

Ensuring TANGOControl System

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Abstract.³

Current use of TANGO is mostly in Synchrotron and a bit further in neutron a neutron source, but the community like to extend this by an explicit request of the industry. Why only extend to industry? Why not also extend it to commercial like can by home automation? Like in those to extensions, many of the issues faced by the developers (even electronic and computer scientist) are quite similar. Perhaps with differences on the time ranges or precisions but with equivalent abstractions.

Spreading the possible uses we are not making more complicated the current TANGO use. By today, TANGO runs in 32 and 64 bit architectures and from very small embedded instruments, up to very big computers with a huge among of CPU and memory available. Then the objective of ensure this *transport layer*, must work in the small case and then will also work in the bigger one.

The goal of ensure TANGO must produce an outcome as similar as the *TLS* is for the web navigation. Must be possible to co-like with non secured access, but with a tendency to a complete transparent ensuring. Perhaps the migration process would be not as fast as we could want, specially due to the introduction of the certificates infrastructure, but as the TANGO installations are contained in the institutions, and upgrade in this way would be like any other upgrade.

It is very important goal to have the TANGO implementation as Free Software, as this paper cryptography outcomes must be to have public access and auditable algorithms and sources.

Keywords: Cryptography⁴, Distributed Systems, Secure engineering.

1 Introduction

- What is TANGO? It's a CORBA based distributed system used for scientific installation control.

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⁴ This big keyword includes proposals over *Public key*, *Elliptic Curves*, *Symmetric algorithms*, *stream cyphers*, *secret sharing* and also *Homomorphic encryption* for databases

- Event system \emptyset MQ(*Zero Message Queue*).
- SARDANA, TAURUS and ATK as *presentation* and *session layers* are in further work.
- What is the meaning of a secure system? What is security in a distributed system?
- TANGO as a Supervisory Control And Data Acquisition (SCADA) and/or Industrial Control System (ICS). TANGO, complemented by ATK, SARDANA, and TAURUS. As will be explained in 7.3, ensure SARDANA, TAURUS and ATK is out of the scope of this paper, except for those things that defines the interaction with TANGO.
- Distributed systems transparencies [1] that TANGO complains, and which are not
- Security threads, policies and mechanisms. Section 2. Go further that the Locking/Access control
- Why to secure it? Trust in a peripheral firewalls is not enough. Often communications between tango installations (different tango-db) requires firewall rules to allow it, but this doesn't allow to filter by agent or by who is allowed to access the information. In practice, what is filtered is an specific computer traffic, but this breaks many of the distributed system transparencies.
- Following the 3 layers structure of a distributed system [1] to identify scenarios in section 2 and solutions in sections 4:
 - Agent authentication in the presentation layer (section 2.1). Possible solutions as the zero-knowledge proof (section 4.1) and Secret Sharing (section 4.2).
 - Domain layer communications protection in section 2.2. Trusted computing with elliptic curves 4.3, data communication with symmetric encryption in section 4.4 and stream cyphering in section 4.5)
 - Ensure Data layer (section 2.3) with homomorphic encryption (section 4.6) in the database.
- The price of the information and the balance between the cost to ensure and the value of the ensured goods. Section 3
- Alert on possible attacks to protect against what is already saw as a security thread, section 5, distinguishing between passive (section 5.1), active (section 5.2) and side channels (section 5.3)
- Any already saw thread should have it countermeasure (section 6) with special interest in intrusion detection (section 6.1)
- final conclusions about ensuring protocols (section 7.1) and IT environmental security (section 7.2)

2 Identifying scenarios

- Non-implemented distributed system transparencies from [1] necessary to ensure a quality service.

- Confidentiality (encryption and authentication): information must be disclosed only *to* the authorized and only *by* the authorized),
- Integrity (authorization): only authorized can set information.
- Auditory: trace who access where (extremely useful for a security breach analysis).
- In terms of security threads, which is more representative from [2] for the current use case? Three may types: *Hospital, Bank, Military Base*. Practical paranoia [3]
- Cryptosystem configuration, security levels and information classification. 3
- Cryptosystem setup reset.
- Setup & Public-Key distribution protocols [2] sec.3.7.2
- Secret Shared schemas for (k,n)-to decrypt or (k,n)-signants 4.2
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2.1 Ensuring presentation layer

- Agent authentication in a distributed system
- Ensuring communication between agents and between those agents with the user interfaces.
 - *Command, Attribute, Properties*: Authenticate who can do the *read* and *write* operations. Encrypted logging who did any change.
 - This can be compared with *RFID* communication between card and readers, but adding communication in between the agents
- multicast, events ($\emptyset MQ$) and the other features that must be secured. Perhaps secret sharing? Secret splitting? section 4.2
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2.2 Ensuring domain layer

- Trusted Computing and Hardware protections
- Ensure logging system
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2.3 Ensuring data layer

- TANGO database access control
- Ensuring between instrumentation and the agents out of the scope of this paper. This is a very dependant on the instrumentation manufacturers. From the iso layer level view, even if the access to the hardware is not networked, the agent communication to the instrumentation is *data link layer* and this paper is focus in *transport* and above layers.
- Homomorphic Encryption for Database access
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3 Security levels

- Security levels: Open or unclassified, confidential, Secret, Top Secret.
- remember the German standard on this levelling and the European commission “*fiche 17*” (“Exchange of EU classified information”)

4 Communication hybrid schema

- Embedded in instrumentation, limited calculation capacity (it must behave indistinguishable if it’s a huge server or an embedded board), limited bandwidth (Don’t increase the current needs significantly): *very good candidate for elliptic curves (section 4.3), generalized Rijndael (section 4.4) and stream cipher (section 4.5)*.
- Public-key to agreed a session key as the usual hybrid systems. This session keys shall be used for symmetric or stream cyphering.
- Session keys refresh.
- Use the Symmetric key to seed a shared PseudoRandomGenerator as a key for a stream cipher of transmitted data and listened data between talkers
- *PseudoRandomGenerator* (PRG), can be use the KeyDerivationFunction (KDF) of the Rijndael or better other possible alternatives
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4.1 Zero-knowledge proof for authentication

- The agents in the distributed system must be authenticated to be sure that they hasn’t been supplanted
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4.2 Secret Sharing and secret splitting

- Multicast and event system. When a event is emitted, many would be subscribed, but encryption must be only made once.
- To allow some one access to some specific data, perhaps it can require the “grant” from more than one agent of the distributed system. That is, to give it the key may (k,n) must act to.
- Authorization units may be bigger than one agent. A (k,n)-signature to have only one to verify for all.

4.3 Elliptic curves for public key

- Set institution set of curves with different sizes for different level of secrecy (or even different curves for a separable sets in the same secrecy level). Isogeny volcanoes [4].
- Capability to reset a curve setup on any of those secrecy levels (section 3)
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4.4 Rijndael generalization for symmetric key

- How to decide the good parameters of Rijndael? (#rounds,#rows,#columns,wordsize of the block and the key) [5]
- Current AES has advantage on 32bit processor implementation, what about 64bits
- AESWrap [6]
- Secrecy levels (section 3)
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4.5 Key Derivation Functions for stream ciphering

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4.6 Homomorphic Encryption

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5 Brainstorming attacks

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5.1 Passive attacks

- Eavesdropping
- Noise to block an alarm transmission
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5.2 Active attacks

- Men-in-the-middle (active attacks) between agents
- Interruption: Break the public face, web site or gui. Kill a vital agent.
- Modification/Fabrication: Supplant agents.
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5.3 Side channel attacks

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6 Attacks countermeasures

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6.1 Intrusion Detection

- Detection and recovery
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7 Conclusions

- All those fields mentioned on this paper require a much further detailed paper each.
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7.1 Protocols

- Protocol layers [7]
- Security architecture patterns
- Trust ring vs. trust tree (institution CA until the leaves)
- Streaming protection systems (specially for DevEncoded transmission of big images when fast acquisitions)
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7.2 Environmental IT Security

- The weakest brick: secure the transmission but store in a plain file system
- Human behaviour and psychology.
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7.3 Further work

- ATK/ TAURUS user authentication using PAM system (or equivalent in non unix-like systems). Any other user interface that can access tango.
- In all the algorithms on this paper this must be taken into account to minimize redesigns.

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