Group 4 Multi Armed Bandits

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 The dilemma of the Gambler when faced with the Multi
 Armed Bandit
- GETTING FORMAL
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 We take a stab at coming up with a (hopefully) optimal solution for the MAB.

- SIMULATIONS

 A practical Implementation of some of the strategies we discussed.
- VARIANTS OF MAB
 Imaginative twists to the original MAB!
- REAL WORLD APPLICATIONS

 MAB is cool and all...but
 where can it actually be
 applied?



Explore-Exploit Dilemma



- You're in Vegas, faced with a row of slot machines
- No idea which machines readily give rewards and which are real stingy.
- After a few pulls, you have a rough guess which machine is looking profitable and which isn't.
- Do you continue to exploit your currently best machine, or explore other machines: maybe they're gold mines waiting to be discovered.





Let's kick in some math. Shall we?



Problem Formulation

- Row of 'k' slot machines(strategies) 1, 2, ..., k
- Reward for playing strategy s is a r.v. θ_s with distribution π_s , mean μ_s
- At each iteration 't'
 - slot s(t) is played according to algorithm ALG
 - Expected reward for this step = $\mathbb{E}_{\theta_{s(t)}}[\theta_{s(t)}] = \mu_{s(t)}$
- Best strategy, $s^* = \operatorname{argmax}_{i \in [k]} \mu_i$
- Regret R(n, ALG) over n iterations

$$R(n, ALG) = B - \sum_{1 \le t \le n} \mathbb{E}_{\theta_{S(t)}} [\theta_{S(t)}]$$

where B is a regret benchmark

$$=\sum_{1\leq t\leq n}\mu_{S^*}-\mu_{S(t)}$$

$$R(n, ALG) = \sum_{1 \le t \le n} \mu_{S^*} - \mu_{S(t)} = \sum_{s \in [n] \setminus S^*} (\mu_{S^*} - \mu_s) \mathbb{E}_{\theta}[T_s(n)]$$

 $T_s(n)$: Number of trials of machine s in a sequence of n trials

From Lai and Robbins [1], we have

$$\lim_{n \to \infty} \frac{R(n, ALG)}{\ln n} \ge \sum_{s \in [n] \setminus s^*} \frac{(\mu_{s^*} - \mu_s)}{D_{KL}(\pi_s, \pi_{s^*})}$$

$$\Rightarrow R(n, ALG) = O(\ln n)$$

• Visualization:

$$\mathbb{E}_{\theta}[T_{s}(n)] \geq \frac{\ln n}{D_{KL}(\pi_{s}, \pi_{s^{*}})}$$

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$$\mathbb{E}_{\theta}[T_s(n)] \ge \frac{\ln n}{D_{KL}(\pi_s, \pi_{s^*})}$$



Biased coin:

$$\mathbb{P}(H) = 0.5 - \epsilon$$

Unbiased coin:

$$\mathbb{P}(H) = 0.5$$

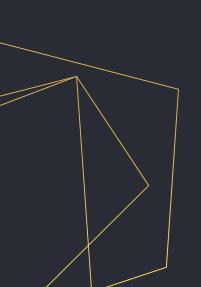
T the number of tosses required to determine with $l(=say\ 0.99)$ surety

$$T \ge O(\frac{1}{\epsilon^2})$$



GAMBLING STRATEGIES

How to get rich playing a multi armed bandit.



1. Epsilon - Greedy Strategy

- We keep track of the average reward we've earned from each arm, as the trials progress.
- 1. Choose the current "best arm" with probability 1- ϵ and the others each with a probability ϵ/n -1
- 1. Update the average reward of the arm chosen based on the outcome

- 1. $P(a(t) \neq a^*) = \varepsilon \text{ as } t \rightarrow \infty$
- Expected Cumulative Regret is of order O(n)
- This algorithm basically says explore with probability ε and exploit with probability
 1-ε

2. Epsilon-decreasing Strategy

- Similar to Epsilon-greedy strategy but the probability of choosing the current best arm is a decreasing function of t
- 1. One such function for epsilon is $\varepsilon(t) = \varepsilon_0/t$
- The exploration phase vanishes after sufficient number of trials

- 1. $P(a(t) \neq a^*) = \varepsilon_0/t \text{ as } t \to \infty$
- The expected cumulative regret is of logarithmic order i.e O(logn)
- 1. If we take $\mathcal{E}(t) = \mathcal{E}_0/t^2$, then the exploration phase vanishes much faster. Too fast for the algorithm to have made a reasonable guess for which is the best arm.

Hoeffding's inequality

$$P(E[\overline{X}] - \overline{X} \ge t) \le e^{-2nt^2}$$
 where n is the number of samples

Principle of Optimism in face of Uncertainty

Certain classes of algorithms for multi armed bandits are based on the principle of optimism in the face of uncertainty. UCB1 is one of the most important algorithm based on it. Suppose you have a black box/oracle that gives you an upper bound on the expected reward for all arms and it gets better each time. This bound reaches the true value after sufficiently large number of trials.

3. Upper Confidence Bound (UCB)

- There is no oracle that can give the upper bound to the expected reward with certainty
- There are probabilistic bounds, that can be derived from Hoeffding's inequality
- 3. It can be seen that $\Pr[R(a) R'(a) > \sqrt{\frac{\log(\frac{1}{\delta})}{2n}}] \le \delta$
- 4. Take δ = 1/n⁴. The bound becomes 2logn/n_a

- 1. For each arm a, calculate $u_a = R'(a) + \sqrt{\frac{2 \log n}{n_a}}$
- 2. Choose the arm with the largest u_a. Let this arm be
- 3. Update n_a= n_a+1 for the arm chosen

4. Update R'(a) =
$$\frac{n_a.R'(a) + reward}{n_a + 1}$$

for the arm chosen

5. Go to step 1

4. Thompson Sampling

- Bayesian approach unlike the previous ones
- Doesn't assume a fixed mean
- The means of each arm are assumed to have some distribution and the distributions get updated based on the rewards
- The arm is chosen based on the value obtained by sampling the distributions
- Usually each arm's mean is assumed to have a beta distribution with parameters x,y different for each arm

- 1. Initialise x = 1, y = 1 for all arms
- 1. Now take k samples from each distribution and calculate the average for each of the arms
- 1. Choose the arm with the highest average and get the reward r
- l. Update : x = x+r , y = y + 1-r for the chosen arm
- 1. Go to step 2

Does β distribution update to another β -distribution??

Yes!

$$f(p)\alpha p^{x-1}(1-p)^{y-1}$$

$$f(p/r=1)\alpha f(r=1/p).f(p) = p.p^{x-1}(1-p)^{y-1} = p^{x}(1-p)^{y-1}$$

$$f(p/r=0)\alpha f(r=0/p).f(p) = (1-p).p^{x-1}(1-p)^{y-1} = p^{x-1}(1-p)^{y}$$





Let's have a computer back up the bold claims we made "in theory".



SOLUTIONS WE'LL SIMULATE

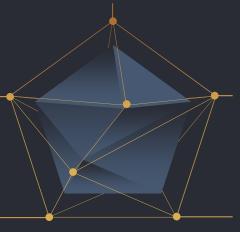
EPSILON DECREASE - 2

The value of epsilon decreases as 1/n²

EPSILON DECREASE - 1

The value of epsilon decreases as 1/n

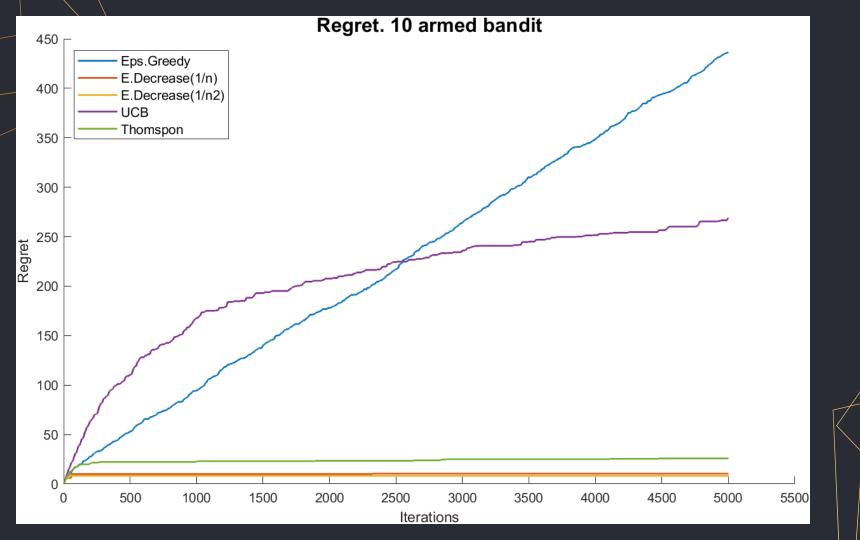
EPSILON GREEDY

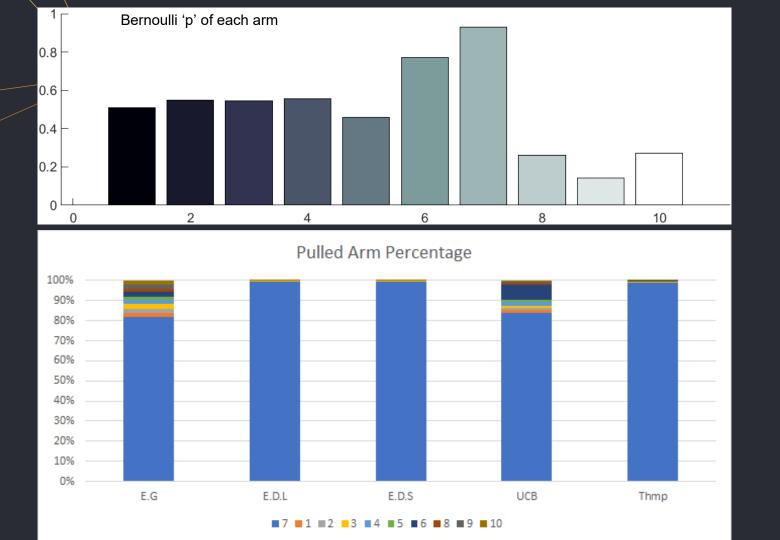


UPPER CONFIDENCE BOUND

THOMPSON SAMPLING

Models the reward of an arm as a Beta distribution





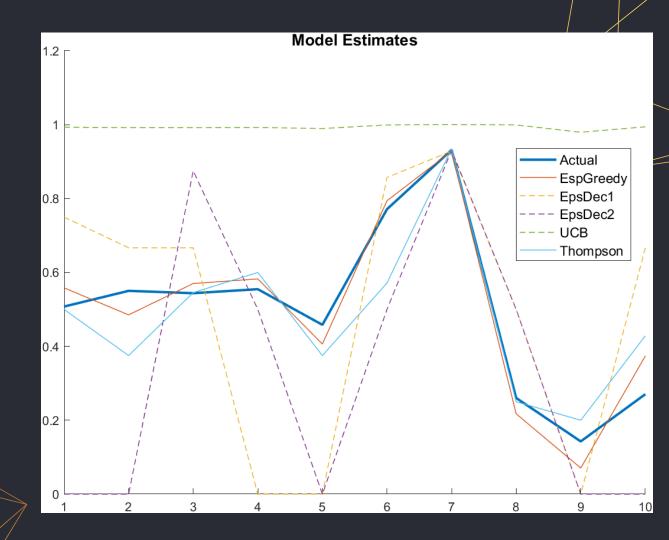
NUMBER OF

ESTIMATING ARM PROBABILITIES

Eps Greedy and Thompson have the best estimates.

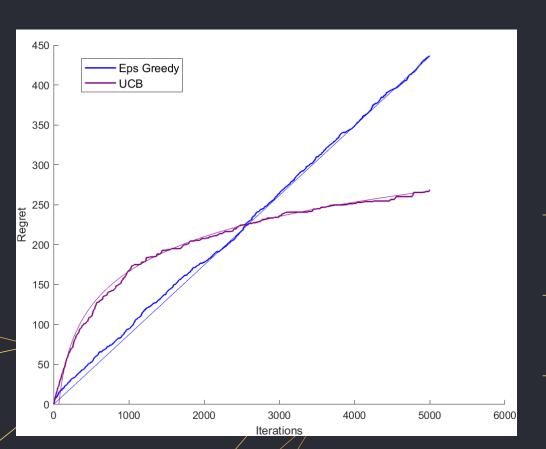
UCB is an upper bound anyway, it's not supposed to model it well.

Though Eps Decay strategies model the distribution as a whole poorly, every model has correctly estimated the probability of the best arm.



THOMPSON SAMPLING

THEORETIC COMPARISON



Very satisfyingly, the results agree almost too well with the theory.

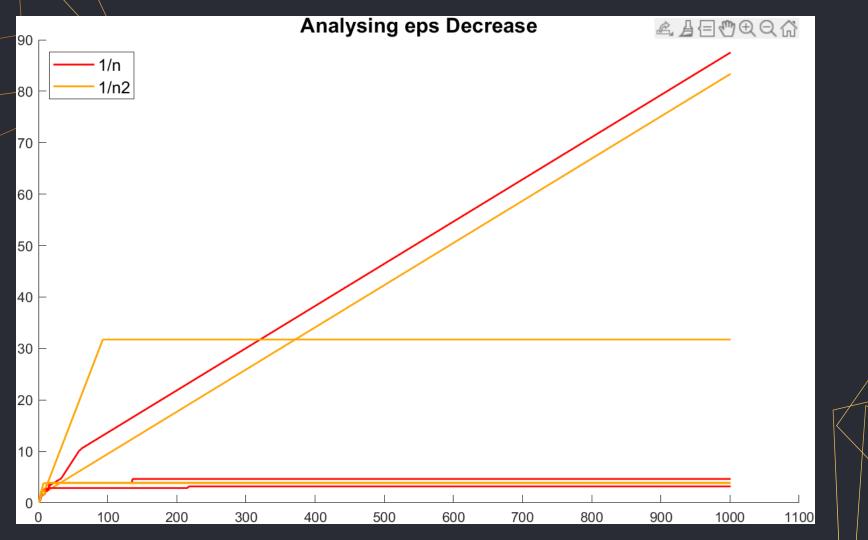
	THEORETIC	MODELLED
EPS GREEDY	O(n)	0.087*n
UCB	O(log.n)	62.6*log(n) - 266

WHAT'S THE DEAL WITH THE EPSILON DECREASE STRATEGIES?

Why on earth do simple epsilon strategies perform so well? These eps. **Decrease** strategies stop exploring pretty quick as epsilon decays to zero. After a while, all they do is exploit. And they perform great. **If** they've managed to find the right arm, that is.

The problem is, the faster you decay, the less time you have to find the right arm. We shall see this in full force in the next slide.



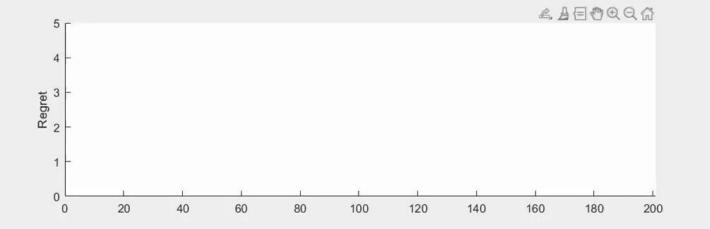


NON-STATIONARY BANDITS

As we have seen, the epsilon decrease strategies have been absolutely destroying our MAB (on their good days). Let's make things a little more challenging.

A non-stationary bandit is one that randomly changes the Bernoulli random variable 'p' for each arm. For simplicity, this random change will happen after every 200 iterations.

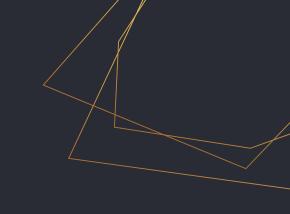
When the eps decrease strategies cannot now "cling onto their best arm", let us see what becomes of them. And of the rest.

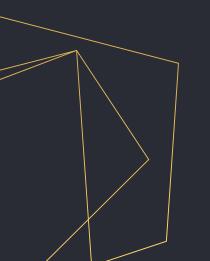












VARIANTS OF M.A.B.

As if the original wasn't bad enough.

ADVERSARIAL BANDIT

• An adversary can control the reward function

COMBINATORIAL BANDIT

- Non linear reward function
- Example: No reward when you win three consecutive turns

DUELLING BANDIT

- Pitching two strategies against each other.
- Play two slot machines, you only know if one is performing better than the other.







ETHICAL CLINICAL TRIALS

- In a treatment of disease
- Evaluation of k treatments on n patients.
- Reward cures and penalize deaths.

INTERNET AD PLACEMENT

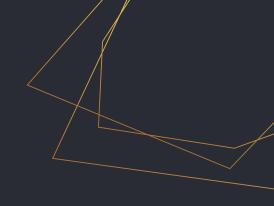
- k Advertisers
- *n* Users visiting the website
- Reward if they click it.

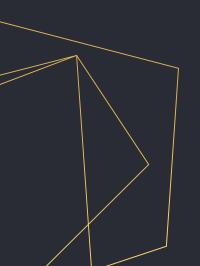
SERVER SELECTION IN NETWORKS

- Client looks for k servers
- For *n* instances try different servers
- Cost is the time delay

PRICING IDENTICAL GOODS FOR SALE

- *k* Pricing strategies
- *n* Markets
- Reward is the profit made.





Thank you!