a) 80 people decide to wait 20 people decided not to wait

$$K = 80 + 20 = 100$$

Initial Entropy at Node 1 is:

$$H\left(\frac{80}{100}, \frac{20}{100}\right) = \frac{-80}{100} \log_2 \frac{80}{100} - \frac{20}{100} \log_2 \frac{20}{100}$$
$$= 0.2575 + 0.46438$$
$$= 10.72188$$

Node A

7est:

weekend ?

yes Node B wait = 20

did not wait = 15

Node C wait = 60

No

did not wait = 5

$$H\left(\frac{20}{35}, \frac{15}{35}\right)$$
 where $k = 20 + 15$

$$= \frac{-20}{35} \log_2 \frac{20}{35} - \frac{15}{35} \log_2 \frac{15}{35}$$

$$= 0.4613 + 0.5238 = \boxed{0.9852}$$

$$H\left(\frac{60}{65}, \frac{5}{65}\right) = -\frac{60}{65} \log_2 \frac{60}{65} - \frac{5}{65} \log_2 \frac{5}{65}$$

$$= -\frac{12}{13} \log \frac{12}{13} - \frac{1}{13} \log_{13} \frac{1}{13}$$

* Information Gain

$$I(E,L) = H(E) - \frac{L}{2} \frac{Ki}{k} H(Ei)$$

$$0.72188 - \frac{35}{100} (0.9852) - \frac{65}{100} (0.3912)$$

As far as from Node A, it is a weekend

At E, since it uses the exact test
whether its a weekend or not, all all ibutes

will be weekend and Mode I will have
no attributes.

Since Tuesday is not a weekend and it is rainy, by looking at the graph, the leaf node will and up in 1 Node F. Output is that he will wait.

According to the question, the leaf node will end up in Node H. The butput for that case is that he will not want.

for 1 Y=3 Y=2 $Root = \frac{-5}{10} log, \frac{5}{10} - \frac{5}{10} log_2 \stackrel{>}{=} \frac{1}{10}$ = 0.5 + 0.5 = 1 $\frac{1011}{H(\frac{3}{3}\frac{9}{3})} = \frac{-3}{3}\log_{2}\frac{3}{3} - \frac{9}{3}\log_{2}\frac{9}{3} = \frac{10}{3}$ for= H(-1,3) = -1 192 - 3 192 - 4 = $(0.25 \times 2) + (0.75 \times 0.415) = [0.81125]$ for 3 $H(\frac{1}{3},\frac{1}{3}) = -\frac{1}{3} \cdot \frac{10}{3} = \frac{10.533}{3} + 0.3851$ = 7/0.9183* Checking for root as B foot = -5 log 5 - 5 log 5 W

$$\frac{for2}{H(\frac{1}{4},\frac{3}{4})} = \frac{-1}{4} \log_2 \frac{1}{4} - \frac{3}{4} \log_2 \frac{3}{4}$$

$$\frac{for2}{H(\frac{3}{4},\frac{1}{4})} = \frac{-3}{4} \log_2 \frac{3}{4} - \frac{1}{4} \log_2 \frac{1}{4}$$

$$= \frac{10 \cdot 81125}{10 \cdot 81125}$$

$$\frac{for3}{H(\frac{1}{2},\frac{1}{2})} = \frac{-1}{2} \log_2 \frac{1}{2} - \frac{1}{2} \log_2 \frac{1}{2}$$

$$+ \frac{(hocking for root as C)}{Root = 1}$$

$$\frac{for1}{H(\frac{1}{5},\frac{4}{5})} = \frac{-3}{5} \log_2 \frac{1}{5} - \frac{1}{5} \log_2 \frac{4}{5}$$

$$= \frac{10 \cdot 7218}{10 \cdot 81125}$$

$$\frac{for2}{H(\frac{3}{4},\frac{1}{4})} = \frac{-3}{4} \log_2 \frac{3}{4} - \frac{1}{4} \log_2 \frac{1}{4}$$

$$= \frac{10 \cdot 81125}{10 \cdot 81125}$$

$$\frac{for3}{H(\frac{1}{4},\frac{1}{4})} = \frac{-3}{4} \log_2 \frac{3}{4} - \frac{1}{4} \log_2 \frac{1}{4}$$

$$= \frac{10 \cdot 81125}{10 \cdot 81125}$$

$$1 - 0 - \frac{4}{10} \left(0.81125 \right) - \frac{3}{10} \left(0.9183 \right)$$

$$= 1 - 0.3245 - 0.27549 = 0.400 L$$

$$\frac{1 - \frac{4}{10} \left(0.81125\right) - \frac{4}{10} \left(0.81125\right) - \frac{2}{10} \left(1\right)}{10 \left(0.81125\right) - \frac{2}{10} \left(1\right)}$$

$$= 1 - 0.3245 - 0.3245 - 0.2 = 0.151$$

$$1 - \frac{5}{10} \left(0.7218 \right) - \frac{4}{10} \left(0.81125 \right) - 0$$

$$= 1 - 0.3609 - 0.3245$$

$$= 10.3146$$

Therefore, A has the most information gain le A is the root.

$$A = 1000$$
, $B = 0$, $C = 0$, $D = 0$

$$frightst Entropy
A = 250, B = 250, C = 250, D = 250$$

$$= > 4 - 250 log_2 250 - 250 log_2 250 l$$

$$I \cdot G = \left[H \left(\frac{250}{1000}, \frac{250}{1000}, \frac{250}{1000} \right) - \left(\frac{250}{1000} \right) - \left(\frac{250}{1000} \right) \right]$$

$$\frac{250}{250}$$
, $\frac{0}{250}$, $\frac{0}{250}$, $\frac{0}{250}$] $\frac{250}{1000}$

* Lowest Information Gain

$$I4 = \left[H \left(\frac{1000}{1000} , \frac{0}{1000} , \frac{0}{1000} \right) - \frac{2}{12} H(\mathcal{E}_i) \right]$$

$$= \boxed{0}$$

TASK4

We will be able to achieve the task by seversing all the outputs. It we henerse, we get an accuracy of 727.

Therefore, un can guarantee achieving Lutter than 60% accusacy.

TASK5

In Praision Tree concept, largest possible rumber of elements for X is 22.

In our case, we have 5 boolean variables, thus, the largest possible purcher of elements for X will be:

$$2^{2^5} = 2^{32}$$