Divide and Conquer

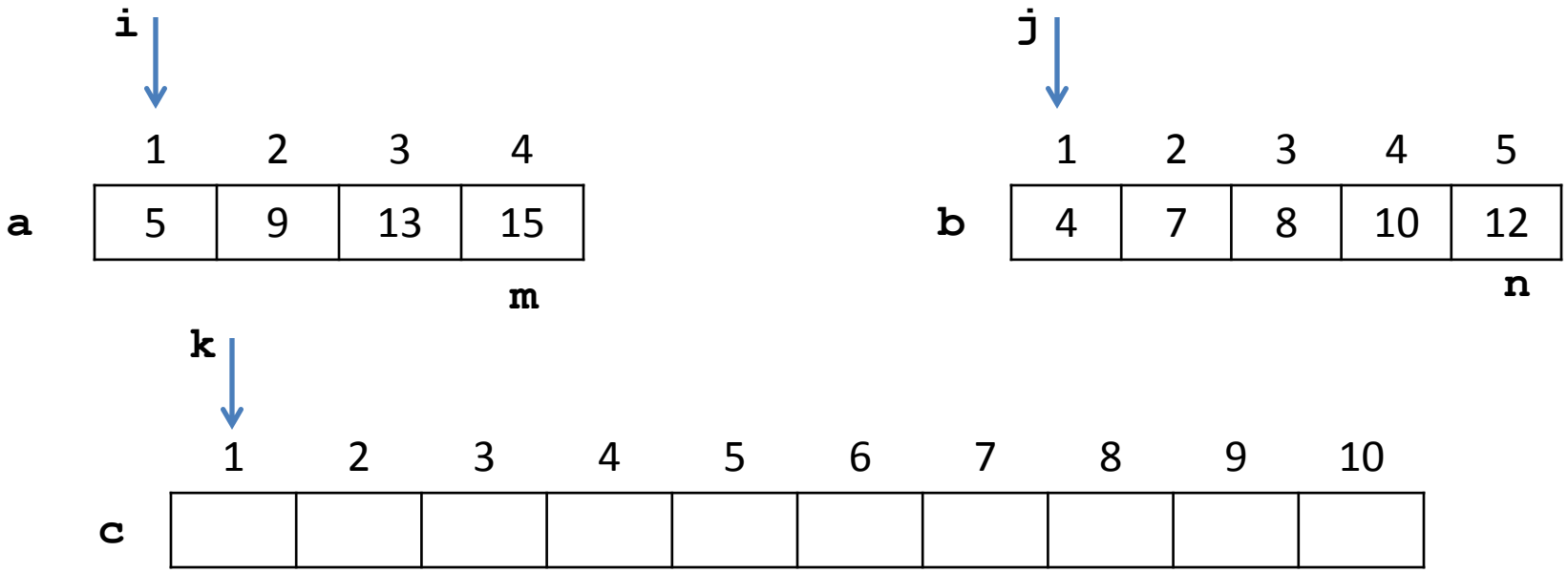
Merge Sort

Divide and Conquer

- Divide and Conquer strategy computes the solution of the problem by computing the solution to the sub-problems and unifying the solutions of the sub-problems as the solution of the given problem.
- Solution of the sub-problem is computed in a similar way till the base case is reached.
- Merge Sort is an example

Merging

- Combining two sorted lists into one sorted list is called as merging.
- Initialize an index variable to the index of the first element of each of the list.
- Initialize an index variable to the index of the first element of the initially empty merged list.



1. Compare a[i] and b[j] and the smallest is copied into the current index location k of the merged list.
2. Increment the source and destination indices of the copy operation
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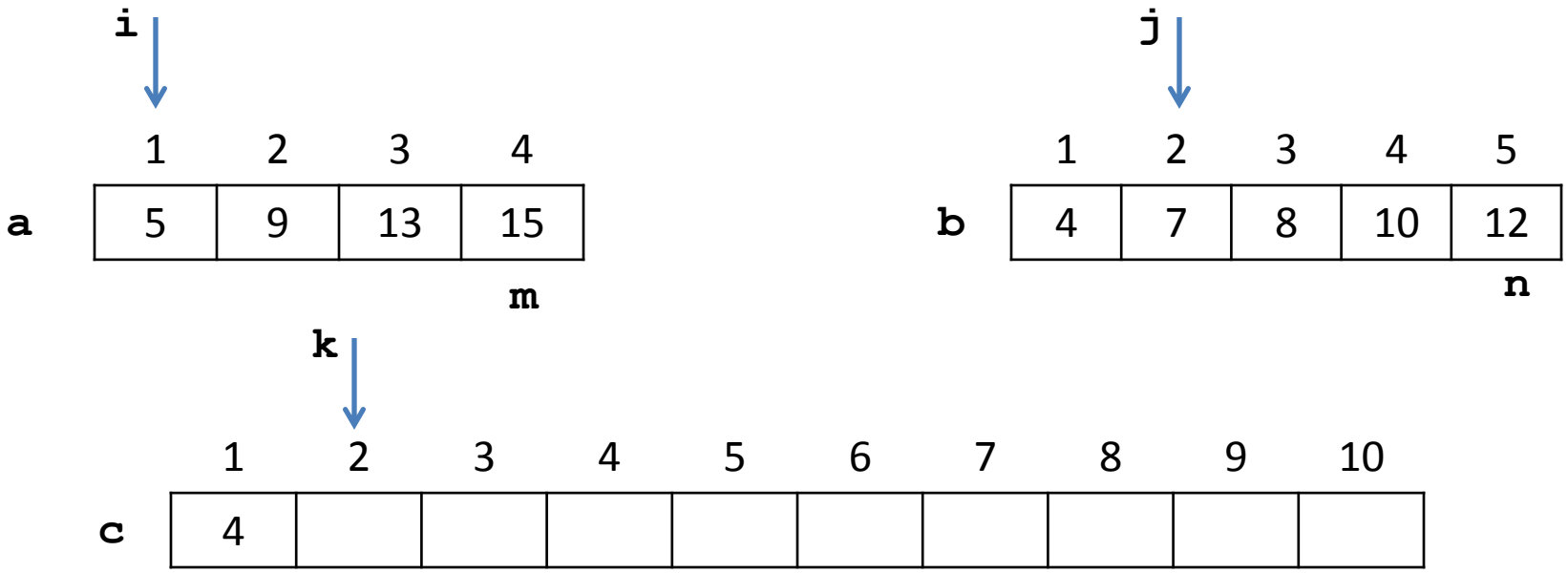
4. k:=k+1;



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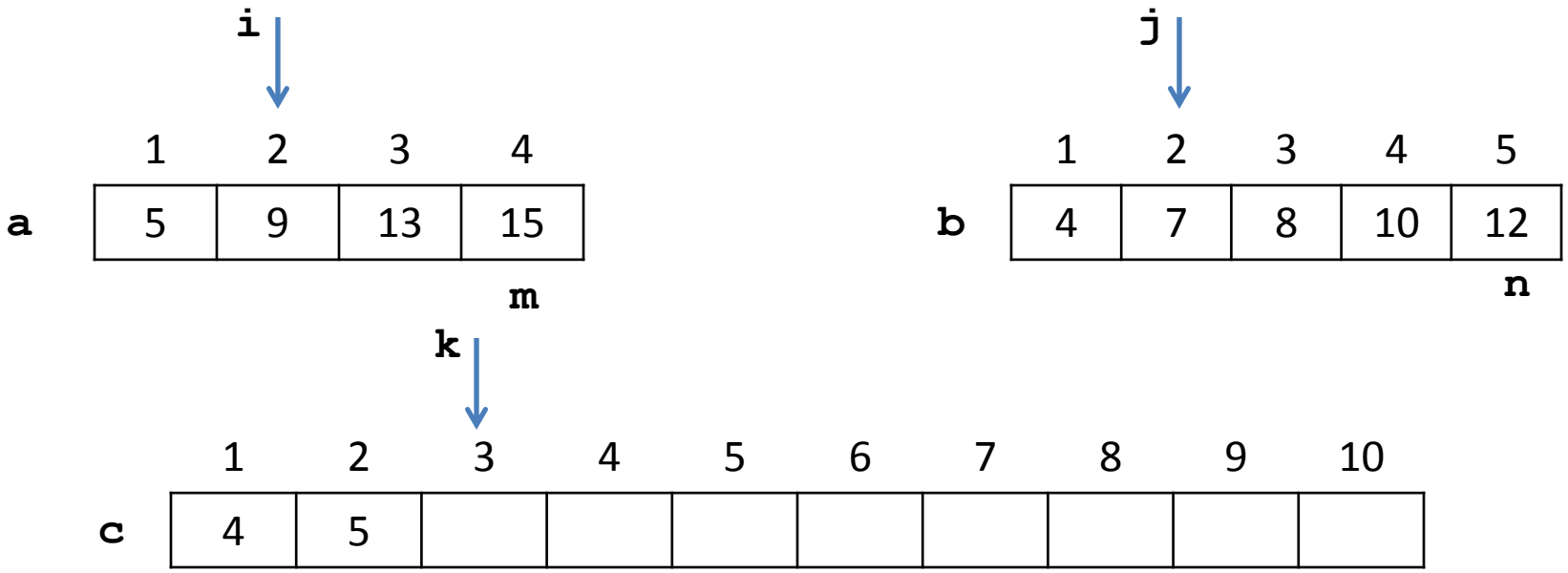

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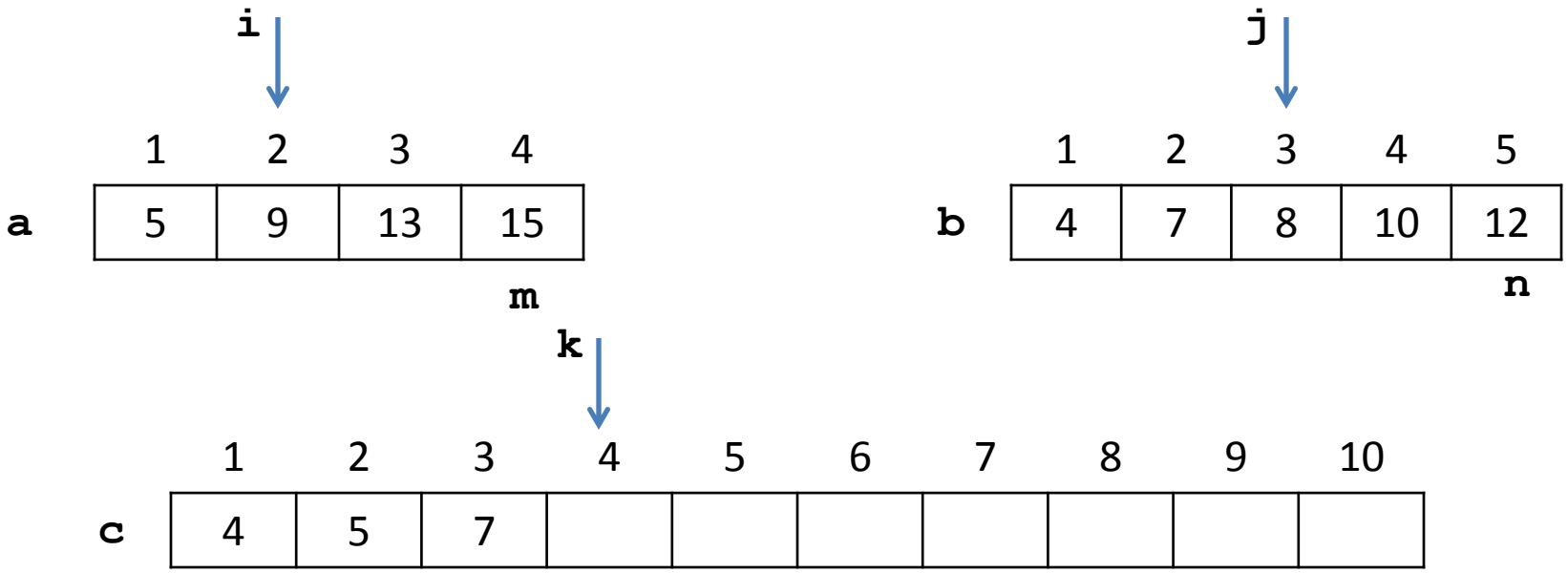

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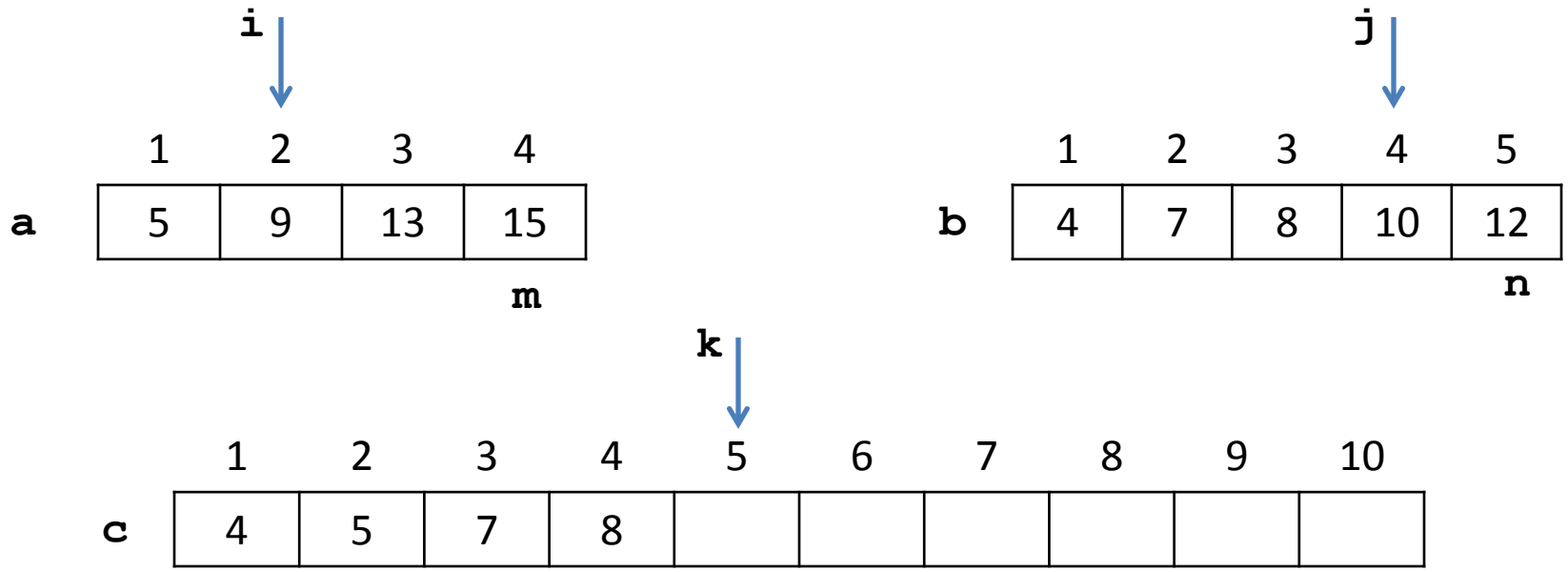

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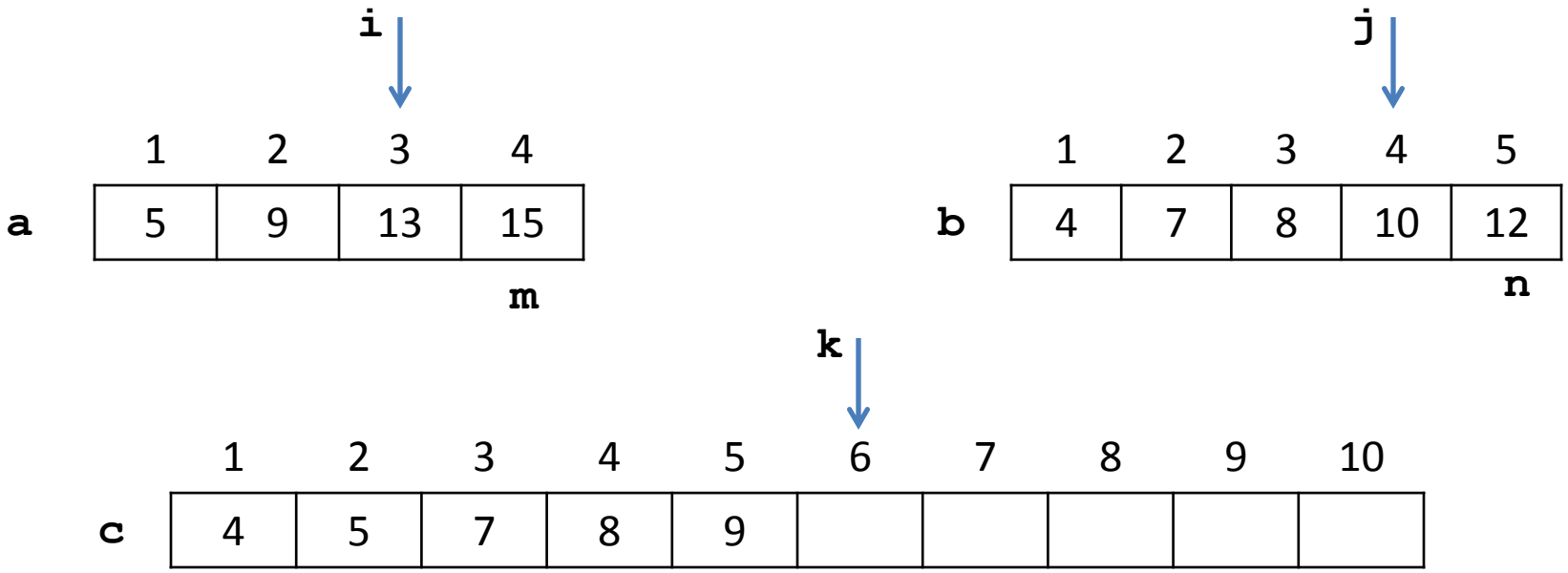

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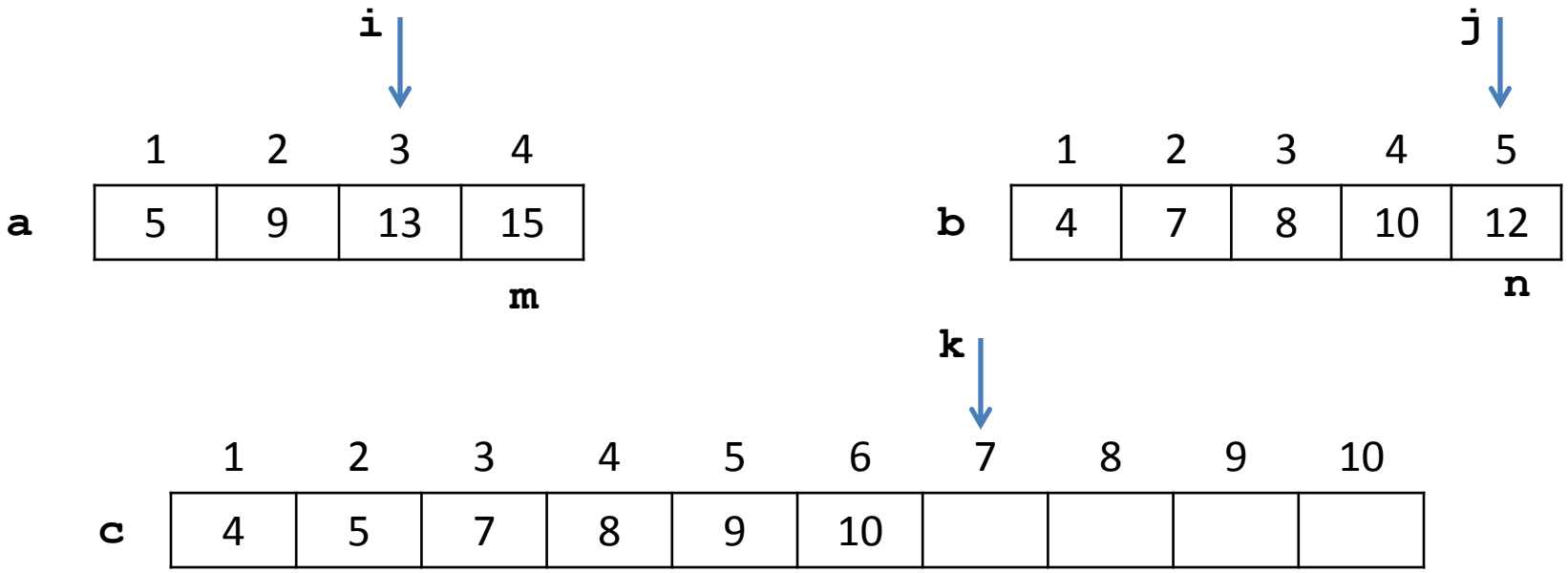

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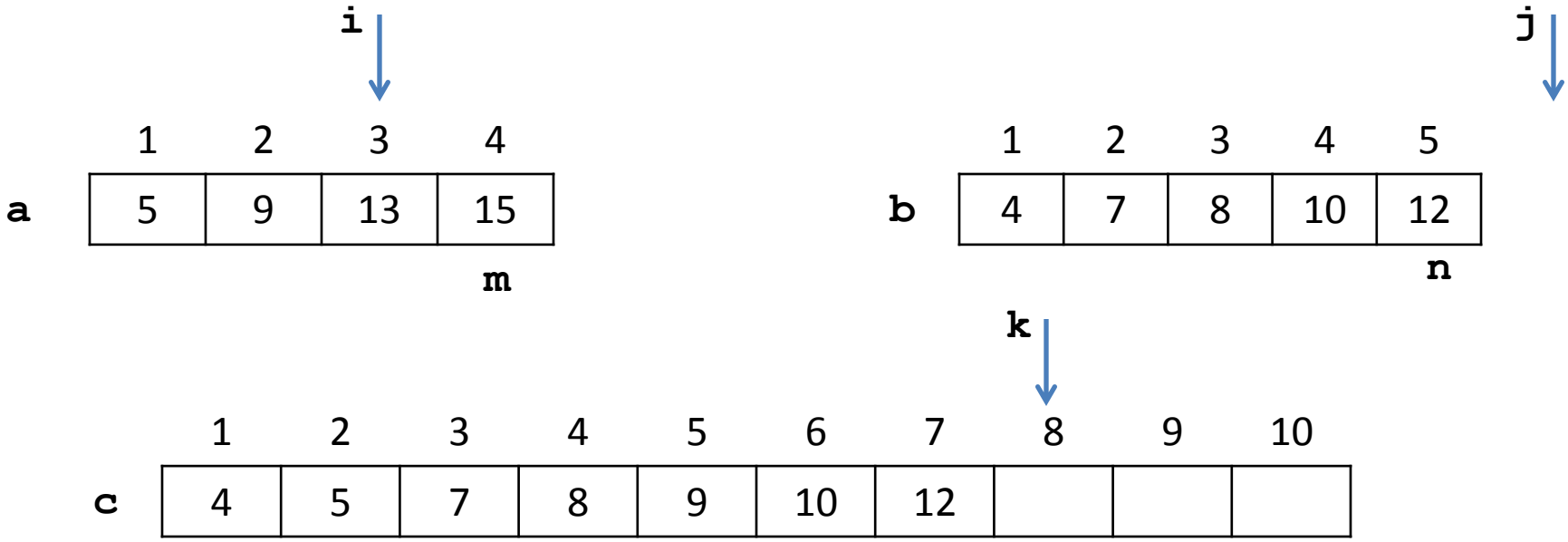
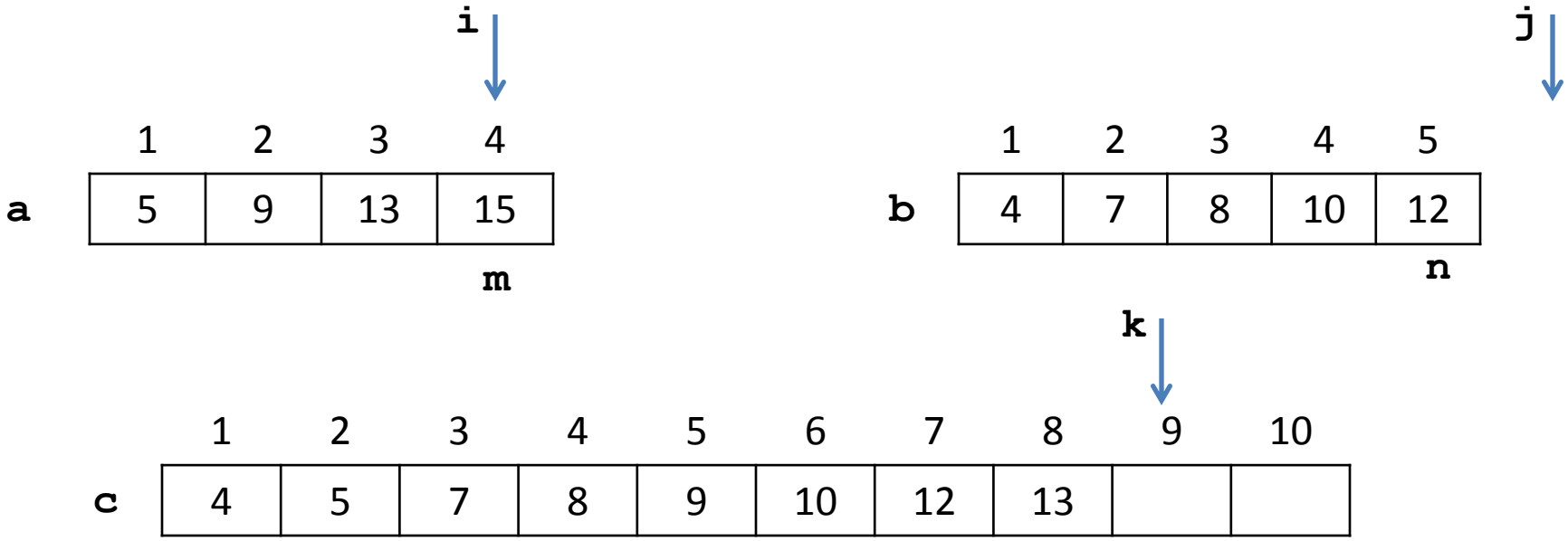


1. When we reach end of one of the lists, elements are copied from the other list, till the end of the source list is reached.



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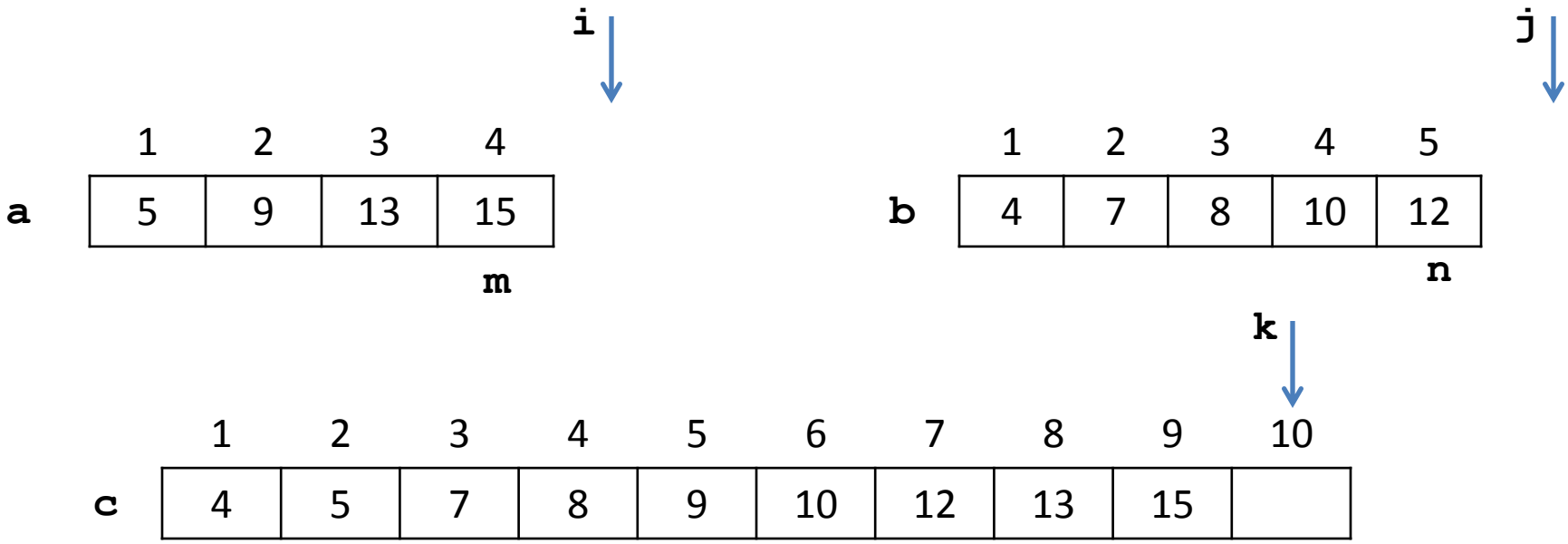
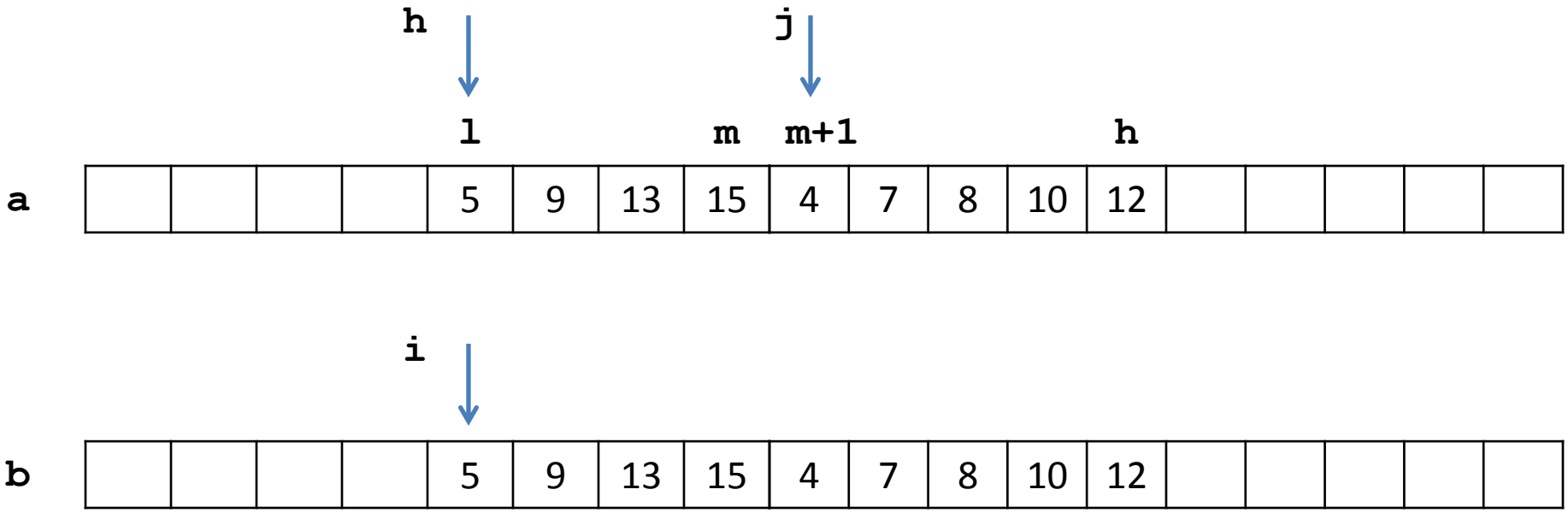


1. When we reach end of one of the lists, elements are copied from the other list, till the end of the source list is reached.
2. Copy the elements from i to m or j to n

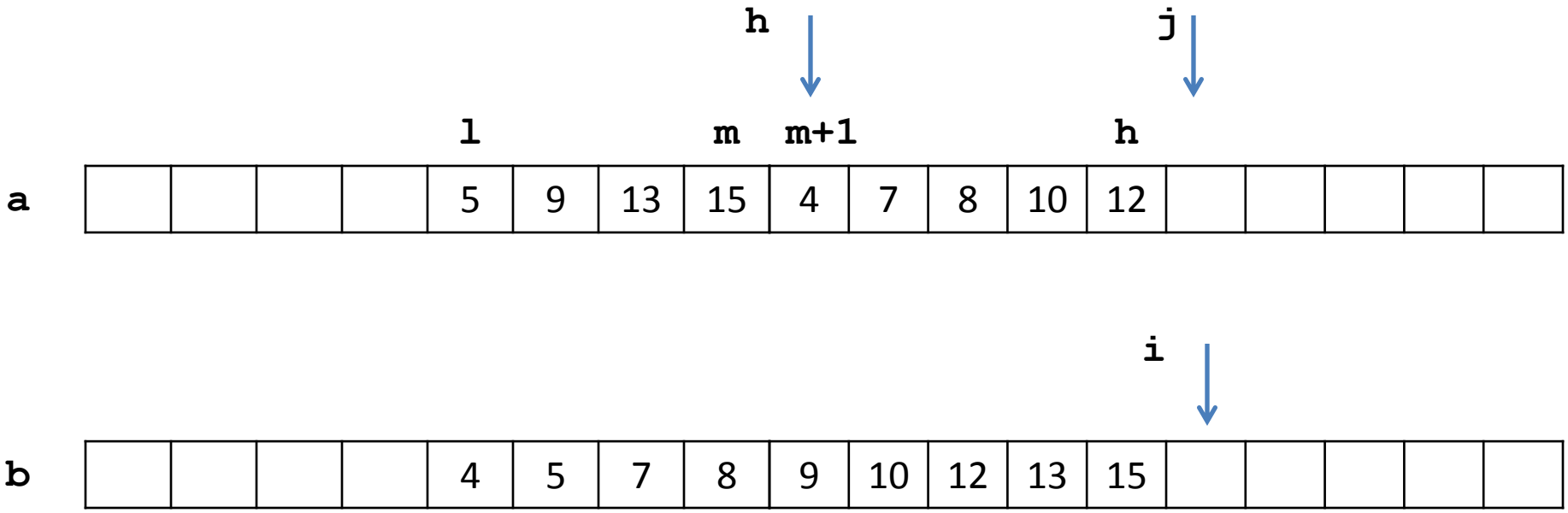
```
i:=1;j:=1;k:=1;
while(i<=m and j<=n) do
{
    if(a[i]<=b[j])then
    {
        c[k]:=a[i];i:=i+1;
    }
    else
    {
        c[k]:=b[j];j:=j+1;
    }
    k:=k+1;
}
for p:= i to m  do
{
    c[k]:=a[p];k:=k+1;
}
for p:= j to n  do
{
    c[k]:=b[p];j:=j+1;
}
```

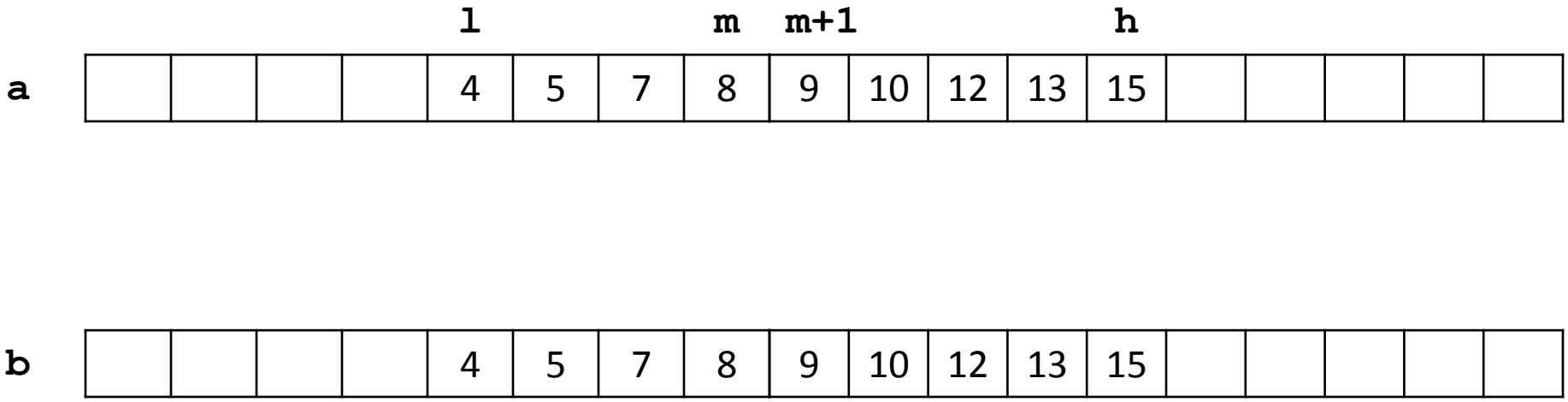
Merging in Merge Sort

- Merging in Merge Sort has a small variation.
- The two lists to be merged are logical partitions in the same array.
- A temporary array is used to store the merged list of the two logical partitions.
- The merged list from the temporary array is copied back.



low(l)
mid(m)
high(h)


h
j
l
m
m+1
h
a
5
9
13
15
4
7
8
10
12
i
b
4
5
7
8
9
10
12
13
15


l
m
m+1
h
a
4
5
7
8
9
10
12
13
15
b
4
5
7
8
9
10
12
13
15

```
Algorithm Merge(low,mid,high)
//a[low:mid] and a[mid+1:high] are two sorted logical
//partitions of global array a
//These two partitions are merged to array b[low:high]
//The elements are copied from b[low:high] to a[low:high]
{
    h:=low;i:=low;j:=mid+1;
    while(h<=mid and j<=high) do
    {
        if(a[h]<=a[j])then
        {
            b[i]:=a[h];h:=h+1;
        }
        else
        {
            b[i]:=a[j];j:=j+1;
        }
        i:=i+1;
    }
```

```
    for k:= h to mid  do
    {
        b[i]:=a[k];i:=i+1;
    }
    for k:= j to high  do
    {
        b[i]:=a[k];i:=i+1;
    }
    for k:= low to high  do
        a[k]:=b[k];
    }

}
```

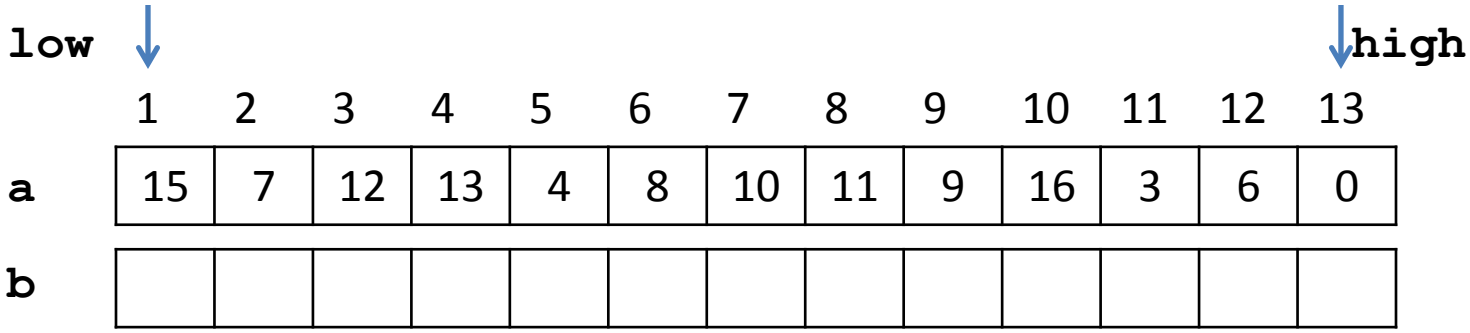
Merge Sort

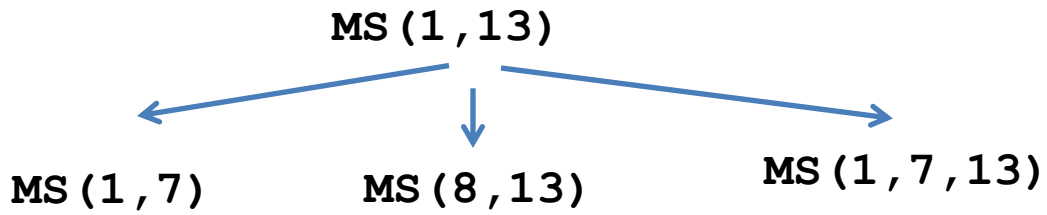
- Given a sequence of **`n`** elements **`a[1:n]`**, the objective is to sort them in non-decreasing order.
- The recursive formulation of Merge Sort is **`MergeSort(low,high)`**:

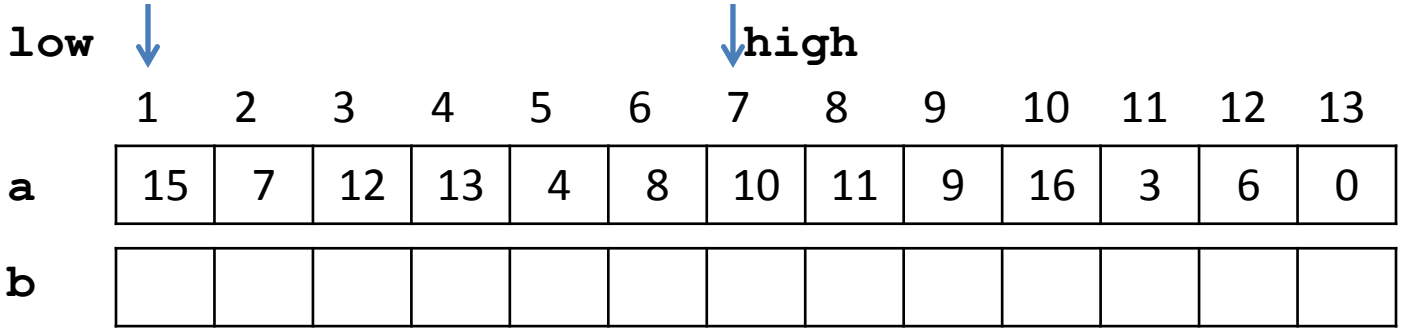
Merge Sort

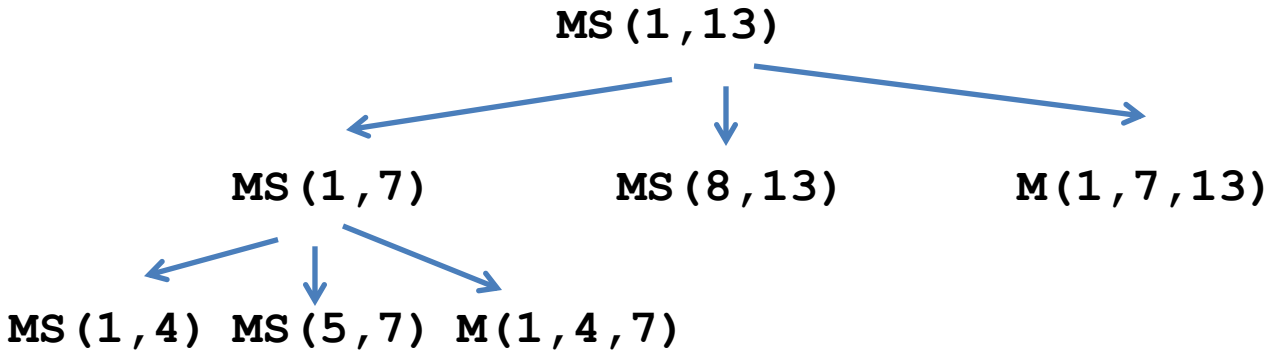
- **`MergeSort(low,high)`**:
- A list with one element is Merge Sorted.(base case)

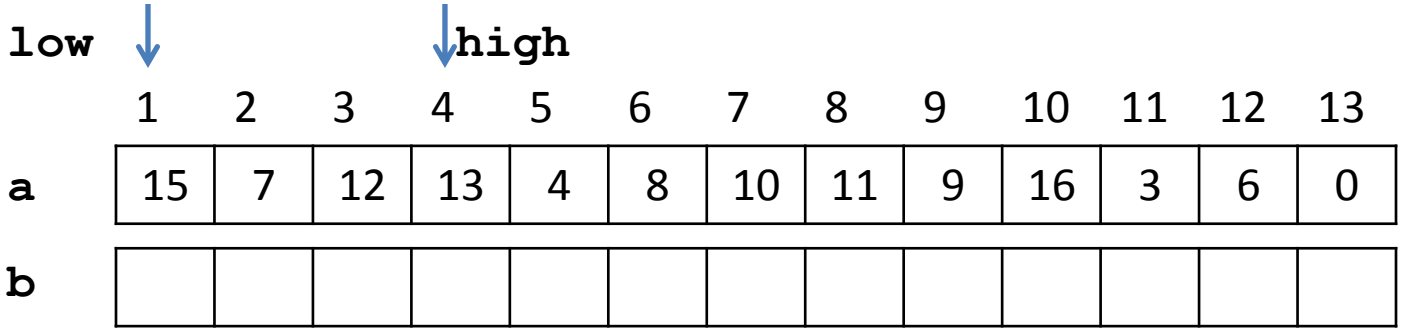
  `if (low<high)then:`
  - Divide the list into two logical partitions.
  - **`mid=(low+high)/2`**
  - Merge Sort each partition.
  - **`MergeSort(low,mid)`**
  - **`MergeSort(mid+1,high)`**
  - Merge two sorted partitions into one.
  - **`Merge(low,mid,high)`**

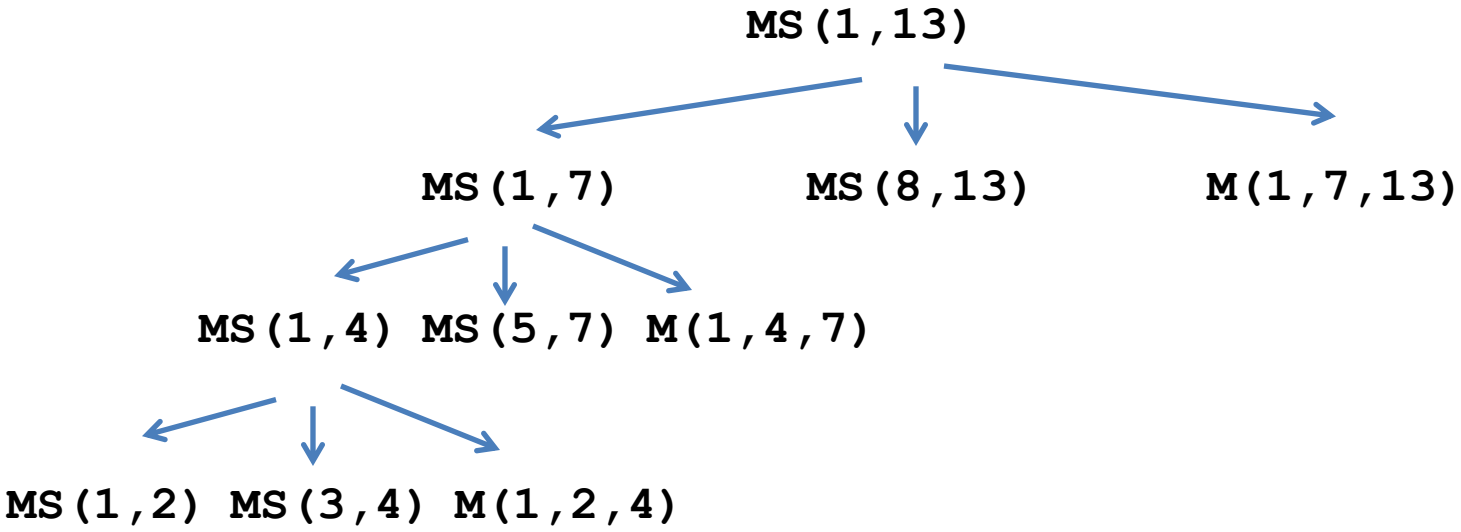

low
high
1 2 3 4 5 6 7 8 9 10 11 12 13
a
15 7 12 13 4 8 10 11 9 16 3 6 0
b

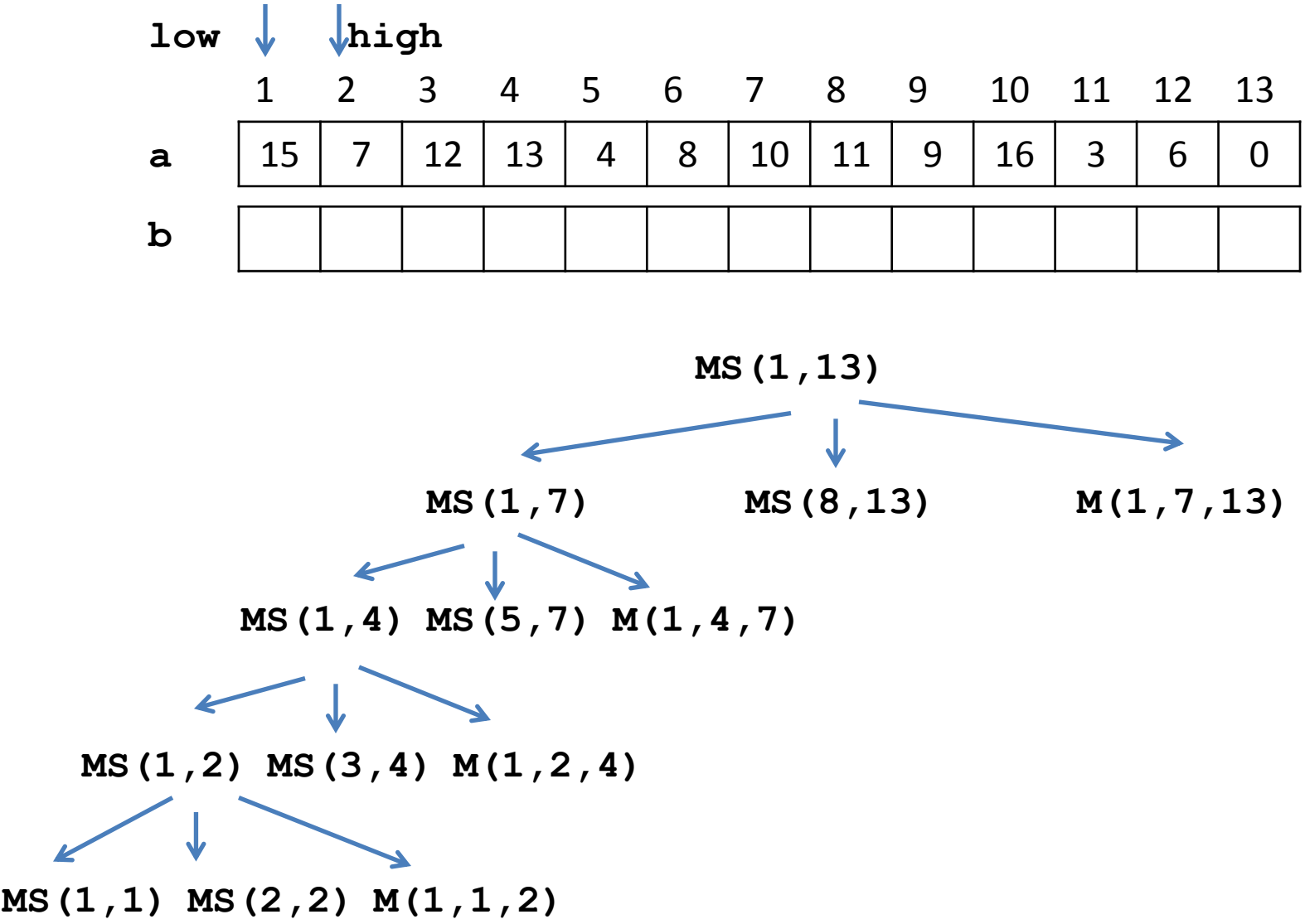

MS(1,13)
MS(1,7)
MS(8,13)
MS(1,7,13)

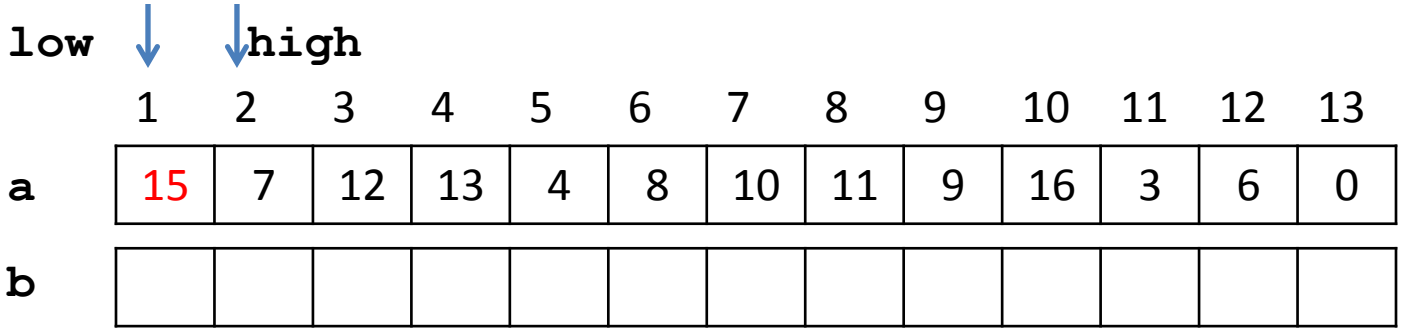

low
high
1 2 3 4 5 6 7 8 9 10 11 12 13
a
15 7 12 13 4 8 10 11 9 16 3 6 0
b


MS(1,13)
MS(1,7)
MS(8,13)
M(1,7,13)
MS(1,4)
MS(5,7)
M(1,4,7)


low
high
1 2 3 4 5 6 7 8 9 10 11 12 13
a
15 7 12 13 4 8 10 11 9 16 3 6 0
b


MS(1,13)
MS(1,7)
MS(8,13)
M(1,7,13)
MS(1,4)
MS(5,7)
M(1,4,7)
MS(1,2)
MS(3,4)
M(1,2,4)


low
high
1 2 3 4 5 6 7 8 9 10 11 12 13
a
15 7 12 13 4 8 10 11 9 16 3 6 0
b
MS(1,13)
MS(1,7)
MS(8,13)
M(1,7,13)
MS(1,4)
MS(5,7)
M(1,4,7)
MS(1,2)
MS(3,4)
M(1,2,4)
MS(1,1)
MS(2,2)
M(1,1,2)


low
high
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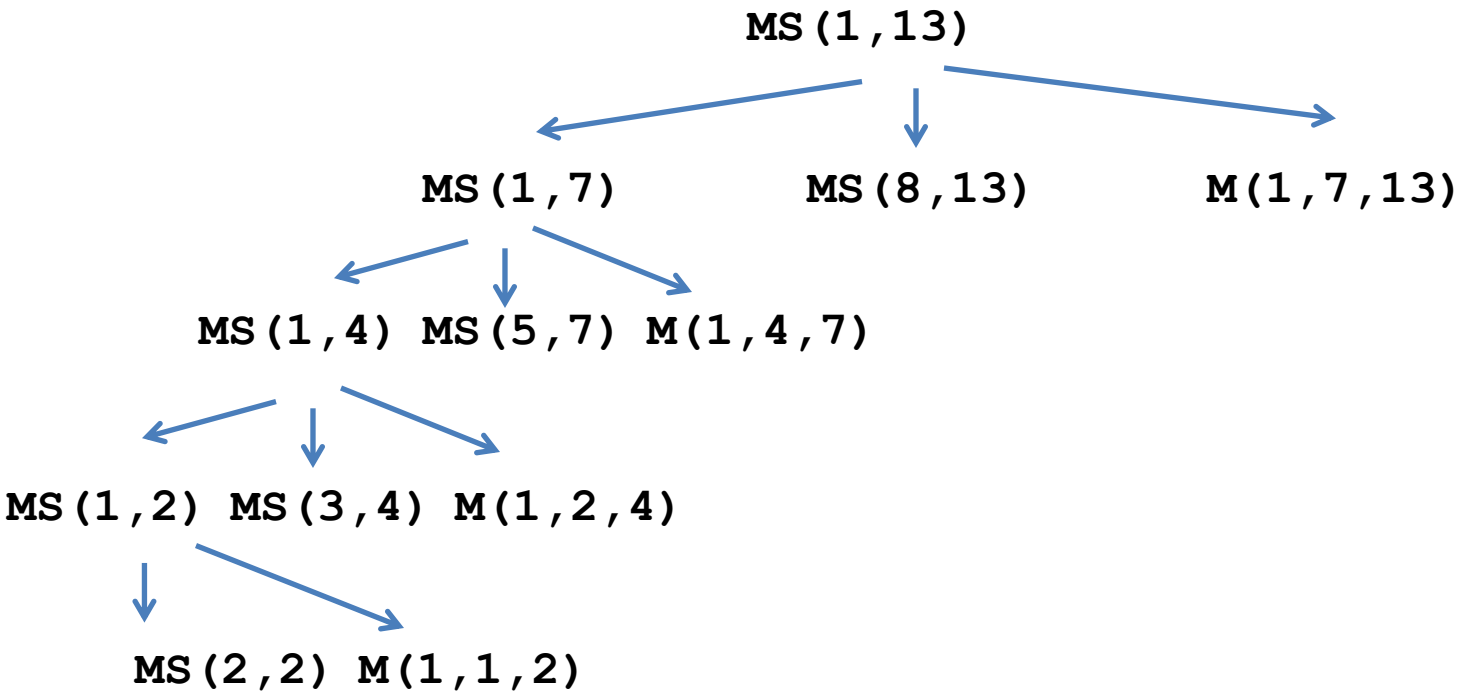
15 7 12 13 4 8 10 11 9 16 3 6 0
b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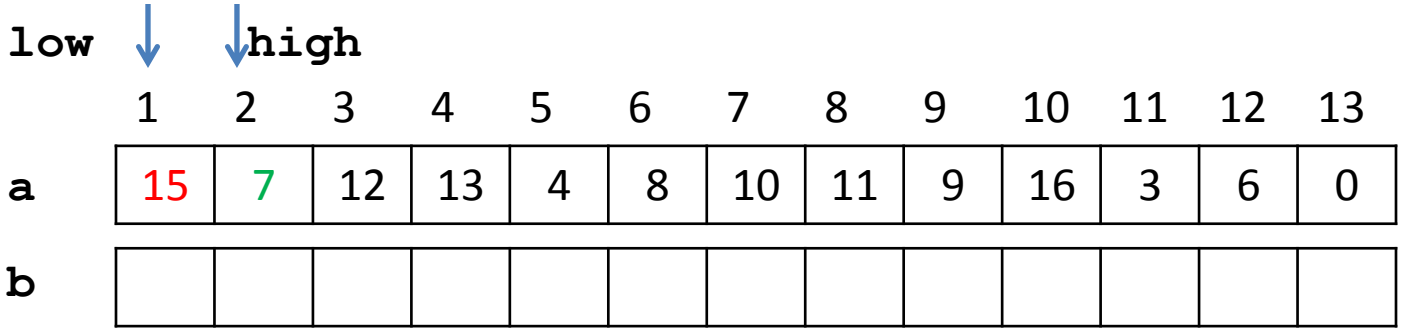
M(1,4,7)
MS(1,2)
MS(3,4)
M(1,2,4)
MS(2,2)
M(1,1,2)


low
high
1 2 3 4 5 6 7 8 9 10 11 12 13
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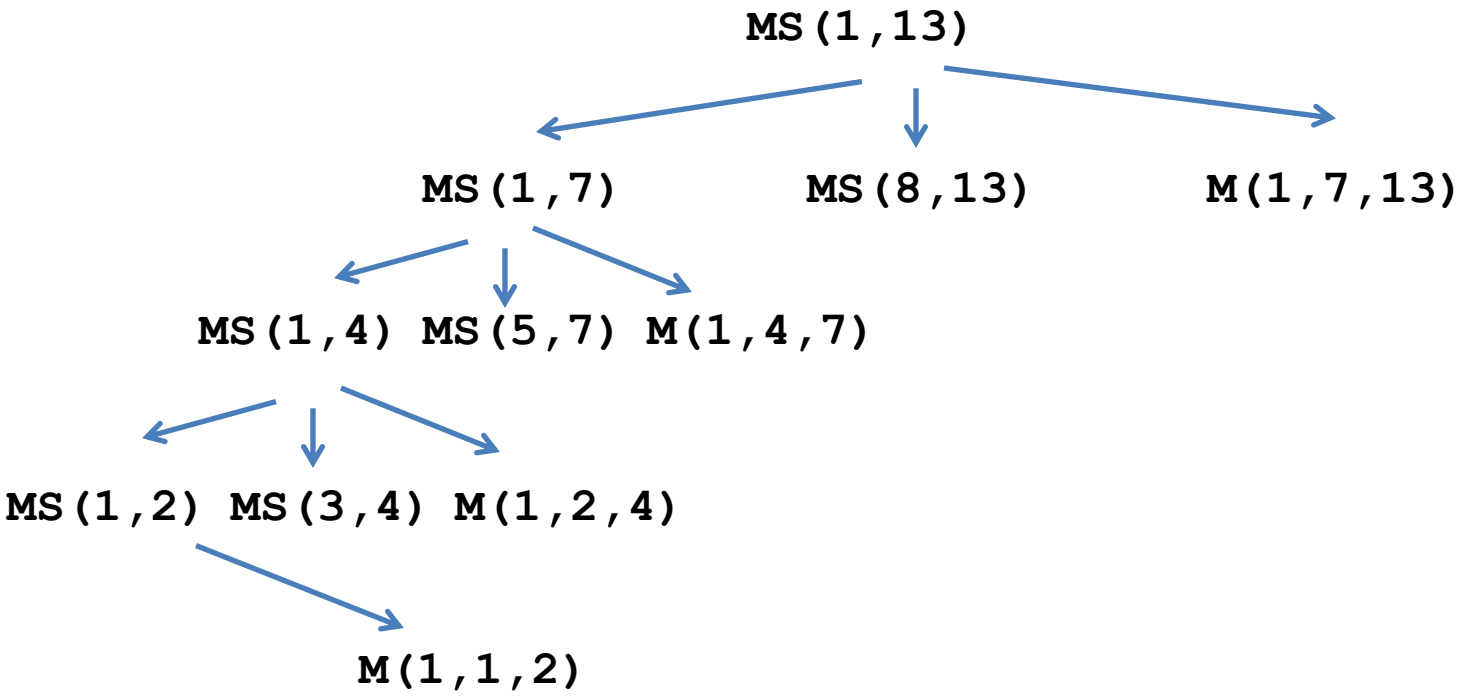
15 7 12 13 4 8 10 11 9 16 3 6 0
b


MS(1,13)
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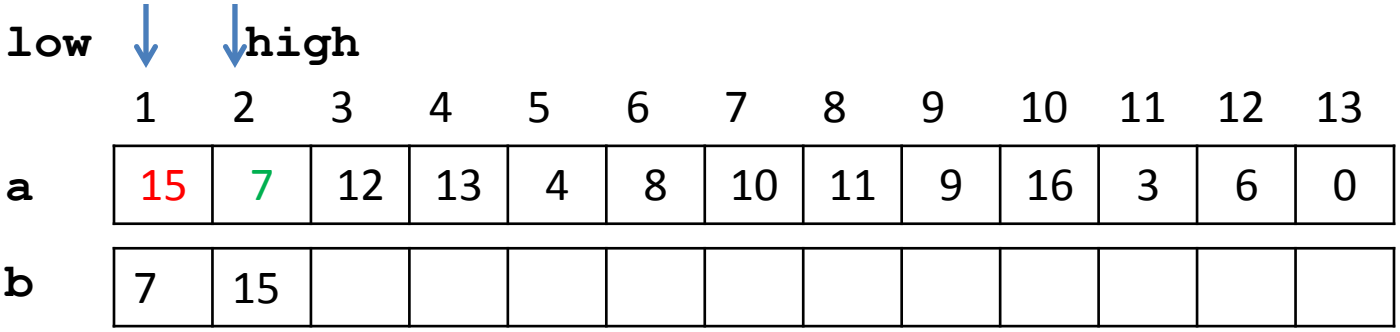
M(1,4,7)
MS(1,2)
MS(3,4)
M(1,2,4)
M(1,1,2)


low
high
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a
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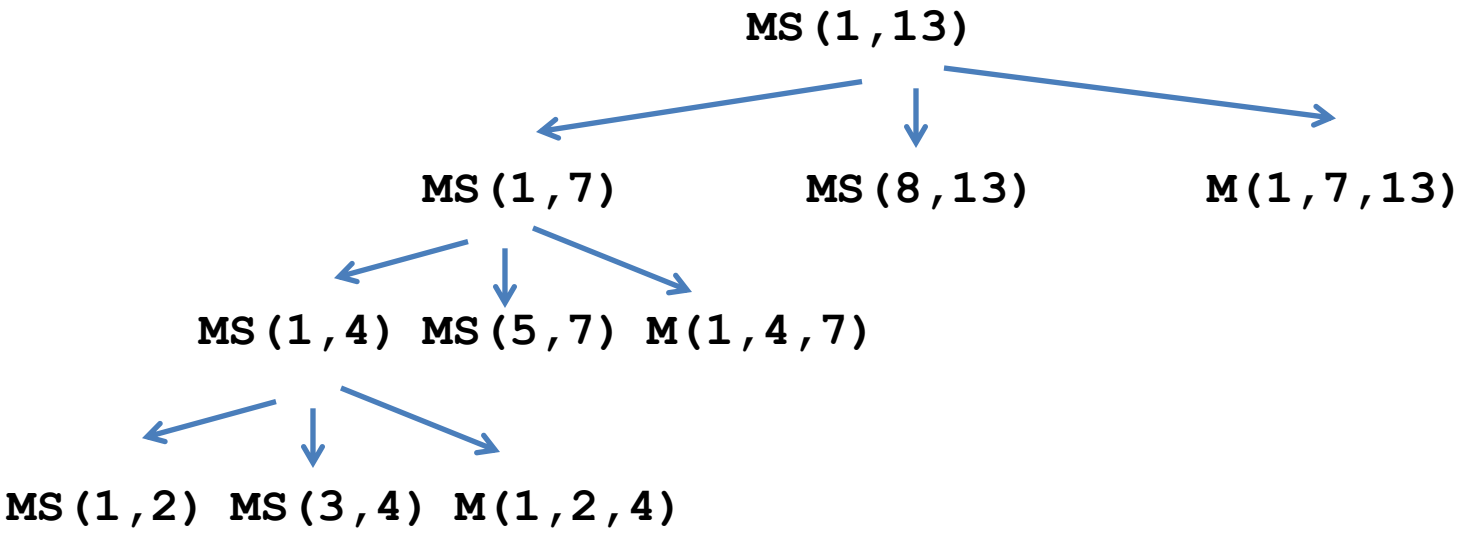
b
7 15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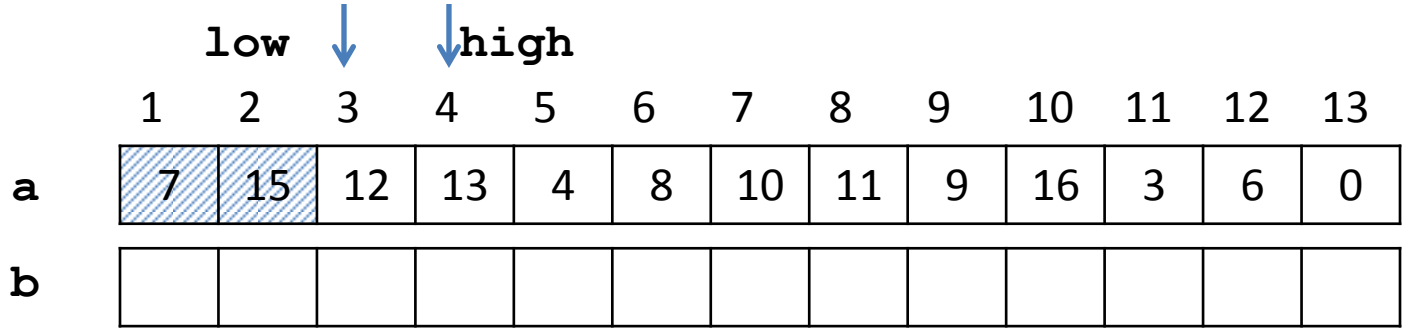
M(1,4,7)
MS(1,2)
MS(3,4)
M(1,2,4)


low
high
1 2 3 4 5 6 7 8 9 10 11 12 13
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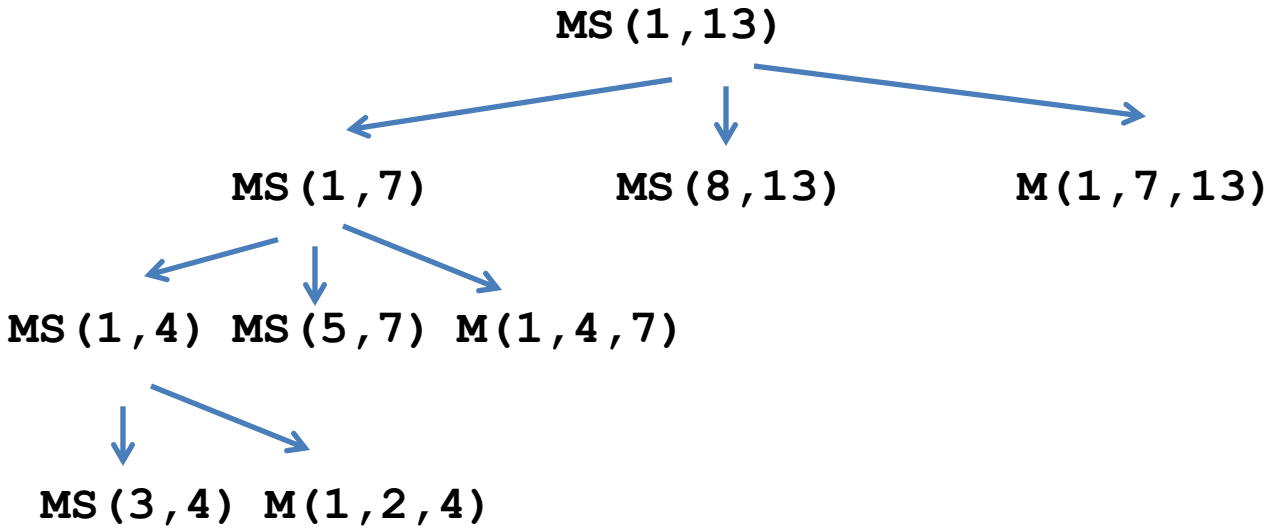
b


MS(1,13)
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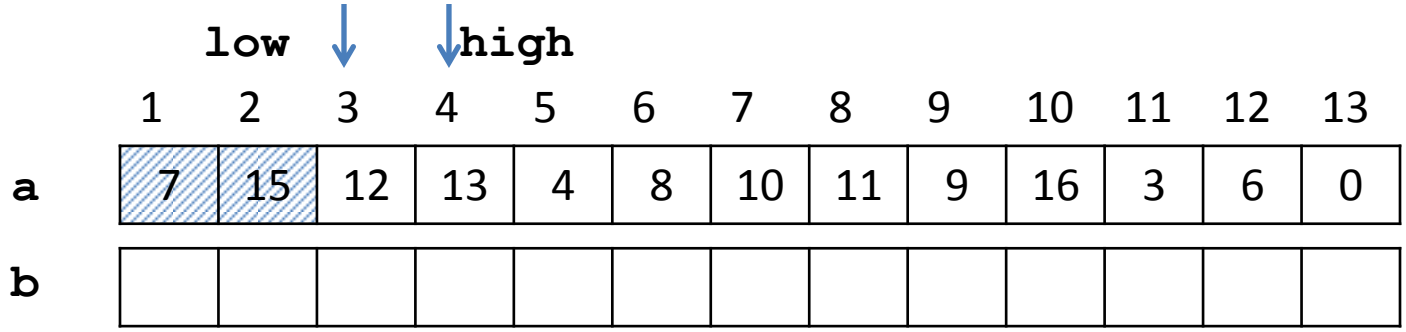
MS(5,7)
M(1,4,7)
MS(3,4)
M(1,2,4)


low
high
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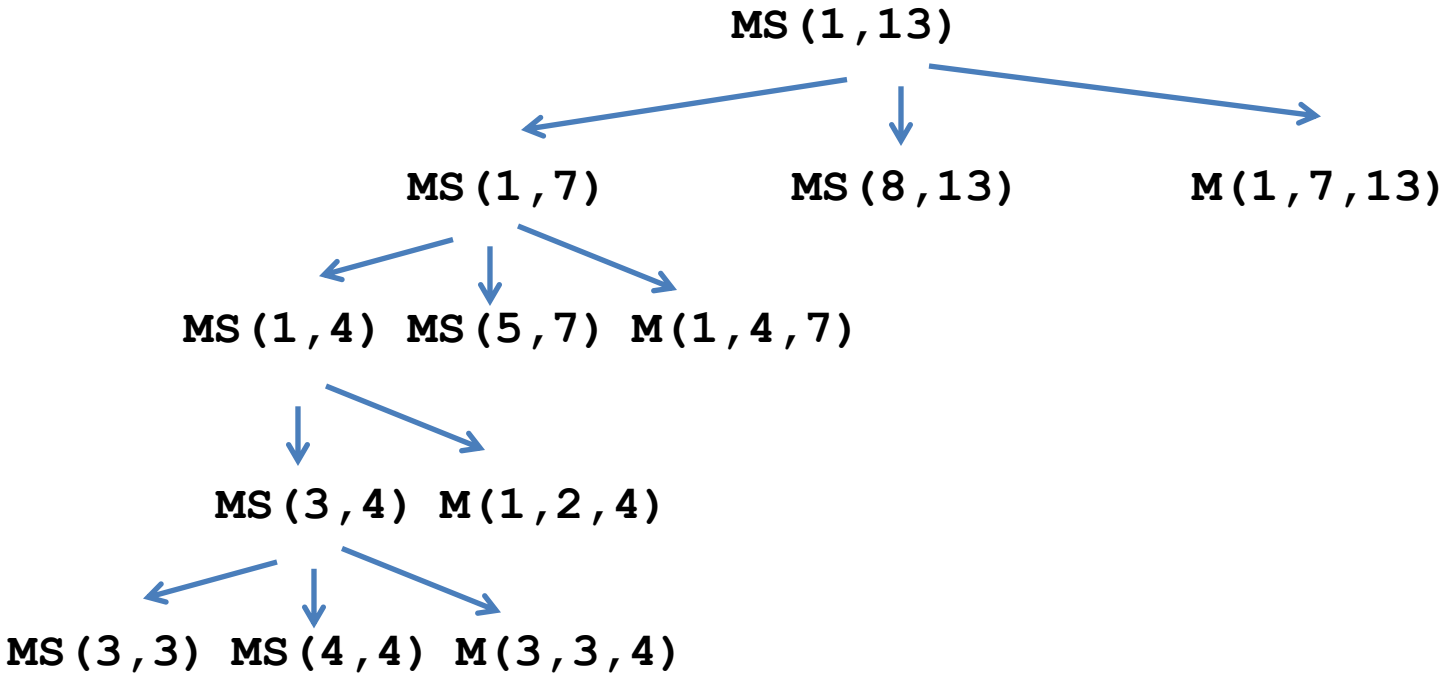
a
7 15 12 13 4 8 10 11 9 16 3 6 0
b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MS(1,4)
MS(5,7)
M(1,4,7)
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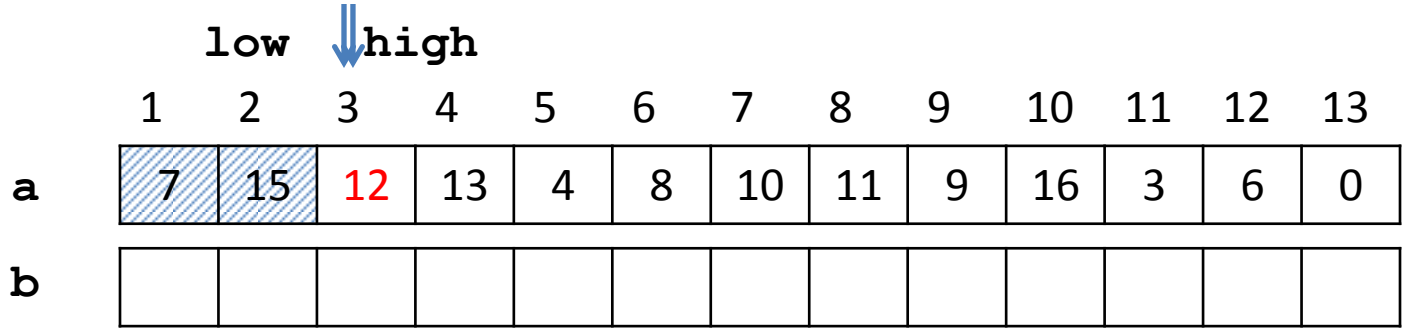
M(1,2,4)
MS(3,3)
MS(4,4)
M(3,3,4)


low
high
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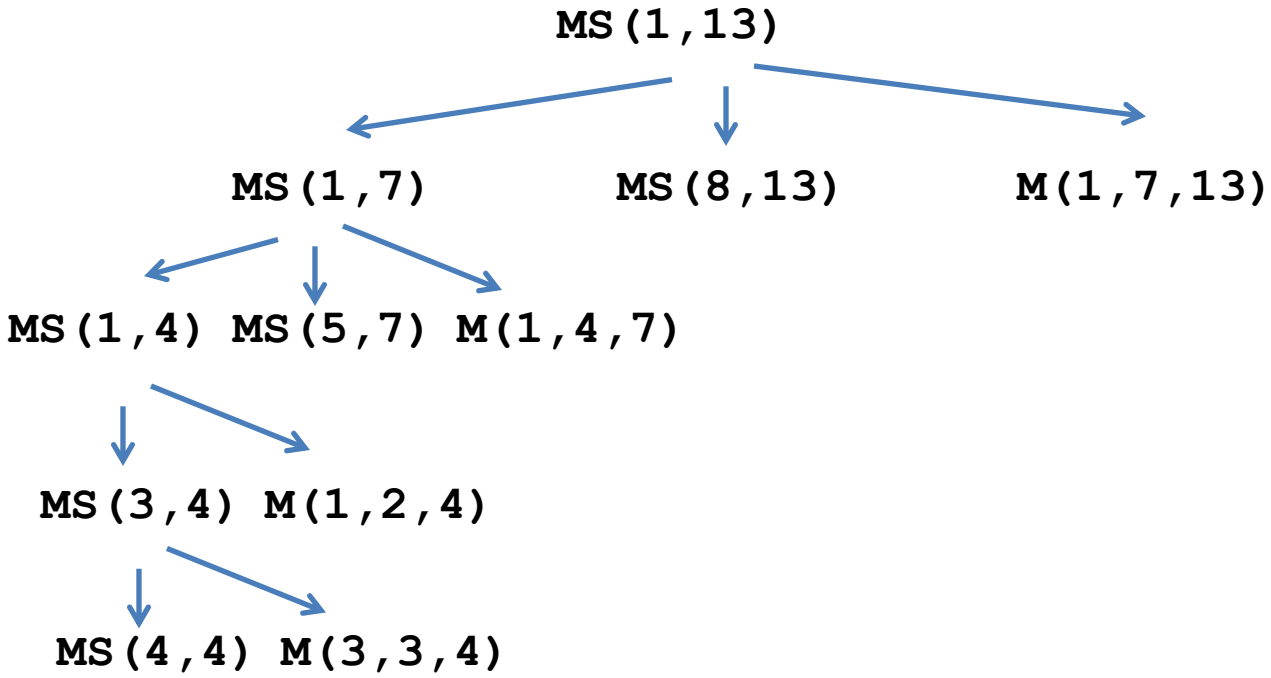
7 15 12 13 4 8 10 11 9 16 3 6 0
b


MS(1,13)
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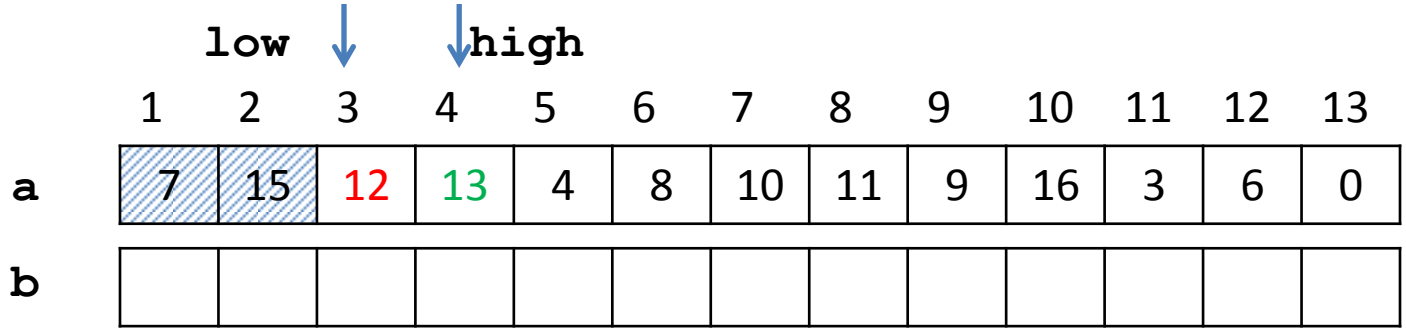
MS(1,4)
MS(5,7)
M(1,4,7)
MS(3,4)
M(1,2,4)
MS(4,4)
M(3,3,4)


low
high
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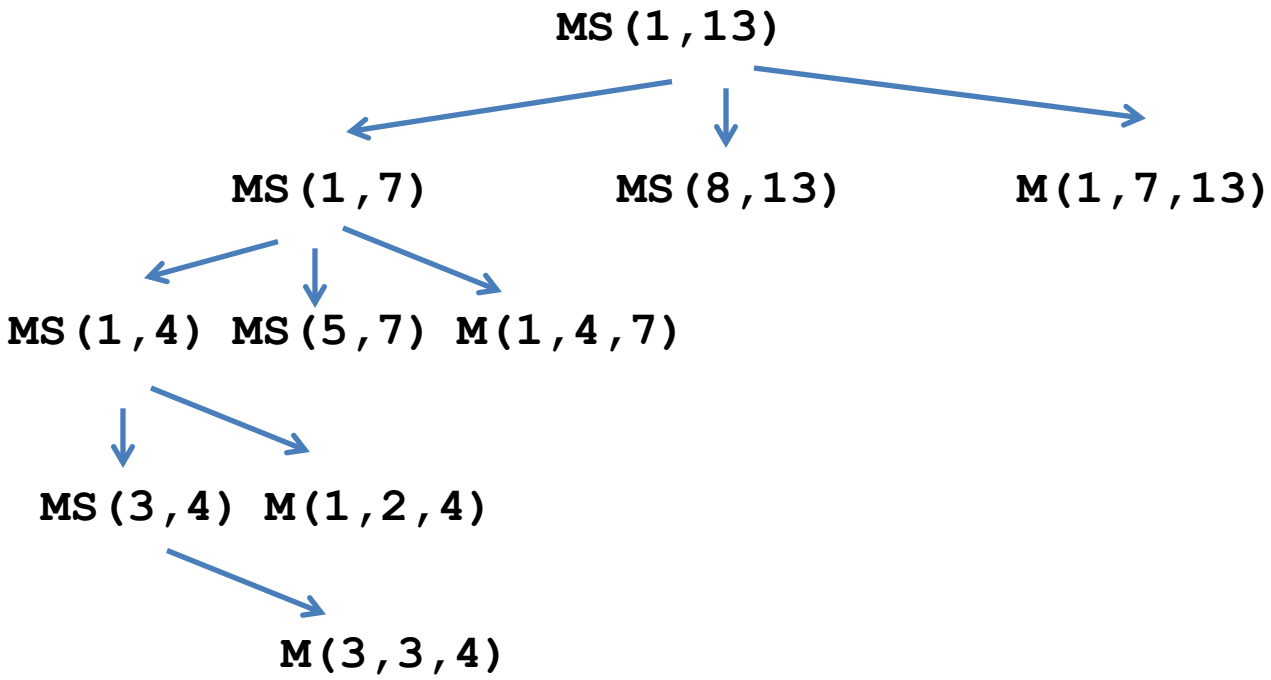
b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MS(8,13)
M(1,7,13)
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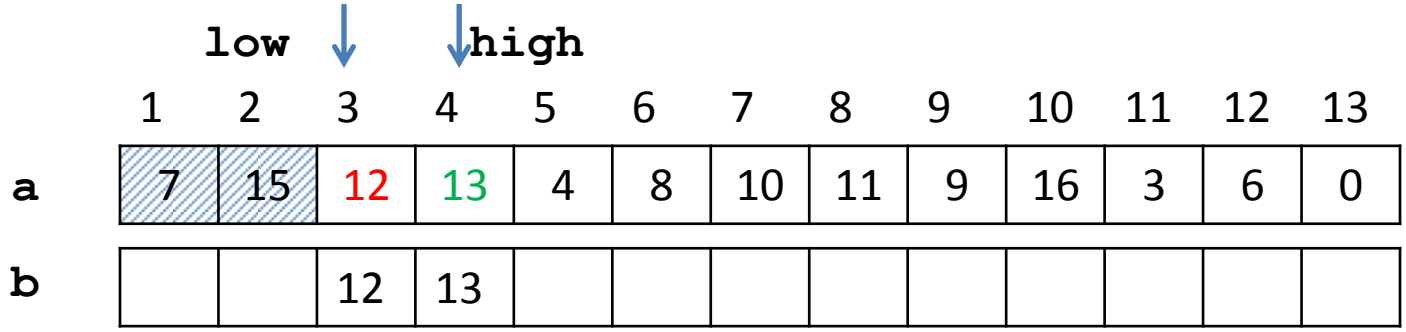
MS(5,7)
M(1,4,7)
MS(3,4)
M(1,2,4)
M(3,3,4)


low
high
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a
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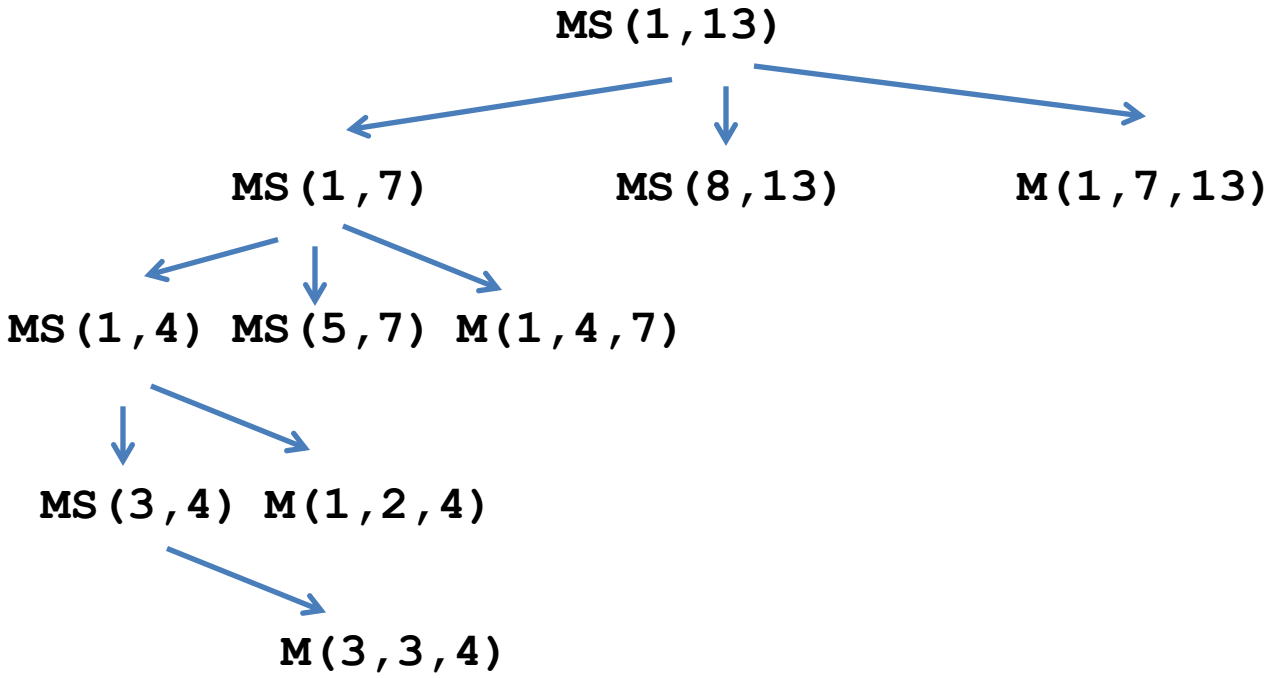
b
12 13


MS(1,13)
MS(1,7)
MS(8,13)
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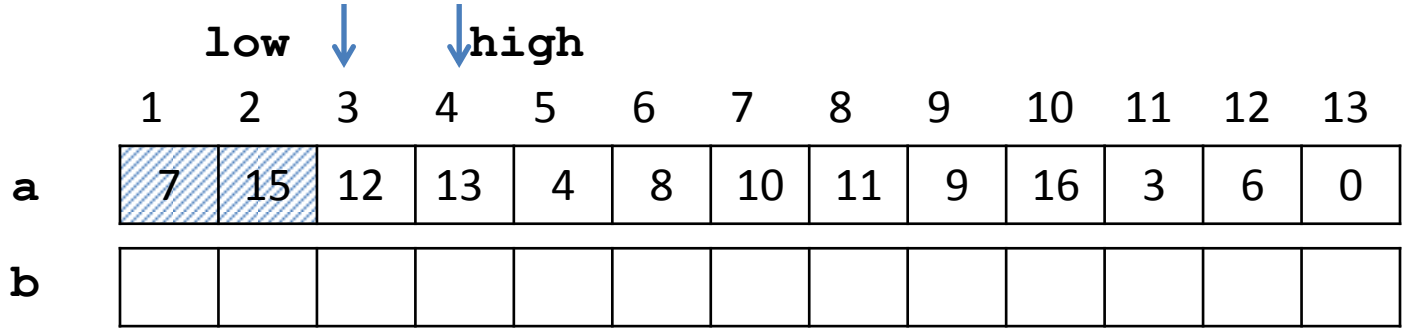
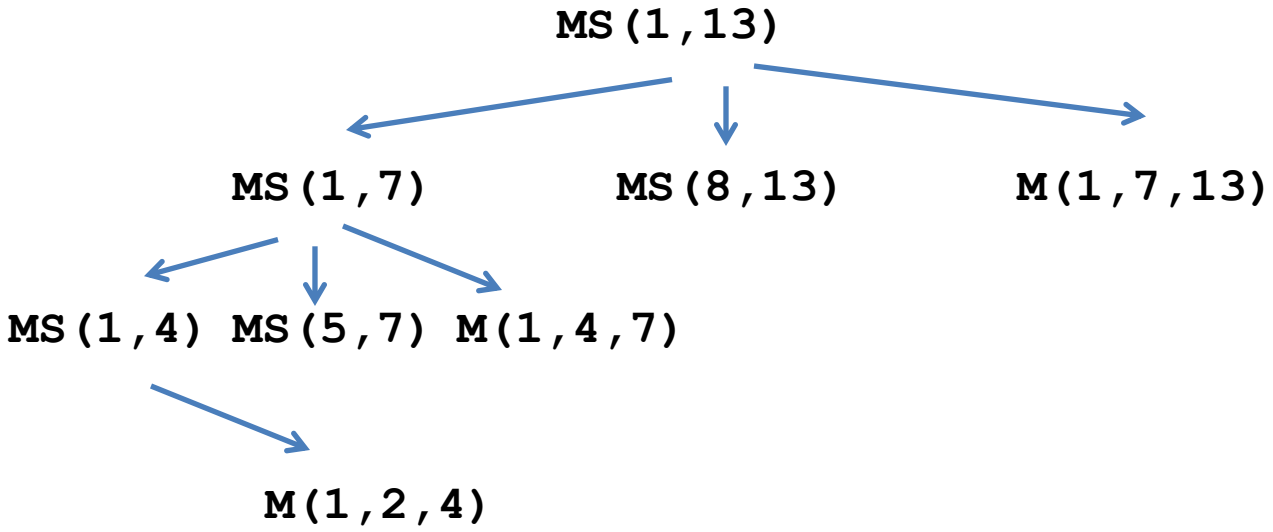
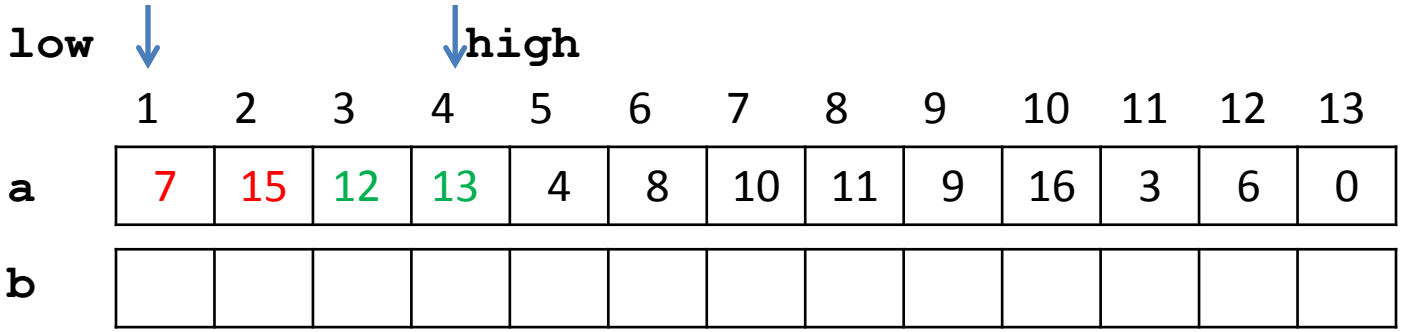
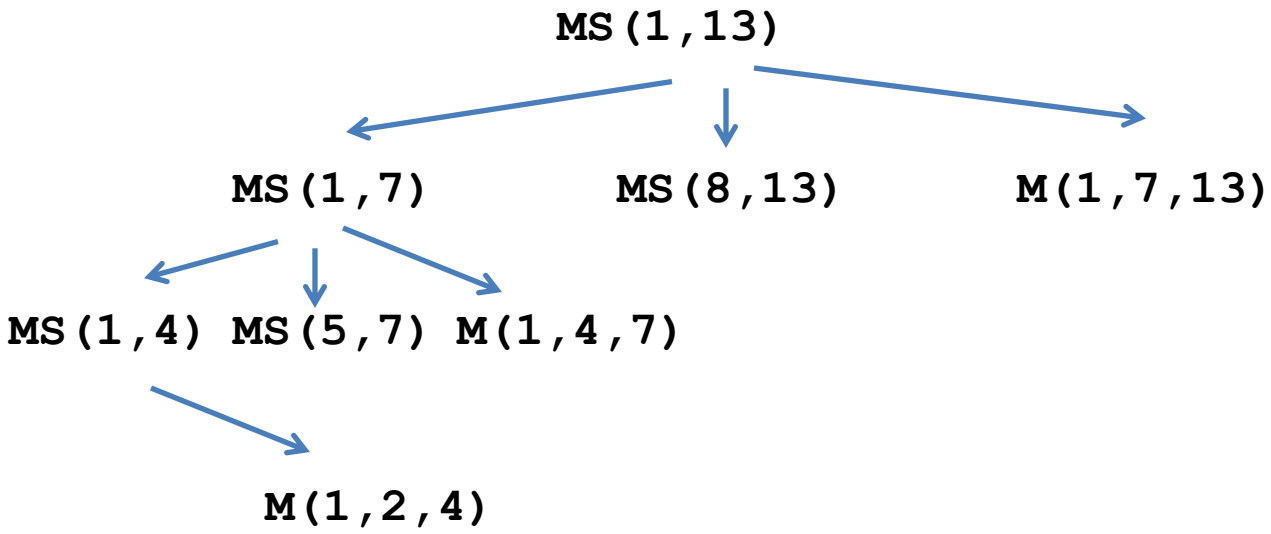
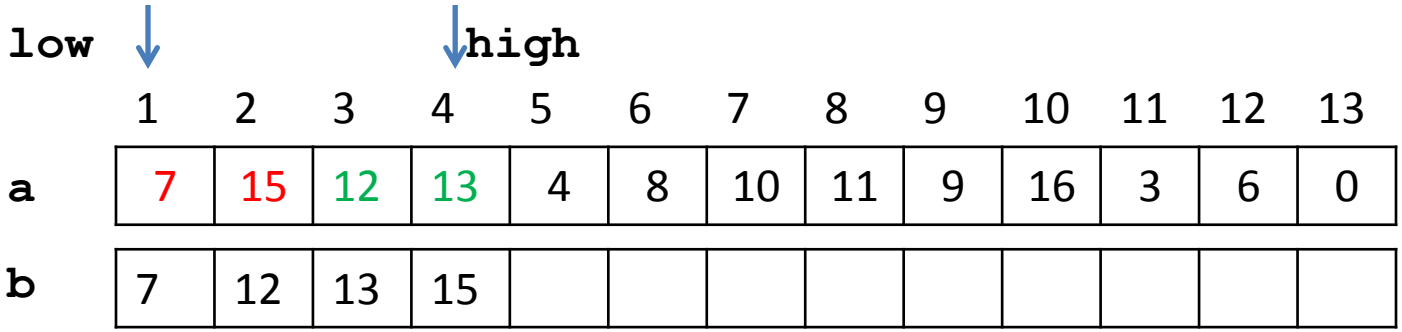
MS(1,4)
MS(5,7)
M(1,4,7)
MS(3,4)
M(1,2,4)
M(3,3,4)

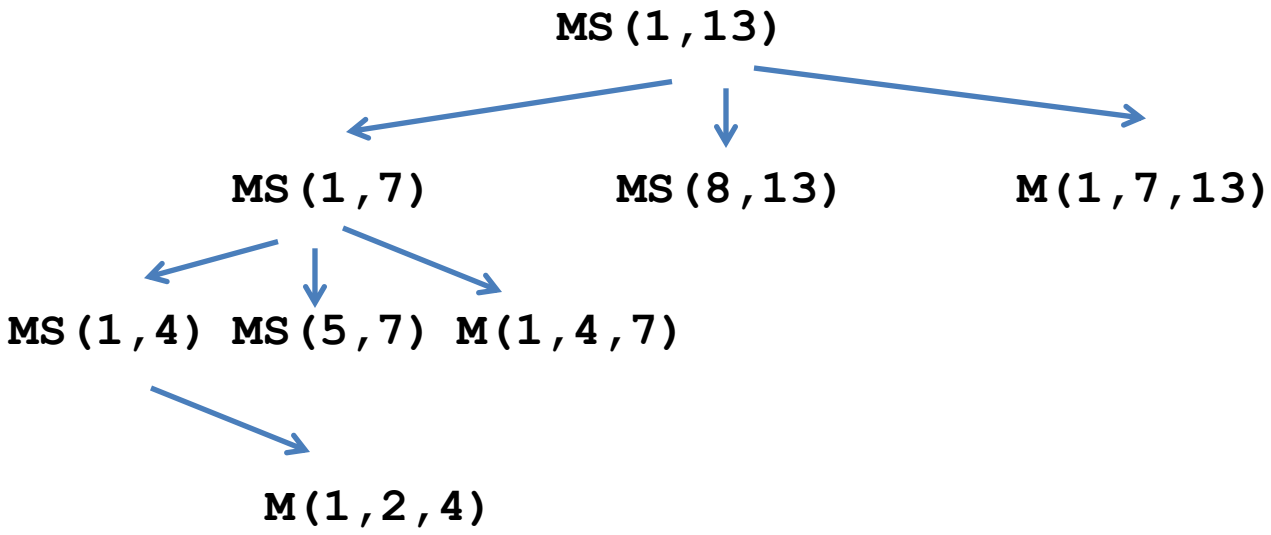

low
high
1 2 3 4 5 6 7 8 9 10 11 12 13
a
7 15 12 13 4 8 10 11 9 16 3 6 0
b

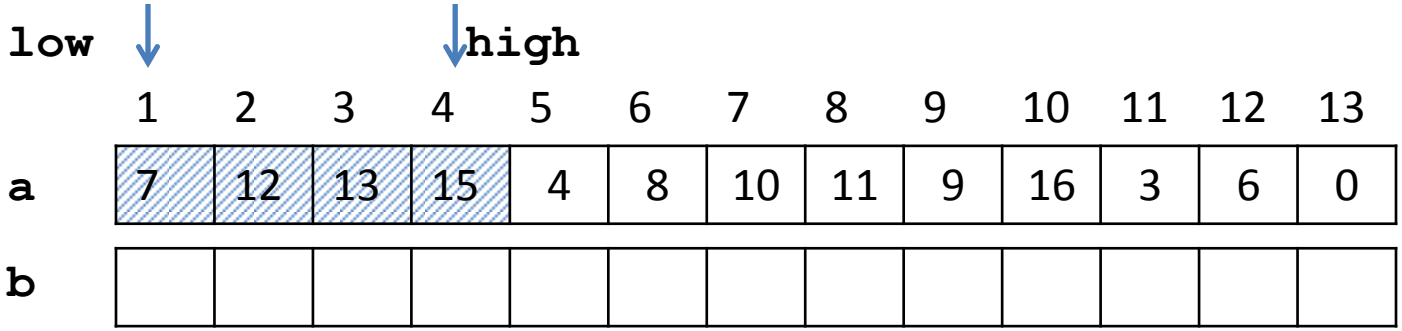

MS(1,13)
MS(1,7)
MS(8,13)
M(1,7,13)
MS(1,4)
MS(5,7)
M(1,4,7)
M(1,2,4)

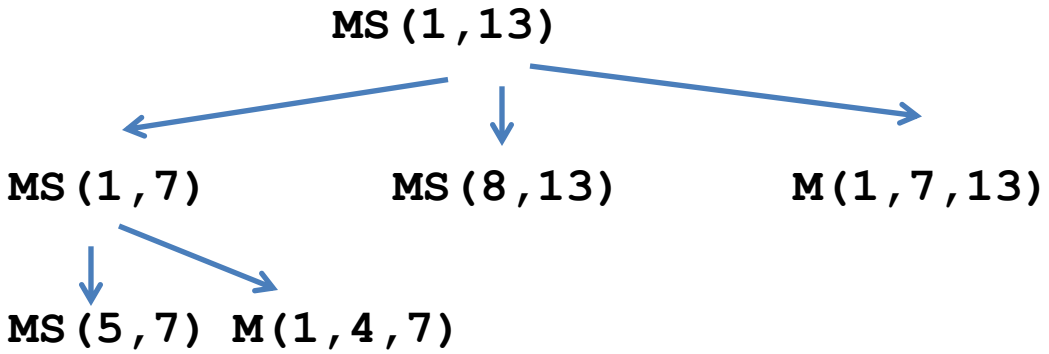

low
high
1 2 3 4 5 6 7 8 9 10 11 12 13
a
7 15 12 13 4 8 10 11 9 16 3 6 0
b

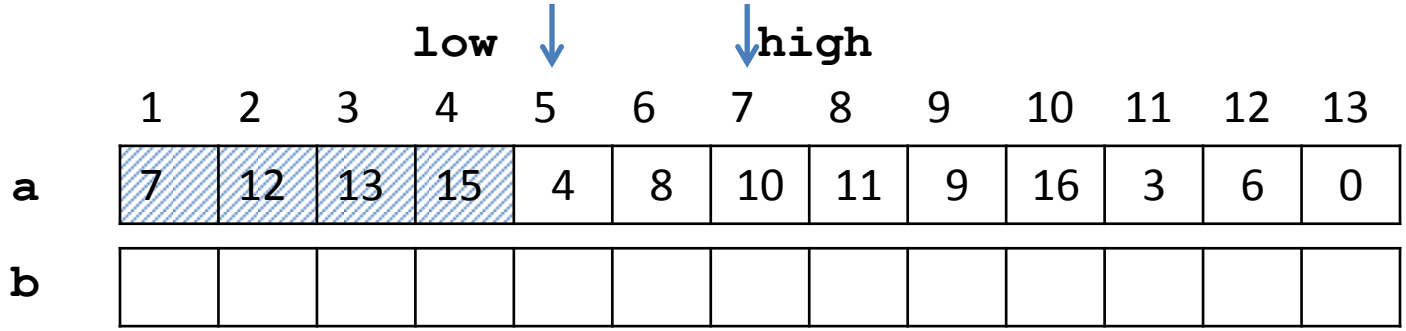

MS(1,13)
MS(1,7)
MS(8,13)
M(1,7,13)
MS(1,4)
MS(5,7)
M(1,4,7)
M(1,2,4)

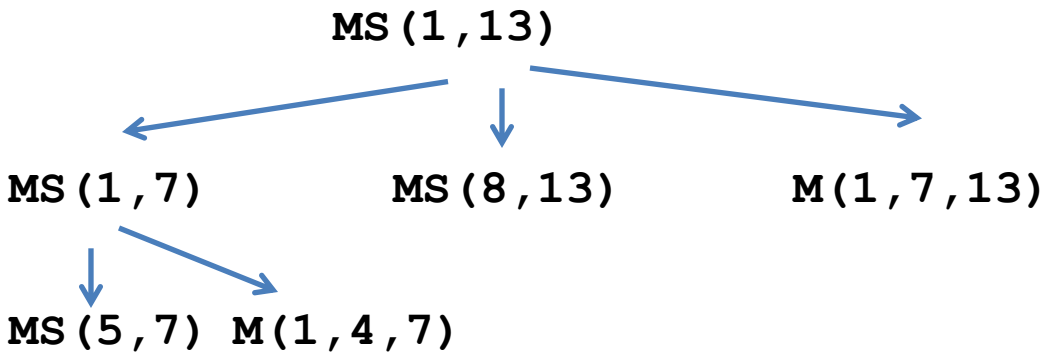

low
high
1 2 3 4 5 6 7 8 9 10 11 12 13
a
7 15 12 13 4 8 10 11 9 16 3 6 0
b
7 12 13 15


MS(1,13)
MS(1,7)
MS(8,13)
M(1,7,13)
MS(1,4)
MS(5,7)
M(1,4,7)
M(1,2,4)


low
high
1 2 3 4 5 6 7 8 9 10 11 12 13
a
7 12 13 15 4 8 10 11 9 16 3 6 0
b


MS(1,13)
MS(1,7)
MS(8,13)
M(1,7,13)
MS(5,7)
M(1,4,7)


low
high
1 2 3 4 5 6 7 8 9 10 11 12 13
a
7 12 13 15 4 8 10 11 9 16 3 6 0
b

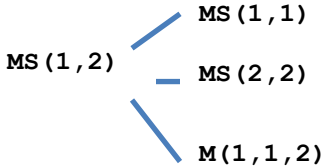

MS(1,13)
MS(1,7)
MS(8,13)
M(1,7,13)
MS(5,7)
M(1,4,7)

MS(1,1)

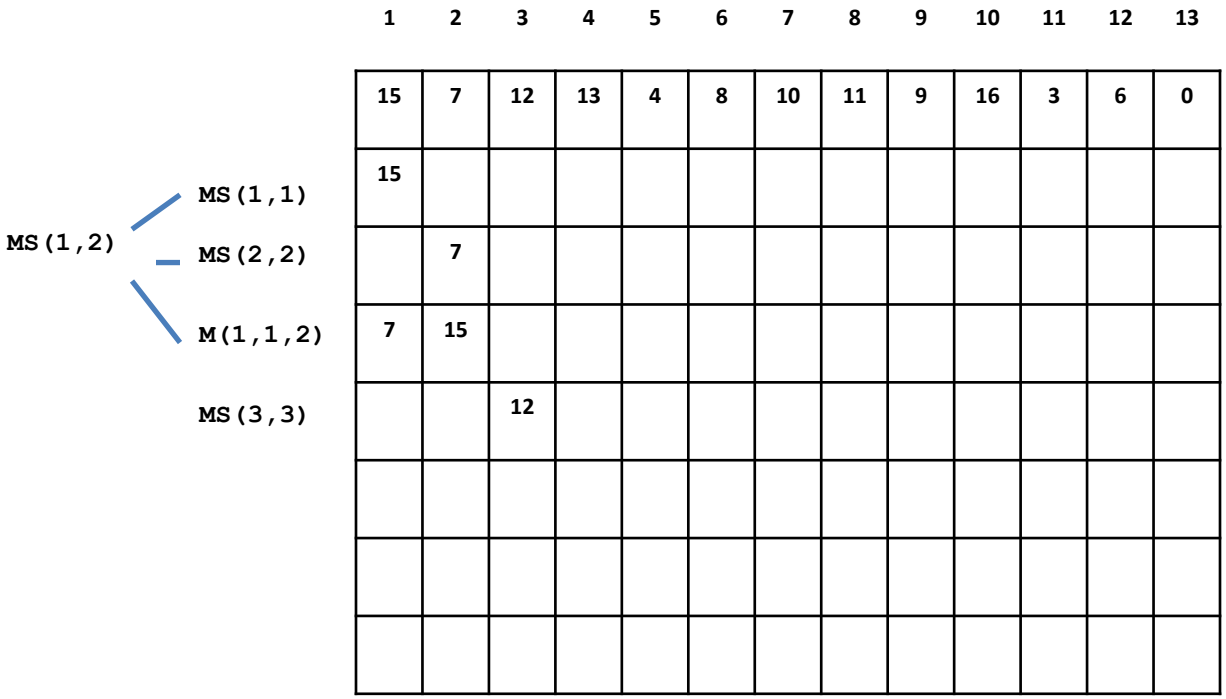
1	2	3	4	5	6	7	8	9	10	11	12	13
15	7	12	13	4	8	10	11	9	16	3	6	0
15												

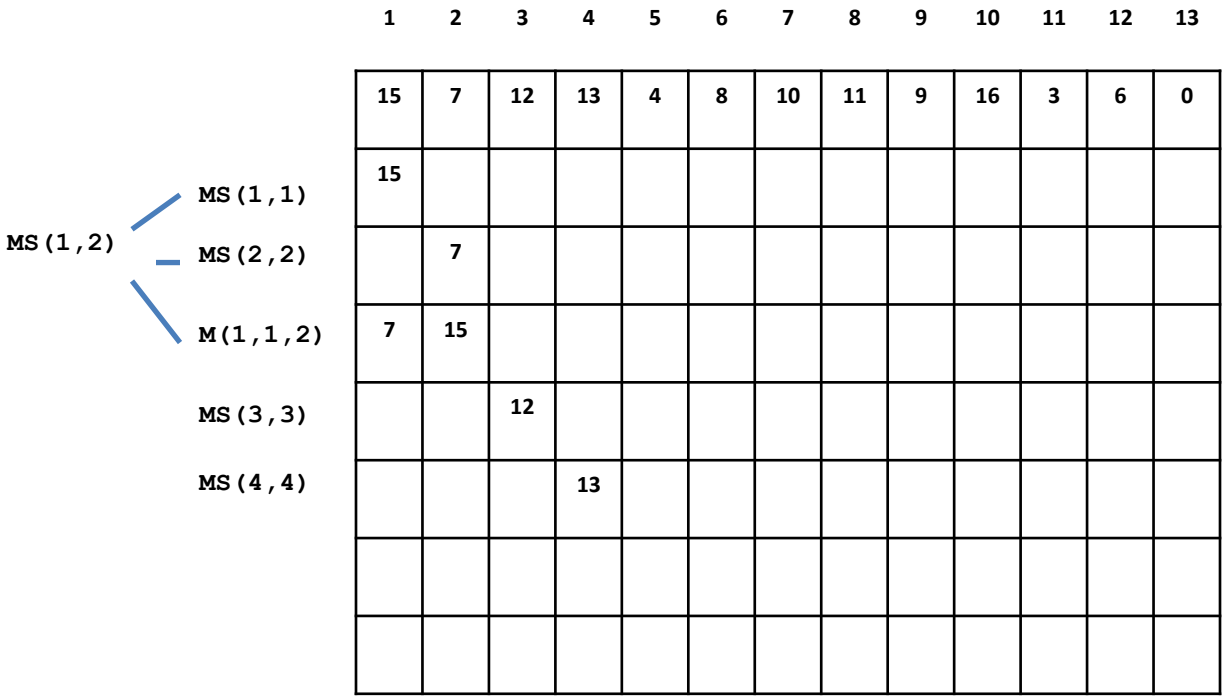
	1	2	3	4	5	6	7	8	9	10	11	12	13
	15	7	12	13	4	8	10	11	9	16	3	6	0
MS(1,1)	15												
MS(2,2)		7											

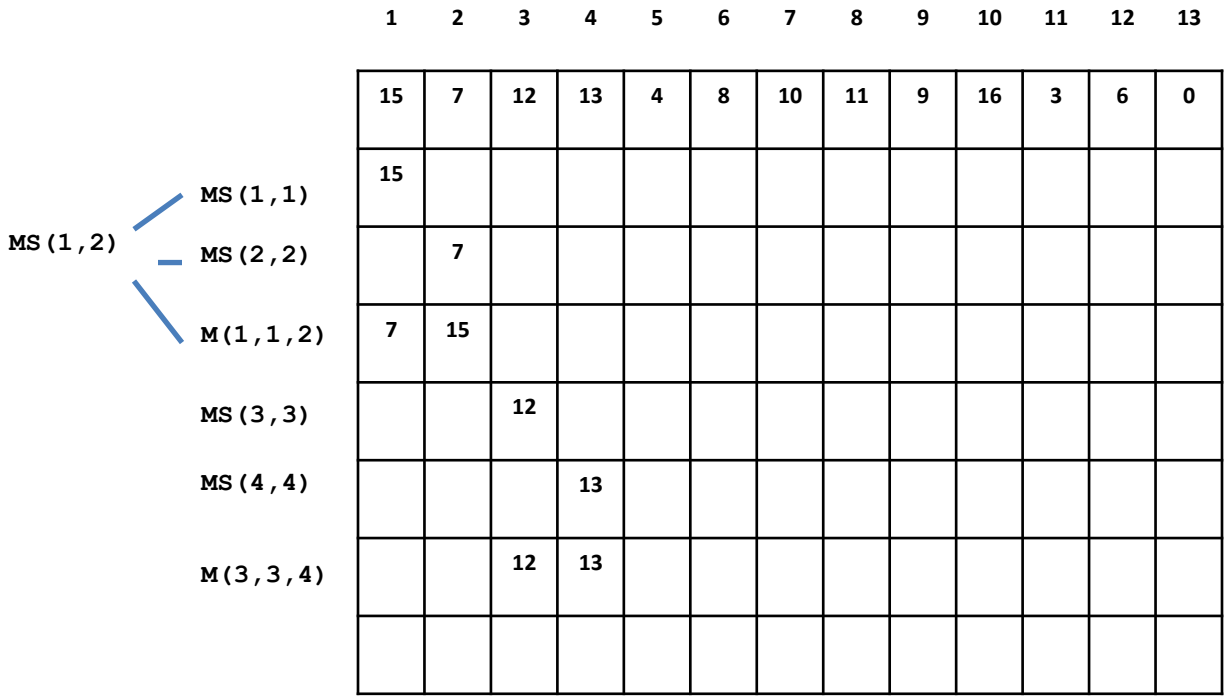
	1	2	3	4	5	6	7	8	9	10	11	12	13
	15	7	12	13	4	8	10	11	9	16	3	6	0
MS(1,1)	15												
MS(2,2)		7											
M(1,1,2)	7	15											

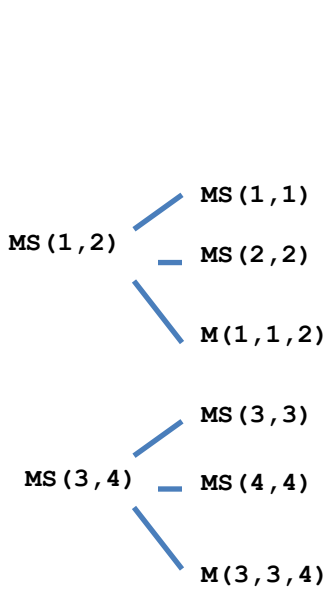


1	2	3	4	5	6	7	8	9	10	11	12	13
15	7	12	13	4	8	10	11	9	16	3	6	0
15												
	7											
7	15											

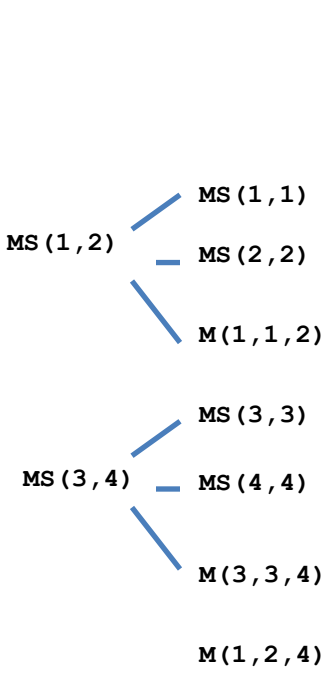

1 2 3 4 5 6 7 8 9 10 11 12 13
15 7 12 13 4 8 10 11 9 16 3 6 0
15
7
7 15
12
MS(1,2)
MS(1,1)
MS(2,2)
M(1,1,2)
MS(3,3)


1 2 3 4 5 6 7 8 9 10 11 12 13
15 7 12 13 4 8 10 11 9 16 3 6 0
15
7
7 15
12
13
MS(1,2)
MS(1,1)
MS(2,2)
M(1,1,2)
MS(3,3)
MS(4,4)

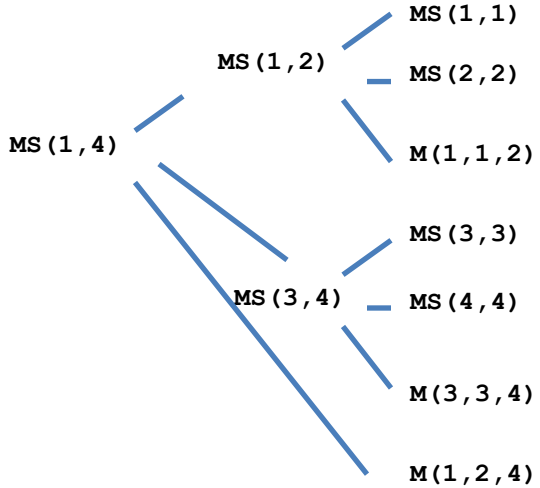

1 2 3 4 5 6 7 8 9 10 11 12 13
15 7 12 13 4 8 10 11 9 16 3 6 0
MS(1,2)
MS(1,1)
MS(2,2)
M(1,1,2)
MS(3,3)
MS(4,4)
M(3,3,4)
15
7
7 15
12
13
12 13



1	2	3	4	5	6	7	8	9	10	11	12	13
15	7	12	13	4	8	10	11	9	16	3	6	0
15												
	7											
7	15											
		12										
			13									
		12	13									



1	2	3	4	5	6	7	8	9	10	11	12	13
15	7	12	13	4	8	10	11	9	16	3	6	0
15												
	7											
7	15											
		12										
			13									
		12	13									
7	12	13	15									



1	2	3	4	5	6	7	8	9	10	11	12	13
15	7	12	13	4	8	10	11	9	16	3	6	0
15												
	7											
7	15											
		12										
			13									
		12	13									
7	12	13	15									

MergeSort

```
Algorithm MergeSort(low,high)
//a[low:high] is a global array to be sorted
//b[low:high] is a temporary array used for merging
{
    if(low<high)then
    {
        mid:=floor(low+high)/2);
        MergeSort(low,mid);
        MergeSort(mid+1,high);
        Merge(low,mid,high);
    }
}
```

MergeSort

- $T(n) = 2T\left(\frac{n}{2}\right) + \theta(n)$
- $a = 2, b = 2, k = 1$
- $\log_b a = \text{k or a} = b^k$
- $T(n) = \theta(nlogn)$

MergeSort

- $T(n) = 2T\left(\frac{n}{2}\right) + n$
- $T(n) = 2(2T\left(\frac{n}{4}\right) + n/2) + n$
- $= 2^2T\left(\frac{n}{2^2}\right) + 2n$
- $= 2^2(2T\left(\frac{n}{2^3}\right) + \frac{n}{2^2}) + 2n$
- $= 2^3T\left(\frac{n}{2^3}\right) + 3n$
- …
- $= 2^kT\left(\frac{\boldsymbol{n}}{\boldsymbol{2^k}}\right) + kn$ $\qquad \frac{\boldsymbol{n}}{\boldsymbol{2^k}} = 1. k = log n$
- $= nT(1) + n\ \log n$
- $= n + n\log n$
- $= \theta(n\log n)$