# UCS 2411 DATABASE LAB

# PHARMACY SUPPLY MANAGEMENT SYSTEM

SUBMITTED BY SRI HARI VK (3122225001138) R SKANTA SAMVARTAN (3122225001137) SHRI VISHAAL (3122225001131)

**Department of Computer Science and Engineering** 

Sri Sivasubramaniya Nadar College of Engineering (An Autonomous Institution, Affiliated to Anna University)

Kalavakkam – 603110

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### **Insights of the Project:**

The project is pharmaceutical management system that tracks drugs from suppliers to pharmacies and customers. The system includes inventory management, order processing, supplier and manufacturer details, and customer information.

### **Tables and Their Attributes**

### 1. Supplier

### -Attributes:

- `supplier\_id`: Unique identifier for the supplier.
- `supplier\_name`: Name of the supplier.
- `phone no`: Contact number of the supplier.
- 'address': Physical address of the supplier.
- 'email': Email address of the supplier.

### -Relationships:

- Supplies drugs to the 'Pharmacy'.

### 2. Pharmacy

#### -Attributes:

- `pharmacy id`: Unique identifier for the pharmacy.
- `pharmacy\_name`: Name of the pharmacy.
- `pharmacy type`: Type of pharmacy (e.g., retail, hospital).
- `license no`: License number of the pharmacy.
- `open hours`: Operating hours of the pharmacy.
- `phone`: Contact number of the pharmacy.
- `address`: Physical address of the pharmacy.
- `administration\_status`: Status of the administration (possibly regulatory compliance status).

### -Relationships:

- Receives supplies from 'Supplier'.

#### 3. Inventory

#### - Attributes:

- `inventory\_id`: Unique identifier for the inventory entry.
- 'location': Storage location within the inventory.
- `batch number`: Batch number of the stored drugs.
- 'invoice date': Date of the invoice for the inventory entry.
- `expiry date`: Expiry date of the stored drugs.
- `quantity`: Quantity of the drugs in the inventory.
- `status`: Current status of the inventory entry.

#### -Relationships:

- Gets drugs from `Supplier`.
- Stores `Drug`.



### 4. Drug

#### -Attributes:

- `drug\_id`: Unique identifier for the drug.
- `drug\_name`: Name of the drug.
- `drug code`: Code associated with the drug.
- `storage\_conditions`: Conditions required for storing the drug.
- `drug\_category`: Category or type of the drug.
- `composition`: Chemical composition of the drug.
- `dosage`: Dosage information for the drug.
- `unit\_price`: Price per unit of the drug.

### -Relationships:

- Stored in `Inventory`.
- Manufactured by `Manufacturer`.

### 5. Manufacturer

### -Attributes:

- `manufacturer id`: Unique identifier for the manufacturer.
- `manufacturer name`: Name of the manufacturer.
- `phone no`: Contact number of the manufacturer.
- `address`: Physical address of the manufacturer.
- `email`: Email address of the manufacturer.

### -Relationships:

- Manufactures `Drug`.

#### 6. Customer

#### -Attributes:

- `cust id`: Unique identifier for the customer.
- `cust name`: Name of the customer.
- `phone`: Contact number of the customer.
- `email`: Email address of the customer.
- `gender`: Gender of the customer.
- `dob`: Date of birth of the customer.
- `address`: Physical address of the customer.

#### -Relationships:

- Places 'Order'.

### 7. Order

#### -Attributes:

- `order id`: Unique identifier for the order.
- `order date`: Date the order was placed.
- `quantity`: Quantity of drugs ordered.
- `tot amt`: Total amount of the order.
- `delivery\_date`: Date the order was delivered.

#### -Relationships:

- Contains `Drug`.
- Placed by `Customer`.

### 8. Payment



#### - Attributes:

- `pay id`: Unique identifier for the payment.
- `pay amt`: Amount paid.
- `pay date`: Date of the payment.

### - Relationships:

- Associated with 'Customer'.

### **Relationships:**

- Supplier to Inventory: Suppliers provide drugs to be stored in the inventory.
- Inventory to Drug: Inventory entries store specific drugs.
- Drug to Manufacturer: Drugs are produced by manufacturers.
- Customer to Order: Customers place orders for drugs.
- Order to Drug: Orders include specific drugs.
- Customer to Payment: Payments are made by customers.

# **ASSUMPTIONS:**

### **Customer Registration:**

New customers are added to the CUSTOMER table with unique USER\_NAME.

Passwords are securely stored (assumed hashed and not plain text).

Placing Orders:

When a customer places an order, a record is created in the ORDERS table with a foreign key reference to USER\_NAME.

Each order contains multiple entries in the ORDER\_LIST table, detailing the DRUG\_ID and quantity ordered.

Order History:

The system can query the ORDERS table to retrieve all orders placed by a specific customer using USER NAME.

The ORDER\_LIST table can be queried to find out all drugs ordered within a particular order.

#### Manufacturer and Drug:

Each drug (DRUG\_ID) is produced by a single manufacturer (MANUFACTURER\_ID). The relationship ensures traceability and accountability for the drug's production. Drug and Inventory:

Each drug can have multiple inventory records, indicating different batches or stocks. This allows tracking of stock levels, expiry dates, and inventory locations.

Drug and Pharmacy:

The single pharmacy maintains a record for each drug, indicating the supplied quantity. This relationship simplifies the tracking of drug availability in the pharmacy. Drug and Orders:

Each drug can be ordered multiple times by different customers.

The ORDER\_LIST table captures these details, linking each drug to its respective orders and quantities.



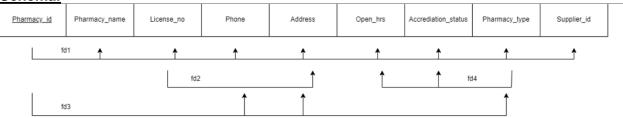
# **IDENTIFICATION OF FD's:**

# 1)Pharmacy:

```
 \begin{array}{l} (\text{A-Pharmacy\_id,B} \rightarrow \text{Pharmacy\_name,C} \rightarrow \text{License\_no,D} \rightarrow \text{Phone,E} \rightarrow \text{address,F} \rightarrow \text{open\_hrs,} \\ \text{G} \rightarrow \text{accreditation\_status,H} \rightarrow \text{Pharmacy\_type,I} \rightarrow \text{Supplier\_id}) \\ \text{Pharmacy\_id} \rightarrow \text{Pharmacy\_name,license\_no,Phone,address,open\_hrs,accreditation\_status,} \\ \text{Pharmacy\_type.} \\ \text{License\_no} \rightarrow \text{address} \\ \text{Pharmacy\_type} \rightarrow \text{accrediation\_status} \\ \text{Pharmacy\_type} \rightarrow \text{accrediation\_status} \\ \text{Pharmacy\_type} \rightarrow \text{open\_hrs} \\ \text{Pharmacy\_id,Pharmacy\_name} \rightarrow \text{phone,Pharmarcy\_type,address} \\ \end{array}
```

- A → B,C,D,E,F,G,H,I
- $\bullet$  C  $\rightarrow$  E
- $H \rightarrow F,G$
- $\bullet$  A,B  $\rightarrow$  D,H,E

### Schema:



### **MINIMAL FDs:**

a)
With A → B
A+={A,B,C,D,H,E,F,G,I}
Without A → B
A+={A}
Therefore A→B is not redundant
b)
With A→C

With  $A\rightarrow C$   $A+=\{A,C,E\}$  Without  $A\rightarrow C$   $A+=\{A\}$ Therefore  $A\rightarrow C$  is not redundant



```
d)
        With A→E
        A+=\{A,B,C,D,E,F,G,H,I\}
        Without A→E
        A \leftarrow \{A,B,C,D,E,F,G,H,I\}
        Therefore the fd is not redundant
  e)
        With C \rightarrow E
        C=\{C,E\}
        Without C \rightarrow E
        C += \{C\}
        Therefore C \rightarrow E is not redundant
 f)
        With H \rightarrow F,G
        H+=\{H,F,G\}
        Without H \rightarrow F,G
        H+=\{H\}
        Therefore H \rightarrow F,G is not redundant
 g)
        With A,B \rightarrow D,H,E
        AB+=\{A,B,D,H,E,C,F,G,I\}
        Without A,B \rightarrow D,H,E
        AB+=\{A,B,C,D,E,F,G,H,I\}
        Therefore A,B → D,H,E is redundant
• The minimal FDs are
        {
                 A \rightarrow B,C,D,H,I
                 C \rightarrow E
                 H \rightarrow F,G
        }
```

- The candidate key is {A}
- PA={A}
- NPA={B,C,D,E,F,G,H,I}

# 2)Supplier:

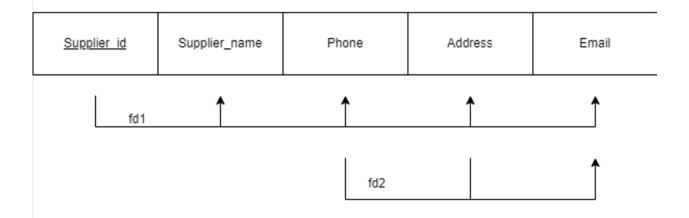
 $\mathsf{A} \to \mathsf{Supplier\_id}, \mathsf{B} \to \mathsf{Supplier\_name}, \mathsf{C} \to \mathsf{phone}, \mathsf{D} \to \mathsf{address}, \mathsf{E} \to \mathsf{email}$ 

# FDs:

- $\bullet$  A  $\rightarrow$  B,C,D,E
- $\bullet$  C,E  $\rightarrow$  D

#### Schema:





### **Minimal FDs:**

i)Decomposition:

- $\bullet \ A \to B$
- $\bullet \ \mathsf{A} \to \mathsf{C}$
- $\bullet \ A \to D$
- $\bullet \ \mathsf{A} \to \mathsf{E}$

$$I)A \rightarrow B$$

With  $A \rightarrow B$ 

 $A \leftarrow \{A,B,C,D,E\}$ 

Without  $A \rightarrow B$ 

 $A \leftarrow \{A,C,D,E\}$ 

Therefore  $A \rightarrow B$  Is not redundant

### ii) $A \rightarrow C$

With  $A \rightarrow C$ 

 $A \leftarrow \{A,B,C,D,E\}$ 

Without  $A \rightarrow C$ 

 $A \leftarrow \{A,B,D,E\}$ 

Therefore  $A \rightarrow C$  is not redundant

# iii) $A \rightarrow D$

With  $A \rightarrow D$ 

 $A \leftarrow \{A,B,C,D,E\}$ 

Without  $A \rightarrow D$ 

 $A \leftarrow \{A,B,D,C,E\}$ 

Therefore  $A \rightarrow D$  is redundant

# $IV)A \rightarrow E$

With  $A \rightarrow E$ 

 $A \leftarrow \{A,B,C,D,E\}$ 

Without  $A \rightarrow E$ 

 $A \leftarrow \{A,B,C,D\}$ 

Therefore  $A \rightarrow E$  is not redundant

2)



```
C,E \rightarrow D

With C,E \rightarrow D

CE+=\{C,E,D\}

Without C,E \rightarrow D

CE+=\{C,E\}

Therefore C,E \rightarrow D is not redundant

MINIMAL FDs:

A \rightarrow B,C,E

C,E \rightarrow D
```

The candidate key for relation is {A} PA={A} NPA={B,C,D,E}

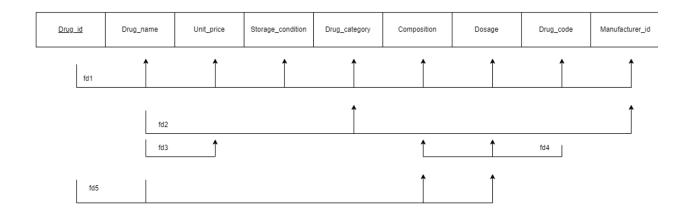
# <u>3)DRUG:</u>

 $A\text{-}Drug\_id, B \to Drug\_name, C \to Unit\_price, D \to Storage\_condition, E \to Drug\_category, F \to Composition, G \to Dosage, H \to Drug\_code, I \to manufacturer\_id$ 

### FDs:

- $\bullet$  A  $\rightarrow$  B,C,D,E,F,G,H,I
- $\bullet$  B  $\rightarrow$  E,I
- $\bullet$  H  $\rightarrow$  G,F
- $\bullet \ B \to C$
- $\bullet$  A,B  $\rightarrow$  G,F

### Schema:



### **MINIMAL FDs:**

```
i)Decomposition:
```

 $A \rightarrow B$ 

 $A \rightarrow C$ 

 $\mathsf{A}\to\mathsf{D}$ 

 $\mathsf{A}\to\mathsf{E}$ 

 $A \rightarrow F$ 

 $A \rightarrow G$ 

 $A \rightarrow H$ 

 $A \rightarrow I$ 

 $\mathsf{B}\to\mathsf{E}$ 

 $\mathsf{B}\to\mathsf{I}$ 

 $H \rightarrow G$ 

 $\mathsf{H} \to \mathsf{F}$ 

 $\mathsf{B} \to \mathsf{C}$ 

1)

a)

With  $A \rightarrow B$ 

 $A+=\{A,B,C,D,E,F,G,H,I\}$ 

Without  $A \rightarrow B$ 

 $A \leftarrow \{A,C,D,E,F,G,H,I\}$ 

 $A \rightarrow B$  is not redundant

b)

With  $A \rightarrow C$ 

 $A+=\{A,B,C,D,E,F,G,H,I\}$ 

Without  $A \rightarrow C$ 

 $A+=\{A,B,C,D,E,F,G,H,I\}$ 

 $A \rightarrow C$  is redundant

c)

With  $A \rightarrow D$ 

 $A+=\{A,B,C,D,E,F,G,H,I\}$ 

Without  $A \rightarrow D$ 

 $A+=\{A,B,C,E,F,G,H,I\}$ 

 $\mathsf{A} \to \mathsf{D}$  is not redundant

d)

With  $A \rightarrow E$ 

 $A+=\{A,B,C,D,E,F,G,H,I\}$ 

Without  $A \rightarrow E$ 

 $A+=\{A,B,C,D,E,F,G,H,I\}$ 

 $A \rightarrow E$  is redundant

```
e)
          With A \rightarrow F
          A \leftarrow \{A,B,C,D,E,F,G,H,I\}
          Without A \rightarrow F
          A \leftarrow \{A,B,C,D,E,F,G,H,I\}
          A \rightarrow F is redundant
f)
          With A \rightarrow G
          A \leftarrow \{A,B,C,D,E,F,G,H,I\}
          Without A \rightarrow G
          A \leftarrow \{A,B,C,D,E,F,G,H,I\}
          A \rightarrow G is redundant
g)
          With A \rightarrow H
          A+=\{A,B,C,D,E,F,G,H,I\}
          Without A \rightarrow H
          A+=\{A,B,C,D,E,F,G,I\}
          A \rightarrow H is not redundant
g1)
          With A \rightarrow I
          A+=\{A,B,C,D,E,F,G,H,I\}
          Without A \rightarrow I
          A \leftarrow \{A,B,C,D,E,F,G,H,I\}
          A \rightarrow I is redundant
h)
          With B \rightarrow E
          B+=\{B,E,C,I\}
          Without B \rightarrow E
          B+=\{B, C, I\}
          Therefore B \to E is not redundant
i)
          With B \rightarrow C
          B+=\{B,C,E,I\}
          Without B \rightarrow C
          B+=\{B,E,I\}
          Therefore B \rightarrow C is not redundant
i1)
          With B \rightarrow I
          B+=\{B,C,E,I\}
          Without B \rightarrow I
          B+=\{B,C,E\}
          Therefore B → I is not redundant
j)
          With H \rightarrow G
          H+=\{H,G,F\}
          Without H \rightarrow G
          H+=\{H,F\}
          Therefore H \rightarrow G is not redundant
k)
          With H \rightarrow F
          H+=\{H,G,F\}
          Without H \rightarrow F
          H+=\{H,G\}
          Therefore H \rightarrow F is not redundant
```



### MINIMAL FDs:

- $\bullet$  A  $\rightarrow$  B,D,H
- $\bullet~H \to F,G$
- B  $\rightarrow$  C,E,I

# Candidate key of the relation is {A}

### Therefore

- PA={A}
- NPA={B,C,D,E,F,G,H,I}

# 4)Inventory:

A-Inventory\_id

B-Batch no

C-Location

**D-Status** 

E-Qty

F-Expiry\_date

G-Supplier\_id

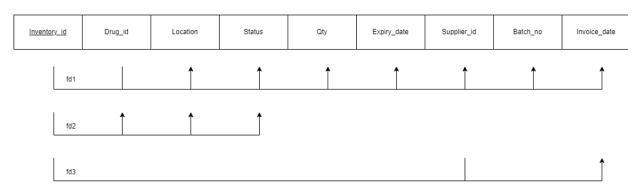
H-Drug\_id

I-Invoice\_date

### FDs:

- A,H → B,C,D,E,F,G,I
- $\bullet \ A \to C, D, B$   $\bullet \ A, G \to I$

#### Schema:



# Decomposing:

 $A,H \rightarrow B$ 

 $A,H \rightarrow C$ 

 $A,H \rightarrow D$ 

 $A,H \rightarrow E$  $A,H \rightarrow F$ 

 $A,H \rightarrow G$ 

 $A,H \rightarrow I$ 

A→C

 $A,G \rightarrow I$ 

```
Minimal FDs:
a)A,H \rightarrow B
        With A, H \rightarrow B
        AH+=\{A,B,C,D,E,F,G,H,I\}
        Without AH \rightarrow B
        AH+=\{A,B,C,D,E,F,G,H,I\}
        Therefore the above FD is redundant
b)
        With AH \rightarrow C
        AH+=\{A,B,C,D,E,F,G,H,I\}
        Without AH \rightarrow C
        AH+=\{A,B,C,D,E,F,G,H,I\}
        Therefore the above FD is redundant
c)
        With AH \rightarrow D
        AD+=\{A,B,C,D,E,F,G,H,I\}
        Without AH \rightarrow D
        AD+=\{A,B,C,D,E,F,G,H,I\}
        Therefore the above FD is redundant
d)
        With AH \rightarrow E
        AH+=\{A,B,C,D,E,F,G,H,I\}
        Without AH \rightarrow E
        AH+=\{A,B,C,D,F,G,H,I\}
        Therefore the above FD is not redundant
e)
        With AH \rightarrow F
        AH+=\{A,B,C,D,E,F,G,H,I\}
        Without AH \rightarrow F
        AH+=\{A,B,C,D,E,G,H,I\}
        Therefore the above fd is not redundant
f)
        With A,H \rightarrow G
        AH+=\{A,B,C,D,E,F,G,H,I\}
        Without AH \rightarrow G
        AH+=\{A,B,C,D,E,F,H,I\}
        Therefore the above fd is not redundant
g)
        With AH \rightarrow I
        AH+=\{A,B,C,D,E,F,G,H,I\}
        Without AH \rightarrow I
        AH+=\{A,B,C,D,E,F,G,H,I\}
        Therefore the above fd is redundant
h)
        With A,G \rightarrow I
        AG+=\{A,G,I\}
        Without A,G \rightarrow I
        AG += \{A,G\}
```



Therefore the above fd is not redundant

# MINIMAL FDs:

- $\bullet$  AH  $\rightarrow$  EFG
- A→B,C,D
- $\bullet \ AG \to I$

The candidate key for the relation is AH. PA={A,H} NPA={B,C,D,E,F,G,I}

# **5)CUSTOMER TABLE:**

### **ATTRIBUTES:**

A: CUST\_ID

B: CUST\_NAME

C: EMAIL

D: ADDRESS

E: PHONE

F: PAY ID

G: PAY\_AMT

H: PAY\_DATE

### **FUNCTIONALS DEPENDENCIES:**

 $A \rightarrow B$  , C , D , E

 $A, B \rightarrow C$ 

 $\mathsf{A} \ , \ \mathsf{B} \to \mathsf{C}$ 

 $A, B \rightarrow B, C$ 

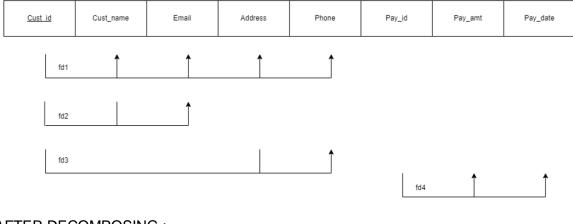
 $A, D \rightarrow E$ 

 $F \rightarrow G$ , H

# Minimal FD's:

REMOVING TRIVIAL FD's : A , B  $\rightarrow$ B, C

### Schema:



```
AFTER DECOMPOSING:
```

```
A \rightarrow B
```

 $A \rightarrow C$ 

 $\mathsf{A}\to\mathsf{D}$ 

 $\mathsf{A}\to\mathsf{E}$ 

$$A, B \rightarrow C$$

A ,  $D \rightarrow E$ 

 $\mathsf{F} \to \mathsf{G}$ 

 $\mathsf{F}\to\mathsf{H}$ 

### CHECKING FOR EXTRANEOUS:

1) A,  $B \rightarrow C$ 

TAKING B AS EXTRANEOUS:

 $A \rightarrow C$ 

 ${A}+ = {A, B, C, D, E}$ 

TAKING A AS EXTRANEOUS:

 $B \rightarrow C$ 

 $\{B\} + = \{B, C\}$ 

SO ,  $A \rightarrow$ ; C IS THE MINIMAL FD.

### 2) A, D $\rightarrow$ E

TAKING A AS EXTRANEOUS:

 $A \rightarrow E$ 

 ${A}+ = {A, B, C, D, E}$ 

TAKING A AS EXTRANEOUS:

 $D \rightarrow E$ 

 $\{D\} + = \{D, E\}$ 

SO,  $A \rightarrow E$  IS THE MINIMAL FD.

# NON - REDUNDANT AND MINIMAL FD's:

 $A \rightarrow B$ 

 $A \rightarrow C$ 

 $\mathsf{A}\to\mathsf{D}$  $A \rightarrow E$ 

 $\mathsf{F}\to\mathsf{G}$ 

 $\mathsf{F}\to\mathsf{H}$ 

FINDING CANDIDATE KEY:

# **6)MANUFACTURERS:**

### **Attributes:**

A: MANUFACTURER\_ID

B: MANUFACTURER\_NAME

C: ADDRESS

D: EMAIL

E: PHONE\_NO

F: INVENTORY\_ID

### **Functional Dependencies:**

Given the attributes, the likely functional dependencies can be inferred as follows:

$$A \rightarrow B, C, D, E, F$$

### Minimal FD's:

 $\mathsf{A}\to\mathsf{B}$ 

 $\mathsf{A} \to \mathsf{C}$ 

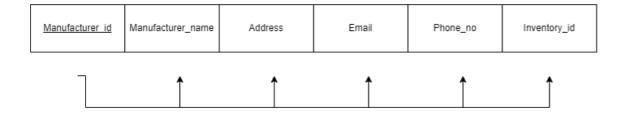
 $\mathsf{A}\to\mathsf{D}$ 

 $\mathsf{A} \to \mathsf{E}$ 

 $\mathsf{A}\to\mathsf{F}$ 

Since the given dependency is already minimal and there are no composite keys, we can proceed to normalization.

### Schema:



### **Candidate Key:**

SSI

{A} (MANUFACTURER\_ID)

# **Prime Attributes:**

{A}

# Non-Prime Attributes:

• {B, C, D, E, F}

# 7)ORDERS:

Let the attributes be:

A: ORDER\_ID

B: ORDER\_DATE

C: DELV\_DATE

D: QUANTITY

E: TOT\_AMT

F: CUST\_ID

G: DRUG\_ID

# **Functional Dependencies:**

AG->DE

AF->BC

A->BCF

# Minimal FD's:

=>Splitting the FD's:

A->B

A->C

A->F



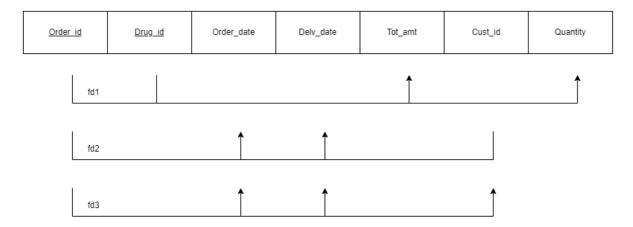
AF->B

AF->C

AG->D

AG->E

### Schema:



# =>Removing Extraneous Attributes:

Extraneous Attributes are:

Since,

A->B , A->C

# =>Checking For Redundancy:

a) A->B

With A->B:

 $A+=\{B,C,D,E,F,G\}$ 

Without A->B:

 $A+=\{C,D,E,F,G\}$ 

A->B IS NOT REDUNDANT.

b) A->C

With A->C:

 $A+=\{B,C,D,E,F,G\}$ 

Without A->C:

 $A+=\{B,D,E,F,G\}$ 

A->C IS NOT REDUNDANT.

c) A->F

With A->F:

 $A+=\{B,C,D,E,F,G\}$ 

Without A->F:

 $A+=\{B,C,D,E,G\}$ 

A->F IS NOT REDUNDANT.

d) AG->D

With AG->D:

 $AG+=\{A,G,D,E\}$ 

Without AG->D:

 $AG+=\{A,G,E\}$ 

AG->D IS NOT REDUNDANT.

e) AG->E

With AG->E:

 $AG+=\{A,G,D,E\}$ 

Without AG->E:

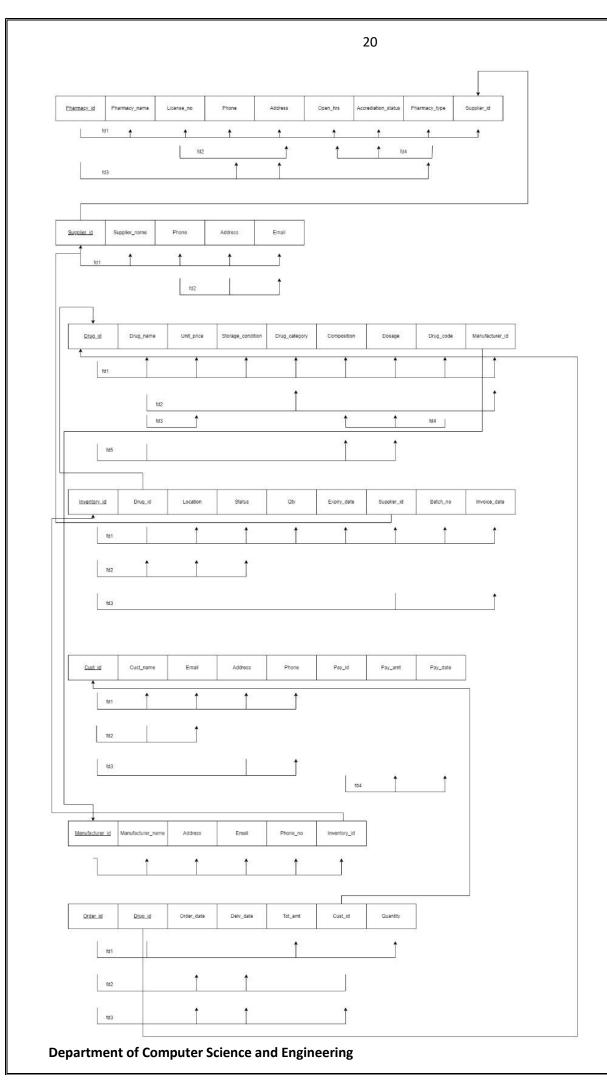
 $AG+=\{A,G,D\}$ 

AG->E IS NOT REDUNDANT.

**CANDIDATE KEY: {A,AG}** 

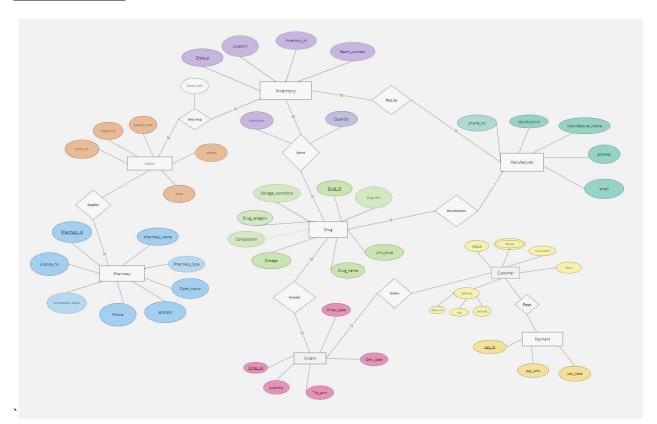
PRIME ATTRIBUTES: {A,G}

**NON-PRIME ATTRIBUTES: {B,C,D,E,F}** 





# **ER DIAGRAM:**





# PHASE 2 NORMALIZATION OF RELATIONS INTO

# **BCNF,3NF**

# 1)Pharmacy:

# **Checking for 1NF:**

• Atomicity is satisfied in this relation relation

Therefore the following relation is in the 1NF

# **Checking for 2NF:**

- No partial dependency should determine a NPA attribute of the relation
- Therefore the given relation satisfies the above condition
- The given relation is in the 2NF

# **Checking for 3NF:**

- No NPA → NPA in the relation
- $\bullet$  C  $\rightarrow$  E
- $H \rightarrow F,G$
- The above FDs violate the 3NF conditions

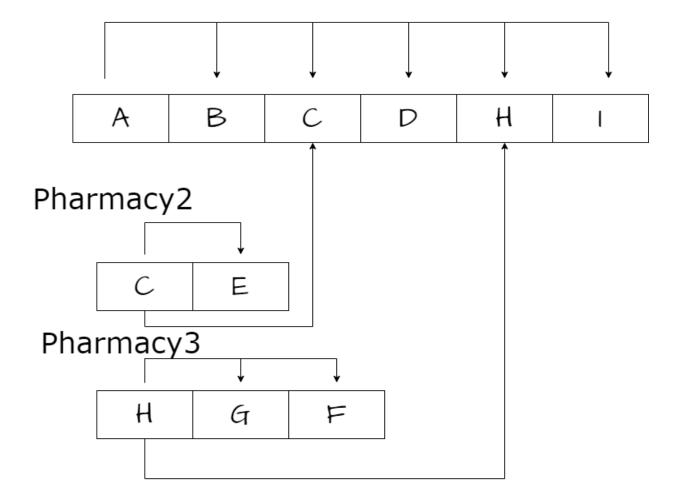
# **Checking for BCNF:**

The condition for BCNF is

- LHS must be a super key
- There is no FDs that violate BCNF



# Pharmacy1



# Pharmacy1:

Candidate key = $\{A\}$ 

FD's= $\{A \rightarrow B,C,D,H,I\}$ 

# Pharmacy2:

Candidate  $key=\{c\}$ 

SSI

$$FD's=\{C \rightarrow E\}$$

# Pharmacy3:

Candidate key={H}

FD's=
$$\{H \rightarrow F,G\}$$

# Pharmacy4:

Candidate  $key = \{A\}$ 

FD's = 
$$\{A \rightarrow C, H\}$$

# 2)Supplier:

# **Checking for 1NF:**

- The atomicity of the relation is satisfied.
- The following FDs are in 1NF relation.

# **Checking for 2NF:**

- No partial dependency should determine a NPA attribute of the relation
- Therefore the given relation satisfies the above condition
- The given relation is in the 2NF

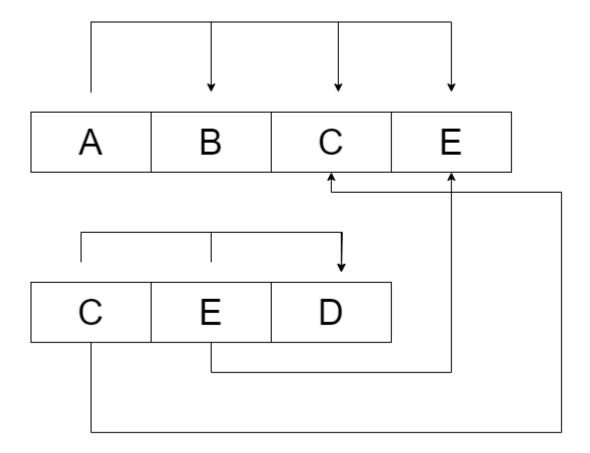
# **Checking for 3NF:**

• No transitive dependency in the relation



• NPA  $\rightarrow$  NPA

The fd which violates the 3NF is C,E  $\rightarrow$  D



# Supplier1:

- Candidate key={A}
- PA={A}
- NPA= $\{B,C,E\}$
- FDs:

 $A \rightarrow B,C,E$ 

# **Supplier2:**

- Candidate Key={C}
- PA={C}
- NPA= $\{E,D\}$
- FDs:

SSI

 $C \rightarrow E,D$ 

# **Checking for BCNF:**

- LHS must be a super key
- The relation satisfies BCNF

# **3)DRUG:**

# **Checking for 1NF:**

- The atomicity of the relation is satisfied.
- The above relation satisfies the condition.

# **Checking for 2NF:**

- No partial dependency should exist in the relation
- The above relation satisfies the condition.

# **Checking for 3NF:**

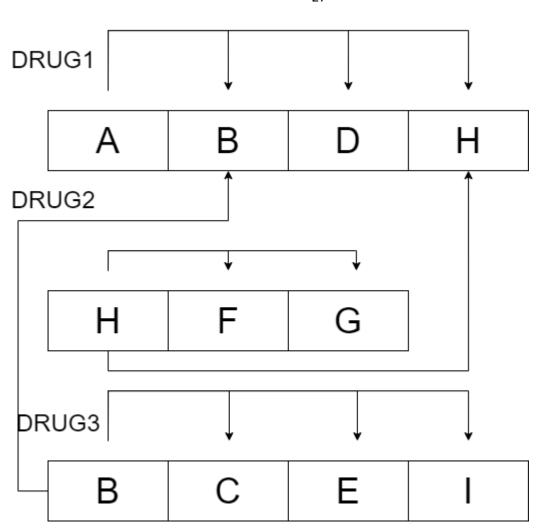
- No transitive dependency should exist in the relation.
- NPA  $\rightarrow$  NPA

 $H \rightarrow F,G$ 

 $B \rightarrow C, E, I$ 

The above two FDs violate the 3NF condition.





The 3NF is satisfied.

# Drug1:

- Candidate key={A}
- PA={A}
- NPA={B,D,H}
- FDs:

$$A \rightarrow B,D,H$$

# Drug2:

- Candidate key={H}
- PA={H}
- NPA= $\{F,G\}$
- FDs:

SSI

# Drug3:

- Candidate key={B}
- PA={B}
- NPA={C,E,I}
- FDs:

$$B \rightarrow C, E, I$$

# **Checking for BCNF:**

- LHS must be a Super-key
- The above relation satisfies the BCNF condition

# 4)Inventory:

# **Checking for 1NF:**

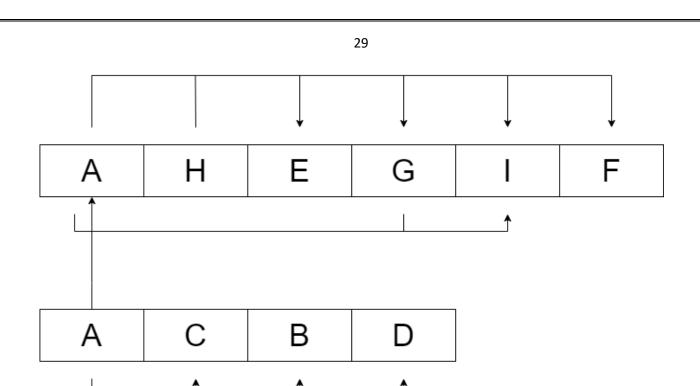
The atomicity of the relation is satisfied

# **Checking for 2NF:**

No partial dependency in the relation

A→BCD Violates the 2NF condition





# **Inventory1:**

- Candidate key={A,H}
- PA={A,H}
- NPA= $\{E,F,G,I\}$
- FDs:

# **Inventory2:**

- Candidate key={A}
- PA={A}
- NPA={CBD}
- FDs:

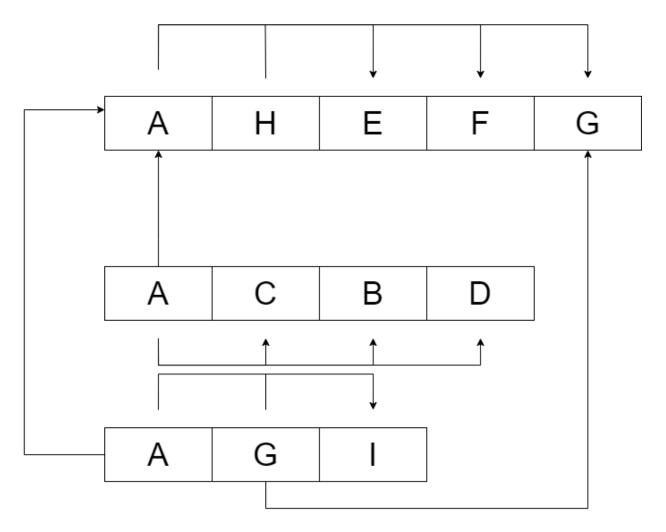
# **Checking for 3NF:**

- PA={AH}
- NPA= $\{E,G,I,F\}$
- Candidate key is AH
- No transitive dependency should exist in the relation
- NPA→NPA



AG→I

violates the condition for 3NF



Therefore 3CNF is satisfied

# **Checking for BCNF:**

- LHS must be a super-key
- The above relation satisfies the condition.



# 5) CUSTOMER TABLE:

### **CHECKING FOR 1NF:**

COMPOSITE AND MULTIVALUED ATTRIBUTE PRESENT

### **DECOMPOSING TO 1 NF:**

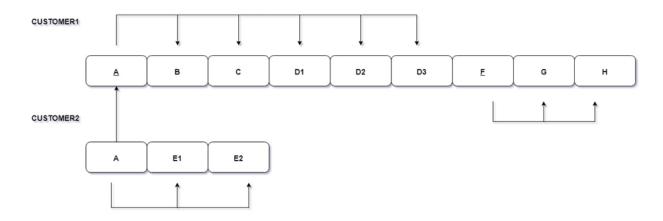
D1: DOOR\_NO

D2: CITY

D3: PINCODE

E1: PHONE\_NUM1

E2: PHONE\_NUM2



### CUSTOMER1 TABLE FDS ARE:

 $A \rightarrow B$ ,  $A \rightarrow C$ ,  $A \rightarrow D1$ ,  $A \rightarrow D2$ ,  $A \rightarrow D3$ ,  $F \rightarrow G$ ,  $F \rightarrow H$ 

CUSTOMER2 TABLE FDS ARE:{A -> E1, E2}

A IS THE CANDIDATE KEY FOR CUSTOMER1 TABLE.

### **CHECKING FOR 2NF:**

PARTIAL DEPENDENCIES ARE PRESENT IN CUSTOMER1 TABLE.

THE FDS WITH PARTIAL DEPENDENCIES

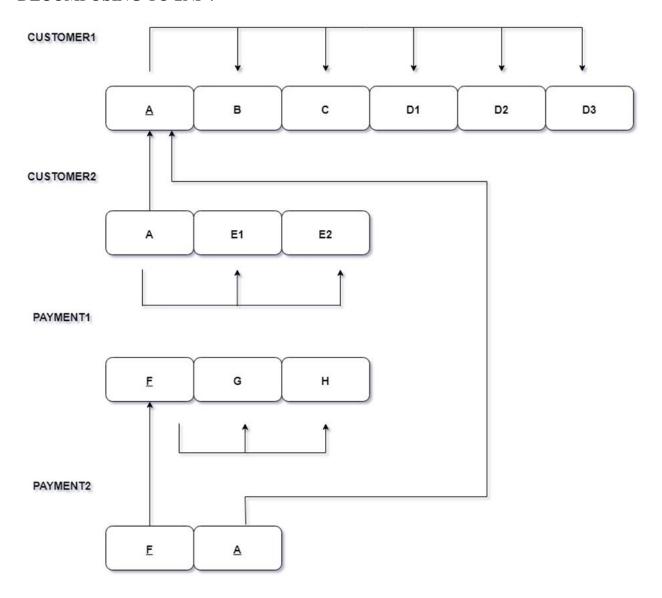
{

{

SSI

```
A \rightarrow B , C , D1 , D2 , D3 F \rightarrow G, H
```

# DECOMPOSING TO 2NF:



CHECKING FOR FURTHER PARTIAL DEPENDENCIES:

IN CUSTOMER1 TABLE THE CANDIDATE KEY IS {A}

THERE IS NO PARTIAL DEPENDENCIES IN CUSTOMER1.

IN PAYMENT1 TABLE THE CANDIDATE KEY ID {F}

THERE IS NO PARTIAL DEPENDENCIES IN PAYMENT1.



### **CHECKING FOR 3NF:**

NO, TRANSITIVE DEPENDENCIES IS PRESENT.

THEREFORE 3NF IS SATISFIED.

### **CHECKING FOR BCNF:**

ALL FDS HAVE SUPER KEY AS THE LEFT ATTRIBUTE.

THEREFORE BCNF IS SATISFIED.

# 6)MANUFACTURERS:

### **1NF (First Normal Form):**

- Condition: The relation with the functional dependency A → B, C, D, E, F is in First Normal Form (1NF) because:
- All attributes contain atomic values.
- Each attribute holds a single value.
- Values in each column are of the same data type.
- Each row is uniquely identified by the primary key MANUFACTURER\_ID, ensuring there are no repeating groups or arrays.

Conclusion: The given relation satisfies all the conditions of 1NF.

### **2NF (Second Normal Form):**

- Condition: A relation is in 2NF if it is in 1NF and all non prim attributes are fully functionally dependent on the entire primary key.
- Explanation:

The relation is in 1NF because all attributes contain only atomic values.

The primary key is A (MANUFACTURER\_ID).

- The functional dependency A → B, C, D, E, F shows that all non prime attributes (B, C, D, E, F) are fully functionally dependent on the primary key A.
- Since A is a single attribute key, there cannot be partial dependencies.



**Conclusion**: The relation is in 2NF because it is in 1NF and all non-prime attributes are fully functionally dependent on the primary key A.

### **3NF (Third Normal Form):**

- Condition: A relation is in 3NF if it is in 2NF and for every functional dependency X
   → Y, either X is a super key or Y is a prime attribute.
- Explanation:

The relation is already in 2NF.

For the dependency  $A \rightarrow B$ , C, D, E, F, the left-hand side A is a super key (since A is the primary key).

There are no transitive dependencies (dependencies where a non-prime attribute depends on another non-prime attribute).

Conclusion: The relation is in 3NF because it is in 2NF and for every functional dependency, the left-hand side is a super key.

# **BCNF** (Boyce-Codd Normal Form):

Condition: A relation is in BCNF if for every non-trivial functional dependency  $X \rightarrow Y$ , X is a super key.

### • Explanation:

The relation is already in 3NF.

The given functional dependency  $A \rightarrow B$ , C, D, E, F shows that the primary key A determines all other attributes.

In this case, A (MANUFACTURER\_ID) is a super key because it uniquely identifies each tuple in the relation.

There are no other functional dependencies where a non-super key determines another attribute.



#### MANUFACTURER

MANUFACTURER_ID	MANUFACURER_NAME	ADDRESS	EMAIL	PHONE_NO	INVENTORY_ID
	<b>↑</b>	<u> </u>	<u> </u>	<u> </u>	<u> </u>

# 7)ORDERS:

# • 1NF(First Normal Form):

The relations with the functional dependency  $A \rightarrow B$ , C, F AND AG->D,E are in First Normal Form (1NF) because:

- All attributes contain atomic values.
- Each attribute holds a single value.
- Values in each column are of the same data type.
- Each row is uniquely identified by the primary key order\_id, (order\_id,drug\_id) ensuring there are no repeating groups or arrays.

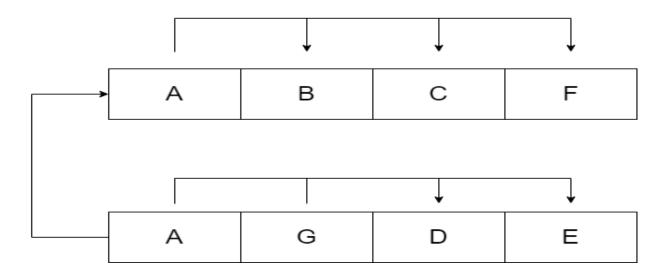
Therefore, the given relation satisfies all the conditions of 1NF.

### • 2NF (Second Normal Form):

- To satisfy Second Normal Form (2NF), the relation must first be in 1NF and also meet the following criteria:
- There should be no partial dependency, which means no non-prime attribute should be functionally dependent on any proper subset of any candidate key of the relation.
- Given:
  - The candidate keys are {A, AG}.
  - Prime attributes are {A, G}.
  - Non-prime attributes are  $\{B, C, D, E, F\}$ .
  - From the given functional dependencies:
    - $\bullet \quad A \to B$
    - $\bullet \quad A \to C$
    - $\bullet \quad A \to F$
- These functional dependencies show that B, C, and F are dependent on A, which is a candidate key by itself, so there is no partial dependency with respect to A.
- However, for the combined candidate key AG:
  - $\bullet \quad AG \to D$
  - $AG \rightarrow E$
  - Since A is a subset of AG and:
    - $\bullet$  A  $\rightarrow$  B
    - $\bullet \quad A \to C$
    - $\bullet \quad A \to F$



- These show that B, C, and F are partially dependent on A (a proper subset of AG), thus indicating a partial dependency, which violates 2NF.
- Therefore, the given relation does not satisfy 2NF due to partial dependencies of non-prime attributes on a part of a candidate key.



### • 3NF (Third Normal Form):

-Condition: A relation is in 3NF if it is in 2NF and for every functional dependency  $X \to Y$ , either X is a super key or Y is a prime attribute.

- Explanation:
- The relation is already in 2NF.
- For  $A \rightarrow B, C, F$
- There are no transitive dependencies (dependencies where a non-prime attribute depends on another non-prime attribute).
- Conclusion: The relation is in 3NF because it is in 2NF and for every functional dependency, the left-hand side is a super key.

### -For AG->DE:

- There are no transitive dependencies (dependencies where a non-prime attribute depends on another non-prime attribute).
- Conclusion: The relation is in 3NF because it is in 2NF and for every functional dependency, the left-hand side is a super key.



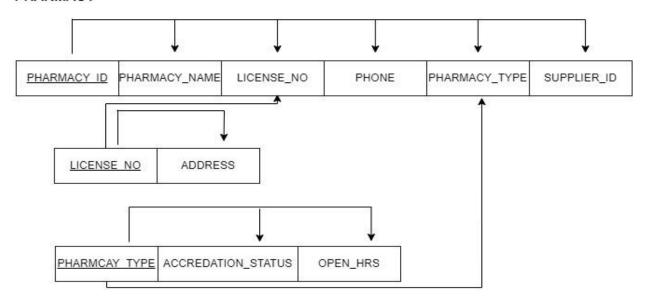
# • BCNF (Boyce-Codd Normal Form):

- Condition: A relation is in BCNF if for every non-trivial functional dependency X → Y,
   X is a superkey.
- Explanation:
- The relation is already in 3NF.
- The given functional dependency  $A \rightarrow B$ , C F and AG ->DE shows that the primary key A and AG determines all other attributes of the same table.

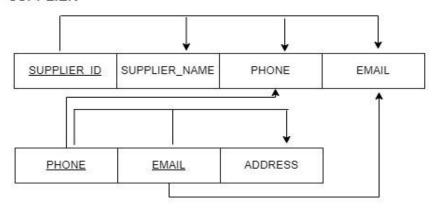


## **FINAL SCHEMA:**

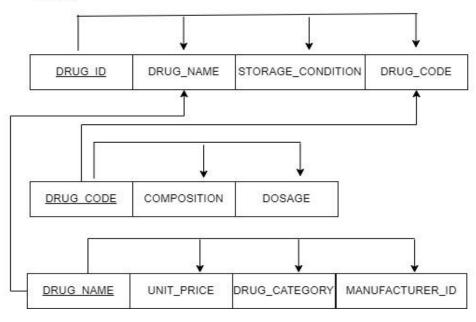
#### PHARMACY



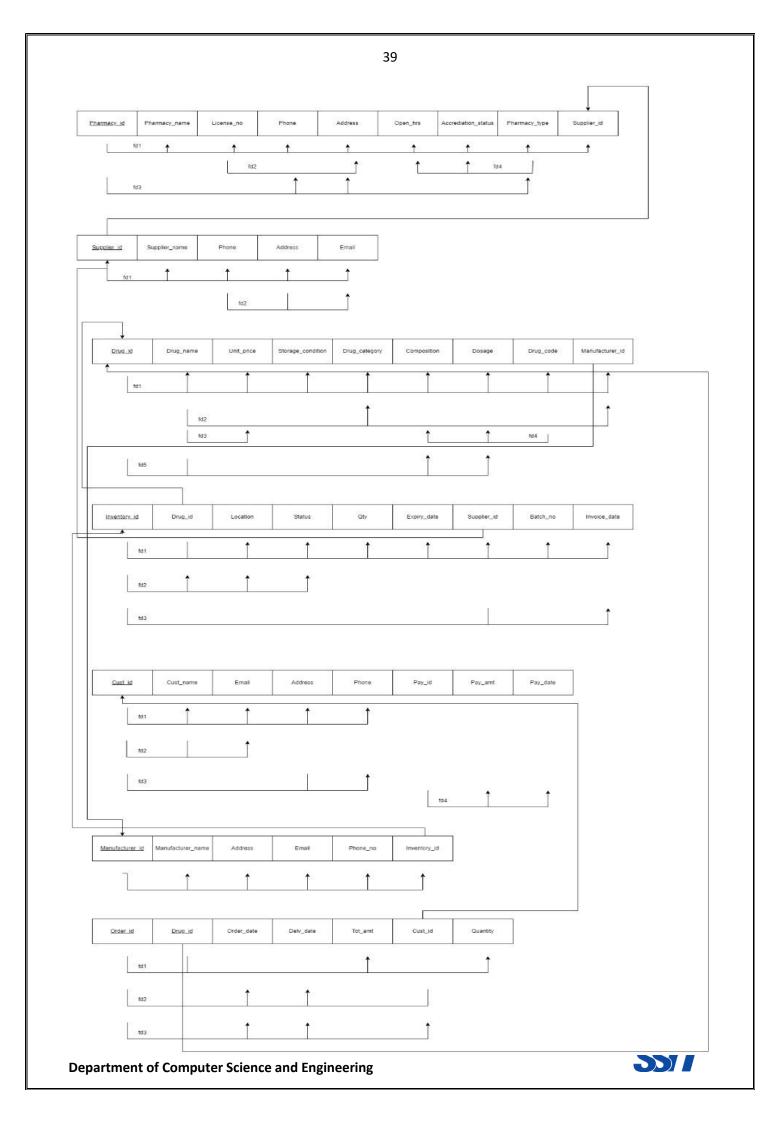
#### SUPPLIER



#### DRUG







## **SQL CODE:**

```
DROP TABLE ORDER LIST;
DROP TABLE ORDERS;
DROP TABLE PHARMACY;
DROP TABLE MANUFACTURER;
DROP TABLE SUPPLIER;
DROP TABLE INVENTORY;
DROP TABLE DRUGS;
DROP TABLE CUSTOMER;
CREATE TABLE CUSTOMER(
    USER_NAME VARCHAR2(100) CONSTRAINT C_ID_PK PRIMARY KEY,
    PASSWORD VARCHAR2(100),
    CUST NAME VARCHAR2(100),
    EMAIL VARCHAR2(100),
    DOOR_NO INT,
    CITY VARCHAR2(100),
    PINCODE NUMERIC(6)
);
CREATE TABLE DRUGS(
    DRUG_ID CHAR(4) CONSTRAINT D_ID_PK PRIMARY KEY,
    DRUG NAME VARCHAR2(100),
    STORAGE_CONDITION VARCHAR2(100),
    DRUG CODE INT,
    UNIT_PRICE INT,
    CONSTRAINT D ID CH CHECK(DRUG ID LIKE 'D%')
);
CREATE TABLE INVENTORY(
    INVENTORY_ID CHAR(4) CONSTRAINT IN_PK PRIMARY KEY,
    DRUG_ID CHAR(4) CONSTRAINT I_FK_D REFERENCES DRUGS(DRUG_ID),
    EXPIRY DATE DATE,
    QTY INT,
    CONSTRAINT I_ID_CH CHECK(INVENTORY_ID LIKE 'I%')
);
CREATE TABLE SUPPLIER(
    SUPPLIER_ID CHAR(4) CONSTRAINT SUPPLIER_PK PRIMARY KEY,
    INVENTORY ID CHAR(4) CONSTRAINT I FK REFERENCES INVENTORY(INVENTORY ID),
    DRUG ID CHAR(4) CONSTRAINT D FK REFERENCES DRUGS(DRUG ID),
    SUPPLIER_NAME VARCHAR2(100),
    PHONE_NO VARCHAR2(10) CONSTRAINT S_PH_U UNIQUE,
    EMAIL VARCHAR2(100),
    CONSTRAINT S_ID_CH CHECK(SUPPLIER_ID LIKE 'S%')
);
```



```
CREATE TABLE MANUFACTURER(
    MANUFACTURER_ID CHAR(4) CONSTRAINT MF_PK PRIMARY KEY,
    MANUFACTURER_NAME VARCHAR2(100),
    DRUG_ID CONSTRAINT D1_FK REFERENCES DRUGS(DRUG_ID),
    ADDRESS VARCHAR2(100),
    PHONE NO NUMERIC CONSTRAINT M PH U UNIQUE,
    EMAIL VARCHAR(100),
    INVENTORY_ID CHAR(4) CONSTRAINT M_FK_I REFERENCES INVENTORY(INVENTORY_ID),
    CONSTRAINT M_ID_CH CHECK(MANUFACTURER_ID LIKE 'M%')
);
CREATE TABLE PHARMACY(
    DRUG_ID CHAR(4) CONSTRAINT D_FK_P REFERENCES DRUGS(DRUG_ID),
    SUPPLIER_ID CHAR(4) CONSTRAINT S_FK_P REFERENCES SUPPLIER(SUPPLIER_ID),
    SUPPLIED_QTY INT,
    PHARMACY_NAME VARCHAR2(100),
    PHONE_NO NUMERIC(10) CONSTRAINT P_PH_U UNIQUE,
    CONSTRAINT P_ID_PK PRIMARY KEY(SUPPLIER_ID,DRUG_ID)
);
CREATE TABLE ORDERS(
    ORDER_ID CHAR(4) CONSTRAINT O_PK PRIMARY KEY,
    USER_NAME VARCHAR2(100) CONSTRAINT O_FK_C REFERENCES CUSTOMER(USER_NAME),
   ORDER_DATE DATE,
    TOTAL_AMT INT,
    CONSTRAINT O_ID_CH CHECK(ORDER_ID LIKE '0%')
);
CREATE TABLE ORDER_LIST(
    ORDER_ID CHAR(4) CONSTRAINT O_FK_OL REFERENCES ORDERS(ORDER_ID),
    DRUG_ID CHAR(4) CONSTRAINT D_FK_O REFERENCES DRUGS(DRUG_ID),
    QTY INT,
    CONSTRAINT CMP_KEY_OL PRIMARY KEY(ORDER_ID,DRUG_ID)
);
INSERT INTO CUSTOMER (USER_NAME, PASSWORD, CUST_NAME, EMAIL, DOOR_NO, CITY,
PINCODE) VALUES
('john_doe', 'password123', 'John Doe', 'john.doe@example.com', 101, 'New York',
10001);
INSERT INTO CUSTOMER (USER_NAME, PASSWORD, CUST_NAME, EMAIL, DOOR_NO, CITY,
PINCODE) VALUES
('jane_smith', 'pass456', 'Jane Smith', 'jane.smith@example.com', 202, 'Los Ange-
les', 90001);
```

```
INSERT INTO CUSTOMER (USER_NAME, PASSWORD, CUST_NAME, EMAIL, DOOR_NO, CITY,
PINCODE) VALUES
('alice_w', 'alice789', 'Alice Wonderland', 'alice.w@example.com', 303, 'Chicago',
60007);
INSERT INTO DRUGS (DRUG ID, DRUG NAME, STORAGE CONDITION, DRUG CODE, UNIT PRICE)
('D001', 'Aspirin', 'Room Temperature', 1001, 10);
INSERT INTO DRUGS (DRUG_ID, DRUG_NAME, STORAGE_CONDITION, DRUG_CODE, UNIT_PRICE)
VALUES
('D002', 'Paracetamol', 'Cool and Dry', 1002, 15);
INSERT INTO DRUGS (DRUG_ID, DRUG_NAME, STORAGE_CONDITION, DRUG_CODE, UNIT_PRICE)
VALUES
('D003', 'Ibuprofen', 'Room Temperature', 1003, 12);
INSERT INTO INVENTORY (INVENTORY_ID, DRUG_ID, EXPIRY_DATE, QTY) VALUES
('I001', 'D001', '31-DEC-2023', 100);
INSERT INTO INVENTORY (INVENTORY ID, DRUG ID, EXPIRY DATE, QTY) VALUES
('I002', 'D002', '30-NOV-2023', 200);
INSERT INTO INVENTORY (INVENTORY_ID, DRUG_ID, EXPIRY_DATE, QTY) VALUES
('I003', 'D003', '31-OCT-2023', 150);
INSERT INTO SUPPLIER (SUPPLIER_ID, INVENTORY_ID, DRUG_ID, SUPPLIER_NAME, PHONE_NO,
EMAIL) VALUES
('S001', 'I001', 'D001', 'Supplier One', '1234567890', 'supplier1@example.com');
INSERT INTO SUPPLIER (SUPPLIER_ID, INVENTORY_ID, DRUG_ID, SUPPLIER_NAME, PHONE_NO,
EMAIL) VALUES
('S002', 'I002', 'D002', 'Supplier Two', '0987654321', 'supplier2@example.com');
INSERT INTO SUPPLIER (SUPPLIER_ID, INVENTORY_ID, DRUG_ID, SUPPLIER_NAME, PHONE_NO,
EMAIL) VALUES
('S003', 'I003', 'D003', 'Supplier Three', '1122334455', 'supplier3@example.com');
INSERT INTO MANUFACTURER (MANUFACTURER_ID, MANUFACTURER_NAME, DRUG_ID, ADDRESS,
PHONE_NO, EMAIL, INVENTORY_ID) VALUES
('M001', 'Manufacturer One', 'D001', '123 Street, City', 1112223333, 'manu1@exam-
ple.com', 'I001');
```



```
INSERT INTO MANUFACTURER (MANUFACTURER_ID, MANUFACTURER_NAME, DRUG_ID, ADDRESS,
PHONE_NO, EMAIL, INVENTORY_ID) VALUES
('M002', 'Manufacturer Two', 'D002', '456 Avenue, City', 4445556666, 'manu2@exam-
ple.com', 'I002');
INSERT INTO MANUFACTURER (MANUFACTURER_ID, MANUFACTURER_NAME, DRUG_ID, ADDRESS,
PHONE NO, EMAIL, INVENTORY ID) VALUES
('M003', 'Manufacturer Three', 'D003', '789 Boulevard, City', 7778889999,
'manu3@example.com', 'I003');
INSERT INTO PHARMACY (DRUG_ID, SUPPLIER_ID, SUPPLIED_QTY, PHARMACY_NAME, PHONE_NO)
VALUES
('D001', 'S001', 50, 'Pharmacy One', 1231231231);
INSERT INTO PHARMACY (DRUG_ID, SUPPLIER_ID, SUPPLIED_QTY, PHARMACY_NAME, PHONE_NO)
VALUES
('D002', 'S002', 75, 'Pharmacy Two', 4564564564);
INSERT INTO PHARMACY (DRUG ID, SUPPLIER ID, SUPPLIED QTY, PHARMACY NAME, PHONE NO)
VALUES
('D003', 'S003', 100, 'Pharmacy Three', 7897897897);
INSERT INTO ORDERS (ORDER_ID, USER_NAME, ORDER_DATE, TOTAL_AMT) VALUES
('0001', 'john_doe', '15-JUN-2023', 150);
INSERT INTO ORDERS (ORDER_ID, USER_NAME, ORDER_DATE, TOTAL_AMT) VALUES
('0002', 'jane_smith', '16-JUN-2023', 200);
INSERT INTO ORDERS (ORDER_ID, USER_NAME, ORDER_DATE, TOTAL_AMT) VALUES
('0003', 'alice_w','17-JUN-2023', 180);
INSERT INTO ORDER LIST (ORDER ID, DRUG ID, QTY) VALUES
('0001', 'D001', 10);
INSERT INTO ORDER_LIST (ORDER_ID, DRUG_ID, QTY) VALUES
('0001', 'D002', 5);
INSERT INTO ORDER_LIST (ORDER_ID, DRUG_ID, QTY) VALUES
('0002', 'D002', 10);
INSERT INTO ORDER_LIST (ORDER_ID, DRUG_ID, QTY) VALUES
('0002', 'D003', 8);
```

```
INSERT INTO ORDER_LIST (ORDER_ID, DRUG_ID, QTY) VALUES
('0003', 'D001', 6);
INSERT INTO ORDER_LIST (ORDER_ID, DRUG_ID, QTY) VALUES
('0003', 'D003', 12);
```



#### **Trigger used:**

Trigger is used to make sure that order date is prior to expiry date.

```
CREATE OR REPLACE TRIGGER trg_check_order_date
BEFORE INSERT ON ORDER_LIST
FOR EACH ROW
DECLARE
    v_order_date ORDERS.ORDER_DATE%TYPE;
    v_expiry_date INVENTORY.EXPIRY_DATE%TYPE;
BEGIN
    -- Retrieve the order date from the ORDERS table
    SELECT ORDER_DATE INTO v_order_date
    FROM ORDERS
   WHERE ORDER_ID = :NEW.ORDER_ID;
    -- Retrieve the expiry date from the INVENTORY table
    SELECT EXPIRY_DATE INTO v_expiry_date
    FROM INVENTORY
   WHERE DRUG_ID = :NEW.DRUG_ID;
    -- Check if the order date is before the expiry date
    IF v_order_date >= v_expiry_date THEN
        RAISE_APPLICATION_ERROR(-20001, 'Order date must be before the expiry date
of the drug.');
   END IF;
END;
```



#### **Procedure used:**

Procedure is used to calculate the bill amount for the order.

```
CREATE OR REPLACE PROCEDURE CALCULATE_TOTAL_AMOUNT (
    p_order_id IN CHAR
) AS
   v_total_amt INT := 0;
BEGIN
    -- Calculate the total amount for the given order ID
    SELECT SUM(o1.QTY * d.UNIT_PRICE)
    INTO v_total_amt
    FROM ORDER LIST ol
    JOIN DRUGS d ON ol.DRUG_ID = d.DRUG_ID
   WHERE ol.ORDER_ID = p_order_id;
    -- Update the total amount in the ORDERS table
   UPDATE ORDERS
    SET TOTAL_AMT = v_total_amt
   WHERE ORDER_ID = p_order_id;
   COMMIT;
END;
```



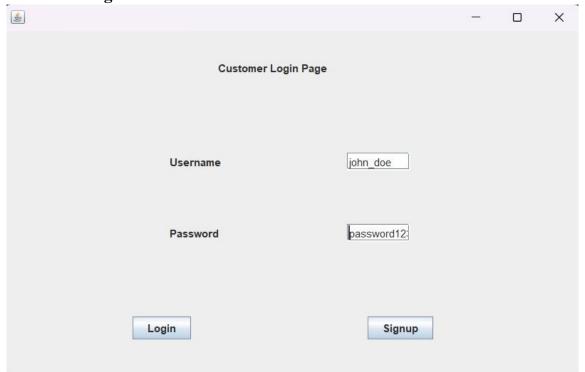
# **OUTPUT SCREENSHOTS:**

### **Login Page:**

This is the basic login page that would be displayed to the user where the user has to enter their login credentials to access the page further.



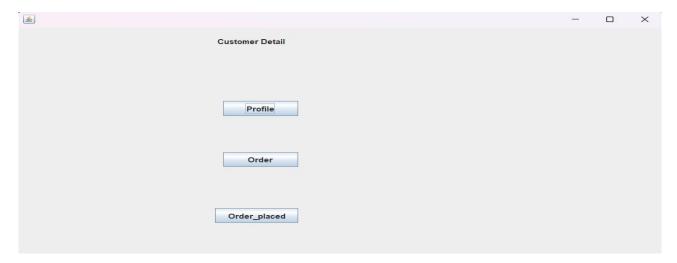
### **Customer Login:**





### Menu Page:

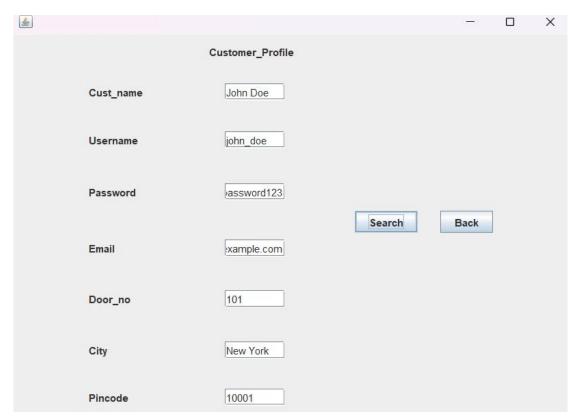
Menu Page will display the various options that are available to the user. The options include gateway to view the profile ,to visit the orders placed, to place new order.



#### **Profile Page:**

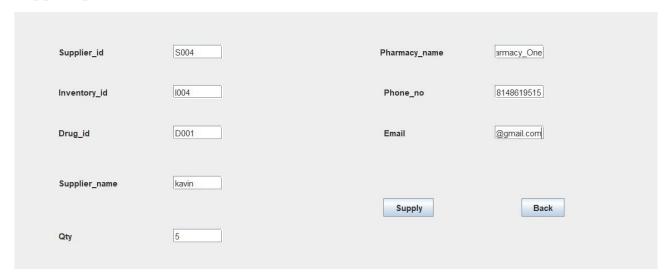
Profile page will display the details of the customer that includes their customer name, user name, password, mail\_id, address details, etc.

#### **Customer Profile:**



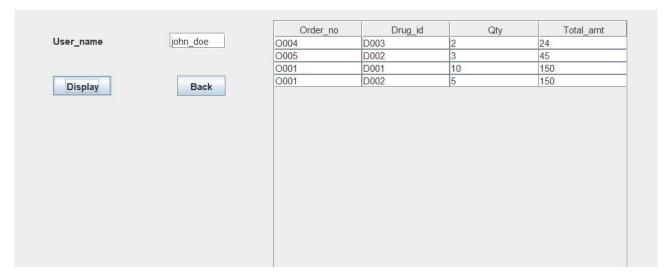


### **Supplier profile:**



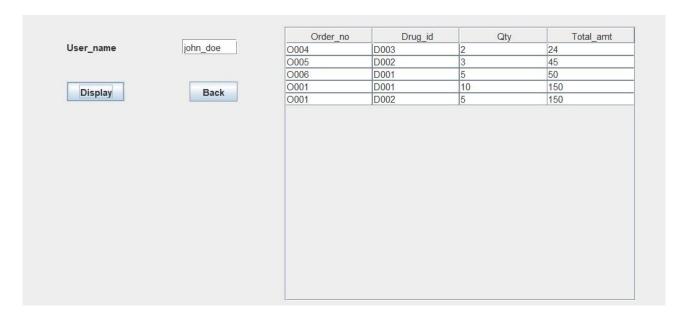
#### **Orders Placed:**

Orders Placed will contain in detail contents of the orders placed by the customer. It provides the details of order\_id,drug id of the drug ordered, quantity orders,total amount.



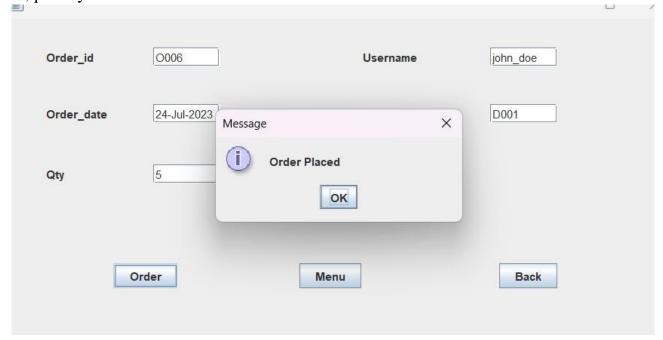
### **Orders List After Placing New order:**





### Placing New order:

Orders page will open the gateway to order new drug where user must give inputs like drug id,quantity to be ordered.

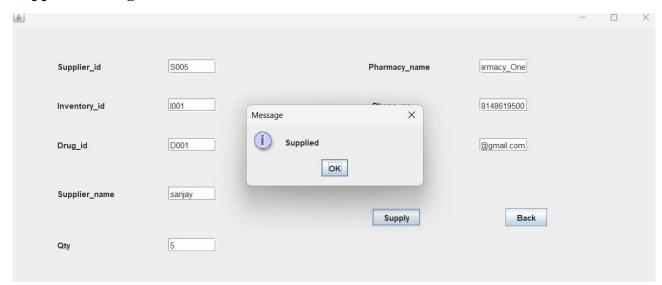




#### **Available Stock:**

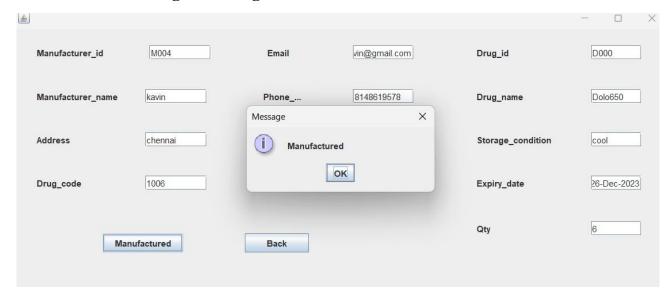


## **Supplier Adding New Stock:**





## **Manufacturer Adding New Drug:**





# **NOVELTY:**

- Deploy the pharmacy management system on a cloud platform to leverage scalability, reliability, and flexibility.
- Use AI to predict future trends, helping the pharmacy to manage inventory more effectively and anticipate demand.
  - Analyze feedback and reviews to continually improve services and user satisfaction
  - Implement advanced security features to protect user data and ensure secure transactions.
  - Integrate the pharmacy management system with external services and platforms for enhanced functionality.

