rating[par2] += rating[par1];
   --cc_count;
   return true;
```

}

7.6 Dijkstra by GM

```
#include<bits/stdc++.h>
using namespace std;
const int N = 3e5 + 9, mod = 998244353;
int n. m:
vector<pair<int, int>> g[N], r[N];
vector<long long> dijkstra(int s, int t,
    vector<int> &cnt) {
   const long long inf = 1e18;
   priority_queue<pair<long long, int>,
       vector<pair<long long, int>>,
       greater<pair<long long, int>>> q;
   vector<long long> d(n + 1, inf);
   vector < bool > vis(n + 1, 0);
   q.push({0, s});
   d[s] = 0;
   cnt.resize(n + 1, 0); // number of shortest
        paths
   cnt[s] = 1;
   while(!q.empty()) {
       auto x = q.top();
       q.pop();
       int u = x.second;
       if(vis[u]) continue;
       vis[u] = 1;
       for(auto y: g[u]) {
       int v = v.first;
       long long w = y.second;
       if(d[u] + w < d[v]) {
          d[v] = d[u] + w;
          q.push({d[v], v});
           cnt[v] = cnt[u];
       } else if(d[u] + w == d[v]) cnt[v] =
           (cnt[v] + cnt[u]) % mod;
       }
   }
   return d;
int u[N], v[N], w[N];
```

```
int32 t main() {
   ios_base::sync_with_stdio(0);
   cin.tie(0);
   int s, t;
   cin >> n >> m >> s >> t;
   for(int i = 1; i <= m; i++) {</pre>
       cin >> u[i] >> v[i] >> w[i]:
       g[u[i]].push_back({v[i], w[i]});
       r[v[i]].push_back({u[i], w[i]});
   vector<int> cnt1, cnt2;
   auto d1 = dijkstra(s, t, cnt1);
   auto d2 = dijkstra(t, s, cnt2);
   long long ans = d1[t];
   for(int i = 1; i <= m; i++) {</pre>
       int x = u[i], y = v[i];
       long long nw = d1[x] + w[i] + d2[y];
       if(nw == ans && 1LL * cnt1[x] * cnt2[y] %
            mod == cnt1[t]) cout << "YES\n";</pre>
       else if(nw - ans + 1 < w[i]) cout << "CAN"</pre>
            " << nw - ans + 1 << ^{\prime}\n';
       else cout << "NO\n";</pre>
   return 0;
}
```

7.7 Eular Path on Undirected Graph

```
/*
all the edges should be in the same connected
    component
#undirected graph: euler path: all degrees are
    even or exactly two of them are odd.
#undirected graph: euler circuit: all degrees are
    even
*
find_euler_tour(u):
    for each node v in adj[u]
        remove edge u, v
        find_euler_tour(b)
```

```
push(u) // Stored in reverse order
//euler path in an undirected graph
//it also finds circuit if it exists
vector<pair<int, int>> g[N];
vector<int> ans:
int done[N];
int vis[N * N]; //number of edges
void dfs(int u) {
 while (done[u] < g[u].size()) {</pre>
   auto e = g[u][done[u]++];
   if (vis[e.second]) continue;
   vis[e.second] = 1;
   dfs(e.first):
 ans.push_back(u);
}
int solve(int n) {
 int edges = 0;
 ans.clear();
 memset(done, 0, sizeof done);
 memset(vis, 0, sizeof vis);
 vector < int > deg(n + 1, 0);
 for (int u = 1; u <= n; u++) {
   for (auto e : g[u]) {
     deg[e.first]++, deg[u]++, edges++;
   }
 int odd = 0. root = 0:
 for (int i = 1; i <= n; i++) {</pre>
   if (deg[i] & 1) odd++, root = i;
 }
 if (odd > 2) return 0;
 if (root == 0) {
   for (int i = 1; i <= n; i++) if (deg[i]) root</pre>
        = i;
 if (root == 0) return 1; //empty graph
 dfs(root);
 if (ans.size() != edges / 2 + 1) return 0;
 reverse(ans.begin(), ans.end());
 return 1;
}
int32_t main() {
```

```
ios_base::sync_with_stdio(0);
cin.tie(0);
int t;
cin >> t:
while (t--) {
  int n. m:
  cin >> n >> m;
  vector<int> deg(n + 1, 0);
  for (int i = 1; i <= m; i++) {</pre>
   int u, v;
   cin >> u >> v;
   g[u].push_back({v, i});
   g[v].push_back({u, i});
    deg[u]++, deg[v]++;
  int sz = m:
  for (int i = 1; i <= n; i++) {</pre>
   if (deg[i] & 1) {
     ++sz:
     g[n + 1].push_back({i, sz});
     g[i].push_back({n + 1, sz});
  }
  int ok = solve(n + 1);
  assert(ok);
  vector\langle int \rangle in(n + 2, 0), out(n + 2, 0);
  for (int i = 0; i + 1 < ans.size(); i++) {</pre>
    if (ans[i] != n + 1 && ans[i + 1] != n + 1) {
     in[ans[i + 1]]++:
     out[ans[i]]++;
   }
  }
  int res = 0;
  for (int i = 1; i <= n; i++) res += in[i] ==</pre>
      out[i]:
  cout << res << '\n';
  for (int i = 0; i + 1 < ans.size(); i++) if</pre>
       (ans[i] != n + 1 && ans[i + 1] != n + 1)
      cout << ans[i] << ' ' << ans[i + 1] <<</pre>
       '\n':
  for (int i = 0; i <= n + 1; i++) g[i].clear();</pre>
return 0;
```

7.8 Minimum Spanning Tree from Each Egde

```
const double inf = 0.0000:
const int 1x = 2e5 + 5;
const int mod = 998244353:
const int hs = 3797;
struct graph {
    int u;
    int v;
   int w;
    int idx;
    bool operator < (const graph ob) const {</pre>
       return w < ob.w;</pre>
};
int cost[lx];
int bap[lx], n, m;
graph temp;
vector<graph> adj;
set<int> e[lx];
int find(int x)
    if(x == bap[x]) return x;
    return bap[x] = find(bap[x]);
void answer(int u, int v, int w)
{
   u = find(u):
   v = find(v);
   if(e[u].size() < e[v].size())</pre>
       swap(u, v);
   bap[v] = u;
    while(e[v].begin() != e[v].end()) {
```

```
if(e[u].find(*e[v].begin()) != e[u].end())
           cost[*e[v].begin()] -= w;
       }
       else e[u].insert(*e[v].begin());
       e[v].erase(e[v].begin());
   }
}
void solve()
{
   cin >> n >> m;
   for(int i = 1; i <= n; i++) bap[i] = i;</pre>
   for(int i = 1; i <= m; i++) {</pre>
       cin >> temp.u >> temp.v >> temp.w;
       temp.idx = i;
       cost[i] = temp.w;
       adj.push_back(temp);
       e[temp.u].insert(i);
       e[temp.v].insert(i);
   }
   sort(adj.begin(), adj.end());
   11 \text{ ans} = 0:
   for(int i = 0; i < m; i++) {</pre>
       int u = adj[i].u;
       int v = adj[i].v;
       int w = adj[i].w;
       int idx = adj[i].idx;
       if(find(u) == find(v)) continue;
       answer(u, v, w);
       ans += w;
   }
   for(int i = 1; i <= m; i++) cout << ans + 1LL</pre>
        * cost[i] << ' ';
   cout << endl;</pre>
```

7.9 Notes

- 1. To make a tree cyclic, add $\frac{\text{number of leaves} + 1}{2}$ number of edges.
- 2. An Euler path exists in an undirected graph if and only if the degree of each node is even. In a directed graph, an Euler path exists if and only if the indegree and outdegree of each node are equal.
- 3. A graph is a biconnected component (block) if it has no articulation points, but may have bridges (like two nodes). A graph is a bridge-tree if it has no bridges, but may have articulation points.
- 4. For an Euler path to exist in an undirected graph, the graph must have exactly two nodes with odd degree. The path will start at one node and end at the other. Additionally, exactly one node should have an indegree equal to its outdegree plus one, and exactly one node should have an outdegree equal to its indegree plus one.
- 5. An Euler path or cycle exists in a graph if and only if the graph is connected.

7.10 Topological Sort

```
vector<vector<int>> adjList;
vector<int> topoLogicalSort(vector<pair<int,
    int>> edges, int nodeCount) {
    vector<int> inDegree(nodeCount + 1, 0);
    vector<vector<int>> adjList(nodeCount + 1);
    for(auto edge : edges) {
        inDegree[edge.second] += 1;
        adjList[edge.first].push_back(edge.second);
    }
    queue<int> nodesWithIndegreeZero;
    for(int i = 0; i <nodeCount; i++) {
        if(inDegree[i] == 0) {
            nodesWithIndegreeZero.push(i);
        }
}</pre>
```

```
}
}
vector<int> topoLogicallySortedNodes;
while (!nodesWithIndegreeZero.empty()) {
    int node = nodesWithIndegreeZero.front();
    topoLogicallySortedNodes.push_back(node);
    nodesWithIndegreeZero.pop();
    for (auto x: adjList[node]) {
        inDegree[x]--;
        if (inDegree[x] == 0) {
            nodesWithIndegreeZero.push(x);
        }
    }
}
return topoLogicallySortedNodes;
}
```

7.11 dijkstra simple

```
// Dijkstra for weighted graph with non-negative
    weights
void dijkstra(int src, vector<vector<pair<int,</pre>
    int>>> &adj, vector<int> &dist)
 int n = adj.size(); // Nodes are 0-based
 dist.assign(n, INT_MAX);
 priority_queue<pair<int, int>, vector<pair<int,</pre>
      int>>, greater<>> pq;
 dist[src] = 0;
 pq.push({0, src});
 while (!pq.empty())
   auto [d, u] = pq.top();
   pq.pop();
   if (d > dist[u])
     continue; // Outdated entry
   for (auto [v, w] : adj[u])
     if (dist[u] + w < dist[v])</pre>
       dist[v] = dist[u] + w;
       pq.push({dist[v], v});
```

```
}
}

// Assumes adj[u] = vector of {neighbor, weight}
// Make sure graph has no negative weights
```

8 G.Tree

8.1 2D Fenwick Tree

```
const 11 sz = 1040:
11 fre[sz][sz], ar[sz][sz];
11 n;
void update(ll x, ll y, ll val) {
   ++x; ++y;
   for(ll i=x; i<=n; i += (i&(-i))) {</pre>
       for(11 j=y; j<=n; j += (j&(-j))) {
           fre[i][j] += val;
       }
   }
}
11 query(ll x, ll y) {
   ++x; ++y;
   11 sum=0:
   for(ll i=x; i>0; i -= (i&(-i))) {
       for(11 j=y; j>0; j -= (j&(-j))) {
           sum += fre[i][j];
       }
   }
   return sum:
ll areaSum(ll x1, ll y1, ll x2, ll y2) {
   ll ans = query(x2, y2)
            - query(x2, y1-1) - query(x1-1, y2)
           + query(x1-1, y1-1);
   return ans;
```

8.2 MST Kruskal's

```
// DSU is required
// sort the edge according to your problem(asc,
// Asc -> Minimum spanning tree
// Dsc -> Maximum spanning tree
   int n, m; cin >> n >> m;
   vector<array<int, 3>> ed;
   for(int i = 1; i <= m; i++){
       int u, v, w; cin >> u >> v >> w;
       ed.push_back({w, u, v});
   sort(ed.begin(), ed.end());
   long long ans = 0;
   dsu d(n);
   for (auto e: ed){
       int u = e[1], v = e[2], w = e[0];
       if (d.same(u, v)) continue;
       ans += w;
       d.merge(u, v);
   }
   cout << ans << '\n';</pre>
```

8.3 MST To find Kth Sortest Number in Range

```
#define ULL
                       unsigned long long
#define ff
                       first
#define ss
                       second
#define pb
                       push_back
#define pii
                       pair<int,int>
#define all(a)
                       a.begin(),a.end()
#define MEM(a,x)
                       memset(a,x,sizeof(a))
#define FOR(i,a,b)
                       for(int i=a;i<=b;i++)</pre>
#define ROF(i,a,b)
                       for(int i=a;i>=b;i--)
#define REP(i,b)
                       for(int i=0;i<b;i++)</pre>
const int N=1e6+5;
pii arr[N];
int flag[N];
vector<int>tree[4*N];
```

```
void build(int node,int L,int R)
{
   if(L==R){
       tree[node].pb(arr[L].ss);
       return;
   }
   int mid=(L+R)/2;
   build(node*2,L,mid);
   build(2*node+1,mid+1,R);
   merge(all(tree[2*node]),all(tree[2*node+1]),back_inse
int query(int node,int L,int R,int l,int r,int
    val)
{
   if(L==R) return tree[node][0];
        cnt=upper_bound(all(tree[2*node]),r)-lower_bound(
   int mid=(L+R)/2;
   if(val>cnt) return
        query(node*2+1,mid+1,R,l,r,val-cnt);
   else return query(node*2,L,mid,1,r,val);
}
int main()
Ł
   int n,q;
   while(scanf("%d %d",&n,&q)!=EOF){
       FOR(i,1,4*n) tree[i].clear();
       FOR(i,1,n) scanf("%d",&arr[i].ff);
       FOR(i,1,n) arr[i].ss=i,flag[i]=arr[i].ff;
       sort(arr+1,arr+n+1);
       build(1,1,n);
       while(q--){
           int 1,r,k;
           scanf("%d %d %d",&l,&r,&k);
           printf("\frac{d}{n},flag[query(1,1,n,l,r,k)]);
```

8.4 Merger Sort Tree Using Segment Tree

```
const int N=1e6+5:
vector<int>V[N],tree[4*N];
void build(int node,int L,int R) {
   if(L==R) {
       sort(all(V[L]));
       tree[node] = V[L];
       return:
   int mid=(L+R)/2;
   build(node*2,L,mid);
   build(2*node+1,mid+1,R);
   merge(all(tree[2*node]),all(tree[2*node+1]),
   back_inserter(tree[node]));
}
int query(int node,int L,int R,int l,int r,int
    val) {
   if(r<L or R<l or tree[node].empty()) return 0;</pre>
   if(1<=L and R<=r) {</pre>
       int cnt=upper_bound(all(tree[node]),val)
       -tree[node].begin();
       return cnt;
   }
   int mid=(L+R)/2;
   int x=query(node*2,L,mid,l,r,val);
   int y=query(node*2+1,mid+1,R,l,r,val);
   return x+y;
}
```

8.5 Minimum Spanning Tree – MST using Prim's Algo

```
int main() {
   int N=5,m=6;
   vector<pair<int,int> > adj[N];
   int parent[N];
   int key[N];
   bool mstSet[N];
   for (int i = 0; i < N; i++)
        key[i] = INT_MAX, mstSet[i] = false;
   priority_queue< pair<int,int>, vector
        <pair<int,int>>, greater<pair<int,int>> >
```

```
pq;
kev[0] = 0;
parent[0] = -1;
pq.push({0, 0});
while(!pq.empty()) {
   int u = pq.top().second;
   pq.pop();
   mstSet[u] = true;
   for (auto it : adj[u]) {
       int v = it.first:
       int weight = it.second;
       if (mstSet[v] == false && weight <</pre>
           kev[v]) {
           parent[v] = u;
           key[v] = weight;
           pq.push({key[v], v});
   }
}
for (int i = 1; i < N; i++)</pre>
   cout << parent[i] << " - " << i <<" \n";
return 0;
```

9 H.Strings

9.1 KMP Algorithm

```
// finding frequency of P in S, as a substring.
// call build_failure_function() first
#define MAX 1000005
int failure[MAX];
// longest prefix that also mactchs current suffix
void build_failure_function(string pattern, int m)
{
   failure[0] = 0;
   failure[1] = 0; // base case

for (int i = 2; i <= pattern.size(); i++)
   f</pre>
```

```
int j = failure[i - 1];
       while (true)
           if (pattern[j] == pattern[i - 1])
              failure[i] = j + 1;
              break:
           if (j == 0)
              failure[i] = 0;
              break;
           i = failure[i]:
}
int frequency_of_P_in_S(string &s, stirng &p)
   build_failure_function(p, p.size());
   int total_mtc = 0;
   int ans = 0;
   for (int i = 0; i < s.size(); i++)</pre>
       if (s[i] == p[total_mtc])
           total mtc++:
       else
           while (total_mtc > 0)
              total_mtc = failure[total_mtc];
              if (s[i] == p[total_mtc])
                  total mtc++:
                  break;
       if (p.size() == total_mtc)
           ans++:
           total_mtc = failure[total_mtc];
```

18

```
}
return ans;
}
```

9.2 LCS

9.3 Lexiographically smallest string by cycle shift

```
int minlex(char[] s) {
   int len = s.size();
   int n = 2*len, i=0, j=1, k=0, a, b;
   while(i+k<n && j+k<n) {</pre>
       a=(i+k>=len)?s[i+k-len]:s[i+k];
       b=(j+k>=len)?s[j+k-len]:s[j+k];
       if(a==b) k++:
       else if(a>b) {
           i=i+k+1;
           if(i<=j) i=j+1;</pre>
           k=0;
       }
       else {
           j=j+k+1;
           if(j<=i) j=i+1;</pre>
           k=0;
       }
   return min(i, j);
```

9.4 Longest Palindromic Subsequence

9.5 Manacher's Algorithm

```
int p[lx];
void manachers_algorithm(string s)
   string t = "#";
   for(char c : s) {
       t += c;
       t += '#':
   }
   t = "-" + t + "?":
   int 1 = 1, r = 1, n = s.size();
   for(int i = 1; i < n; i++) {</pre>
       p[i] = max(0, min(r - i, p[1 + r - i]));
       while(s[i - p[i]] == s[i + p[i]]) {
           ++p[i];
       if(i + p[i] > r) {
          1 = i - p[i];
          r = i + p[i];
   }
```

9.6 Trie

```
struct Node {
   Node *next[26];
   int frequency;
   Node() {
       frequency = 0;
       for (int i = 0; i < 26; i++) next[i] =</pre>
           nullptr;
   }
};
void addString(Node *root, string s) {
   Node *cur = root;
   for (char c : s) {
       if (cur->next[c-'a'] == nullptr)
            cur->next[c-'a'] = new Node();
       cur = cur->next[c-'a'];
       cur->frequency++;
   }
int queryString(Node *root, string s) {
   Node *cur = root;
   for (char c : s) {
       if (cur->next[c-'a'] == nullptr) return 0;
       cur = cur->next[c-'a'];
   return cur->frequency;
}
int main()
   ios::sync_with_stdio(false);
   cin.tie(0);
   int n, q;
   cin >> n >> q;
   Node *root = new Node();
   for (int i = 0; i < n; i++) {</pre>
       string s;
       cin >> s;
       addString(root, s);
   for (int i = 0; i < q; i++) {</pre>
       string s;
       cin >> s;
```

```
cout << queryString(root, s) << "\n";
}
return 0;
}</pre>
```

9.7 Z function

```
int z[lx]:
void Z_Algorithm(string s)
   z[0] = 0:
   int 1 = 0, r = 0, n = s.size();
   for(int i = 1; i < n; i++) {</pre>
       if(i <= r) {
          z[i] = min(z[i-1], r-i+1); //
               Mirror Index of i will be i - 1
       }
       while(i + z[i] < n and s[z[i]] == s[i +
           z[i]]) {
          ++z[i];
       }
       if(i + z[i] - 1 > r) {
          1 = i:
          r = i + z[i] - 1;
      }
   }
}
```

9.8 string hashing

```
while (k) {
       if (k \& 1) ans = (long long) ans * n % mod;
       n = (long long) n * n % mod;
       k >>= 1;
   return ans:
// extra mod 999999989, 1e9+9, 1e9+7
const int MOD1 = 127657753, MOD2 = 987654319;
const int p1 = 137, p2 = 277;
int ip1, ip2;
pair<int, int> pw[N], ipw[N];
void prec() {
   pw[0] = \{1, 1\};
   for (int i = 1; i < N; i++) {</pre>
       pw[i].first = 1LL * pw[i - 1].first * p1 %
            MOD1:
       pw[i].second = 1LL * pw[i - 1].second * p2
           % MOD2:
   }
   ip1 = power(p1, MOD1 - 2, MOD1);
   ip2 = power(p2, MOD2 - 2, MOD2);
   ipw[0] = \{1, 1\};
   for (int i = 1; i < N; i++) {</pre>
       ipw[i].first = 1LL * ipw[i - 1].first *
            ip1 % MOD1;
       ipw[i].second = 1LL * ipw[i - 1].second *
           ip2 % MOD2;
   }
struct Hashing {
   int n;
   string s; // 0 - indexed
   vector<pair<int, int>> hs; // 1 - indexed
   Hashing() {}
   Hashing(string _s) {
       n = _s.size();
       s = _s;
       hs.emplace_back(0, 0);
       for (int i = 0; i < n; i++) {</pre>
       pair<int, int> p;
       p.first = (hs[i].first + 1LL * pw[i].first
            * s[i] % MOD1) % MOD1;
```

```
p.second = (hs[i].second + 1LL *
           pw[i].second * s[i] % MOD2) % MOD2;
       hs.push_back(p);
       }
   pair<int, int> get_hash(int 1, int r) { // 1
        - indexed
       assert(1 <= 1 && 1 <= r && r <= n):
       pair<int, int> ans;
       ans.first = (hs[r].first - hs[l - 1].first
           + MOD1) * 1LL * ipw[l - 1].first %
       ans.second = (hs[r].second - hs[1 -
           1].second + MOD2) * 1LL * ipw[1 -
           1].second % MOD2;
       return ans:
   pair<int, int> get_hash() {
       return get_hash(1, n);
};
```

10 I.Bit

10.1 Equation $* a+b = a \oplus b + 2(a \& b)$

```
 * a+b = a-b+a\&b 
* a\oplus b = a|b-a\&b 
* k_{th} \text{ bit is set in x iff } x \bmod 2^{k-1} > 2^k 
* k_{th} \text{ bit is set in x iff } x \bmod 2^{k-1} - x \bmod 2^k \neq 0 (=2^k) 
* n \bmod 2^i = n\&(2^i-1) 
* 0.1\oplus 2\oplus 3\oplus \cdots \oplus (4k-1) = 0 \text{ for any } k \geq 0
```

10.2 Maximum And Pair

```
Int checkbit(int p, int n) {
   cnt = 0;
   for(i=0; i<n; i++) {</pre>
```

```
if(p&(ai==p)cnt++;
}
return cnt;
}
for(b=32; b>=0; b--) {
   cnt = checkbit(res |<<b, n);
   if(cnt >=2) res|=1<<b;
}</pre>
```

10.3 Maximum XOR of all subsequence

```
#define INT_BITS 32
int maxSubarrayXOR(int set[], int n) {
   for (int i = INT_BITS-1; i >= 0; i--) {
       int maxInd = index:
       int maxEle = INT MIN:
       for (int j = index; j < n; j++) {
           if ( (set[j] & (1 << i)) != 0 &&</pre>
               set[j] > maxEle )
               maxEle = set[j], maxInd = j;
       }
       if (maxEle == INT_MIN) continue;
       swap(set[index], set[maxInd]);
       maxInd = index;
       for (int j=0; j<n; j++) {</pre>
           if (j != maxInd && (set[j] & (1 << i))</pre>
               set[j] = set[j] ^ set[maxInd];
       }
       index++;
   }
   int res = 0;
   for (int i = 0; i < n; i++)</pre>
       res ^= set[i];
   return res;
```

10.4 Minimum XOR Operation

Minimum xor pair - Trie

- 1. Sort the array.
- 2. Find the minimum of $ar[i] \oplus ar[i+1]$, where \oplus denotes the XOR operation.

Minimum XOR subarray - Trie Minimum XOR of OR and AND in an array

1. Minimum XOR of $(x \lor y) \oplus (x \land y)$ is equal to the minimum XOR of all pairs in the array, where \lor denotes the OR operation and \land denotes the AND operation.

Sum of XOR of all elements of subsets

1. $2^{n-1} \times (OR \text{ of whole array})$, where n is the number of elements in the array.

Maximum OR pair

= Maximum element ∨ next greater element (excluding itself)

Sum of OR of all subsets

= $\sum_{i=0}^{n-1} a_i \cdot 2^{n-1}$, where a_i is the *i*-th element in the array and n is the number of elements.

10.5 Number of Subarrays with XOR 0

```
#include<bits/stdc++.h>
using namespace std;

const int N = 1e6 + 9;
int a[N], p[N];

int32_t main() {
   ios_base::sync_with_stdio(0);
   cin.tie(0);
   int n; cin >> n;
   for (int i = 1; i <= n; i++) {
      cin >> a[i];
      p[i] = p[i - 1] ^ a[i];
   }

map<int, int> mp;
long long ans = 0;
mp[p[0]]++; // don't forget to add this (why?)
```

```
for (int i = 1; i <= n; i++) {
   ans += mp[p[i]];
   mp[p[i]]++;
}
cout << ans << '\n';
return 0;
}</pre>
```

10.6 Sum of XOR of All subset in Array

```
int xorSum(int arr[], int n) {
  int bits = 0;
  for (int i=0; i < n; ++i) bits |= arr[i];
  int ans = bits * pow(2, n-1);
  return ans;
}</pre>
```

10.7 Sum of all and of all subset

```
ans = 0;
for(i=0; i<32; i++) {
   cnt=0;
   for(j=0; j<n; j++) {
       if(aj&(1<<i)cnt++;
   }
   subsets = (1<<cnt)-1;
   subsets = subset^(1<<i) ans += subset;
}</pre>
```

11 J.Misc

11.1 Base conversion

```
#pragma once
string const DIGITS = "0123456789ABCDEF";
```

```
unsigned long long int
    basis_string_to_number(string &s, int b) {
   unsigned long long int result = OULL;
   for (char d : s) {
       result = b * result
               + (find(DIGITS.begin(),
                   DIGITS.end(), d) -
                   DIGITS.begin());
   }
   return result;
string number_to_basis_string(unsigned long long
    int n, int b) {
   vector<char> ds:
   do {
       ds.push_back(DIGITS[n % b]);
       n = n / b;
   } while (n != 0);
   return string(ds.rbegin(), ds.rend());
}
```

11.2 Equation Solve

```
// a1*x + b1*y + c1*z = d1 | a2*x + b2*y + c2*z =
    d2 \mid a3*x + b3*y + c3*z = d3
bool eq3(double &x, double &y, double &z,
        double a1, double b1, double c1, double p,
        double a2, double b2, double c2, double q,
        double a3, double b3, double c3, double
            r) {
   double detr = a1 * (b2 * c3 - c2 * b3) - b1 *
        (a2 * c3 - c2 * a3) + c1 * (a2 * b3 - b2)
        * a3):
   if (fabs(detr) < 1e-9) return false;</pre>
   x = (p * (b2 * c3 - c2 * b3) + q * (c1 * b3 -
        b1 * c3) + r * (b1 * c2 - c1 * b2)) /
        detr:
   y = (p * (c2 * a3 - a2 * c3) + q * (a1 * c3 -
        c1 * a3) + r * (c1 * a2 - a1 * c2)) /
   z = (p * (a2 * b3 - b2 * a3) + q * (b1 * a3 -
        a1 * b3) + r * (a1 * b2 - b1 * a2)) /
```

11.3 LCS - longest common subsequences

11.4 LIS - Longest Increasing Subsequence

```
// LIS - Longest increasing subsequence
// O(n*logn)
const int INF=INT_MAX;
vector<int> LIS(vector<int> & inp){
   int n=inp.size();
   vector<int> a(n+2, INF);
```

```
vector<int> res(n,1);
a[0]=-INF;
int lis=0;
for(int i=0; i<n; i++){
    int low = lower_bound(a.begin(), a.end(),
        inp[i])-a.begin();
    a[low] = inp[i];
    lis = low;
    res[i] = lis;
}
return res;
}</pre>
```

11.5 Mex of all Subarray

```
#include<bits/stdc++.h>
using namespace std:
const int N = 1e5 + 9, inf = 1e9;
struct ST {
   int t[4 * N]:
   ST() {}
   void build(int n, int b, int e) {
       t[n] = 0:
       if (b == e) {
          return;
       int mid = (b + e) >> 1, l = n << 1, r = 1
           | 1;
       build(1, b, mid);
       build(r, mid + 1, e);
       t[n] = min(t[1], t[r]);
   void upd(int n, int b, int e, int i, int x) {
       if (b > i || e < i) return;</pre>
       if (b == e && b == i) {
          t[n] = x;
          return;
       int mid = (b + e) >> 1, l = n << 1, r = 1
       upd(1, b, mid, i, x);
       upd(r, mid + 1, e, i, x);
```

```
t[n] = min(t[1], t[r]);
   }
    int get_min(int n, int b, int e, int i, int
        j) {
       if (b > j || e < i) return inf;</pre>
       if (b >= i && e <= j) return t[n];</pre>
       int mid = (b + e) >> 1, l = n << 1, r = 1
       int L = get_min(1, b, mid, i, j);
       int R = get_min(r, mid + 1, e, i, j);
       return min(L, R);
    int get_mex(int n, int b, int e, int i) { //
        mex of [i... cur id]
       if (b == e) return b;
       int mid = (b + e) >> 1, l = n << 1, r = 1
       if (t[1] >= i) return get_mex(r, mid + 1,
            e, i);
       return get_mex(1, b, mid, i);
   }
} t;
int a[N], f[N];
int32_t main() {
    int n;
    cin >> n;
    for (int i = 1; i <= n; i++) {</pre>
       cin >> a[i];
       --a[i]:
    }
    t.build(1, 0, n);
    set<array<int, 3>> seg; // for cur_id = i,
        [x[0]...i], [x[0] + 1...i], ...
        [x[1]...i] has mex x[2]
    for (int i = 1; i <= n; i++) {</pre>
       int x = a[i];
       int r = min(i - 1, t.get_min(1, 0, n, 0, x))
            - 1));
       int 1 = t.get_min(1, 0, n, 0, x) + 1;
       if (1 <= r) {
           auto it = seg.lower_bound(\{1, -1, -1\});
           while (it != seg.end() && (*it)[1] <=</pre>
               r) {
               auto x = *it;
```

```
it = seg.erase(it);
       }
   }
   t.upd(1, 0, n, x, i);
   for (int j = r; j >= 1; ) {
       int m = t.get_mex(1, 0, n, j);
       int L = \max(1, t.get_min(1, 0, n, 0,
           m) + 1):
       f[m] = 1;
       seg.insert({L, j, m});
       j = L - 1;
   int m = !a[i];
   seg.insert({i, i, m});
   f[m] = 1;
}
int ans = 0;
while (f[ans]) ++ans;
cout << ans + 1 << '\n';
return 0;
```

11.6 Mo's Algorithm

```
const int M = 1e9+7;
int n, q;
vector<int> v;
struct query{int l,r,idx;};
int block;
bool comp1(query p,query q){
   if (p.1 / block != q.1 / block) {
       if(p.l==q.1) return p.r<q.r;</pre>
       return p.l < q.l;</pre>
   }
   return (p.l / block & 1) ? (p.r < q.r) : (p.r
        > q.r);
vector<int> fre(1000001, 0);
int no_of_distinct=0;
void add(int idx){
   fre[v[idx]]++;
   if(fre[v[idx]]==1) no_of_distinct++;
```

```
void rmv(int idx){
   fre[v[idx]]--;
   if(fre[v[idx]]==0) no_of_distinct--;
int get_answer(){
   return no_of_distinct;
void mos_algorithm(int n, vector<query>&queries){
   vector<int> answers(queries.size());
   block = (int)sqrt(n);
   sort(queries.begin(), queries.end(),comp1);
   int cur_1 = 0;
   int cur r = -1:
   for (query q : queries) {
       while (cur_1 > q.1) {cur_1--; add(cur_1);}
       while (cur_r < q.r) {cur_r++; add(cur_r);}</pre>
       while (cur_1 < q.1) {rmv(cur_1);cur_1++;}</pre>
       while (cur_r > q.r) {rmv(cur_r);cur_r--;}
       answers[q.idx] = get_answer();
   for(int i:answers) {cout<<i<"\n";}</pre>
}
void srfahad2021(){
   cin>>n;
   v.assign(n, 0);
   for(int i=0: i<n: i++){</pre>
       cin>>v[i];
   vector<query> queries;
   cin>>q;
   for(int i=0; i<q; i++){</pre>
       int 1, r; cin>>l>>r;
       1--; r--;
       queries.push_back({1, r, i});
   mos_algorithm(q, queries);
int32_t main(){
   fast();
   int test=1;
   // cin>>test:
   for(int i=1; i<=test; i++){</pre>
       // cout<<"Case "<<i<": ";
```

```
srfahad2021();
}
return 0;
}
```

11.7 Nim Game 2d

```
int main() {
   int t;
   scanf("%d",&t);
   FOR(tc,1,t) {
       int r,c;
       scanf("%d %d",&r,&c);
       int nim=0:
       FOR(i,1,r) {
          FOR(j,1,c) {
              int tmp;
              scanf("%d",&tmp);
              if(((r-i)+(c-j))%2) {
                  nim^=tmp;
              }
          }
       if(nim) printf("Case %d: win\n",tc);
       else printf("Case %d: lose\n",tc);
   }
}
```

11.8 Number of Subsegment Equal to K

```
int a[lx], t[lx];
int n, k;
map<ll, int> mp;
int o[lx], Plus[lx], Minus[lx], cnt[lx];
int l, r;
ll p[lx];
const int block = 320;
struct Query {
```

```
int 1, r, idx;
       bool operator < (Query q) const {</pre>
               int b1 = 1 / block;
               int b2 = q.1 / block;
               if(b1 == b2) return r < q.r;
               return b1 < b2;</pre>
};
11 \operatorname{res}[1x], ans = 0;
11 query(int x, int y)
       while(x < 1) {
               --1:
               ans += cnt[Plus[1]];
               cnt[o[1]]++:
       while(r < y) {</pre>
               ++r;
               ans += cnt[Minus[r]];
               cnt[o[r]]++;
       }
       while(1 < x) {
               cnt[o[1]]--;
               ans -= cnt[Plus[1]];
               1++:
       while(y < r) {</pre>
               cnt[o[r]]--;
               ans -= cnt[Minus[r]];
       }
       return ans;
}
void solve()
       cin >> n >> k;
       for(int i = 1; i <= n; i++) cin >> t[i];
       for(int i = 1; i <= n; i++) cin >> a[i];
       for(int i = 1; i <= n; i++) {</pre>
               if(t[i] == 1) p[i] = p[i - 1] +
                    a[i];
               else p[i] = p[i - 1] - a[i];
```

```
for(int i = 0; i <= n; i++) {</pre>
       mp[p[i]];
       mp[p[i] + k];
       mp[p[i] - k];
}
int idx = 0;
for(auto &it : mp) {
       it.second = ++idx;
for(int i = 0; i <= n; i++) {</pre>
       o[i] = mp[p[i]];
       Plus[i] = mp[p[i] + k];
       Minus[i] = mp[p[i] - k];
cnt[o[0]]++:
1 = 0; r = 0;
int gr;
vector<Query> Q;
cin >> qr;
for(int i = 0; i < qr; i++) {</pre>
       int x, y;
       cin >> x >> y;
       Query q;
       q.1 = x - 1;
       q.r = y;
       q.idx = i;
       Q.push_back(q);
}
sort(Q.begin(), Q.end());
for(auto it : Q) {
       res[it.idx] = query(it.1, it.r);
for(int i = 0; i < qr; i++) cout << res[i]</pre>
    << endl;
```

24

11.9 Ternary Search

```
double ternary_search(double 1, double r) {
   double eps = 1e-9; //set the error limit here
```

```
while (r - 1 > eps) {
    double m1 = 1 + (r - 1) / 3;
    double m2 = r - (r - 1) / 3;
    double f1 = f(m1);//evaluates the function
        at m1
    double f2 = f(m2); //evaluates the
        function at m2
    if (f1 < f2)
        1 = m1;
    else
        r = m2;
}
return f(1);//return the maximum of f(x) in
    [1, r]</pre>
```

```
}
```

11.10 Tower of Hanoi