SIGNED/UNSIGNED INTEGER ARITHMETIC IN C

 $\begin{tabular}{ll} $Vineel\ Kovvuri \\ $http://vineelkovvuri.com \end{tabular}$

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1 Introduction

This article is about understanding how integer conversions happen in C language. C standard defines the integer conversion rules not specific to any architecture. It makes things more complicated for the programmers to try and understand them.

Why do we need integer conversions at all? The answer is simple we need to have single type for the expression. Let's say we have an expression <expr1><op><expr2>when expr1 and expr2 are different types, we want the resulting expression from this to have one type.

The C99 specification defines below set of integer types called 'standard integer types'. Interestingly, the sizes of these integer types are not defined at all. It only defines the minimum supported size. For example, int need not be 8 bytes long on x64 platforms, the only definition for int is it should have at least 16 bits and similarly long should have at least 32 bits it need not be 8 bytes long. Depending on the platform and the processor architecture the ABI (Application Binary Interface) and the 64bit programming model will define the size of these basic types. Windows x64 follows LLP64(meaning only 'long long' and pointer size is 64 bit wide), So below are the sizes of the standard types that we are sticking to in this article.

Type	Size
signed char	1 bytes
unsigned char	1 bytes
signed short	2 bytes
unsigned short	2 bytes
signed int	4 bytes
unsigned int	4 bytes
signed long	4 bytes
unsigned long	4 bytes
signed long long	8 bytes
unsigned long long	8 bytes

Also, The specification leaves other aspects of C language definition undefined and this made room for optimization for the compilers. For example, the result of signed arithmetic leading to overflow/underflow is not defined because the specification predates to the machine architectures when 2's complement representation for -ve numbers is not universal, Even though virtually every architecture now uses 2's complement representation for -ve numbers. Hence the result of unsigned overflow/underflow is well defined but not signed overflow/underflow!

2 How signed-ness is represented in the hardware?

Processors do have the concept of signed/unsigned, but unlike in C language, where this information is baked into the variable type definition, processor registers do not hold the type/signed/unsigned information. After all, registers are just placeholders of data. But the instructions themselves do have the notion of signed/unsigned, That is why we have signed 'imul', unsigned 'mul', signed (JL/JG) vs unsigned(JB/JA) jump instructions. This is an important distinction between programming languages and machine code. It helps in understanding how high-level code gets translated to underlying machine code. Instructions which produce/consume signed data understand that -ve numbers should be in 2's complement form. So if a register (8bit) has 11110110 it's up to the instruction to treat that data as either -10 (signed) or 246 (unsigned). If an operation results in a -ve result say -5 then destination register will be written 11111011(2's complement representation of -5). 'add' and 'sub' instructions themselves are not affected by signed and unsigned numbers because of modulo arithmetic.

3 How signed-ness is interpreted in assembly?

In one of the above paragraphs, we said, processor registers does not hold any signed/unsigned type information with them and it's up to the instructions to interpret the data as either signed or unsigned.

But the interesting thing about x86/x64 processors is, they provide a flag register(EFLAGS) where some of the bits are called Status Flags. These bits represent the status of what happened with the previous instruction.

The most important flags of these for our discussion are

- 1. Carry flag Set if an arithmetic operation generates a carry or a borrow out of the most-significant bit of the result; cleared otherwise. This flag indicates an overflow condition for unsigned-integer arithmetic.
- 2. Sign flag Set equal to the most significant bit of the result, which is the sign bit of a signed integer. (0 indicates a positive value and 1 indicates a negative value.)
- 3. Overflow flag Set if the integer result is **too large a positive number or too small a negative number (excluding the sign-bit) to fit in the destination operand** cleared otherwise. This flag indicates an overflow condition for signed-integer (two's complement) arithmetic

Of the above three status flags, Carry flag is obvious but other two flags need some explanation. Sign flag is set irrespective of whether you are treating the numbers as positive or negative, All it cares is if the result's MSB is set or not.

For example in below code, 'sub' can treat 0x81 as -127 or +129. So the result can also be either interpreted as -128 or +128, But the processor after executing 'sub' indicates this possibility in Sign Flag. So operations (JA/JB) who want to treat the result an unsigned will ignore this sign flag whereas operations (like JG/JL) who treat it as signed will take this sign flag into account.

```
mov al, 0x81
sub al, 0x1 // This triggers sign flag
```

Overflow flag is a little different. First of all, the Overflow flag is not set when an operation results in an overflow of the number that can be represented in a register. Like for example, if we try to add 1 to 0xff does not set this bit.

```
mov al, 0xFF // 0b111111111

add al, 0x1 // This will not trigger Overflow flag
```

Overflow flag is mainly meant for indicating if there is a chance of result if interpreted as signed will cross beyond what can be represented in the register. for example if we take a byte(al) it can contain value from 0b00000000(0) to 0b11111111(255) as unsigned or 0b10000000(-128) to 0b011111111(127) as signed number. Now overflow flag is set when a result goes beyond 0b011111111 i.e., too large a positive number and less than 0b10000000 i.e., too small a negative number

```
mov al, 0x7e // This is always positive add al, 0x2 // This will trigger overflow flag
```

Let us see how overflow flag is used to determine the below conditional when var1 and var2 are signed numbers.

```
if (var1 < var2) {
    printf("var1 < var2");
}</pre>
```

The above code is translated to cmp followed by jl. cmp does var1-var2. And jl says execute printf if SF!= OF. Now lets see the combinations why SF needs to be not equal to OF

```
var1 = -125;
var2 = 10;
```

```
var1-var2 => -135 => which cannot be represented in a byte meaning the sign
bit of the byte will not be set but the overflow bit gets set. So SF=0 and
OF=1

var1 = -120;
var2 = -100;
var1-var2 => -20 => this can be represented in a byte and its -ve so SF = 1
but because it is still under the range of singed numbers representable in a
byte OF = 0
```

That is why jl is defined as SF != OF, When this is true it means var1 <var2 even if var1 and var2 are signed(meaning -ve numbers)

4 Signed vs Unsigned integers in C

Signed numbers are the way -ve numbers are represented and unsigned numbers are the way in which +ve numbers are represented. If let's say a data type has 1 byte of storage then the possible bit representation will be from 0b0000000(0) - 0b11111111(255). Now, these 255 valid slots can be interpreted as all positive numbers or divide the range into two halves and call one half of the slot as positive and other as negative numbers. There are other alternate representations for -ve number but almost all systems use 2's complement notation to represent -ve numbers. The take away here is if we include -ve number in the 255 slots then we will represent from -128 to 127. According to 2's complement notation -128 is represented as 0b10000000(0x80) whereas 127 is represented as 0b011111111 (0x7F).

5 Integer Arithmetic Conversions in C

C language defines below set of rules to convert the arguments in an expression.

5.1 Rank

Every integer type has an integer conversion rank defined as follows:

- 1. No two signed integer types shall have the same rank, even if they have the same representation.
- 2. The rank of a signed integer type shall be greater than the rank of any signed integer type with less precision.
- 3. The rank of long long int shall be greater than the rank of long int, which shall be greater than the rank of int, which shall be greater than the rank of short int, which shall be greater than the rank of signed char.
- 4. The rank of any unsigned integer type shall equal the rank of the corresponding signed integer type, if any.
- 5. The rank of any standard integer type shall be greater than the rank of any extended integer type with the same width.
- 6. The rank of char shall equal the rank of signed char and unsigned char.
- 7. The rank of Bool shall be less than the rank of all other standard integer types.
- 8. The rank of any enumerated type shall equal the rank of the compatible integer type (see 6.7.2.2).
- 9. The rank of any extended signed integer type relative to another extended signed integer type with the same precision is implementation-defined, but still subject to the other rules for determining the integer conversion rank.
- 10. For all integer types T1, T2, and T3, if T1 has greater rank than T2 and T2 has greater rank than T3, then T1 has greater rank than T3.

At a high level the gist of above rules is as follows

```
Rank(signed char) == Rank(unsigned char) < Rank(signed short) == Rank(unsigned of short) < Rank(signed int) == Rank(unsigned int) < Rank(signed long) == Rank(unsigned long long) == Rank(unsigned long long)
```

5.2 Integer Promotions

The following may be used in an expression wherever an int or unsigned int may be used:

- An object or expression with an integer type whose integer conversion rank is less than or equal to the rank of int and unsigned int.
- A bit-field of type Bool, int, signed int, or unsigned int.
- 1a) If an int can represent all values of the original type, the value is converted to an int;
- 1b) otherwise, it is converted to an unsigned int. These are called the integer promotions. All other types are unchanged by the integer promotions.

5.3 Usual arithmetic conversions

- 2a) If both operands have the same type, then no further conversion is needed.
- 2b) Otherwise, if both operands have signed integer types or both have unsigned integer types, the operand with the type of lesser integer conversion rank is converted to the type of the operand with greater rank.
- 2c) Otherwise, if the operand that has unsigned integer type has rank greater or equal to the rank of the type of the other operand, then the operand with signed integer type is converted to the type of the operand with unsigned integer type.
- 2d) Otherwise, if the type of the operand with signed integer type can represent all of the values of the type of the operand with unsigned integer type, then the operand with unsigned integer type is converted to the type of the operand with a signed integer type.
- 2e) Otherwise, both operands are converted to the unsigned integer type corresponding to the type of the operand with signed integer type.

5.3.1 Rule 2c interpretation

The gist of 2c is to make sense of expressions like z = a+b where a is singed number and b is unsigned number. Now a 4 byte signed int is added to 8 byte unsigned long long, according to this rule 4 byte signed int will become 8 byte unsigned long long. The important point to remember here is conversion of signed int to unsigned long long, The way the compiler does this is by sign extend 4 byte signed number so that its absolute value will not change i.e.., -10 remain -10 whether it is a 4 byte or 8 byte. But according to these rules the expression will have the type unsigned long long so statements like below will result in unexpected behavior. Since the var1+var2 is an unsigned expression the compiler will enforce it by using a unsigned jump instruction called jae instead of jge instruction.

Whenever a signed data type has to be converted to unsigned data type the absolute value remains unchanged because of sign extension will be done using movex instruction. So if a variable is type signed short with value -10, when this variable is sign extended to unsigned int the value of even after the conversion remains -10 when interpreted as singed int because the binary represent for -10 for 2 byte storage and 4 byte storage yields the same.

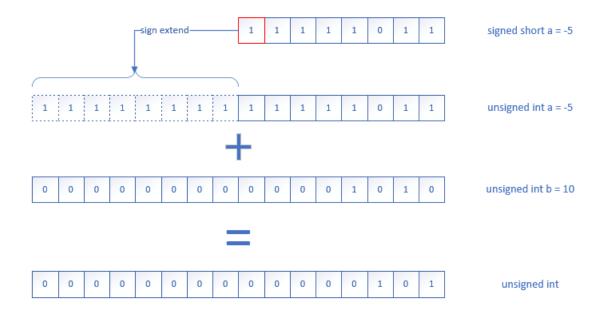


Figure 1: signed number converted to unsigned numbers

godbolt's compiler explorer is very helpful in understanding how the compiler translates expressions and apply the above rules. Clang AST view especially helps in graphically looking at how these promotions/conversion happen

```
signed int var1 = -100;
unsigned long long var2 = 10;
if (var1 + var2 < 0) { //resulting unsigned long long expression will never be

→ lessthan 0
printf("This will never get printed");
}</pre>
```

```
$LN4:
                                                      ; 00000038H
                rsp, 56
        sub
                DWORD PTR var1$[rsp], -100
        mov
                                                            ; fffffffffffff9cH
                QWORD PTR var2$[rsp], 10
        mov
                rax, DWORD PTR var1$[rsp]
        movsxd
                rax, QWORD PTR var2$[rsp]
        add
                rax, rax
        test
                SHORT $LN2@main
                                      //This is an unsigned jump
        jae
                rcx, OFFSET FLAT: $SG4416
        lea
        call
                printf
$LN2@main:
        xor
                eax, eax
        add
                rsp, 56
                                                      ; 00000038H
        ret
main
        ENDP
```

```
Fragment of clang's AST for above program

|-BinaryOperator <line:7:9, col:23> 'bool' '<'
| |-BinaryOperator <col:9, col:16> 'unsigned long long' '+'
| | |-ImplicitCastExpr <col:9> 'unsigned long long' <IntegralCast>
| | | `-ImplicitCastExpr <col:9> 'int' <LValueToRValue>
| | | `-DeclRefExpr <col:9> 'int' lvalue Var 0x55ea65e030e0 'var1' 'int'
```

These rules can only dictate what need to be done for each expression in a statement one at a time. It cannot determine the type of the complete statement at once because of this we might end up with unexpected results if we are not careful enough about how the type is propagating. for example, in var $= \exp 1 + \exp 2 \exp 3$ the rules can only tell how $\exp 2 \exp 3$ interact and how that result interacts with $\exp 2 \exp 3$ s new type may not gel well with $\exp 1$ and might result in an unexpected result.

5.3.2 Rule 2d interpretation

Rule 2d tries to make sense when you have two variables with small the unsigned number and signed number are present in an expression and the singed storage can hold the unsigned number entirely, then by this rule, the unsigned number gets converted to signed number.

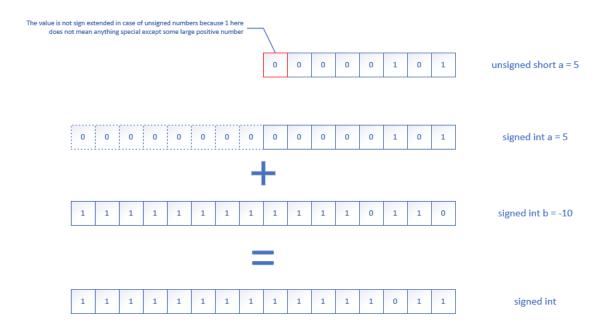


Figure 2: unsigned number converted to signed numbers

For example:

```
unsigned int var1 = 100;
signed long long var2 = -10;
if (var1 + var2 < 0) { //resulting in signed long long expression
    printf("This gets printed if var1 + var2 is < 0");
}</pre>
```

```
      sub
      rsp, 56
      ; 00000038H

      mov
      DWORD PTR var1$[rsp], 100
      ; 00000064H

      mov
      QWORD PTR var2$[rsp], -10

      mov
      eax, DWORD PTR var1$[rsp]

      add
      rax, QWORD PTR var2$[rsp]

      test
      rax, rax
```

```
SHORT $LN2@main
                                                  // This is a signed jump
        jge
                rcx, OFFSET FLAT: $SG4869
        lea
        call
                printf
$LN2@main:
                eax, eax
        xor
        add
                rsp, 56
                                                       ; 00000038H
        ret
main
        ENDP
```

```
Fragment of clang's AST for above program

|-BinaryOperator <lience: 7:9, col: 23> 'bool' '<'
| |-BinaryOperator <col: 9, col: 16> 'long long' '+'
| | |-ImplicitCastExpr <col: 9> 'long long' <IntegralCast>
| | | `-ImplicitCastExpr <col: 9> 'unsigned int' <LValueToRValue>
| | | `-DeclRefExpr <col: 9> 'unsigned int' lvalue Var 0x56305e1200e0
| 'var1' 'unsigned int'
| | `-ImplicitCastExpr <col: 16> 'long long' <LValueToRValue>
| | `-DeclRefExpr <col: 16> 'long long' lvalue Var 0x56305e1201a8 'var2'
| 'long long'
| `-ImplicitCastExpr <col: 23> 'long long' <IntegralCast>
| `-IntegerLiteral <col: 23> 'int' 0
```

5.3.3 Rule 2e interpretation

And finally, Rule 2e is applied only between unsigned int and signed long because neither 2c,2d fits these type on LLP64 model. One interesting thing we can observe if read between the lines is, the resultant type has unsignedness but the type is of signed variable and also both operand types are converted unlike other rules. For example:

```
unsigned int var1 = 100; // this gets converted to unsigned long
signed long var2 = -10; // this will also converted to unsigned long
if (var1 + var2 < 0) { //resulting type is unsigned long
    printf("This will not be printed");
}</pre>
```

```
; 00000038H
        sub
                rsp, 56
                DWORD PTR var1$[rsp], 100
                                              ; 00000064H
        mov
                DWORD PTR var2$[rsp], -10
        mov
                eax, DWORD PTR var2$[rsp]
        mov
                ecx, DWORD PTR var1$[rsp]
        mov
        add
                ecx, eax
        mov
                eax, ecx
        test
                eax, eax
                                              // This is an unsigned jump
        jae
                SHORT $LN2@main
                rcx, OFFSET FLAT: $SG4869
        lea
        call
                printf
$LN2@main:
                eax, eax
                                                      : 00000038H
        add
                rsp, 56
        ret
                0
        ENDP
main
```

```
Fragment of clang's AST for above program
```

5.4 Integer conversion matrix on LLP64 Programming Model

Below table summaries the rules applied for each expression combination

LLP64 Programming Model		Size(MS VC x64)	1	1	. 2	2	4	4	4	4	8	8
		Rank	1	1	. 2	2	3	3	4	4	5	5
		4	signed	unsigned	signed	unsigned	signed	unsigned	signed	unsigned	signed	unsigned
Size(MS VC x64)	Rank	<type1> oper <type2></type2></type1>	char	char	short	short	int	int	long	long	long long	long long
1	1	signed char	1a	1a	1a	1a	1a	1b	2b	2c	2b	2c
1	1	unsigned char		1a	1a	1a	1a	1b	2d	2b	2d	2b
2	2	signed short			1a	1a	1a	1b	2b	2c	2b	2c
2	2	unsigned short				1a	1a	1b	2d	2b	2d	2b
4	3	signed int					2a	1b	2b	2c	2b	2c
4	3	unsigned int						2a	2e	2b	2d	2b
4	4	signed long							2a	2c	2b	2c
4	4	unsigned long								2a	2d	2b
8	5	signed long long									2a	2c
8	5	unsigned long long										2a

Figure 3: Rules applied for each combination of expression

Below table summaries the resulting data type of the expression

		Size(MS VC x64)	1	1	2	2	4	4	4	4	8	8
LLP64 Programr	ming Model	Rank			2				4	4	5	
		Num	signed	unsigned		_	_	unsigned	signed	unsigned		unsigned
Size(MS VC x64)	Rank	<type1> oper <type2></type2></type1>	char	char	short	short	_		long	long	long long	long long
,			signed	signed	signed	signed	signed	unsigned	_	unsigned		unsigned
1	1	signed char	int	int	int	int	int	int	long	long	long long	long long
				signed	signed	signed	signed	unsigned	signed	unsigned	signed	unsigned
1	1	unsigned char		int	int	int	int	int	long	long	long long	long long
					signed	signed	signed	unsigned	signed	unsigned	signed	unsigned
2	2	signed short			int	int	int	int	long	long	long long	long long
						signed	signed	unsigned	signed	unsigned	signed	unsigned
2	2	unsigned short				int	int	int	long	long	long long	long long
							signed	unsigned	signed	unsigned	signed	unsigned
4	3	signed int					int	int	long	long	long long	long long
								unsigned	unsigned	unsigned	signed	unsigned
4	3	unsigned int						int	long	long	long long	long long
									signed	unsigned	signed	unsigned
4	4	signed long							long	long	long long	long long
										unsigned	signed	unsigned
4	4	unsigned long								long	long long	long long
											signed	unsigned
8	5	signed long long									long long	long long
												unsigned
8	5	unsigned long long										long long

Figure 4: Resulting type of the expression

6 References

- 1. INT02-C. Understand integer conversion rules
- 2. C99 Standard

- 3. Should I use Signed or Unsigned Ints In C? (Part 1)
- 4. Should I use Signed or Unsigned Ints In C? (Part 2)
- 5. X86 Opcode and Instruction Reference
- 6. Why is imul used for multiplying unsigned numbers?
- 7. System V Application Binary Interface AMD64 Architecture Processor Supplement
- 8. Why did the Win64 team choose the LLP64 model?
- 9. Abstract Data Models

7 Binary representation

7.1 unsigned numbers

	_		_
0 0x0 00000000	54 0x36 00110110	108 0x6c 01101100	162 0xa2 10100010
1 0x1 00000001	55 0x37 00110111	109 0x6d 01101101	163 0xa3 10100011
2 0x2 00000010	56 0x38 00111000	110 0x6e 01101110	164 0xa4 10100100
3 0x3 00000011	57 0x39 00111001	111 0x6f 01101111	165 0xa5 10100101
4 0x4 00000100	58 0x3a 00111010	112 0x70 01110000	166 0xa6 10100110
5 0x5 00000101	59 0x3b 00111011	113 0x71 01110001	167 0xa7 10100111
6 0x6 00000110	60 0x3c 00111100	114 0x72 01110010	168 0xa8 10101000
7 0x7 00000111	61 0x3d 00111101	115 0x73 01110011	169 0xa9 10101001
8 0x8 00001000	62 0x3e 00111110	116 0x74 01110100	170 0xaa 10101010
9 0x9 00001001	63 0x3f 00111111	117 0x75 01110101	171 0xab 10101011
10 0xa 00001010	64 0x40 01000000	118 0x76 01110110	172 0xac 10101100
11 0xb 00001011	65 0x41 01000001	119 0x77 01110111	173 0xad 10101101
12 0xc 00001100	66 0x42 01000010	120 0x78 01111000	174 0xae 10101110
13 0xd 00001101	67 0x43 01000011	121 0x79 01111001	175 0xaf 10101111
14 0xe 00001110	68 0x44 01000100	122 0x7a 01111010	176 0xb0 10110000
15 0xf 00001111	69 0x45 01000101	123 0x7b 01111011	177 0xb1 10110001
16 0x10 00010000	70 0x46 01000110	124 0x7c 01111100	178 0xb2 10110010
17 0x11 00010001	71 0x47 01000111	125 0x7d 01111101	179 0xb3 10110011
18 0x12 00010010	72 0x48 01001000	126 0x7e 01111110	180 0xb4 10110100
19 0x13 00010011	73 0x49 01001001	127 0x7f 01111111	181 0xb5 10110101
20 0x14 00010100	74 0x4a 01001010	128 0x80 10000000	182 0xb6 10110110
21 0x15 00010101	75 0x4b 01001011	129 0x81 10000001	183 0xb7 10110111
22 0x16 00010110	76 0x4c 01001100	130 0x82 10000010	184 0xb8 10111000
23 0x17 00010111	77 0x4d 01001101	131 0x83 10000011	185 0xb9 10111001
24 0x18 00011000	78 0x4e 01001110	132 0x84 10000100	186 0xba 10111010
25 0x19 00011001	79 0x4f 01001111	133 0x85 10000101	187 0xbb 10111011
26 0x1a 00011010	80 0x50 01010000	134 0x86 10000110	188 0xbc 10111100
27 0x1b 00011011	81 0x51 01010001	135 0x87 10000111	189 0xbd 10111101
28 0x1c 00011100	82 0x52 01010010	136 0x88 10001000	190 0xbe 10111110
29 0x1d 00011101	83 0x53 01010011	137 0x89 10001001	191 0xbf 10111111
30 0x1e 00011110	84 0x54 01010100	138 0x8a 10001010	192 0xc0 11000000
31 0x1f 00011111	85 0x55 01010101	139 0x8b 10001011	193 0xc1 11000001
32 0x20 00100000	86 0x56 01010110	140 0x8c 10001100	194 0xc2 11000010
33 0x21 00100001	87 0x57 01010111	141 0x8d 10001101	195 0xc3 11000011
34 0x22 00100010	88 0x58 01011000	142 0x8e 10001110	196 0xc4 11000100
35 0x23 00100011	89 0x59 01011001	143 0x8f 10001111	197 0xc5 11000101
36 0x24 00100100	90 0x5a 01011010	144 0x90 10010000	198 0xc6 11000110 199 0xc7 11000111
37 0x25 00100101 38 0x26 00100110	91 0x5b 01011011 92 0x5c 01011100	145 0x91 10010001 146 0x92 10010010	200 0xc8 1100111
39 0x27 00100111	93 0x5d 01011101	147 0x93 10010011	200 0xc8 11001000 201 0xc9 11001001
40 0x28 0010111 40 0x28 00101000	94 0x5e 01011101	148 0x94 10010100	202 0xca 11001010
41 0x29 00101000	95 0x5f 01011111	149 0x95 10010101	203 0xcb 11001011
42 0x2a 00101010 42 0x2a 00101010	96 0x60 01100000	150 0x96 10010110	204 0xcc 11001100
43 0x2b 00101011	97 0x61 01100001	151 0x97 10010111	205 0xcd 11001101
44 0x2c 00101100	98 0x62 01100010	152 0x98 1001111000	206 0xce 11001101
45 0x2d 00101101	99 0x63 01100011	153 0x99 10011001	207 0xcf 11001111
46 0x2e 00101110	100 0x64 01100100	154 0x9a 10011010	208 0xd0 11010000
47 0x2f 00101111	101 0x65 01100101	155 0x9b 10011011	209 0xd1 11010001
48 0x30 00110000	102 0x66 01100110	156 0x9c 10011100	210 0xd2 11010010
49 0x31 00110001	103 0x67 01100111	157 0x9d 10011101	211 0xd3 11010011
50 0x32 00110010	104 0x68 01101000	158 0x9e 10011110	212 0xd4 11010100
51 0x33 00110011	105 0x69 01101001	159 0x9f 10011111	213 0xd5 11010101
52 0x34 00110100	106 0x6a 01101010	160 0xa0 10100000	214 0xd6 11010110
53 0x35 00110101	107 0x6b 01101011	161 0xa1 10100001	215 0xd7 11010111
	1	1	I

216 0xd8 11011000	226 0xe2 11100010	236 0xec 11101100	246 0xf6 11110110
217 0xd9 11011001	227 0xe3 11100011	237 0xed 11101101	247 0xf7 11110111
218 0xda 11011010	228 0xe4 11100100	238 0xee 11101110	248 0xf8 11111000
219 0xdb 11011011	229 0xe5 11100101	239 0xef 11101111	249 0xf9 11111001
220 0xdc 11011100	230 0xe6 11100110	240 0xf0 11110000	250 0xfa 11111010
221 0xdd 11011101	231 0xe7 11100111	241 0xf1 11110001	251 0xfb 11111011
222 0xde 11011110	232 0xe8 11101000	242 0xf2 11110010	252 0xfc 11111100
223 0xdf 11011111	233 0xe9 11101001	243 0xf3 11110011	253 0xfd 11111101
224 0xe0 11100000	234 0xea 11101010	244 0xf4 11110100	254 0xfe 11111110
225 0xe1 11100001	235 0xeb 11101011	245 0xf5 11110101	255 0xff 11111111

7.2 signed numbers

S			
-128 0x80 10000000	-72 0xb8 10111000	-16 0xf0 11110000	40 0x28 00101000
-127 0x81 10000001	-71 0xb9 10111001	-15 0xf1 11110001	41 0x29 00101001
-126 0x82 10000010	-70 0xba 10111010	-14 0xf2 11110010	42 0x2a 00101010
-125 0x83 10000011	-69 0xbb 10111011	-13 0xf3 11110011	43 0x2b 00101011
-124 0x84 10000100	-68 0xbc 10111100	-12 0xf4 11110100	44 0x2c 00101100
-123 0x85 10000101	-67 0xbd 10111101	-11 0xf5 11110101	45 0x2d 00101101
-122 0x86 10000110	-66 0xbe 10111110	-10 0xf6 11110110	46 0x2e 00101110
-121 0x87 10000111	-65 0xbf 10111111	-9 0xf7 11110111	47 0x2f 00101111
-120 0x88 10001000	-64 0xc0 11000000	-8 0xf8 11111000	48 0x30 00110000
-119 0x89 10001001	-63 0xc1 11000001	-7 0xf9 11111001	49 0x31 00110001
-118 0x8a 10001010	-62 0xc2 11000010	-6 0xfa 11111010	50 0x32 00110010
-117 0x8b 10001011	-61 0xc3 11000011	-5 0xfb 11111011	51 0x33 00110011
-116 0x8c 10001100	-60 0xc4 11000100	-4 0xfc 11111100	52 0x34 00110100
-115 0x8d 10001101	-59 0xc5 11000101	-3 0xfd 11111101	53 0x35 00110101
-114 0x8e 10001110	-58 0xc6 11000110	-2 0xfe 11111110	54 0x36 00110110
-113 0x8f 10001111	-57 0xc7 11000111	-1 0xff 11111111	55 0x37 00110111
-112 0x90 10010000	-56 0xc8 11001000	0 0x0 00000000	56 0x38 00111000
-111 0x91 10010001	-55 0xc9 11001001	1 0x1 00000001	57 0x39 00111001
-110 0x92 10010010	-54 0xca 11001010	2 0x2 00000010	58 0x3a 00111010
-109 0x93 10010011	-53 0xcb 11001011	3 0x3 00000011	59 0x3b 00111011
-108 0x94 10010100	-52 0xcc 11001100	4 0x4 00000100	60 0x3c 00111100
-107 0x95 10010101	-51 0xcd 11001101	5 0x5 00000101	61 0x3d 00111101
-106 0x96 10010110	-50 0xce 11001110	6 0x6 00000110	62 0x3e 00111110
-105 0x97 10010111	-49 0xcf 11001111	7 0x7 00000111	63 0x3f 00111111
-104 0x98 10011000	-48 0xd0 11010000	8 0x8 00001000	64 0x40 01000000
-103 0x99 10011001	-47 0xd1 11010001	9 0x9 00001001	65 0x41 01000001
-102 0x9a 10011010	-46 0xd2 11010010	10 0xa 00001010	66 0x42 01000010
-101 0x9b 10011011 -100 0x9c 10011100	-45 0xd3 11010011 -44 0xd4 11010100	11 0xb 00001011 12 0xc 00001100	67 0x43 01000011 68 0x44 01000100
-99 0x9d 10011101	-43 0xd5 11010101	13 0xd 00001100	69 0x45 01000100
-98 0x9e 10011101	-42 0xd6 11010110	14 0xe 00001101	70 0x46 01000101
-97 0x9f 10011111	-41 0xd7 11010111	15 0xf 00001110	71 0x47 01000111
-96 0xa0 10100000	-40 0xd8 11011000	16 0x10 00010000	72 0x48 01001000
-95 0xa1 10100001	-39 0xd9 11011001	17 0x11 00010001	73 0x49 01001001
-94 0xa2 10100010	-38 0xda 11011010	18 0x12 00010010	74 0x4a 01001010
-93 0xa3 10100011	-37 0xdb 11011011	19 0x13 00010011	75 0x4b 01001011
-92 0xa4 10100100	-36 0xdc 11011100	20 0x14 00010100	76 0x4c 01001100
-91 0xa5 10100101	-35 0xdd 11011101	21 0x15 00010101	77 0x4d 01001101
-90 0xa6 10100110	-34 0xde 11011110	22 0x16 00010110	78 0x4e 01001110
-89 0xa7 10100111	-33 0xdf 11011111	23 0x17 00010111	79 0x4f 01001111
-88 0xa8 10101000	-32 0xe0 11100000	24 0x18 00011000	80 0x50 01010000
-87 0xa9 10101001	-31 0xe1 11100001	25 0x19 00011001	81 0x51 01010001
-86 0xaa 10101010	-30 0xe2 11100010	26 0x1a 00011010	82 0x52 01010010
-85 0xab 10101011	-29 0xe3 11100011	27 0x1b 00011011	83 0x53 01010011
-84 0xac 10101100	-28 0xe4 11100100	28 0x1c 00011100	84 0x54 01010100
-83 0xad 10101101	-27 0xe5 11100101	29 0x1d 00011101	85 0x55 01010101
-82 0xae 10101110	-26 0xe6 11100110	30 0x1e 00011110	86 0x56 01010110
-81 0xaf 10101111	-25 0xe7 11100111	31 0x1f 00011111	87 0x57 01010111
-80 0xb0 10110000	-24 0xe8 11101000	32 0x20 00100000	88 0x58 01011000
-79 0xb1 10110001	-23 0xe9 11101001	33 0x21 00100001	89 0x59 01011001
-78 0xb2 10110010	-22 0xea 11101010	34 0x22 00100010	90 0x5a 01011010
-77 0xb3 10110011	-21 0xeb 11101011	35 0x23 00100011	91 0x5b 01011011
-76 0xb4 10110100	-20 0xec 11101100	36 0x24 00100100	92 0x5c 01011100
-75 0xb5 10110101	-19 0xed 11101101	37 0x25 00100101	93 0x5d 01011101
-74 0xb6 10110110	-18 0xee 11101110	38 0x26 00100110	94 0x5e 01011110
-73 0xb7 10110111	-17 0xef 11101111	39 0x27 00100111	95 0x5f 01011111
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96 0x60 01100000	104 0x68 01101000	112 0x70 01110000	120 0x78 01111000
97 0x61 01100001	105 0x69 01101001	113 0x71 01110001	121 0x79 01111001
98 0x62 01100010	106 0x6a 01101010	114 0x72 01110010	122 0x7a 01111010
99 0x63 01100011	107 0x6b 01101011	115 0x73 01110011	123 0x7b 01111011
100 0x64 01100100	108 0x6c 01101100	116 0x74 01110100	124 0x7c 01111100
101 0x65 01100101	109 0x6d 01101101	117 0x75 01110101	125 0x7d 01111101
102 0x66 01100110	110 0x6e 01101110	118 0x76 01110110	126 0x7e 01111110
103 0x67 01100111	111 0x6f 01101111	119 0x77 01110111	127 0x7f 01111111